# Computer Systems

- Computer System Components
- Key Concepts
- Performance

## Types of Computing Device

Microprocessor: General purpose processor, mainly used in desktop computers and servers.

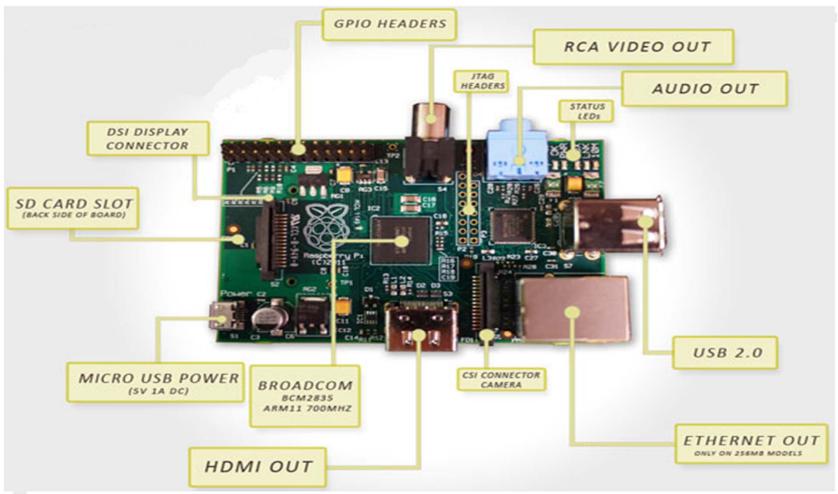
Microcontroller: Lots of peripheral support, includes functional parts to reduce the external parts count, real-time applications.

Digital Signal Processor (DSP): Processor optimized for Multiply Accumulate (MAC) operations. Processes data in real time.

EEE225/NJP

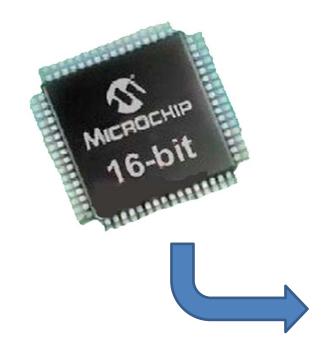
### ARM 11 – 32 bit RISC microprocessor

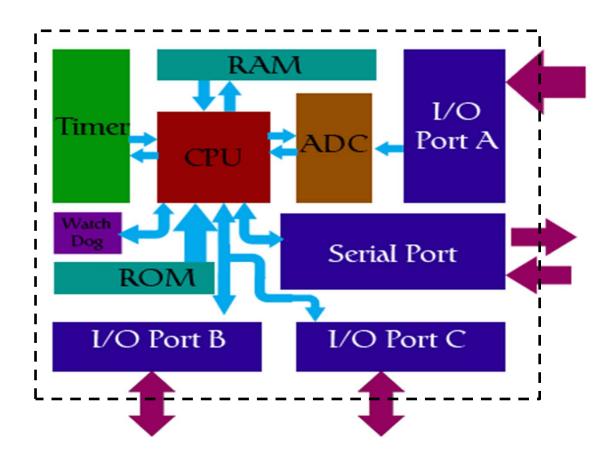
#### Raspberry Pi - an ARM GNU/Linux box for £29



