

# Optimal Approximation Rate of ReLU Networks in terms of Width and Depth\*

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## Abstract

This paper concentrates on the approximation power of deep feed-forward neural networks in terms of width and depth. It is proved by construction that ReLU networks with width  $\mathcal{O}(\max\{d\lfloor N^{1/d} \rfloor, N+2\})$  and depth  $\mathcal{O}(L)$  can approximate a Hölder continuous function on  $[0, 1]^d$  with an approximation rate  $\mathcal{O}(\lambda\sqrt{d}(N^2L^2\log_3(N+2))^{-\alpha/d})$ , where  $\alpha \in (0, 1]$  and  $\lambda > 0$  are Hölder order and constant, respectively. Such a rate is optimal up to a constant in terms of width and depth separately, while existing results are only nearly optimal without the log factor in the approximation rate. More generally, for an arbitrary continuous function  $f$  on  $[0, 1]^d$ , the approximation rate becomes  $\mathcal{O}(\sqrt{d}\omega_f((N^2L^2\log_3(N+2))^{-1/d}))$ , where  $\omega_f(\cdot)$  is the modulus of continuity. We also extend our analysis to any continuous function  $f$  on a bounded set. Particularly, if ReLU networks with depth 31 and width  $\mathcal{O}(N)$  is used to approximate one-dimensional Lipschitz continuous functions on  $[0, 1]$  with a constant  $\lambda > 0$ , the approximation rate in terms of the total number of parameters,  $W$ , becomes  $\mathcal{O}(\frac{\lambda}{W \ln W})$ , which has not been discovered in the literature.

**Key words.** Deep ReLU Networks; Hölder Continuity; Optimal Approximation Theory; Bit Extraction; VC-dimension.

## 1 Introduction

Over the past few decades, the expressiveness of neural networks has been widely studied from many points of view, e.g. in terms of combinatorics [18], topology [4], Vapnik-Chervonenkis (VC) dimension [3, 8, 21], fat-shattering dimension [1, 12], information theory [20], classical approximation theory [2, 5, 7, 9, 13, 15, 22, 22, 23, 24, 25, 26, 28, 29], optimization [10, 11, 19], etc. The error analysis of neural networks consists of three parts: the approximation error, the optimization error, and the generalization error. This paper focuses on the approximation error for ReLU networks.

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The approximation errors of feed-forward neural networks with various activation functions have been studied for different types of functions, e.g., smooth functions [6, 14, 15, 16, 27], piecewise smooth functions [20], band-limited functions [17], continuous functions [23, 24, 25, 28]. In [23], it was shown that a ReLU network with width  $C_1(d) \cdot N$  and depth  $C_2(d) \cdot L$  can attain an approximation error  $C_3(d) \cdot \omega_f(N^{-2/d} L^{-2/d})$  to approximate a continuous function  $f$  on  $[0, 1]^d$ , where  $C_1(d)$ ,  $C_2(d)$ , and  $C_3(d)$  are three constants in  $d$  with explicit formulas to specify their values, and  $\omega_f(\cdot)$  is the modulus of continuity of  $f \in C([0, 1]^d)$  defined via

$$\omega_f(r) := \sup \{ |f(\mathbf{x}) - f(\mathbf{y})| : \mathbf{x}, \mathbf{y} \in [0, 1]^d, \|\mathbf{x} - \mathbf{y}\|_2 \leq r \}, \quad \text{for any } r \geq 0.$$

Such an approximation rate is optimal in terms of  $N$  and  $L$  up to a logarithmic term and the corresponding optimal approximation theory is still open. To address this open problem, we provide a constructive proof in this paper to show that ReLU networks of width  $\mathcal{O}(N)$  and depth  $\mathcal{O}(L)$  can approximate an arbitrary continuous function  $f$  on  $[0, 1]^d$  with an optimal approximation error  $\mathcal{O}(\omega_f((N^2 L^2 \ln N)^{-\alpha/d}))$  in terms of  $N$  and  $L$ . As shown by our main result, Theorem 1.1 below, the approximation rate obtained here admits explicit formulas to specify its prefactors when  $\omega_f(\cdot)$  is known.

**Theorem 1.1.** *Given a continuous function  $f \in C([0, 1]^d)$ , for any  $N \in \mathbb{N}^+$ ,  $L \in \mathbb{N}^+$ , and  $p \in [1, \infty]$ , there exists a function  $\phi$  implemented by a ReLU network with width  $C_1 \max\{d\lfloor N^{1/d} \rfloor, N + 2\}$  and depth  $11L + C_2$  such that*

$$\|f - \phi\|_{L^p([0, 1]^d)} \leq 131\sqrt{d}\omega_f\left((N^2 L^2 \log_3(N + 2))^{-1/d}\right),$$

where  $C_1 = 16$  and  $C_2 = 18$  if  $p \in [1, \infty)$ ;  $C_1 = 3^{d+3}$  and  $C_2 = 18 + 2d$  if  $p = \infty$ .

Note that  $3^{d+3} \max\{d\lfloor N^{1/d} \rfloor, N + 2\} \leq 3^{d+3} \max\{dN, 3N\} \leq 3^{d+4}dN$ . Given any  $\tilde{N}, \tilde{L} \in \mathbb{N}^+$  with  $\tilde{N} \geq 3^{d+4}d$  and  $\tilde{L} \geq 29 + 2d$ , there exist  $N, L \in \mathbb{N}^+$  such that

$$3^{d+4}dN \leq \tilde{N} < 3^{d+4}d(N + 1) \quad \text{and} \quad 11L + 18 + 2d \leq \tilde{L} < 11(L + 1) + 18 + 2d.$$

It follows that

$$N \geq \frac{N + 1}{3} > \frac{\tilde{N}}{3^{d+5}d} \quad \text{and} \quad L \geq \frac{L + 1}{2} > \frac{1}{2} \cdot \frac{\tilde{L} - 18 - 2d}{11} = \frac{\tilde{L} - 18 - 2d}{22}.$$

Then we have an immediate corollary from Theorem 1.1.

**Corollary 1.2.** *Given a continuous function  $f \in C([0, 1]^d)$ , for any  $\tilde{N}, \tilde{L} \in \mathbb{N}^+$  with  $\tilde{N} \geq 3^{d+4}d$  and  $\tilde{L} \geq 29 + 2d$ , there exists a function  $\phi$  implemented by a ReLU network with width  $\tilde{N}$  and depth  $\tilde{L}$  such that*

$$\|f - \phi\|_{L^\infty([0, 1]^d)} \leq 131\sqrt{d}\omega_f\left(\left(\left(\frac{\tilde{N}}{3^{d+5}d}\right)^2 \left(\frac{\tilde{L} - 18 - 2d}{22}\right)^2 \log_3\left(\frac{\tilde{N}}{3^{d+5}d} + 2\right)\right)^{-1/d}\right).$$

As a special case of Theorem 1.1 for explicit error characterization, let us take Hölder continuous functions as an example. Let  $\text{Hölder}([0, 1]^d, \alpha, \lambda)$  denote the space of Hölder continuous functions on  $[0, 1]^d$  of order  $\alpha \in (0, 1]$  with a Hölder constant  $\lambda > 0$ . We have an immediate corollary of Theorem 1.1 as follows.

**Corollary 1.3.** *Given a Hölder continuous function  $f \in \text{Hölder}([0, 1]^d, \alpha, \lambda)$ , for any  $N \in \mathbb{N}^+$ ,  $L \in \mathbb{N}^+$ , and  $p \in [1, \infty]$ , there exists a function  $\phi$  implemented by a ReLU network with width  $C_1 \max\{d\lfloor N^{1/d} \rfloor, N + 2\}$  and depth  $11L + C_2$  such that*

$$\|f - \phi\|_{L^p([0, 1]^d)} \leq 131\lambda\sqrt{d}(N^2L^2\log_3(N + 2))^{-\alpha/d},$$

where  $C_1 = 16$  and  $C_2 = 18$  if  $p \in [1, \infty)$ ;  $C_1 = 3^{d+3}$  and  $C_2 = 18 + 2d$  if  $p = \infty$ .

To better illustrate the importance of our theory, we summarize our key contributions as follows.

(1) Upper bound: We provide a quantitative and non-asymptotic approximation rate  $131\sqrt{d}\omega_f((N^2L^2\log_3(N + 2))^{-1/d})$  in terms of width  $\mathcal{O}(N)$  and depth  $\mathcal{O}(L)$  for any  $f \in C([0, 1]^d)$  in Theorem 1.1.

(1.1) This approximation error analysis can be extended to  $f \in C(E)$  for any  $E \subseteq [-R, R]^d$  with  $R > 0$  as we shall see later in Theorem 2.5.

(1.2) In the case of one-dimensional Lipschitz continuous functions on  $[0, 1]$  with a constant  $\lambda > 0$ , the approximation rate in Theorem 1.1 becomes  $\mathcal{O}(\frac{\lambda}{W \ln W})$  for ReLU networks with 31 hidden layers and  $\mathcal{O}(W)$  parameters. To the best of our knowledge, the approximation rate  $\mathcal{O}(\frac{1}{W \ln W})$  is better than existing known results for approximating Lipschitz continuous functions on  $[0, 1]$ .

(2) Lower bound: Through the VC-dimension bounds of ReLU networks given in [8], we show that the approximation rate  $131\lambda\sqrt{d}(N^2L^2\log_3(N + 2))^{-\alpha/d}$  in terms of width  $\mathcal{O}(N)$  and depth  $\mathcal{O}(L)$  for Hölder( $[0, 1]^d, \alpha, \lambda$ ) is optimal as follows.

(2.1) When the width is fixed, both the approximation upper and lower bounds take the form of  $CL^{-2\alpha/d}$  for a positive constant  $C$ .

(2.2) When the depth is fixed, both the approximation upper and lower bounds take the form of  $C(N^2 \ln N)^{-\alpha/d}$  for a positive constant  $C$ .

We would like to point out that if  $N$  and  $L$  vary simultaneously, the rate is optimal in the  $N$ - $L$  plane except for a small region as shown in Figure 1. See Section 2.3 for a detailed discussion. The earlier result in [23] provides a nearly optimal approximation error that has a gap (a logarithmic term) between the lower and upper bounds. It is technically challenging to match the upper bound with the lower bound. Compared to the nearly optimal rate  $19\lambda\sqrt{d}N^{-2\alpha/d}L^{-2\alpha/d}$  for Hölder continuous functions in Hölder( $[0, 1]^d, \alpha, \lambda$ ) in [23], this paper achieves the optimal rate  $131\lambda\sqrt{d}(N^2L^2\log_3(N + 2))^{-\alpha/d}$  using more technical and sophisticated construction. For example, a novel bit extraction technique different to that in [3] is proposed, and new ReLU networks are constructed to approximate step functions more efficiently than those in [23]. The optimal result obtained in this paper could also be extended to other functions spaces, leading to better understanding of deep network approximation.

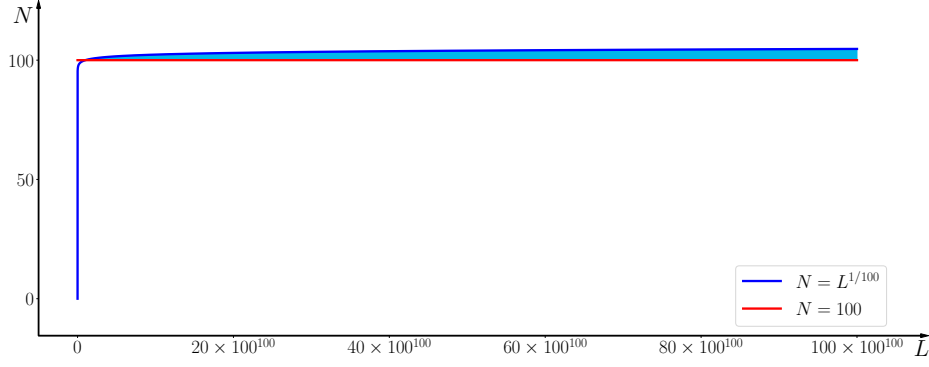


Figure 1: Our rate is optimal except for the region marked in cyan characterized by  $\{(N, L) \in \mathbb{N}^+ : C_1 \leq N \leq L^{C_2}\}$  for two positive constants  $C_1 = 100, C_2 = 1/100$ .

The error analysis of deep learning is to estimate approximation, generalization, and optimization errors. Here, we give a brief discussion, the interested reader can find more details in [15]. Let  $\phi(\mathbf{x}; \boldsymbol{\theta})$  denote a function computed by a network parameterized with  $\boldsymbol{\theta}$ . Given a target function  $f$ , the final goal is to find the expected risk minimizer

$$\boldsymbol{\theta}_{\mathcal{D}} := \arg \min_{\boldsymbol{\theta}} R_{\mathcal{D}}(\boldsymbol{\theta}), \quad \text{where } R_{\mathcal{D}}(\boldsymbol{\theta}) := \mathbb{E}_{\mathbf{x} \sim U(\mathcal{X})} [\ell(\phi(\mathbf{x}; \boldsymbol{\theta}), f(\mathbf{x}))]$$

with a loss function  $\ell(\cdot, \cdot)$  and an unknown data distribution  $U(\mathcal{X})$ .

In practice, for given samples  $\{(\mathbf{x}_i, f(\mathbf{x}_i))\}_{i=1}^n$ , the goal of supervised learning is to identify the empirical risk minimizer

$$\boldsymbol{\theta}_{\mathcal{S}} := \arg \min_{\boldsymbol{\theta}} R_{\mathcal{S}}(\boldsymbol{\theta}), \quad \text{where } R_{\mathcal{S}}(\boldsymbol{\theta}) := \frac{1}{n} \sum_{i=1}^n \ell(\phi(\mathbf{x}_i; \boldsymbol{\theta}), f(\mathbf{x}_i)).$$

In fact, one could only get a numerical minimizer  $\boldsymbol{\theta}_{\mathcal{N}}$  via a numerical optimization method. The discrepancy between the target function and the learned function  $\phi(\mathbf{x}; \boldsymbol{\theta}_{\mathcal{N}})$  is measured by  $R_{\mathcal{D}}(\boldsymbol{\theta}_{\mathcal{N}})$ , which is bounded by

$$R_{\mathcal{D}}(\boldsymbol{\theta}_{\mathcal{N}}) \leq \underbrace{R_{\mathcal{D}}(\boldsymbol{\theta}_{\mathcal{D}})}_{\text{Approximation error}} + \underbrace{[R_{\mathcal{S}}(\boldsymbol{\theta}_{\mathcal{N}}) - R_{\mathcal{S}}(\boldsymbol{\theta}_{\mathcal{S}})]}_{\text{Optimization error}} + \underbrace{[R_{\mathcal{D}}(\boldsymbol{\theta}_{\mathcal{N}}) - R_{\mathcal{S}}(\boldsymbol{\theta}_{\mathcal{N}})] + [R_{\mathcal{S}}(\boldsymbol{\theta}_{\mathcal{D}}) - R_{\mathcal{D}}(\boldsymbol{\theta}_{\mathcal{D}})]}_{\text{Generalization error}}.$$

This paper deals with the approximation error of ReLU networks for continues functions and gives an upper bound of  $R_{\mathcal{D}}(\boldsymbol{\theta}_{\mathcal{D}})$  which is optimal up to a constant. Note that the approximation error analysis given here is independent of data samples and deep learning algorithms. However, the analysis of optimization and generalization errors do depend on data samples, deep learning algorithms, models, etc.

The rest of this paper is organized as follows. In Section 2, we prove Theorem 1.1 by assuming Theorem 2.1 is true, show the optimality of Theorem 1.1, and extend our analysis to continuous functions defined on any bounded set. Next, Theorem 2.1 is proved in Section 3 based on Proposition 3.1 and 3.2, the proofs of which can be found in Section 4. Finally, Section 5 concludes this paper with a short discussion.

## 2 Theoretical analysis

In this section, we first prove Theorem 1.1 and discuss its optimality by assuming Theorem 2.1 is true. Next, we extend our analysis to general continuous functions defined on any bounded set in  $\mathbb{R}^d$ . Notations throughout this paper is summarized in Section 2.1.

### 2.1 Notations

Let us summarize all basic notations used in this paper as follows.

