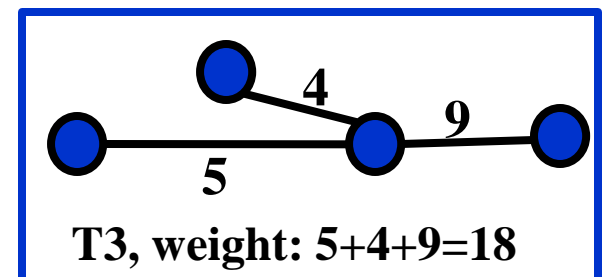
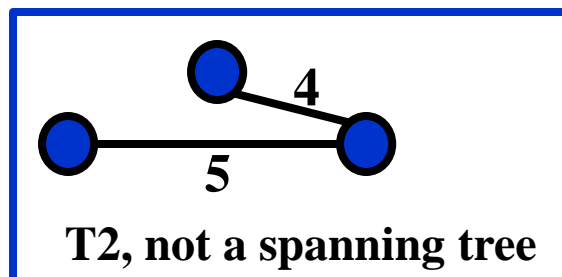
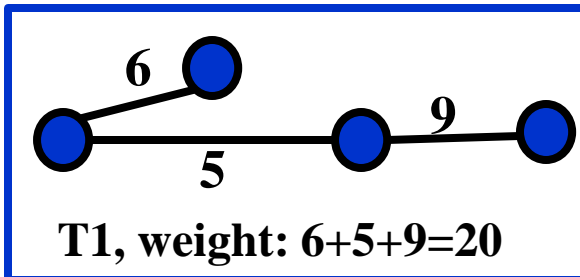
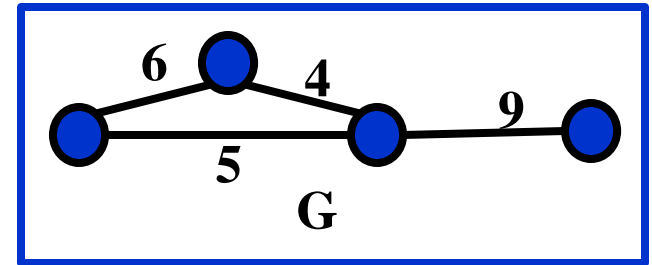


Lecture 10:

Minimum Spanning Tree (MST)

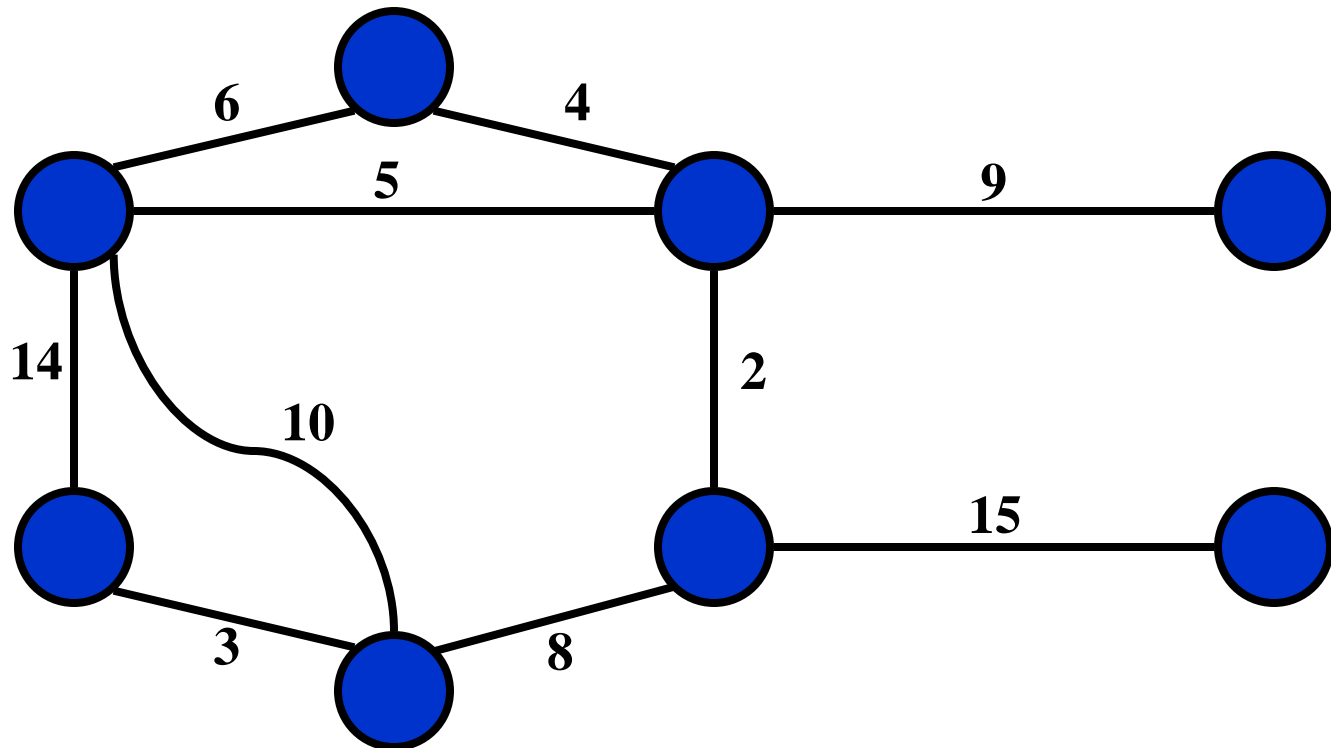
Spanning Tree and Minimum Spanning Tree

- **Spanning tree** of a connected graph: A tree that connects *all* nodes of the graph.
- **Minimum Spanning tree** of a connected weighted graph: A tree that connects all nodes of the graph with total edge weight lowest possible.
- **Example:** For the graph G,
 - T1, T3 are spanning trees
 - T2 is not a spanning tree, because it does not connect all nodes of G
 - T1 is not minimum, because its total weight is 20
 - T2 is minimum with lowest possible total weight 19

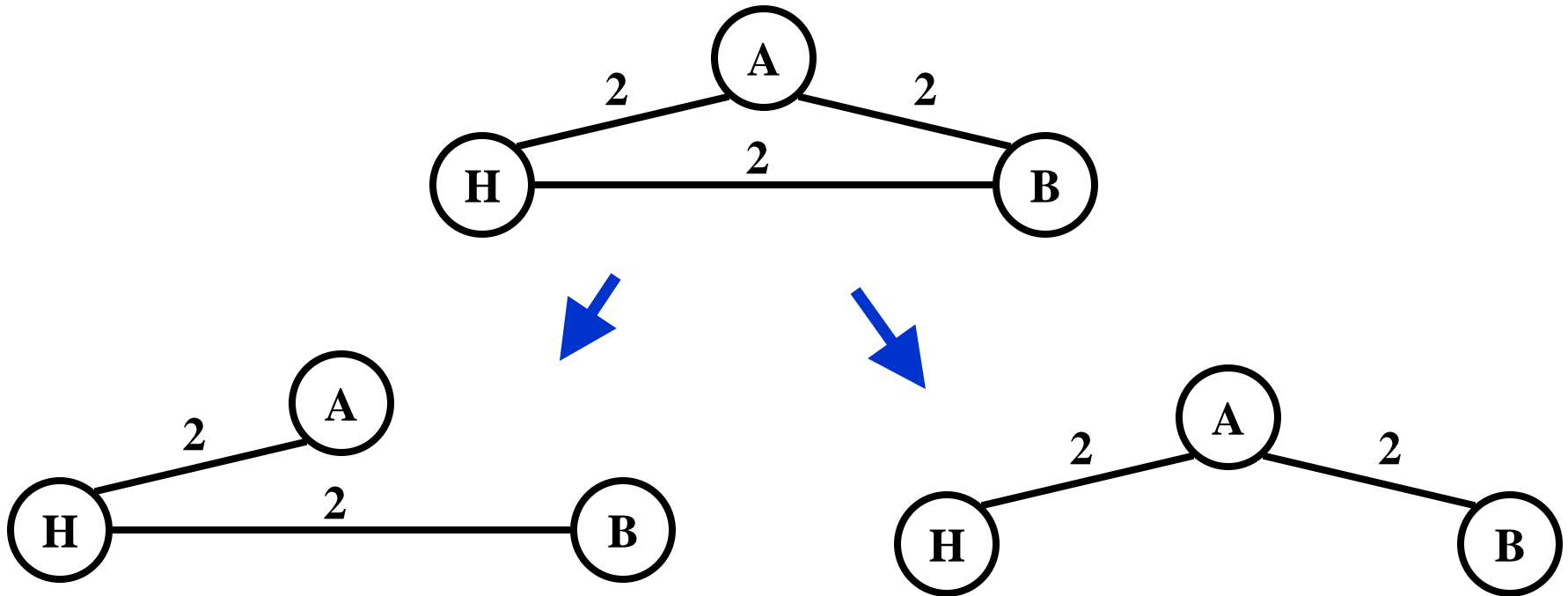


Minimum Spanning Tree

- Problem: given a connected, undirected, weighted graph, find a *minimum spanning tree*

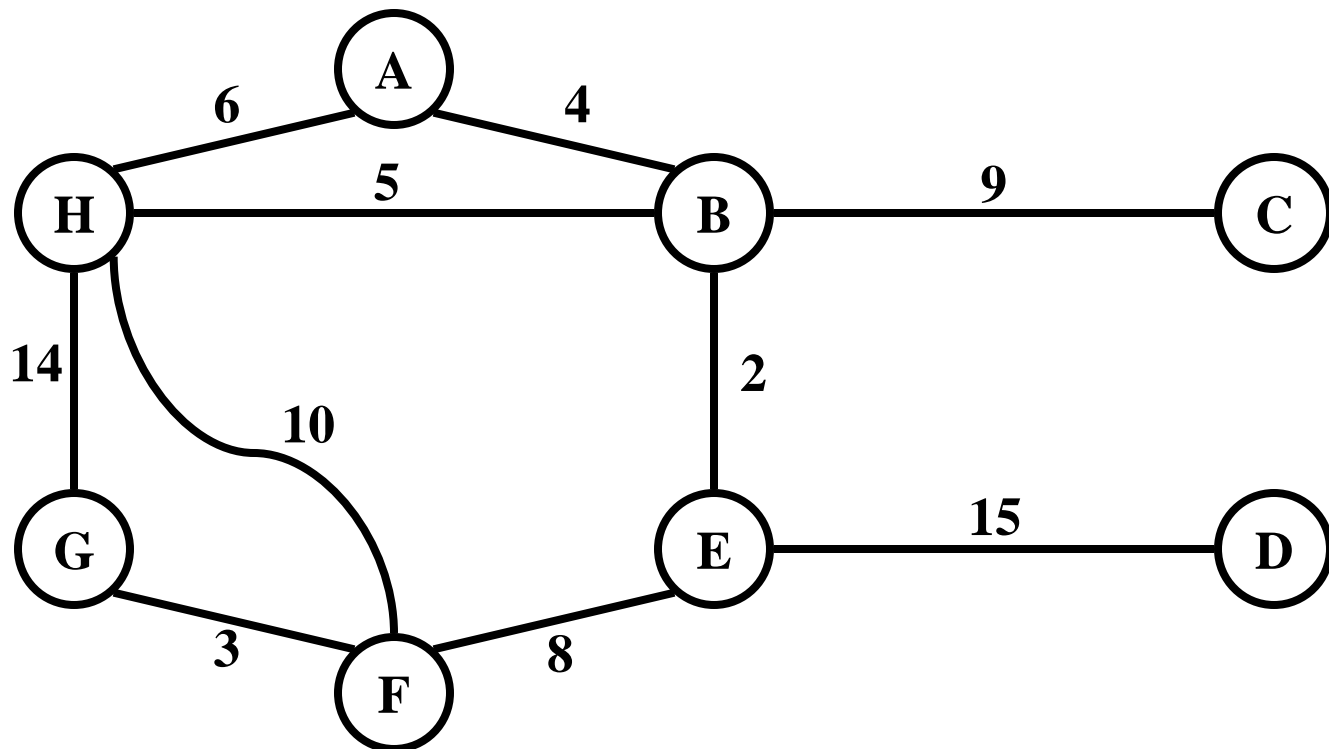


Mora Than One Minimum Spanning Tree Possible



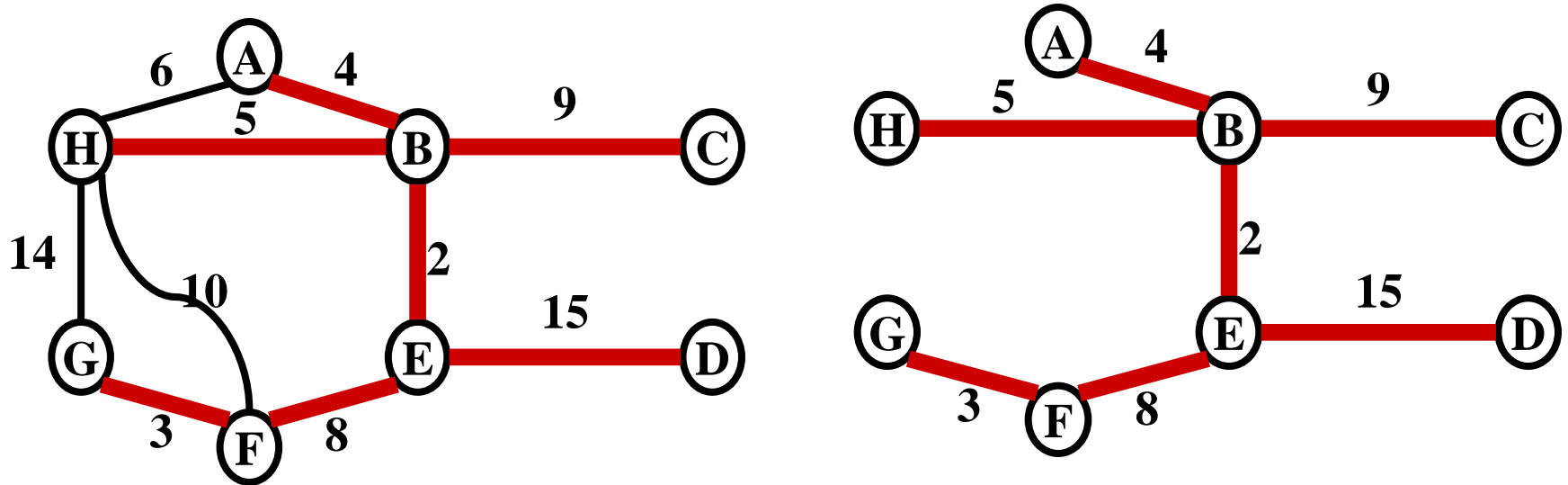
Minimum Spanning Tree

- Which edges form the minimum spanning tree (MST) of the graph below?



