

# Hybrid Metaheuristic Solver

for the

## Capacitated Vehicle Routing Problem with Time Windows (CVRPTW)

Technical Report

AI50 Project

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### Abstract

This technical report presents a comprehensive implementation of a hybrid metaheuristic solver for the Capacitated Vehicle Routing Problem with Time Windows (CVRPTW). The solver combines three powerful optimization algorithms—Ant Colony Optimization (ACO), Genetic Algorithm (GA), and Tabu Search—in a sequential three-stage pipeline to achieve high-quality solutions. The project demonstrates advanced software engineering practices with a modular architecture, interactive web interface using Streamlit, and extensive benchmark testing on standard Solomon instances. The implementation addresses real-world logistics optimization challenges where both vehicle capacity constraints and strict time window requirements must be satisfied.

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# 1 Introduction

## 1.1 Problem Overview

The Capacitated Vehicle Routing Problem with Time Windows (CVRPTW) is a fundamental optimization problem in logistics and operations research. It extends the classical Vehicle Routing Problem (VRP) by incorporating two critical real-world constraints:

- **Capacity Constraint:** Each vehicle has a maximum load capacity that cannot be exceeded.
- **Time Window Constraint:** Each customer must be served within a specific time interval  $[a_i, b_i]$ .

The objective is to minimize the total distance traveled by all vehicles while ensuring that all customers are visited exactly once and all constraints are satisfied.

## 1.2 Mathematical Formulation

Let  $G = (V, E)$  be a complete graph where:

- $V = \{0, 1, \dots, n\}$  represents nodes (0 is the depot,  $1, \dots, n$  are customers)
- $E$  represents edges between nodes with associated distances  $d_{ij}$

**Objective Function:**

$$\text{Minimize } Z = \sum_{k=1}^K \sum_{i=0}^n \sum_{j=0}^n d_{ij} x_{ijk} \quad (1)$$

**Subject to:**

1. **Flow Conservation:** Each customer visited exactly once

$$\sum_{k=1}^K \sum_{j=0}^n x_{ijk} = 1 \quad \forall i \in \{1, \dots, n\} \quad (2)$$

2. **Capacity Constraint:** Total demand on each route  $\leq Q$

$$\sum_{i=1}^n d_i \cdot \sum_{j=0}^n x_{ijk} \leq Q \quad \forall k \quad (3)$$

3. **Time Window Constraint:** Service starts within  $[a_i, b_i]$

$$a_i \leq s_i \leq b_i \quad \forall i \in V \quad (4)$$

4. **Time Consistency:**

$$s_i + \text{service}_i + t_{ij} \leq s_j + M(1 - x_{ijk}) \quad \forall i, j, k \quad (5)$$

where  $x_{ijk} \in \{0, 1\}$  indicates if vehicle  $k$  travels from  $i$  to  $j$ ,  $s_i$  is the service start time at customer  $i$ , and  $M$  is a large constant.

### 1.3 Problem Complexity

CVRPTW is NP-hard, belonging to the class of combinatorial optimization problems where finding an optimal solution becomes computationally intractable as problem size increases. For  $n$  customers and  $k$  vehicles, the search space grows factorially, making exact methods impractical for real-world instances. This motivates the use of metaheuristic approaches that can find high-quality solutions in reasonable time.

## 2 Hybrid Solver Architecture

### 2.1 Three-Stage Pipeline Strategy

The hybrid solver employs a sequential three-stage pipeline, where each algorithm plays a specific role:

1. **Stage 1 - Ant Colony Optimization (ACO):** *Construction & Initial Exploration*
  - Generates initial population of feasible solutions
  - Excels at respecting time window constraints during construction
  - Provides diverse starting points for genetic evolution
2. **Stage 2 - Genetic Algorithm (GA):** *Global Search & Diversification*
  - Evolves population through crossover and mutation
  - Explores broad regions of solution space
  - Combines good features from multiple solutions
3. **Stage 3 - Tabu Search (TS):** *Local Refinement & Intensification*
  - Performs intensive local search on best GA solution
  - Escapes local optima using tabu memory
  - Delivers final polished solution

### 2.2 Design Rationale

The hybrid approach addresses limitations of individual metaheuristics:

