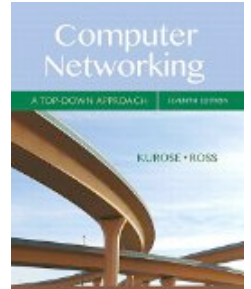


# COMP 375: Lecture 17



- **News & Notes:**
  - Project #2 demos scheduling soon
  - Project #3 due Fri, March 16
- **Reading (Fri, Feb. 9)**
  - Review Sections 3.4.{2-4}

Section 3.4

# **RELIABLE DATA TRANSFER**

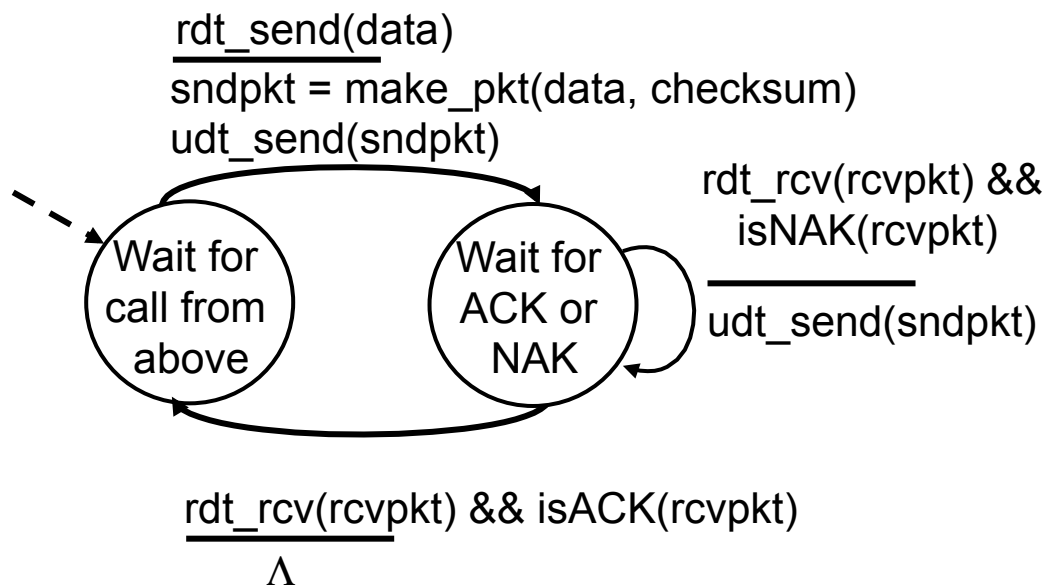
**A**utomatic **R**epeat **R**e**Q**uest (ARQ) protocols are similar to the protocol you use for cell phones.



