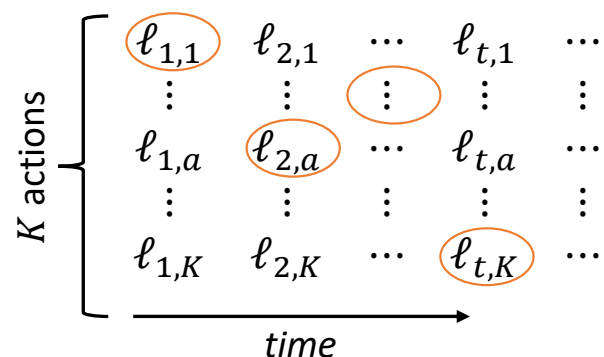
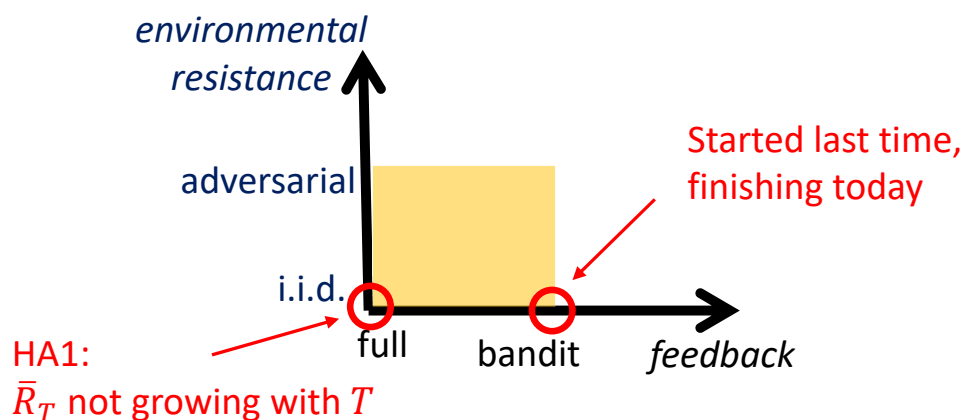


Stochastic Bandits

The UCB algorithm

Yevgeny Seldin

Quick recap of the last lecture



- Regret: $R_T = \sum_{t=1}^T \ell_{t,A_t} - \min_a \sum_{t=1}^T \ell_{t,a}$
- Expected regret: $\mathbb{E}[R_T] = \mathbb{E}[\sum_{t=1}^T \ell_{t,A_t}] - \mathbb{E}[\min_a \sum_{t=1}^T \ell_{t,a}]$
- Pseudo-regret: $\bar{R}_T = \mathbb{E}[\sum_{t=1}^T \ell_{t,A_t}] - \min_a \mathbb{E}[\sum_{t=1}^T \ell_{t,a}] = \mathbb{E}[\sum_{t=1}^T \ell_{t,A_t}] - T\mu^*$
 $= \sum_{a=1}^K \Delta(a) \mathbb{E}[N_T(a)]$

Lower Confidence Bound (LCB) algorithm for losses
 (Originally Upper Confidence Bound (UCB) for rewards)
 (“Optimism in the face of uncertainty” approach)

- Define $L_t^{CB}(a) = \hat{\mu}_{t-1}(a) - \sqrt{\frac{3 \ln t}{2N_{t-1}(a)}}$ lower confidence bound
 - (We will show that with high probability $L_t^{CB}(a) \leq \mu(a)$ for all t)

- LCB Algorithm:

- Play each arm once
- For $t = K + 1, K + 2, \dots$:
 - Play $A_t = \arg \min_a L_t^{CB}(a)$

- No knowledge of T
- No knowledge of Δ
- Works for any K

Rewards \leftrightarrow Losses

$$\begin{aligned}\ell_{t,a} &= 1 - r_{t,a} \\ r_{t,a} &= 1 - \ell_{t,a}\end{aligned}$$