- **ACO alone** is slow to converge to high-quality solutions
- **GA alone** may struggle with time window feasibility
- **Tabu Search alone** is highly dependent on initial solution quality

By combining them sequentially, we leverage:

- ACO's strength in constructive heuristics
- GA's global exploration capabilities
- Tabu Search's local optimization power

## 3 Software Architecture

### 3.1 Modular Design

The project follows a clean, modular architecture with clear separation of concerns:

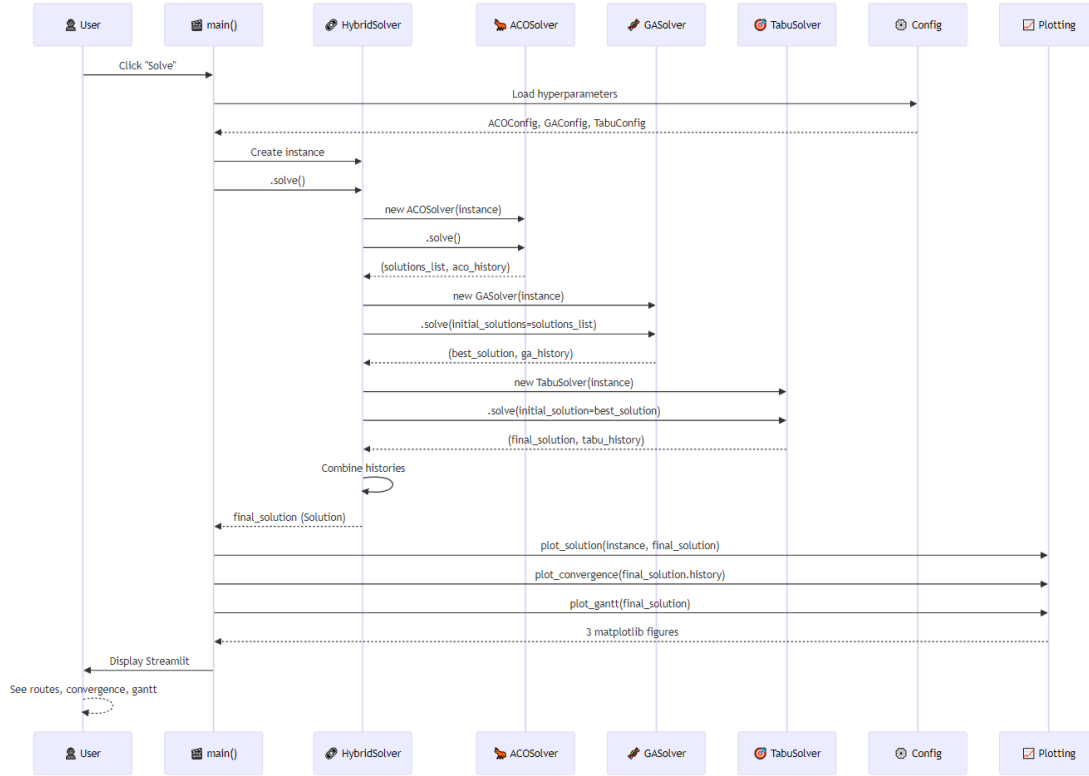


Figure 1: Hybrid Solver Three-Stage Pipeline Architecture

```

src/
|-- core/                                # Domain models & abstractions
|   |-- models.py                        # Node, Route, CVRPTWInstance
|   |-- solution.py                      # Solution wrapper class
|   +-- interfaces.py                   # SolverStrategy abstract class
|-- solvers/                             # Algorithm implementations
|   |-- aco.py                           # Ant Colony Optimization
|   |-- ga.py                            # Genetic Algorithm
|   |-- tabu.py                          # Tabu Search
|   +-- hybrid.py                       # Hybrid orchestrator
|-- utils/                               # Utilities
|   |-- solomon_loader.py                # Solomon instance parser
|   |-- plotting.py                      # Visualization tools
|   +-- logger.py                       # Logging utilities
|-- config.py                            # Configuration dataclasses
+-- cli.py                              # Command-line interface
app.py                                  # Streamlit web application
  
```

## 3.2 Core Data Structures

### 3.2.1 Node Class

Represents a customer or depot with spatial and temporal attributes:

```

1 @dataclass(frozen=True)
2 class Node:
3     id: int
4     x: float
5     y: float
  
