Carnegie Mellon University

12 Query Execution – Part I

Intro to Database Systems
15-445/15-645
Fall 2020

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ADMINISTRIVIA

Homework #3 is due Sun Oct 18th @ 11:59pm

Mid-Term Exam is Wed Oct 21st
→ Morning Session: 9:00am ET
→ Afternoon Session: 3:20pm ET

Project #2 is due Sun Oct 25th @ 11:59pm

PROJECTS

Write your own tests.

Practice defensive programming.

Profile your code to find performance problems.

Do not use Gradescope for debugging.

Do not directly email TAs for help.

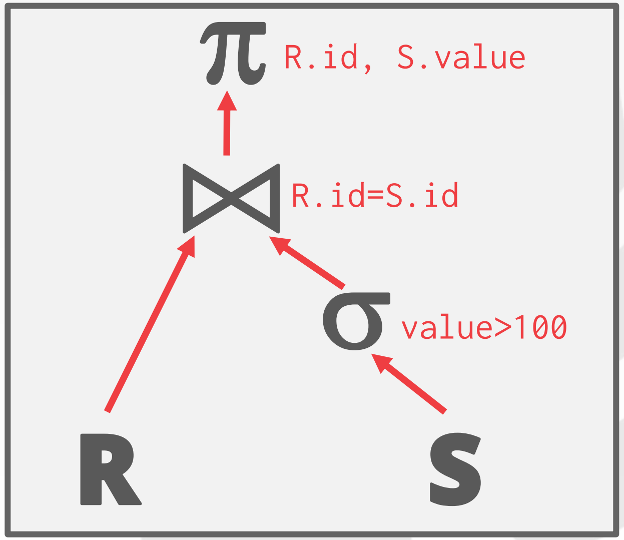
QUERY PLAN

The operators are arranged in a tree.

Data flows from the leaves of the tree up towards the root.

The output of the root node is the result of the query.

```
SELECT R.id, S.cdate
  FROM R JOIN S
    ON R.id = S.id
 WHERE S.value > 100
```



TODAY'S AGENDA

Processing Models

Access Methods

Modification Queries

Expression Evaluation

PROCESSING MODEL

A DBMS's **processing model** defines how the system executes a query plan.

→ Different trade-offs for different workloads.

Approach #1: Iterator Model

Approach #2: Materialization Model

Approach #3: Vectorized / Batch Model

ITERATOR MODEL

Each query plan operator implements a **Next** function.

→ On each invocation, the operator returns either a single tuple or a null marker if there are no more tuples.

→ The operator implements a loop that calls next on its children to retrieve their tuples and then process them.

Also called **Volcano** or **Pipeline** Model.

ITERATOR MODEL

Next()

```
for t in child.Next():
  emit(projection(t))
```

Next()

```
for t₁ in left.Next():
  buildHashTable(t₁)
for t₂ in right.Next():
  if probe(t₂): emit(t₁⋈t₂)
```

Next()

```
for t in child.Next():
  if evalPred(t): emit(t)
```

Next()

```
for t in R:
  emit(t)
```

Next()

```
for t in S:
  emit(t)
```

```
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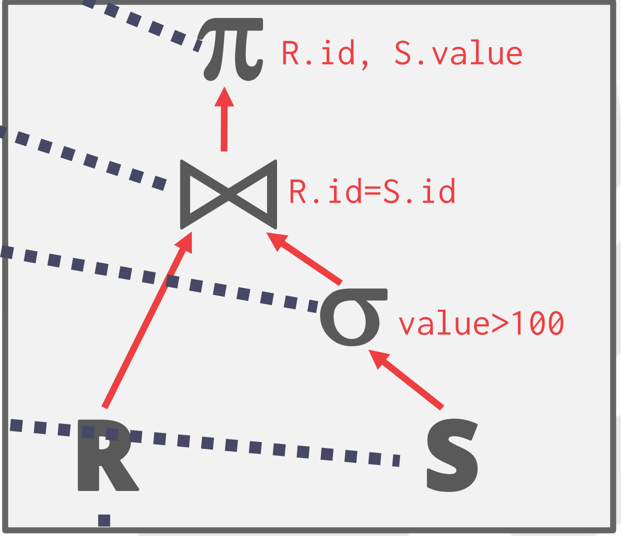
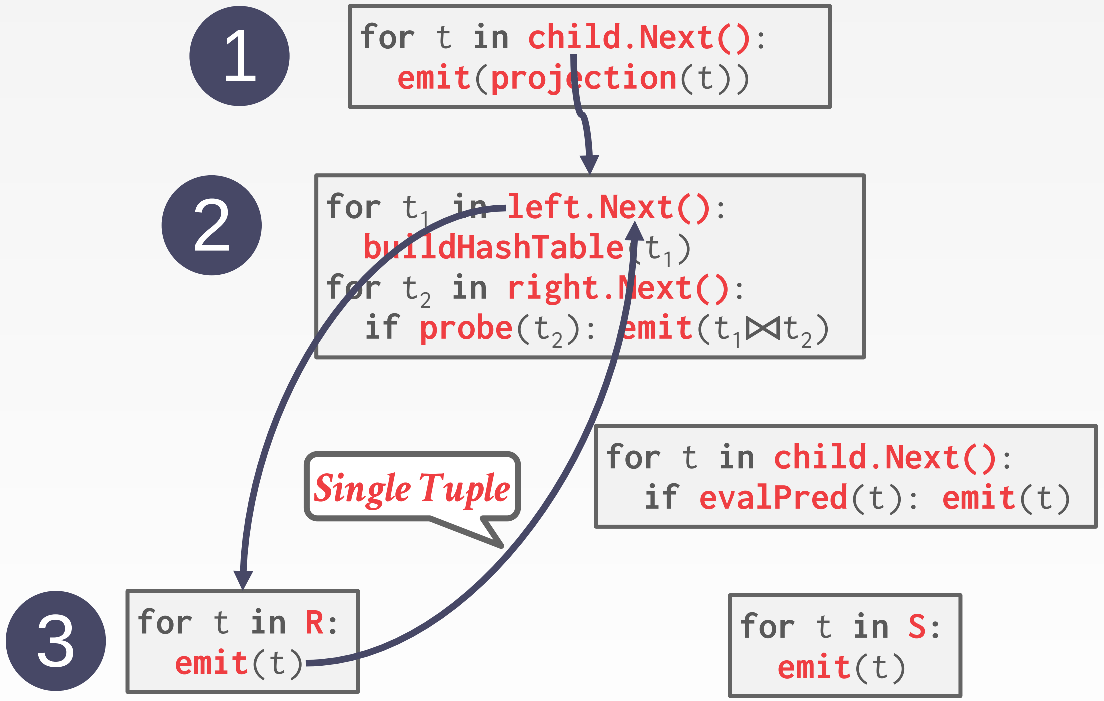


ITERATOR MODEL



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```
for t in child.Next():
  if evalPred(t): emit(t)
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3
```
for t in R:
  emit(t)
```

