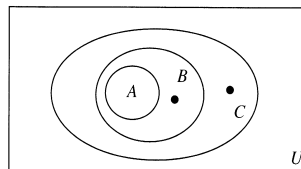


19.  $\forall x \exists y \exists z (y \neq z \wedge \forall w (P(w, x) \leftrightarrow (w = y \vee w = z)))$   
 21. **a)**  $\neg \exists x P(x)$  **b)**  $\exists x (P(x) \wedge \forall y (P(y) \rightarrow y = x))$   
**c)**  $\exists x_1 \exists x_2 (P(x_1) \wedge P(x_2) \wedge x_1 \neq x_2 \wedge \forall y (P(y) \rightarrow (y = x_1 \vee y = x_2)))$  **d)**  $\exists x_1 \exists x_2 \exists x_3 (P(x_1) \wedge P(x_2) \wedge P(x_3) \wedge x_1 \neq x_2 \wedge x_1 \neq x_3 \wedge x_2 \neq x_3 \wedge \forall y (P(y) \rightarrow (y = x_1 \vee y = x_2 \vee y = x_3)))$  **23.** Suppose that  $\exists x (P(x) \rightarrow Q(x))$  is true. Then either  $Q(x_0)$  is true for some  $x_0$ , in which case  $\forall x P(x) \rightarrow \exists x Q(x)$  is true; or  $P(x_0)$  is false for some  $x_0$ , in which case  $\forall x P(x) \rightarrow \exists x Q(x)$  is true. Conversely, suppose that  $\exists x (P(x) \rightarrow Q(x))$  is false. That means that  $\forall x (P(x) \wedge \neg Q(x))$  is true, which implies  $\forall x P(x)$  and  $\forall x (\neg Q(x))$ . This latter proposition is equivalent to  $\neg \exists x Q(x)$ . Thus,  $\forall x P(x) \rightarrow \exists x Q(x)$  is false. **25.** No **27.**  $\forall x \forall y \exists z T(x, y, z)$ , where  $T(x, y, z)$  is the statement that student  $x$  has taken class  $y$  in department  $z$ , where the domains are the set of students in the class, the set of courses at this university, and the set of departments in the school of mathematical sciences **29.**  $\exists! x \exists! y T(x, y)$  and  $\exists x \forall z ((\exists y \forall w (T(z, w) \leftrightarrow w = y)) \leftrightarrow z = x)$ , where  $T(x, y)$  means that student  $x$  has taken class  $y$  and the domain is all students in this class **31.**  $P(a) \rightarrow Q(a)$  and  $Q(a) \rightarrow R(a)$  by universal instantiation; then  $\neg Q(a)$  by modus tollens and  $\neg P(a)$  by modus tollens **33.** We give a proof by contradiction and show that if  $\sqrt{x}$  is rational, then  $x$  is rational, assuming throughout that  $x \geq 0$ . Suppose that  $\sqrt{x} = p/q$  is rational,  $q \neq 0$ . Then  $x = (\sqrt{x})^2 = p^2/q^2$  is also rational ( $q^2$  is again nonzero). **35.** We can give a constructive proof by letting  $m = 10^{500} + 1$ . Then  $m^2 = (10^{500} + 1)^2 > (10^{500})^2 = 10^{1000}$ . **37.** 23 cannot be written as the sum of eight cubes. **39.** 223 cannot be written as the sum of 36 fifth powers.



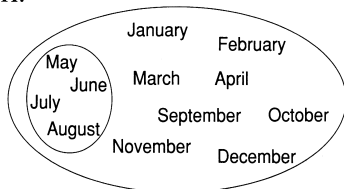
**15.** Suppose that  $x \in A$ . Because  $A \subseteq B$ , this implies that  $x \in B$ . Because  $B \subseteq C$ , we see that  $x \in C$ . Because  $x \in A$  implies that  $x \in C$ , it follows that  $A \subseteq C$ . **17. a)** 1 **b)** 1 **c)** 2 **d)** 3 **19. a)**  $\{\emptyset, \{a\}\}$  **b)**  $\{\emptyset, \{a\}, \{b\}, \{a, b\}\}$  **c)**  $\{\emptyset, \{a\}, \{\{a\}\}, \{\emptyset, \{a\}\}\}$  **21. a)** 8 **b)** 16 **c)** 2 **23. a)**  $\{(a, y), (b, y), (c, y), (d, y), (a, z), (b, z), (c, z), (d, z)\}$  **b)**  $\{(y, a), (y, b), (y, c), (y, d), (z, a), (z, b), (z, c), (z, d)\}$  **25.** The set of triples  $(a, b, c)$ , where  $a$  is an airline and  $b$  and  $c$  are cities **27.**  $\emptyset \times A = \{(x, y) \mid x \in \emptyset \text{ and } y \in A\} = \emptyset = \{(x, y) \mid x \in A \text{ and } y \in \emptyset\} = A \times \emptyset$  **29. mn** **31.** The elements of  $A \times B \times C$  consist of 3-tuples  $(a, b, c)$ , where  $a \in A$ ,  $b \in B$ , and  $c \in C$ , whereas the elements of  $(A \times B) \times C$  look like  $((a, b), c)$ —ordered pairs, the first coordinate of which is again an ordered pair. **33. a)** The square of a real number is never  $-1$ . True **b)** There exists an integer whose square is 2. False **c)** The square of every integer is positive. False **d)** There is a real number equal to its own square. True **35. a)**  $\{-1, 0, 1\}$  **b)**  $\mathbb{Z} - \{0, 1\}$  **c)**  $\emptyset$  **37.** We must show that  $\{\{a\}, \{a, b\}\} = \{\{c\}, \{c, d\}\}$  if and only if  $a = c$  and  $b = d$ . The “if” part is immediate. So assume these two sets are equal. First, consider the case when  $a \neq b$ . Then  $\{\{a\}, \{a, b\}\}$  contains exactly two elements, one of which contains one element. Thus,  $\{\{c\}, \{c, d\}\}$  must have the same property, so  $c \neq d$  and  $\{c\}$  is the element containing exactly one element. Hence,  $\{a\} = \{c\}$ , which implies that  $a = c$ . Also, the two-element sets  $\{a, b\}$  and  $\{c, d\}$  must be equal. Because  $a = c$  and  $a \neq b$ , it follows that  $b = d$ . Second, suppose that  $a = b$ . Then  $\{\{a\}, \{a, b\}\} = \{\{a\}\}$ , a set with one element. Hence,  $\{\{c\}, \{c, d\}\}$  has only one element, which can happen only when  $c = d$ , and the set is  $\{\{c\}\}$ . It then follows that  $a = c$  and  $b = d$ . **39.** Let  $S = \{a_1, a_2, \dots, a_n\}$ . Represent each subset of  $S$  with a bit string of length  $n$ , where the  $i$ th bit is 1 if and only if  $a_i \in S$ . To generate all subsets of  $S$ , list all  $2^n$  bit strings of length  $n$  (for instance, in increasing order), and write down the corresponding subsets.