- Matrices are denoted by bold uppercase letters. For instance,  $\mathbf{A} \in \mathbb{R}^{m \times n}$  is a real matrix of size  $m \times n$ , and  $\mathbf{A}^T$  denotes the transpose of  $\mathbf{A}$ . Vectors are denoted as bold lowercase letters. For example,  $\mathbf{v} = [v_1, \dots, v_d]^T = \begin{bmatrix} v_1 \\ \vdots \\ v_d \end{bmatrix} \in \mathbb{R}^d$  is a column vector with  $\mathbf{v}(i) = v_i$  being the  $i$ -th element. Besides, “[” and “]” are used to partition matrices (vectors) into blocks, e.g.,  $\mathbf{A} = \begin{bmatrix} \mathbf{A}_{11} & \mathbf{A}_{12} \\ \mathbf{A}_{21} & \mathbf{A}_{22} \end{bmatrix}$ .
- For any  $p \in [1, \infty)$ , the  $p$ -norm (or  $\ell^p$ -norm) of a vector  $\mathbf{x} = [x_1, x_2, \dots, x_d]^T \in \mathbb{R}^d$  is defined by
$$\|\mathbf{x}\|_p := (|x_1|^p + |x_2|^p + \dots + |x_d|^p)^{1/p}.$$
- For any  $x \in \mathbb{R}$ , let  $\lfloor x \rfloor := \max\{n : n \leq x, n \in \mathbb{Z}\}$  and  $\lceil x \rceil := \min\{n : n \geq x, n \in \mathbb{Z}\}$ .
- Assume  $\mathbf{n} \in \mathbb{N}^d$ , then  $f(\mathbf{n}) = \mathcal{O}(g(\mathbf{n}))$  means that there exists positive  $C$  independent of  $\mathbf{n}$ ,  $f$ , and  $g$  such that  $f(\mathbf{n}) \leq Cg(\mathbf{n})$  when all entries of  $\mathbf{n}$  go to  $+\infty$ .
- For any  $\theta \in [0, 1)$ , suppose its binary representation is  $\theta = \sum_{\ell=1}^{\infty} \theta_{\ell} 2^{-\ell}$  with  $\theta_{\ell} \in \{0, 1\}$ , we introduce a special notation  $\text{bin}0.\theta_1\theta_2\cdots\theta_L$  to denote the  $L$ -term binary representation of  $\theta$ , i.e.,  $\text{bin}0.\theta_1\theta_2\cdots\theta_L := \sum_{\ell=1}^L \theta_{\ell} 2^{-\ell}$ .
- Let  $\mu(\cdot)$  denote the Lebesgue measure.
- Let  $1_S$  be the characteristic function on a set  $S$ , i.e.,  $1_S$  is equal to 1 on  $S$  and 0 outside  $S$ .
- Let  $|S|$  denote the size of a set  $S$ , i.e., the number of all elements in  $S$ .
- The set difference of two sets  $A$  and  $B$  is denoted by  $A \setminus B := \{x : x \in A, x \notin B\}$ .
- Given any  $K \in \mathbb{N}^+$  and  $\delta \in (0, \frac{1}{K})$ , define a trifling region  $\Omega([0, 1]^d, K, \delta)$  of  $[0, 1]^d$  as

$$\Omega([0, 1]^d, K, \delta) := \bigcup_{j=1}^d \left\{ \mathbf{x} = [x_1, x_2, \dots, x_d]^T \in [0, 1]^d : x_j \in \bigcup_{k=1}^{K-1} \left( \frac{k}{K} - \delta, \frac{k}{K} \right) \right\}. \quad (2.1)$$

In particular,  $\Omega([0, 1]^d, K, \delta) = \emptyset$  if  $K = 1$ . See Figure 2 for two examples of trifling regions.

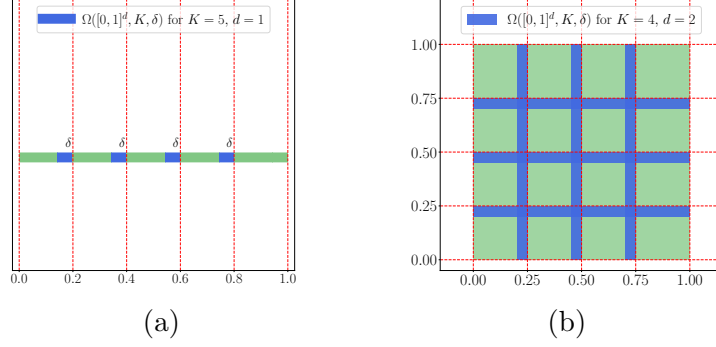


Figure 2: Two examples of trifling regions. (a)  $K = 5, d = 1$ . (b)  $K = 4, d = 2$ .

- Let  $\text{H\"older}([0, 1]^d, \alpha, \lambda)$  denote the space of H\"older continuous functions on  $[0, 1]^d$  of order  $\alpha \in (0, 1]$  with a H\"older constant  $\lambda > 0$ .
- For a continuous piecewise linear function  $f(x)$ , the  $x$  values where the slope changes are typically called **breakpoints**.
- Let  $\text{CPwL}(\mathbb{R}, n)$  denote the space that consists of all continuous piecewise linear functions with at most  $n$  breakpoints on  $\mathbb{R}$ .
- Let  $\sigma : \mathbb{R} \rightarrow \mathbb{R}$  denote the rectified linear unit (ReLU), i.e.  $\sigma(x) = \max\{0, x\}$ . With a slight abuse of notation, we define  $\sigma : \mathbb{R}^d \rightarrow \mathbb{R}^d$  as  $\sigma(\mathbf{x}) = \begin{bmatrix} \max\{0, x_1\} \\ \vdots \\ \max\{0, x_d\} \end{bmatrix}$  for any  $\mathbf{x} = [x_1, \dots, x_d]^T \in \mathbb{R}^d$ .
- We will use  $\mathcal{NN}$  to denote a function implemented by a ReLU network for short and use Python-type notations to specify a class of functions implemented by ReLU networks with several conditions, e.g.,  $\mathcal{NN}(c_1; c_2; \dots; c_m)$  is a set of functions implemented by ReLU networks satisfying  $m$  conditions given by  $\{c_i\}_{1 \leq i \leq m}$ , each of which may specify the number of inputs (`#input`), the number of outputs (`#output`), the total number of neurons in all hidden layers (`#neuron`), the number of hidden layers (`depth`), the total number of parameters (`#parameter`), and the width in each hidden layer (`widthvec`), the maximum width of all hidden layers (`width`), etc. For example, if  $\phi \in \mathcal{NN}(\text{\#input} = 2; \text{widthvec} = [100, 100]; \text{\#output} = 1)$ , then  $\phi$  is a functions satisfies
  - $\phi$  maps from  $\mathbb{R}^2$  to  $\mathbb{R}$ .
  - $\phi$  can be implemented by a ReLU network with two hidden layers and the number of nodes in each hidden layer is 100.
- For any function  $\phi \in \mathcal{NN}(\text{\#input} = d; \text{widthvec} = [N_1, N_2, \dots, N_L]; \text{\#output} = 1)$ , if we set  $N_0 = d$  and  $N_{L+1} = 1$ , then the architecture of the network implementing  $\phi$  can be briefly described as follows:

$$\mathbf{x} = \tilde{\mathbf{h}}_0 \xrightarrow{\mathbf{W}_0, \mathbf{b}_0} \mathbf{h}_1 \xrightarrow{\sigma} \tilde{\mathbf{h}}_1 \dots \xrightarrow{\mathbf{W}_{L-1}, \mathbf{b}_{L-1}} \mathbf{h}_L \xrightarrow{\sigma} \tilde{\mathbf{h}}_L \xrightarrow{\mathbf{W}_L, \mathbf{b}_L} \mathbf{h}_{L+1} = \phi(\mathbf{x}),$$

where  $\mathbf{W}_i \in \mathbb{R}^{N_{i+1} \times N_i}$  and  $\mathbf{b}_i \in \mathbb{R}^{N_{i+1}}$  are the weight matrix and the bias vector in the  $i$ -th affine linear transform  $\mathcal{L}_i$  in  $\phi$ , respectively, i.e.,

$$\mathbf{h}_{i+1} = \mathbf{W}_i \cdot \tilde{\mathbf{h}}_i + \mathbf{b}_i =: \mathcal{L}_i(\tilde{\mathbf{h}}_i), \quad \text{for } i = 0, 1, \dots, L,$$

and

$$\tilde{\mathbf{h}}_i = \sigma(\mathbf{h}_i), \quad \text{for } i = 1, \dots, L.$$

In particular,  $\phi$  can be represented in a form of function compositions as follows

$$\phi = \mathcal{L}_L \circ \sigma \circ \mathcal{L}_{L-1} \circ \sigma \circ \dots \circ \sigma \circ \mathcal{L}_1 \circ \sigma \circ \mathcal{L}_0,$$

which has been illustrated in Figure 3.

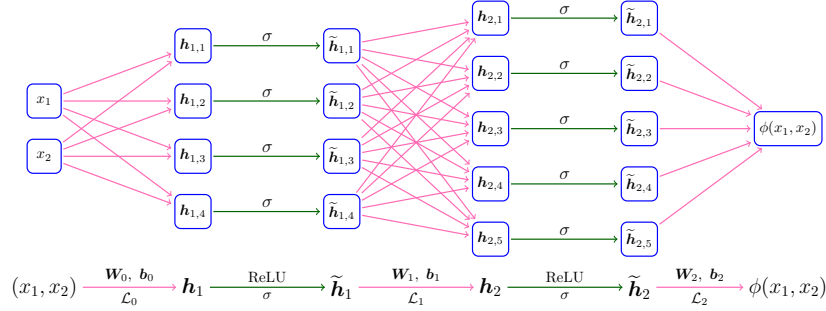


Figure 3: An example of a ReLU network with width 5 and depth 2.

- The expression “a network with width  $N$  and depth  $L$ ” means
  - The maximum width of this network for all **hidden** layers is no more than  $N$ .
  - The number of **hidden** layers of this network is no more than  $L$ .

## 2.2 Proof of Theorem 1.1

The key point is to construct piecewise constant functions to approximate continuous functions in the proof. However, it is impossible to construct a piecewise constant function implemented by a ReLU network due to the continuity of ReLU networks. Thus, we introduce the trifling region  $\Omega([0, 1]^d, K, \delta)$ , defined in Equation (2.1), and use ReLU networks to implement piecewise constant functions outside the trifling region. To prove Theorem 1.1, we first introduce a weaker variant of Theorem 1.1, showing how to construct ReLU networks to pointwisely approximate continuous functions except for the trifling region.

**Theorem 2.1.** *Given a function  $f \in C([0, 1]^d)$ , for any  $N \in \mathbb{N}^+$  and  $L \in \mathbb{N}^+$ , there exists a function  $\phi$  implemented by a ReLU network with width  $\max\{8d\lfloor N^{1/d} \rfloor + 3d, 16N + 30\}$  and depth  $11L + 18$  such that  $\|\phi\|_{L^\infty(\mathbb{R}^d)} \leq |f(\mathbf{0})| + \omega_f(\sqrt{d})$  and*

$$|f(\mathbf{x}) - \phi(\mathbf{x})| \leq 130\sqrt{d}\omega_f\left(\left(N^2L^2\log_3(N+2)\right)^{-1/d}\right), \quad \text{for any } \mathbf{x} \in [0, 1]^d \setminus \Omega([0, 1]^d, K, \delta),$$

where  $K = \lfloor N^{1/d} \rfloor^2 \lfloor L^{1/d} \rfloor^2 \lfloor \log_3(N+2) \rfloor^{1/d}$  and  $\delta$  is an arbitrary number in  $(0, \frac{1}{3K}]$ .



206 With Theorem 2.1 that will be proved in Section 3, we can easily prove Theorem  
 207 1.1 for the case  $p \in [1, \infty)$ . To attain the rate in  $L^\infty$ -norm, we need to control the  
 208 approximation error in the trifling region. To this end, we introduce a theorem to deal  
 209 with the approximation inside the trifling region  $\Omega([0, 1]^d, K, \delta)$ .

210 **Theorem 2.2** (Theorem 3.7 of [29] or Theorem 2.1 of [15]). *Given any  $\varepsilon > 0$ ,  $N, L, K \in$   
 211  $\mathbb{N}^+$ , and  $\delta \in (0, \frac{1}{3K}]$ , assume  $f$  is a continuous function in  $C([0, 1]^d)$  and  $\tilde{\phi}$  can be  
 212 implemented by a ReLU network with width  $N$  and depth  $L$ . If*

$$213 \quad |f(\mathbf{x}) - \tilde{\phi}(\mathbf{x})| \leq \varepsilon, \quad \text{for any } \mathbf{x} \in [0, 1]^d \setminus \Omega([0, 1]^d, K, \delta),$$

214 *then there exists a function  $\phi$  implemented by a new ReLU network with width  $3^d(N + 4)$   
 215 and depth  $L + 2d$  such that*

$$216 \quad |f(\mathbf{x}) - \phi(\mathbf{x})| \leq \varepsilon + d \cdot \omega_f(\delta), \quad \text{for any } \mathbf{x} \in [0, 1]^d.$$

217 Now we are ready to prove Theorem 1.1 by assuming Theorem 2.1 is true, which  
 218 will be proved later in Section 3.

219 *Proof of Theorem 1.1.* We may assume  $f$  is not a constant function since it is a trivial  
 220 case. Then  $\omega_f(r) > 0$  for any  $r > 0$ . Let us first consider the case  $p \in [1, \infty)$ . Set  
 221  $K = \lfloor N^{1/d} \rfloor^2 \lfloor L^{1/d} \rfloor^2 \lfloor \log_3(N + 2) \rfloor^{1/d}$  and choose a small  $\delta \in (0, \frac{1}{3K}]$  such that

$$222 \quad \begin{aligned} Kd\delta(2|f(\mathbf{0})| + 2\omega_f(\sqrt{d}))^p &= \lfloor N^{1/d} \rfloor^2 \lfloor L^{1/d} \rfloor^2 \lfloor \log_3(N + 2) \rfloor^{1/d} d\delta(2|f(\mathbf{0})| + 2\omega_f(\sqrt{d}))^p \\ &\leq \left( \omega_f\left((N^2 L^2 \log_3(N + 2))^{-1/d}\right) \right)^p. \end{aligned}$$

223 By Theorem 2.1, there exists a function  $\phi$  implemented by a ReLU network with width

$$224 \quad \max\{8d\lfloor N^{1/d} \rfloor + 3d, 16N + 30\} \leq 16 \max\{d\lfloor N^{1/d} \rfloor, N + 2\}$$

225 and depth  $11L + 18$  such that  $\|\phi\|_{L^\infty(\mathbb{R}^d)} \leq |f(\mathbf{0})| + \omega_f(\sqrt{d})$  and

$$226 \quad |f(\mathbf{x}) - \phi(\mathbf{x})| \leq 130\sqrt{d}\omega_f\left((N^2 L^2 \log_3(N + 2))^{-1/d}\right), \quad \text{for any } \mathbf{x} \in [0, 1]^d \setminus \Omega([0, 1]^d, K, \delta),$$

227 It follows from  $\mu(\Omega([0, 1]^d, K, \delta)) \leq Kd\delta$  and  $\|f\|_{L^\infty([0, 1]^d)} \leq |f(\mathbf{0})| + \omega_f(\sqrt{d})$  that

$$\begin{aligned} 228 \quad \|f - \phi\|_{L^p([0, 1]^d)}^p &= \int_{\Omega([0, 1]^d, K, \delta)} |f(\mathbf{x}) - \phi(\mathbf{x})|^p d\mathbf{x} + \int_{[0, 1]^d \setminus \Omega([0, 1]^d, K, \delta)} |f(\mathbf{x}) - \phi(\mathbf{x})|^p d\mathbf{x} \\ &\leq Kd\delta(2|f(\mathbf{0})| + 2\omega_f(\sqrt{d}))^p + \left(130\sqrt{d}\omega_f\left((N^2 L^2 \log_3(N + 2))^{-1/d}\right)\right)^p \\ &\leq \left(\omega_f\left((N^2 L^2 \log_3(N + 2))^{-1/d}\right)\right)^p + \left(130\sqrt{d}\omega_f\left((N^2 L^2 \log_3(N + 2))^{-1/d}\right)\right)^p \\ &\leq \left(131\sqrt{d}\omega_f\left((N^2 L^2 \log_3(N + 2))^{-1/d}\right)\right)^p. \end{aligned}$$



229 Hence,  $\|f - \phi\|_{L^p([0,1]^d)} \leq 131\sqrt{d}\omega_f\left((N^2L^2\log_3(N+2))^{-1/d}\right)$ .

230 Next, let us discuss the case  $p = \infty$ . Set  $K = \lfloor N^{1/d} \rfloor^2 \lfloor L^{1/d} \rfloor^2 \lfloor \log_3(N+2) \rfloor^{1/d}$  and  
 231 choose a small  $\delta \in (0, \frac{1}{3K}]$  such that

$$232 \quad d \cdot \omega_f(\delta) \leq \omega_f\left((N^2L^2\log_3(N+2))^{-1/d}\right).$$

233 By Theorem 2.1, there exists a function  $\tilde{\phi}$  implemented by a ReLU network with width  
 234  $\max\{8d\lfloor N^{1/d} \rfloor + 3d, 16N + 30\}$  and depth  $11L + 18$  such that

$$235 \quad |f(\mathbf{x}) - \tilde{\phi}(\mathbf{x})| \leq 130\sqrt{d}\omega_f\left((N^2L^2\log_3(N+2))^{-1/d}\right) =: \varepsilon,$$

236 for any  $\mathbf{x} \in [0,1]^d \setminus \Omega([0,1]^d, K, \delta)$ . By Theorem 2.2, there exists a function  $\phi$  imple-  
 237 mented by a ReLU network with width

$$238 \quad 3^d \left( \max\{8d\lfloor N^{1/d} \rfloor + 3d, 16N + 30\} + 4 \right) \leq 3^{d+3} \max\{d\lfloor N^{1/d} \rfloor, N + 2\}$$

239 and depth  $11L + 18 + 2d$  such that

$$240 \quad |f(\mathbf{x}) - \phi(\mathbf{x})| \leq \varepsilon + d \cdot \omega_f(\delta) \leq 131\sqrt{d}\omega_f\left((N^2L^2\log_3(N+2))^{-1/d}\right), \quad \text{for any } \mathbf{x} \in [0,1]^d.$$

241 So we finish the proof. □

## 242 2.3 Optimality

243 This section will show that the approximation rates in Theorem 1.1 and Corollary 1.3  
 244 are optimal and there is no room to improve for the function class  $\text{Hölder}([0,1]^d, \alpha, \lambda)$ .  
 245 Therefore, the approximation rate for the whole continuous functions space in terms of  
 246 width and depth in Theorem 1.1 cannot be improved. A typical method to characterize  
 247 the optimal approximation theory of neural networks is to study the connection between  
 248 the approximation error and Vapnik–Chervonenkis (VC) dimension [15, 23, 27, 28, 29].  
 249 This method relies on the VC-dimension upper bound given in [8]. In this paper, we  
 250 adopt this method with several modifications to simplify the proof.

