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Group 4 Topic # 26 Payroll Management DBMS		
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### Assignment 8 Goal:

See the diagram examples posted for A7 (4 marks)

**Lab 9) Assignment 8: Normalization 3NF/ BCNF by Algorithm:** (4 marks). For A8 to change to 3NF you need to add additional FDs to be able to breakdown at least one table by Bernstein's Algorithm and show the details for making that table 3NF. Then verify all tables whether or not they are in 3NF or BCNF as shown in class. You can explain why each table is BCNF by showing FDs. You need to use BCNF algorithm (by adding FDs if needed) and showing the details for at least one table to BCNF. Note that during assignments 8 you use algorithm for 3NF or BCNF by both Bernstein's Algorithm for 3NF and BCNF algorithms each of them at least for one table. It is recommended to start creating Java/web based UI

### FOLLOWING FDs for the Code:

For the EMPLOYEE table:

- EMPLOYEE\_ID → NAME
- EMPLOYEE\_ID → DESIGNATION\_ID

For the SALARY table:

- SALARY\_ID → EMPLOYEE\_ID
- SALARY\_ID → AMOUNT

For the PAYMENT table:

- PAYMENT\_ID → EMPLOYEE\_ID
- PAYMENT\_ID → AMOUNT
- PAYMENT\_ID → DATE\_RECEIVED

For the TAX table:

- TAX\_ID → EMPLOYEE\_ID
- TAX\_ID → TAX\_AMOUNT

For the DEDUCTION table:

- DEDUCTION\_ID → EMPLOYEE\_ID
- DEDUCTION\_ID → DEDUCTION\_AMOUNT

- DEDUCTION\_ID → REASON

### USING THE BernStien's Algorithm:

As Employee and Salary tables (FD):

- EMPLOYEE\_ID → NAME
- EMPLOYEE\_ID → DESIGNATION\_ID

For the SALARY table:

- EMPLOYEE\_ID → AMOUNT
- For the EMPLOYEE table, EMPLOYEE\_ID is the candidate key.
- For the SALARY table, EMPLOYEE\_ID uniquely determines AMOUNT.

### DECOMPOSITION TABLE:

```
CREATE TABLE EMPLOYEE (
  EMPLOYEE_ID VARCHAR2(100) NOT NULL PRIMARY KEY,
  NAME VARCHAR2(100) NOT NULL,
  DESIGNATION_ID VARCHAR2(100) REFERENCES DESIGNATION(DESIGNATION_ID)
);
```

```
CREATE TABLE SALARY (
  SALARY_ID VARCHAR2(100) NOT NULL PRIMARY KEY,
  EMPLOYEE_ID VARCHAR2(100) REFERENCES EMPLOYEE(EMPLOYEE_ID),
  AMOUNT NUMBER(10,2) NOT NULL CHECK (AMOUNT >= 0)
);
```

### TO check whether if it's 3NF or BCNF:

- For 3NF, every non-prime attribute must be non-transitively dependent on every key.
- For BCNF, for every FD  $X \rightarrow Y$ , X must be a superkey.

### FIRST TABLE (DESIGNATION TABLE):

- FD: DESIGNATION\_ID → TITLE
- DESIGNATION\_ID is the primary key.
- No non-prime attributes exist (all attributes are part of a key).

- Conclusion: The DESIGNATION table is in BCNF (and therefore also in 3NF) because the only FD has a superkey on the left-hand side.

## **SECOND TABLE (EMPLOYEE TABLE):**

- FDs:  $EMPLOYEE\_ID \rightarrow NAME$ ,  $EMPLOYEE\_ID \rightarrow DESIGNATION\_ID$
- $EMPLOYEE\_ID$  is the primary key.
- Both  $NAME$  and  $DESIGNATION\_ID$  are non-prime attributes and are fully functionally dependent on the primary key.
- There are no transitive dependencies.
- Conclusion: The EMPLOYEE table is in BCNF (and therefore also in 3NF) because every FD has a superkey ( $EMPLOYEE\_ID$ ) on the left-hand side.

## **THIRD TABLE (SALARY TABLE):**

- FD:  $EMPLOYEE\_ID \rightarrow AMOUNT$
- Initially, this might seem incorrect because  $SALARY\_ID$  is the primary key, not  $EMPLOYEE\_ID$ . However, if  $EMPLOYEE\_ID$  is unique (which is suggested by the UNIQUE constraint in the provided SQL), then it can be considered an alternate key.

## **PAYMENT, TAX, and DEDUCTION Tables:**

- $PAYMENT\_ID \rightarrow EMPLOYEE\_ID, AMOUNT, DATE\_RECEIVED$
- $TAX\_ID \rightarrow EMPLOYEE\_ID, TAX\_AMOUNT$
- $DEDUCTION\_ID \rightarrow EMPLOYEE\_ID, DEDUCTION\_AMOUNT, REASON$

All three tables are in BCNF (and therefore also in 3NF) because, for all FDs, the left-hand side is a superkey.

**You need to use BCNF algorithm (by adding FDs if needed) and showing the details for at least one table to BCNF:**

$DEPARTMENT\_ID \rightarrow AMOUNT$

```
CREATE TABLE SALARY (
  SALARY_ID VARCHAR2(100) NOT NULL PRIMARY KEY,
  EMPLOYEE_ID VARCHAR2(100) UNIQUE REFERENCES EMPLOYEE(EMPLOYEE_ID),
  DEPARTMENT_ID VARCHAR2(100),
  AMOUNT NUMBER(10,2) NOT NULL CHECK (AMOUNT >= 0)
);
```

And the functional dependencies are:

$SALARY\_ID \rightarrow EMPLOYEE\_ID, DEPARTMENT\_ID, AMOUNT$   
 $EMPLOYEE\_ID \rightarrow DEPARTMENT\_ID$   
 $DEPARTMENT\_ID \rightarrow AMOUNT$

**Converting this to BCNF:**

**Decompose the Table:**

decompose the SALARY table into two tables to remove the violation:

- SALARY\_BANDS table: This table will store the relationship between DEPARTMENT\_ID and AMOUNT.

```
CREATE TABLE SALARY_BANDS (  
    DEPARTMENT_ID VARCHAR2(100) NOT NULL PRIMARY KEY,  
    AMOUNT NUMBER(10,2) NOT NULL CHECK (AMOUNT >= 0)  
);
```

Which connects to employee salary:

```
CREATE TABLE EMPLOYEE_SALARY (  
    SALARY_ID VARCHAR2(100) NOT NULL PRIMARY KEY,  
    EMPLOYEE_ID VARCHAR2(100) UNIQUE REFERENCES EMPLOYEE(EMPLOYEE_ID),  
    DEPARTMENT_ID VARCHAR2(100) REFERENCES SALARY_BANDS(DEPARTMENT_ID)  
);
```

After this decomposition:

- SALARY\_BANDS satisfies BCNF because there's only one FD, and DEPARTMENT\_ID is a superkey.
- EMPLOYEE\_SALARY satisfies BCNF because SALARY\_ID is a superkey and there are no non-key attributes.

The new schema for the salary table is now in BCNF.