```
for t in S:
  emit(t)
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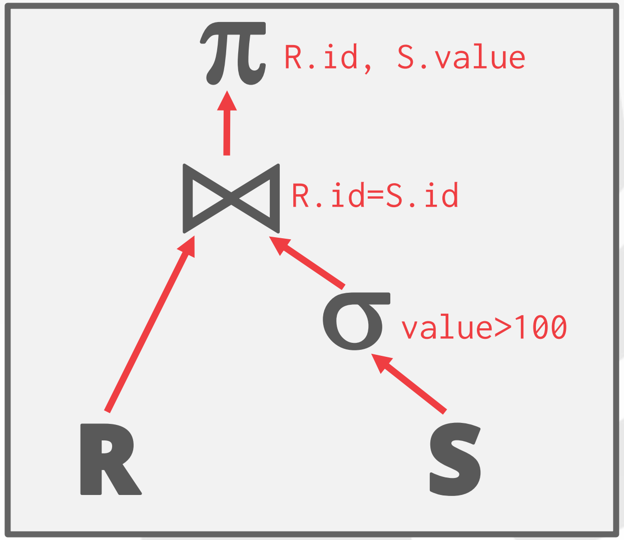
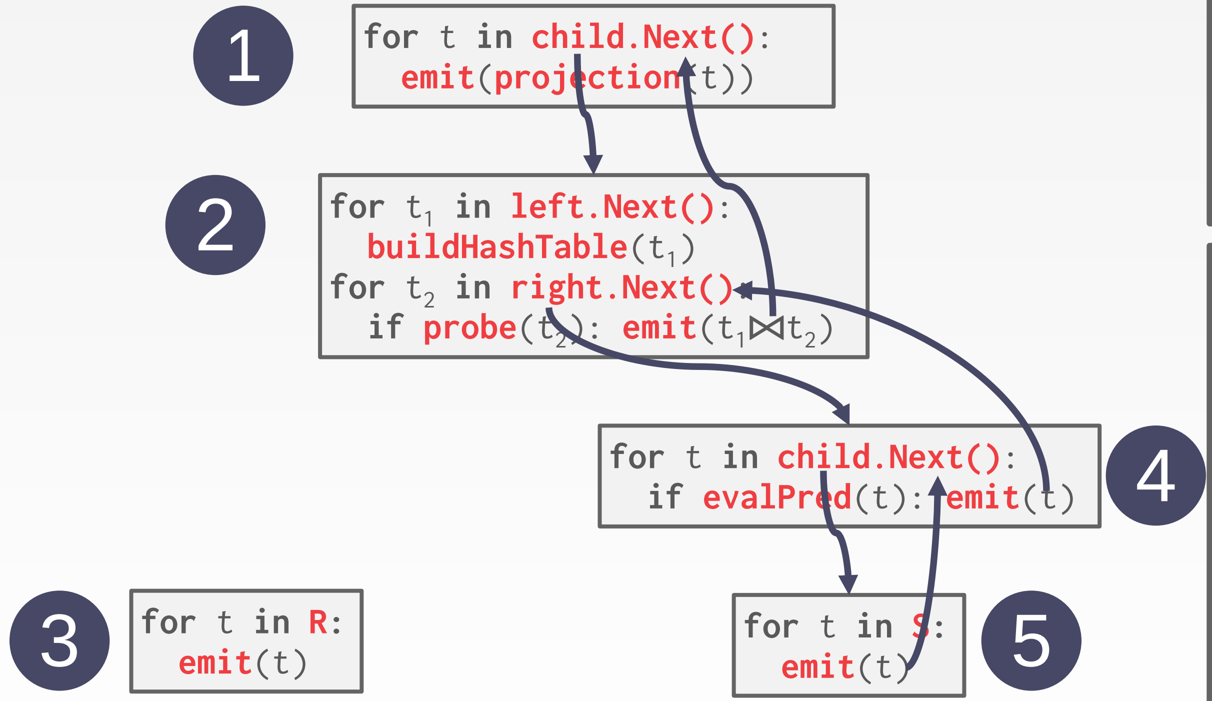


ITERATOR MODEL



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  if probe(t2): emit(t1⋈t2)
```

3
```
for t in R:
  emit(t)
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4
```
for t in child.Next():
  if evalPred(t): emit(t)
```

5
```
for t in S:
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```

```
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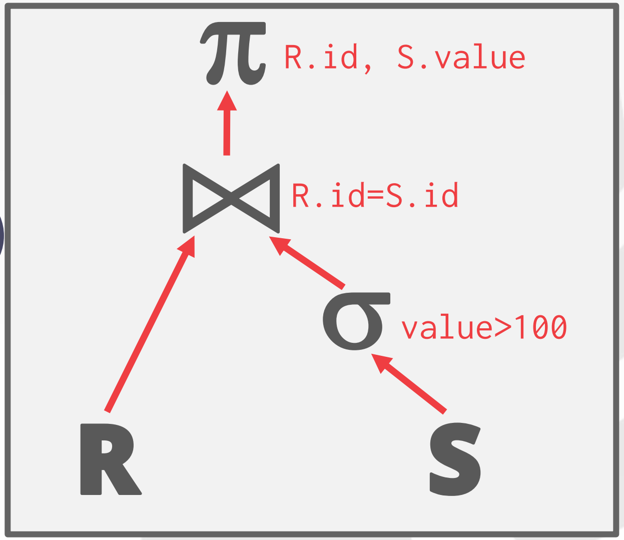

ITERATOR MODEL

This is used in almost every DBMS. Allows for tuple pipelining.

Some operators must block until their children emit all their tuples.
→ Joins, Subqueries, Order By

Output control works easily with this approach.

MATERIALIZATION MODEL

Each operator processes its input all at once and then emits its output all at once.

- → The operator "materializes" its output as a single result.
- → The DBMS can push down hints into to avoid scanning too many tuples.
- → Can send either a materialized row or a single column.

The output can be either whole tuples (NSM) or subsets of columns (DSM)

MATERIALIZATION MODEL

```
out = [ ]
for t in child.Output():
  out.add(projection(t))
return out
```

```
out = [ ]
for t₁ in left.Output():
  buildHashTable(t₁)
for t₂ in right.Output():
  if probe(t₂): out.add(t₁⋈t₂)
return out
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```
out = [ ]
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out = [ ]
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```

```
out = [ ]
for t in S:
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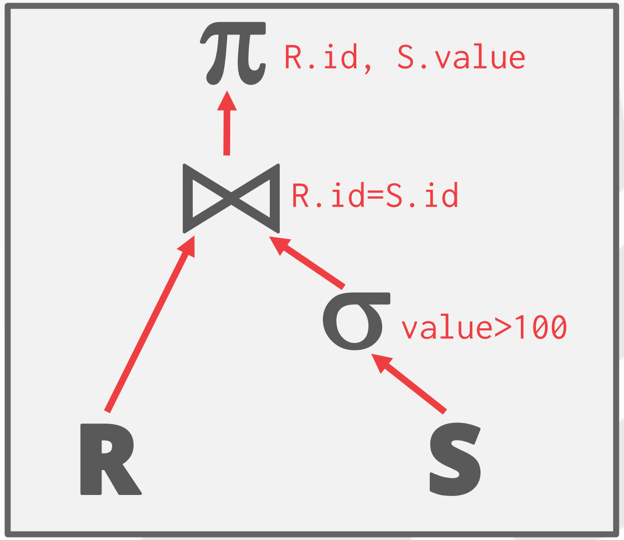

