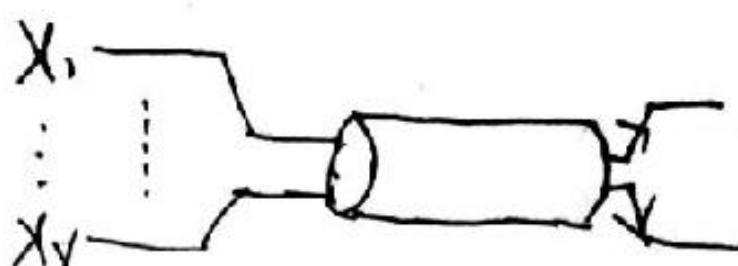


Multiplexing: A / B

Part 3/4

Consider the model as follows:



the nontrivial aspect of the problem is that \$v\$ is a random variable.

Assume the users are active independently, then we have \$v\$ is Binomial \$(N, p) / \text{Bin}(N, p)\$

① Gaussian Random Variable (the binomial distribution is well approximated by Gaussian)

1) standard normal RV: $f_W(x) = \frac{1}{\sqrt{2\pi}} \exp\{-\frac{x^2}{2}\}$ $x \in \mathbb{R} \Leftrightarrow W \sim \mathcal{N}(0, 1)$

2) $X \sim \mathcal{N}(\mu, \sigma^2)$ we have $X = \mu + \sigma W$ so we have pdf: $f_X(x) = \frac{1}{\sqrt{2\pi}\sigma} \exp\{-\frac{(x-\mu)^2}{2\sigma^2}\}$

② CLT (Central Limit Theory) Let $\{X(n), n \geq 1\}$ be RV with $E(X(n)) = \mu$ and $\text{var}(X(n)) = \sigma^2$ then as $n \rightarrow \infty$.

$$\frac{X(1) + \dots + X(n) - n\mu}{\sigma\sqrt{n}} \Rightarrow \mathcal{N}(0, 1) \quad (\Rightarrow \text{means convergence in distribution})$$

So we have $X(n) \xrightarrow{\text{d.s.}} X$ implies $X(n) \xrightarrow{P} X$ implies $X(n) \Rightarrow X$

③ Confidence interval

$$\left[\mu_n - 1.65 \frac{\hat{\sigma}_n}{\sqrt{n}}, \mu_n + 1.65 \frac{\hat{\sigma}_n}{\sqrt{n}} \right] = 95\% \quad \left[\mu_n - 2 \frac{\hat{\sigma}_n}{\sqrt{n}}, \mu_n + 2 \frac{\hat{\sigma}_n}{\sqrt{n}} \right] = 95\%$$

$$\text{where } \hat{\sigma}_n^2 = \frac{n}{n-1} \left\{ \frac{\sum_{m=1}^n X(m)^2}{n} - \mu_n^2 \right\} \quad \mu_n = \frac{\sum_{m=1}^n X(m)}{n}$$

④ Buffer: at any time \$n\$, a transmission completes with probability \$\mu\$ and a new packet arrives with probability \$\lambda\$, independently of the past, \$X_n\$ is a Markov chain

$$p_2 = \lambda(1-\mu) (X(n) \rightarrow X(n+1)), \quad p_0 = \mu(1-\lambda) (X(n+1) \rightarrow X(n)), \quad p_1 = 1 - p_0 - p_2 (X(n) \rightarrow X(n))$$

so we have balance equations:

$$\begin{cases} z(0) = (1-p_2)z(0) + p_0z(1) \\ z(n) = p_2z(n-1) + p_1z(n) + p_0z(n+1) \quad 1 \leq n \leq N-1 \\ z(N) = p_2z(N-1) + (1-p_0)z(N) \end{cases} \Rightarrow \begin{cases} z(i) = z(0)p^i, i=0,1,\dots,N \text{ where } p := \frac{p_2}{p_0} \\ z(0) = \left[\sum_{i=0}^N p^i \right]^{-1} = \frac{1-p}{1-p^{N+1}} \end{cases}$$

$$E(X) = \sum_{i=0}^N i z(i) = z(0) \sum_{i=0}^N i p^i \approx \frac{p}{1-p} = \frac{p_2}{p_0 - p_2} = \frac{\lambda(1-\mu)}{\mu-\lambda} \quad (\text{average packet number})$$