## CHAPTER 2

### Section 2.1

1. **a)**  $\{-1, 1\}$  **b)**  $\{1, 2, 3, 4, 5, 6, 7, 8, 9, 10, 11\}$  **c)**  $\{0, 1, 4, 9, 16, 25, 36, 49, 64, 81\}$  **d)**  $\emptyset$  **3. a)** Yes **b)** No **c)** No **5. a)** Yes **b)** No **c)** Yes **d)** No **e)** No **f)** No **7. a)** False **b)** False **c)** False **d)** True **e)** False **f)** False **g)** True **9. a)** True **b)** True **c)** False **d)** True **e)** True **f)** False

11.



**13.** The dots in certain regions indicate that those regions are not empty.

### Section 2.2

- 1. a)** The set of students who live within one mile of school and who walk to classes **b)** The set of students who live within one mile of school or who walk to classes (or who do both) **c)** The set of students who live within one mile of school but do not walk to classes **d)** The set of students who walk to classes but live more than one mile away from school **3. a)**  $\{0, 1, 2, 3, 4, 5, 6\}$  **b)**  $\{3\}$  **c)**  $\{1, 2, 4, 5\}$  **d)**  $\{0, 6\}$  **5.  $\bar{A} = \{x \mid \neg(x \in A)\} = \{x \mid \neg(\neg(x \in A))\} = \{x \mid x \in A\} = A$**  **7. a)**  $A \cup U = \{x \mid x \in A \vee x \in U\} = \{x \mid x \in A \vee \mathbf{T}\} = \{x \mid \mathbf{T}\} = U$  **b)**  $A \cap \emptyset = \{x \mid x \in A \wedge x \in \emptyset\} = \{x \mid x \in$

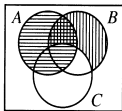
$A \cap F = \{x \mid F\} = \emptyset$  9. a)  $A \cup \bar{A} = \{x \mid x \in A \vee x \notin A\} = U$  b)  $A \cap \bar{A} = \{x \mid x \in A \wedge x \notin A\} = \emptyset$  11. a)  $A \cup B = \{x \mid x \in A \vee x \in B\} = \{x \mid x \in B \vee x \in A\} = B \cup A$  b)  $A \cap B = \{x \mid x \in A \wedge x \in B\} = \{x \mid x \in B \wedge x \in A\} = B \cap A$  13. Suppose  $x \in A \cap (A \cup B)$ . Then  $x \in A$  and  $x \in A \cup B$  by the definition of intersection. Because  $x \in A$ , we have proved that the left-hand side is a subset of the right-hand side. Conversely, let  $x \in A$ . Then by the definition of union,  $x \in A \cup B$  as well. Therefore  $x \in A \cap (A \cup B)$  by the definition of intersection, so the right-hand side is a subset of the left-hand side. 15. a)  $x \in (A \cup B) \equiv x \notin (A \cup B) \equiv \neg(x \in A \vee x \in B) \equiv \neg(x \in A) \wedge \neg(x \in B) \equiv x \notin A \wedge x \notin B \equiv x \in \bar{A} \wedge x \in \bar{B} \equiv x \in \bar{A} \cap \bar{B}$

$A$	$B$	$A \cup B$	$\overline{(A \cup B)}$	$\bar{A}$	$\bar{B}$	$\bar{A} \cap \bar{B}$
1	1	1	0	0	0	0
1	0	1	0	0	1	0
0	1	1	0	1	0	0
0	0	0	1	1	1	1

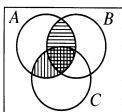
17. a)  $x \in A \cap B \cap \bar{C} \equiv x \notin A \cap B \cap C \equiv x \notin A \vee x \notin B \vee x \notin \bar{C} \equiv x \in \bar{A} \vee x \in \bar{B} \vee x \in C \equiv x \in \bar{A} \cup \bar{B} \cup C$

$A$	$B$	$C$	$A \cap B \cap C$	$\overline{(A \cap B \cap C)}$	$\bar{A}$	$\bar{B}$	$\bar{C}$	$\bar{A} \cup \bar{B} \cup \bar{C}$
1	1	1	1	0	0	0	0	0
1	1	0	0	1	0	0	1	1
1	0	1	0	1	0	1	0	1
1	0	0	0	1	0	1	1	1
0	1	1	0	1	1	0	0	1
0	1	0	0	1	1	0	1	1
0	0	1	0	1	1	1	0	1
0	0	0	0	1	1	1	1	1

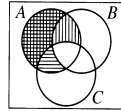
19. Both sides equal  $\{x \mid x \in A \wedge x \notin B\}$ . 21.  $x \in A \cup (B \cup C) \equiv (x \in A) \vee (x \in (B \cup C)) \equiv (x \in A) \vee (x \in B \vee x \in C) \equiv (x \in A \vee x \in B) \vee (x \in C) \equiv x \in (A \cup B) \cup C$  23.  $x \in A \cup (B \cap C) \equiv (x \in A) \vee (x \in (B \cap C)) \equiv (x \in A) \vee (x \in B \wedge x \in C) \equiv (x \in A \vee x \in B) \wedge (x \in A \vee x \in C) \equiv x \in (A \cup B) \cap (A \cup C)$  25. a)  $\{4, 6\}$  b)  $\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, 10\}$  c)  $\{4, 5, 6, 8, 10\}$  d)  $\{0, 2, 4, 5, 6, 7, 8, 9, 10\}$  27. a) The double-shaded portion is the desired set.



b) The desired set is the entire shaded portion.



c) The desired set is the entire shaded portion.