MATERIALIZATION MODEL

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out = [ ]
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  out.add(projection(t))
return out
```

2

```
out = [ ]
for t1 in left.Output():
  buildHashTable(t1)
for t2 in right.Output():
  if probe(t2): out.add(t1⋈t2)
return out
```

```
out = [ ]
for t in child.Output():
  if evalPred(t): out.add(t)
return out
```

All Tuples

3

```
out = [ ]
for t in R:
  out.add(t)
return out
```

```
out = [ ]
for t in S:
  out.add(t)
return out
```

```
SELECT R.id, S.cdate
  FROM R JOIN S
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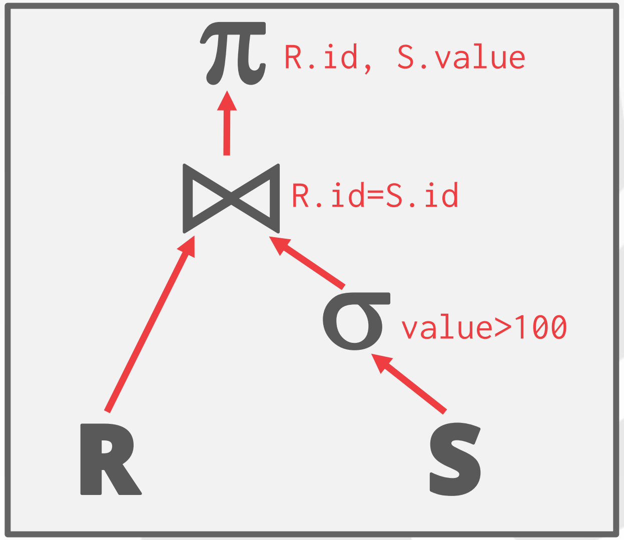
 WHERE S.value > 100
```



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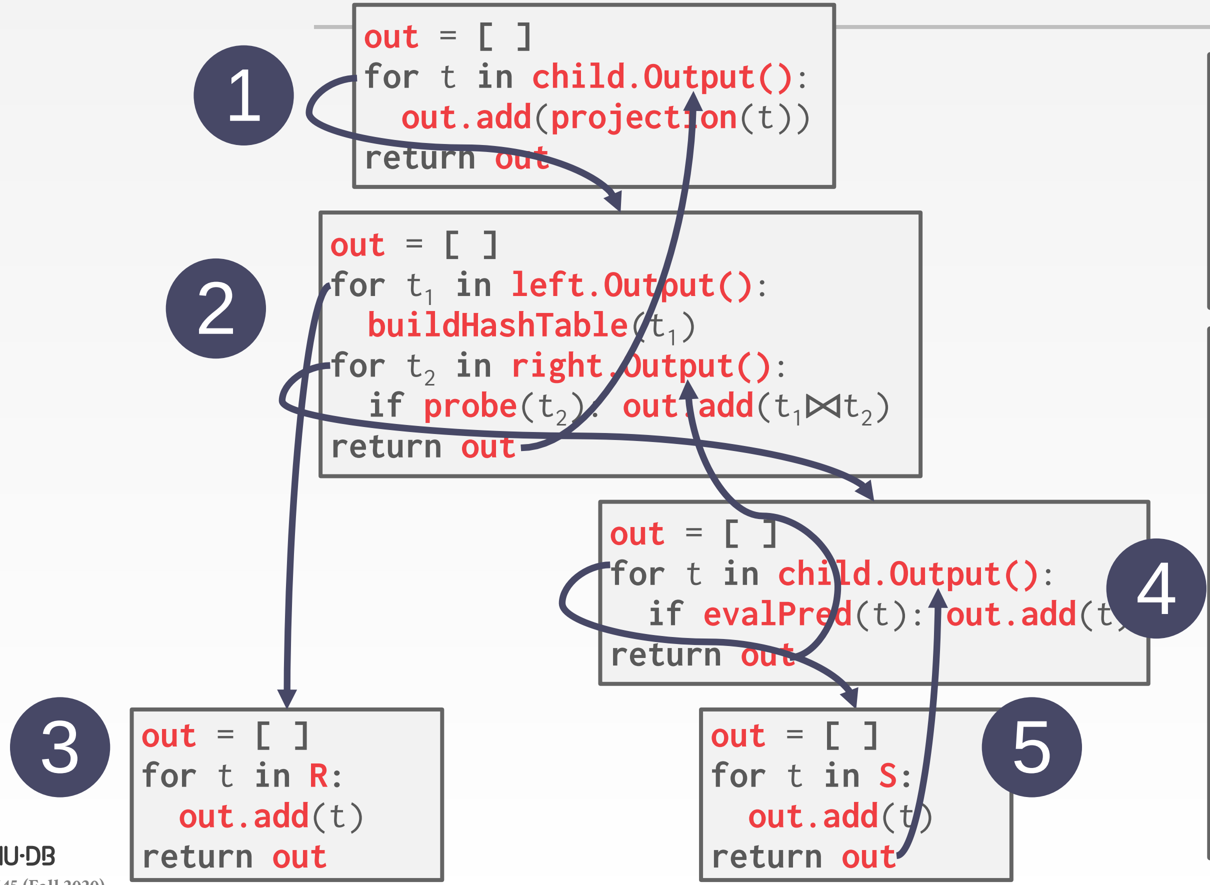

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return out
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out = [ ]
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```

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out = [ ]
for t in S:
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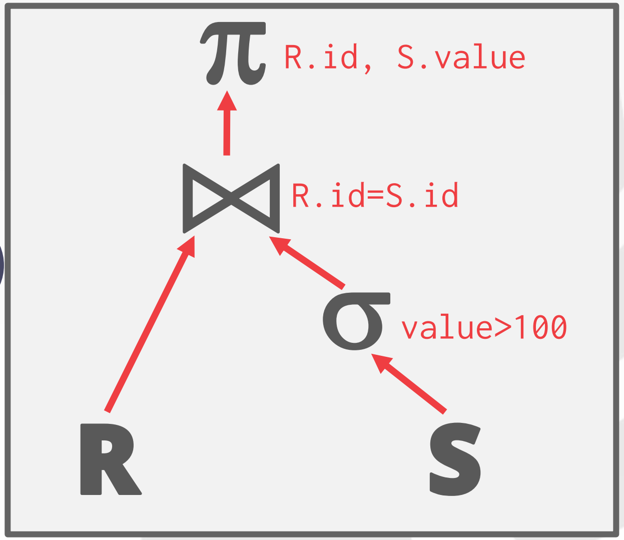

MATERIALIZATION MODEL

Better for OLTP workloads because queries only access a small number of tuples at a time.
→ Lower execution / coordination overhead.
→ Fewer function calls.

Not good for OLAP queries with large intermediate results.

VECTORIZATION MODEL

Like the Iterator Model where each operator implements a **Next** function in this model.

Each operator emits a **batch** of tuples instead of a single tuple.

→ The operator's internal loop processes multiple tuples at a time.

→ The size of the batch can vary based on hardware or query properties.

VECTORIZATION MODEL