- Theorem:

$$\bar{R}_T \leq 6 \sum_{a: \Delta(a) > 0} \frac{\ln T}{\Delta(a)} + \left(1 + \frac{\pi^2}{3}\right) \sum_a \Delta(a)$$

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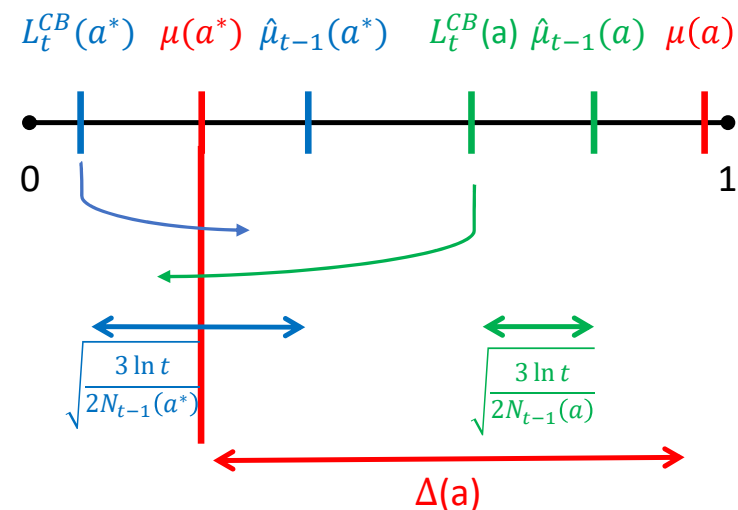
- Theorem:

$$\bar{R}_T \leq 6 \sum_{a: \Delta(a) > 0} \frac{\ln T}{\Delta(a)} + \left(1 + \frac{\pi^2}{3}\right) \sum_a \Delta(a)$$

- Proof:

- $\bar{R}_T = \sum_{a=1}^K \Delta(a) \mathbb{E}[N_T(a)]$
- When can we play $a \neq a^*$?
- Bound the expected number of times $L_t^{CB}(a) \leq L_t^{CB}(a^*)$

Proof



- $\bar{R}_t(a) = \sum_a \Delta(a) \mathbb{E}[N_T(a)]$

- $L_t^{CB}(a) = \hat{\mu}_{t-1}(a) - \sqrt{\frac{3 \ln t}{2N_{t-1}(a)}}$

- Bound the expected number of times $L_t^{CB}(a) \leq L_t^{CB}(a^*)$

- The expected number of times $L_t^{CB}(a) \leq L_t^{CB}(a^*)$ is bounded by

1. The expected number of times $L_t^{CB}(a^*) \geq \mu(a^*)$

2. Plus expected the number of times $L_t^{CB}(a) \leq \mu(a^*)$

Proof continued

1. The expected number of times $L_t^{CB}(a^*) \geq \mu(a^*)$ is bounded by

The expected number of times $\hat{\mu}_{t-1}(a^*) \geq \mu(a^*) + \sqrt{\frac{3 \ln t}{2N_{t-1}(a^*)}}$

(\longleftrightarrow “breaks” – the confidence bound fails)