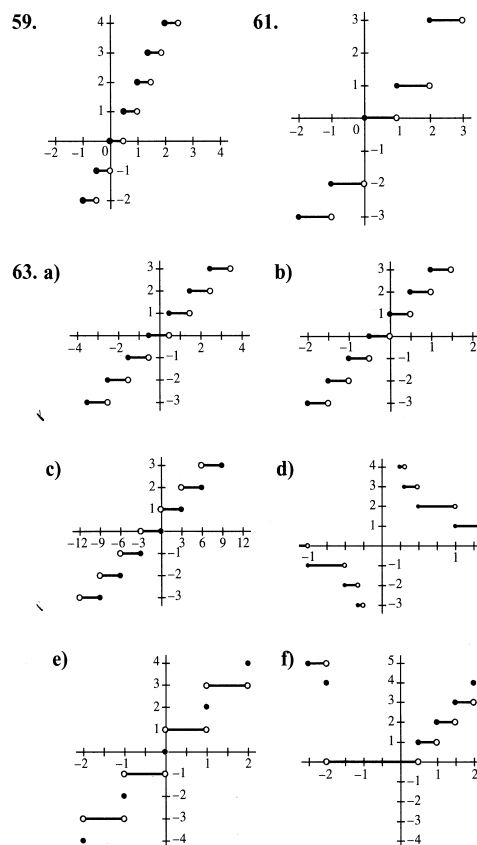
29. a)  $B \subseteq A$  b)  $A \subseteq B$  c)  $A \cap B = \emptyset$  d) Nothing, because this is always true e)  $A = B$  31.  $A \subseteq B \equiv \forall x(x \in A \rightarrow x \in B) \equiv \forall x(x \notin B \rightarrow x \notin A) \equiv \forall x(x \in \bar{B} \rightarrow x \in \bar{A}) \equiv \bar{B} \subseteq \bar{A}$  33. The set of students who are computer science majors but not mathematics majors or who are mathematics majors but not computer science majors 35. An element is in  $(A \cup B) - (A \cap B)$  if it is in the union of  $A$  and  $B$  but not in the intersection of  $A$  and  $B$ , which means that it is in either  $A$  or  $B$  but not in both  $A$  and  $B$ . This is exactly what it means for an element to belong to  $A \oplus B$ . 37. a)  $A \oplus A = (A - A) \cup (A - A) = \emptyset \cup \emptyset = \emptyset$  b)  $A \oplus \emptyset = (A - \emptyset) \cup (\emptyset - A) = A \cup \emptyset = A$  c)  $A \oplus U = (A - U) \cup (U - A) = \emptyset \cup \bar{A} = \bar{A}$  d)  $A \oplus \bar{A} = (A - \bar{A}) \cup (\bar{A} - A) = A \cup \bar{A} = U$  39.  $B = \emptyset$  41. Yes 43. Yes 45. a)  $\{1, 2, 3, \dots, n\}$  b)  $\{1\}$  47. a)  $A_n$  b)  $\{0, 1\}$  49. a)  $\mathbb{Z}, \{-1, 0, 1\}$  b)  $\mathbb{Z} - \{0\}, \emptyset$  c)  $\mathbb{R}, [-1, 1]$  d)  $[1, \infty), \emptyset$  51. a)  $\{1, 2, 3, 4, 7, 8, 9, 10\}$  b)  $\{2, 4, 5, 6, 7\}$  c)  $\{1, 10\}$  53. The bit in the  $i$ th position of the bit string of the difference of two sets is 1 if the  $i$ th bit of the first string is 1 and the  $i$ th bit of the second string is 0, and is 0 otherwise. 55. a) 11 1110 0000 0000 0000 0000 0000  $\vee$  01 1100 1000 0000 0100 0101 0000 = 11 1110 1000 0000 0100 0101 0000, representing  $\{a, b, c, d, e, g, p, t, v\}$  b) 11 1110 0000 0000 0000 0000 0000  $\wedge$  01 1100 1000 0000 0100 0101 0000 = 01 1100 0000 0000 0000 0000 0000, representing  $\{b, c, d\}$  c) (11 1110 0000 0000 0000 0000  $\vee$  00 0110 0110 0001 1000 0110 0110)  $\wedge$  (01 1100 1000 0000 0100 0101 0000  $\vee$  00 1010 0010 0000 1000 0010 0111) = 11 1110 0110 0001 1000 0110 0110  $\wedge$  01 1110 1010 0000 1100 0111 0111 = 01 1110 0010 0000 1000 0110 0110, representing  $\{b, c, d, e, i, o, t, u, x, y\}$  d) 11 1110 0000 0000 0000 0000 0000  $\vee$  01 1100 1000 0000 0100 0101 0000  $\vee$  00 1010 0010 0000 1000 0010 0111  $\vee$  00 0110 0110 0001 1000 0110 0110 = 11 1110 1110 0001 1100 0111 0111, representing  $\{a, b, c, d, e, g, h, i, n, o, p, t, u, v, x, y, z\}$  57. a)  $\{1, 2, 3, \{1, 2, 3\}\}$  b)  $\{\emptyset\}$  c)  $\{\emptyset, \{\emptyset\}\}$  d)  $\{\emptyset, \{\emptyset\}\}$  59. a)  $\{3 \cdot a, 3 \cdot b, 1 \cdot c, 4 \cdot d\}$  b)  $\{2 \cdot a, 2 \cdot b\}$  c)  $\{1 \cdot a, 1 \cdot c\}$  d)  $\{1 \cdot b, 4 \cdot d\}$  e)  $\{5 \cdot a, 5 \cdot b, 1 \cdot c, 4 \cdot d\}$  61.  $\bar{F} = \{0.4 \text{ Alice}, 0.1 \text{ Brian}, 0.6 \text{ Fred}, 0.9 \text{ Oscar}, 0.5 \text{ Rita}\}$ ,  $\bar{R} = \{0.6 \text{ Alice}, 0.2 \text{ Brian}, 0.8 \text{ Fred}, 0.1 \text{ Oscar}, 0.3 \text{ Rita}\}$  63.  $F \cap R = \{0.4 \text{ Alice}, 0.8 \text{ Brian}, 0.2 \text{ Fred}, 0.1 \text{ Oscar}, 0.5 \text{ Rita}\}$