Minimum Spanning Tree

- Answer:



- Total weight: $3+8+4+5+2+9+15=46$. This is smallest total weight.

Two Algorithm

- Krushkal's algorithm
- Prim's algorithm

Kruskal's Algorithm

- Take the minimum weight edge one after another if there is no cycle
- Use disjoint set data structure
- MakeSet() for each element
- Maintain two sets:
 - tree edges
 - other remaining edges
- Use FindSet() to check for cycle (If same set, then must be cycle, because if you connect any two vertices of one set, then it makes a cycle.)

Symbols used:

E: Edge set, V: Vertex set, T: Resulting tree, n: number of vertices

Kruskal's Algorithm

```
Kruskal()
```

```
{
```

```
  T =  $\emptyset$ ;
```

```
  for each v  $\in$  V
```

```
    MakeSet(v);
```

```
  sort E by increasing edge weight w
```

```
  for each (u,v)  $\in$  E (in sorted order)
```

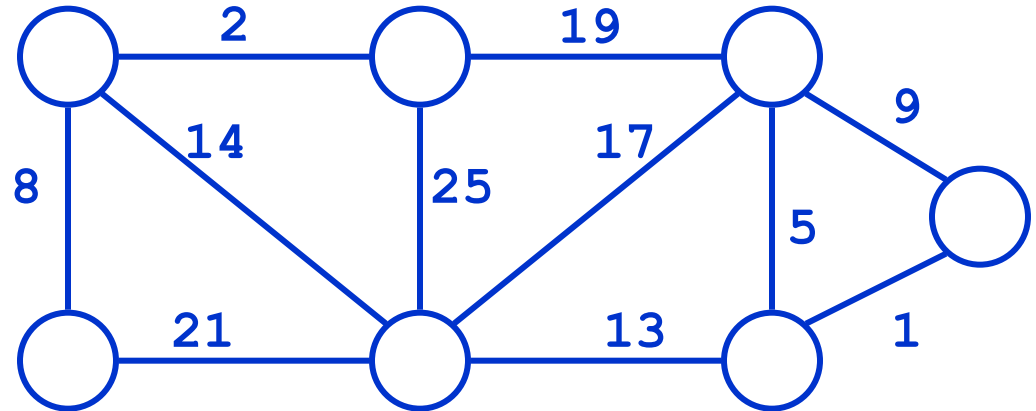
```
    if FindSet(u)  $\neq$  FindSet(v) //no cycle
```

```
      T = T  $\cup$  {(u,v)};
```

```
      Union(FindSet(u), FindSet(v));
```

```
}
```

Run the algorithm:



Kruskal's Algorithm

Kruskal()

{

$T = \emptyset;$

for each $v \in V$

MakeSet(v);

sort E by increasing edge weight w

for each $(u, v) \in E$ (in sorted order)

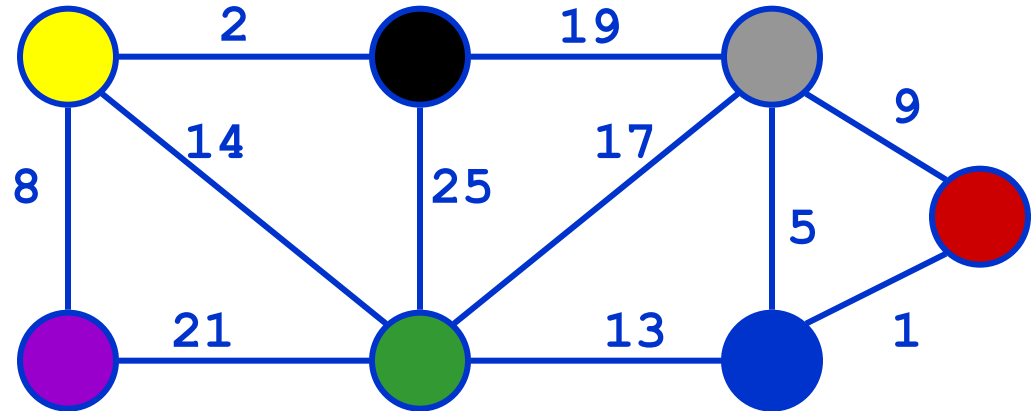
if FindSet(u) \neq FindSet(v) //no cycle

$T = T \cup \{(u, v)\};$

Union(FindSet(u), FindSet(v));

}

Run the algorithm:



Kruskal's Algorithm

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Kruskal()
```

```
{
```

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    for each v  $\in$  V
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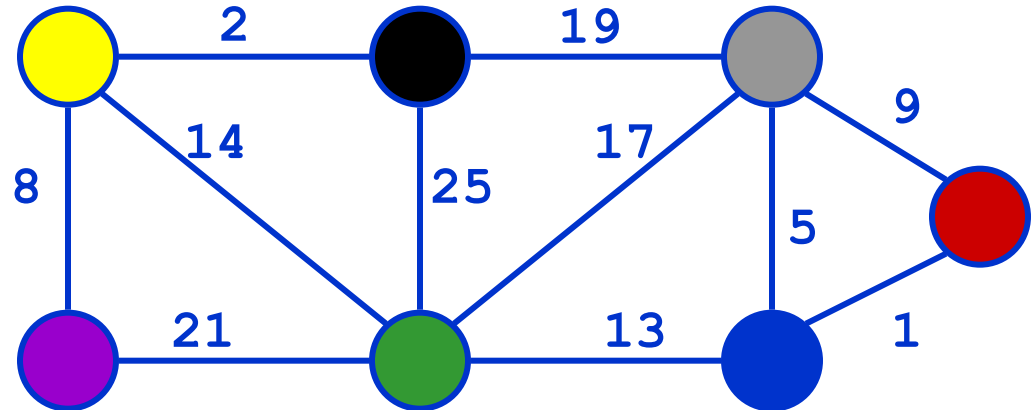
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            Union(FindSet(u), FindSet(v));
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}
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Run the algorithm:



Kruskal's Algorithm

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Kruskal()
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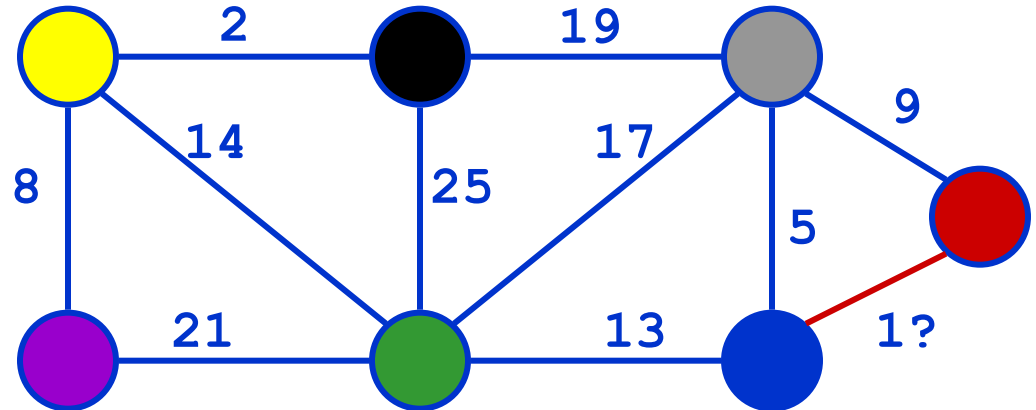
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      T = T  $\cup$  {(u,v)};
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      Union(FindSet(u), FindSet(v));
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}
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Run the algorithm:



Kruskal's Algorithm

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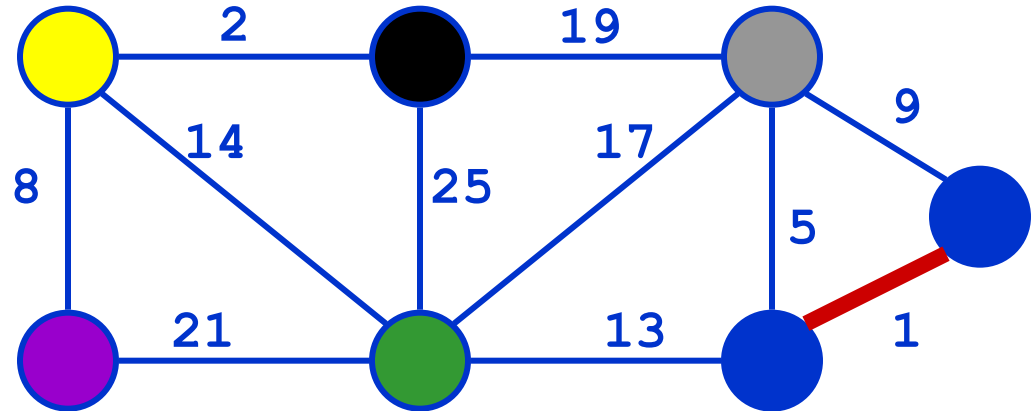
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Run the algorithm:



Kruskal's Algorithm

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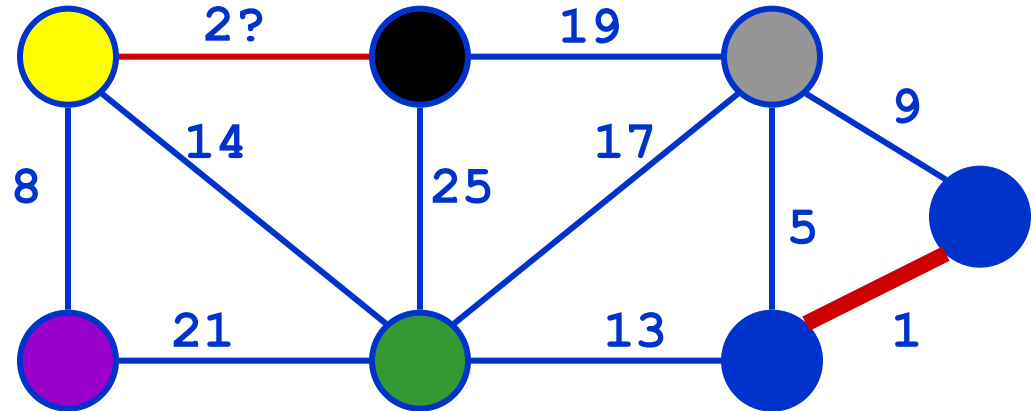
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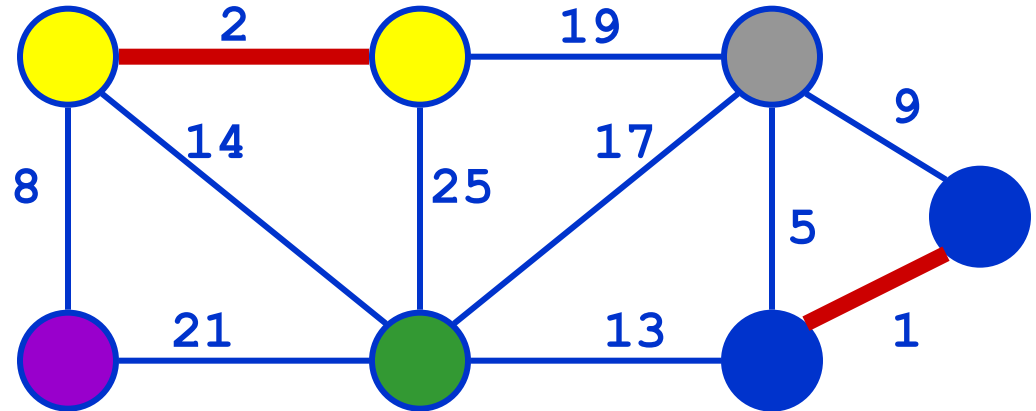
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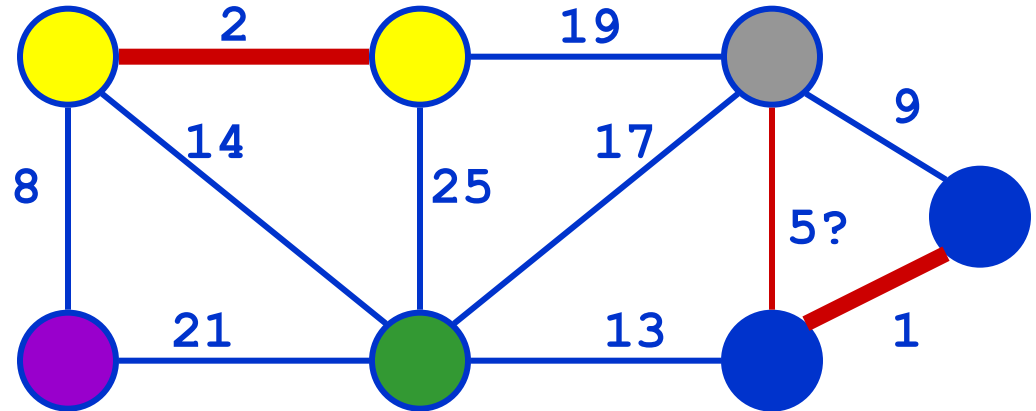
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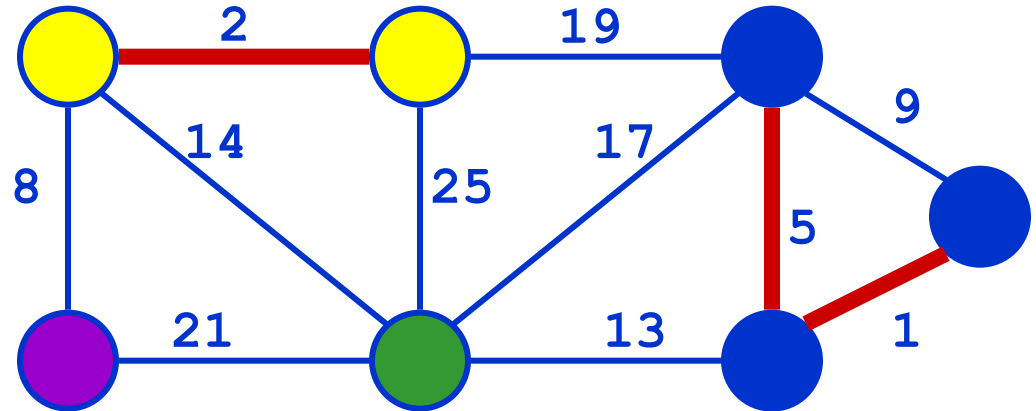
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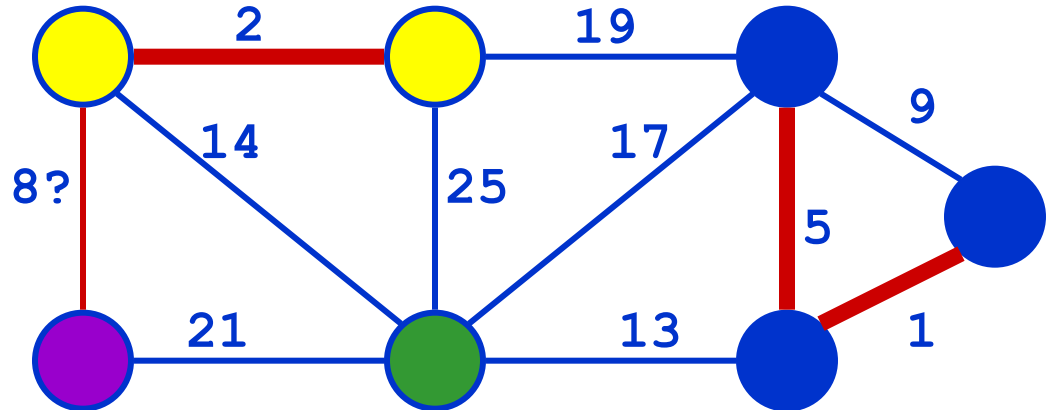
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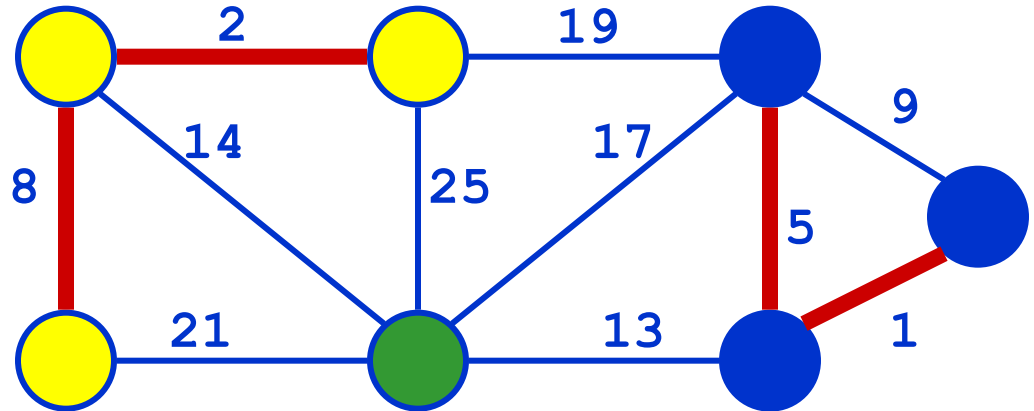
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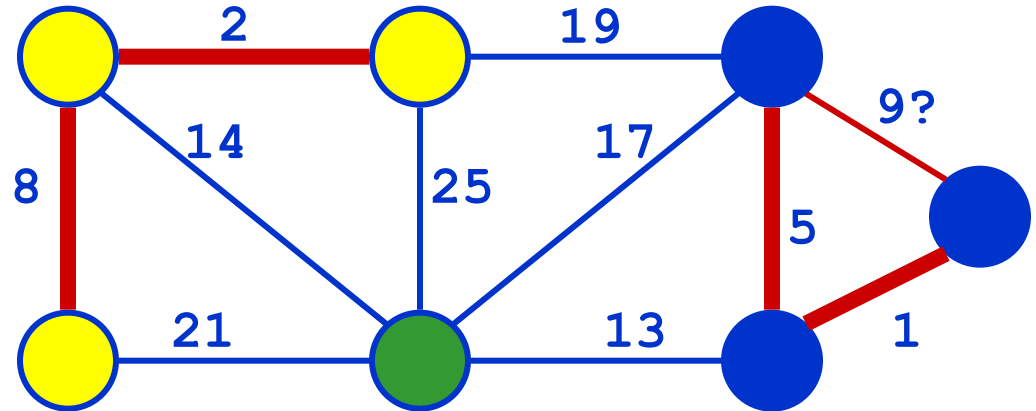
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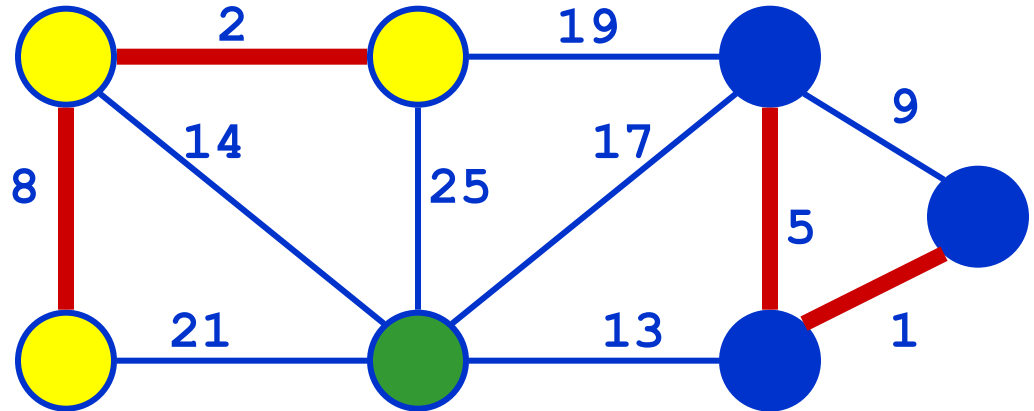
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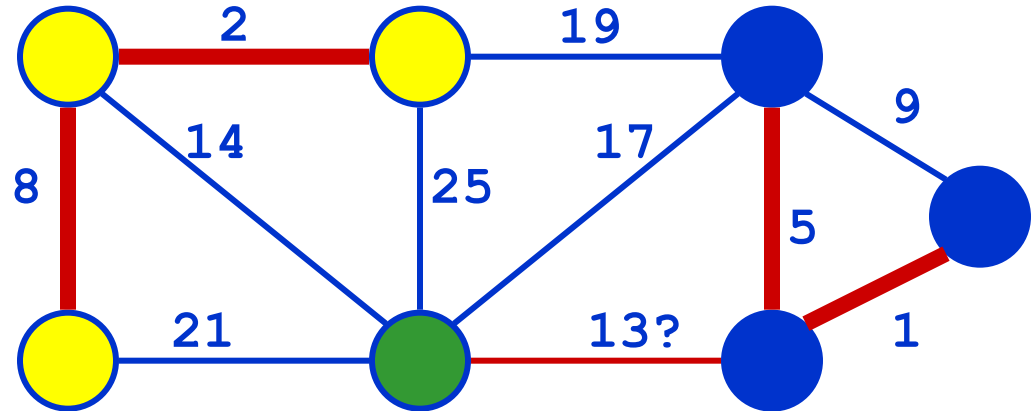
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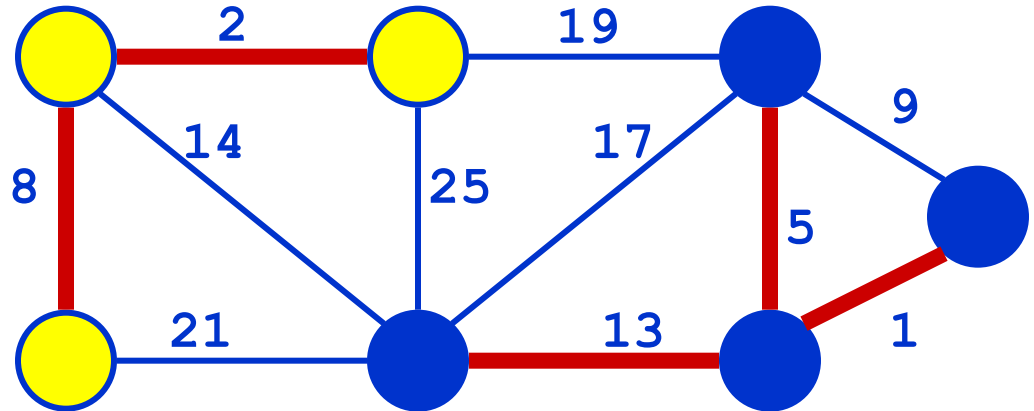
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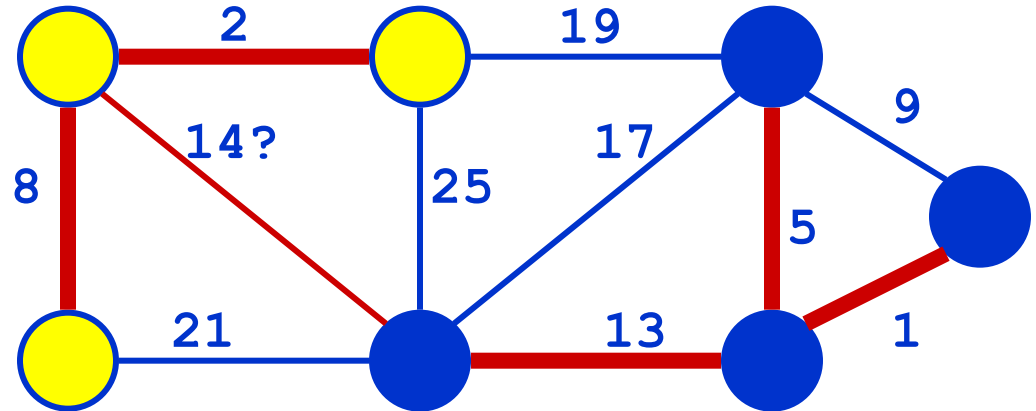
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      T = T  $\cup$  {(u,v)};
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      Union(FindSet(u), FindSet(v));
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}
```

Run the algorithm:



Kruskal's Algorithm

```
Kruskal()
```

```
{
```

```
  T =  $\emptyset$ ;
```

```
  for each  $v \in V$ 
```

```
    MakeSet( $v$ );
```

```
  sort E by increasing edge weight w
```

```
  for each  $(u,v) \in E$  (in sorted order)
```

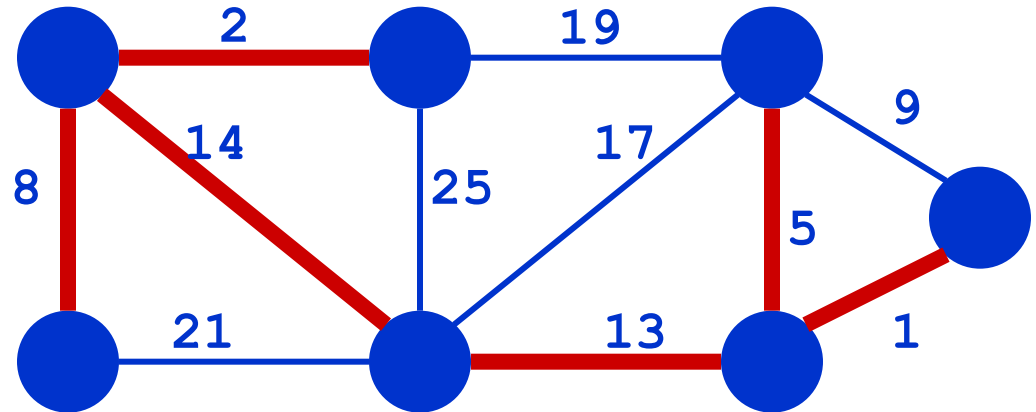
```
    if FindSet( $u$ )  $\neq$  FindSet( $v$ ) //no cycle
```

```
      T = T  $\cup$  { $\{u,v\}$ };
```

```
      Union(FindSet( $u$ ), FindSet( $v$ ));
```

```
}
```