251 Let us first present the definitions of VC-dimension and related concepts. Let  $H$  be  
 252 a class of functions mapping from a general domain  $\mathcal{X}$  to  $\{0,1\}$ . We say  $H$  shatters the  
 253 set  $\{\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_m\} \subseteq \mathcal{X}$  if

$$254 \quad \left| \left\{ [h(\mathbf{x}_1), h(\mathbf{x}_2), \dots, h(\mathbf{x}_m)]^T \in \{0,1\}^m : h \in H \right\} \right| = 2^m,$$

255 where  $|\cdot|$  denotes the size of a set. This equation means, given any  $\theta_i \in \{0,1\}$  for  
 256  $i = 1, 2, \dots, m$ , there exists  $h \in H$  such that  $h(\mathbf{x}_i) = \theta_i$  for all  $i$ .

257 For any  $m \in \mathbb{N}^+$ , we define the growth function of  $H$  as

$$258 \quad \Pi_H(m) := \max_{\mathbf{x}_1, \mathbf{x}_2, \dots, \mathbf{x}_m \in \mathcal{X}} \left| \left\{ [h(\mathbf{x}_1), h(\mathbf{x}_2), \dots, h(\mathbf{x}_m)]^T \in \{0,1\}^m : h \in H \right\} \right|.$$

**Definition 2.3** (VC-dimension). Let  $H$  be a class of functions from  $\mathcal{X}$  to  $\{0, 1\}$ . The VC-dimension of  $H$ , denoted by  $\text{VCDim}(H)$ , is the size of the largest shattered set, namely,  $\text{VCDim}(H) := \sup\{m \in \mathbb{N}^+ : \Pi_H(m) = 2^m\}$ .

Let  $\mathcal{F}$  be a class of functions from  $\mathcal{X}$  to  $\mathbb{R}$ . The VC-dimension of  $\mathcal{F}$ , denoted by  $\text{VCDim}(\mathcal{F})$ , is defined by  $\text{VCDim}(\mathcal{F}) := \text{VCDim}(\mathcal{T} \circ \mathcal{F})$ , where

$$\mathcal{T}(t) := \begin{cases} 1, & t \geq 0, \\ 0, & t < 0 \end{cases} \quad \text{and} \quad \mathcal{T} \circ \mathcal{F} := \{\mathcal{T} \circ f : f \in \mathcal{F}\}.$$

In particular, the expression “VC-dimension of a network (architecture)” means the VC-dimension of the function set that consists of all functions implemented by this network (architecture).

We remark that one may also define  $\text{VCDim}(\mathcal{F})$  as  $\text{VCDim}(\mathcal{F}) := \text{VCDim}(\tilde{\mathcal{T}} \circ \mathcal{F})$ , where

$$\tilde{\mathcal{T}}(t) := \begin{cases} 1, & t > 0, \\ 0, & t \leq 0 \end{cases} \quad \text{and} \quad \tilde{\mathcal{T}} \circ \mathcal{F} := \{\tilde{\mathcal{T}} \circ f : f \in \mathcal{F}\}.$$

Note that function spaces generated by networks are closed under linear transformation. Thus, these two definitions of VC-dimension are equivalent.

The theorem below, Theorem 2.4, reveals the connection between VC-dimension and approximation rate.

**Theorem 2.4.** Assume  $\mathcal{F}$  is a function set with all elements defined on  $[0, 1]^d$ . For any  $\varepsilon \in (0, 2/9)$ , if

$$\inf_{\phi \in \mathcal{F}} \|\phi - f\|_{L^\infty([0, 1]^d)} \leq \varepsilon, \quad \text{for any } f \in \text{H\"older}([0, 1]^d, \alpha, 1), \quad (2.2)$$

then  $\text{VCDim}(\mathcal{F}) \geq (9\varepsilon)^{-d/\alpha}$ .

This theorem demonstrates the connection between VC-dimension of  $\mathcal{F}$  and the approximation rate using elements of  $\mathcal{F}$  to approximate functions in  $\text{H\"older}([0, 1]^d, \alpha, \lambda)$ . To be precise, the VC-dimension of  $\mathcal{F}$  determines an approximation rate lower bound  $\text{VCDim}(\mathcal{F})^{-\alpha/d}/9$ , which is the best possible approximation rate. Denote the best approximation error of functions in  $\text{H\"older}([0, 1]^d, \alpha, 1)$  approximated by ReLU networks with width  $N$  and depth  $L$  as

$$\mathcal{E}_{\alpha, d}(N, L) := \sup_{f \in \text{H\"older}([0, 1]^d, \alpha, 1)} \left( \inf_{\phi \in \mathcal{NN}(\text{width} \leq N; \text{depth} \leq L)} \|\phi - f\|_{L^\infty([0, 1]^d)} \right),$$

We have three remarks listed below.

(i) A large VC-dimension cannot guarantee a good approximation rate. For example, it is easy to verify that

$$\text{VCDim}\left(\{f : f(x) = \cos(ax), a \in \mathbb{R}\}\right) = \infty.$$

However, functions in  $\{f : f(x) = \cos(ax), a \in \mathbb{R}\}$  cannot approximate H\"older continuous functions well.

- (ii) A large VC-dimension is necessary for a good approximation rate, because the best possible approximation rate is controlled by an expression of VC-dimension, as shown in Theorem 2.4. For example, Theorem 6 and 8 of [8] implies that

$$\text{VCDim}(\mathcal{N}(\text{width} \leq N; \text{depth} \leq L)) \leq \min \left\{ \mathcal{O}(N^2 L^2 \ln(NL)), \mathcal{O}(N^3 L^2) \right\},$$

deducing

$$\underbrace{C_1(\alpha, d) \left( \min\{N^2 L^2 \ln(NL), N^3 L^2\} \right)^{-\alpha/d}}_{\text{implied by Theorem 2.4}} \leq \mathcal{E}_{\alpha, d}(N, L) \leq \underbrace{C_2(\alpha, d) \left( N^2 L^2 \ln N \right)^{-\alpha/d}}_{\text{implied by Corollary 1.3}}, \quad \textcircled{1} \quad (2.3)$$

for any  $N, L \in \mathbb{N}^+$  with  $N \geq 2$ , where  $C_1(\alpha, d)$  and  $C_2(\alpha, d)$  are two positive constants determined by  $s, d$ , and  $C_2(s, d)$  can be explicitly expressed.

- When  $L = L_0$  is fixed, Equation 2.3 implies

$$C_1(\alpha, d, L_0) (N^2 \ln N)^{-\alpha/d} \leq \mathcal{E}_{\alpha, d}(N, L_0) \leq C_2(\alpha, d, L_0) (N^2 \ln N)^{-\alpha/d},$$

where  $C_1(\alpha, d, L_0)$  and  $C_2(\alpha, d, L_0)$  are two position constant determined by  $\alpha, d, L_0$ .

- When  $N = N_0$  is fixed, Equation 2.3 implies

$$C_1(\alpha, d, N_0) L^{-2\alpha/d} \leq \mathcal{E}_{\alpha, d}(N_0, L) \leq C_2(\alpha, d, N_0) L^{-2\alpha/d},$$

where  $C_1(\alpha, d, N_0)$  and  $C_2(\alpha, d, N_0)$  are two position constant determined by  $\alpha, d, N_0$ .

- It is easy to verify that Equation (2.3) is tight in

$$\{(N, L) \in \mathbb{N}^+ : C_3(\alpha, d) \leq N \leq L^{C_4(\alpha, d)}\}$$

for two positive constants  $C_3 = C_3(\alpha, d), C_4 = C_4(\alpha, d)$ . See Figure 1 for an illustration for the case  $C_3 = 100$  and  $C_4 = 1/100$ .

Finally, let us present the detailed proof of Theorem 2.4.

*Proof of Theorem 2.4.* Recall that the VC-dimension of a function set is defined as the size of the largest set of points that this class of functions can shatter. So our goal is to find a subset of  $\mathcal{F}$  to shatter  $\mathcal{O}(\varepsilon^{-d/\alpha})$  points in  $[0, 1]^d$ , which can be divided into two steps.

- Construct  $\{f_\chi : \chi \in \mathcal{B}\} \subseteq \text{Hölder}([0, 1]^d, \alpha, 1)$  that scatters  $\mathcal{O}(\varepsilon^{-d/\alpha})$  points, where  $\mathcal{B}$  is a set defined later.
- Design  $\phi_\chi \in \mathcal{F}$ , for each  $\chi \in \mathcal{B}$ , based on  $f_\chi$  and Equation (2.2) such that  $\{\phi_\chi : \chi \in \mathcal{B}\} \subseteq \mathcal{F}$  also shatters  $\mathcal{O}(\varepsilon^{-d/\alpha})$  points.

---

<sup>①</sup>To make this equation valid for any  $N, L \in \mathbb{N}^+$  with  $N \geq 2$ , one needs to choose  $C_1(\alpha, d)$  and  $C_2(\alpha, d)$  carefully based on Theorem 2.4 and Corollary 1.3.

321 The details of these two steps can be found below.

322 **Step 1:** Construct  $\{f_\chi : \chi \in \mathcal{B}\} \subseteq \text{Hölder}([0, 1]^d, \alpha, 1)$  that scatters  $\mathcal{O}(\varepsilon^{-d/\alpha})$  points.

323 Let  $K = \lfloor (9\varepsilon/2)^{-1/\alpha} \rfloor \in \mathbb{N}^+$  and divide  $[0, 1]^d$  into  $K^d$  non-overlapping sub-cubes  
 324  $\{Q_\beta\}_\beta$  as follows:

$$325 \quad Q_\beta := \{\mathbf{x} = [x_1, x_2, \dots, x_d]^T \in [0, 1]^d : x_i \in [\frac{\beta_i}{K}, \frac{\beta_i+1}{K}], \ i = 1, 2, \dots, d\},$$

326 for any index vector  $\beta = [\beta_1, \beta_2, \dots, \beta_d]^T \in \{0, 1, \dots, K-1\}^d$ .

327 Define a function  $\zeta_Q$  on  $[0, 1]^d$  corresponding to  $Q = Q(\mathbf{x}_0, \eta) \subseteq [0, 1]^d$  such that:

- 328 •  $\zeta_Q(\mathbf{x}_0) = (\eta/2)^\alpha/2$ ;
- 329 •  $\zeta_Q(\mathbf{x}) = 0$  for any  $\mathbf{x} \notin Q \setminus \partial Q$ , where  $\partial Q$  is the boundary of  $Q$ ;
- 330 •  $\zeta_Q$  is linear on the line that connects  $\mathbf{x}_0$  and  $\mathbf{x}$  for any  $\mathbf{x} \in \partial Q$ .

331 Define

$$332 \quad \mathcal{B} := \{\chi : \chi \text{ is a map from } \{0, 1, \dots, K-1\}^d \text{ to } \{-1, 1\}\}.$$

333 For each  $\chi \in \mathcal{B}$ , we define

$$334 \quad f_\chi(\mathbf{x}) := \sum_{\beta \in \{0, 1, \dots, K-1\}^d} \chi(\beta) \zeta_{Q_\beta}(\mathbf{x}),$$

335 where  $\zeta_{Q_\beta}(\mathbf{x})$  is the associated function introduced just above. It is easy to check that  
 336  $\{f_\chi : \chi \in \mathcal{B}\} \subseteq \text{Hölder}([0, 1]^d, \alpha, 1)$  can shatter  $K^d = \mathcal{O}(\varepsilon^{-d/\alpha})$  points in  $[0, 1]^d$ .

337 **Step 2:** Construct  $\{\phi_\chi : \chi \in \mathcal{B}\}$  that also scatters  $\mathcal{O}(\varepsilon^{-d/\alpha})$  points.

338 By Equation (2.2), for each  $\chi \in \mathcal{B}$ , there exists  $\phi_\chi \in \mathcal{F}$  such that

$$339 \quad \|\phi_\chi - f_\chi\|_{L^\infty([0, 1]^d)} \leq \varepsilon + \varepsilon/81.$$

340 Let  $\mu(\cdot)$  denote the Lebesgue measure of a set. Then, for each  $\chi \in \mathcal{B}$ , there exists  
 341  $\mathcal{H}_\chi \subseteq [0, 1]^d$  with  $\mu(\mathcal{H}_\chi) = 0$  such that

$$342 \quad |\phi_\chi(\mathbf{x}) - f_\chi(\mathbf{x})| \leq \frac{82}{81}\varepsilon, \quad \text{for any } \mathbf{x} \in [0, 1]^d \setminus \mathcal{H}_\chi.$$

343 Set  $\mathcal{H} = \cup_{\chi \in \mathcal{B}} \mathcal{H}_\chi$ , then we have  $\mu(\mathcal{H}) = 0$  and

$$344 \quad |\phi_\chi(\mathbf{x}) - f_\chi(\mathbf{x})| \leq \frac{82}{81}\varepsilon, \quad \text{for any } \chi \in \mathcal{B} \text{ and } \mathbf{x} \in [0, 1]^d \setminus \mathcal{H}. \quad (2.4)$$

345 Since  $Q_\beta$  has a sidelength  $\frac{1}{K} = \frac{1}{\lfloor (9\varepsilon/2)^{-1/\alpha} \rfloor}$ , we have, for each  $\beta \in \{0, 1, \dots, K-1\}^d$  and  
 346 any  $\mathbf{x} \in \frac{1}{10}Q_\beta$ <sup>②</sup>,

$$347 \quad |f_\chi(\mathbf{x})| = |\zeta_{Q_\beta}(\mathbf{x})| \geq \frac{9}{10}|\zeta_{Q_\beta}(\mathbf{x}_{Q_\beta})| = \frac{9}{10}\left(\frac{1}{2\lfloor (9\varepsilon/2)^{-1/\alpha} \rfloor}\right)^\alpha/2 \geq \frac{81}{80}\varepsilon, \quad (2.5)$$

---

<sup>②</sup>  $\frac{1}{10}Q_\beta$  denotes the closed cube whose sidelength is 1/10 of that of  $Q_\beta$  and which shares the same center of  $Q_\beta$ .

348 where  $\mathbf{x}_{Q_\beta}$  is the center of  $Q_\beta$ .

349 Note that  $(\frac{1}{10}Q_\beta) \setminus \mathcal{H}$  is not empty, since  $\mu((\frac{1}{10}Q_\beta) \setminus \mathcal{H}) > 0$  for each  $\beta \in \{0, 1, \dots, K - 1\}^d$ . Together with Equation (2.4) and (2.5), there exists  $\mathbf{x}_\beta \in (\frac{1}{10}Q_\beta) \setminus \mathcal{H}$  such that, for  
 351 each  $\beta \in \{0, 1, \dots, K - 1\}^d$  and each  $\chi \in \mathcal{B}$ ,

$$352 \quad |f_\chi(\mathbf{x}_\beta)| \geq \frac{81}{80}\varepsilon > \frac{82}{81}\varepsilon \geq |f_\chi(\mathbf{x}_\beta) - \phi_\chi(\mathbf{x}_\beta)|,$$

353 Hence,  $f_\chi(\mathbf{x}_\beta)$  and  $\phi_\chi(\mathbf{x}_\beta)$  have the same sign for each  $\chi \in \mathcal{B}$  and  $\beta \in \{0, 1, \dots, K - 1\}^d$ . Then  $\{\phi_\chi : \chi \in \mathcal{B}\}$  shatters  $\{\mathbf{x}_\beta : \beta \in \{0, 1, \dots, K - 1\}^d\}$  since  $\{f_\chi : \chi \in \mathcal{B}\}$  shatters  
 355  $\{\mathbf{x}_\beta : \beta \in \{0, 1, \dots, K - 1\}^d\}$ . Therefore,

$$356 \quad \text{VCDim}(\mathcal{F}) \geq \text{VCDim}(\{\phi_\chi : \chi \in \mathcal{B}\}) \geq K^d = \lfloor (9\varepsilon/2)^{-1/\alpha} \rfloor^d \geq (9\varepsilon)^{-d/\alpha}, \quad (2.6)$$

357 where the last inequality comes from the fact  $\lfloor x \rfloor \geq x/2 \geq x/(2^{1/\alpha})$  for any  $x \in [1, \infty)$  and  
 358  $\alpha \in (0, 1]$ . So we finish the proof.  $\square$

## 359 2.4 Approximation in irregular domain

360 We extend our analysis to general continuous functions defined on any irregular  
 361 bounded set in  $\mathbb{R}^d$ . The key idea is to extend the target function to a hypercube while  
 362 preserving the modulus of continuity. For a general set  $E \subseteq \mathbb{R}^d$ , the modulus of continuity  
 363 of  $f \in C(E)$  is defined via

$$364 \quad \omega_f^E(r) := \sup \{ |f(\mathbf{x}) - f(\mathbf{y})| : \mathbf{x}, \mathbf{y} \in E, \|\mathbf{x} - \mathbf{y}\|_2 \leq r \}, \quad \text{for any } r \geq 0.$$

365 In particular,  $\omega_f(\cdot)$  is short of  $\omega_f^E(\cdot)$  in the case of  $E = [0, 1]^d$ . Then, Theorem 1.1 can  
 366 be generalized to  $f \in C(E)$  for any bounded set  $E \subseteq [-R, R]^d$  with  $R > 0$ , as shown in  
 367 the following theorem.

368 **Theorem 2.5.** *Given a continuous function  $f \in C(E)$  with  $E \subseteq [-R, R]^d$  and  $R > 0$ , for  
 369 any  $N \in \mathbb{N}^+$ ,  $L \in \mathbb{N}^+$ , and  $p \in [1, \infty]$ , there exists a function  $\phi$  implemented by a ReLU  
 370 network with width  $C_1 \max \{d \lfloor N^{1/d} \rfloor, N + 2\}$  and depth  $11L + C_2$  such that*

$$371 \quad \|f - \phi\|_{L^p(E)} \leq 131(2R)^{d/p} \sqrt{d} \omega_f^E \left( 2R(N^2 L^2 \log_3(N + 2))^{-1/d} \right),$$

372 where  $C_1 = 16$  and  $C_2 = 18$  if  $p \in [1, \infty)$ ;  $C_1 = 3^{d+3}$  and  $C_2 = 18 + 2d$  if  $p = \infty$ .