## Section 2.3

1. a)  $f(0)$  is not defined. b)  $f(x)$  is not defined for  $x < 0$ . c)  $f(x)$  is not well-defined because there are two

distinct values assigned to each  $x$ . **3. a)** Not a function **b)** A function **c)** Not a function **5. a)** Domain the set of bit strings; range the set of integers **b)** Domain the set of bit strings; range the set of even nonnegative integers **c)** Domain the set of bit strings; range the set of nonnegative integers not exceeding 7 **d)** Domain the set of positive integers; range the set of squares of positive integers  $= \{1, 4, 9, 16, \dots\}$  **7. a)** Domain  $\mathbb{Z}^+ \times \mathbb{Z}^+$ ; range  $\mathbb{Z}^+$  **b)** Domain  $\mathbb{Z}^+$ ; range  $\{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$  **c)** Domain the set of bit strings; range  $\mathbb{N}$  **d)** Domain the set of bit strings; range  $\mathbb{N}$  **9. a)** 1 **b)** 0 **c)** 0 **d)** -1 **e)** 3 **f)** -1 **g)** 2 **h)** 1 **11.** Only the function in part (a) **13.** Only the functions in parts (a) and (d) **15. a)** Onto **b)** Not onto **c)** Onto **d)** Not onto **e)** Onto **17. a)** The function  $f(x)$  with  $f(x) = 3x + 1$  when  $x \geq 0$  and  $f(x) = -3x + 2$  when  $x < 0$  **b)**  $f(x) = |x| + 1$  **c)** The function  $f(x)$  with  $f(x) = 2x + 1$  when  $x \geq 0$  and  $f(x) = -2x$  when  $x < 0$  **d)**  $f(x) = x^2 + 1$  **19. a)** Yes **b)** No **c)** Yes **d)** No **21.** Suppose that  $f$  is strictly decreasing. This means that  $f(x) > f(y)$  whenever  $x < y$ . To show that  $g$  is strictly increasing, suppose that  $x < y$ . Then  $g(x) = 1/f(x) < 1/f(y) = g(y)$ . Conversely, suppose that  $g$  is strictly increasing. This means that  $g(x) < g(y)$  whenever  $x < y$ . To show that  $f$  is strictly decreasing, suppose that  $x < y$ . Then  $f(x) = 1/g(x) > 1/g(y) = f(y)$ . **23.** Many answers are possible. One example is  $f(x) = 17$ . **25.** The function is not one-to-one, so it is not invertible. On the restricted domain, the function is the identity function on the nonnegative real numbers,  $f(x) = x$ , so it is its own inverse. **27. a)**  $f(S) = \{0, 1, 3\}$  **b)**  $f(S) = \{0, 1, 3, 5, 8\}$  **c)**  $f(S) = \{0, 8, 16, 40\}$  **d)**  $f(S) = \{1, 12, 33, 65\}$  **29. a)** Let  $x$  and  $y$  be distinct elements of  $A$ . Because  $g$  is one-to-one,  $g(x)$  and  $g(y)$  are distinct elements of  $B$ . Because  $f$  is one-to-one,  $f(g(x)) = (f \circ g)(x)$  and  $f(g(y)) = (f \circ g)(y)$  are distinct elements of  $C$ . Hence,  $f \circ g$  is one-to-one. **b)** Let  $y \in C$ . Because  $f$  is onto,  $y = f(b)$  for some  $b \in B$ . Now because  $g$  is onto,  $b = g(x)$  for some  $x \in A$ . Hence,  $y = f(b) = f(g(x)) = (f \circ g)(x)$ . It follows that  $f \circ g$  is onto. **31.** No. For example, suppose that  $A = \{a\}$ ,  $B = \{b, c\}$ , and  $C = \{d\}$ . Let  $g(a) = b$ ,  $f(b) = d$ , and  $f(c) = d$ . Then  $f$  and  $f \circ g$  are onto, but  $g$  is not. **33.**  $(f + g)(x) = x^2 + x + 3$ ,  $(fg)(x) = x^3 + 2x^2 + x + 2$  **35.**  $f$  is one-to-one because  $f(x_1) = f(x_2) \rightarrow ax_1 + b = ax_2 + b \rightarrow ax_1 = ax_2 \rightarrow x_1 = x_2$ .  $f$  is onto because  $f((y - b)/a) = y$ .  $f^{-1}(y) = (y - b)/a$ . **37.** Let  $f(1) = a$ ,  $f(2) = a$ . Let  $S = \{1\}$  and  $T = \{2\}$ . Then  $f(S \cap T) = f(\emptyset) = \emptyset$ , but  $f(S) \cap f(T) = \{a\} \cap \{a\} = \{a\}$ . **39. a)**  $\{x \mid 0 \leq x < 1\}$  **b)**  $\{x \mid -1 \leq x < 2\}$  **c)**  $\emptyset$  **41.**  $f^{-1}(\overline{S}) = \{x \in A \mid f(x) \notin S\} = \{x \in A \mid f(x) \in \overline{S}\} = \overline{f^{-1}(S)}$  **43.** Let  $x = \lfloor x \rfloor + \epsilon$ , where  $\epsilon$  is a real number with  $0 \leq \epsilon < 1$ . If  $\epsilon < \frac{1}{2}$ , then  $\lfloor x \rfloor - 1 < x - \frac{1}{2} < \lfloor x \rfloor$ , so  $\lfloor x - \frac{1}{2} \rfloor = \lfloor x \rfloor - 1$  and this is the integer closest to  $x$ . If  $\epsilon \geq \frac{1}{2}$ , then  $\lfloor x \rfloor < x - \frac{1}{2} < \lfloor x \rfloor + 1$ , so  $\lfloor x - \frac{1}{2} \rfloor = \lfloor x \rfloor$  and this is the integer closest to  $x$ . If  $\epsilon = \frac{1}{2}$ , then  $\lfloor x - \frac{1}{2} \rfloor = \lfloor x \rfloor$ , which is the smaller of the two integers that surround  $x$  and are the same distance from  $x$ . **45.** Write the real number  $x$  as  $\lfloor x \rfloor + \epsilon$ , where  $\epsilon$  is a real number with  $0 \leq \epsilon < 1$ . Because  $\epsilon = x - \lfloor x \rfloor$ , it follows that  $0 \leq -\lfloor x \rfloor < 1$ . The first two in-