Run the algorithm:



Kruskal's Algorithm

```
Kruskal()
```

```
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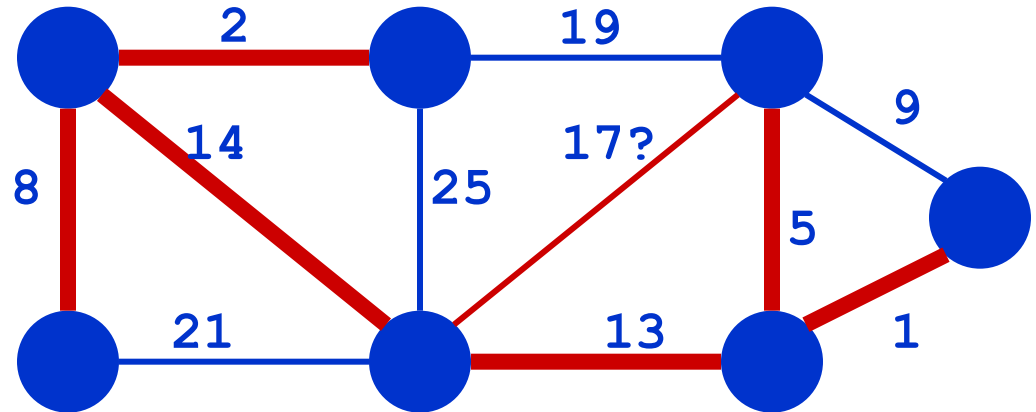
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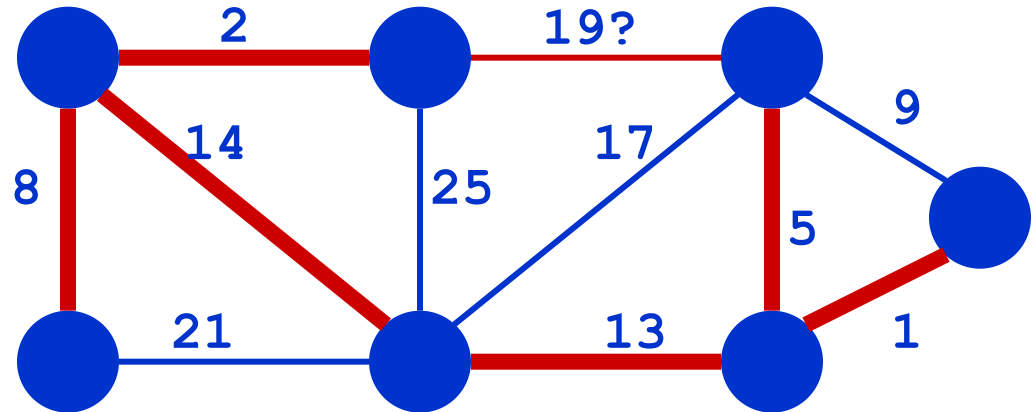
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Kruskal's Algorithm

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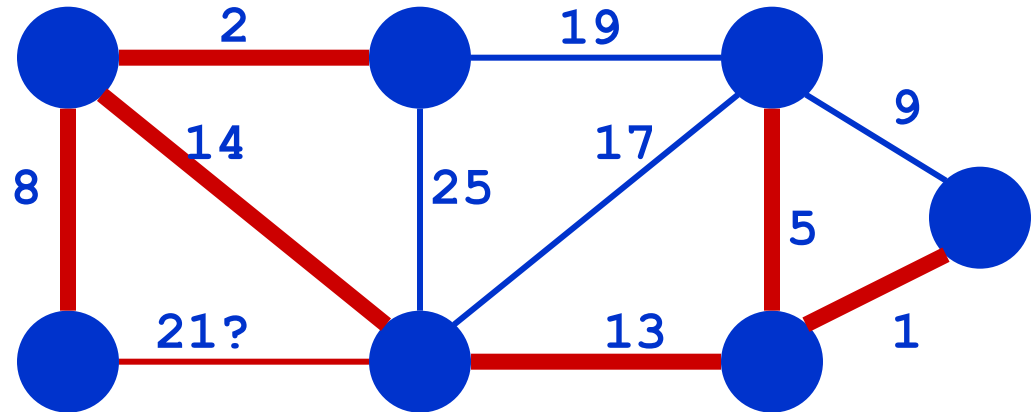
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Run the algorithm:



Kruskal's Algorithm

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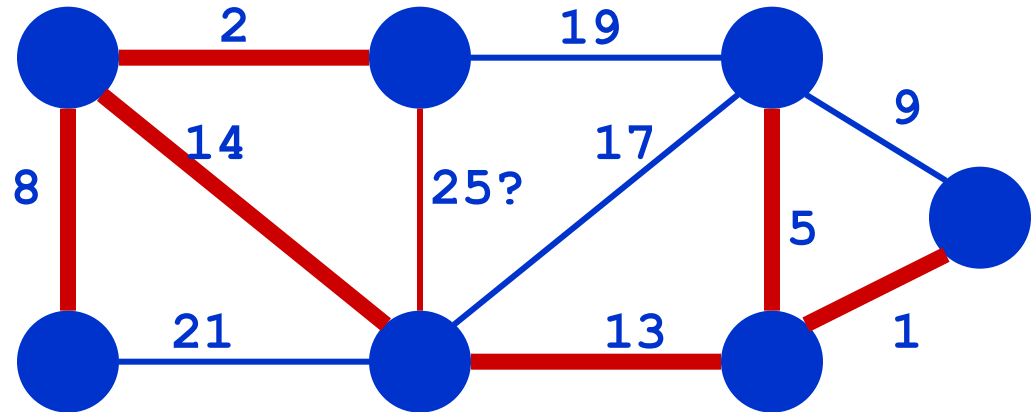
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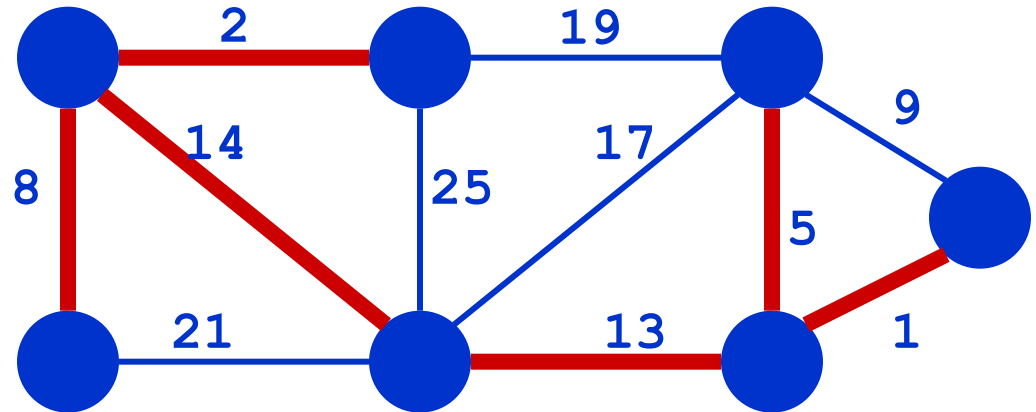
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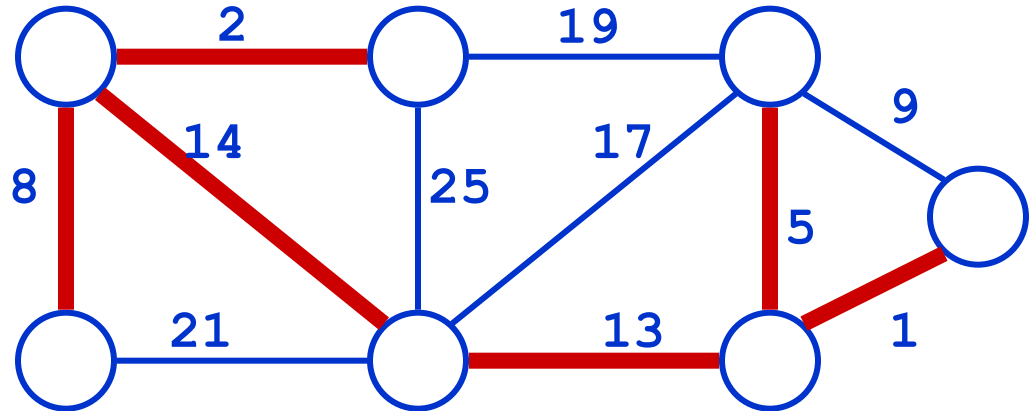
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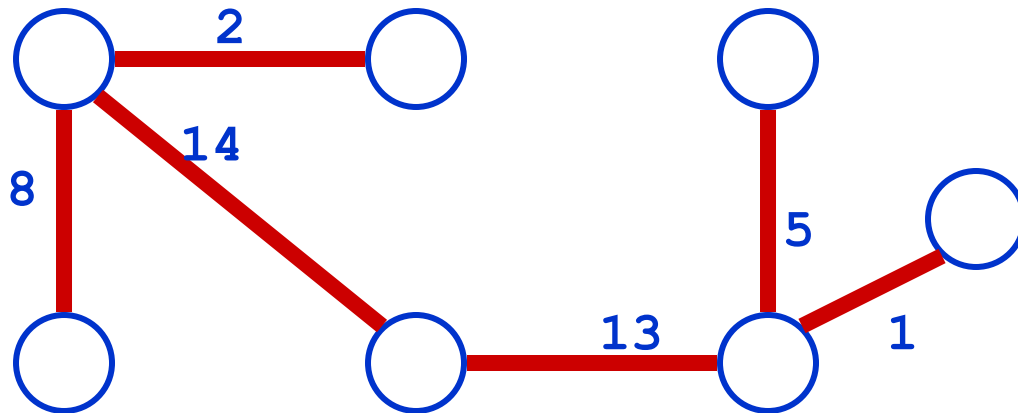
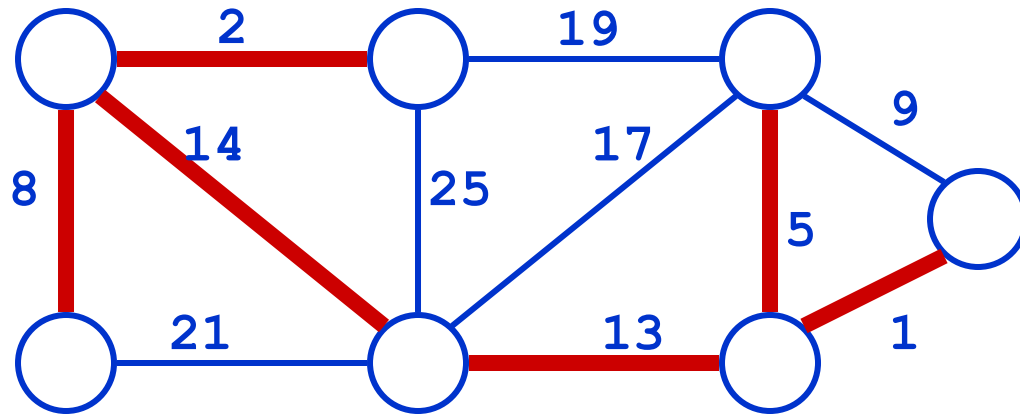
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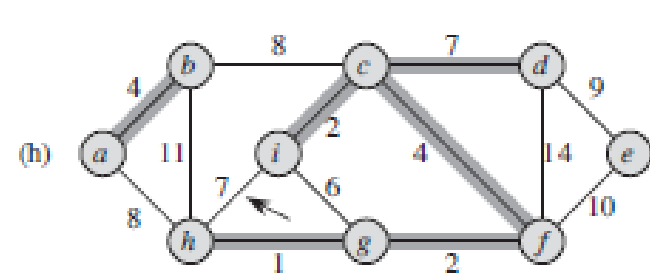
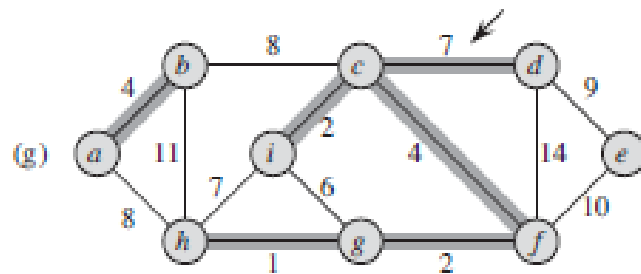
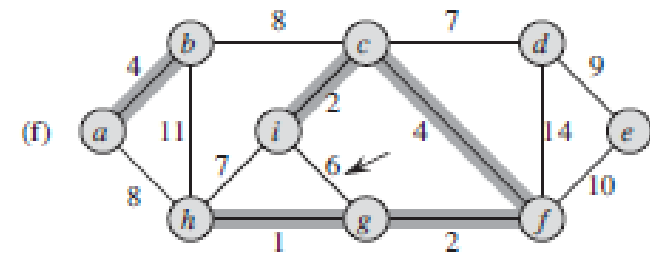
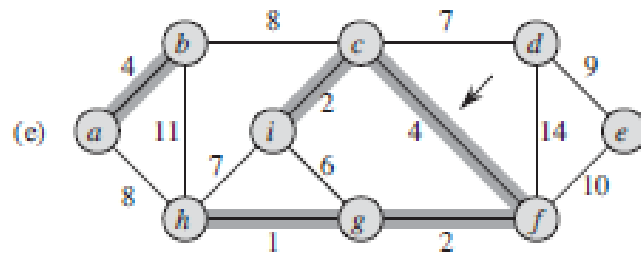
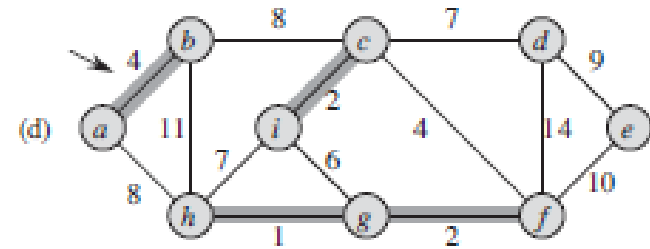
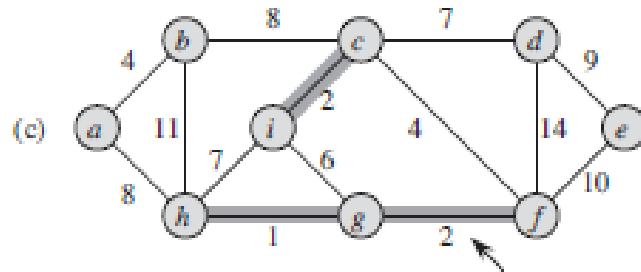
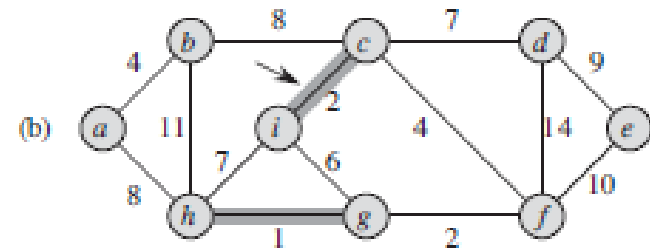
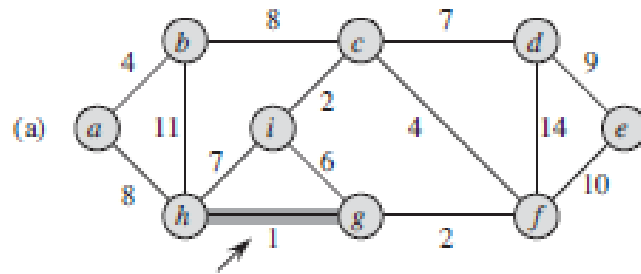
Run the algorithm:



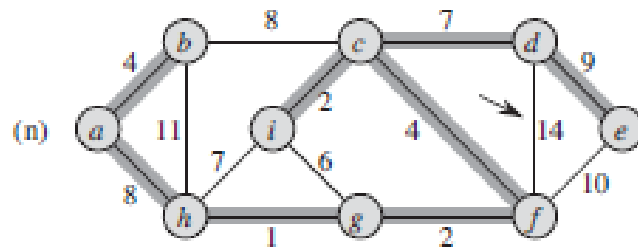
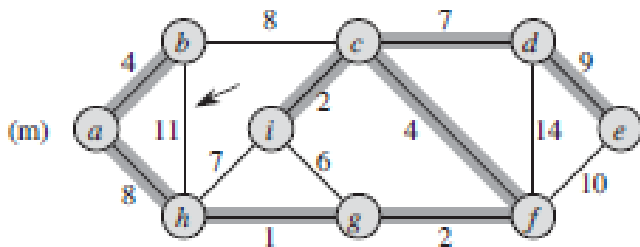
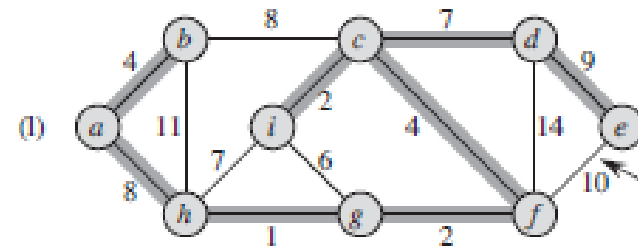
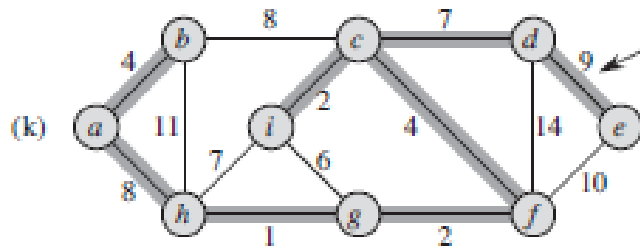
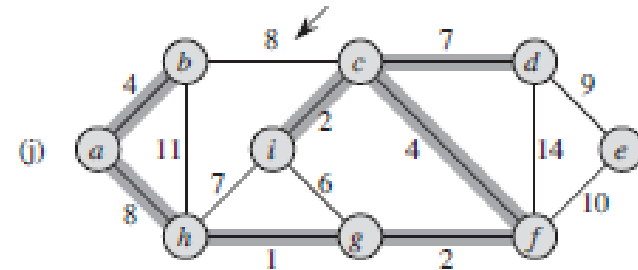
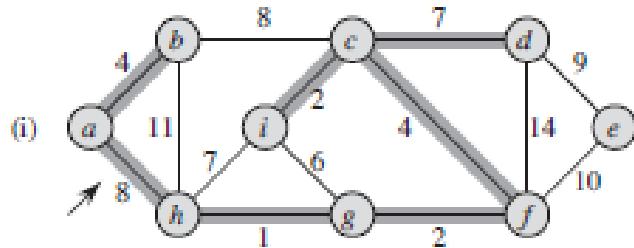
Kruskal's Algorithm



Another Example: Krushkal's Algorithm







Another Example: Krushkal's Algorithm



Cost of Kruskal's Algorithm