2. The expected the number of times $L_t^{CB}(a) \leq \mu(a^*)$ is bounded by

2.1 The expected number of times $\hat{\mu}_{t-1}(a) \leq \mu(a) - \sqrt{\frac{3 \ln t}{2N_{t-1}(a)}}$

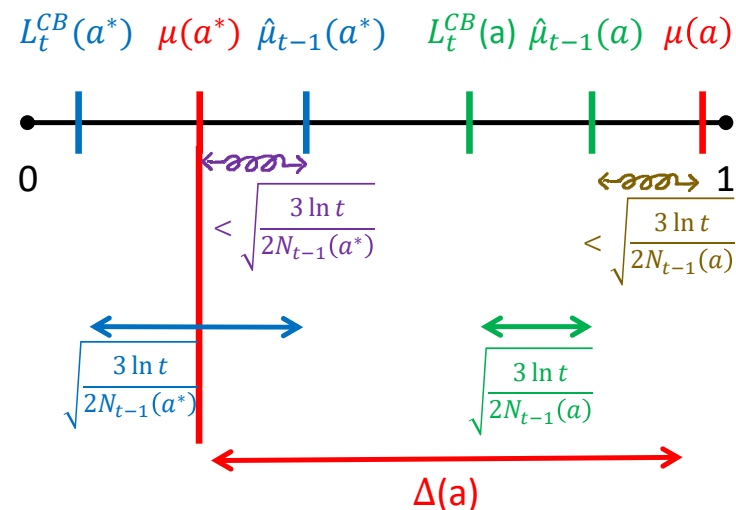
(\longleftrightarrow “breaks” – the confidence bound fails)

2.2 Plus the number of times $\sqrt{\frac{3 \ln t}{2N_{t-1}(a)}} + \sqrt{\frac{3 \ln t}{2N_{t-1}(a)}} = \sqrt{\frac{6 \ln t}{N_{t-1}(a)}} \geq \Delta(a)$

(The confidence bound holds, but it is not tight enough)

$$\Rightarrow N_{t-1}(a) \leq \frac{6 \ln t}{\Delta(a)^2} \leq \frac{6 \ln T}{\Delta(a)^2}$$

- Mid-summary: $\mathbb{E}[N_T(a)] \leq \frac{6 \ln T}{\Delta(a)^2} + 1 + \mathbb{E}[1.] + \mathbb{E}[2.1]$



Proof continued

- $\mathbb{E}[N_T(a)] \leq \frac{6 \ln T}{\Delta(a)^2} + 1 + \mathbb{E}[\text{~~blue~~}] + \mathbb{E}[\text{~~orange~~}]$
- Let $F(a^*)$ be the expected number of times $\hat{\mu}_{t-1}(a^*) \geq \mu(a^*) + \sqrt{\frac{3 \ln t}{2N_{t-1}(a^*)}}$
- Bound $\mathbb{P}\left(\hat{\mu}_{t-1}(a^*) - \mu(a^*) \geq \sqrt{\frac{3 \ln t}{2N_{t-1}(a^*)}}\right)$ — $N_{t-1}(a^*)$ is a random variable dependent on $\hat{\mu}_t(a^*)$!

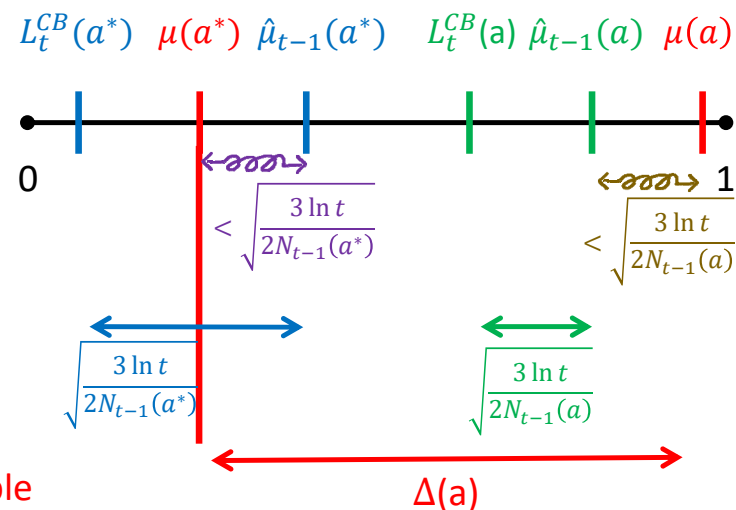
- Idea: break dependent events into independent events and take a union bound
- Introduce X_1, \dots, X_T r.v. with the same distribution as ℓ_{t,a^*}

