

Access control

What is access control

- Security functionality ensuring that authorized users can do what they are permitted to do
 - Hardware
 - Operating systems
 - Web servers
 - Databases
- Preserve confidentiality, integrity, availability
- OS provides access to resources
 - CPU, memory, files, devices, network
- Needs access control to
 - Protect OS from applications
 - Protect applications from each other

Subjects and Objects

- Subject: active entity needing to access resources
- Object: passive entity accessed by subjects
 - Files, devices
- Subjects access objects: they perform actions on objects



- Access control
 - Define what operations subjects can perform on objects
 - Permissions

Example

Subjects :

Alice and Bob: doctors

Tom and Paul: nurses

Sarah: secretary

Objects :

patients

documents: medical or administrative

Actions :

consult, modify the medical doc of a patient

create a doc

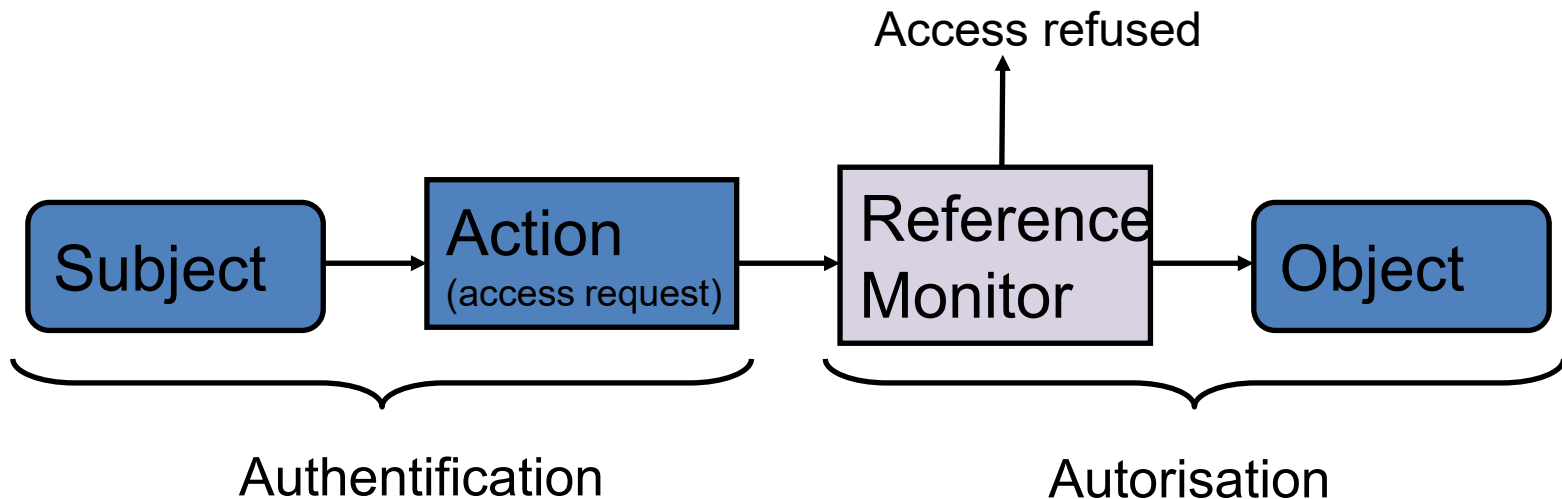
We can specify :

what is permitted for an object

or what a subject can do

Access Control

- Reference Monitor
 - Decision process filtering the access demands
 - Guard controlling whether a subject can access an object



Principle of Least Privilege

- Each subject can access only the resources necessary to perform its task
 - The nurses cannot modify docs of patients
 - They can read them, or add information

Access control classification

- Discretionary Access Control (DAC)
- Mandatory Access Control (MAC)
- Role-Based Access Control (RBAC)
- Task-Based Access Control (TBAC)
- ...