```
Kruskal()
```

```
{  
    T =  $\emptyset$ ;  
    for each v  $\in$  V   $O(n)$   
        MakeSet(v);   $O(E \log E)$   
    sort E by increasing edge weight w  
    for each (u,v)  $\in$  E (in sorted order)  
        if FindSet(u)  $\neq$  FindSet(v)   $O(1)$   
            T = T  $\cup$  {{u,v}};  
            Union(FindSet(u), FindSet(v));  
    }  Total E Union():  $O(E \log E)$ 
```

Total Cost: $O(n) + O(E \log E) + O(E \log E) = O(E \log E)$

E: Number of edges, and $E \geq n$

Prim's Algorithm

1. Initially, all node have weight ∞
2. Keep all of them in a queue (heap)
3. Take minimum node from heap (ExtractMin())
4. For each neighbor of this node, if edge weight from this node is smaller than current weight , then update weight (DecreaseKey()) and add that edge in the tree.
5. Repeat step 3, 4 until finished.

n: number of vertex

Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$   
     $\text{weight}[u] = \infty$ ;  
 $Q = \text{all node}$ ;  
 $\text{weight}[\text{root}] = 0$ ;  
 $p[\text{root}] = \text{NULL}$ ;  
while ( $Q$  not empty)
```

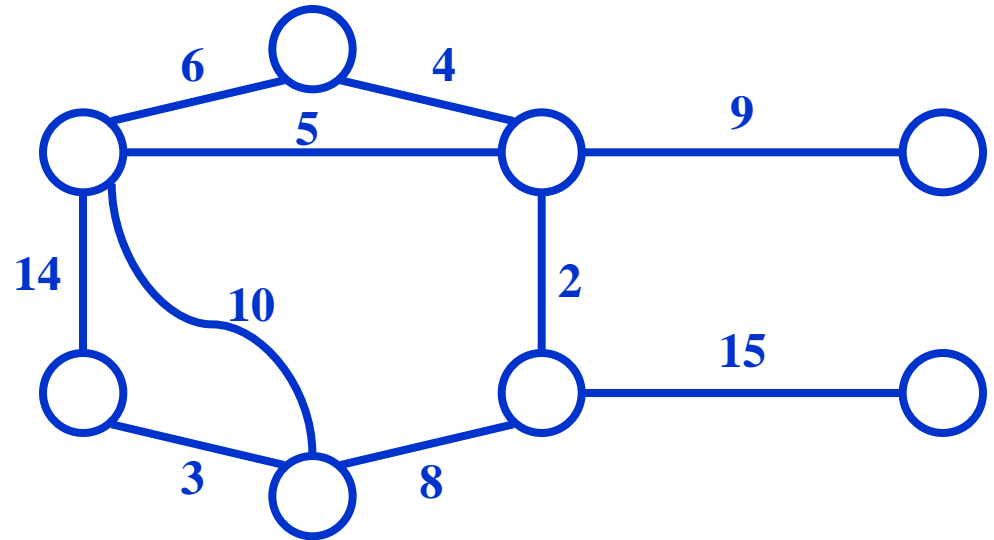
```
     $u = \text{ExtractMin}(Q)$  ;
```

```
    for each  $v \in \text{Adjacent}[u]$ 
```

```
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
```

```
             $p[v] = u$ ;
```

```
             $\text{weight}[v] = w(u, v)$  ;
```



Run on example graph

Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
```

```
     $\text{weight}[u] = \infty$ ;
```

```
Q = all node;
```

```
 $\text{weight}[\text{root}] = 0$ ;
```

```
p[root] = NULL;
```

```
while (Q not empty)
```

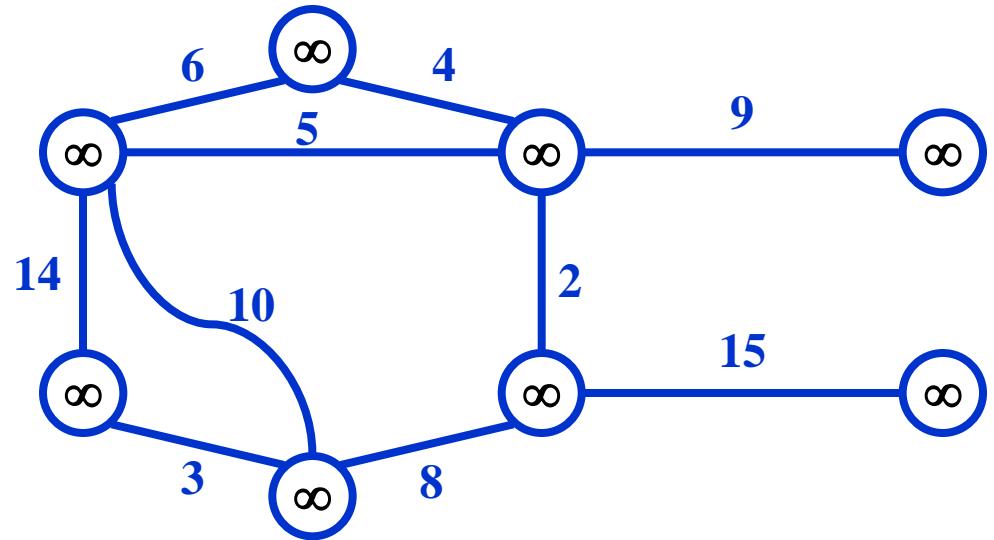
```
     $u = \text{ExtractMin}(Q)$ ;
```

