

ICS 2020 Problem Sheet #2

Problem 2.1:

a)

The Big O notation (Landau notation) definition :

$$O(g) = \{f \mid \exists k \in \mathbb{N}. f \leq k \cdot g\}$$

$$\text{for } t_1(n) = 5n^2 + 16$$

$$5n^2 + 16 \leq 5n^2 + 16$$

$$\Rightarrow 5n^2 + 16 \leq 5n^2 + 16n^2 \quad \text{for } n_0 = 0 \text{ and } n > n_0$$

$$\Rightarrow 5n^2 + 16 \leq 21n^2 \quad // \quad 21 \neq 0 \quad // \quad n_0 = 0 \text{ and } n > n_0$$

then $t_1 \in O(n^2)$.

$$\text{for } t_2(n) = 6n^3 + n^2 + 18$$

$$6n^3 + n^2 + 18 \leq 6n^3 + n^2 + 18$$

$$\Rightarrow 6n^3 + n^2 + 18 \leq 6n^3 + n^3 + 18n^3 \quad \text{for } n_0 = 0 \text{ and } n > n_0$$

$$\Rightarrow 6n^3 + n^2 + 18 \leq 25n^3 \quad // \quad 25 \neq 0 \quad // \quad n_0 = 0 \text{ and } n > n_0$$

then $t_2 \in O(n^3)$.

b)

Logically, the entire program belongs to The big O set $O(n^3)$:

$$O(1) \subset O(\log_2(n)) \subset O(n) \subset O(n \log_2(n)) \subset O(n^2) \subset O(n^k) \subset O(n^i)$$

with $k > 2$ and $i > 1$

we conclude that:

$$O(n^2) \subset O(n^k) \Rightarrow O(n^2) \subset O(n^3)$$

c)

We have that $f_1 \in O(g_1)$ and $f_2 \in O(g_2)$, then :

$$\exists k_1 \in \mathbb{N}. f_1 \leq k_1 \cdot g_1 \text{ and } \exists k_2 \in \mathbb{N}. f_2 \leq k_2 \cdot g_2$$

$$\Rightarrow f_1 + f_2 \leq k_1 \cdot g_1 + k_2 \cdot g_2$$

$$\Rightarrow f_1 + f_2 \leq k_1 \cdot \max\{g_1, g_2\} + k_2 \cdot \max\{g_1, g_2\}$$

$$\Rightarrow f_1 + f_2 \leq (k_1 + k_2) \cdot \max\{g_1, g_2\} \quad (k_1 + k_2) \in \mathbb{N}$$

then $(f_1 + f_2) \in O(\max\{g_1, g_2\})$

Problem 2.2:

n

$$\text{let } \sum_{k=1}^n (2k - 1)^2 = X_n$$

$$k=1$$

In order to prove that $X_n = n(2n - 1)(2n + 1)/3$

We must show that the case is right for x_1 (x_1 here is our first element here) :

$$x_1 = (2 \cdot 1 - 1)^2 = 1 \text{ and for } n = 1 : n(2n - 1)(2n + 1)/3 = 1 \cdot (2 - 1) \cdot (2 + 1)/3 = 1$$

Now we suppose that: $X_n = n(2n - 1)(2n + 1)/3$

and prove that $X_{n+1} = (n+1)(2n + 1)(2n + 3)/3$

$$X_{n+1} = (n+1)(2n + 1)(2n + 3)/3$$

(Here we got the element $n+1$ out of X_{n+1})

$$\leftrightarrow X_n + (2(n+1)-1)^2 = (n+1)(2n + 1)(2n + 3)/3$$

(here we replaced X_n by $n(2n - 1)(2n + 1)/3$

$$\leftrightarrow n(2n - 1)(2n + 1)/3 + (2(n+1)-1)^2 = (n+1)(2n + 1)(2n + 3)/3$$

$$\leftrightarrow n(2n - 1)(2n + 1)/3 + (2n+1)^2 = (n+1)(2n + 1)(2n + 3)/3$$

$$\leftrightarrow (2n + 1)(n(2n - 1)/3 + (2n+1)) = (n+1)(2n + 1)(2n + 3)/3 \text{ since } n \neq -1/2$$

$$\leftrightarrow n(2n - 1)/3 + 2n+1 = (n+1)(2n + 3)/3$$

$$\leftrightarrow n(2n - 1) + 6n+3 = (n+1)(2n + 3)$$

$$\leftrightarrow 2n^2 - n + 6n+3 = 2n^2+3n+2n+3$$

$$\leftrightarrow 2n^2 + 5n+3 = 2n^2+5n+3 \text{ which is true}$$

We can now conclude that $X_n = n(2n - 1)(2n + 1)/3$

Problem 2.3:

a)

a list comprehension that include all factors of 210 would be :

```
[ x | x <- [1..210], 210 `mod` x == 0]
```

The program would then print :

```
[1,2,3,5,6,7,10,14,15,21,30,35,42,70,105,210]
```

b)

that list would be :

```
[ (x,y,z) | x <- [1..100], y <- [1..100], z <- [1..100], x*x + y*y == z*z, if (x,y,z)==(y,x,z)
then (x,y,z) 'rem' (y,x,z)]
```

When I tried it the program didn't return anything so here is a list that would do the job without removing the duplicate elements :

```
[ (x,y,z) | x <- [1..100], y <- [1..100], z <- [1..100], x*x + y*y == z*z]
```