EEE225/NJP

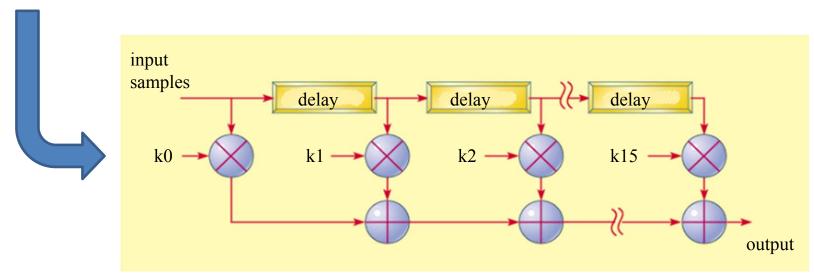
#### **PIC Microcontrollers**



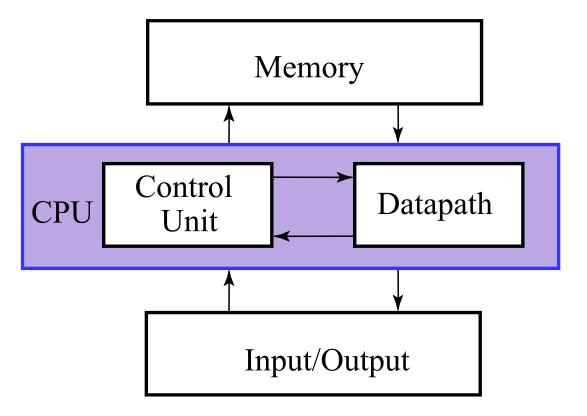


#### **DSP – Texas Instruments TMS320**





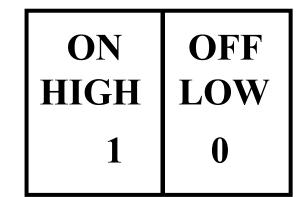
### Simple Model



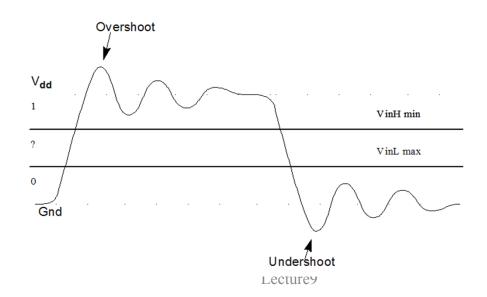
- A computer takes input and acts on it according to stored instructions (program)
- To store instructions, we need some memory
- To process instructions, we need a processor (Central Processing Unit or CPU)
- To transact with the computer we need input and output

### Representation of information

- Binary signal (two states)
- Achieved using transistors switches
- Low voltage = 0, High voltage = 1



 $\bullet$  Voltage level,  $\boldsymbol{V}_{DD}$  , depends upon the chosen technology



EEE225/NJP

#### **Combinations of bits**

A binary digit is called a 'bit' and can take the values 0 or 1

- 1 bit 2 alternatives **0,1**
- 2 bits 4 alternatives **00**, **01**, **10**, **11**
- 3 bits 8 alternatives
- 4 bits 16 alternatives
- 8 bits  $2^8 = 256$  alternatives
- n bits 2<sup>n</sup> alternatives

Information words are built up from binary digits to form a word of any desired length. The length depends on the application. A group of 8 bits is called a byte.

e.g. 4 bit word ABCD = 0000, 0001, 0010, 0011, ..., 1111

## Binary Numbers

The position of the bit in a binary number indicates its weight which is a power of two. A binary number can be converted to decimal by adding the weights of all bits that are 1.

$$1101_2 = (1 \times 2^3) + (1 \times 2^2) + (0 \times 2^1) + (1 \times 2^0)$$

$$= (1 \times 8) + (1 \times 4) + (0 \times 2) + (1 \times 1) = 8 + 4 + 1 = 13_{10}$$

The left-most bit is known as the most significant bit or MSB. The right-most bit is known as the least significant bit or LSB.

### Binary to Hexadecimal Conversion

Hex: there are 16 digits

0,1,2,3,4,5,6,7,8,9,A,B,C,D,E,F

each digit corresponds to four binary digits as  $2^4 = 16$ 

```
0000 0001 0010 0011
0100 0101 0110 0111
1000 1001 1010 1011
1100 1101 1110 1111
```

### Binary to Hexadecimal Conversion

To convert from binary to hex, group the binary number in fours and assign the corresponding hex digit to each group.

$$(110 \ 1010 \ 0101 \ 1111)_2 = (6A5F)_{16}$$
6 A 5 F

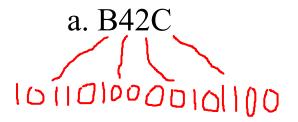
#### convert from binary to hex:

- a. 10010111001101
- b. 1010011111101001

#### convert from hex to binary:

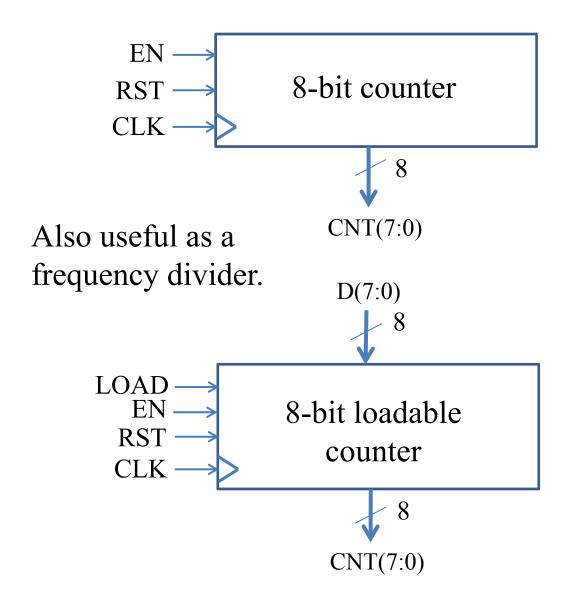
- a. B42C
- b. EFC

convert from binary to hex:









decimal	binary	hex
0	$0\ 0\ 0\ 0\ 0$	0
1	00001	1
2	00010	2
3	00011	3
4	00100	4
5	00101	5
6	0 0 1 1 0	6
7	0 0 1 1 1	7
8	01000	8
9	0 1 0 0 1	9
10	01010	A
11	0 1 0 1 1	В
12	0 1 1 0 0	C
13	0 1 1 0 1	D
14	0 1 1 1 0	E
15	0 1 1 1 1	F
16	10000	10
17	10001	11
18	10010	12

EEE225/NJP

Lecture9

#### **Non-Numeric Information**

Instructions and non-numeric data are a problem. So we represent both of these by numbers.