```
    for each  $v \in \text{Adjacent}[u]$ 
```

```
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
```

```
            p[v] = u;
```

```
             $\text{weight}[v] = w(u, v)$ ;
```



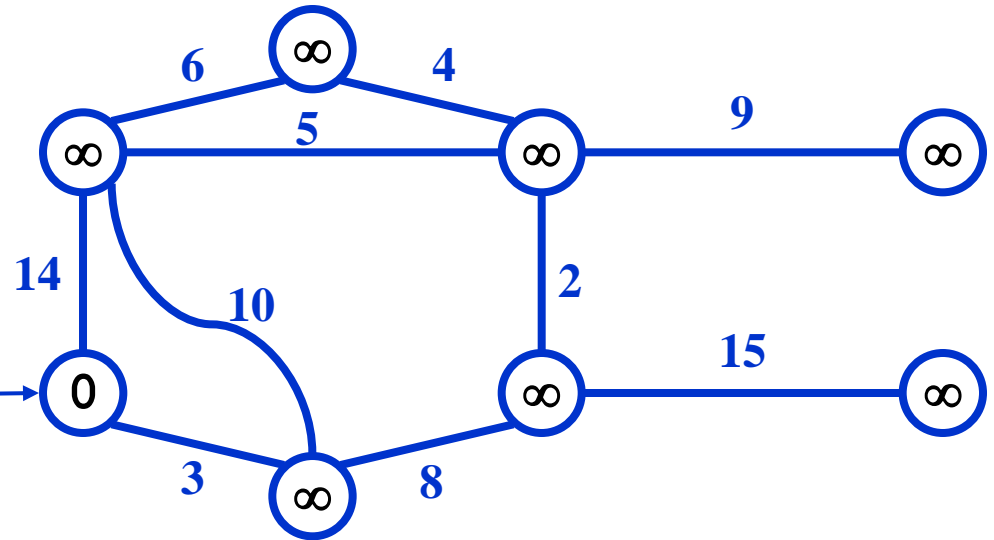
Run on example graph

Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
     $\text{weight}[u] = \infty$ ;
 $Q = \text{all node}$ ;
 $\text{weight}[\text{root}] = 0$ ;  $r \rightarrow$ 
 $p[\text{root}] = \text{NULL}$ ;
while ( $Q$  not empty)
     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```



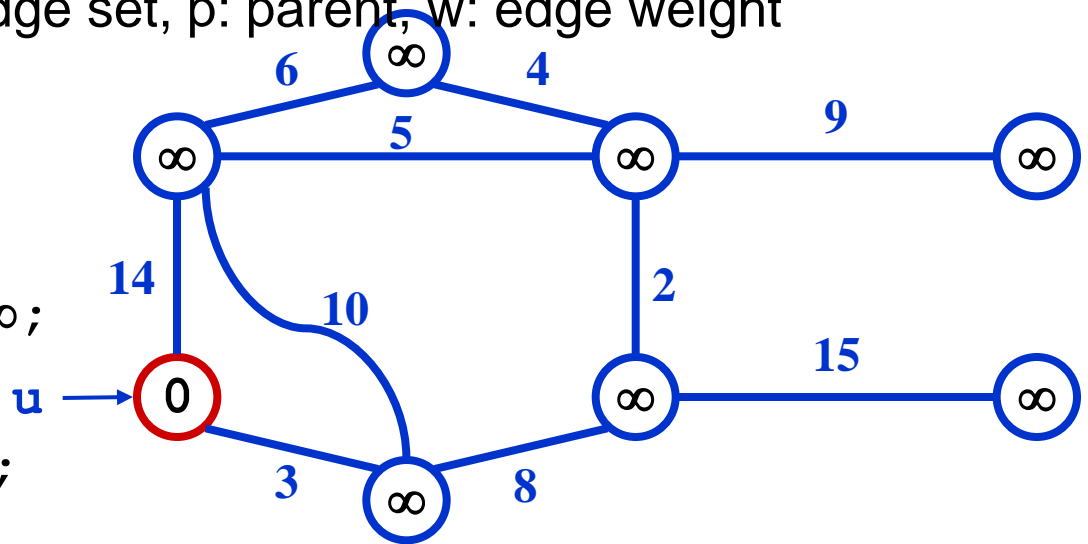
Pick a start vertex r

Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
     $\text{weight}[u] = \infty$ ;
 $Q = \text{all node}$ ;
 $\text{weight}[\text{root}] = 0$ ;
 $p[\text{root}] = \text{NULL}$ ;
while ( $Q$  not empty)
     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```



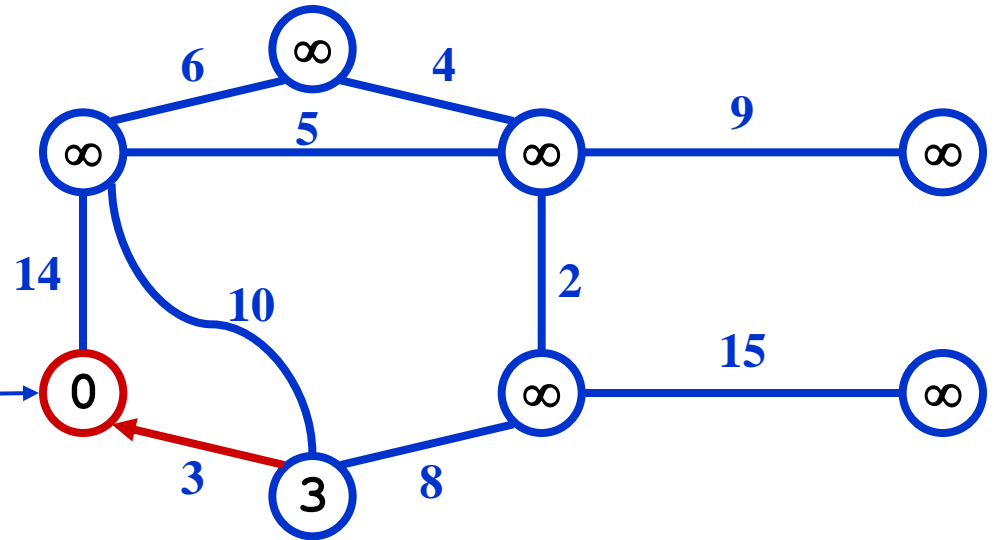
Red vertices have been removed from Q

Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
     $\text{weight}[u] = \infty$ ;
 $Q = \text{all node}$ ;
 $\text{weight}[\text{root}] = 0$ ;
 $p[\text{root}] = \text{NULL}$ ;
while ( $Q$  not empty)
     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```



Red arrows indicate parent pointers

Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
```

```
weight[u] =  $\infty$ ;
```

Q = all node;

```
weight[root] = 0; u
```

```
p[root] = NULL;
```

```
while (Q not empty)
```

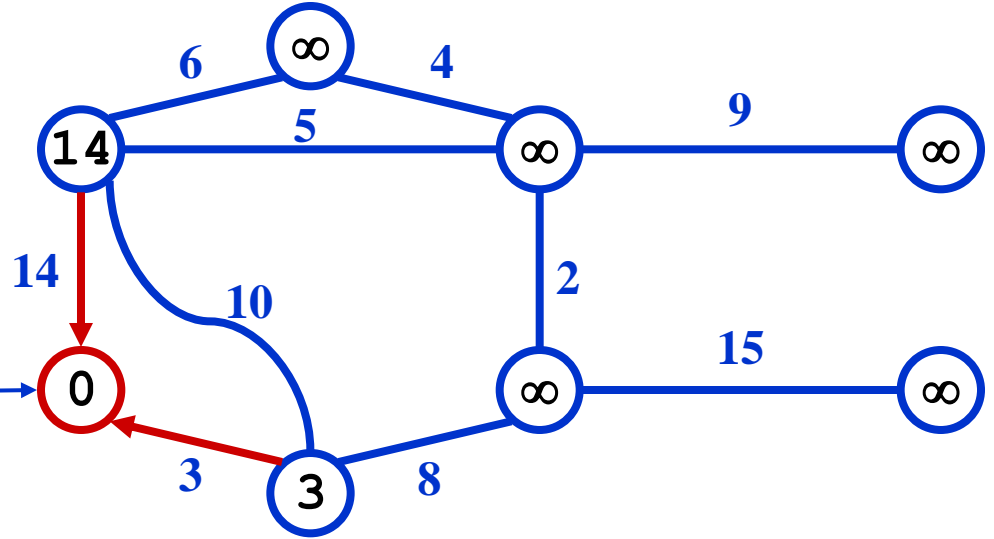
```
u = ExtractMin(Q);
```

```
for each  $v \in \text{Adjacent}[u]$ 
```

```
if (v ∈ Q and w(u,v) < weight[v])
```

p[v] = u;

```
weight[v] = w(u,v);
```

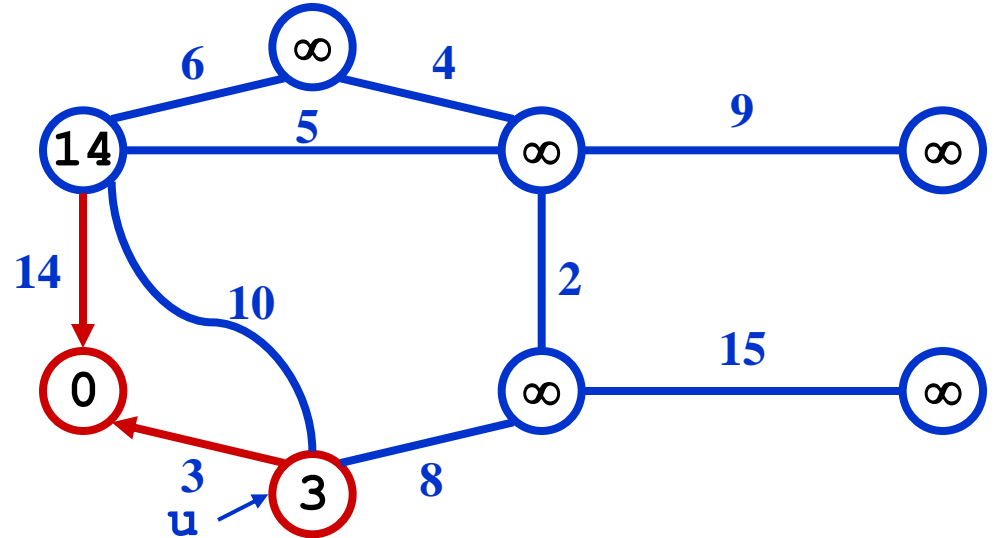


Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
     $\text{weight}[u] = \infty$ ;
 $Q = \text{all node}$ ;
 $\text{weight}[\text{root}] = 0$ ;
 $p[\text{root}] = \text{NULL}$ ;
while ( $Q$  not empty)
     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```

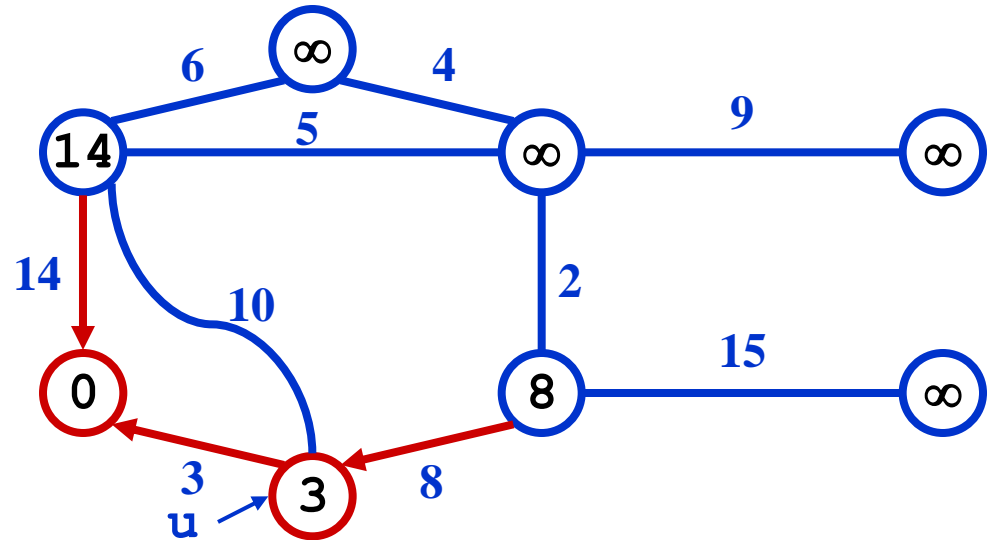


Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
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     $\text{weight}[u] = \infty$ ;
 $Q = \text{all node}$ ;
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 $p[\text{root}] = \text{NULL}$ ;
while ( $Q$  not empty)
     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```

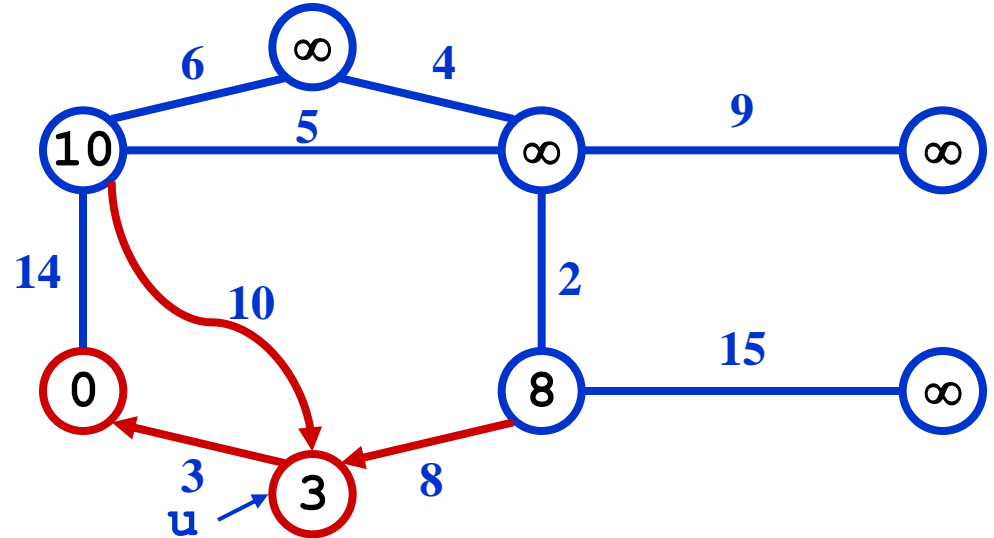


Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
     $\text{weight}[u] = \infty$ ;
 $Q = \text{all node}$ ;
 $\text{weight}[\text{root}] = 0$ ;
 $p[\text{root}] = \text{NULL}$ ;
while ( $Q$  not empty)
     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```



Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

for all node u

```
weight[u] =  $\infty$ ;
```

```
Q = all node;
```

```
weight[root] = 0;
```

```
p[root] = NULL;
```

```
while (Q not empty)
```