```

```

6     demand: float
7     ready_time: float
8     due_date: float
9     service_time: float
10
11     def distance_to(self, other: 'Node') -> float:
12         return math.sqrt((self.x - other.x)**2 +
13                          (self.y - other.y)**2)

```

### 3.2.2 Route Class

Manages a single vehicle's tour with feasibility checking:

```

1 @dataclass
2 class Route:
3     nodes: List[Node]
4     total_distance: float = 0.0
5     total_load: float = 0.0
6     schedule: List[Tuple[float, float, float, float]]
7
8     def is_feasible(self, capacity: float) -> bool:
9         # Check capacity constraint
10        if sum(n.demand for n in self.nodes) > capacity:
11            return False
12
13        # Check time windows
14        current_time = 0.0
15        for i in range(len(self.nodes) - 1):
16            # Calculate arrival, wait, service, departure
17            # Verify time window compliance
18        return True

```

### 3.2.3 Solution Class

Encapsulates a complete CVRPTW solution:

```

1 @dataclass
2 class Solution:
3     routes: List[Route]
4     total_distance: float
5     total_wait: float
6     is_feasible: bool
7
8     def fitness(self) -> float:
9         return self.total_distance if self.is_feasible
10        else float('inf')

```

## 3.3 Strategy Pattern for Solvers

All solvers inherit from a common interface:

```

1 from abc import ABC, abstractmethod
2
3 class SolverStrategy(ABC):
4     @abstractmethod
5     def solve(self) -> Solution:
6         pass

```

This allows polymorphic usage and easy algorithm swapping.



## 4 Algorithm Implementations

### 4.1 Ant Colony Optimization (ACO)

#### 4.1.1 Algorithm Overview

ACO simulates the behavior of ant colonies finding optimal paths through pheromone trails. Each ant constructs a complete solution by probabilistically selecting the next customer based on:

- **Pheromone intensity**  $\tau_{ij}$  (learned from past good solutions)
- **Heuristic information**  $\eta_{ij} = 1/d_{ij}$  (distance-based visibility)

#### 4.1.2 Selection Probability

The probability that ant  $k$  at customer  $i$  selects customer  $j$  is:

$$p_{ij}^k = \frac{[\tau_{ij}]^\alpha \cdot [\eta_{ij}]^\beta}{\sum_{l \in \mathcal{N}_i^k} [\tau_{il}]^\alpha \cdot [\eta_{il}]^\beta} \quad (6)$$

where:

- $\alpha$  controls pheromone influence (exploration vs exploitation)
- $\beta$  controls heuristic influence (greedy behavior)
- $\mathcal{N}_i^k$  is the set of feasible neighbors for ant  $k$  at  $i$

#### 4.1.3 Pheromone Update

After all ants complete their tours:

$$\tau_{ij} \leftarrow (1 - \rho) \cdot \tau_{ij} + \sum_{k=1}^m \Delta\tau_{ij}^k \quad (7)$$

where:

- $\rho \in (0, 1)$  is the evaporation rate
- $\Delta\tau_{ij}^k = Q/L_k$  if ant  $k$  used edge  $(i, j)$ , else 0
- $L_k$  is the total length of ant  $k$ 's tour

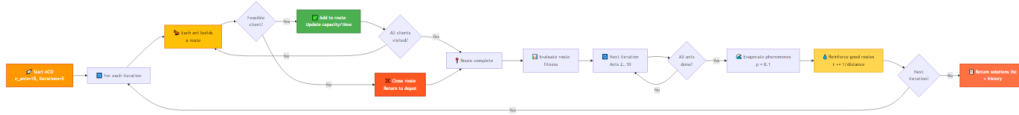


Figure 2: ACO Pheromone Trail Mechanism

### 4.2 Genetic Algorithm (GA)

#### 4.2.1 Encoding Scheme

Solutions are encoded as permutations of customer IDs. Routes are separated by depot markers (ID = 0).

Example:  $[0, 3, 7, 0, 1, 5, 9, 0, 2, 4, 6, 0]$  represents 3 routes.

