

# JOHN BOLTON ON OBAMA'S INTERNET HANDOVER: 'WITHIN TEN YEARS, THE INTERNET AS WE KNOW IT WILL END'

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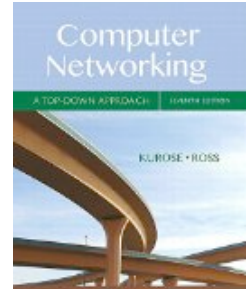
by JOHN HAYWARD | 22 Sep 2016

2,410

On Thursday's Breitbart News  
Daily on SiriusXM, former

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# COMP 375: Lecture 10



- **News & Notes:**

- Class / Office Hours cancelled Wed - Fri
- Project #2 due Wed, Feb. 28

- **Reading (Mon, Feb. 26)**

- Sections 3.1 – 3.3

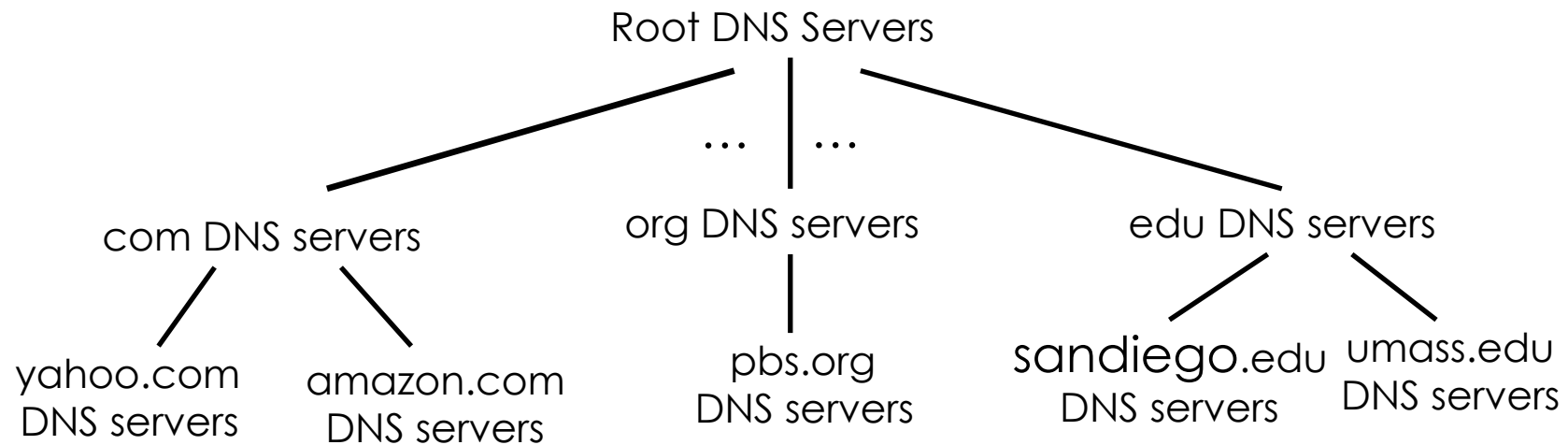
Section 2.4

# **DOMAIN NAME SERVICE (DNS)**

A DNS record contains a TTL that informs how long it is cached.

- *What is the trade-off involved with setting TTL to a large value rather than a small one.*