Discretionary Access Control (DAC)

- Owner determines access rights
 - Typically identity-based access control
 - Owner specifies other users who have access
- A subject can pass information onto anyother subject
- In some cases, access rights may be transferred
- Users are in charge of access permissions
 - Most systems use this
- Unix file system

Discretionary Access Control (DAC)

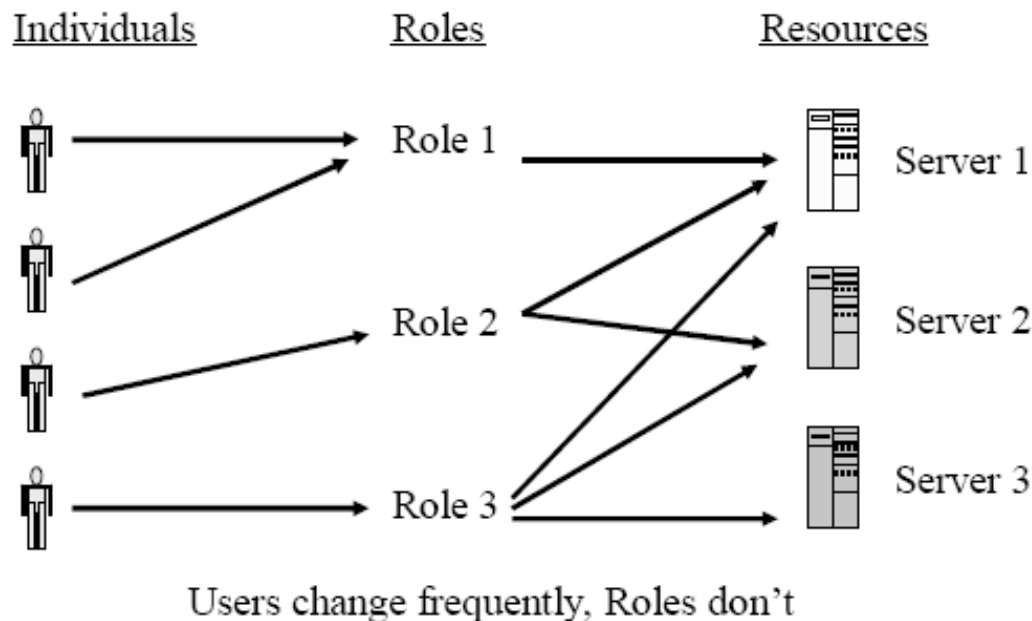
- Suppose we can find the owner of each object
- Violates the principle of least privilege
- Access permission may be transferred without informing owner
 - E.g., A grants read permission to B on f, B copies f to f', B is owner of f', and transmits read permission to C
 - A is not informed
- DAC can limit access for good users
- Not adapted in sensible systems

Mandatory Access Control (MAC)

- Motivation
 - Control information flow even in sensible environment
- Non discretionary
 - Permissions fixed for each subject
- Rules specify granting of access
 - Also called rule-based access control
- Administrators are in charge of access permissions

Role-Based Access Control (RBAC)

- Access decisions depend on roles
 - Administrators define roles for various job functions
 - Each role contains permissions to perform certain operations
 - Users are assigned one or more roles
- Allow enforcement of both MAC & DAC



Access Control Matrix

- Primary abstraction for protection in computer security
- Rows: domains (subjects or groups of subjects)
- Columns: objects
- Each entry in the matrix represents an access right of a domain on an object

	bill.doc	edit.exe	fun.com
<i>Alice</i>	\emptyset	{execute}	{execute, read}
<i>Bob</i>	{read, write}	{execute}	{execute, read, write}

Access Control Matrix: Examples

- Modelling a programming language
 - S: procedures/modules
 - O: procedures/variables
 - A: execution of functions
- Modelling a local area network
 - S: stations
 - O: stations
 - A: protocols

	Compt	Inc_ctr	Dcr_ctr
Inc_ctr	+		
Dcr_ctr	-		
Manager		call	call

hosts	Station1	Station2	Station3
Station1	own	ftp	ftp, mail
Station2		own, ftp, nfs	ftp, nfs
Station3	mail	ftp	own, ftp, nfs

Bell LaPadula (BLP) Model

- Designed for the military: US airforce
- Based on U.S. military classification levels
- Ensures confidentiality
- Multi-level mandatory access control

Multilevel access control

- Entities are assigned security levels
- Totally ordered
 - TS>S>C>NC
 - R>Prop>Sens>Pub

Military
domain

Top Secret (TS)

Secret (S)

Confidentiel (C)

Non Classifié (NC)

Commercial
domaine

Restrained (R)