```
out = [ ]
for t in child.Next():
  out.add(projection(t))
  if |out|>n: emit(out)
```

```
out = [ ]
for t1 in left.Next():
  buildHashTable(t1)
for t2 in right.Next():
  if probe(t2): out.add(t1⋈t2)
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out = [ ]
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for t in R:
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out = [ ]
for t in S:
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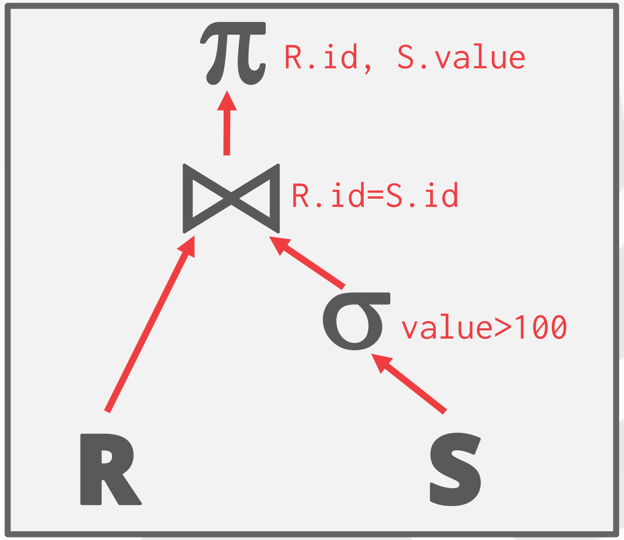
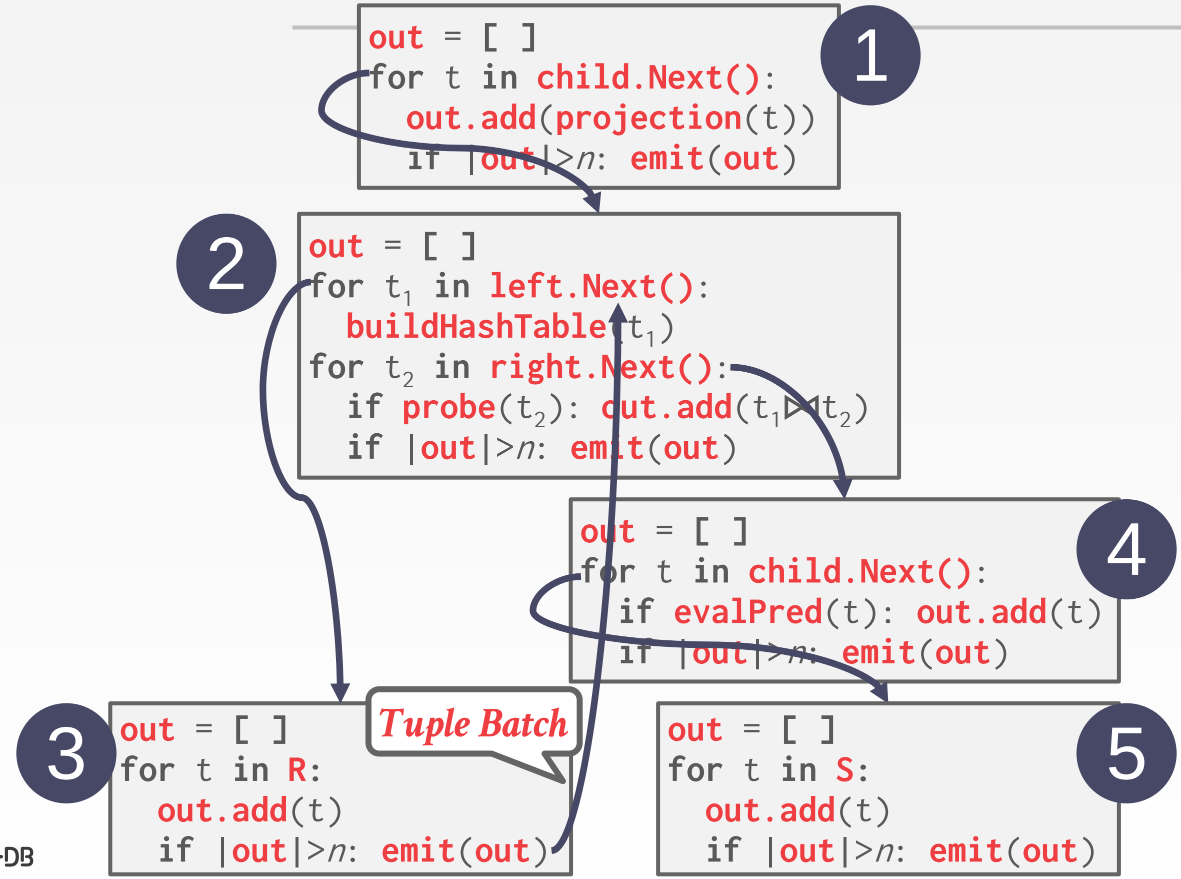


VECTORIZATION MODEL



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```

3
```
out = [ ]
for t in R:
  out.add(t)
  if |out|>n: emit(out)
```

Tuple Batch

4
```
out = [ ]
for t in child.Next():
  if evalPred(t): out.add(t)
  if |out|>n: emit(out)
```

5
```
out = [ ]
for t in S:
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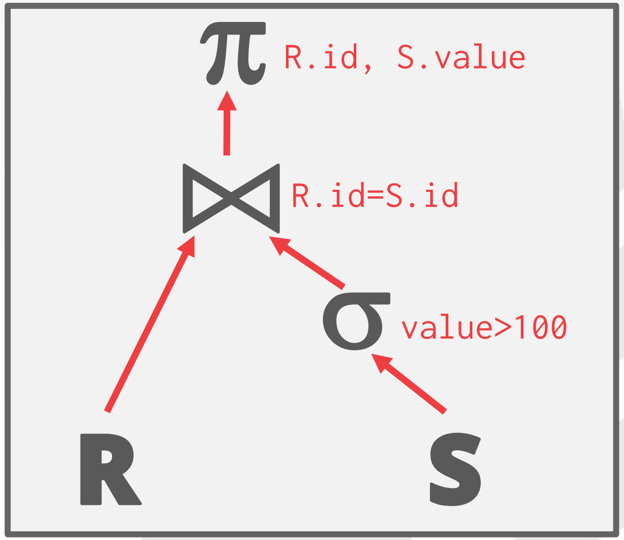

VECTORIZATION MODEL

Ideal for OLAP queries because it greatly reduces the number of invocations per operator.

Allows for operators to use vectorized (SIMD) instructions to process batches of tuples.

PLAN PROCESSING DIRECTION

Approach #1: Top-to-Bottom
→ Start with the root and "pull" data up from its children.
→ Tuples are always passed with function calls.

Approach #2: Bottom-to-Top
→ Start with leaf nodes and push data to their parents.
→ Allows for tighter control of caches/registers in pipelines.

ACCESS METHODS

An <u>**access method**</u> is a way that the DBMS can access the data stored in a table.
→ Not defined in relational algebra.

Three basic approaches:
→ Sequential Scan
→ Index Scan
→ Multi-Index / "Bitmap" Scan