373 *Proof.* Given any  $f \in C(E)$ , by Lemma 4.2 of [23] via setting  $S = \mathbb{R}^d$ , there exists  
 374  $g \in C(\mathbb{R}^d)$  such that

- 375 •  $g(\mathbf{x}) = f(\mathbf{x})$  for any  $\mathbf{x} \in E \subseteq [-R, R]^d$ ;
- 376 •  $\omega_g^S(r) = \omega_f^E(r)$  for any  $r \geq 0$ .

377 Define

$$378 \quad \tilde{g}(\mathbf{x}) := g(2R\mathbf{x} - R), \quad \text{for any } \mathbf{x} \in \mathbb{R}^d.$$

By applying Theorem 1.1 to  $\tilde{g} \in C([0, 1]^d)$ , there exists a function  $\tilde{\phi}$  implemented by a ReLU network with width  $C_1 \max\{d\lfloor N^{1/d} \rfloor, N + 2\}$  and depth  $11L + C_2$  such that

$$\|\tilde{\phi} - \tilde{g}\|_{L^p([0, 1]^d)} \leq 131\sqrt{d}\omega_{\tilde{g}}\left((N^2L^2\log_3(N + 2))^{-1/d}\right),$$

where  $C_1 = 16$  and  $C_2 = 18$  if  $p \in [1, \infty)$ ;  $C_1 = 3^{d+3}$  and  $C_2 = 18 + 2d$  if  $p = \infty$ .

Recall that  $f(\mathbf{x}) = g(\mathbf{x}) = \tilde{g}(\frac{\mathbf{x}+R}{2R})$  for any  $\mathbf{x} \in E \subseteq [-R, R]^d$  and

$$\omega_{\tilde{g}}(r) \leq \omega_{\tilde{g}}^S(r) = \omega_g^S(2Rr) = \omega_f^E(2Rr), \quad \text{for any } r \geq 0.$$

Define  $\phi(\mathbf{x}) := \tilde{\phi}(\frac{\mathbf{x}+R}{2R}) = \tilde{\phi} \circ \mathcal{L}(\mathbf{x})$  for any  $\mathbf{x} \in \mathbb{R}^d$ , where  $\mathcal{L}$  is an affine linear map given by  $\mathcal{L}(\mathbf{x}) = \frac{\mathbf{x}+R}{2R}$ . Clearly,  $\phi$  can be implemented by a ReLU network with width  $C_1 \max\{d\lfloor N^{1/d} \rfloor, N + 2\}$  and depth  $11L + C_2$ , where  $C_1 = 16$  and  $C_2 = 18$  if  $p \in [1, \infty)$ ;  $C_1 = 3^{d+3}$  and  $C_2 = 18 + 2d$  if  $p = \infty$ . Moreover, for any  $\mathbf{x} \in E \subseteq [-R, R]^d$ , we have  $\frac{\mathbf{x}+R}{2R} \in [0, 1]^d$ , implying

$$\begin{aligned} \|\phi - f\|_{L^p(E)} &= \|\phi - g\|_{L^p(E)} = \|\tilde{\phi} \circ \mathcal{L} - \tilde{g} \circ \mathcal{L}\|_{L^p(E)} \\ &\leq \|\tilde{\phi} \circ \mathcal{L} - \tilde{g} \circ \mathcal{L}\|_{L^p([-R, R]^d)} = (2R)^{d/p} \|\tilde{\phi} - \tilde{g}\|_{L^p([0, 1]^d)} \\ &\leq 131(2R)^{d/p} \sqrt{d}\omega_{\tilde{g}}\left((N^2L^2\log_3(N + 2))^{-1/d}\right) \\ &\leq 131(2R)^{d/p} \sqrt{d}\omega_f^E\left(2R(N^2L^2\log_3(N + 2))^{-1/d}\right). \end{aligned}$$

With the discussion above, we have proved Theorem 2.5.  $\square$

## 3 Proof of Theorem 2.1

We will prove Theorem 2.1 in this section. We first present the key ideas in Section 3.1. The detailed proof is presented in Section 3, based on two propositions in Section 3.1, the proofs of which can be founded in Section 4.

### 3.1 Key ideas of proving Theorem 2.1

Given an arbitrary  $f \in C([0, 1]^d)$ , our goal is to construct an almost piecewise constant function  $\phi$  implemented by a ReLU network to approximate  $f$  well. To this end, we introduce a piecewise constant function  $f_p \approx f$  serving as an intermediate approximant in our construction in the sense that

$$f \approx f_p \text{ on } [0, 1]^d \quad \text{and} \quad f_p \approx \phi \text{ on } [0, 1]^d \setminus \Omega([0, 1]^d, K, \delta).$$

The approximation in  $f \approx f_p$  is a simple and standard technique in constructive approximation. The most technical part is to design a deep ReLU network with the desired width and depth to implement a function  $\phi$  with  $\phi \approx f_p$  outside  $\Omega([0, 1]^d, K, \delta)$ . See Figure 4 for an illustration. The introduction of the trifling region is to ease the construction of  $\phi$ , which is a continuous piecewise linear function, to approximate the discontinuous function  $f_p$  by removing the difficulty near discontinuous points, essentially smoothing  $f_p$  by restricting the approximation domain in  $[0, 1]^d \setminus \Omega([0, 1]^d, K, \delta)$ .

Now let us discuss the detailed steps of construction.

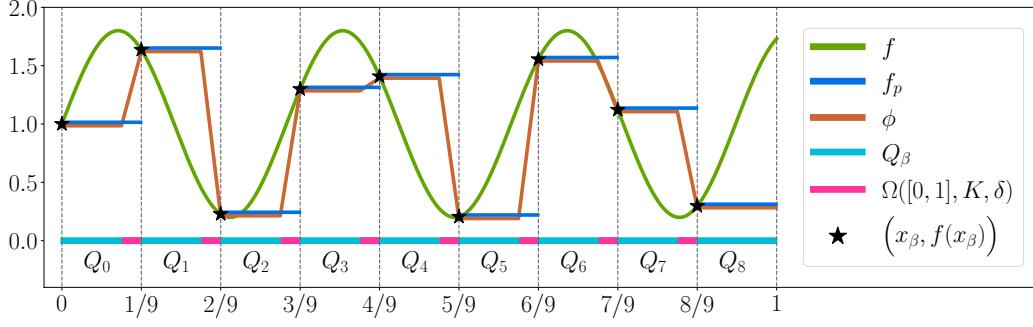


Figure 4: An illustration of  $f$ ,  $f_p$ ,  $\phi$ ,  $x_\beta$ ,  $Q_\beta$ , and the trifling region  $\Omega([0, 1]^d, K, \delta)$  in the one-dimensional case for  $\beta \in \{0, 1, \dots, K-1\}^d$ , where  $K = N^2 L^2 \log_3(N+2)$  and  $d = 1$  with  $N = 1$  and  $L = 3$ .  $f$  is the target function;  $f_p$  is the piecewise constant function approximating  $f$ ;  $\phi$  is a function, implemented by a ReLU network, approximating  $f$ ; and  $x_\beta$  is a representative of  $Q_\beta$ . The measure of  $\Omega([0, 1]^d, K, \delta)$  can be arbitrarily small as we shall see in the proof of Theorem 1.1.

(1) First, divide  $[0, 1]^d$  into a union of important regions  $\{Q_\beta\}_\beta$  and the trifling region  $\Omega([0, 1]^d, K, \delta)$ , where each  $Q_\beta$  is associated with a representative  $\mathbf{x}_\beta \in Q_\beta$  such that  $f(\mathbf{x}_\beta) = f_p(\mathbf{x}_\beta)$  for each index vector  $\beta \in \{0, 1, \dots, K-1\}^d$ , where  $K = \mathcal{O}((N^2 L^2 \ln N)^{1/d})$  is the partition number per dimension (see Figure 6 for examples for  $d = 1$  and  $d = 2$ ).

(2) Next, we design a vector function  $\Phi_1(\mathbf{x})$  constructed via

$$\Phi_1(\mathbf{x}) = [\phi_1(x_1), \phi_1(x_2), \dots, \phi_1(x_d)]^T$$

to project the whole cube  $Q_\beta$  to a  $d$ -dimensional index  $\beta$  for each  $\beta$ , where each one-dimensional function  $\phi_1$  is a step function implemented by a ReLU network.

(3) The third step is to solve a point fitting problem. To be precise, we construct a function  $\phi_2$  implemented by a ReLU network to map  $\beta$  approximately to  $f_p(\mathbf{x}_\beta) = f(\mathbf{x}_\beta)$ . Then  $\phi_2 \circ \Phi_1(\mathbf{x}) = \phi_2(\beta) \approx f_p(\mathbf{x}_\beta) = f(\mathbf{x}_\beta)$  for any  $\mathbf{x} \in Q_\beta$  and each  $\beta$ , implying  $\phi := \phi_2 \circ \Phi_1 \approx f_p \approx f$  on  $[0, 1]^d \setminus \Omega([0, 1]^d, K, \delta)$ . We would like to point out that we only need to care about the values of  $\phi_2$  at a set of points  $\{0, 1, \dots, K-1\}^d$  in the construction of  $\phi_2$  according to our design  $\phi = \phi_2 \circ \Phi_1$  as illustrated in Figure 5. Therefore, it is not necessary to care about the values of  $\phi_2$  sampled outside the set  $\{0, 1, \dots, K-1\}^d$ , which is a key point to ease the design of a ReLU network to implement  $\phi_2$  as we shall see later.

Finally, we discuss how to implement  $\Phi_1$  and  $\phi_2$  by deep ReLU networks with width  $\mathcal{O}(N)$  and depth  $\mathcal{O}(L)$  using two propositions as we shall prove in Section 4.2 and 4.3 later. We first construct a ReLU network with desired width and depth by Proposition 3.1 to implement a one-dimensional step function  $\phi_1$ . Then  $\Phi_1$  can be attained via defining

$$\Phi_1(\mathbf{x}) = [\phi_1(x_1), \phi_1(x_2), \dots, \phi_1(x_d)]^T, \quad \text{for any } \mathbf{x} = [x_1, x_2, \dots, x_d]^T \in \mathbb{R}^d.$$



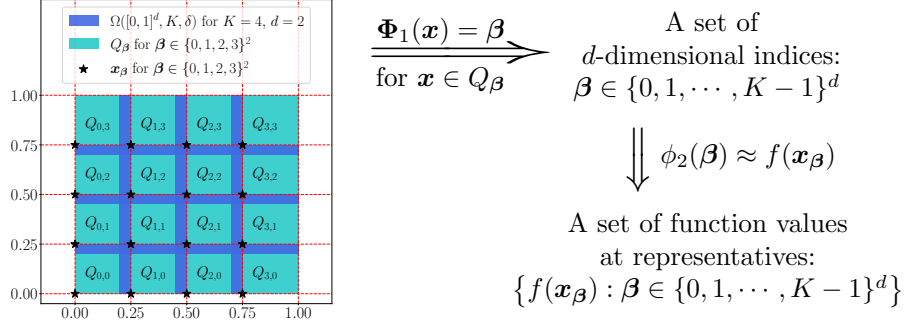


Figure 5: An illustration of the desired function  $\phi = \phi_2 \circ \Phi_1$ . Note that  $\phi \approx f$  on  $[0,1]^d \setminus \Omega([0,1]^d, K, \delta)$ , since  $\phi(\mathbf{x}) = \phi_2 \circ \Phi_1(\mathbf{x}) = \phi_2(\beta) \approx f(\mathbf{x}_{\beta})$  for any  $\mathbf{x} \in Q_{\beta}$  and each  $\beta \in \{0,1,\dots,K-1\}^d$ .

**Proposition 3.1.** For any  $N, L, d \in \mathbb{N}^+$  and  $\delta \in (0, \frac{1}{3K}]$  with

$$K = \lfloor N^{1/d} \rfloor^2 \lfloor L^{2/d} \rfloor \lfloor n^{1/d} \rfloor, \quad \text{where } n = \lfloor \log_3(N+2) \rfloor,$$

there exists a one-dimensional function  $\phi$  implemented by a ReLU network with width  $8\lfloor N^{1/d} \rfloor + 3$  and depth  $2\lfloor L^{1/d} \rfloor + 5$  such that

$$\phi(x) = k, \quad \text{if } x \in \left[\frac{k}{K}, \frac{k+1}{K} - \delta \cdot 1_{\{k \leq K-2\}}\right] \text{ for } k = 0, 1, \dots, K-1.$$

The setting  $K = \lfloor N^{1/d} \rfloor^2 \lfloor L^{1/d} \rfloor^2 \lfloor n^{1/d} \rfloor = \mathcal{O}(N^{2/d} L^{2/d} n^{1/d})$  is not neat here, but it is very convenient for later use. The construction of  $\phi_2$  is a direct result of Proposition 3.2 below, the proof of which relies on the bit extraction technique in [3].

**Proposition 3.2.** Given any  $\varepsilon > 0$  and arbitrary  $N, L, J \in \mathbb{N}^+$  with  $J \leq N^2 L^2 \lfloor \log_3(N+2) \rfloor$ , assume  $y_j \geq 0$  for  $j = 0, 1, \dots, J-1$  are samples with

$$|y_j - y_{j-1}| \leq \varepsilon, \quad \text{for } j = 1, 2, \dots, J-1.$$

Then there exists  $\phi \in \mathcal{NN}(\#input = 1; \text{width} \leq 16N + 30; \text{depth} \leq 6L + 10; \#output = 1)$  such that

$$(i) \quad |\phi(j) - y_j| \leq \varepsilon \text{ for } j = 0, 1, \dots, J-1.$$

$$(ii) \quad 0 \leq \phi(x) \leq \max\{y_j : j = 0, 1, \dots, J-1\} \text{ for any } x \in \mathbb{R}.$$

With the above propositions ready, let us prove Theorem 2.1 in Section 3.

### 3.2 Constructive proof

We essentially construct an almost piecewise constant function implemented by a ReLU network with width  $\mathcal{O}(N)$  and depth  $\mathcal{O}(L)$  to approximate  $f$ . We may assume  $f$  is not a constant function since it is a trivial case. Then  $\omega_f(r) > 0$  for any  $r > 0$ . It is clear that  $|f(\mathbf{x}) - f(\mathbf{0})| \leq \omega_f(\sqrt{d})$  for any  $\mathbf{x} \in [0,1]^d$ . Define  $\tilde{f} = f - f(\mathbf{0}) + \omega_f(\sqrt{d})$ , then  $0 \leq \tilde{f}(\mathbf{x}) \leq 2\omega_f(\sqrt{d})$  for any  $\mathbf{x} \in [0,1]^d$ .

Let  $M = N^2 L$ ,  $n = \lfloor \log_3(N+2) \rfloor$ ,  $K = \lfloor N^{1/d} \rfloor^2 \lfloor L^{1/d} \rfloor^2 \lfloor n^{1/d} \rfloor$ , and  $\delta$  be an arbitrary number in  $(0, \frac{1}{3K}]$ . The proof can be divided into four steps as follows:

1. Normalize  $f$  as  $\tilde{f}$ , divide  $[0, 1]^d$  into a union of sub-cubes  $\{Q_\beta\}_{\beta \in \{0,1,\dots,K-1\}^d}$  and the trifling region  $\Omega([0, 1]^d, K, \delta)$ , and denote  $\mathbf{x}_\beta$  as the vertex of  $Q_\beta$  with minimum  $\|\cdot\|_1$  norm;
2. Construct a sub-network to implement a vector function  $\Phi_1$  projecting the whole cube  $Q_\beta$  to the  $d$ -dimensional index  $\beta$  for each  $\beta$ , i.e.,  $\Phi_1(\mathbf{x}) = \beta$  for all  $\mathbf{x} \in Q_\beta$ ;
3. Construct a sub-network to implement a function  $\phi_2$  mapping the index  $\beta$  approximately to  $\tilde{f}(\mathbf{x}_\beta)$ . This core step can be further divided into three sub-steps:
  - 3.1. Construct a sub-network to implement  $\psi_1$  bijectively mapping the index set  $\{0, 1, \dots, K-1\}^d$  to an auxiliary set  $\mathcal{A}_1 \subseteq \{\frac{j}{2K^d} : j = 0, 1, \dots, 2K^d\}$  defined later (see Figure 7 for an illustration);
  - 3.2. Determine a continuous piecewise linear function  $g$  with a set of breakpoints  $\mathcal{A}_1 \cup \mathcal{A}_2 \cup \{1\}$  satisfying: 1) assign the values of  $g$  at breakpoints in  $\mathcal{A}_1$  based on  $\{\tilde{f}(\mathbf{x}_\beta)\}_\beta$ , i.e.,  $g \circ \psi_1(\beta) = \tilde{f}(\mathbf{x}_\beta)$ ; 2) assign the values of  $g$  at breakpoints in  $\mathcal{A}_2 \cup \{1\}$  to reduce the variation of  $g$  for applying Proposition 3.2;
  - 3.3. Apply Proposition 3.2 to construct a sub-network to implement a function  $\psi_2$  approximating  $g$  well on  $\mathcal{A}_1 \cup \mathcal{A}_2 \cup \{1\}$ . Then the desired function  $\phi_2$  is given by  $\phi_2 = \psi_2 \circ \psi_1$  satisfying  $\phi_2(\beta) = \psi_2 \circ \psi_1(\beta) \approx g \circ \psi_1(\beta) = \tilde{f}(\mathbf{x}_\beta)$ ;
4. Construct the final target network to implement the desired function  $\phi$  such that  $\phi(\mathbf{x}) = \phi_2 \circ \Phi_1(\mathbf{x}) + f(\mathbf{0}) - \omega_f(\sqrt{d}) \approx \tilde{f}(\mathbf{x}_\beta) + f(\mathbf{0}) - \omega_f(\sqrt{d}) = f(\mathbf{x}_\beta)$  for  $\mathbf{x} \in Q_\beta$ .