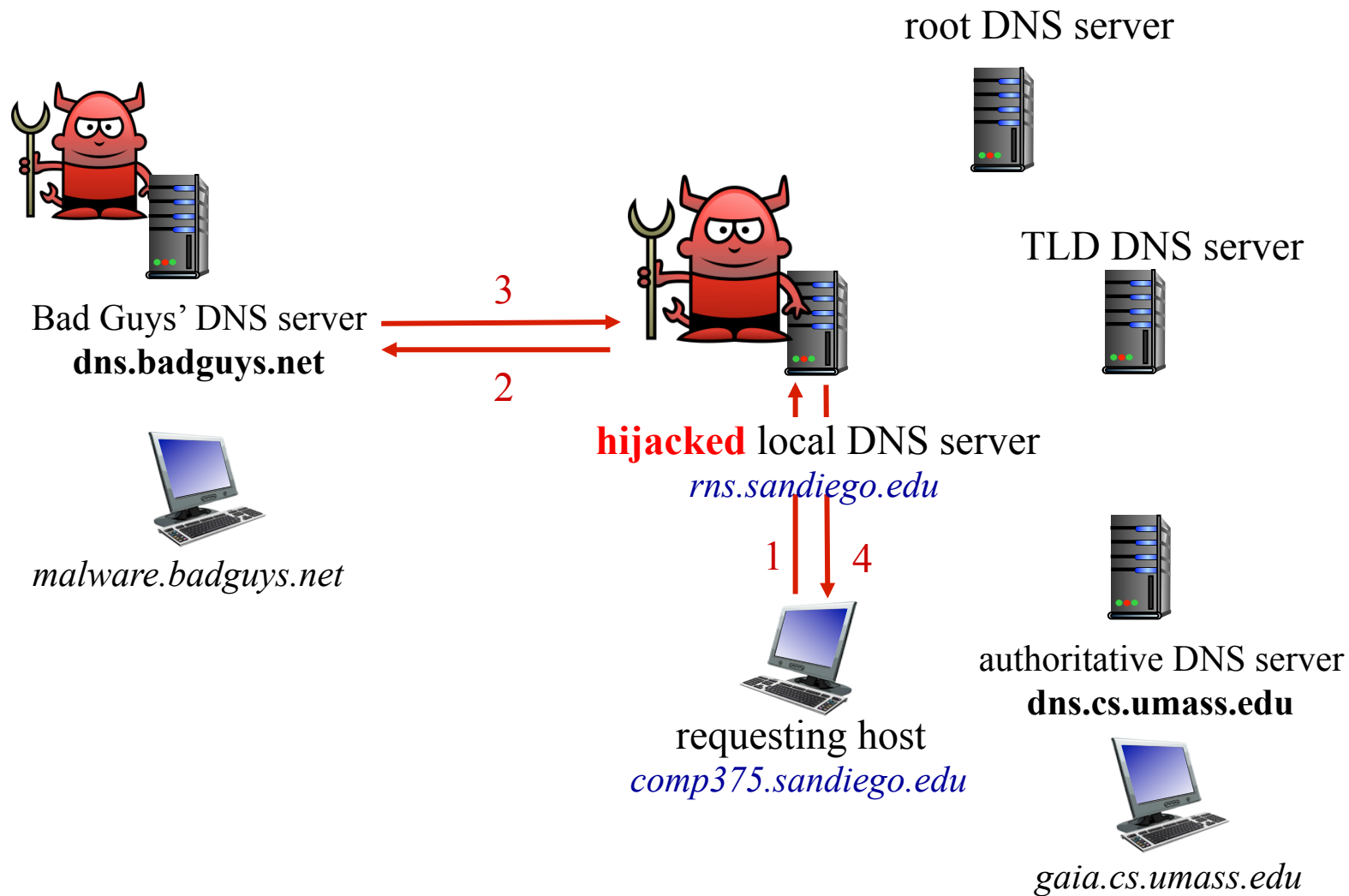
What would change if we wanted to add a new domain (e.g. DrSat.com)?