(Average Delay) $W = \sum_{k=0}^{\infty} \frac{k+1}{\mu} \phi(k)$ where $\phi(k) = z(k)(1-\mu) + z(k+1)\mu$ substitute it, we have

$$W = \frac{1-\mu}{\mu-\lambda} = \frac{1}{\lambda} E(X)$$

It is tempting to conclude that the probability that a packet finds k packet in the buffer upon its arrival is $P[X_{n+1} = k+1 | A_n = 1] = P[X_n = k | A_n = 1] = P[X_n = k] = z_k(k)$

Little's Law) $L = \lambda W$

L is the average number of customers in a system, λ is the average arrival rate of customer, and W is the average time that a customer spend in the system.

⑤ Multiple Access

the fraction of time that exactly one device transmits is $P(X(n)=1) = Np(1-p)^{N-1}$

the maximum over p of this success rate occurs when $p = 1/N \Rightarrow \lambda^* = (1 - \frac{1}{N})^{N-1} \approx \frac{1}{e}$

⑥ characteristic Functions

$$\phi_X(u) = E(e^{iuX}) \quad u \in \mathbb{R} \Leftrightarrow \phi_X(u) = \int_{-\infty}^{\infty} e^{iuX} f_X(x) dx$$

let $X = \sum_{i=1}^N U_i$ Then $\phi_X(u) = e^{-\frac{u^2}{2}}$ (proof: p 60)

⑦ Moments of $N(0,1)$

$$\begin{aligned} \phi_X(u) &= E(e^{iuX}) = E\left(\sum_{n=0}^{\infty} \frac{1}{n!} (iuX)^n\right) = \sum_{n=0}^{\infty} \frac{1}{n!} (iu)^n E(X^n) \\ \text{while } \phi_X(u) &= e^{-\frac{u^2}{2}} = \sum_{m=0}^{\infty} \frac{1}{m!} \left(-\frac{u^2}{2}\right)^m = \sum_{m=0}^{\infty} \frac{1}{m!} \left(-\frac{1}{2}\right)^m u^{2m} \end{aligned} \quad \left. \begin{array}{l} \text{match the coefficients} \\ \text{of } u^{2m}, \text{ and we can find} \end{array} \right\}$$

$$\frac{1}{(2m)!} i^{2m} E(X^{2m}) = \frac{1}{m!} \left(-\frac{1}{2}\right)^m \Rightarrow E(X^{2m}) = \frac{(2m)!}{m! 2^m}, \quad E(X^{2m+1}) = 0$$

⑧ 1) Poisson as a limit of Binomial: $P(\lambda) \sim B(n, \lambda/n)$

2) Exponential as limit of Geometric: $\text{Exp}(\lambda) \sim G(\lambda/n)$

⑨ Error Function $Q(x) = P(X > x)$ where $X \sim N(0,1)$

Bounds on EF: $\frac{x}{1+x^2} \frac{1}{\sqrt{2\pi}} \exp\left\{-\frac{x^2}{2}\right\} \leq Q(x) \leq \frac{1}{x\sqrt{2\pi}} \exp\left\{-\frac{x^2}{2}\right\} \quad \forall x > 0$

⑩ Adaptive Multiple Access

$$p(n+1) = \begin{cases} p(n) & \text{if } X(n)=1 \\ ap(n) & \text{if } X(n)=1 \\ \min\{bp(n), 1\} & \text{if } X(n)=0 \end{cases}$$