equalities,  $x - 1 < \lfloor x \rfloor$  and  $\lfloor x \rfloor \leq x$ , follow directly. For the other two inequalities, write  $x = \lfloor x \rfloor - \epsilon'$ , where  $0 \leq \epsilon' < 1$ . Then  $0 \leq \lfloor x \rfloor - x < 1$ , and the desired inequality follows. **47. a)** If  $x < n$ , because  $\lfloor x \rfloor \leq x$ , it follows that  $\lfloor x \rfloor < n$ . Suppose that  $x \geq n$ . By the definition of the floor function, it follows that  $\lfloor x \rfloor \geq n$ . This means that if  $\lfloor x \rfloor < n$ , then  $x < n$ . **b)** If  $n < x$ , then because  $x \leq \lceil x \rceil$ , it follows that  $n \leq \lceil x \rceil$ . Suppose that  $n \geq x$ . By the definition of the ceiling function, it follows that  $\lceil x \rceil \leq n$ . This means that if  $n < \lceil x \rceil$ , then  $n < x$ . **49.** If  $n$  is even, then  $n = 2k$  for some integer  $k$ . Thus,  $\lfloor n/2 \rfloor = \lfloor k \rfloor = k = n/2$ . If  $n$  is odd, then  $n = 2k + 1$  for some integer  $k$ . Thus,  $\lfloor n/2 \rfloor = \lfloor k + \frac{1}{2} \rfloor = k = (n - 1)/2$ . **51.** Assume that  $x \geq 0$ . The left-hand side is  $\lceil -x \rceil$  and the right-hand side is  $-\lfloor x \rfloor$ . If  $x$  is an integer, then both sides equal  $-x$ . Otherwise, let  $x = n + \epsilon$ , where  $n$  is a natural number and  $\epsilon$  is a real number with  $0 \leq \epsilon < 1$ . Then  $\lceil -x \rceil = \lceil -n - \epsilon \rceil = -n$  and  $-\lfloor x \rfloor = -\lfloor n + \epsilon \rfloor = -n$  also. When  $x < 0$ , the equation also holds because it can be obtained by substituting  $-x$  for  $x$ . **53.**  $\lceil b \rceil - \lfloor a \rfloor - 1$  **55. a)** 1 **b)** 3 **c)** 126 **d)** 3600 **57. a)** 100 **b)** 256 **c)** 1030 **d)** 30,200



**g)** See part (a). **65.**  $f^{-1}(y) = (y - 1)^{1/3}$  **67. a)**  $f_{A \cap B}(x) = 1 \leftrightarrow x \in A \cap B \leftrightarrow x \in A \text{ and } x \in B \leftrightarrow f_A(x) = 1 \text{ and } f_B(x) = 1$