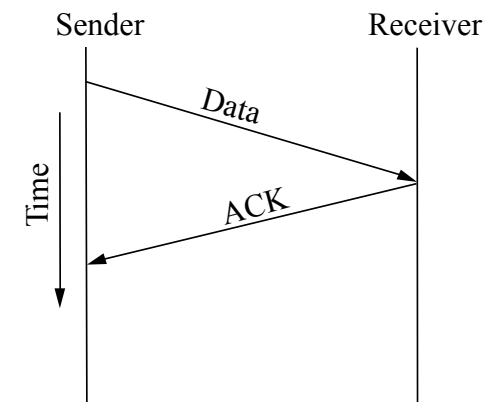
PHOTO: CARL PENDLE/GETTY IMAGES

# We'll focus on two ways to view ARQ protocols: FSMs and timelines.

## Finite State Machine (FSM)

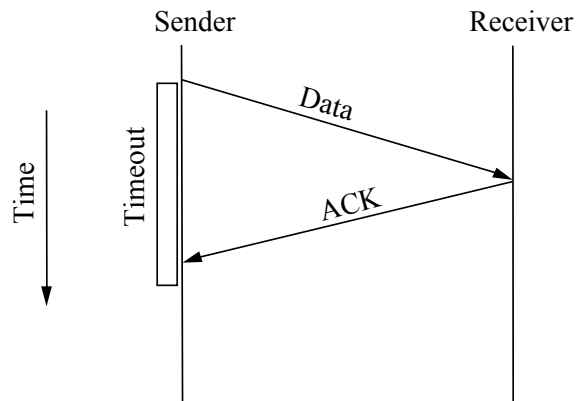


## Timeline

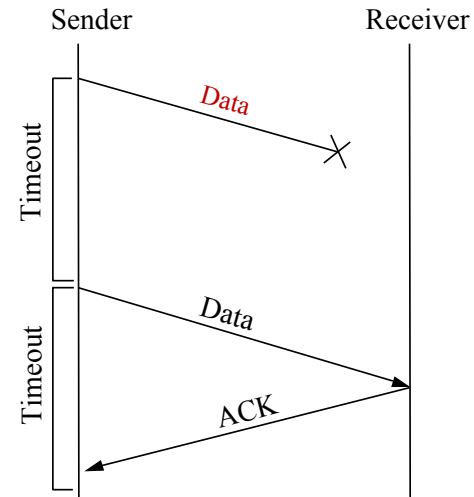


In ARQ, receiver sends ACKs when it receives data.

**Scenario 1:  
No Data Loss**

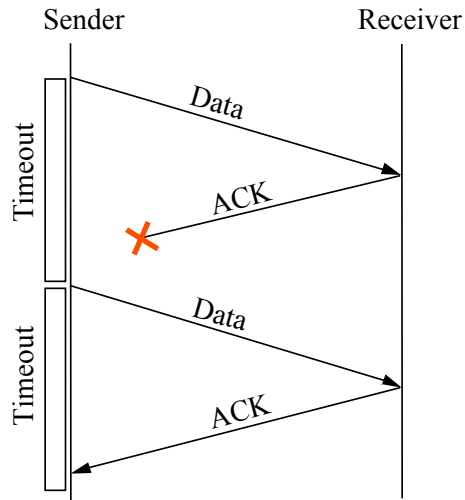


**Scenario 2:  
Data Loss**

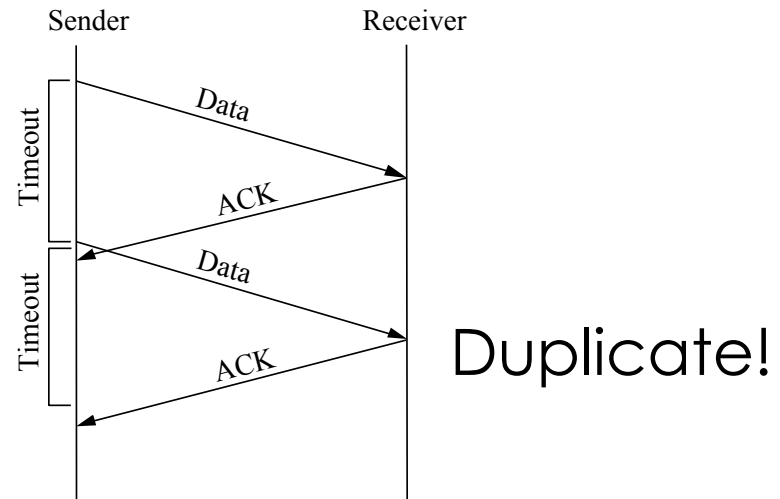


ACKs can get lost/corrupted too,  
which leads to duplication.

### Scenario 3: ACK Loss



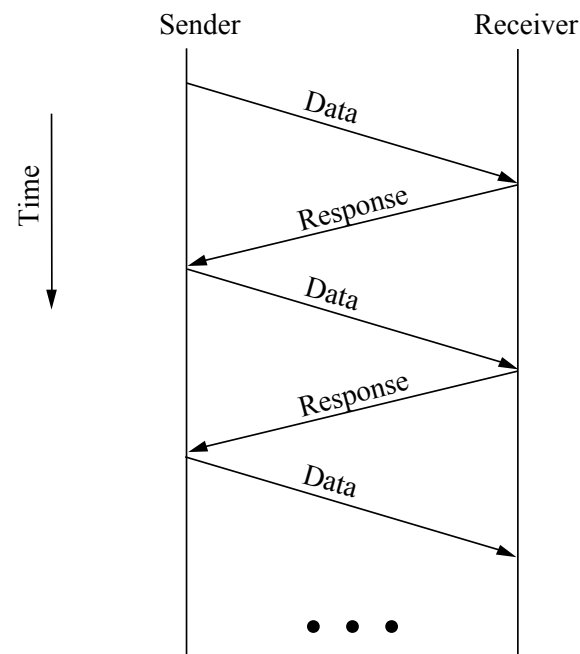
### Scenario 4: Slow ACK



ARQ protocols differ on when they send new messages and what they do after data loss.