The details of these steps can be found below.

**Step 1:** Divide  $[0, 1]^d$  into  $\{Q_\beta\}_{\beta \in \{0,1,\dots,K-1\}^d}$  and  $\Omega([0, 1]^d, K, \delta)$ .

Define  $\mathbf{x}_\beta := \beta/K$  and

$$Q_\beta := \left\{ \mathbf{x} = [x_1, \dots, x_d]^T \in [0, 1]^d : x_i \in \left[ \frac{\beta_i}{K}, \frac{\beta_i+1}{K} - \delta \cdot 1_{\{\beta_i \leq K-2\}} \right], i = 1, \dots, d \right\}$$

for each  $d$ -dimensional index  $\beta = [\beta_1, \dots, \beta_d]^T \in \{0, 1, \dots, K-1\}^d$ . Recall that  $\Omega([0, 1]^d, K, \delta)$  is the trifling region defined in Equation (2.1). Apparently,  $\mathbf{x}_\beta$  is the vertex of  $Q_\beta$  with minimum  $\|\cdot\|_1$  norm and

$$[0, 1]^d = \left( \cup_{\beta \in \{0,1,\dots,K-1\}^d} Q_\beta \right) \cup \Omega([0, 1]^d, K, \delta),$$

see Figure 6 for illustrations.

**Step 2:** Construct  $\Phi_1$  mapping  $\mathbf{x} \in Q_\beta$  to  $\beta$ .

By Proposition 3.1, there exists  $\phi_1 \in \mathcal{NN}(\text{width} \leq 8\lfloor N^{1/d} \rfloor + 3; \text{depth} \leq 2\lfloor L^{1/d} \rfloor + 5)$  such that

$$\phi_1(x) = k, \quad \text{if } x \in \left[ \frac{k}{K}, \frac{k+1}{K} - \delta \cdot 1_{\{k \leq K-2\}} \right] \text{ for } k = 0, 1, \dots, K-1.$$

It follows that  $\phi_1(x_i) = \beta_i$  if  $\mathbf{x} = [x_1, x_2, \dots, x_d]^T \in Q_\beta$  for each  $\beta = [\beta_1, \beta_2, \dots, \beta_d]^T$ .

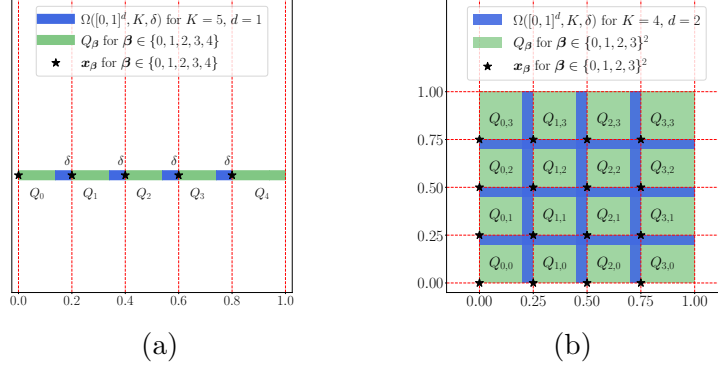


Figure 6: Illustrations of  $\Omega([0, 1]^d, K, \delta)$ ,  $Q_\beta$ , and  $\mathbf{x}_\beta$  for  $\beta \in \{0, 1, \dots, K-1\}^d$ . (a)  $K = 5$  and  $d = 1$ . (b)  $K = 4$  and  $d = 2$ .

By defining

$$\Phi_1(\mathbf{x}) := [\phi_1(x_1), \phi_1(x_2), \dots, \phi_1(x_d)]^T, \quad \text{for any } \mathbf{x} = [x_1, x_2, \dots, x_d]^T \in \mathbb{R}^d,$$

we have  $\Phi_1(\mathbf{x}) = \beta$  if  $\mathbf{x} \in Q_\beta$  for  $\beta \in \{0, 1, \dots, K-1\}^d$ .

**Step 3:** Construct  $\phi_2$  mapping  $\beta$  approximately to  $\tilde{f}(\mathbf{x}_\beta)$ .

The construction of the sub-network implementing  $\phi_2$  is essentially based on Proposition 3.2. To meet the requirements of applying Proposition 3.2, we first define two auxiliary set  $\mathcal{A}_1$  and  $\mathcal{A}_2$  as

$$\mathcal{A}_1 := \left\{ \frac{i}{K^{d-1}} + \frac{k}{2K^d} : i = 0, 1, \dots, K^{d-1}-1 \quad \text{and} \quad k = 0, 1, \dots, K-1 \right\}$$

and

$$\mathcal{A}_2 := \left\{ \frac{i}{K^{d-1}} + \frac{K+k}{2K^d} : i = 0, 1, \dots, K^{d-1}-1 \quad \text{and} \quad k = 0, 1, \dots, K-1 \right\}.$$

Clearly,  $\mathcal{A}_1 \cup \mathcal{A}_2 \cup \{1\} = \left\{ \frac{j}{2K^d} : j = 0, 1, \dots, 2K^d \right\}$  and  $\mathcal{A}_1 \cap \mathcal{A}_2 = \emptyset$ . See Figure 6 for an illustration of  $\mathcal{A}_1$  and  $\mathcal{A}_2$ . Next, we further divide this step into three sub-steps.

**Step 3.1:** Construct  $\psi_1$  bijectively mapping  $\{0, 1, \dots, K-1\}^d$  to  $\mathcal{A}_1$ .

Inspired by the binary representation, we define

$$\psi_1(\mathbf{x}) := \frac{x_d}{2K^d} + \sum_{i=1}^{d-1} \frac{x_i}{K^i}, \quad \text{for any } \mathbf{x} = [x_1, x_2, \dots, x_d]^T \in \mathbb{R}^d. \quad (3.1)$$

Then  $\psi_1$  is a linear function bijectively mapping the index set  $\{0, 1, \dots, K-1\}^d$  to

$$\begin{aligned} & \left\{ \frac{\beta_d}{2K^d} + \sum_{i=1}^{d-1} \frac{\beta_i}{K^i} : \beta \in \{0, 1, \dots, K-1\}^d \right\} \\ &= \left\{ \frac{i}{K^{d-1}} + \frac{k}{2K^d} : i = 0, 1, \dots, K^{d-1}-1 \quad \text{and} \quad k = 0, 1, \dots, K-1 \right\} = \mathcal{A}_1. \end{aligned}$$

**Step 3.2:** Construct  $g$  to satisfy  $g \circ \psi_1(\beta) = \tilde{f}(\mathbf{x}_\beta)$  and to meet the requirements of applying Proposition 3.2.

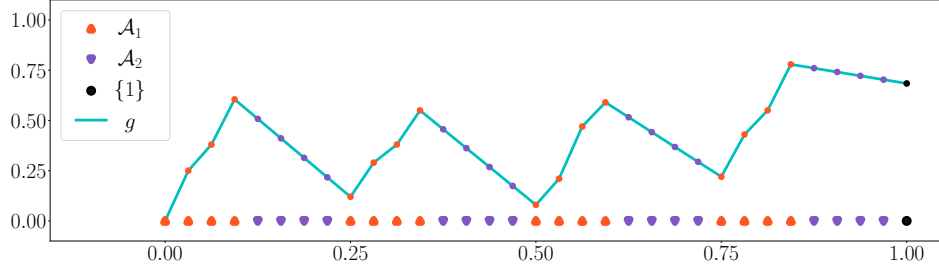


Figure 7: An illustration of  $\mathcal{A}_1$ ,  $\mathcal{A}_2$ ,  $\{1\}$ , and  $g$  for  $d = 2$  and  $K = 4$ .

Let  $g : [0, 1] \rightarrow \mathbb{R}$  be a continuous piecewise linear function with a set of breakpoints  $\{\frac{j}{2K^d} : j = 0, 1, \dots, 2K^d\} = \mathcal{A}_1 \cup \mathcal{A}_2 \cup \{1\}$  and the values of  $g$  at these breakpoints satisfy the following properties:

- The values of  $g$  at the breakpoints in  $\mathcal{A}_1$  are set as

$$g(\psi_1(\beta)) = \tilde{f}(\mathbf{x}_\beta), \quad \text{for any } \beta \in \{0, 1, \dots, K-1\}^d; \quad (3.2)$$

- At the breakpoint 1, let  $g(1) = \tilde{f}(\mathbf{1})$ , where  $\mathbf{1} = [1, 1, \dots, 1]^T \in \mathbb{R}^d$ ;
- The values of  $g$  at the breakpoints in  $\mathcal{A}_2$  are assigned to reduce the variation of  $g$ , which is a requirement of applying Proposition 3.2. Note that

$$\left\{ \frac{i}{K^{d-1}} - \frac{K+1}{2K^d}, \frac{i}{K^{d-1}} \right\} \subseteq \mathcal{A}_1 \cup \{1\}, \quad \text{for } i = 1, 2, \dots, K^{d-1},$$

implying the values of  $g$  at  $\frac{i}{K^{d-1}} - \frac{K+1}{2K^d}$  and  $\frac{i}{K^{d-1}}$  have been assigned for  $i = 1, 2, \dots, K^{d-1}$ . Thus, the values of  $g$  at the breakpoints in  $\mathcal{A}_2$  can be successfully assigned by letting  $g$  linear on each interval  $[\frac{i}{K^{d-1}} - \frac{K+1}{2K^d}, \frac{i}{K^{d-1}}]$  for  $i = 1, 2, \dots, K^{d-1}$ , since  $\mathcal{A}_2 \subseteq \cup_{i=1}^{K^{d-1}} [\frac{i}{K^{d-1}} - \frac{K+1}{2K^d}, \frac{i}{K^{d-1}}]$ .

Apparently, such a function  $g$  exists (see Figure 7 for an example) and satisfies

$$\left| g\left(\frac{j}{2K^d}\right) - g\left(\frac{j-1}{2K^d}\right) \right| \leq \max \left\{ \omega_f\left(\frac{1}{K}\right), \omega_f(\sqrt{d})/K \right\} \leq \omega_f\left(\frac{\sqrt{d}}{K}\right), \quad \text{for } j = 1, 2, \dots, 2K^d,$$

and

$$0 \leq g\left(\frac{j}{2K^d}\right) \leq 2\omega_f(\sqrt{d}), \quad \text{for } j = 0, 1, \dots, 2K^d.$$

**Step 3.3:** Construct  $\psi_2$  approximating  $g$  well on  $\mathcal{A}_1 \cup \mathcal{A}_2 \cup \{1\}$ .

Note that

$$2K^d = 2\left(\lfloor N^{1/d} \rfloor^2 \lfloor L^{1/d} \rfloor^2 \lfloor n^{1/d} \rfloor\right)^d \leq 2(N^2 L^2 n) \leq N^2 \lceil \sqrt{2}L \rceil^2 \lceil \log_3(N+2) \rceil.$$

By Proposition 3.2 (set  $y_j = g(\frac{j}{2K^2})$  and  $\varepsilon = \omega_f(\frac{\sqrt{d}}{K}) > 0$  therein), there exists

$$\tilde{\psi}_2 \in \mathcal{NN}(\#input = 1; \text{ width } \leq 16N + 30; \text{ depth } \leq 6\lceil \sqrt{2}L \rceil + 10; \#output = 1)$$

532 such that

$$533 \quad |\tilde{\psi}_2(j) - g(\frac{j}{2K^d})| \leq \omega_f(\frac{\sqrt{d}}{K}), \quad \text{for } j = 0, 1, \dots, 2K^d - 1,$$

534 and

$$535 \quad 0 \leq \tilde{\psi}_2(x) \leq \max\{g(\frac{j}{2K^d}) : j = 0, 1, \dots, 2K^d - 1\} \leq 2\omega_f(\sqrt{d}), \quad \text{for any } x \in \mathbb{R}.$$

536 By defining  $\psi_2(x) := \tilde{\psi}_2(2K^d x)$  for any  $x \in \mathbb{R}$ , we have  $\psi_2 \in \mathcal{NN}(\#input = 1; \text{width} \leq$   
 537  $16N + 30; \text{depth} \leq 6\lceil\sqrt{2}L\rceil + 10; \#output = 1)$ ,

$$538 \quad 0 \leq \psi_2(x) = \tilde{\psi}_2(2K^d x) \leq 2\omega_f(\sqrt{d}), \quad \text{for any } x \in \mathbb{R}, \quad (3.3)$$

539 and

$$540 \quad |\psi_2(\frac{j}{2K^d}) - g(\frac{j}{2K^d})| = |\tilde{\psi}_2(j) - g(\frac{j}{2K^d})| \leq \omega_f(\frac{\sqrt{d}}{K}), \quad \text{for } j = 0, 1, \dots, 2K^d - 1. \quad (3.4)$$

541 Let us end Step 3 by defining the desired function  $\phi_2$  as  $\phi_2 := \psi_2 \circ \psi_1$ . Note that  
 542  $\psi_1 : \mathbb{R}^d \rightarrow \mathbb{R}$  is a linear function and  $\psi_2 \in \mathcal{NN}(\#input = 1; \text{width} \leq 16N + 30; \text{depth} \leq$   
 543  $6\lceil\sqrt{2}L\rceil + 10; \#output = 1)$ . Thus,  $\phi_2 \in \mathcal{NN}(\#input = 1; \text{width} \leq 16N + 30; \text{depth} \leq$   
 544  $6\lceil\sqrt{2}L\rceil + 10; \#output = 1)$ . By Equation (3.2) and (3.4), we have

$$545 \quad |\phi_2(\beta) - \tilde{f}(\mathbf{x}_\beta)| = |\psi_2(\psi_1(\beta)) - g(\psi_1(\beta))| \leq \omega_f(\frac{\sqrt{d}}{K}), \quad (3.5)$$

546 for any  $\beta \in \{0, 1, \dots, K - 1\}^d$ . Equation (3.3) and  $\phi_2 = \psi_2 \circ \psi_1$  implies

$$547 \quad 0 \leq \phi_2(\mathbf{x}) \leq 2\omega_f(\sqrt{d}), \quad \text{for any } \mathbf{x} \in \mathbb{R}^d. \quad (3.6)$$

548 **Step 4:** Construct the final network to implement the desired function  $\phi$ .

549 Define  $\phi := \phi_2 \circ \Phi_1 + f(\mathbf{0}) - \omega_f(\sqrt{d})$ . Since  $\phi_1 \in \mathcal{NN}(\text{width} \leq 8\lceil N^{1/d} \rceil + 3; \text{depth} \leq$   
 550  $2\lceil L^{1/d} \rceil + 5)$ , we have  $\Phi_1 \in \mathcal{NN}(\#input = d; \text{width} \leq 8d\lceil N^{1/d} \rceil + 3d; \text{depth} \leq 2L +$   
 551  $5; \#output = d)$ . It follows from the fact  $\lceil\sqrt{2}L\rceil \leq \lceil\frac{3}{2}L\rceil \leq \frac{3}{2}L + \frac{1}{2}$  that  $6\lceil\sqrt{2}L\rceil + 10 \leq 9L + 13$ ,  
 552 implying

$$553 \quad \begin{aligned} \phi_2 &\in \mathcal{NN}(\#input = 1; \text{width} \leq 16N + 30; \text{depth} \leq 6\lceil\sqrt{2}L\rceil + 10; \#output = 1) \\ &\subseteq \mathcal{NN}(\#input = 1; \text{width} \leq 16N + 30; \text{depth} \leq 9L + 13; \#output = 1). \end{aligned}$$

554 Thus,  $\phi = \phi_2 \circ \Phi_1 + f(\mathbf{0}) - \omega_f(\sqrt{d})$  is in

$$555 \quad \mathcal{NN}(\text{width} \leq \max\{8d\lceil N^{1/d} \rceil + 3d, 16N + 30\}; \text{depth} \leq (2L + 5) + (9L + 13) = 11L + 18).$$

556 Now let us estimate the approximation error. Note that  $f = \tilde{f} + f(\mathbf{0}) - \omega_f(\sqrt{d})$ . By  
 557 Equation (3.5), for any  $\mathbf{x} \in Q_\beta$  and  $\beta \in \{0, 1, \dots, K - 1\}^d$ , we have

$$\begin{aligned} |f(\mathbf{x}) - \phi(\mathbf{x})| &= |\tilde{f}(\mathbf{x}) - \phi_2(\Phi_1(\mathbf{x}))| = |\tilde{f}(\mathbf{x}) - \phi_2(\beta)| \\ &\leq |\tilde{f}(\mathbf{x}) - \tilde{f}(\mathbf{x}_\beta)| + |\tilde{f}(\mathbf{x}_\beta) - \phi_2(\beta)| \\ 558 \quad &\leq \omega_f(\frac{\sqrt{d}}{K}) + \omega_f(\frac{\sqrt{d}}{K}) \leq 2\omega_f\left(64\sqrt{d}\left(N^2L^2\log_3(N+2)\right)^{-1/d}\right), \end{aligned}$$

where the last inequality comes from the fact

$$K = \lfloor N^{1/d} \rfloor^2 \lfloor L^{1/d} \rfloor^2 \lfloor n^{1/d} \rfloor \geq \frac{N^{2/d} L^{2/d} n^{1/d}}{32} = \frac{N^{2/d} L^{2/d} \lfloor \log_3(N+2) \rfloor^{1/d}}{32} \geq \frac{(N^2 L^2 \log_3(N+2))^{1/d}}{64},$$

for any  $N, L \in \mathbb{N}^+$ . Recall the fact  $\omega_f(j \cdot r) \leq j \cdot \omega_f(r)$  for any  $j \in \mathbb{N}^+$  and  $r \in [0, \infty)$ .