```
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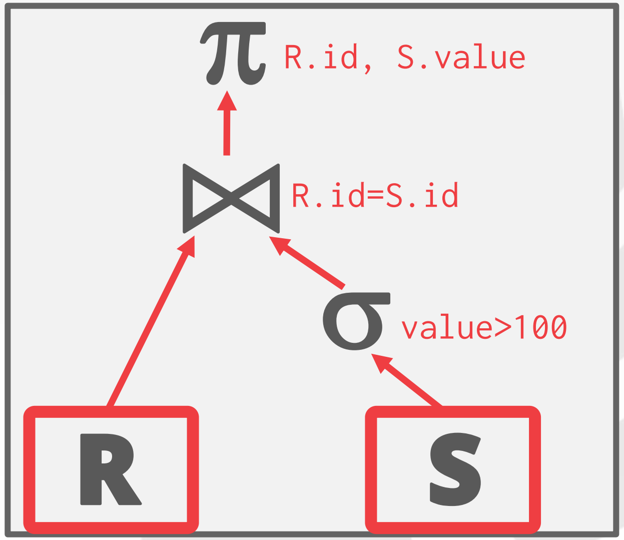


SEQUENTIAL SCAN

For each page in the table:
→ Retrieve it from the buffer pool.
→ Iterate over each tuple and check whether to include it.

```
for page in table.pages:
  for t in page.tuples:
    if evalPred(t):
      // Do Something!
```

The DBMS maintains an internal **cursor** that tracks the last page / slot it examined.

SEQUENTIAL SCAN: OPTIMIZATIONS

This is almost always the worst thing that the DBMS can do to execute a query.

Sequential Scan Optimizations:
→ Prefetching
→ Buffer Pool Bypass
→ Parallelization
→ Heap Clustering
→ Zone Maps
→ Late Materialization

ZONE MAPS

Pre-computed aggregates for the attribute values in a page. DBMS checks the zone map first to decide whether it wants to access the page.

```
SELECT * FROM table
 WHERE val > 600
```

Original Data

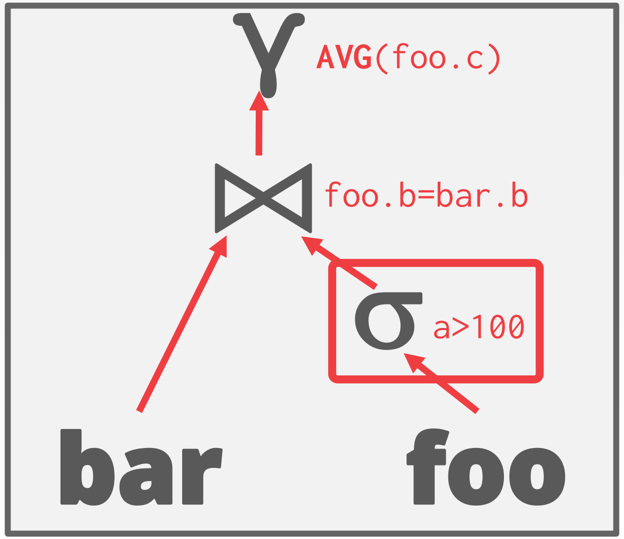
val
100
200
300
400
400

Zone Map

type	val
MIN	100
MAX	400
AVG	280
SUM	1400
COUNT	5

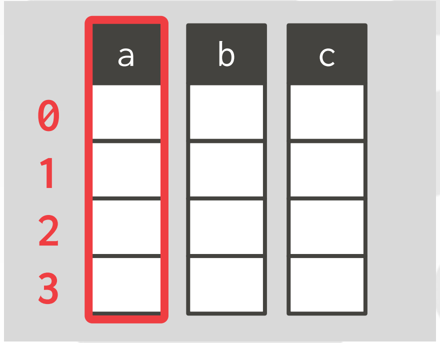
LATE MATERIALIZATION

DSM DBMSs can delay stitching together tuples until the upper parts of the query plan.



Offsets

```
SELECT AVG(foo.c)
  FROM foo JOIN bar
    ON foo.b = bar.b
 WHERE foo.a > 100
```



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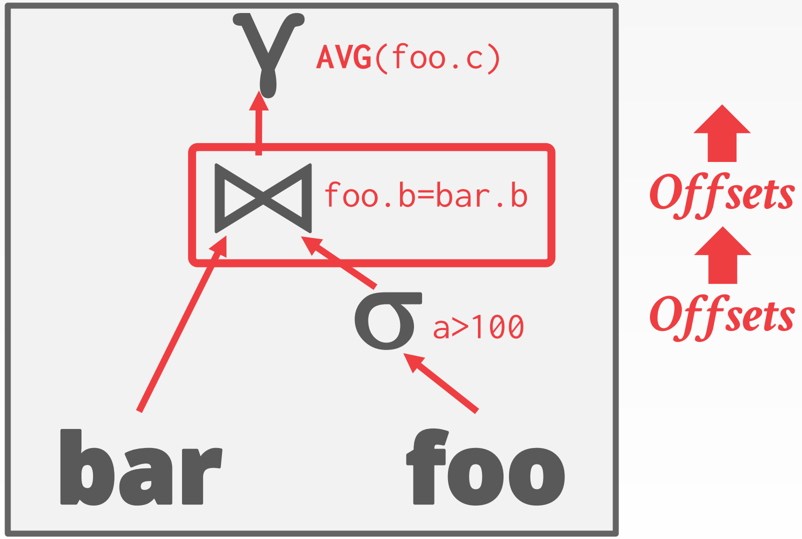

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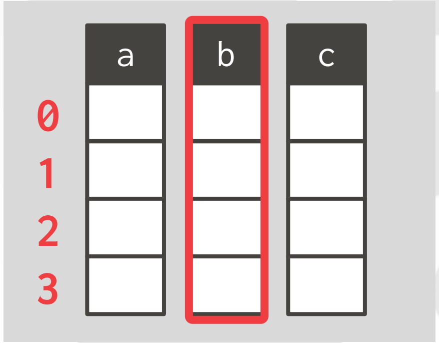