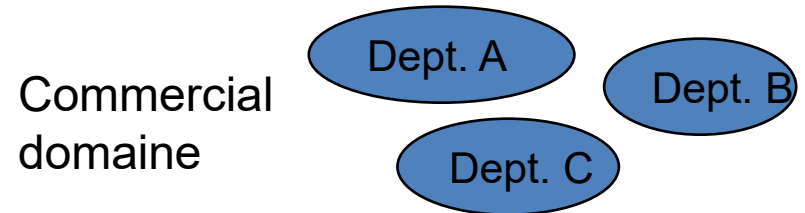
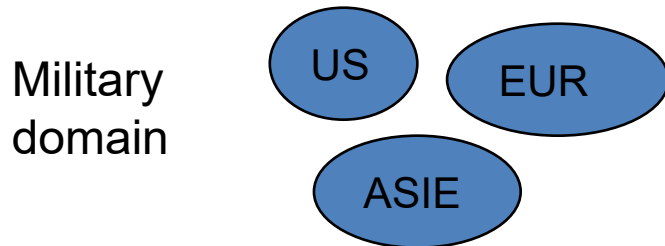
Proprietary (Prop)

Sensible (Sens)

Public (Pub)

Multilevel access control

- Entities are assigned categories
- Principle of need to know
- Partially ordered

