

# Tallinn University of Technology

DEPARTMENT OF SOFTWARE SCIENCE

SAMPLE EXAM PAPER 2019

**Advanced Programming**

**ITT8060**

Time allowed TWO hours 30 minutes

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**Answer ALL FOUR questions**

No calculators, mobile phones or other electronic devices capable of storing or retrieving text may be used.

A print text book  
(Real World Functional Programming, Functional Programming Using F#, or Expert F#)  
is allowed.

**DO NOT open the examination paper until instructed to do so**

Name:	
Student ID:	
Marks (to be filled by teaching staff):	

Please circle **A**, **B**, **C** or **D** according to which of them best matches the answer. In case there are multiple correct answers you should choose the best one. Only a single circle is considered to be the correct answer. In case you make a mistake, cross out the answer and write clearly next to the question what the answer is.

- a. The expression `List.filter (fun (x,y) -> x>y) [1,2;2,3;3,4;5,0]` returns  
 A. `[5,0]` B. `[1;2;3;5]` C. `[2;3;4;0]` D. `[1,2;2,3;3,4;5,0]`
- b. The value of `0 |> List.fold (>>) id [(+) 1; (*) 3; (-) 1]` is  
 A. 2 B. -11 C. -2 D. None of the above
- (Given the type of the function `(|>)` is `('a -> ('a -> 'b) -> 'b)` and the type of the function `(>>)` is `((('a -> 'b) -> ('b -> 'c) -> 'a -> 'c))`)
- c. Evaluating the expression `(lazy (1,2))` will return  
 A. type error B. 2 C. 1 D. None of the above
- d. Evaluating the expression `Option.bind Some (None: bool option)` returns,  
 given the type of `Option.bind` is `((('a -> 'b option) -> 'a option -> 'b option)`  
 A. `Some false` B. `false` C. `None` D. `Some Some false`
- e. The type of the expression  
`printfn "nice day!"`  
 is  
 A. `string*int` B. `string` C. `unit` D. `int`
- f. A function of type `'a list -> 'int list -> 'a` can be applied given the first argument is of type  
 A. `int list` B. `string list list` C. all of the above D. none of the above
- g. The type of the following function is

```
let rec c b = match b with
| []      -> []
| d :: e -> (c [(List.head e)])
```

- A. `'a list -> unit` B. `'a list -> 'b list` C. `'a list -> 'a list` D. `int`
- h. The type of the expression `let f = fun x -> (fun y -> x+y)` matches the type of the expression  
 A. `let g x y = x * y` B. `let g x,y = x + y` C. Both of the above D. None of the above

i. Evaluation of the following code will result in

```
let sleepWorkflow = async {
  printfn "starting"
  do! Async.Sleep 2000
  printfn "finished"
}
```

- A. A value of type `Async<unit>` being created;      B. “starting” and “finishing” being printed;      C. A delay of 2000 ms;      D. None of the above.

j. The function `f` defined below

```
let rec f x =
  seq{
    yield x
    yield! f (x)
  }
```

- A. has type `'a -> 'b list.`      B. has type `int -> int list.`      C. has type `'a -> seq<'a>.`      D. None of the above.

k. Given that the type of `List.reduce` is `(('a -> 'a -> 'a) -> 'a list -> 'a)`, the expression `List.reduce (*) [1]` will evaluate to

- A. 1      B. [1]      C. runtime error;      D. none of the above.

l. Given that the type of `List.collect` is `(('a -> 'b list) -> 'a list -> 'b list)`, the expression `List.collect (fun x -> [x + x]) [2; 3; 5]` will evaluate to

- A. `[[4]; [9]; [25]]`      B. `[30]`      C. `[4; 9; 25]`      D. None of the above.

m. Given the declaration

```
let rec g x y =
  match y with
  | [] -> 1
  | h :: t -> x h + g x t
```

- A. The type of `g` is `(int -> int) -> int list -> int`      B. `g` is tail recursive      C. All of the above      D. None of the above

## Question 2: Partial functions

A function of type `'a -> 'b option` can be thought of as a function of type `'a -> 'b` that happens to be partial. This means that there may be values of type `'a` where this function is undefined (we cannot produce a value of type `'b`). The type `'a -> 'b option` precisely says that given an `'a` this function may produce a value of type `'b` (failure to produce an output is represented by `None`).

There are library functions `map` and `bind` for `Option` with the given types:

`Option.map : ('a -> 'b) -> 'a option -> 'b option`

`Option.bind : ('a -> 'b option) -> 'a option -> 'b option`

`Option.map` can be used to apply an ordinary function `('a -> 'b)` to an optional value to get an optional result.

`Option.bind` can be used to apply a partial function `('a -> 'b option)` to an optional value to get an optional result.

- a. Implement a function `apply` that allows to apply a partial function to an optional value to get an optional result:

`apply : ('a -> 'b) option -> 'a option -> 'b option`

For example, the result of `apply (Some id) (Some 3)` must be `Some 3`.

(5 points)

(i) Evaluate `apply (Some ((+) 1)) None` (2 points)

(ii) Evaluate `apply (Some List.head) Some [3;2;1]` (2 points)

- b. Implement the function

`sequence : 'a option list -> 'a list option`

so that given a list `xs` of optional values it evaluates to

- `Some xs'` when all of the optional values in `xs` were `Some` and `xs'` is a list of the values inside
- `None` when there was at least one `None` in `xs`.

Evaluating `sequence [Some 1; Some 2]` should evaluate to `Some [1; 2]`.