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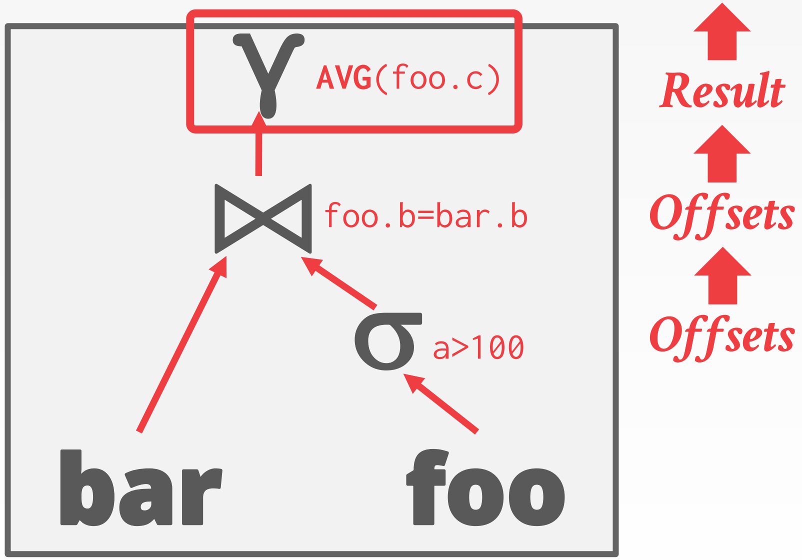

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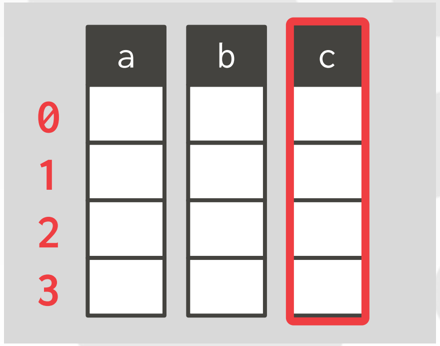


INDEX SCAN

The DBMS picks an index to find the tuples that the query needs.

Lecture 14

Which index to use depends on:
→ What attributes the index contains
→ What attributes the query references
→ The attribute's value domains
→ Predicate composition
→ Whether the index has unique or non-unique keys

INDEX SCAN

Suppose that we a single table with 100 tuples and two indexes:

→ Index #1: **age**
→ Index #2: **dept**

```
SELECT * FROM students
 WHERE age < 30
   AND dept = 'CS'
   AND country = 'US'
```

Scenario #1

There are 99 people under the age of 30 but only 2 people in the CS department.

Scenario #2

There are 99 people in the CS department but only 2 people under the age of 30.

MULTI-INDEX SCAN

If there are multiple indexes that the DBMS can use for a query:

→ Compute sets of record ids using each matching index.
→ Combine these sets based on the query's predicates (union vs. intersect).
→ Retrieve the records and apply any remaining predicates.

Postgres calls this Bitmap Scan.

MULTI-INDEX SCAN

With an index on **age** and an index on **dept**,

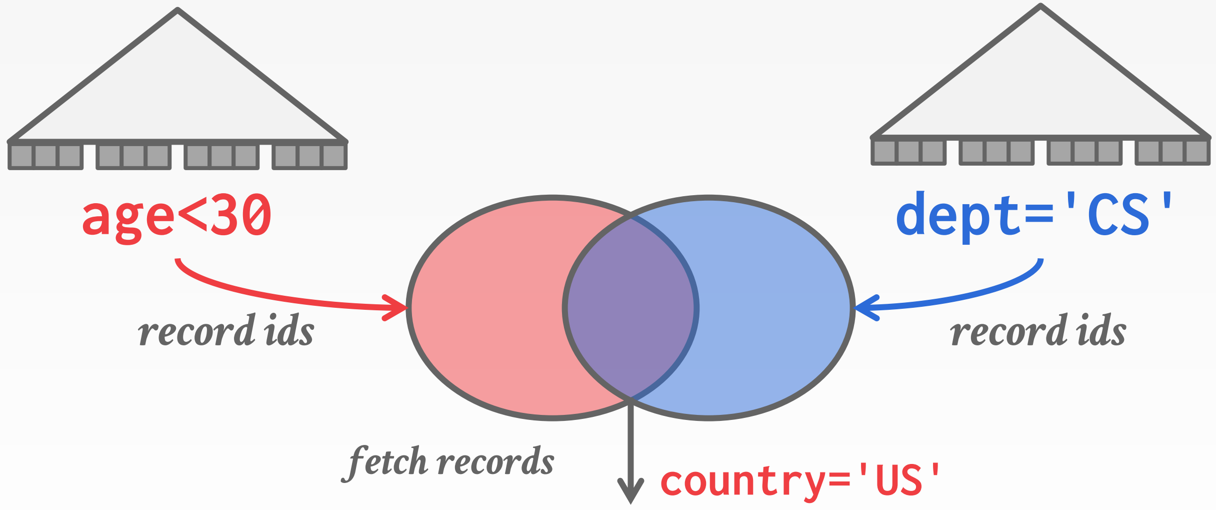
- → We can retrieve the record ids satisfying **age<30** using the first,
- → Then retrieve the record ids satisfying **dept='CS'** using the second,
- → Take their intersection
- → Retrieve records and check **country='US'**.

```
SELECT * FROM students
 WHERE age < 30
   AND dept = 'CS'
   AND country = 'US'
```

MULTI-INDEX SCAN

Set intersection can be done with bitmaps, hash tables, or Bloom filters.

```
SELECT * FROM students
 WHERE age < 30
   AND dept = 'CS'
   AND country = 'US'
```



MODIFICATION QUERIES

Operators that modify the database (**INSERT**, **UPDATE**, **DELETE**) are responsible for checking constraints and updating indexes.

UPDATE/**DELETE**:
→ Child operators pass Record Ids for target tuples.
→ Must keep track of previously seen tuples.

INSERT:
→ **Choice #1**: Materialize tuples inside of the operator.
→ **Choice #2**: Operator inserts any tuple passed in from child operators.

UPDATE QUERY PROBLEM

```
for t in child.Next():
  removeFromIndex(idx_salary, t.salary, t)
  updateTuple(t.salary = t.salary + 1000)
  insertIntoIndex(idx_salary, t.salary, t)
```

```
for t in people:
  emit(t)
```

```
CREATE INDEX idx_salary
    ON people (salary);
```

```
UPDATE people
   SET salary = salary + 100
 WHERE salary < 1000
```

Index(people.salary)

UPDATE QUERY PROBLEM

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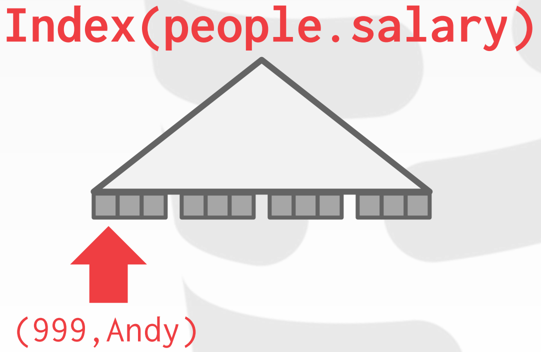