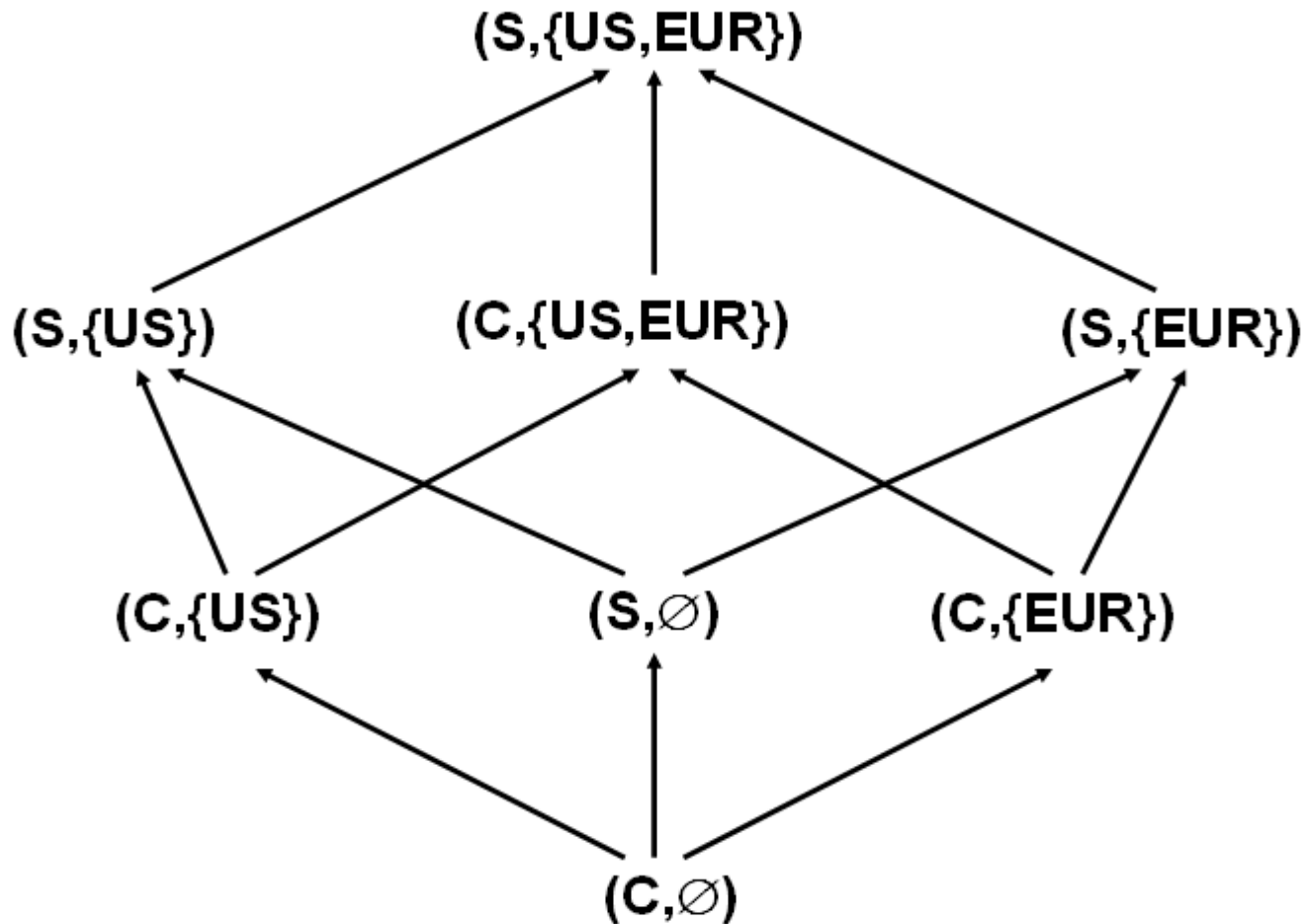


Security labels

- Label of security: $LS = Lev \times P(Cat)$
 - $(TS, \{US, EU\})$
- Partially ordered relationship: dom (dominate)
 - $(n1, C1) \text{ dom } (n2, C2) \text{ iff } n1 \geq n2 \text{ and } C2 \subseteq C1$
 - Example : $(TS, \{US, EUR\}) \text{ dom } (S, \{US\})$
- (LS, dom) is a lattice
 - admits an upper and lower-bound

Security labels

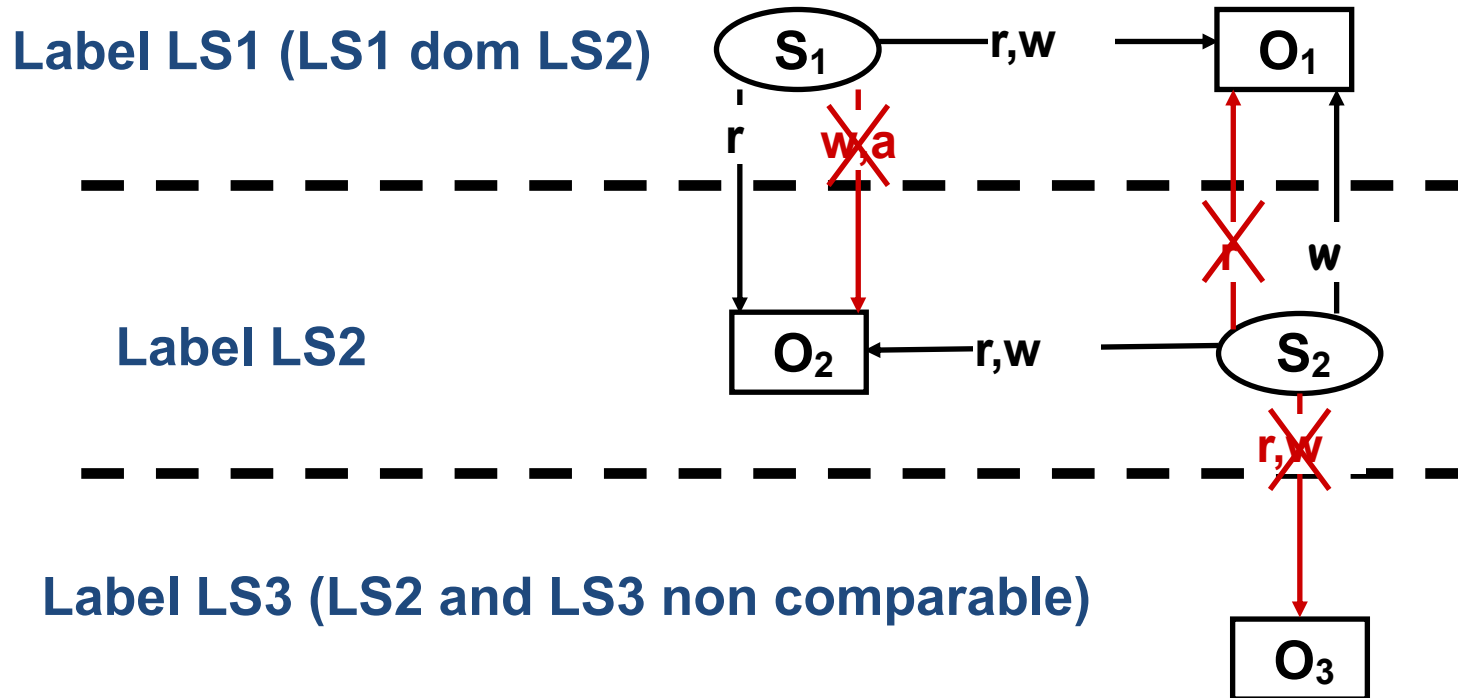
- Lattice induced by security labels: $\text{Lev}=\{C,S\}$, $\text{Cat}=\{US,EUR\}$