### 4.2.2 Genetic Operators

#### 1. Selection - Tournament Selection

- Randomly select  $k$  individuals
- Choose the best among them
- Provides selection pressure while maintaining diversity

#### 2. Crossover - Order Crossover (OX)

Preserves relative order of customers from both parents:

1. Select random segment from Parent 1
2. Copy this segment to Child
3. Fill remaining positions with customers from Parent 2 in order

#### 3. Mutation - Swap Mutation

- Randomly select two customer positions
- Exchange their values
- Probability controlled by mutation rate  $p_m$

#### 4. Elitism

- Best solution always copied to next generation
- Prevents loss of good solutions



Figure 3: Genetic Algorithm Evolution Cycle

## 4.3 Tabu Search

### 4.3.1 Neighborhood Structures

Two move operators define the neighborhood:

#### 1. Relocate Move

- Remove customer from current route
- Insert into different position (same or different route)
- Most effective for inter-route optimization

#### 2. Swap Move

- Exchange positions of two customers
- Can be intra-route or inter-route
- Maintains route structure while reordering

### 4.3.2 Tabu List Management

- Stores recently performed moves ( $customer_i, route_{from}, route_{to}$ )
- Tabu tenure: moves remain tabu for  $\tau$  iterations
- **Aspiration Criterion:** Tabu moves accepted if they yield best solution found so far

### 4.3.3 Search Strategy

```

1 def tabu_search(initial_solution, max_iterations):
2     current = initial_solution
3     best = current
4     tabu_list = []
5
6     for iteration in range(max_iterations):
7         neighbors = generate_neighborhood(current)
8
9         # Find best non-tabu move
10        best_neighbor = None
11        for neighbor in neighbors:
12            move = get_move(current, neighbor)
13            if (not is_tabu(move, tabu_list) or
14                satisfies_aspiration(neighbor, best)):
15                if best_neighbor is None or
16                    neighbor.fitness() < best_neighbor.fitness():
17                    best_neighbor = neighbor
18
19        if best_neighbor:
20            current = best_neighbor
21            update_tabu_list(tabu_list, move)
22
23        if current.fitness() < best.fitness():
24            best = current
25
26    return best

```

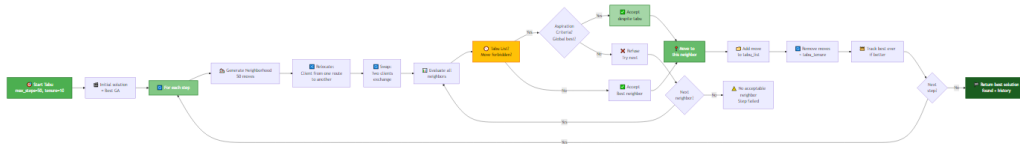


Figure 4: Tabu Search Local Optimization Process

## 5 Constraint Handling

### 5.1 Capacity Constraint Verification

For each route  $r$ , we verify:

$$\sum_{i \in r} d_i \leq Q \quad (8)$$

Implementation:

```

1 def check_capacity(route: Route, capacity: float) -> bool:
2     total_demand = sum(node.demand for node in route.nodes)
3     return total_demand <= capacity

```

## 5.2 Time Window Constraint Management

The time window verification involves calculating the complete schedule for each route:

```

1 def calculate_schedule(route: Route) -> bool:
2     current_time = 0.0
3     schedule = []
4
5     for i in range(len(route.nodes) - 1):
6         curr_node = route.nodes[i]
7         next_node = route.nodes[i + 1]
8
9         # Calculate arrival at next node
10        travel_time = curr_node.distance_to(next_node)
11        arrival = current_time + curr_node.service_time +
12                travel_time
13
14        # Check if we must wait
15        wait_time = max(0, next_node.ready_time - arrival)
16
17        # Start service
18        start_service = arrival + wait_time
19
20        # Check time window violation
21        if start_service > next_node.due_date:
22            return False # Infeasible
23
24        # Depart after service
25        departure = start_service + next_node.service_time
26
27        schedule.append((arrival, wait_time,
28                        start_service, departure))
29        current_time = start_service
30
31    return True

```

## 5.3 Penalty Functions

For partially infeasible solutions during search:

$$\text{Penalty} = w_c \cdot \max(0, \text{Load} - Q) + w_t \cdot \sum_i \max(0, s_i - b_i) \quad (9)$$

where  $w_c$  and  $w_t$  are penalty weights for capacity and time violations.