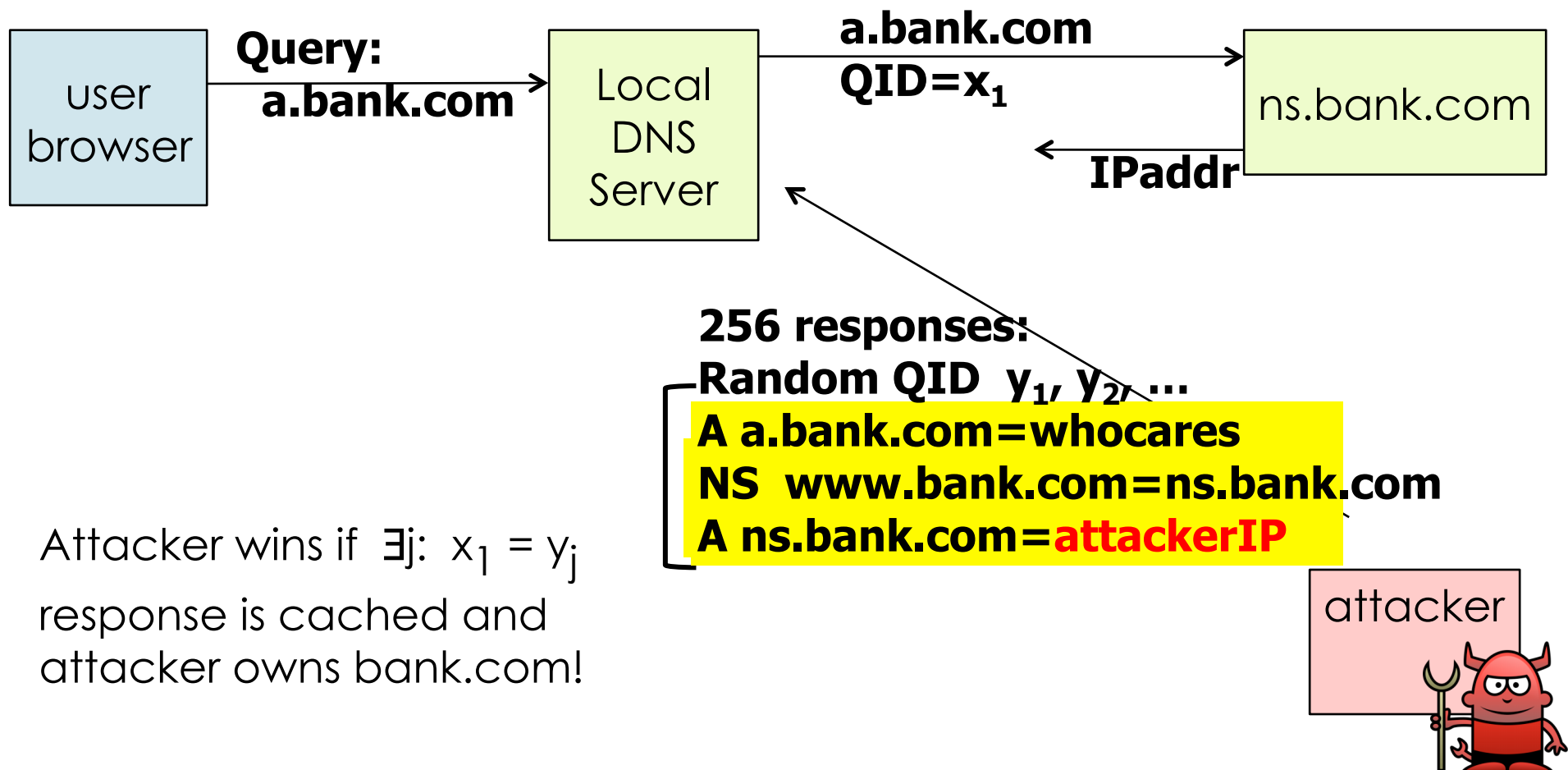
## What aspect of DNS seems like it is particularly insecure?

- |           |                                                                            |
|-----------|----------------------------------------------------------------------------|
| <b>A.</b> | UDP sockets accept packets from any host on the entire Internet.           |
| <b>B.</b> | The TTL allows a mapping to persist for new queries after it has been set. |
| <b>C.</b> | Attackers can control when recursive DNS servers look up specific domains. |
| <b>D.</b> | Responses can have additional “helpful information” appended.              |