Bell LaPadula (BLP) Model

- LS(S): label of security of subject
- LS(O): label of security of object
- Simple Security Condition: **NO READ UP**
 - S can read O iff $LS(S) \text{ dom } LS(O)$
- *-Property (star-property, ou confinement): **NO WRITE DOWN**
 - S can write O iff $LS(O) \text{ dom } LS(S)$

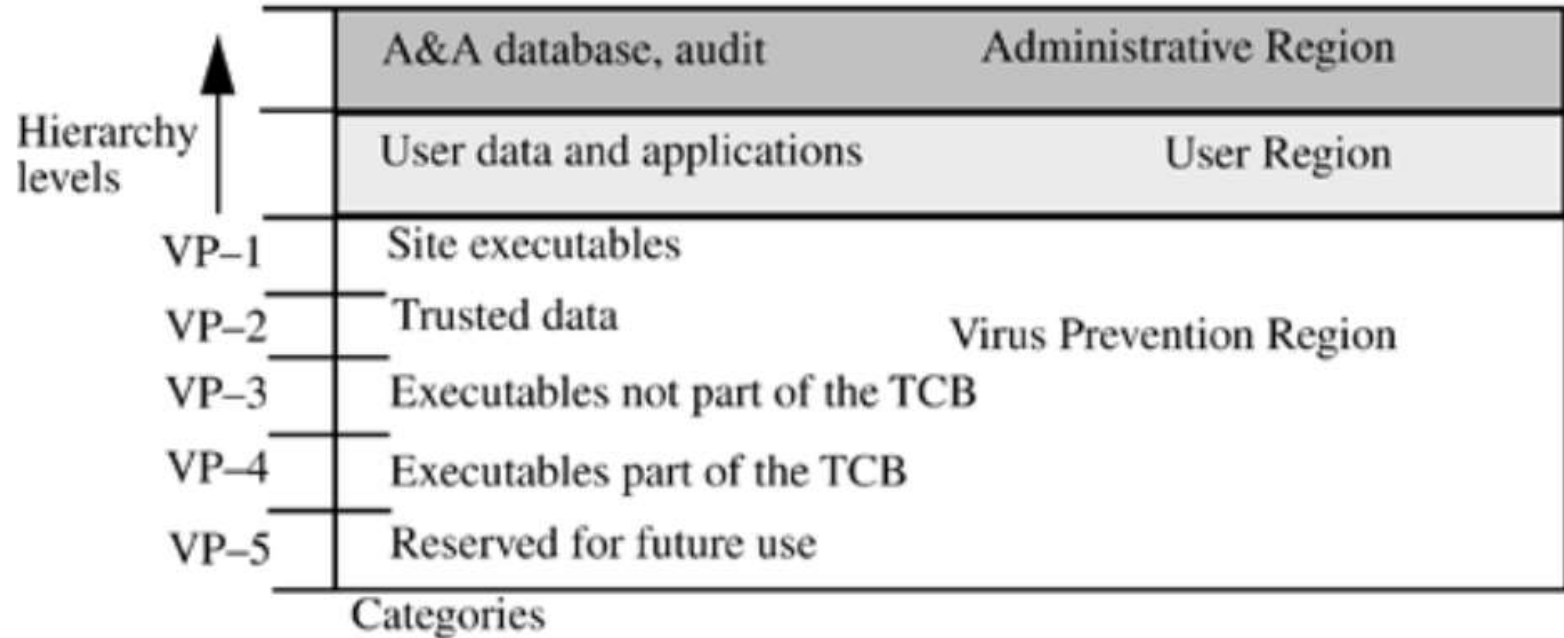
Bell LaPadula (BLP) Model



Test

- Consider the system of a hospital
 - S: doctors, nurses
 - O: prescriptions, documents
 - Doctors can read and write prescriptions and docs
 - Nurses can read and write prescriptions, but only add in docs
- Prove we cannot model the system using BLP model
- Consider an extended BLP model
 - Each S has two labels (L1,L2) with $L2 \text{ dom } L1$
 - Consider S with labels (L1,L2) and O with label L
 - S can read O iff $L2 \text{ dom } L$
 - S can write O iff $L \text{ dom } L1$
 - Prove that the extended BLP model can model the system
- Prove that we can create an information flow between two objects with non-comparable labels

Unix Data General B2 system



- S: users & processes; O: files, directories
 - Apply BLP model
- Why UR is between AR and VPR?
- Write-up is not allowed, why?
- S can create files in directory D iff $LS(S)=LS(D)$, why?
 - Problems with /tmp and /var/mail: multilevel directory