- Let $\bar{\mu}_s = \frac{1}{s} \sum_{i=1}^s X_i$

- $\mathbb{P}\left(\hat{\mu}_{t-1}(a^*) - \mu(a^*) \geq \sqrt{\frac{3 \ln t}{2N_{t-1}(a^*)}}\right) \leq \mathbb{P}\left(\exists s \in \{1, \dots, t\}: \bar{\mu}_s - \mu(a^*) \geq \sqrt{\frac{\ln t^3}{2s}}\right)$

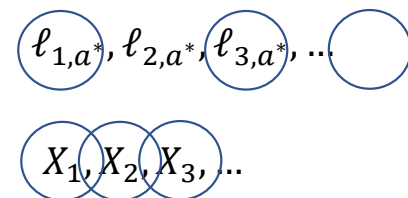
$$\begin{aligned} & \stackrel{\text{union}}{\leq} \sum_{s=1}^t \mathbb{P}\left(\bar{\mu}_s - \mu(a^*) \geq \sqrt{\frac{\ln t^3}{2s}}\right) \\ & \stackrel{\text{Hoeffding}}{\leq} \sum_{s=1}^t \frac{1}{t^3} = \frac{1}{t^2} \end{aligned}$$

- $\mathbb{E}[F(a^*)] = \sum_{t=1}^{\infty} \mathbb{P}\left(L_t^{CB}(a^*) \geq \mu(a^*)\right) \leq \sum_{t=1}^{\infty} \frac{1}{t^2} \leq \frac{\pi^2}{6}$



Hoeffding:

$$\mathbb{P}\left(\frac{1}{n} \sum_{i=1}^n Z_i - \mu \geq \sqrt{\frac{\ln \frac{1}{\delta}}{2n}}\right) \leq \delta$$



Proof summary

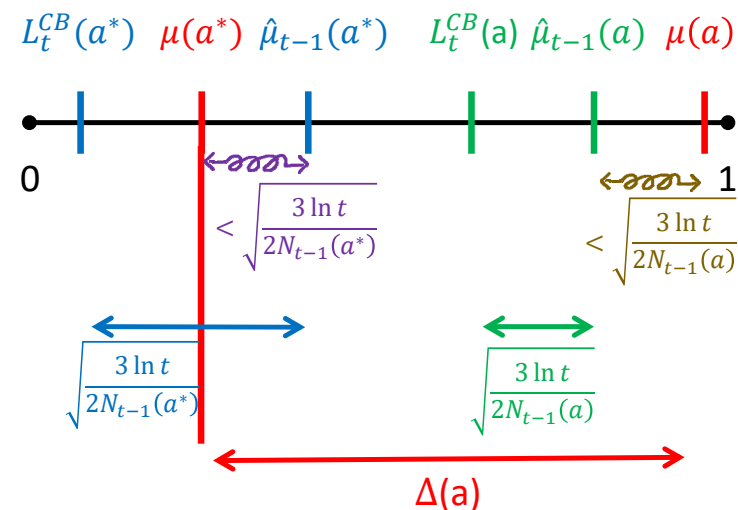
- $\bar{R}_t(a) = \sum_a \Delta(a) \mathbb{E}[N_T(a)]$

- $\mathbb{E}[N_T(a)] \leq \underbrace{\frac{6 \ln T}{\Delta(a)^2}}_{\text{The time it takes for confidence intervals to start working}} + 1 + \underbrace{\frac{\pi^2}{6} + \frac{\pi^2}{6}}_{\text{The expected number of times confidence intervals fail}}$