# DNS is vulnerable when someone hijacks a DNS server.

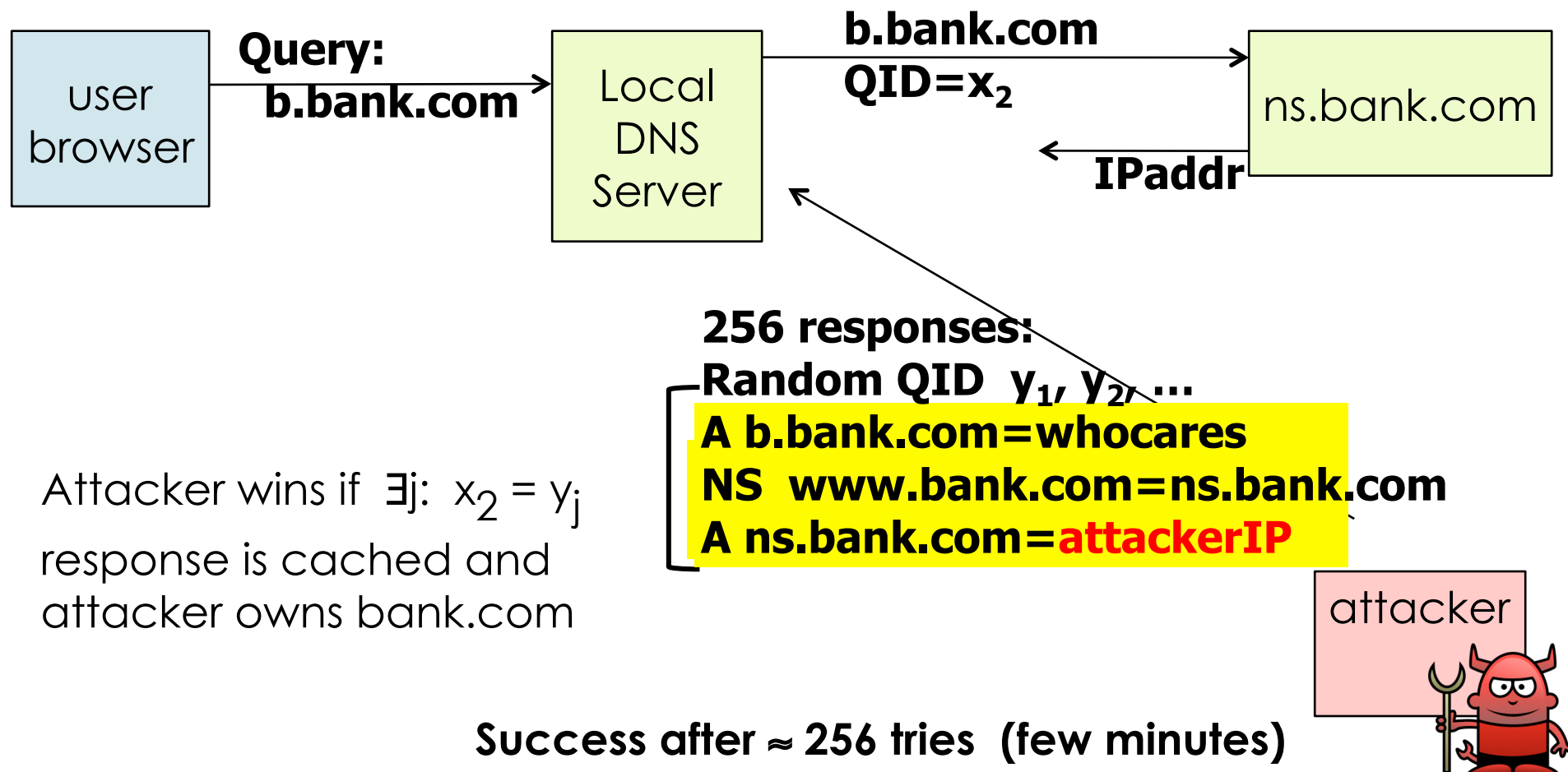


DNS caches can be poisoned by giving bogus additional information.





# Flooding Local DNS with random requests increases likelihood of success.



There are several proposals for improving DNS security.