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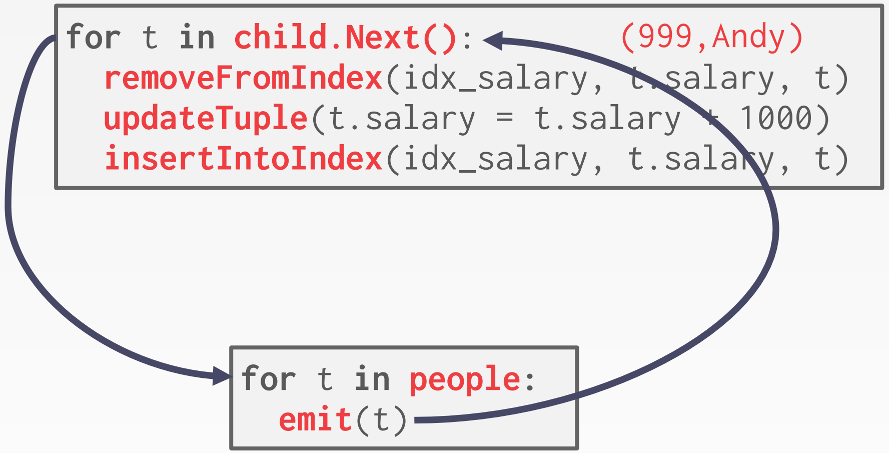

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```

(999,Andy)

```
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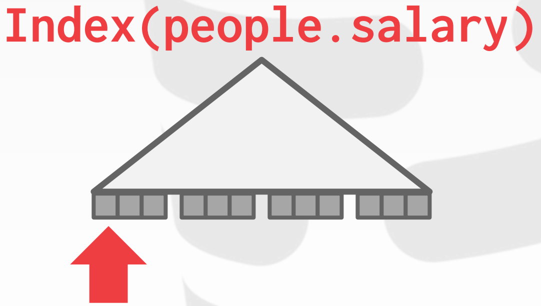

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for t in people:
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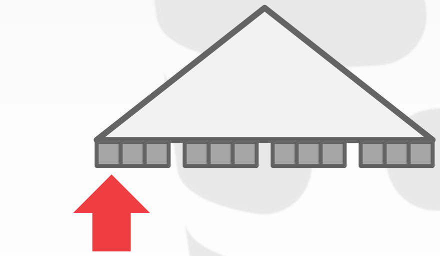
```

Index(people.salary)



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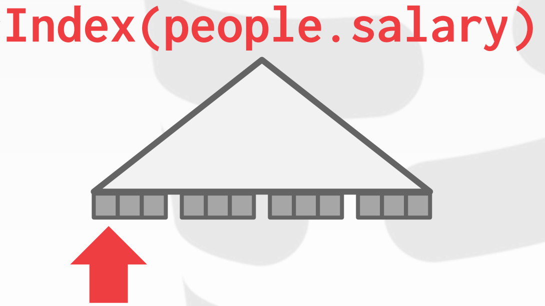

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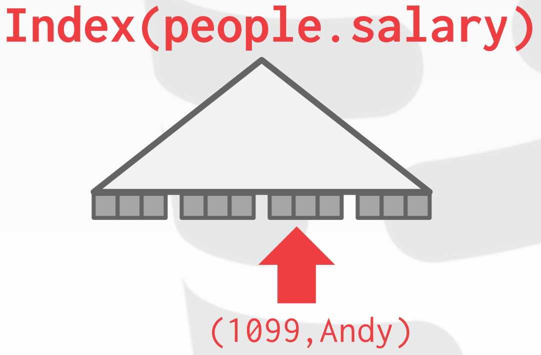

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CREATE INDEX idx_salary
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```

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UPDATE people
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```
for t in child.Next():          (1099,Andy)
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  insertIntoIndex(idx_salary, t.salary, t)
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for t in people:
  emit(t)
```
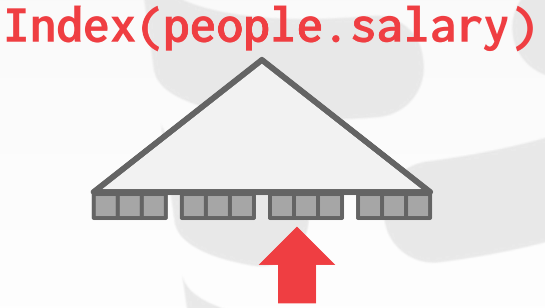



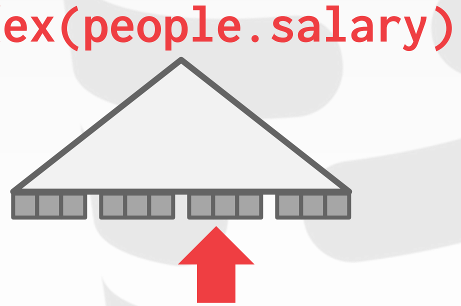
UPDATE QUERY PROBLEM

```
for t in child.Next():          (1199,Andy)
  removeFromIndex(idx_salary, t.salary, t)
  updateTuple(t.salary = t.salary + 1000)
  insertIntoIndex(idx_salary, t.salary, t)
```

```
for t in people:
  emit(t)
```

```
CREATE INDEX idx_salary
    ON people (salary);
```

```
UPDATE people
   SET salary = salary + 100
 WHERE salary < 1000
```



HALLOWEEN PROBLEM

Anomaly where an update operation changes the physical location of a tuple, which causes a scan operator to visit the tuple multiple times.
→ Can occur on clustered tables or index scans.

First discovered by IBM researchers while working on System R on Halloween day in 1976.

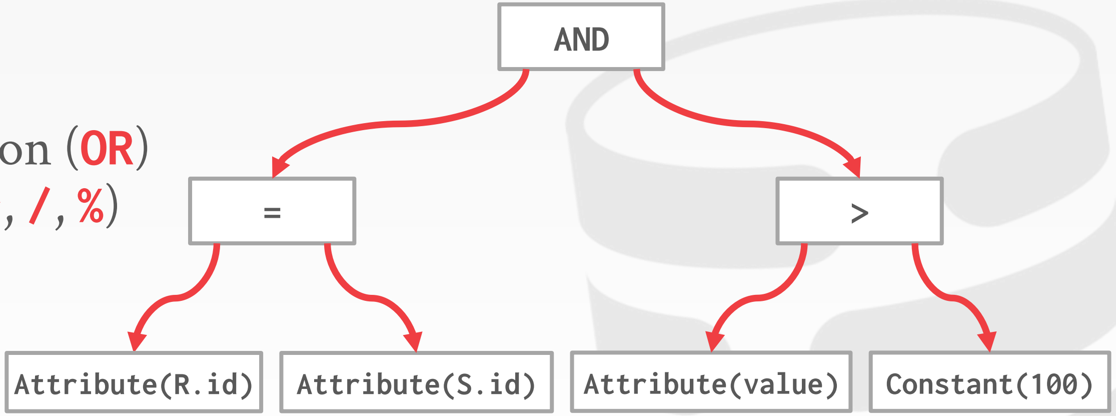
EXPRESSION EVALUATION

The DBMS represents a **WHERE** clause as an **expression tree**.

The nodes in the tree represent different expression types:

→ Comparisons (**=**, **<**, **>**, **!=**)
→ Conjunction (**AND**), Disjunction (**OR**)
→ Arithmetic Operators (**+**, **-**, *****, **/**, **%**)
→ Constant Values
→ Tuple Attribute References

