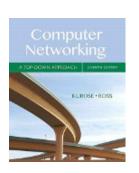
#### COMP 375: Lecture 15



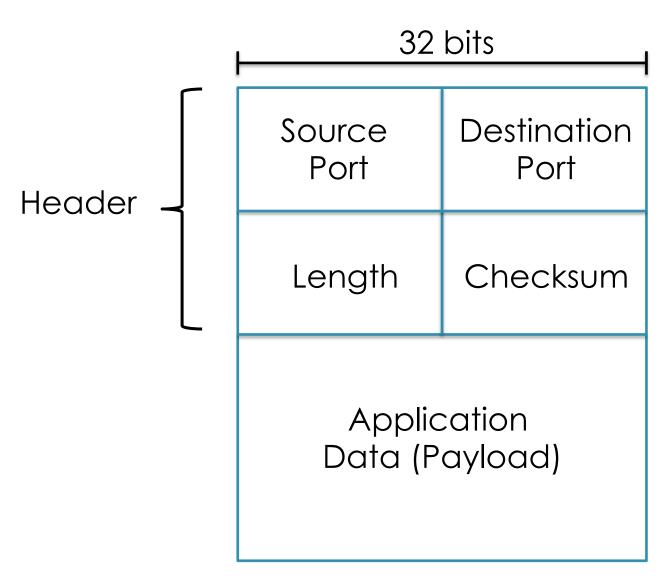
#### News & Notes:

- Midterm #1 in class Monday
- Project #2 due @ 10PM
  - 1 project buck for 24 hour extension, 2 project bucks extends until Tuesday @ 10PM
- Cool new website: www.hackterms.com
- Reading (Wed, Mar. 7)
  - (Tentative) Sections 3.4.{2-4}

Section 3.3



### UDP is **connectionless**, with **best effort** service and **minimal overhead**.



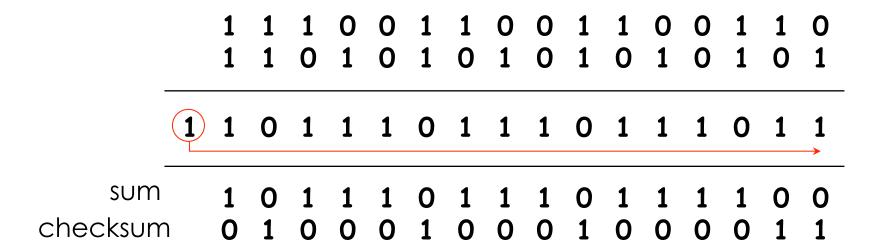
## UDP's **checksum** is designed to detect transmission errors.

- Checksum Calculation:
  - Divide data into 16-bit chunks
  - Calculate the sum of all of these chunks
  - 3. Take the 1's complement of sum

## One's Complement is simply the bitwise not (~) operator.



### What is the checksum of the UDP message 0xE666D555?



**Note:** When adding numbers, a carryout from the most significant bit needs to be added to the result.

## Verifying the checksum follows a similar process to calculating it.

- Validation Algorithm:
  - Divide received data into 16-bit chunks
  - Calculate the sum of all these chunks
  - 3. Add the sum to the received checksum
    - If result is not 0xFFFF, there were errors
    - > Otherwise, ...?

Does sum + checksum == 0xFFFF guarantee there were no transmission errors?

A. Yes!

B No!

## There are three classes of programs that benefit from using UDP.

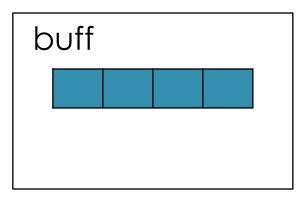
- 1. Latency sensitive
- 2. Error correction unnecessary
- 3. Communicating with *lots* of others

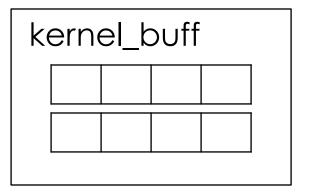
What if you want something more reliable than UDP, but faster/not as full featured as TCP?

- A. Sorry, you're out of luck. 🕾
- B. Write your own transport protocol.
- C Add in the features you want at the application layer.