Therefore, for any  $\mathbf{x} \in \bigcup_{\beta \in \{0,1,\dots,K-1\}^d} Q_\beta = [0,1]^d \setminus \Omega([0,1]^d, K, \delta)$ , we have

$$\begin{aligned} |f(\mathbf{x}) - \phi(\mathbf{x})| &\leq 2\omega_f\left(64\sqrt{d}(N^2 L^2 \log_3(N+2))^{-1/d}\right) \\ &\leq 2\lfloor 64\sqrt{d} \rfloor \omega_f\left((N^2 L^2 \log_3(N+2))^{-1/d}\right) \\ &\leq 130\sqrt{d} \omega_f\left((N^2 L^2 \log_3(N+2))^{-1/d}\right). \end{aligned}$$

It remains to show the upper bound of  $\phi$ . By Equation (3.6) and  $\phi = \phi_2 \circ \Phi_1 + f(\mathbf{0}) - \omega_f(\sqrt{d})$ , it holds that  $\|\phi\|_{L^\infty(\mathbb{R}^d)} \leq |f(\mathbf{0})| + \omega_f(\sqrt{d})$ . Thus, we finish the proof.

## 4 Proofs of propositions in Section 3

In this section, we will prove the propositions in Section 3. We first introduce several basic results of ReLU networks. Next, we prove Proposition 3.1 and 3.2 based on these basic results.

### 4.1 Basic results of ReLU networks

To simplify the proofs of two propositions in Section 3, we introduce three lemmas below, which are basic results of ReLU networks

**Lemma 4.1.** *For any  $N_1, N_2 \in \mathbb{N}^+$ , given  $N_1(N_2 + 1) + 1$  samples  $(x_i, y_i) \in \mathbb{R}^2$  with  $x_0 < x_1 < \dots < x_{N_1(N_2+1)}$  and  $y_i \geq 0$  for  $i = 0, 1, \dots, N_1(N_2+1)$ , there exists  $\phi \in \mathcal{NN}(\#input = 1; \text{widthvec} = [2N_1, 2N_2 + 1]; \#output = 1)$  satisfying the following conditions.*

- (i)  $\phi(x_i) = y_i$  for  $i = 0, 1, \dots, N_1(N_2 + 1)$ .
- (ii)  $\phi$  is linear on each interval  $[x_{i-1}, x_i]$  for  $i \notin \{(N_2 + 1)j : j = 1, 2, \dots, N_1\}$ .

**Lemma 4.2.** *Given any  $N, L, d \in \mathbb{N}^+$ , it holds that*

$$\begin{aligned} &\mathcal{NN}(\#input = d; \text{widthvec} = [N, NL]; \#output = 1) \\ &\subseteq \mathcal{NN}(\#input = d; \text{width} \leq 2N + 2; \text{depth} \leq L + 1; \#output = 1). \end{aligned}$$

Lemma 4.1 is a part of Theorem 3.2 in [29] or Lemma 2.2 in [22]. Lemma 4.1 is Theorem 3.1 in [29] or Lemma 3.4 in [22].

**Lemma 4.3.** *For any  $n \in \mathbb{N}^+$ , it holds that*

$$\text{CPwL}(\mathbb{R}, n) \subseteq \mathcal{NN}(\#input = 1; \text{widthvec} = [n + 1]; \#output = 1). \quad (4.1)$$

584 *Proof.* We use the mathematics induction to prove Equation (4.1). First, consider the  
 585 case  $n = 1$ . Given any  $f \in \text{CPwL}(\mathbb{R}, n)$ , there exist  $a_1, a_2, x_0 \in \mathbb{R}$  such that

$$586 \quad f(x) = \begin{cases} a_1(x - x_0) + f(x_0), & \text{if } x \geq x_0, \\ a_2(x_0 - x) + f(x_0), & \text{if } x < x_0. \end{cases}$$

587 Thus,  $f(x) = a_1\sigma(x - x_0) + a_2\sigma(x_0 - x) + f(x_0)$  for any  $x \in \mathbb{R}$ , implying  $f \in \mathcal{NN}(\#input =$   
 588  $1; \text{widthvec} = [2]; \#output = 1)$ . Thus, Equation (4.1) holds for  $n = 1$ .

589 Now assume Equation (4.1) holds for  $n = k \in \mathbb{N}^+$ , we would like to show it is also  
 590 true for  $n = k + 1$ . Given any  $f \in \text{CPwL}(\mathbb{R}, k + 1)$ , we may assume the biggest breakpoint  
 591 of  $f$  is  $x_0$  since it is trivial for the case that  $f$  has no breakpoint. Denote the slopes of  
 592 the linear pieces left and right next to  $x_0$  by  $a_1$  and  $a_2$ , respectively. Define

$$593 \quad \tilde{f}(x) := f(x) - (a_2 - a_1)\sigma(x - x_0), \quad \text{for any } x \in \mathbb{R}.$$

594 Then  $\tilde{f}$  has at most  $k$  breakpoints. By the induction hypothesis, we have

$$595 \quad \tilde{f} \in \text{CPwL}(\mathbb{R}, k) \subseteq \mathcal{NN}(\#input = 1; \text{widthvec} = [k + 1]; \#output = 1).$$

596 Thus, there exist  $w_{0,j}, b_{0,j}, w_{1,j}, b_1$  for  $j = 1, 2, \dots, k + 1$  such that

$$597 \quad \tilde{f}(x) = \sum_{j=1}^{k+1} w_{1,j}\sigma(w_{0,j}x + b_{0,j}) + b_1, \quad \text{for any } x \in \mathbb{R}.$$

598 Therefore, for any  $x \in \mathbb{R}$ , we have

$$599 \quad f(x) = (a_2 - a_1)\sigma(x - x_0) + \tilde{f}(x) = (a_2 - a_1)\sigma(x - x_0) + \sum_{j=1}^k w_{1,j}\sigma(w_{0,j}x + b_{0,j}) + b_1,$$

600 implying  $f \in \mathcal{NN}(\#input = 1; \text{widthvec} = [k + 1]; \#output = 1)$ . Thus, Equation (4.1)  
 601 holds for  $k + 1$ , which means we finish the induction process. So we complete the proof.  $\square$

## 602 4.2 Proof of Proposition 3.1

603 Now, let us present the detailed proof of Proposition 3.1. Denote  $K = \widetilde{M} \cdot \widetilde{L}$ , where  
 604  $\widetilde{M} = \lfloor N^{1/d} \rfloor^2 \lfloor L^{1/d} \rfloor$ ,  $n = \lfloor \log_3(N + 2) \rfloor$ , and  $\widetilde{L} = \lfloor L^{1/d} \rfloor \lfloor n^{1/d} \rfloor$ . Consider the sample set

$$605 \quad \begin{aligned} & \{(1, \widetilde{M} - 1), (2, 0)\} \cup \left\{ \left( \frac{m}{\widetilde{M}}, m \right) : m = 0, 1, \dots, \widetilde{M} - 1 \right\} \\ & \cup \left\{ \left( \frac{m+1}{\widetilde{M}} - \delta, m \right) : m = 0, 1, \dots, \widetilde{M} - 2 \right\}. \end{aligned}$$

606 Its size is

$$607 \quad 2\widetilde{M} + 1 = 2\lfloor N^{1/d} \rfloor^2 \lfloor L^{1/d} \rfloor + 1 = \lfloor N^{1/d} \rfloor \cdot \left( (2\lfloor N^{1/d} \rfloor \lfloor L^{1/d} \rfloor - 1) + 1 \right) + 1.$$

608 By Lemma 4.1 (set  $N_1 = \lfloor N^{1/d} \rfloor$  and  $N_2 = 2\lfloor N^{1/d} \rfloor \lfloor L^{1/d} \rfloor - 1$  therein), there exists

$$609 \quad \begin{aligned} \phi_1 & \in \mathcal{NN}(\text{widthvec} = [2\lfloor N^{1/d} \rfloor, 2(2\lfloor N^{1/d} \rfloor \lfloor L^{1/d} \rfloor - 1) + 1]) \\ & = \mathcal{NN}(\text{widthvec} = [2\lfloor N^{1/d} \rfloor, 4\lfloor N^{1/d} \rfloor \lfloor L^{1/d} \rfloor - 1]) \end{aligned}$$

610 such that



611 •  $\phi_1(\frac{\widetilde{M}-1}{\widetilde{M}}) = \phi_1(1) = \widetilde{M} - 1$  and  $\phi_1(\frac{m}{\widetilde{M}}) = \phi_1(\frac{m+1}{\widetilde{M}} - \delta) = m$  for  $m = 0, 1, \dots, \widetilde{M} - 2$ .

612 •  $\phi_1$  is linear on  $[\frac{\widetilde{M}-1}{\widetilde{M}}, 1]$  and each interval  $[\frac{m}{\widetilde{M}}, \frac{m+1}{\widetilde{M}} - \delta]$  for  $m = 0, 1, \dots, \widetilde{M} - 2$ .

613 Then, for  $m = 0, 1, \dots, \widetilde{M} - 1$ , we have

$$614 \quad \phi_1(x) = m, \quad \text{for any } x \in [\frac{m}{\widetilde{M}}, \frac{m+1}{\widetilde{M}} - \delta \cdot 1_{\{m \leq \widetilde{M}-2\}}]. \quad (4.2)$$

615 Now consider the another sample set

$$616 \quad \left\{ (\frac{1}{\widetilde{M}}, \widetilde{L} - 1), (2, 0) \right\} \cup \left\{ (\frac{\ell}{\widetilde{M}\widetilde{L}}, \ell) : \ell = 0, 1, \dots, \widetilde{L} - 1 \right\} \\ \cup \left\{ (\frac{\ell+1}{\widetilde{M}\widetilde{L}} - \delta, \ell) : \ell = 0, 1, \dots, \widetilde{L} - 2 \right\}.$$

617 Its size is

$$618 \quad 2\widetilde{L} + 1 = 2\lfloor L^{1/d} \rfloor \lfloor n^{1/d} \rfloor + 1 = \lfloor n^{1/d} \rfloor \cdot ((2\lfloor L^{1/d} \rfloor - 1) + 1) + 1.$$

619 By Lemma 4.1 (set  $N_1 = \lfloor n^{1/d} \rfloor$  and  $N_2 = 2\lfloor L^{1/d} \rfloor - 1$  therein), there exists

$$620 \quad \phi_2 \in \mathcal{NN}(\text{widthvec} = [2\lfloor n^{1/d} \rfloor, 2(2\lfloor L^{1/d} \rfloor - 1) + 1]) \\ = \mathcal{NN}(\text{widthvec} = [2\lfloor n^{1/d} \rfloor, 4\lfloor L^{1/d} \rfloor - 1])$$

621 such that

622 •  $\phi_2(\frac{\widetilde{L}-1}{\widetilde{M}\widetilde{L}}) = \phi_2(\frac{1}{\widetilde{M}}) = \widetilde{L} - 1$  and  $\phi_2(\frac{\ell}{\widetilde{M}\widetilde{L}}) = \phi_2(\frac{\ell+1}{\widetilde{M}\widetilde{L}} - \delta) = \ell$  for  $\ell = 0, 1, \dots, \widetilde{L} - 2$ .

623 •  $\phi_2$  is linear on  $[\frac{\widetilde{L}-1}{\widetilde{M}\widetilde{L}}, \frac{1}{\widetilde{M}}]$  and each interval  $[\frac{\ell}{\widetilde{M}\widetilde{L}}, \frac{\ell+1}{\widetilde{M}\widetilde{L}} - \delta]$  for  $\ell = 0, 1, \dots, \widetilde{L} - 2$ .

624 It follows that, for  $m = 0, 1, \dots, \widetilde{M} - 1$  and  $\ell = 0, 1, \dots, \widetilde{L} - 1$ ,

$$625 \quad \phi_2(x - \frac{m}{\widetilde{M}}) = \ell, \quad \text{for any } x \in [\frac{m\widetilde{L}+\ell}{\widetilde{M}\widetilde{L}}, \frac{m\widetilde{L}+\ell+1}{\widetilde{M}\widetilde{L}} - \delta \cdot 1_{\{\ell \leq \widetilde{L}-2\}}]. \quad (4.3)$$

626  $K = \widetilde{M} \cdot \widetilde{L}$  implies any  $k \in \{0, 1, \dots, K-1\}$  can be unique represented by  $k = m\widetilde{L} + \ell$  for  
627  $m = 0, 1, \dots, \widetilde{M} - 1$  and  $\ell = 0, 1, \dots, \widetilde{L} - 1$ . Then the desired function  $\phi$  can be implemented  
628 by a ReLU network shown in Figure 8.

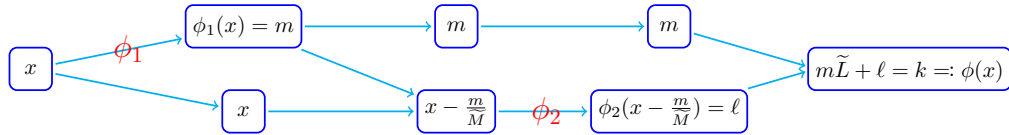


Figure 8: An illustration of the network architecture implementing  $\phi$  based on Equation (4.2) and (4.3) for  $x \in [\frac{k}{K}, \frac{k}{K} - \delta \cdot 1_{\{k \leq K-2\}}] = [\frac{m\widetilde{L}+\ell}{\widetilde{M}\widetilde{L}}, \frac{m\widetilde{L}+\ell+1}{\widetilde{M}\widetilde{L}} - \delta \cdot 1_{\{m \leq \widetilde{M}-2 \text{ or } \ell \leq \widetilde{L}-2\}}]$ , where  $k = m\widetilde{L} + \ell$  for  $m = 0, 1, \dots, \widetilde{M} - 1$  and  $\ell = 0, 1, \dots, \widetilde{L} - 1$ .

629 Clearly,

$$630 \quad \phi(x) = k, \quad \text{if } x \in [\frac{k}{K}, \frac{k}{K} - \delta \cdot 1_{\{k \leq K-2\}}], \quad \text{for any } k \in \{0, 1, \dots, K-1\}.$$

631 By Lemma 4.2, we have

$$632 \quad \begin{aligned} \phi_1 &\in \mathcal{NN}(\#input = 1; \text{widthvec} = [2\lfloor N^{1/d} \rfloor, 4\lfloor N^{1/d} \rfloor \lfloor L^{1/d} \rfloor - 1]; \#output = 1) \\ &\subseteq \mathcal{NN}(\#input = 1; \text{width} \leq 8\lfloor N^{1/d} \rfloor + 2; \text{depth} \leq \lfloor L^{1/d} \rfloor + 1; \#output = 1) \end{aligned}$$

633 and

$$634 \quad \begin{aligned} \phi_2 &\in \mathcal{NN}(\#input = 1; \text{widthvec} = [2\lfloor n^{1/d} \rfloor, 4\lfloor L^{1/d} \rfloor - 1]; \#output = 1) \\ &\subseteq \mathcal{NN}(\#input = 1; \text{width} \leq 8\lfloor n^{1/d} \rfloor + 2; \text{depth} \leq \lfloor L^{1/d} \rfloor + 1; \#output = 1). \end{aligned}$$

635 Recall that  $n = \lfloor \log_3(N + 2) \rfloor \leq N$ . It follows from Figure 8 that  $\phi$  can be implemented  
636 by a ReLU network with width

$$637 \quad \max \{8\lfloor N^{1/d} \rfloor + 2 + 1, 8\lfloor n^{1/d} \rfloor + 2 + 1\} = 8\lfloor N^{1/d} \rfloor + 3$$

638 and depth

$$639 \quad (\lfloor L^{1/d} \rfloor + 1) + 2 + (\lfloor L^{1/d} \rfloor + 1) + 1 = 2\lfloor L^{1/d} \rfloor + 5.$$

640 So we finish the proof.

### 641 4.3 Proof of Proposition 3.2

642 The proof of Proposition 3.2 is based on the bit extraction technique in [3, 8]. In  
643 fact, we modify this technique to extract the sum of many bits rather than one bit and  
644 this modification can be summarized in Lemma 4.4 and 4.5 below.

645 **Lemma 4.4.** *For any  $n \in \mathbb{N}^+$ , there exists a function  $\phi$  in*

$$646 \quad \mathcal{NN}(\#input = 2; \text{width} \leq (n + 1)2^{n+1}; \text{depth} \leq 3; \#output = 1)$$

647 *such that: Given any  $\theta_j \in \{0, 1\}$  for  $j = 1, 2, \dots, n$ , we have*

$$648 \quad \phi(\text{bin } 0.\theta_1\theta_2\cdots\theta_n, i) = \sum_{j=1}^i \theta_j, \quad \text{for any } i \in \{0, 1, 2, \dots, n\}. \quad \textcircled{3}$$

649 *Proof.* Define  $\theta = \text{bin } 0.\theta_1\theta_2\cdots\theta_n$ . Clearly,

$$650 \quad \theta_j = \lfloor 2^j \theta \rfloor / 2 - \lfloor 2^{j-1} \theta \rfloor, \quad \text{for any } j \in \{1, 2, \dots, n\}.$$

651 We shall use a ReLU network to replace  $\lfloor \cdot \rfloor$ . Let  $g \in \text{CPwL}(\mathbb{R}, 2^{n+1} - 2)$  be the function  
652 matching the set of samples

$$653 \quad \bigcup_{k=0}^{2^n-1} \{(k, k), (k + 1 - \delta, k)\}, \quad \text{where } \delta = 2^{-(n+1)}.$$

654 Then  $g(x) = \lfloor x \rfloor$  for any  $x \in \bigcup_{k=0}^{2^n-1} [k, k + 1 - \delta]$ . Note that

$$655 \quad 2^j \theta \in \bigcup_{k=0}^{2^n-1} [k, k + 1 - \delta], \quad \text{for any } j \in \{1, 2, \dots, n\}.$$

---

<sup>③</sup>By convention,  $\sum_{j=n}^m a_j = 0$  if  $n > m$ , no matter what  $a_j$  is for each  $j$ .