```
SELECT R.id, S.cdate
  FROM R JOIN S
    ON R.id = S.id
 WHERE S.value > 100
```

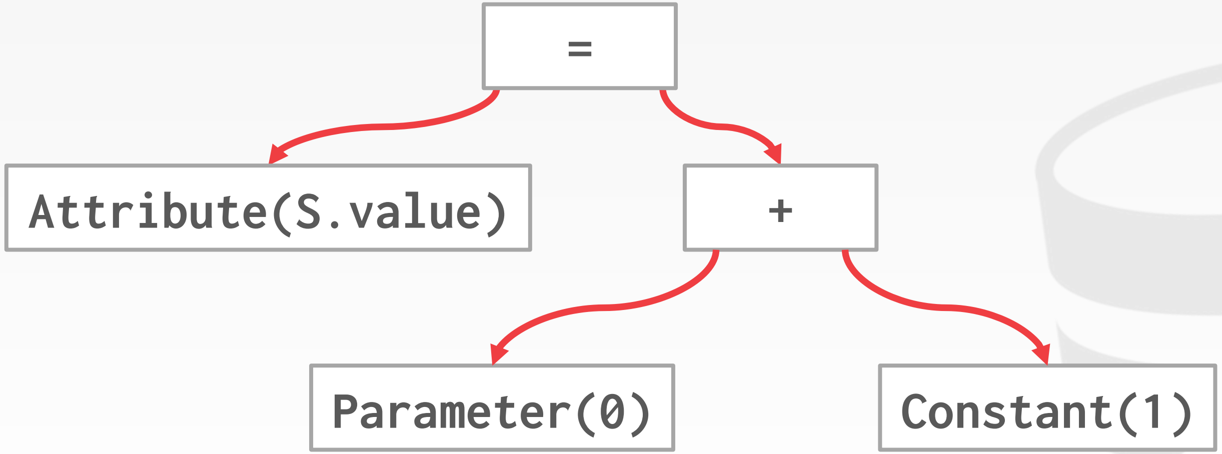


EXPRESSION EVALUATION

```
SELECT * FROM S
 WHERE B.value = ? + 1
```

Execution Context

```
Current Tuple        Query Parameters    Table Schema
(123, 1000)          (int:999)           S→(int:id, int:value)
```

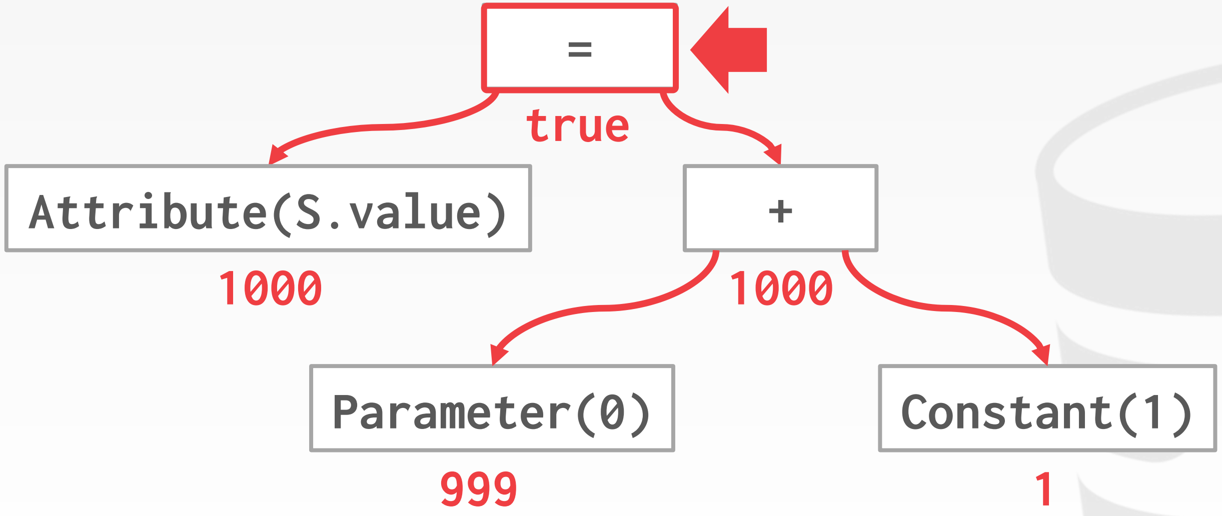


EXPRESSION EVALUATION

```
SELECT * FROM S
 WHERE B.value = ? + 1
```

Execution Context

Current Tuple	Query Parameters	Table Schema
(123, 1000)	(int:999)	S→(int:id, int:value)



EXPRESSION EVALUATION

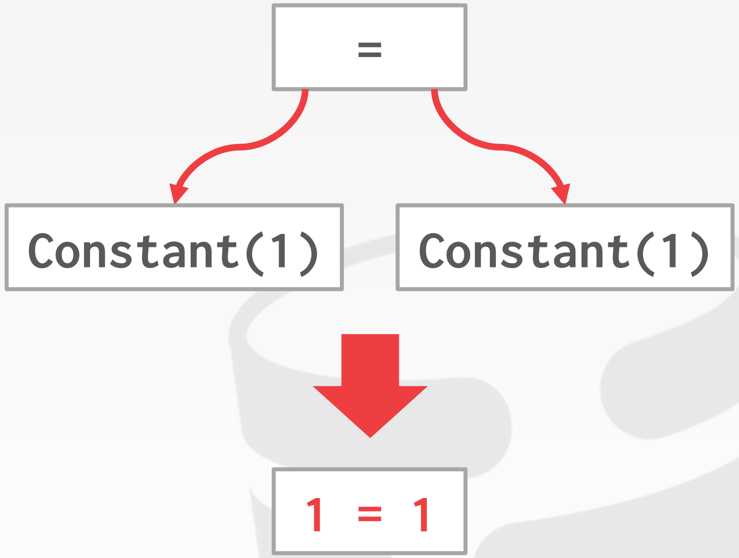
Evaluating predicates in this manner is slow.

→ The DBMS traverses the tree and for each node that it visits it must figure out what the operator needs to do.

Consider the predicate "`WHERE 1=1`"

A better approach is to just evaluate the expression directly.

→ Think JIT compilation



CONCLUSION

The same query plan be executed in multiple ways.

(Most) DBMSs will want to use an index scan as much as possible.

Expression trees are flexible but slow.

NEXT CLASS

Parallel Query Execution