$f_B(x) = 1 \Leftrightarrow f_A(x)f_B(x) = 1$     **b)**  $f_{A \cup B}(x) = 1 \Leftrightarrow x \in A \cup B \Leftrightarrow x \in A \text{ or } x \in B \Leftrightarrow f_A(x) = 1 \text{ or } f_B(x) = 1 \Leftrightarrow f_A(x) + f_B(x) - f_A(x)f_B(x) = 1$     **c)**  $f_{\bar{A}}(x) = 1 \Leftrightarrow x \in \bar{A} \Leftrightarrow x \notin A \Leftrightarrow f_A(x) = 0 \Leftrightarrow 1 - f_A(x) = 1$     **d)**  $f_{A \oplus B}(x) = 1 \Leftrightarrow x \in A \oplus B \Leftrightarrow (x \in A \text{ and } x \notin B) \text{ or } (x \notin A \text{ and } x \in B) \Leftrightarrow f_A(x) + f_B(x) - 2f_A(x)f_B(x) = 1$     **69. a)** True; because  $\lfloor x \rfloor$  is already an integer,  $\lceil \lfloor x \rfloor \rceil = \lfloor x \rfloor$ .    **b)** False;  $x = \frac{1}{2}$  is a counterexample.    **c)** True; if  $x$  or  $y$  is an integer, then by property 4b in Table 1, the difference is 0. If neither  $x$  nor  $y$  is an integer, then  $x = n + \epsilon$  and  $y = m + \delta$ , where  $n$  and  $m$  are integers and  $\epsilon$  and  $\delta$  are positive real numbers less than 1. Then  $m + n < x + y < m + n + 2$ , so  $\lceil x + y \rceil$  is either  $m + n + 1$  or  $m + n + 2$ . Therefore, the given expression is either  $(n + 1) + (m + 1) - (m + n + 1) = 1$  or  $(n + 1) + (m + 1) - (m + n + 2) = 0$ , as desired.    **d)** False;  $x = \frac{1}{4}$  and  $y = 3$  is a counterexample.    **e)** False;  $x = \frac{1}{2}$  is a counterexample.    **71. a)** If  $x$  is a positive integer, then the two sides are equal. So suppose that  $x = n^2 + m + \epsilon$ , where  $n^2$  is the largest perfect square less than  $x$ ,  $m$  is a nonnegative integer, and  $0 < \epsilon \leq 1$ . Then both  $\sqrt{x}$  and  $\sqrt{\lfloor x \rfloor} = \sqrt{n^2 + m}$  are between  $n$  and  $n + 1$ , so both sides equal  $n$ .    **b)** If  $x$  is a positive integer, then the two sides are equal. So suppose that  $x = n^2 - m - \epsilon$ , where  $n^2$  is the smallest perfect square greater than  $x$ ,  $m$  is a nonnegative integer, and  $\epsilon$  is a real number with  $0 < \epsilon \leq 1$ . Then both  $\sqrt{x}$  and  $\sqrt{\lfloor x \rfloor} = \sqrt{n^2 - m}$  are between  $n - 1$  and  $n$ . Therefore, both sides of the equation equal  $n$ .    **73. a)** Domain is  $\mathbf{Z}$ ; codomain is  $\mathbf{R}$ ; domain of definition is the set of nonzero integers; the set of values for which  $f$  is undefined is  $\{0\}$ ; not a total function.    **b)** Domain is  $\mathbf{Z}$ ; codomain is  $\mathbf{Z}$ ; domain of definition is  $\mathbf{Z}$ ; set of values for which  $f$  is undefined is  $\emptyset$ ; total function.    **c)** Domain is  $\mathbf{Z} \times \mathbf{Z}$ ; codomain is  $\mathbf{Q}$ ; domain of definition is  $\mathbf{Z} \times (\mathbf{Z} - \{0\})$ ; set of values for which  $f$  is undefined is  $\mathbf{Z} \times \{0\}$ ; not a total function.    **d)** Domain is  $\mathbf{Z} \times \mathbf{Z}$ ; codomain is  $\mathbf{Z}$ ; domain of definition is  $\mathbf{Z} \times \mathbf{Z}$ ; set of values for which  $f$  is undefined is  $\emptyset$ ; total function.    **e)** Domain is  $\mathbf{Z} \times \mathbf{Z}$ ; codomain is  $\mathbf{Z}$ ; domain of definitions is  $\{(m, n) \mid m > n\}$ ; set of values for which  $f$  is undefined is  $\{(m, n) \mid m \leq n\}$ ; not a total function.    **75. a)** By definition, to say that  $S$  has cardinality  $m$  is to say that  $S$  has exactly  $m$  distinct elements. Therefore we can assign the first object to 1, the second to 2, and so on. This provides the one-to-one correspondence.    **b)** By part (a), there is a bijection  $f$  from  $S$  to  $\{1, 2, \dots, m\}$  and a bijection  $g$  from  $T$  to  $\{1, 2, \dots, m\}$ . Then the composition  $g^{-1} \circ f$  is the desired bijection from  $S$  to  $T$ .    **77.** It is clear from the formula that the range of values the function takes on for a fixed value of  $m + n$ , say  $m + n = x$ , is  $(x - 2)(x - 1)/2 + 1$  through  $(x - 2)(x - 1)/2 + (x - 1)$ , because  $m$  can assume the values  $1, 2, 3, \dots, (x - 1)$  under these conditions, and the first term in the formula is a fixed positive integer when  $m + n$  is fixed. To show that this function is one-to-one and onto, we merely need to show that the range of values for  $x + 1$  picks up precisely where the range of values for  $x$  left off, i.e., that  $f(x - 1, 1) + 1 = f(1, x)$ . We have  $f(x - 1, 1) + 1 = \frac{(x-2)(x-1)}{2} + (x - 1) + 1 = \frac{x^2 - x + 2}{2} = \frac{(x-1)x}{2} + 1 = f(1, x)$ .