# The kernel (OS) buffers data before sending it out over the network.

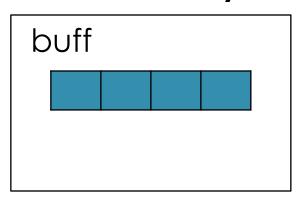
#### **Process' Memory**



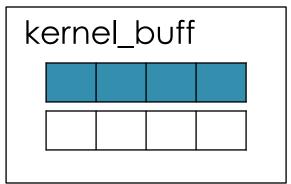


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#### **Process' Memory**



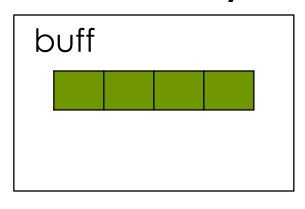


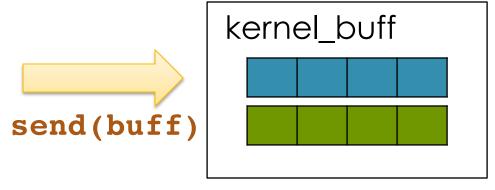




## By waiting, can reduce number of network transmissions.

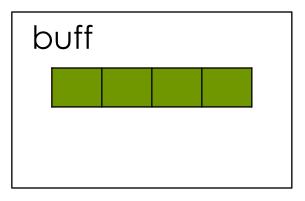
#### **Process' Memory**

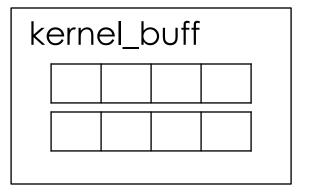




## By waiting, can reduce number of network transmissions.

#### **Process' Memory**







## TCP blocks (waits) if kernel buffer is full. Should UDP do the same?

- A. Yes. It needs to and therefore it should.
- **B. Yes**. It doesn't need to, but it might be useful.
- C. No. It does not need to and should not do so.

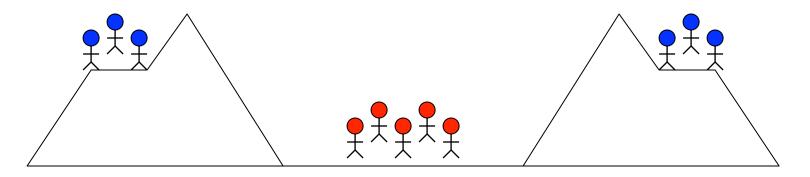
## Project tip: Split your HTTP responses into multiple sends.

```
void sendFile(int sock, path file)
{
    size_t fsize = fs::file_size(file);
    sendOKHeader(sock, fsize);
    readAndSendFileData(sock, file);
}
```

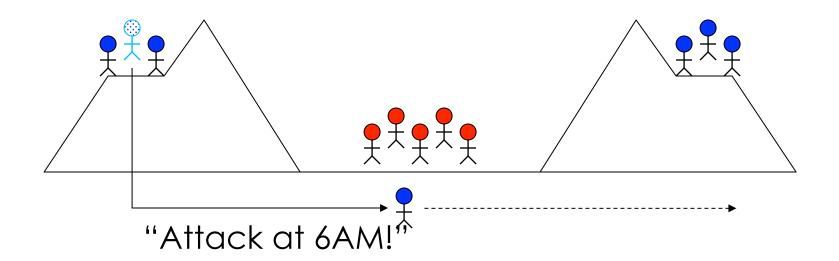
#### Section 3.4

#### RELIABLE DATA TRANSFER

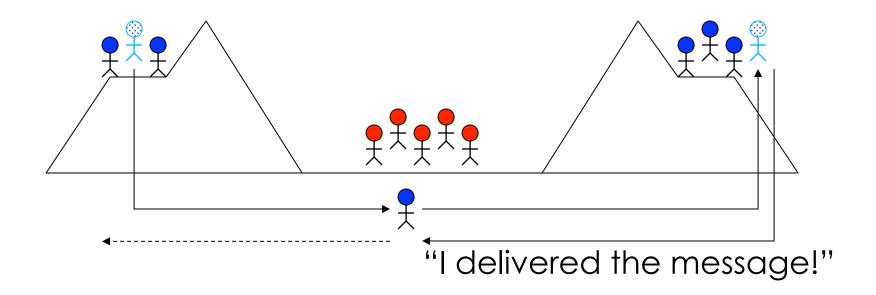
In the **Two Generals problem** two army divisions surround enemy, but must jointly attack to be successful.



### The divisions can only communicate via a messenger, who may not make it.



## To be sure the messenger made it, we await confirmation.



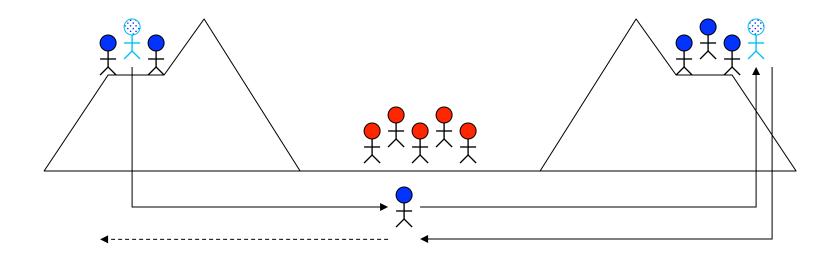
In the "two generals problem", can the two armies reliably coordinate their attack?

A. Yes!

**B**, No!

Be ready to explain why!

This problem shows us we can improve success rate, but can't make guarantees.



# Give up? No way! We're pretty good at engineering solutions!





### We have several tools for implementing reliable communication.

#### Our Concerns:

- Message corruption
- Message duplication
- Message loss
- Message reordering
- Performance

#### Our Toolbox:

- Checksums
- > Timeouts
- ACKs & NACKs
- Sequence numbering
- Pipelining