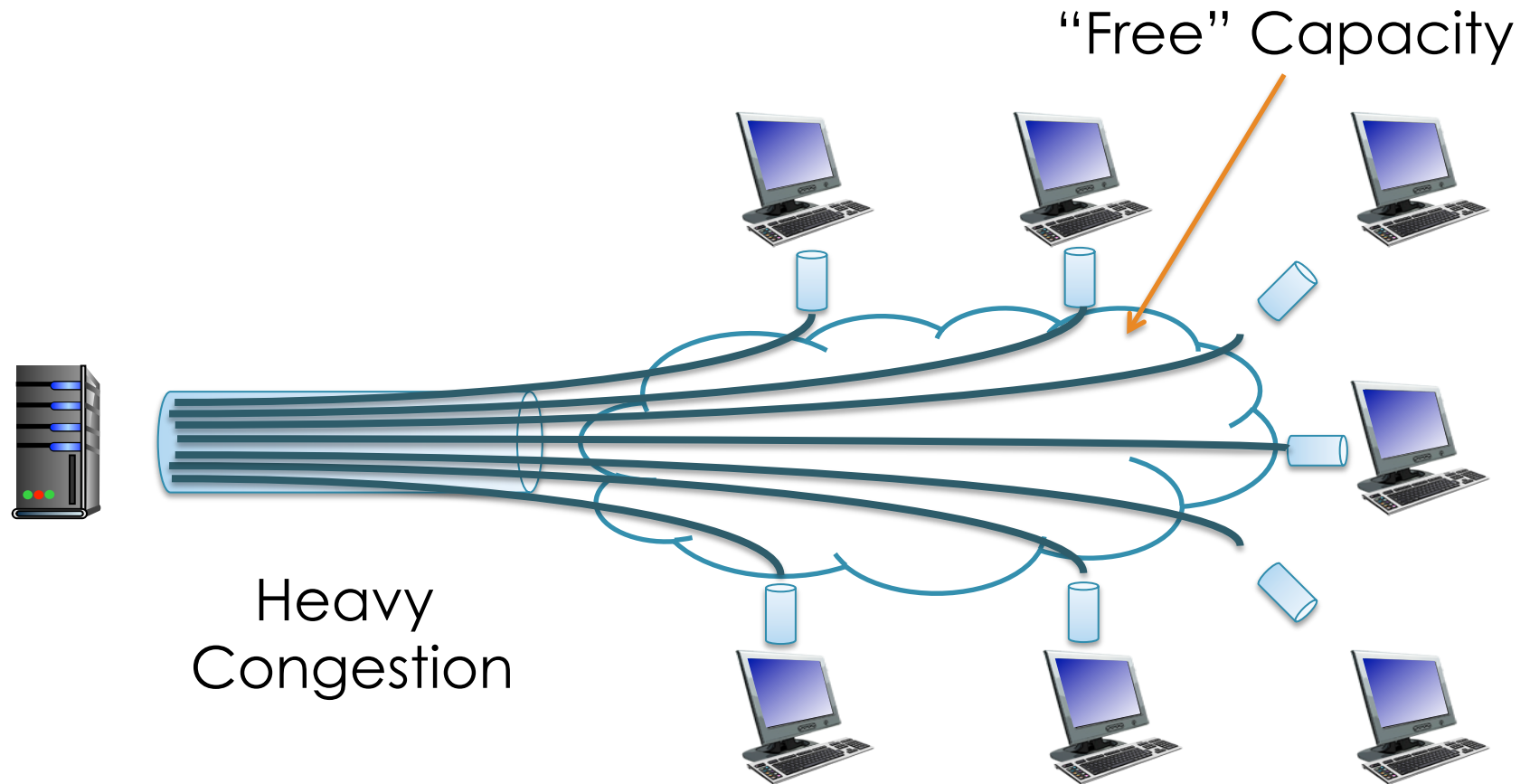
- Use authentication (DNSSEC)
- Reduce probability of success
  - Increase size of query ID
  - Randomize UDP “Source” Port
  - Ask each query twice