## Section 2.4

**1. a)** 3    **b)** -1    **c)** 787    **d)** 2639    **3. a)**  $a_0 = 2, a_1 = 3, a_2 = 5, a_3 = 9$     **b)**  $a_0 = 1, a_1 = 4, a_2 = 27, a_3 = 256$   
**c)**  $a_0 = 0, a_1 = 0, a_2 = 1, a_3 = 1$     **d)**  $a_0 = 0, a_1 = 1, a_2 = 2, a_3 = 3$     **5. a)** 2, 5, 8, 11, 14, 17, 20, 23, 26, 29  
**b)** 1, 1, 1, 2, 2, 2, 3, 3, 3, 4    **c)** 1, 1, 3, 3, 5, 5, 7, 7, 9, 9  
**d)** -1, -2, -2, 8, 88, 656, 4912, 40064, 362368, 3627776  
**e)** 3, 6, 12, 24, 48, 96, 192, 384, 768, 1536    **f)** 1, 1, 2, 3, 5, 8, 13, 21, 34, 55    **g)** 1, 2, 2, 3, 3, 3, 3, 4, 4, 4    **h)** 3, 3, 5, 4, 4, 3, 5, 5, 4, 3    **7.** Each term could be twice the previous term; the  $n$ th term could be obtained from the previous term by adding  $n - 1$ ; the terms could be the positive integers that are not multiples of 3; there are infinitely many other possibilities.    **9. a)** One 1 and one 0, followed by two 1s and two 0s, followed by three 1s and three 0s, and so on; 1, 1, 1  
**b)** The positive integers are listed in increasing order with each even positive integer listed twice; 9, 10, 10.    **c)** The terms in odd-numbered locations are the successive powers of 2; the terms in even-numbered locations are all 0; 32, 0, 64.  
**d)**  $a_n = 3 \cdot 2^{n-1}$ ; 384, 768, 1536    **e)**  $a_n = 15 - 7(n - 1) = 22 - 7n$ ; -34, -41, -48    **f)**  $a_n = (n^2 + n + 4)/2$ ; 57, 68, 80    **g)**  $a_n = 2n^3$ ; 1024, 1458, 2000    **h)**  $a_n = n! + 1$ ; 362881, 3628801, 39916801    **11.** Among the integers  $1, 2, \dots, a_n$ , where  $a_n$  is the  $n$ th positive integer not a perfect square, the nonsquares are  $a_1, a_2, \dots, a_n$  and the squares are  $1^2, 2^2, \dots, k^2$ , where  $k$  is the integer with  $k^2 < n + k < (k + 1)^2$ . Consequently,  $a_n = n + k$ , where  $k^2 < a_n < (k + 1)^2$ . To find  $k$ , first note that  $k^2 < n + k < (k + 1)^2$ , so  $k^2 + 1 \leq n + k \leq (k + 1)^2 - 1$ . Hence,  $(k - \frac{1}{2})^2 + \frac{3}{4} = k^2 - k + 1 \leq n \leq k^2 + k = (k + \frac{1}{2})^2 - \frac{1}{4}$ . It follows that  $k - \frac{1}{2} < \sqrt{n} < k + \frac{1}{2}$ , so  $k = \lfloor \sqrt{n} \rfloor$  and  $a_n = n + k = n + \lfloor \sqrt{n} \rfloor$ .    **13. a)** 20  
**b)** 11    **c)** 30    **d)** 511    **15. a)** 1533    **b)** 510    **c)** 4923  
**d)** 9842    **17. a)** 21    **b)** 78    **c)** 18    **d)** 18  
**19.**  $\sum_{j=1}^n (a_j - a_{j-1}) = a_n - a_0$     **21. a)**  $n^2$     **b)**  $n(n + 1)/2$     **23.** 15150    **25.**  $\frac{n(n+1)(2n+1)}{3} + \frac{n(n+1)}{2} + (n + 1)(m - (n + 1)^2 + 1)$ , where  $n = \lfloor \sqrt{m} \rfloor - 1$     **27. a)** 0    **b)** 1680  
**c)** 1    **d)** 1024    **29.** 34    **31. a)** Countable, -1, -2, -3, -4, ...    **b)** Countable, 0, 2, -2, 4, -4, ...    **c)** Uncountable  
**d)** Countable, 0, 7, -7, 14, -14, ...    **33. a)** Countable: match  $n$  with the string of  $n$  1s.    **b)** Countable. To find a correspondence, follow the path in Example 20, but omit fractions in the top three rows (as well as continuing to omit fractions not in lowest terms).    **c)** Uncountable  
**d)** Uncountable    **35.** Assume that  $A - B$  is countable. Then, because  $A = (A - B) \cup B$ , the elements of  $A$  can be listed in a sequence by alternating elements of  $A - B$  and elements of  $B$ . This contradicts the uncountability of  $A$ .    **37.** Assume that  $B$  is countable. Then the elements of  $B$  can be listed as  $b_1, b_2, b_3, \dots$ . Because  $A$  is a subset of  $B$ , taking the subsequence of  $\{b_n\}$  that contains the terms that are in  $A$  gives a listing of the elements of  $A$ . Because  $A$  is uncountable, this is impossible.    **39.** We are given bijections  $f$  from  $A$  to  $B$  and  $g$  from  $C$  to  $D$ . Then the function from  $A \times C$  to  $B \times D$  that sends  $(a, c)$  to  $(f(a), g(c))$  is a bijection.    **41.** Suppose that  $A_1, A_2, A_3, \dots$  are countable sets. Because  $A_i$  is countable,

we can list its elements in a sequence as  $a_{i1}, a_{i2}, a_{i3}, \dots$ . The elements of the set  $\bigcup_{i=1}^n A_i$  can be listed by listing all terms  $a_{ij}$  with  $i + j = 2$ , then all terms  $a_{ij}$  with  $i + j = 3$ , then all terms  $a_{ij}$  with  $i + j = 4$ , and so on. **43.** There are a finite number of bit strings of length  $m$ , namely,  $2^m$ . The set of all bit strings is the union of the sets of bit strings of length  $m$  for  $m = 0, 1, 2, \dots$ . Because the union of a countable number of countable sets is countable (see Exercise 41), there are a countable number of bit strings. **45.** For any finite alphabet there are a finite number of strings of length  $n$ , whenever  $n$  is a positive integer. It follows by the result of Exercise 41 that there are only a countable number of strings from any given finite alphabet. Because the set of all computer programs in a particular language is a subset of the set of all strings of a finite alphabet, which is a countable set by the result from Exercise 36, it is itself a countable set. **47.** Exercise 45 shows that there are only a countable number of computer programs. Consequently, there are only a countable number of computable functions. Because, as Exercise 46 shows, there are an uncountable number of functions, not all functions are computable.