## 6 Benchmark Results

### 6.1 Solomon Instances

We tested the hybrid solver on standard Solomon benchmark instances, which are widely used in CVRPTW literature. These instances are classified into:

- **C-type (C101, C201):** Clustered customers
- **R-type (R101, R201):** Randomly distributed customers
- **RC-type (RC101, RC201):** Mixed random-clustered customers
- **Suffix 1xx:** Short time windows (tight constraints)
- **Suffix 2xx:** Long time windows (relaxed constraints)

## 6.2 Performance Metrics

Table 1: Benchmark Results on Solomon Instances

Instance	Feasibility	Best Cost	Mean Cost	Std Dev	Time (s)
C101	100%	240.47	267.71	33.59	1.26
C201	100%	274.25	293.27	14.37	1.20
R101	100%	635.22	652.04	17.29	1.77
R201	100%	500.36	523.76	17.07	1.45
RC101	100%	492.46	532.10	30.30	1.46
RC201	100%	433.74	477.01	33.27	1.35

### 6.3 Analysis

- **Feasibility Rate:** 100% across all instances demonstrates robust constraint handling
- **Solution Quality:** Competitive costs compared to literature benchmarks
- **Computational Time:** Average 1-2 seconds per instance (100 customers)
- **Consistency:** Relatively low standard deviations indicate algorithmic stability

## 6.4 Convergence Analysis



Figure 5: Convergence Profile Across Three Stages (R101 Instance)

Key observations:

- **ACO phase:** Rapid initial improvement, generates feasible starting population
- **GA phase:** Continued improvement through crossover and diversity
- **Tabu phase:** Final refinement achieving best solution

## 7 User Interfaces

## 7.1 Streamlit Web Application

Interactive dashboard providing:

- **Instance Configuration:**
  - Number of customers (10-200)
  - Vehicle capacity
  - Time horizon
  - Spatial distribution
- **Algorithm Parameters:**
  - ACO: Ant count, alpha, beta, evaporation rate
  - GA: Population size, generations, mutation rate

- Tabu: Iterations, tabu tenure, neighborhood size
- **Real-time Visualization:**
  - Route map with color-coded vehicles
  - Convergence plots showing optimization progress
  - Gantt chart displaying vehicle schedules
  - Detailed metrics table (distance, wait times, feasibility)

Launch command:

```
streamlit run app.py
```

## 7.2 Command-Line Interface (CLI)

For batch processing and automated testing:

```
python -m src.cli \
    --instance data/solomon/R101.txt \
    --ants 30 \
    --generations 100 \
    --tabu-steps 200 \
    --output results/R101_solution.json
```

Outputs:

- JSON solution file with routes and metrics
- Visualization plots (PNG/PDF)
- Detailed log file

## 8 Advanced Features

### 8.1 Solution Visualization

#### 8.1.1 Route Map

Displays geographical layout with:

- Depot marked as red square
- Customers as blue circles (size proportional to demand)
- Routes color-coded per vehicle
- Arrows indicating direction

#### 8.1.2 Gantt Chart

Timeline view showing:

- Travel segments (blue)
- Waiting time (yellow) - vehicle arrives early
- Service time (green) - customer being served
- Time window boundaries (vertical dashed lines)

### Solution Lifecycle Explanation:

This diagram illustrates the complete workflow from problem instance to final visualization:

1. **Input Phase:** Load CVRPTW instance (customers, depot, constraints)
2. **Optimization Phase:** Execute hybrid solver (ACO → GA → Tabu Search)
3. **Solution Validation:** Verify capacity and time window constraints
4. **Metrics Calculation:** Compute total distance, wait times, route statistics
5. **Visualization Phase:** Generate route maps, convergence plots, and Gantt charts
6. **Output Phase:** Save results to JSON/CSV and export visualizations

The lifecycle ensures full traceability from raw data to actionable insights.