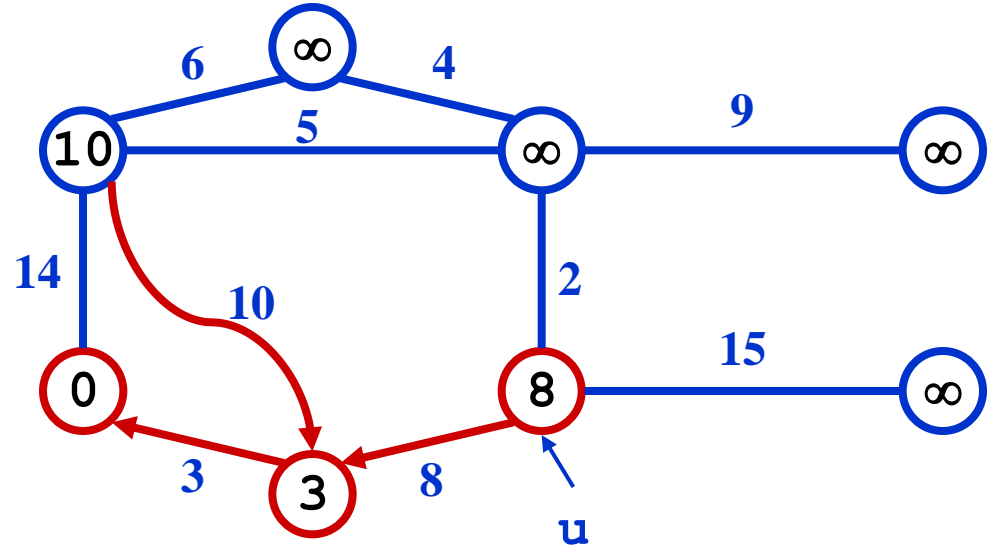
```
u = ExtractMin(Q);
```

```
for each  $v \in \text{Adjacent}[u]$ 
```

```
if (v ∈ Q and w(u, v) < weight[v])
```

p[v] = u;

```
weight[v] = w(u, v);
```

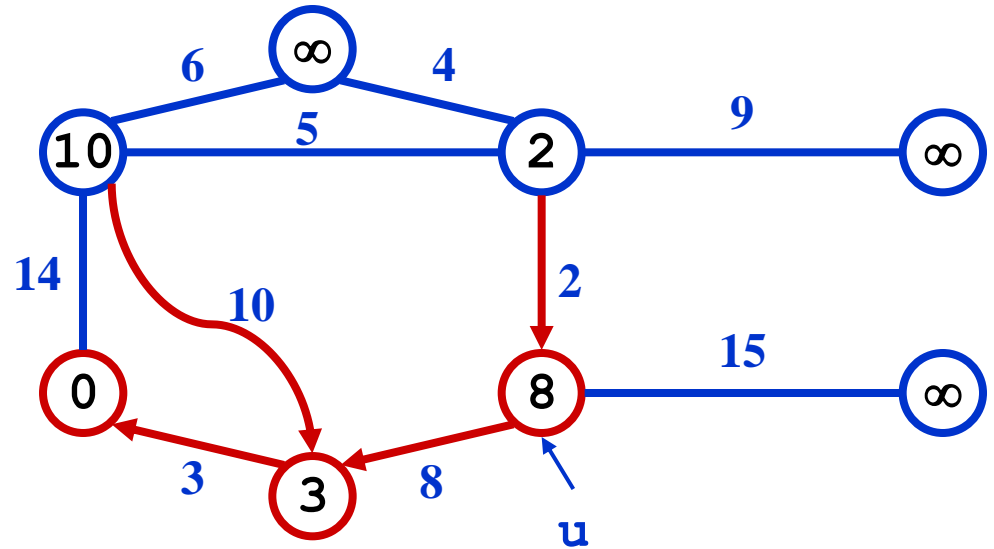


Prim's Algorithm

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MST-Prim

```
for all node  $u$ 
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 $p[\text{root}] = \text{NULL}$ ;
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     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
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```

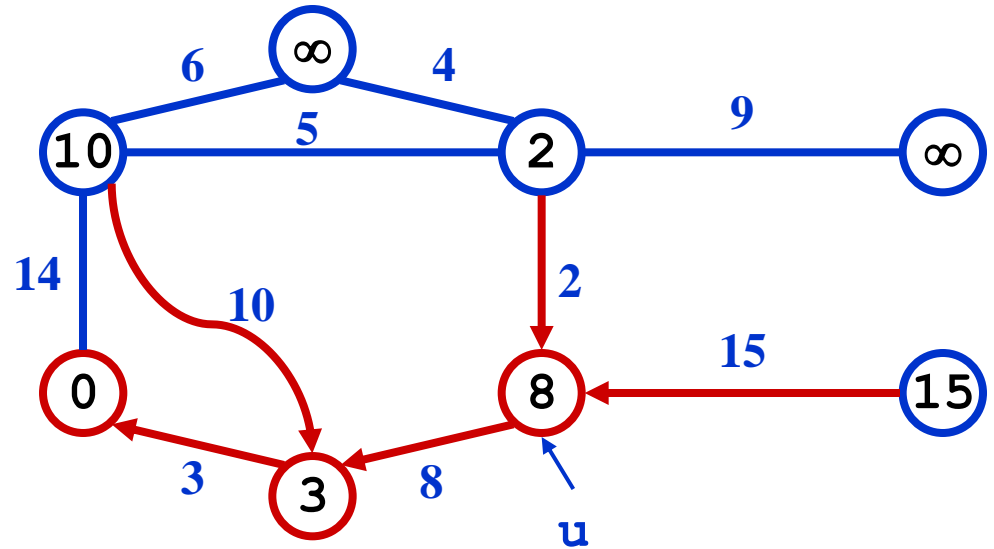


Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
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 $p[\text{root}] = \text{NULL}$ ;
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     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```

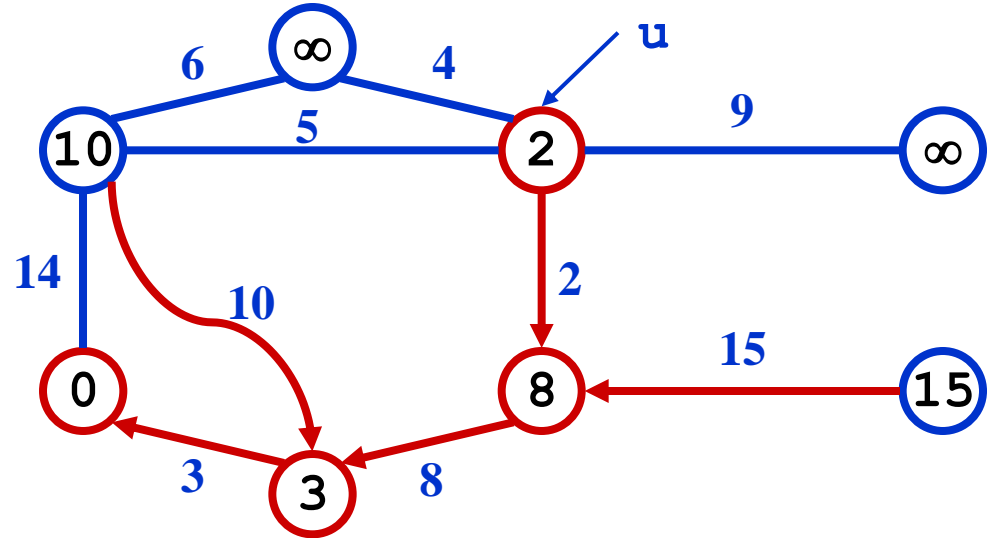


Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
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        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```

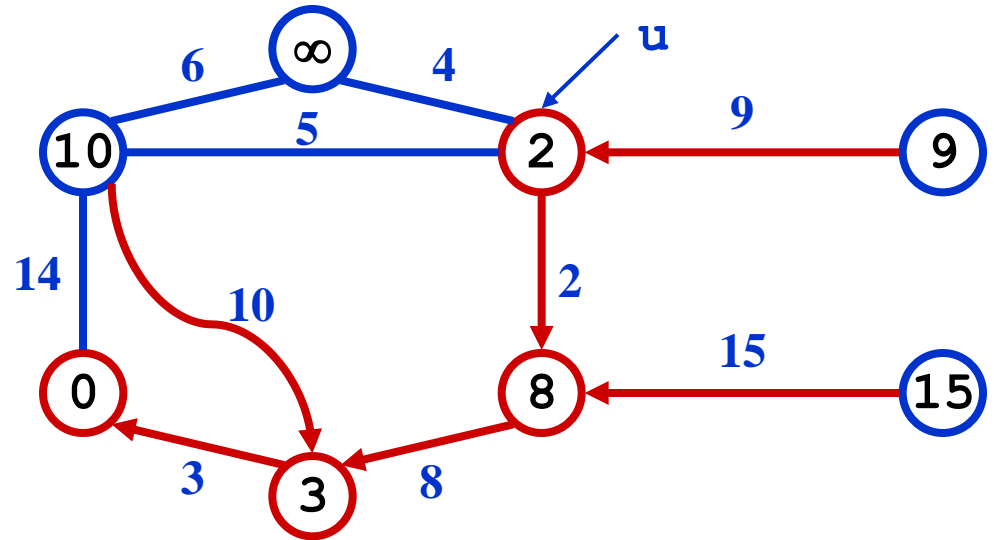


Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
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     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```

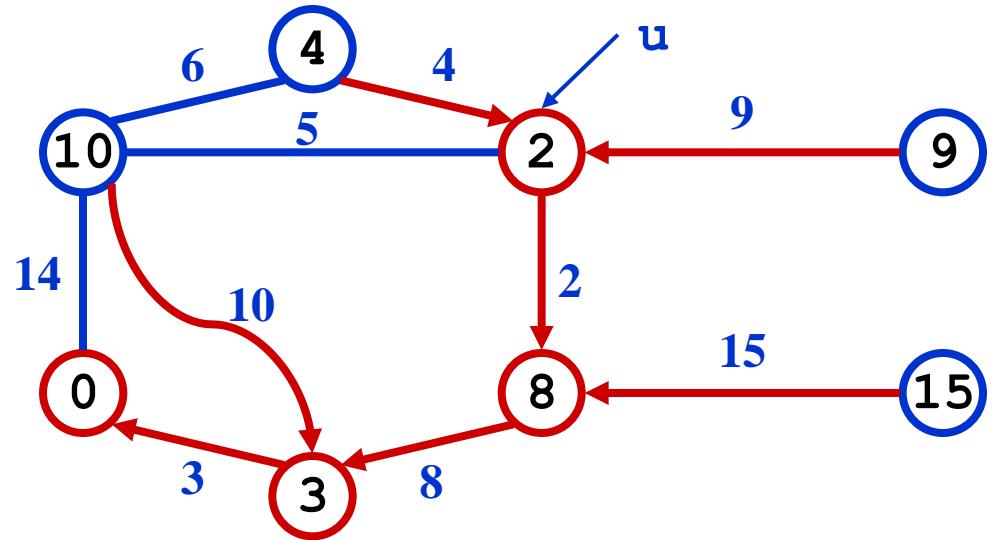


Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
     $\text{weight}[u] = \infty$ ;
 $Q = \text{all node}$ ;
 $\text{weight}[\text{root}] = 0$ ;
 $p[\text{root}] = \text{NULL}$ ;
while ( $Q$  not empty)
     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```

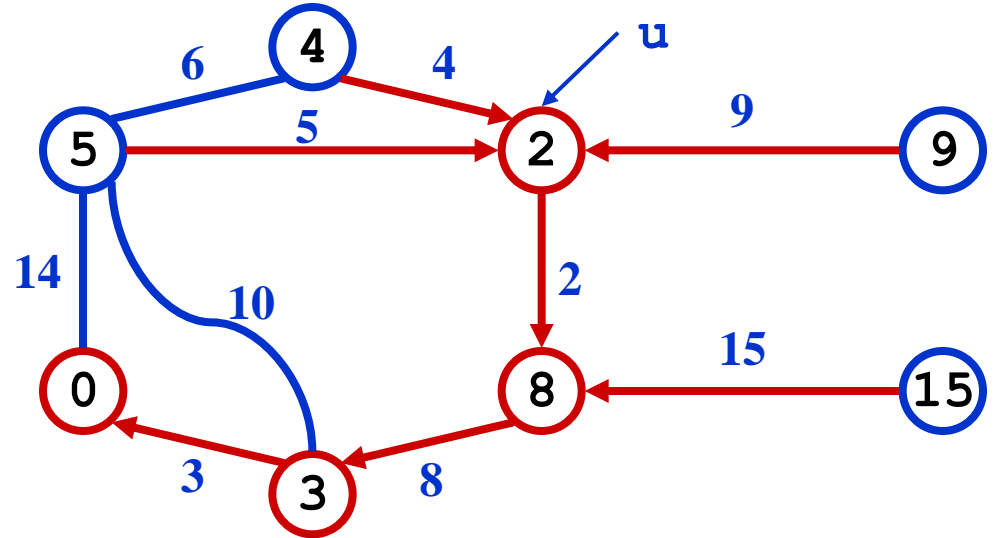


Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
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 $Q = \text{all node}$ ;
 $\text{weight}[\text{root}] = 0$ ;
 $p[\text{root}] = \text{NULL}$ ;
while ( $Q$  not empty)
     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```

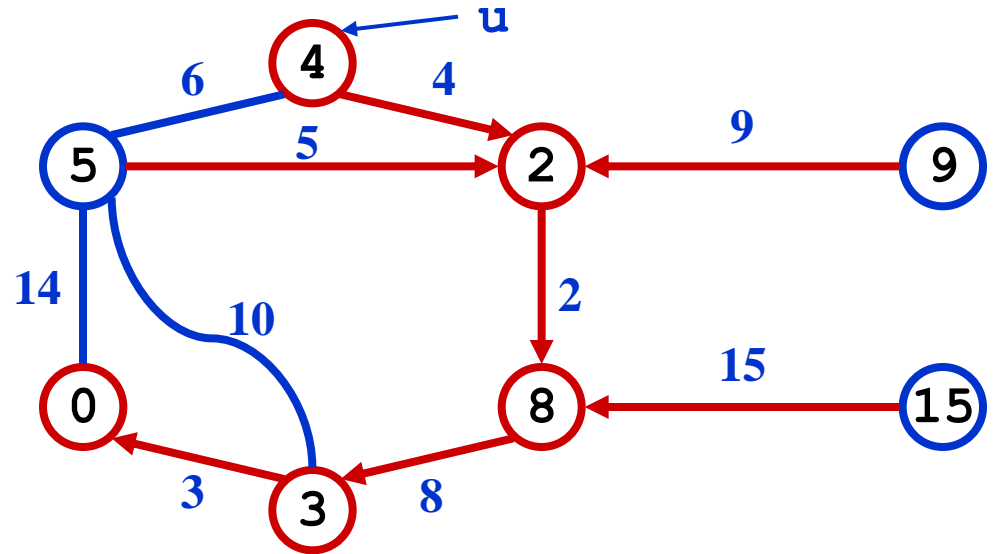


Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node u
    weight[u] =  $\infty$ ;
Q = all node;
weight[root] = 0;
p[root] = NULL;
while (Q not empty)
    u = ExtractMin(Q);
    for each v  $\in$  Adjacent[u]
        if (v  $\in$  Q and w(u,v) < weight[v])
            p[v] = u;
            weight[v] = w(u,v);
```