### Supplementary Exercises

1. a)  $\bar{A}$  b)  $A \cap B$  c)  $A - B$  d)  $\bar{A} \cap \bar{B}$  e)  $A \oplus B$   
 3. Yes 5.  $A - (A - B) = A - (A \cap \bar{B}) = A \cap (A \cap \bar{B}) = A \cap (\bar{A} \cup B) = (A \cap \bar{A}) \cup (A \cap B) = \emptyset \cup (A \cap B) = A \cap B$   
 7. Let  $A = \{1\}$ ,  $B = \emptyset$ ,  $C = \{1\}$ . Then  $(A - B) - C = \emptyset$ , but  $A - (B - C) = \{1\}$ . 9. No. For example, let  $A = B = \{a, b\}$ ,  $C = \emptyset$ , and  $D = \{a\}$ . Then  $(A - B) - (C - D) = \emptyset - \emptyset = \emptyset$ , but  $(A - C) - (B - D) = \{a, b\} - \{b\} = \{a\}$ .  
 11. a)  $|\emptyset| \leq |A \cap B| \leq |A| \leq |A \cup B| \leq |U|$  b)  $|\emptyset| \leq |A - B| \leq |A \oplus B| \leq |A \cup B| \leq |A| + |B|$  13. a) Yes, no b) Yes, no c)  $f$  has inverse with  $f^{-1}(a) = 3$ ,  $f^{-1}(b) = 4$ ,  $f^{-1}(c) = 2$ ,  $f^{-1}(d) = 1$ ;  $g$  has no inverse. 15. Let  $f(a) = f(b) = 1$ ,  $f(c) = f(d) = 2$ ,  $S = \{a, c\}$ ,  $T = \{b, d\}$ . Then  $f(S \cap T) = f(\emptyset) = \emptyset$ , but  $f(S) \cap f(T) = \{1, 2\} \cap \{1, 2\} = \{1, 2\}$ . 17. Let  $x \in A$ . Then  $S_f(\{x\}) = \{f(y) \mid y \in \{x\}\} = \{f(x)\}$ . By the same reasoning,  $S_g(\{x\}) = \{g(x)\}$ . Because  $S_f = S_g$ , we can conclude that  $\{f(x)\} = \{g(x)\}$ , and so necessarily  $f(x) = g(x)$ . 19. The equation is true if and only if the sum of the fractional parts of  $x$  and  $y$  is less than 1. 21. The equation is true if and only if either both  $x$  and  $y$  are integers, or  $x$  is not an integer but the sum of the fractional parts of  $x$  and  $y$  is less than or equal to 1. 23. If  $x$  is an integer, then  $\lfloor x \rfloor + \lfloor m - x \rfloor = x + m - x = m$ . Otherwise, write  $x$  in terms of its integer and fractional parts:  $x = n + \epsilon$ , where  $n = \lfloor x \rfloor$  and  $0 < \epsilon < 1$ . In this case  $\lfloor x \rfloor + \lfloor m - x \rfloor = \lfloor n + \epsilon \rfloor + \lfloor m - n - \epsilon \rfloor = n + m - n - 1 = m - 1$ . 25. Write  $n = 2k + 1$  for some integer  $k$ . Then  $n^2 = 4k^2 + 4k + 1$ , so  $n^2/4 = k^2 + k + \frac{1}{4}$ . Therefore,  $\lceil n^2/4 \rceil = k^2 + k + 1$ . But  $(n^2 + 3)/4 = (4k^2 + 4k + 1 + 3)/4 = k^2 + k + 1$ . 27. Let  $x = n + (r/m) + \epsilon$ , where  $n$  is an integer,  $r$  is a nonnegative integer less than  $m$ , and  $\epsilon$  is a real number with  $0 \leq \epsilon < 1/m$ . The left-hand side is  $\lfloor nm + r + m\epsilon \rfloor = nm + r$ . On the right-hand side, the terms  $\lfloor x \rfloor$  through  $\lfloor x +$

$(m + r - 1)/m \rfloor$  are all just  $n$  and the terms from  $\lfloor x + (m - r)/m \rfloor$  on are all  $n + 1$ . Therefore, the right-hand side is  $(m - r)n + r(n + 1) = nm + r$ , as well. 29. 101 31.  $a_1 = 1$ ;  $a_{2n+1} = n \cdot a_{2n}$  for all  $n > 0$ ; and  $a_{2n} = n + a_{2n-1}$  for all  $n > 0$ . The next four terms are 5346, 5353, 37471, and 37479.

## CHAPTER 3

### Section 3.1

1.  $max := 1, i := 2, max := 8, i := 3, max := 12, i := 4, i := 5, i := 6, i := 7, max := 14, i := 8, i := 9, i := 10, i := 11$   
 3. **procedure** *sum*( $a_1, \dots, a_n$ : integers)  
      $sum := a_1$   
     **for**  $i := 2$  **to**  $n$   
          $sum := sum + a_i$   
     {*sum* has desired value}  
 5. **procedure** *duplicates*( $a_1, a_2, \dots, a_n$ : integers in nondecreasing order)  
      $k := 0$  {this counts the duplicates}  
      $j := 2$   
     **while**  $j \leq n$   
         **begin**  
             **if**  $a_j = a_{j-1}$  **then**  
                 **begin**  
                      $k := k + 1$   
                      $c_k := a_j$   
                     **while** ( $j \leq n$  and  $a_j = c_k$ )  
                          $j := j + 1$   
                 **end**  
              $j := j + 1$   
         **end** { $c_1, c_2, \dots, c_k$  is the desired list}  
 7. **procedure** *last even location*( $a_1, a_2, \dots, a_n$ : integers)  
      $k := 0$   
     **for**  $i := 1$  **to**  $n$   
         **if**  $a_i$  is even **then**  $k := i$   
     **end** { $k$  is the desired location (or 0 if there are no evens)}  
 9. **procedure** *palindrome check*( $a_1 a_2 \dots a_n$ : string)  
      $answer := \text{true}$   
     **for**  $i := 1$  **to**  $\lfloor n/2 \rfloor$   
         **if**  $a_i \neq a_{n+1-i}$  **then**  $answer := \text{false}$   
     **end** {*answer* is true iff string is a palindrome}  
 11. **procedure** *interchange*( $x, y$ : real numbers)  
      $z := x$   
      $x := y$   
      $y := z$   
 The minimum number of assignments needed is three.  
 13. Linear search:  $i := 1, i := 2, i := 3, i := 4, i := 5, i := 6, i := 7, location := 7$ ; binary search:  $i := 1, j := 8, m := 4, i := 5, m := 6, i := 7, m := 7, j := 7, location := 7$   
 15. **procedure** *insert*( $x, a_1, a_2, \dots, a_n$ : integers)  
     {the list is in order:  $a_1 \leq a_2 \leq \dots \leq a_n$ }  
      $a_{n+1} := x + 1$   
      $i := 1$