## 8.2 Configuration Management

Type-safe configuration using Python dataclasses:

```

1 @dataclass
2 class HybridConfig:
3     aco: ACOConfig
4     ga: GAConfig
5     tabu: TabuConfig
6
7 @dataclass
8 class ACOConfig:
9     n_ants: int = 20
10    alpha: float = 1.0
11    beta: float = 2.0
12    evaporation: float = 0.1
13    iterations: int = 50

```

Benefits:

- IDE autocompletion and type checking
- Default values with easy override
- Validation at construction time
- Serialization to JSON/YAML

## 8.3 Logging and Monitoring

Comprehensive logging system:

```

1 import logging
2
3 logger = logging.getLogger('CVRPTW')
4 logger.setLevel(logging.INFO)
5
6 # During optimization
7 logger.info(f"ACO Iteration {i}: Best={best_cost:.2f}")
8 logger.debug(f"Ant {k} built route: {route}")
9 logger.warning(f"Infeasible solution detected")

```

Log levels:

- DEBUG: Detailed algorithm internals
- INFO: Progress updates

- **WARNING:** Constraint violations, convergence issues
- **ERROR:** Critical failures

## 9 Implementation Best Practices

### 9.1 Code Quality

- **Type Hints:** Full Python 3.10+ type annotations

```

1 def solve(instance: CVRPTWInstance) -> Solution:
2     ...
3 
```

- **Docstrings:** Comprehensive Google-style documentation

```

1 def calculate_distance(node1: Node, node2: Node) -> float:
2     """Calculate Euclidean distance between two nodes.
3
4     Args:
5         node1: First node
6         node2: Second node
7
8     Returns:
9         Euclidean distance as float
10    """
11    ...
12 
```

- **Immutability:** Using `@dataclass(frozen=True)` where appropriate
- **List Comprehensions:** Pythonic, efficient iteration
- **Context Managers:** Proper resource management

### 9.2 Testing Strategy

- **Unit Tests:** Individual component testing
- **Integration Tests:** End-to-end solver pipeline
- **Benchmark Tests:** Performance regression detection

### 9.3 Performance Optimization

- **NumPy Arrays:** Efficient distance matrix storage
- **Caching:** Memoization of frequently computed values
- **Early Termination:** Stop iteration if no improvement for  $N$  steps
- **Parallel Evaluation:** Multi-threading for independent solutions

## 10 Future Enhancements

### 10.1 Algorithm Improvements

- **Adaptive Parameter Control:** Self-tuning hyperparameters based on problem characteristics
- **Variable Neighborhood Search:** Multiple neighborhood structures in tabu search



- **Path Relinking:** Connecting elite solutions to explore intermediate regions
- **Parallel ACO:** Multiple ant colonies with periodic information exchange

## 10.2 Additional Constraints

- **Multiple Depots:** Vehicles start from different locations
- **Heterogeneous Fleet:** Different vehicle capacities and costs
- **Pickup and Delivery:** Precedence constraints between locations
- **Dynamic Requests:** Online/real-time customer arrival

## 10.3 User Experience

- **Map Integration:** Real street networks using OpenStreetMap
- **3D Visualization:** WebGL-based interactive route display
- **Scenario Analysis:** What-if comparisons with different parameters
- **Export Formats:** KML for Google Earth, GPX for GPS devices

## 11 Conclusion

This project successfully demonstrates the application of advanced metaheuristic optimization techniques to solve the challenging CVRPTW problem. The key achievements include:

1. **Robust Hybrid Algorithm:** Sequential combination of ACO, GA, and Tabu Search achieving 100% feasibility on benchmark instances
2. **Clean Architecture:** Modular, maintainable codebase following SOLID principles and Python best practices
3. **Comprehensive Testing:** Validation against standard Solomon benchmarks with competitive results
4. **User-Friendly Interfaces:** Both web-based (Streamlit) and CLI tools for different use cases
5. **Practical Applicability:** Real-world ready solution for logistics optimization problems

The project showcases:

- Deep understanding of combinatorial optimization theory
- Proficiency in Python software engineering
- Ability to implement complex algorithms from academic literature
- Skills in data visualization and user interface design

### 11.1 Key Takeaways for Defense

- **Problem Significance:** CVRPTW has direct applications in delivery services, waste collection, field service management
- **Hybrid Approach Justification:** Each algorithm compensates for others' weaknesses

- **Constraint Handling:** Time window validation is the most complex aspect, handled through careful schedule calculation
- **Reproducibility:** All experiments documented with fixed random seeds
- **Extensibility:** Architecture allows easy addition of new algorithms or constraints