Thus,

$$\theta_j = \lfloor 2^j \theta \rfloor / 2 - \lfloor 2^{j-1} \theta \rfloor = g(2^j \theta) / 2 - g(2^{j-1} \theta), \quad \text{for any } j \in \{1, 2, \dots, n\}. \quad (4.4)$$

It is easy to design a ReLU network to output  $\theta_1, \theta_2, \dots, \theta_n$  by Equation (4.4) when using  $\theta = \text{bin}0.\theta_1\theta_2\cdots\theta_n$  as the input. However, it is highly non-trivial to construct a ReLU network to output  $\sum_{j=1}^i \theta_j$  with another input  $i$ , since many operations like multiplication and comparison are not allowed in designing ReLU networks. Now let us establish a formula to represent  $\sum_{j=1}^i \theta_j$  in a form of a ReLU FNN as follows.

Define  $\mathcal{T}(n) := \sigma(n+1) - \sigma(n) = \begin{cases} 1, & n \geq 0 \\ 0, & n < 0 \end{cases}$  for any integer  $n$ . Then, by Equation (4.4) and the fact  $x_1 x_2 = \sigma(x_1 + x_2 - 1)$  for any  $x_1, x_2 \in \{0, 1\}$ , we have, for  $i = 0, 1, 2, \dots, n$ ,

$$\begin{aligned} \sum_{j=1}^i \theta_j &= \sum_{j=1}^n \theta_j \cdot \mathcal{T}(i-j) = \sum_{j=1}^n \theta_j \cdot (\sigma(i-j+1) - \sigma(i-j)) \\ &= \sum_{j=1}^n \sigma(\theta_j + \sigma(i-j+1) - \sigma(i-j) - 1) \\ &= \sum_{j=1}^n \sigma(g(2^j \theta) / 2 - g(2^{j-1} \theta) + \sigma(i-j+1) - \sigma(i-j) - 1). \end{aligned}$$

Define

$$z_{i,j} := \sigma(g(2^j \theta) / 2 - g(2^{j-1} \theta) + \sigma(i-j+1) - \sigma(i-j) - 1), \quad (4.5)$$

for any  $i, j \in \{1, 2, \dots, n\}$ . Then the goal is to design  $\phi$  satisfying

$$\phi(\theta, i) = \sum_{j=1}^i \theta_j = \sum_{j=1}^n z_{i,j}, \quad \text{for any } i \in \{0, 1, 2, \dots, n\}. \quad (4.6)$$

See Figure 9 for the network architecture implementing the desired function  $\phi$ .

By Lemma 4.3, we have

$$g \in \text{CPwL}(\mathbb{R}, 2^{n+1} - 2) \subseteq \mathcal{NN}(\#input = 1; \text{widthvec} = [2^{n+1} - 1]; \#output = 1),$$

implying

$$g(2^j \cdot) \in \text{CPwL}(\mathbb{R}, 2^{n+1} - 2) \subseteq \mathcal{NN}(\#input = 1; \text{widthvec} = [2^{n+1} - 1]; \#output = 1),$$

for any  $j = 0, 1, 2, \dots, n$ . Clearly, the network in Figure 9 has width  $(n+1)(2^{n+1} - 1) + (n+1) = (n+1)2^{n+1}$  and depth 3. So we finish the proof.  $\square$

**Lemma 4.5.** *For any  $n, L \in \mathbb{N}^+$ , there exists a function  $\phi$  in*

$$\mathcal{NN}(\#input = 2; \text{width} \leq (n+3)2^{n+1} + 4; \text{depth} \leq 4L + 2; \#output = 1)$$

*such that: Given any  $\theta_j \in \{0, 1\}$  for  $j = 1, 2, \dots, Ln$ , we have*

$$\phi(\text{bin}0.\theta_1\theta_2\cdots\theta_{Ln}, k) = \sum_{j=1}^k \theta_j, \quad \text{for any } k \in \{1, 2, \dots, Ln\}.$$

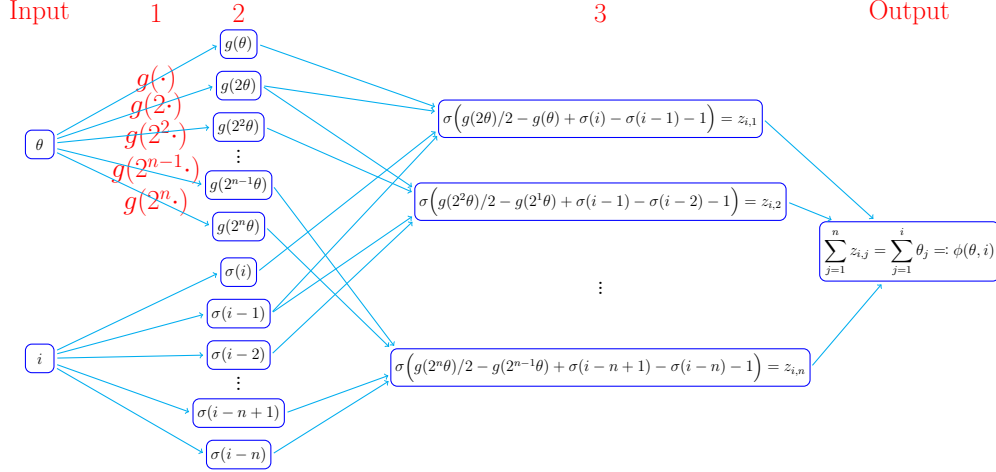


Figure 9: An illustration of the network implementing the desired function  $\phi$  with the input  $[\theta, i]^T = [\text{bin}0.\theta_1\theta_2\cdots\theta_n, i]^T$  for any  $i \in \{0, 1, 2, \dots, n\}$  and  $\theta_1, \theta_2, \dots, \theta_n \in \{0, 1\}$ .  $g(2^j \cdot)$  can be implemented by a one-hidden-layer network with width  $2^{n+1} - 1$  for each  $j \in \{0, 1, \dots, n\}$ . The red numbers above the architecture indicate the order of hidden layers. The network architecture is essentially determined by Equation (4.5) and (4.6), which are valid no matter what  $\theta_1, \theta_2, \dots, \theta_n \in \{0, 1\}$ . Thus, the desired function  $\phi$  is independent of  $\theta_1, \theta_2, \dots, \theta_n \in \{0, 1\}$ . We omit ReLU ( $\sigma$ ) for a neuron if its output is non-negative without ReLU. Such a simplification are applied to similar figures in this paper.

681 *Proof.* Let  $g_1 \in \text{CPwL}(\mathbb{R}, 2^{n+1} - 2)$  be the function matching the set of samples

$$682 \quad \bigcup_{i=0}^{2^n-1} \{(i, i), (i+1-\delta, i)\}, \quad \text{where } \delta = 2^{-(L_{n+1})}.$$

683 Then  $g_1(x) = \lfloor x \rfloor$  for any  $x \in \bigcup_{i=0}^{2^n-1} [i, i+1-\delta]$ . Note that

$$684 \quad 2^n \cdot \text{bin}0.\theta_{\ell_{n+1}}\cdots\theta_{L_n} \in \bigcup_{i=0}^{2^n-1} [i, i+1-\delta], \quad \text{for any } \ell \in \{0, 1, \dots, L-1\}.$$

685 Thus, for any  $\ell \in \{0, 1, \dots, L-1\}$ , we have

$$686 \quad \text{bin}0.\theta_{\ell_{n+1}}\cdots\theta_{\ell_{n+n}} = \frac{\lfloor 2^n \cdot \text{bin}0.\theta_{\ell_{n+1}}\cdots\theta_{L_n} \rfloor}{2^n} = \frac{g_1(2^n \cdot \text{bin}0.\theta_{\ell_{n+1}}\cdots\theta_{L_n})}{2^n}. \quad (4.7)$$

687 Define  $g_2(x) := 2^n x - g_1(2^n x)$  for any  $x \in \mathbb{R}$ . Then  $g_2 \in \text{CPwL}(\mathbb{R}, 2^{n+1} - 2)$  and

$$688 \quad \begin{aligned} & \text{bin}0.\theta_{(\ell+1)_{n+1}}\cdots\theta_{L_n} = 2^n \left( \text{bin}0.\theta_{\ell_{n+1}}\cdots\theta_{L_n} - \text{bin}0.\theta_{\ell_{n+1}}\cdots\theta_{\ell_{n+n}} \right) \\ & = 2^n \left( \text{bin}0.\theta_{\ell_{n+1}}\cdots\theta_{L_n} - \frac{g_1(2^n \cdot \text{bin}0.\theta_{\ell_{n+1}}\cdots\theta_{L_n})}{2^n} \right) = g_2(\text{bin}0.\theta_{\ell_{n+1}}\cdots\theta_{L_n}). \end{aligned} \quad (4.8)$$

689 By Lemma 4.4, there exists

$$690 \quad \phi_1 \in \mathcal{NN}(\# \text{input} = 2; \text{width} \leq (n+1)2^{n+1}; \text{depth} \leq 3; \# \text{output} = 1)$$

such that: For any  $\xi_1, \xi_2, \dots, \xi_n \in \{0, 1\}$ , we have

$$\phi_1(\text{bin}0.\xi_1\xi_2\cdots\xi_n, i) = \sum_{j=1}^i \xi_j, \quad \text{for } i = 0, 1, 2, \dots, n.$$

It follows that

$$\phi_1(\text{bin}0.\theta_{\ell n+1}\theta_{\ell n+2}\cdots\theta_{\ell n+n}, i) = \sum_{j=1}^i \theta_{\ell n+j}, \quad \text{for } \ell = 0, 1, \dots, L-1 \text{ and } i = 0, 1, \dots, n. \quad (4.9)$$

Define  $\phi_{2,\ell}(x) := \min\{\sigma(x - \ell n), n\}$  for any  $x \in \mathbb{R}$  and  $\ell \in \{0, 1, \dots, L-1\}$ . For any  $k \in \{1, 2, \dots, Ln\}$ , there exists  $k_1 \in \{0, 1, \dots, L-1\}$  and  $k_2 \in \{1, 2, \dots, n\}$  such that  $k = k_1 n + k_2$ , implying

$$\begin{aligned} \sum_{i=1}^k \theta_i &= \sum_{i=1}^{k_1 n + k_2} \theta_i = \sum_{\ell=0}^{k_1-1} \left( \sum_{j=1}^n \theta_{\ell n+j} \right) + \sum_{\ell=k_1}^{k_1} \left( \sum_{j=1}^{k_2} \theta_{\ell n+j} \right) + \sum_{\ell=k_1+1}^{L-1} \left( \sum_{j=1}^0 \theta_{\ell n+j} \right) \\ &= \sum_{\ell=0}^{L-1} \left( \min\{\sigma(k - \ell n), n\} \sum_{j=1}^n \theta_{\ell n+j} \right) = \sum_{\ell=0}^{L-1} \left( \sum_{j=1}^n \theta_{\ell n+j} \right). \end{aligned} \quad (4.10)$$

Then, the desired function  $\phi$  can be implemented by the network architecture in Figure 10.

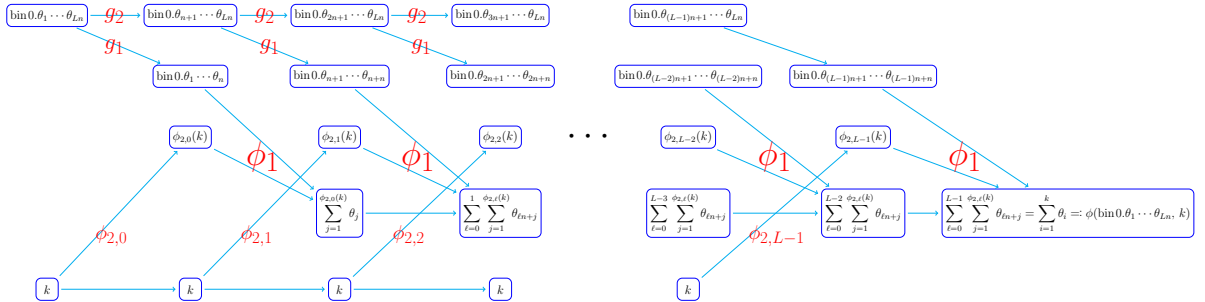


Figure 10: An illustration of the network implementing the desired function  $\phi$  with the input  $[\text{bin}0.\theta_1\theta_2\cdots\theta_{Ln}, k]^T$  for any  $k \in \{1, 2, \dots, Ln\}$  and  $\theta_1, \theta_2, \dots, \theta_{Ln} \in \{0, 1\}$ . The network architecture is essentially determined by Equation (4.7), (4.8), (4.9), and (4.10), which are valid no matter what  $\theta_1, \theta_2, \dots, \theta_{Ln} \in \{0, 1\}$ . Thus, the desired function  $\phi$  is independent of  $\theta_1, \theta_2, \dots, \theta_{Ln} \in \{0, 1\}$ . We omit ReLU ( $\sigma$ ) for a neuron if its output is non-negative without ReLU.

By Lemma 4.3, we have

$$g_1, g_2 \in \text{CPwL}(\mathbb{R}, 2^{n+1} - 2) \subseteq \mathcal{NN}(\# \text{input} = 1; \text{widthvec} = [2^{n+1} - 1]; \# \text{output} = 1).$$

Recall that  $\phi_1 \in \mathcal{NN}(\text{width} \leq (n+1)2^{n+1}; \text{depth} \leq 3)$ . As shown in Figure 11,  $\phi_{2,\ell}(x) \in \mathcal{NN}(\text{width} \leq 4; \text{depth} \leq 2)$  for  $\ell = 0, 1, \dots, L-1$ . Therefore, the network in Figure 10 has width

$$(2^{n+1} - 1) + (2^{n+1} - 1) + (n+1)2^{n+1} + 1 + 4 + 1 = (n+3)2^{n+1} + 4$$

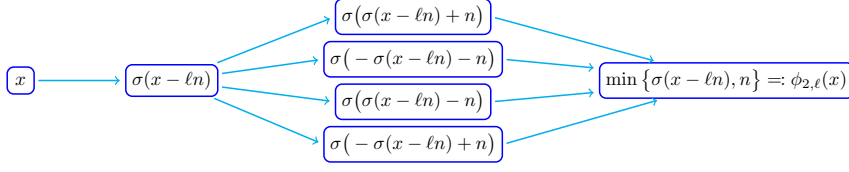


Figure 11: An illustration of the network implementing the desired function  $\phi_{2,\ell}$  for each  $\ell \in \{0, 1, \dots, L-1\}$ , based on  $\min\{x, n\} = \frac{1}{2}(\sigma(x - n) - \sigma(-x - n) - \sigma(x + n) - \sigma(-x - n))$ .

707 and depth

$$708 \quad 2 + L(1 + 3) = 4L + 2.$$

709 So we finish the proof. □

710 Next, we introduce Lemma 4.6 to map indices to the partial sum of given bits.

711 **Lemma 4.6.** *Given any  $N, L \in \mathbb{N}^+$  and arbitrary  $\theta_{m,k} \in \{0, 1\}$  for  $m = 0, 1, \dots, M-1$  and*  
 712  *$k = 0, 1, \dots, Ln-1$ , where  $M = N^2L$  and  $n = \lfloor \log_3(N+2) \rfloor$ , there exists*

$$713 \quad \phi \in \mathcal{NN}(\#input = 2; \text{ width} \leq 6N + 14; \text{ depth} \leq 5L + 4; \#output = 1)$$

714 such that

$$715 \quad \phi(m, k) = \sum_{j=0}^k \theta_{m,j}, \quad \text{for } m = 0, 1, \dots, M-1 \text{ and } k = 0, 1, \dots, Ln-1.$$

716 *Proof.* Define

$$717 \quad y_m := \text{bin}0.\theta_{m,0}\theta_{m,1}\dots\theta_{m,Ln-1}, \quad \text{for } m = 0, 1, \dots, M-1.$$

718 Consider the sample set  $\{(m, y_m) : m = 0, 1, \dots, M\}$ , whose cardinality is

$$719 \quad M + 1 = N((NL - 1) + 1) + 1.$$

720 By Lemma 4.1 (set  $N_1 = N$  and  $N_2 = NL - 1$  therein), there exists

$$721 \quad \begin{aligned} \phi_1 &\in \mathcal{NN}(\#input = 1; \text{ widthvec} = [2N, 2(NL - 1) + 1]; \#output = 1) \\ &= \mathcal{NN}(\#input = 1; \text{ widthvec} = [2N, 2NL - 1]; \#output = 1) \end{aligned}$$

722 such that

$$723 \quad \phi_1(m) = y_m, \quad \text{for } m = 0, 1, \dots, M-1.$$

724 By Lemma 4.4, there exists

$$725 \quad \phi_2 \in \mathcal{NN}(\#input = 2; \text{ width} \leq (n + 3)2^{n+1} + 4; \text{ depth} \leq 4L + 2; \#output = 1)$$

726 such that, for any  $\xi_1, \xi_2, \dots, \xi_{Ln} \in \{0, 1\}$ , we have

$$727 \quad \phi_2(\text{bin}0.\xi_1\xi_2\dots\xi_{Ln}, k) = \sum_{j=1}^k \xi_j, \quad \text{for } k = 1, 2, \dots, Ln.$$

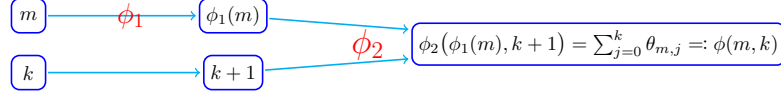


Figure 12: An illustration of the network implementing the desired function  $\phi$  for  $m = 0, 1, \dots, M-1$  and  $k = 0, 1, \dots, Ln-1$ .