- Non-pipelined:
  - Stop-and-wait
- Pipelined
  - Go-back-N (GBN)
  - Selective Repeat (SR)

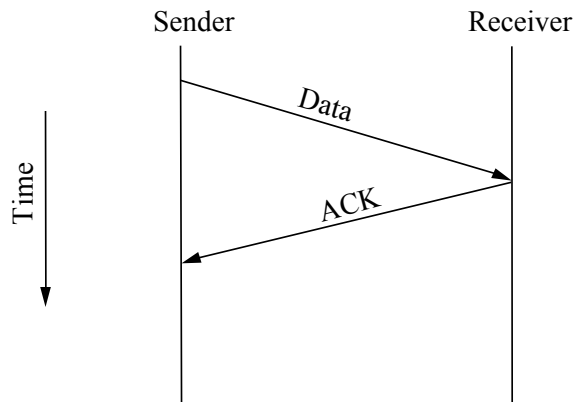
In **Stop-and-Wait**, sender only sends subsequent data after getting response.



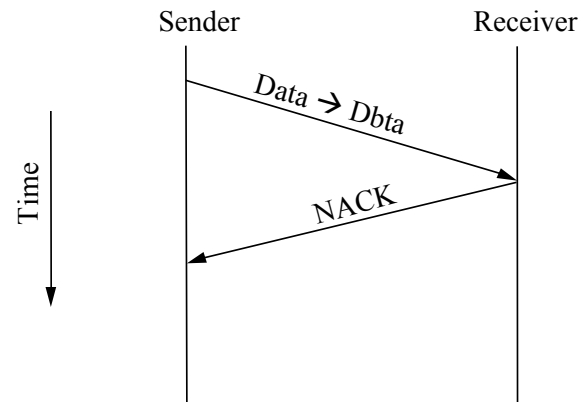


We can use checksums to detect errors, replying with ACK or NACK.

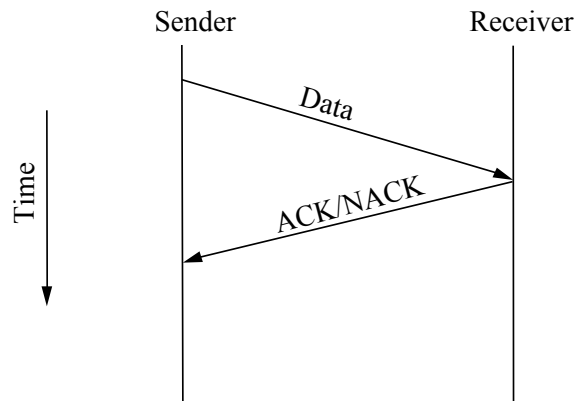
### Scenario 5: Uncorrupted Data



### Scenario 6: Corrupted Data

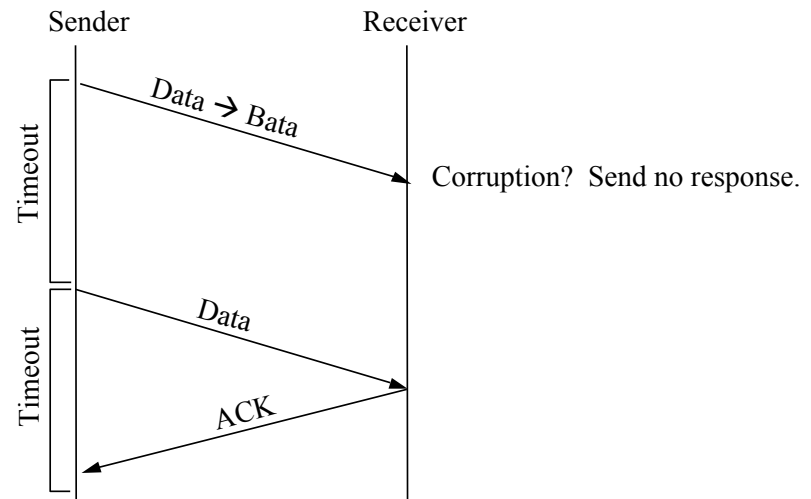
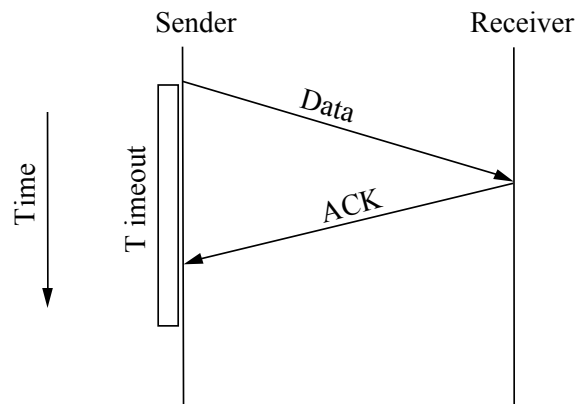