Section 2.5

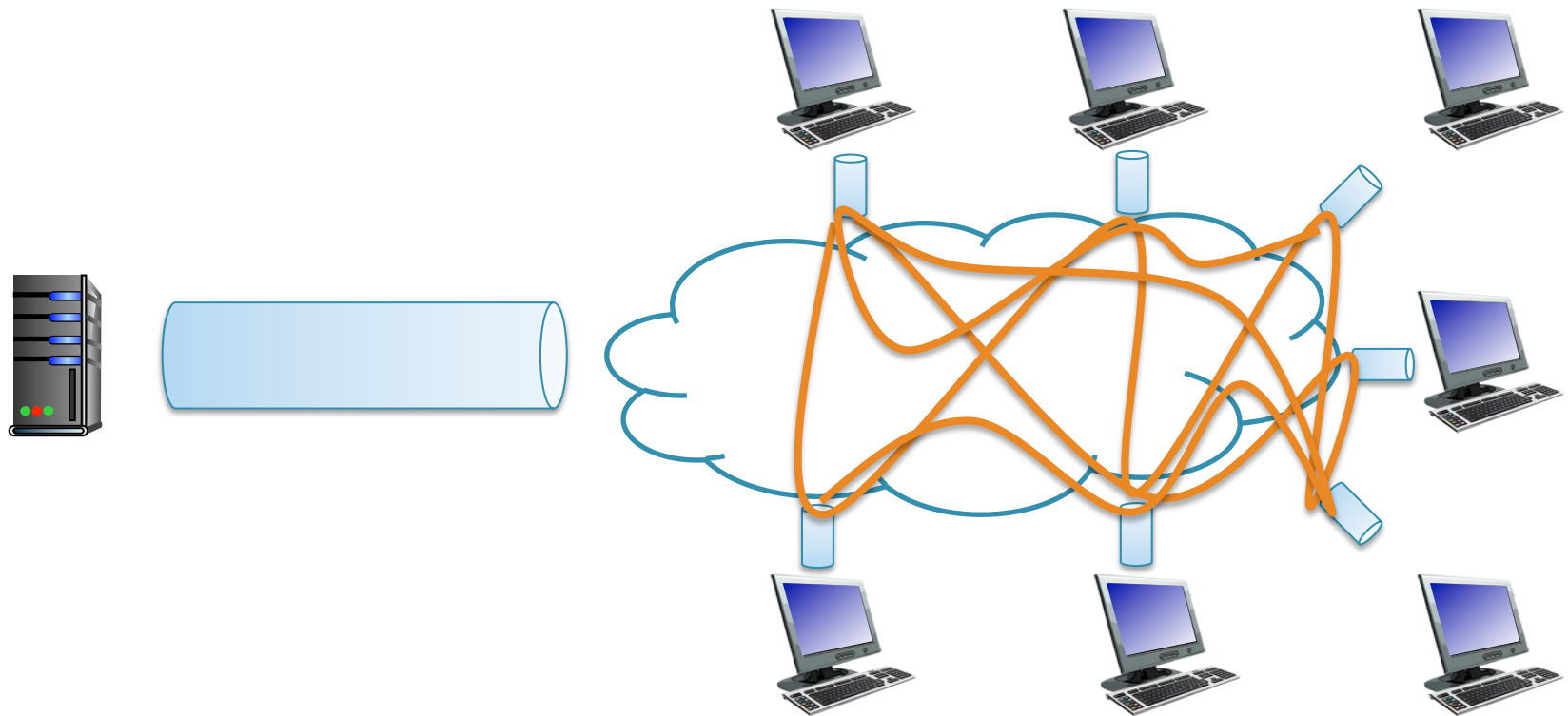
# **PEER-TO-PEER FILE DISTRIBUTION**

What is the best way to distribute a file to a large number of people?

In the Client/Server model, a server has a heavier burden than clients.

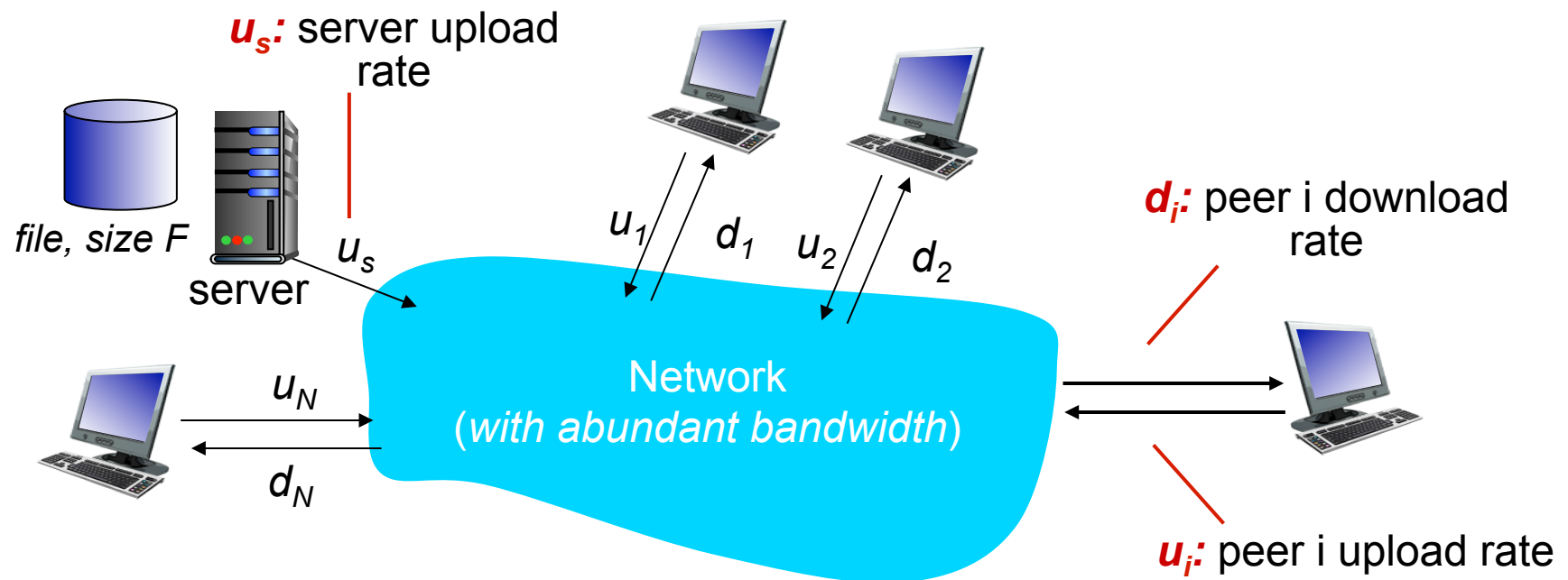


In the P2P model, everyone shares the burden.