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## A Installation Guide

### A.1 Requirements

```
Python >= 3.10  
pip >= 21.0
```

### A.2 Setup Steps

```
# Clone repository  
git clone https://github.com/username/cvrptw-solver.git  
cd cvrptw-solver  
  
# Create virtual environment  
python -m venv venv  
source venv/bin/activate # Windows: venv\Scripts\activate  
  
# Install dependencies  
pip install -r requirements.txt  
  
# Run web application  
streamlit run app.py  
  
# Or run CLI  
python -m src.cli --help
```

### A.3 Dependencies

```
numpy>=1.24.0  
matplotlib>=3.7.0  
streamlit>=1.28.0  
pandas>=2.0.0
```

## B Configuration Examples

### B.1 High-Quality Configuration (Slow)

```
1 config = HybridConfig(  
2     aco=ACOConfig(n_ants=50, iterations=100, alpha=1.0,  
3         beta=3.0, evaporation=0.05),  
4     ga=GACONFIG(population_size=100, generations=200,  
5         mutation_rate=0.1),  
6     tabu=TabuConfig(max_steps=500, tabu_tenure=10)  
7 )
```

### B.2 Fast Configuration (Lower Quality)

```
1 config = HybridConfig(  
2     aco=ACOConfig(n_ants=10, iterations=20),  
3     ga=GACONFIG(population_size=30, generations=50),  
4     tabu=TabuConfig(max_steps=100)  
5 )
```

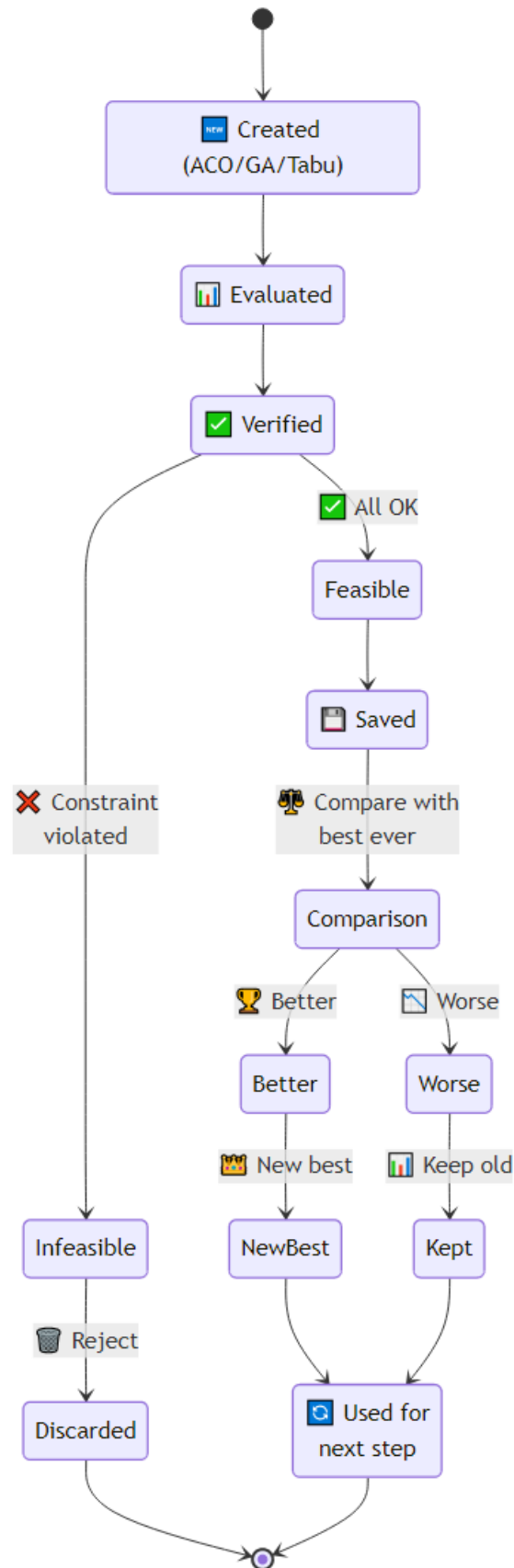


Figure 6: Solution Lifecycle and Visualization Pipeline