An instruction is a number which the processor interprets as a particular, unique instruction.

Non-numeric data such as text is handled similarly. For example ASCII (American Standard for Character Information Interchange) codes represent each character by a number:

$$'0' = 48$$
 $'A' = 65$ 
 $' = 32$ 









EEE225/NJP

Lecture9

### Some approximate figures for aircraft performance:

Aircraft	Passengers	Range (miles)	Speed (mph)	Passenger throughput (passengers x mph)
Boeing 777	375	4630	610	228,750
Boeing 747	470	4150	610	286,700
Concorde	132	4000	1350	178,200
DC-8-50	146	8720	544	79,424

Response Time: the response time or execution time is the total time for the computer to complete a task.

**Throughput**: the throughput or bandwidth is the number of tasks completed per unit time.

EEE225/NJP Lecture9

#### Performance in relation to Execution Time

For a computer X, simplifying execution time to CPU execution time:

Performance<sub>X</sub> = 
$$\frac{1}{\text{Execution time}_{X}}$$

For two computers X and Y, if the performance of X is greater than the performance of Y

So if we say that X is *n* times faster than Y

$$\frac{\text{Performance}_{X}}{\text{Performance}_{Y}} = \frac{\text{Execution time}_{Y}}{\text{Execution time}_{X}} = n$$

### Relative Performance Example

If computer X runs a program in 10 seconds and computer Y runs the same program in 15 seconds, how much faster is X than Y?

### Relative Performance Example

If computer X runs a program in 10 seconds and computer Y runs the same program in 15 seconds, how much faster is X than Y?

$$\frac{\text{Performance}_{X}}{\text{Performance}_{Y}} = \frac{\text{Execution time}_{Y}}{\text{Execution time}_{X}} =$$

### Relative Performance Example

If computer X runs a program in 10 seconds and computer Y runs the same program in 15 seconds, how much faster is X than Y?

$$\frac{\text{Performance}_{X}}{\text{Performance}_{Y}} = \frac{\text{Execution time}_{Y}}{\text{Execution time}_{X}} = \frac{15}{10} = 1.5$$

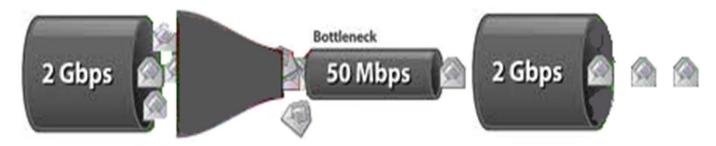
Computer X is 1.5 times faster than computer Y.

To improve performance we need to decrease execution time.

### How can we improve performance?

Find the bottleneck.

#### CPU limited task



Input ---> Processing ---> Output



I/O limited task

#### **CPU Performance**

If we limit ourselves to CPU performance, we obtain the expression for CPU execution time.

Clock rate:  $1 \text{ GHz} = 10^9 \text{ cycles / s (cycle time } 10^{-9} \text{ s} = 1 \text{ ns})$  $200 \text{ MHz} = 200 \times 10^6 \text{ cycles / s (cycle time } = 5 \text{ ns})$ Clock period

#### Instruction Performance

The previous analysis did not take into consideration the number of instructions in the program..

Average clock cycles per instruction is often abbreviated to CPI

CPU time = Instructions count x CPI x Clock cycle time

#### Example

Suppose we have two computers which have the same instruction set architecture.

Computer A has a clock cycle time of 250ps and a CPI of 2.0 for a certain program. Computer B has a clock cycle time of 500ps and a CPI of 1.2 for the same program. Which computer is faster for this program and by how much?

Now we can compute the CPU time for each computer.

CPU time<sub>A</sub> = 
$$I \times 2.0 \times 250 \text{ps} = 500 \times I \text{ ps}$$

We see that computer A is faster. How much faster?

Now we can compute the CPU time for each computer.

We see that computer A is faster. How much faster?

$$\frac{\text{Performance}_{A}}{\text{Performance}_{B}}^{A} = \frac{\text{Execution time}_{B}}{\text{Execution time}_{A}} = \frac{600 \text{ x I ps}}{500 \text{ x I ps}} = 1.2$$

Computer A is 1.2 times faster than computer B for this program.