Determining file distribution time will require knowledge of transfer rates of all entities.

How much time to distribute a file (size:  $F$ ) from one server to  $N$  peers?

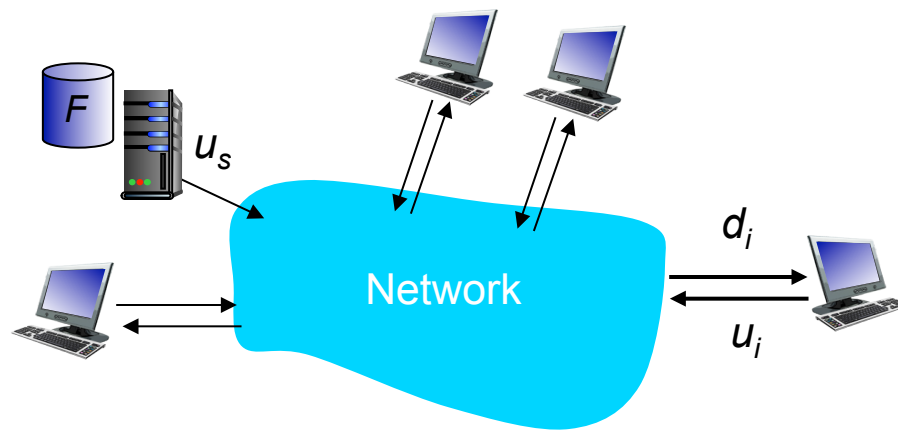


What dictates the **minimum** time for a server to distribute a file of size  $F$  to  $N$  clients?

- |           |                                                                                                                                                                     |
|-----------|---------------------------------------------------------------------------------------------------------------------------------------------------------------------|
| A.        | The MAX of the time it takes any client to download the file.                                                                                                       |
| <b>B.</b> | The MAX of the time it takes for the server to upload $N \cdot F$ bits, <b>and</b> the time it takes the client with the lowest download rate to download the file. |
| C.        | The MIN of the time it takes for the server to upload $N \cdot F$ bits, <b>and</b> the time it takes the client with the lowest download rate to download the file. |



Time to distribute with C/S usually limited by upload rate of the server.

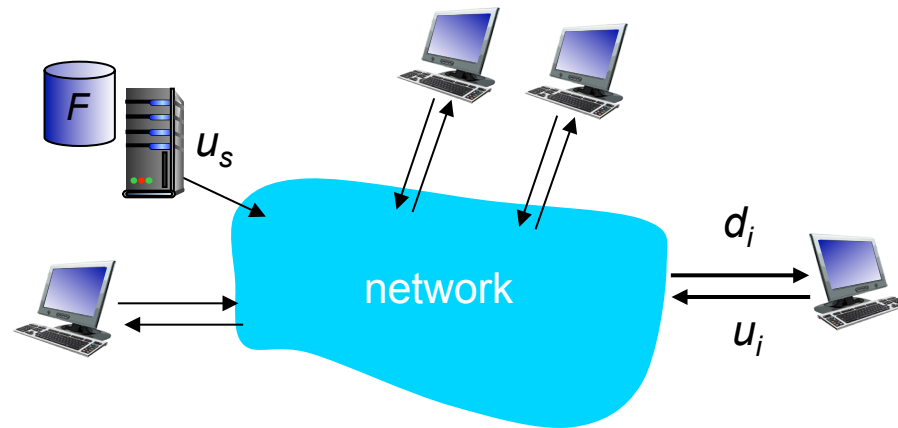


$$D_{c-s} \geq \max\{NF/u_s, F/d_{min}\}$$

What dictates the **minimum** time to distribute the same file in a P2P network?