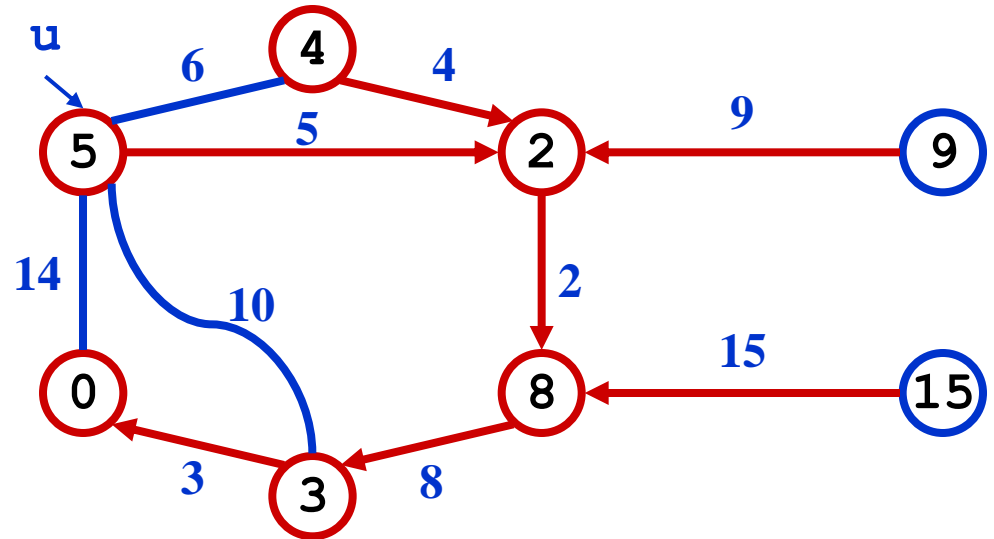


Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
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 $Q = \text{all node}$ ;
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while ( $Q$  not empty)
     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```

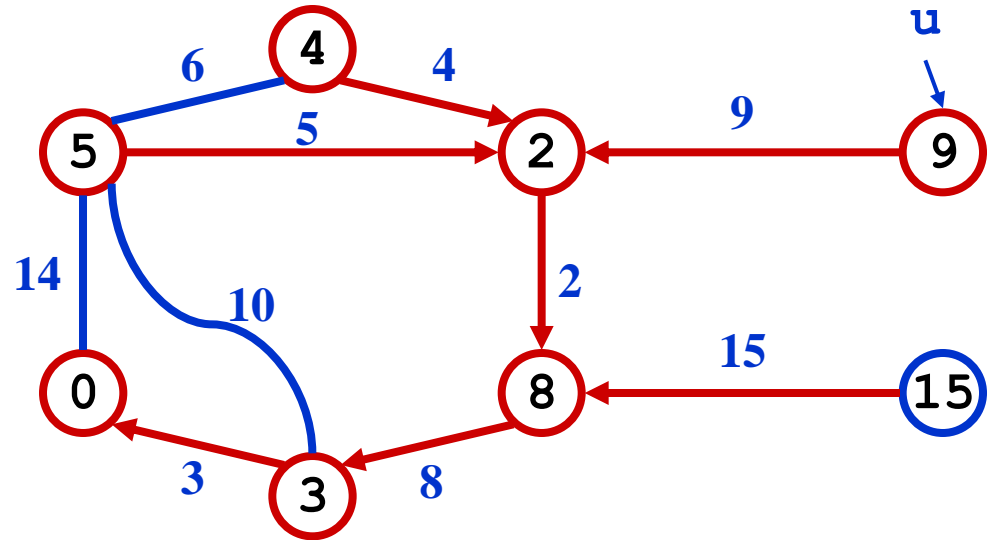


Prim's Algorithm

Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
     $\text{weight}[u] = \infty$ ;
 $Q = \text{all node}$ ;
 $\text{weight}[\text{root}] = 0$ ;
 $p[\text{root}] = \text{NULL}$ ;
while ( $Q$  not empty)
     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```

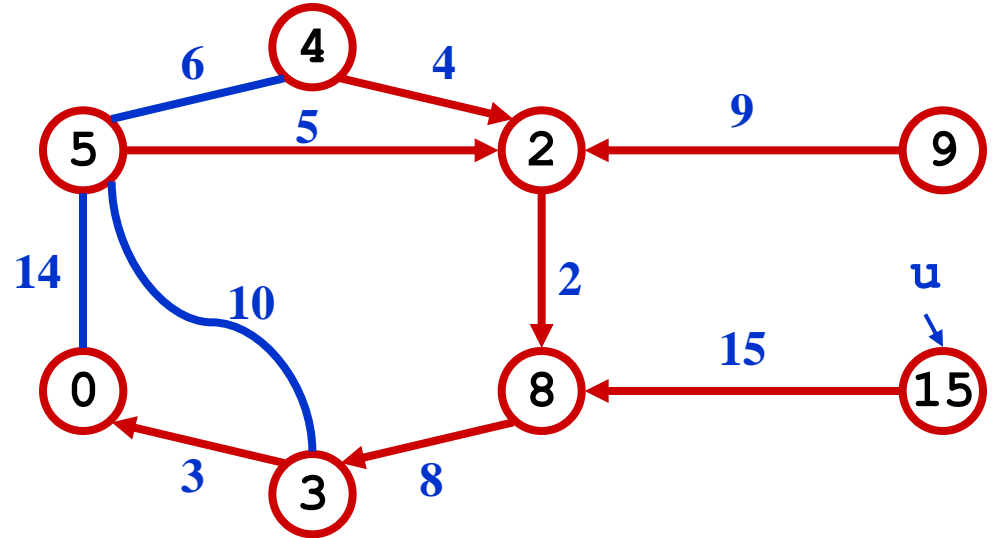


Prim's Algorithm

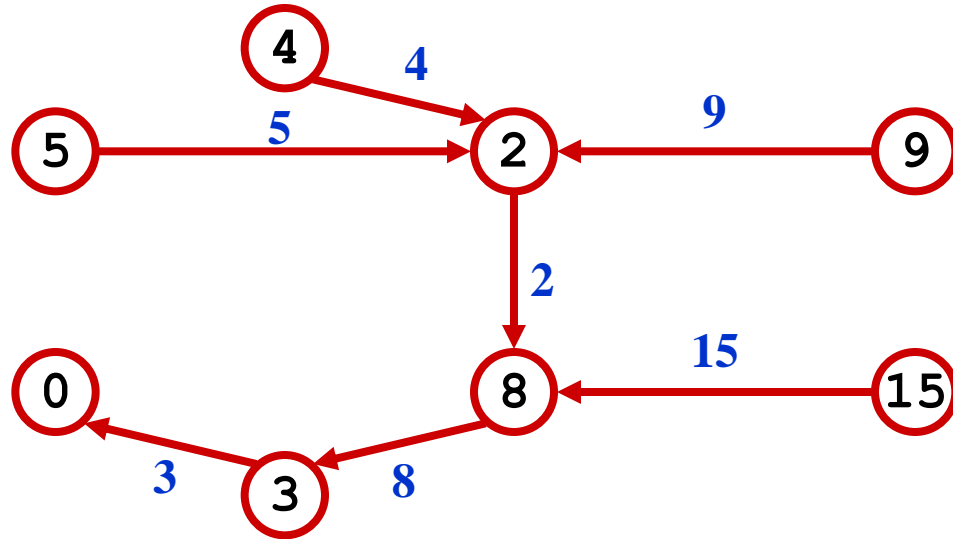
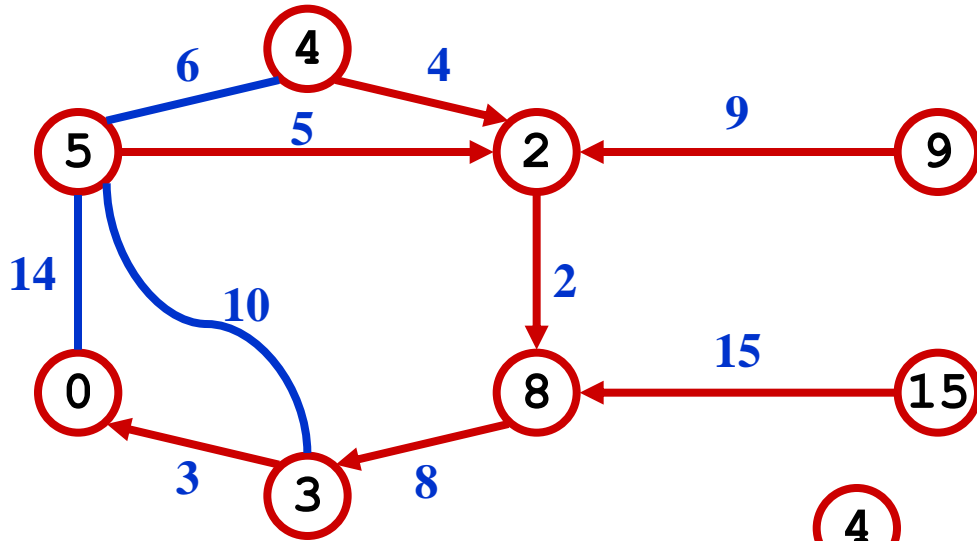
Q: heap, V: vertex set, E: Edge set, p: parent, w: edge weight

MST-Prim

```
for all node  $u$ 
     $\text{weight}[u] = \infty$ ;
 $Q = \text{all node}$ ;
 $\text{weight}[\text{root}] = 0$ ;
 $p[\text{root}] = \text{NULL}$ ;
while ( $Q$  not empty)
     $u = \text{ExtractMin}(Q)$ ;
    for each  $v \in \text{Adjacent}[u]$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u$ ;
             $\text{weight}[v] = w(u, v)$ ;
```



Prim's Algorithm



Cost of Prim's Algorithm

MST-Prim

for each node u $\leftarrow O(n)$

$\text{weight}[u] = \infty;$

$Q = \text{all node};$ $\leftarrow \text{BuildHeap} = O(n \log n)$

$\text{weight}[\text{root}] = 0;$

$p[\text{root}] = \text{NULL};$

while (Q not empty) $\leftarrow n \text{ times}$

$u = \text{ExtractMin}(Q);$ $\leftarrow \text{HeapExtractMin} = O(\log n)$

 for each $v \in \text{Adjacent}[u]$ $\leftarrow \text{Total for all nodes} = E \text{ times}$

 if ($v \in Q$ and $w(u, v) < \text{weight}[v]$)



$p[v] = u;$





$\text{weight}[v] = w(u, v);$ $\leftarrow \text{HeapDecreaseKey} = O(\log n)$

Worst Case Total for this part: $n \cdot (O(\log n) + E \cdot O(\log n)) = O(n \log n) + O(nE \log n) = O(nE \log n)$

Cost of Prim's Algorithm

MST-Prim

```
for each node  $u$    $O(n)$   
     $\text{weight}[u] = \infty;$   
 $Q = \text{all node};$    $\text{BuildHeap} = O(n \log n)$   
 $\text{weight}[\text{root}] = 0;$   
 $p[\text{root}] = \text{NULL};$   


$\text{while } (Q \text{ not empty})$    $n \text{ times}$   
     $u = \text{ExtractMin}(Q);$    $\text{HeapExtractMin} = O(\log n)$   
    for each  $v \in \text{Adjacent}[u]$    $\text{Total for all nodes} = E \text{ times}$   
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )  
             $p[v] = u;$   
             $\text{weight}[v] = w(u, v);$    $\text{HeapDecreaseKey} = O(\log n)$


```



But, amortized cost (aggregate method): How many total ExtractMin()?

Answer: n . How many total DecreaseKey()? **Answer:** E .

Cost of Prim's Algorithm

MST-Prim

```
for each node  $u$   $\leftarrow O(n)$ 
     $\text{weight}[u] = \infty;$ 
 $Q = \text{all node};$   $\leftarrow \text{BuildHeap} = O(n \log n)$ 
 $\text{weight}[\text{root}] = 0;$ 
 $p[\text{root}] = \text{NULL};$ 
while ( $Q$  not empty)  $\leftarrow n \text{ times}$ 
     $u = \text{ExtractMin}(Q);$   $\leftarrow \text{HeapExtractMin} = O(\log n)$ 
    for each  $v \in \text{Adjacent}[u]$   $\leftarrow \text{Total for all nodes} = E \text{ times}$ 
        if ( $v \in Q$  and  $w(u, v) < \text{weight}[v]$ )
             $p[v] = u;$ 
             $\text{weight}[v] = w(u, v);$   $\leftarrow \text{HeapDecreaseKey} = O(\log n)$ 
```

So, total cost: $\text{BuildHeap}() + n \text{ ExtractMin} + E \text{ DecreaseKey} = n \cdot O(\log n) + E \cdot O(n \log n) = O(n \log n) + O(E \log n) = \mathbf{O(E \log n)}$, because $E \geq n$