Evaluating `sequence [Some 1; None; Some 2]` should evaluate to `None`.

(7 points)

- c. Implement the function

`sequence' : 'a option list -> 'a list option`

that behaves according to the specification of `sequence` but is implemented tail recursively.

(7 points)

- d. Give the type of the value that results evaluating the expression:

`[Some [1;2;3]; Some [3;4] ; None] |> sequence`

(2 points)

### Question 3: Expression trees

Given the following type definitions:

```
type VName = string
```

```
type Expr = Var   of VName
          | Neg   of Expr
          | And   of (Expr * Expr)
          | Or    of (Expr * Expr)
```

```
type Assignment = (VName * bool) list
```

And the definition of the function `lookup : v:VName -> vs:Assignment -> bool` that looks up a bool value corresponding to the particular variable name `VName`:

- a. Define the function `interpret : e:Expr -> vs:Assignment -> bool` that will determine if an expression `e` evaluates to true or false given the assignments to the variables in `vs`. `Neg` represents negation (`not`), `And` represents conjunction (`&&`) and `Or` represents disjunction (`||`).

For example,

```
interpret (And (Or (Var "X1", Var "X2"), Neg (Var "X2"))) ["X1",true; "X2",false]
```

should evaluate to `true`.

(5 points)

- b. Define a function

```
variables : e:Expr -> VName list
```

that returns all distinct names of the Boolean variables used in the expression. The results should contain distinct values, i.e. each name at most once.

Hint: consider using `List.sort : ('a list -> 'a list)` when `'a : comparison` and/or

`List.groupBy : (('a -> 'b) -> 'a list -> ('b * 'a list) list)` when `'b : equality` library functions.

For example, `variables (And (Or (Var "X1", Var "X2"), Neg (Var "X2")))` should evaluate to `["X1"; "X2"]`.

(7 points)

- c. Consider the following type of arithmetic expressions:

```
type Expr =
  | Const of int
  | Sum   of Expr * Expr
  | Diff  of Expr * Expr
  | Prod  of Expr * Expr
```

A value `Const n` represents a primitive expression that denotes the number `n`, while values `Sum e1 e2`, `Diff e1 e2`, and `Prod e1 e2`, represent sums, differences, products, respectively.

Write a function `eval : Expr -> int` that evaluates the given expression.

(3 points)

- d. Write a function `commute : Expr -> Expr` that swaps the arguments of all sums and products in the given expression (which should not change the result of the expression).

(4 points)

- e. Polish notation is a notation for expressions where the operators come before their arguments and no parentheses are used. For example, the expressions `3 + 5`, `3 + (4 * 5)`, and `(3 + 4) * 5` are written `+ 3 5`, `+ 3 * 4 5`, and `* + 3 4 5`, respectively, when using Polish notation.

Write a function `pn : Expr -> string` that turns the given expression into its representation in Polish notation.

(6 points)

## Question 4: Euler method for numeric approximation

The Euler method can be used to numerically approximate solutions to first-order ordinary differential equations (ODEs) with a given initial value. The ODE has to be provided in the following form:

$$dy(t)/dt = f(t, y(t))$$

with an initial value  $y(t_0) = y_0$ . This can be numerically approximated with a formula

$$y_{n+1} = y_n + hf(t_n, y_n)$$

An F# implementation of the function is given as `eulerStep` below. We can use the implementation to approximate the Newton's law of cooling where the cooling constant is 0.05 and target temperature is 22.0° C.

The function `newtonNext1` will compute the next pair of `t` and `y` where the returned `y` element in the pair represents the temperature one second later (the time instance is represented by `t`) and target temperature being 22.0° C.

```
let eulerStep f (h : float) t y =
    ((t + h), (y + h * (f t y)))
```

```
let newton t y = -0.05 * (y - 22.0)
```

```
let newtonNext1 p =
    let (t0,y0) = p
    eulerStep newton 1.0 t0 y0
```

- a. Write a function `euler : (float * float) -> (float*float) seq` that for any given pair `(t,y)` computes the infinite sequence consisting of the pairs `(t,y)`, `newtonNext1 (t,y)`, `newtonNext1 (newtonNext1 (t,y))`, and so on. Use sequence expressions in your implementation.

For initial values, where  $y > 22.0$ , the sequence will represent the approximation of Newton's law of cooling of an object down to 22.0° C.

(4 points)

- b. Write a function `euler' : (float*float) -> (float*float) seq` that for any given pair `(t,y)` computes the same sequence as `euler (t,y)`. Use the `Seq.unfold` function in your implementation. The type of `Seq.unfold` is `(('a -> ('b * 'a) option) -> 'a -> seq<'b>)`.

(4 points)

- c. The cooling curve is interesting until a small margin  $\epsilon$  from the target temperature. We want to get the part of the cooling approximation sequence, where the value of `y` is more than  $\epsilon$  larger than the target value (22.0 in this case).

Write a function `coolingApprox : float -> (float * float) seq -> (float * float) list` that turns a given sequence into its prefix that ends as soon as  $y \leq 22.0 + \epsilon$ . The  $\epsilon$  is given as the first argument to the `coolingApprox` function. Hint: access the sequence step-by-step e.g. by accessing `Seq.head` and `Seq.tail` and comparing the `y` value to the margin.

(7 points)

- d. Write a function `coolingApprox' : float -> (float * float) seq -> (float * float) list` that for any given `epsilon` and sequence `s` computes the same list as `coolingApprox epsilon s`. Make sure that all recursive calls in the definition of `coolingApprox'` are tail calls.

(10 points)