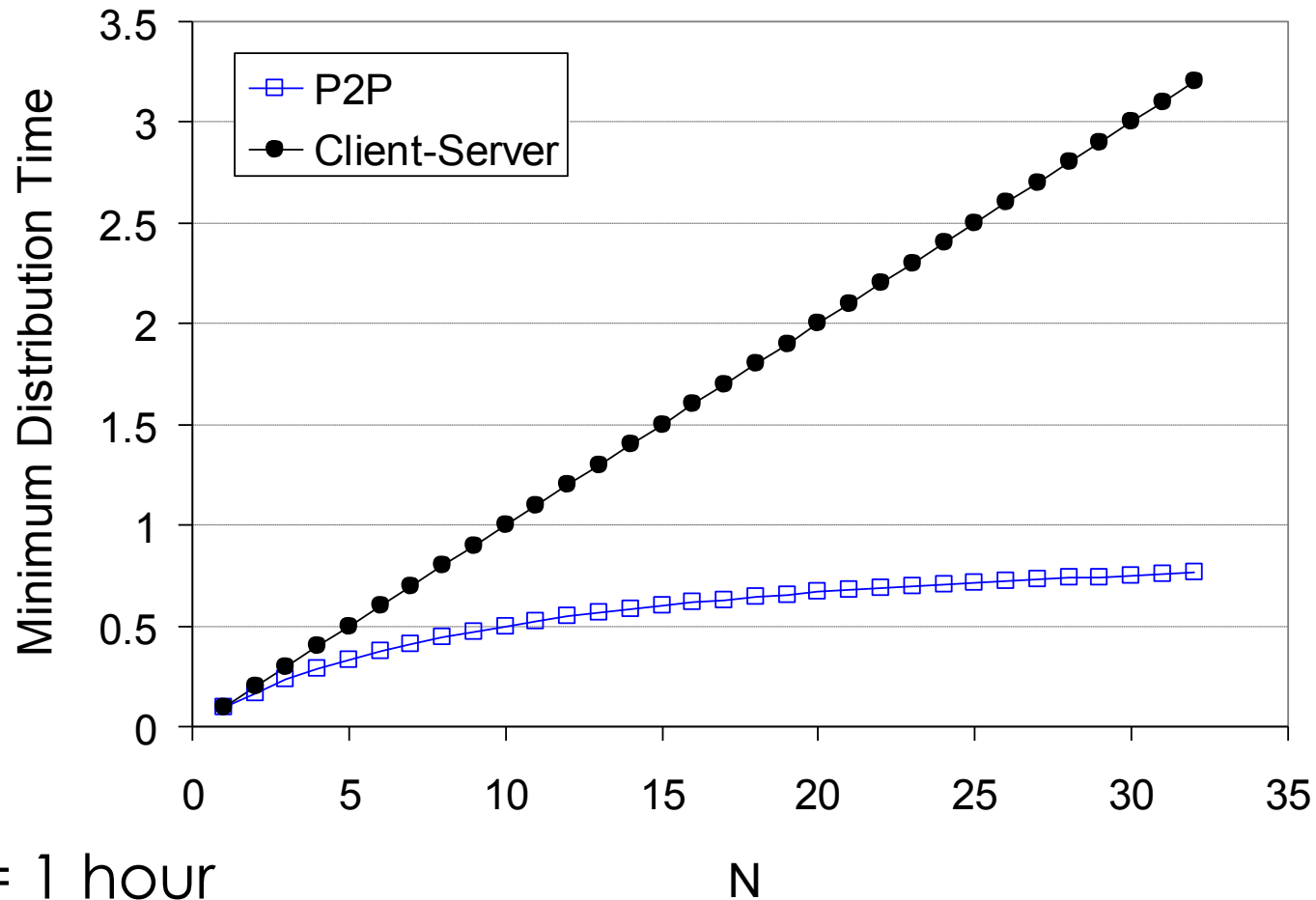
- |    |                                                                               |
|----|-------------------------------------------------------------------------------|
| A. | $N \cdot F$ divided by the total upload bandwidth of peers and server.        |
| B. | MAX of answer A and the time it takes the original server to upload the file. |
| C. | MAX of answer B and the download speed of the <i>slowest</i> peer             |

P2P distribution time *usually* limited by aggregate upload bandwidth.



$$D_{P2P} \geq \max\{F/u_s, F/d_{min}, NF/(u_s + \sum u_i)\}$$

C/S increases linearly with time, P2P is logarithmic.



- $F/u_i = 1$  hour
- $u_s = 10 * u_i$
- $d_{min} \geq u_s$