It follows that, for any  $\xi_0, \xi_1, \dots, \xi_{Ln-1} \in \{0, 1\}$ , we have

$$\phi_2(\text{bin} 0.\xi_0\xi_1\dots\xi_{Ln-1}, k+1) = \sum_{j=0}^k \xi_j, \quad \text{for } k = 0, 1, \dots, Ln-1.$$

Thus, for  $m = 0, 1, \dots, M-1$  and  $k = 0, 1, \dots, Ln-1$ , we have

$$\phi_2(\phi_1(m), k+1) = \phi_2(y_m, k+1) = \phi_2(0.\theta_{m,0}\theta_{m,1}\dots\theta_{m,Ln-1}, k+1) = \sum_{j=0}^k \theta_{m,j}.$$

Hence, the desired function  $\phi$  can be implemented by the network shown in Figure 12. By Lemma 4.2,  $\phi_1 \in \mathcal{NN}(\text{widthvec} = [2N, 2NL-1]) \subseteq \mathcal{NN}(\text{width} \leq 4N+2; \text{depth} \leq L+1)$ . It holds that

$$(n+3)2^{n+1} + 4 \leq 6 \cdot (3^n) + 2 = 6 \cdot (3^{\lfloor \log_3(N+2) \rfloor}) + 2 \leq 6(N+2) + 2 = 6N+14,$$

implying

$$\begin{aligned} \phi_2 &\in \mathcal{NN}(\#input = 2; \text{width} \leq (n+3)2^{n+1} + 4; \text{depth} \leq 4L+2; \#output = 1) \\ &\subseteq \mathcal{NN}(\#input = 2; \text{width} \leq 6N+14; \text{depth} \leq 4L+2; \#output = 1). \end{aligned}$$

Therefore, the network in Figure 12 is with width  $\max\{(4N+2)+1, 6N+14\} = 6N+14$  and depth  $(4L+2)+1+(L+1) = 5L+4$ . So we finish the proof.  $\square$

Next, we apply Lemma 4.6 to prove Lemma 4.7 below, which is a key intermediate conclusion to prove Proposition 3.2.

**Lemma 4.7.** For any  $\varepsilon > 0$  and  $N, L \in \mathbb{N}^+$ , denote  $M = N^2L$  and  $n = \lfloor \log_3(N+2) \rfloor$ . Assume  $y_{m,k} \geq 0$  for  $m = 0, 1, \dots, M-1$  and  $k = 0, 1, \dots, Ln-1$  are samples with

$$|y_{m,k} - y_{m,k-1}| \leq \varepsilon, \quad \text{for } m = 0, 1, \dots, M-1 \quad \text{and} \quad k = 1, 2, \dots, Ln-1.$$

Then there exists  $\phi \in \mathcal{NN}(\#input = 2; \text{width} \leq 16N+30; \text{depth} \leq 5L+7; \#output = 1)$  such that

(i)  $|\phi(m, k) - y_{m,k}| \leq \varepsilon$  for  $m = 0, 1, \dots, M-1$  and  $k = 0, 1, \dots, Ln-1$ ;

(ii)  $0 \leq \phi(x_1, x_2) \leq \max\{y_{m,k} : m = 0, 1, \dots, M-1 \text{ and } k = 0, 1, \dots, Ln-1\}$  for any  $x_1, x_2 \in \mathbb{R}$ .



750 *Proof.* Define

$$751 \quad a_{m,k} := \lfloor y_{m,k}/\varepsilon \rfloor, \quad \text{for } m = 0, 1, \dots, M-1 \quad \text{and} \quad k = 0, 1, \dots, Ln-1.$$

752 We will construct a function implemented by a ReLU network to map the index  $(m, k)$   
 753 to  $a_{m,k}\varepsilon$  for  $m = 0, 1, \dots, M-1$  and  $k = 0, 1, \dots, Ln-1$ .

754 Define  $b_{m,0} := 0$  and  $b_{m,k} := a_{m,k} - a_{m,k-1}$  for  $m = 0, 1, \dots, M-1$  and  $k = 1, 2, \dots, Ln-1$ .  
 755 Since  $|y_{m,k} - y_{m,k-1}| \leq \varepsilon$  for all  $m$  and  $k$ , we have  $b_{m,k} \in \{-1, 0, 1\}$ . Hence, there exist  $c_{m,k}$   
 756 and  $d_{m,k} \in \{0, 1\}$  such that  $b_{m,k} = c_{m,k} - d_{m,k}$ , which implies

$$\begin{aligned} 757 \quad a_{m,k} &= a_{m,0} + \sum_{j=1}^k (a_{m,j} - a_{m,j-1}) = a_{m,0} + \sum_{j=1}^k b_{m,j} = a_{m,0} + \sum_{j=0}^k b_{m,j} \\ &= a_{m,0} + \sum_{j=0}^k c_{m,j} - \sum_{j=0}^k d_{m,j}, \end{aligned}$$

758 for  $m = 0, 1, \dots, M-1$  and  $k = 0, 1, \dots, Ln-1$ .

759 Consider the sample set

$$760 \quad \{(m, a_{m,0}) : m = 0, 1, \dots, M-1\} \cup \{(M, 0)\}.$$

761 Its size is  $M+1 = N \cdot ((NL-1)+1) + 1$ , by Lemma 4.1 (set  $N_1 = N$  and  $N_2 = NL-1$   
 762 therein), there exists

$$763 \quad \psi_1 \in \mathcal{NN}(\text{widthvec} = [2N, 2(NL-1)+1]) = \mathcal{NN}(\text{widthvec} = [2N, 2NL-1])$$

764 such that

$$765 \quad \psi_1(m) = a_{m,0}, \quad \text{for } m = 0, 1, \dots, M-1.$$

766 By Lemma 4.6, there exist  $\psi_2, \psi_3 \in \mathcal{NN}(\text{width} \leq 6N+14; \text{depth} \leq 5L+4)$  such that

$$767 \quad \psi_2(m, k) = \sum_{j=0}^k c_{m,j} \quad \text{and} \quad \psi_3(m, k) = \sum_{j=0}^k d_{m,j},$$

768 for  $m = 0, 1, \dots, M-1$  and  $k = 0, 1, \dots, Ln-1$ . Hence, it holds that

$$769 \quad a_{m,k} = a_{m,0} + \sum_{j=0}^k c_{m,j} - \sum_{j=0}^k d_{m,j} = \psi_1(m) + \psi_2(m, k) - \psi_3(m, k), \quad (4.11)$$

770 for  $m = 0, 1, \dots, M-1$  and  $k = 0, 1, \dots, Ln-1$ .

771 Define

$$772 \quad y_{\max} := \max\{y_{m,k} : m = 0, 1, \dots, M-1 \quad \text{and} \quad k = 0, 1, \dots, Ln-1\}.$$

773 Then the desired function can be implemented by two sub-networks shown in Figure 13.

774 By Lemma 4.2,

$$\begin{aligned} 775 \quad \psi_1 &\in \mathcal{NN}(\#input = 1; \text{widthvec} = [2N, 2NL-1]; \#output = 1) \\ &\subseteq \mathcal{NN}(\#input = 1; \text{width} \leq 4N+2; \text{depth} \leq L+1; \#output = 1). \end{aligned}$$

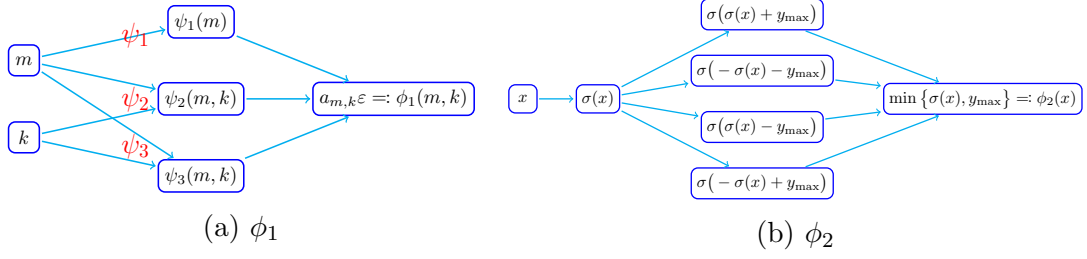


Figure 13: Illustrations of two sub-networks implementing the desired function  $\phi = \phi_2 \circ \phi_1$  for  $m = 0, 1, \dots, M-1$  and  $k = 0, 1, \dots, Ln-1$ , based on Equation (4.11) and the fact  $\min\{x_1, x_2\} = \frac{x_1+x_2-|x_1-x_2|}{2} = \frac{\sigma(x_1+x_2)-\sigma(-x_1-x_2)-\sigma(x_1-x_2)-\sigma(-x_1+x_2)}{2}$ .

Recall that  $\psi_2, \psi_3 \in \mathcal{NN}$ (width  $\leq 6N + 14$ ; depth  $\leq 5L + 4$ ). Thus,  $\phi_1 \in \mathcal{NN}$ (width  $\leq (4N + 2) + 2(6N + 14) = 16N + 30$ ; depth  $\leq (5L + 4) + 1 = 5L + 5$ ) as shown in Figure 13. And it is clear that  $\phi_2 \in \mathcal{NN}$ (width  $\leq 4$ ; depth  $\leq 2$ ), implying  $\phi = \phi_2 \circ \phi_1 \in \mathcal{NN}$ (width  $\leq 16N + 30$ ; depth  $\leq (5L + 5) + 2 = 5L + 7$ ).

Clearly,  $0 \leq \phi(x_1, x_2) \leq y_{\max}$  for any  $x_1, x_2 \in \mathbb{R}$ , since  $\phi(x_1, x_2) = \phi_2 \circ \phi_1(x_1, x_2) = \max\{\sigma(\phi_1(x_1, x_2)), y_{\max}\}$ .

Note that  $0 \leq a_{m,k}\varepsilon = \lfloor y_{m,k}/\varepsilon \rfloor \varepsilon \leq y_{\max}$ . Then we have  $\phi(m, k) = \phi_2 \circ \phi_1(m, k) = \phi_2(a_{m,k}\varepsilon) = \max\{\sigma(a_{m,k}\varepsilon), y_{\max}\} = a_{m,k}\varepsilon$ . Therefore,

$$|\phi(m, k) - y_{m,k}| = |a_{m,k}\varepsilon - y_{m,k}| = \left| \lfloor y_{m,k}/\varepsilon \rfloor \varepsilon - y_{m,k} \right| \leq \varepsilon,$$

for  $m = 0, 1, \dots, M-1$  and  $k = 0, 1, \dots, Ln-1$ . Hence, we finish the proof.  $\square$

Finally, we apply Lemma 4.7 to prove Proposition 3.2.

*Proof of Proposition 3.2.* Denote  $M = N^2L$ ,  $n = \lfloor \log_3(N + 2) \rfloor$ , and  $\widehat{L} = Ln$ . We may assume  $J = MLn = M\widehat{L}$  since we can set  $y_{J-1} = y_J = y_{J+1} = \dots = y_{M\widehat{L}-1}$  if  $J < M\widehat{L}$ .

Consider the sample set

$$\{(m\widehat{L}, m) : m = 0, 1, \dots, M\} \cup \{(m\widehat{L} + \widehat{L} - 1, m) : m = 0, 1, \dots, M-1\}.$$

Its size is  $2M + 1 = N \cdot ((2NL - 1) + 1) + 1$ . By Lemma 4.1 (set  $N_1 = N$  and  $N_2 = NL - 1$  therein), there exist

$$\phi_1 \in \mathcal{NN}(\text{widthvec} = [2N, 2(2NL - 1) + 1]) = \mathcal{NN}(\text{widthvec} = [2N, 4NL - 1])$$

such that

- $\phi_1(M\widehat{L}) = M$  and  $\phi_1(m\widehat{L}) = \phi_1(m\widehat{L} + \widehat{L} - 1) = m$  for  $m = 0, 1, \dots, M-1$ .
- $\phi_1$  is linear on each interval  $[m\widehat{L}, m\widehat{L} + \widehat{L} - 1]$  for  $m = 0, 1, \dots, M-1$ .

It follows that

$$\phi_1(j) = m, \quad \text{and} \quad j - \widehat{L}\phi_1(j) = k, \quad \text{where } j = m\widehat{L} + k, \quad (4.12)$$

for  $m = 0, 1, \dots, M-1$  and  $k = 0, 1, \dots, \widehat{L} - 1$ .

Note that any number  $j$  in  $\{0, 1, \dots, J-1\}$  can be uniquely indexed as  $j = m\widehat{L} + k$  for  $m = 0, 1, \dots, M-1$  and  $k = 0, 1, \dots, \widehat{L}-1$ . So we can denote  $y_j = y_{m\widehat{L}+k}$  as  $y_{m,k}$ . Then by Lemma 4.7, there exists  $\phi_2 \in \mathcal{NN}$ (width  $\leq 12N + 8$ ; depth  $\leq 3L + 6$ ) such that

$$|\phi_2(m, k) - y_{m,k}| \leq \varepsilon, \quad \text{for } m = 0, 1, \dots, M-1 \quad \text{and} \quad k = 0, 1, \dots, \widehat{L}-1, \quad (4.13)$$

and

$$0 \leq \phi_2(x_1, x_2) \leq y_{\max}, \quad \text{for any } x_1, x_2 \in \mathbb{R}, \quad (4.14)$$

where  $y_{\max} := \max\{y_{m,k} : m = 0, 1, \dots, M-1 \text{ and } k = 0, 1, \dots, \widehat{L}-1\} = \max\{y_j : j = 0, 1, \dots, M\widehat{L}-1\}$ .

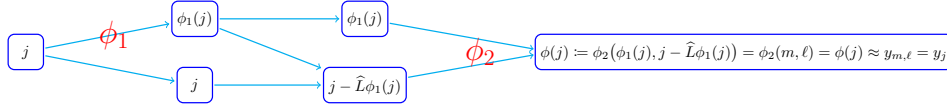


Figure 14: An illustration of the ReLU network implementing the desired function  $\phi$  based Equation (4.12). The index  $j \in \{0, 1, \dots, M\widehat{L}-1\}$  is uniquely represented by  $j = m\widehat{L} + k$  for  $m = 0, 1, \dots, M-1$  and  $k = 0, 1, \dots, \widehat{L}-1$ .

By Lemma 4.2,

$$\begin{aligned} \phi_1 &\in \mathcal{NN}(\#input = 1; \text{widthvec} = [2N, 4NL - 1]; \#output = 1) \\ &\subseteq \mathcal{NN}(\#input = 1; \text{width} \leq 8N + 2; \text{depth} \leq L + 1; \#output = 1). \end{aligned}$$

Recall that  $\phi_2 \in \mathcal{NN}$ (width  $\leq 16N + 30$ ; depth  $\leq 5L + 7$ ). So  $\phi \in \mathcal{NN}$ (width  $\leq 16N + 30$ ; depth  $\leq (L + 1) + 2 + (5L + 7) = 6L + 10$ ) as shown in Figure 14.

Equation (4.14) implies

$$0 \leq \phi(x) \leq y_{\max}, \quad \text{for any } x \in \mathbb{R},$$

since  $\phi$  is given by  $\phi(x) = \phi_2(\phi_1(x), x - L\phi_1(x))$ .

Represent  $j \in \{0, 1, \dots, M\widehat{L}-1\}$  via  $j = m\widehat{L} + k$  for  $m = 0, 1, \dots, M-1$  and  $k = 0, 1, \dots, \widehat{L}-1$ . Then, by Equation (4.13), we have

$$|\phi(j) - y_j| = |\phi_2(\phi_1(j), j - L\phi_1(j)) - y_j| = |\phi_2(m, k) - y_{m,k}| \leq \varepsilon,$$

for any  $j \in \{0, 1, \dots, M\widehat{L}-1\} = \{0, 1, \dots, J-1\}$ . So we finish the proof.  $\square$

We would like to remark that the key idea in the proof of Proposition 3.2 is the bit extraction technique in Lemma 4.5, which allows us to store  $Ln$  bits in a binary number  $\text{bin}0.\theta_1\theta_2\cdots\theta_{Ln}$  and extract each bit  $\theta_i$ . The extraction operator can be efficiently carried out via a deep ReLU neural network demonstrating the power of depth.

## 5 Conclusion and future work

This paper aims at a quantitative and optimal approximation rate for ReLU networks in terms of the width and depth to approximate continuous functions. It is shown by construction that ReLU networks with width  $\mathcal{O}(N)$  and depth  $\mathcal{O}(L)$  can approximate an arbitrary continuous function on  $[0, 1]^d$  with an approximation rate  $\mathcal{O}(\omega_f((N^2 L^2 \ln N)^{-1/d}))$ . By connecting the approximation property to VC-dimension, we prove that such a rate is optimal for Hölder continuous functions on  $[0, 1]^d$  in terms of the width and depth separately, and hence this rate is also optimal for the whole continuous function class. We also extend our analysis to general continuous functions on any bounded set in  $\mathbb{R}^d$ . We would like to remark that our analysis was based on the fully connected feed-forward neural networks and the ReLU activation function. It would be very interesting to extend our conclusions to neural networks with other types of architectures (e.g., convolutional neural networks) and activation functions (e.g., tanh and sigmoid functions).

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