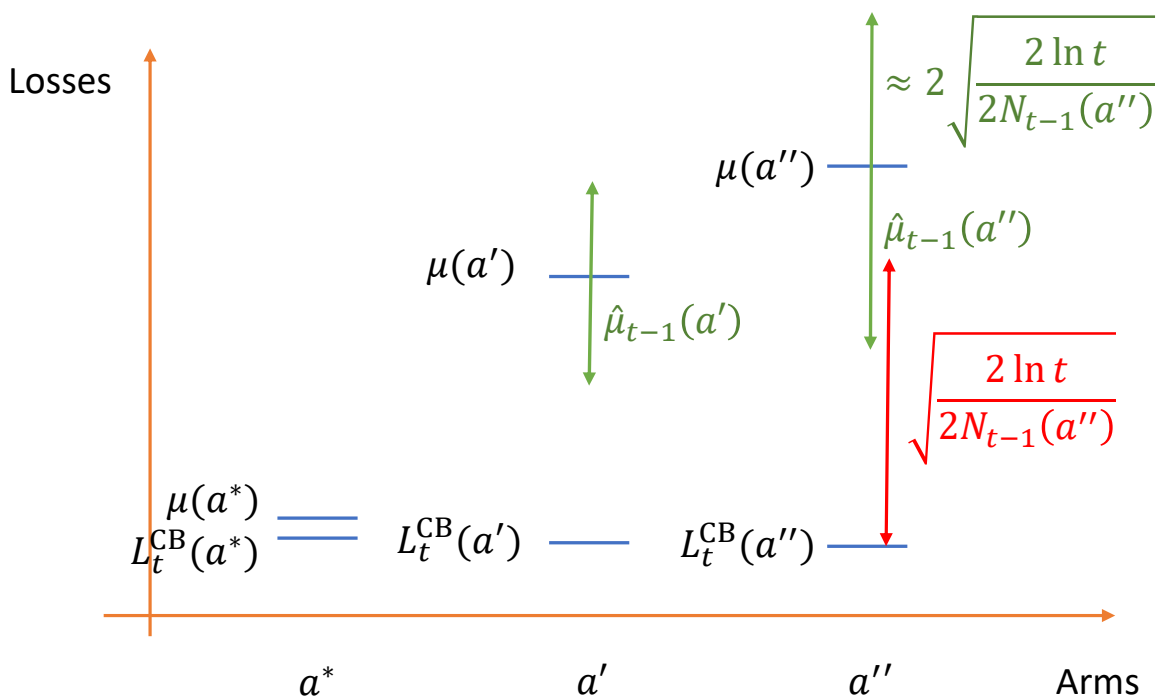
- $\bar{R}_T \leq 6 \sum_{a: \Delta(a) > 0} \frac{\ln T}{\Delta(a)} + \left(1 + \frac{\pi^2}{3}\right) \sum_a \Delta(a)$

- Home assignment:

- Take $L_t^{CB}(a) = \hat{\mu}_{t-1}(a) - \sqrt{\frac{2 \ln t}{2N_{t-1}(a)}}$ (instead of $L_t^{CB}(a) = \hat{\mu}_{t-1}(a) - \sqrt{\frac{3 \ln t}{2N_{t-1}(a)}}$; i.e. confidence $\frac{1}{t^2}$ instead $\frac{1}{t^3}$)
- Show $\bar{R}_T \leq 4 \sum_{a: \Delta(a) > 0} \frac{\ln T}{\Delta(a)} + (2 \ln T + 3) \sum_a \Delta(a)$



LCB algorithm dynamics (with $L_t^{CB}(a) = \hat{\mu}_{t-1}(a) - \sqrt{\frac{2 \ln t}{2N_{t-1}(a)}}$)



- Confidence interval of the played arm shrinks ($N_{t-1}(a)$ grows)
- Confidence intervals of all other arms grow ($\ln t$ grows)
- \Rightarrow all LCBs are roughly at the same level
- Most of the time $L_t^{CB}(a^*) \leq \mu(a^*)$
- a^* is played a lot, so $L_t^{CB}(a^*)$ is very close to $\mu(a^*)$
- All other arms are played just enough to keep $\sqrt{\frac{2 \ln t}{2N_{t-1}(a)}} = \theta(\Delta(a))$, i.e. $N_t(a) = \theta\left(\frac{\ln t}{\Delta(a)^2}\right)$