# Could we do this with just ACKs or just NACKs?

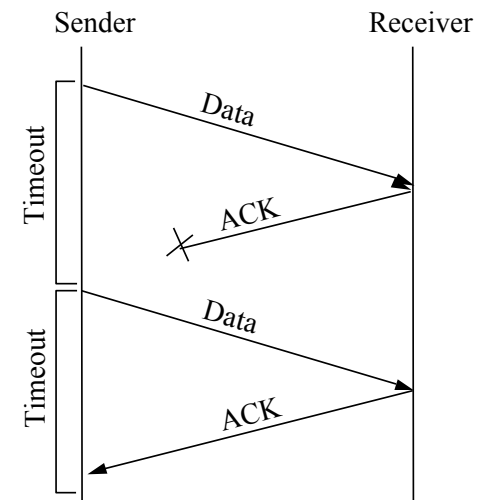
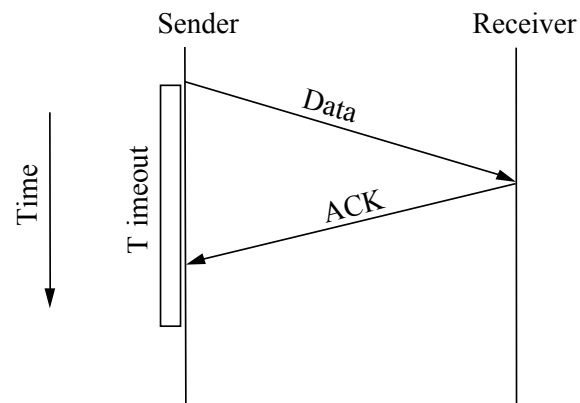


- |           |   |
|-----------|---|
| A.        | <b>No</b> , we need them both.  |
| <b>B.</b> | <b>Yes</b> , we could do without one of them, but we'd need some other mechanism. |
| C.        | <b>Yes</b> , we could get by without one of them.                                 |

With timeouts, we don't need NACK, but we can't avoid ACKs.

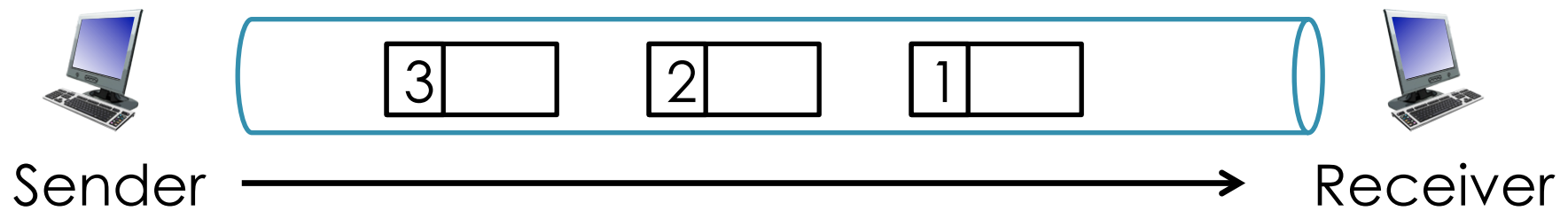


Adding timeouts might create new problems for us to worry about. How many? Examples?



- |    |                            |
|----|----------------------------|
| A. | No new problems            |
| B. | One new problem            |
| C. | Two new problems           |
| D. | More than two new problems |

Adding **sequence numbers** helps us handle duplication.



When using stop-and-wait, we only need a 1-bit sequence number (i.e. 0 or 1).

*Why is that the case?*

- |    |   |
|----|---|
| A. | We only have two possible types of packets, original and duplicate.   |
| B. | Since this is stop-and-wait, we only have one not-yet-ACKed packet in flight.   |
| C. | Having more than two sequence numbers would blow up the number of states in our finite state machine to an unreasonable size. |