

Online Learning with Committees and Its Application

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Online prediction with experts, repeated game playing and convex optimization

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A simple policy

The final decision is made with the **majority voting** from the expert committee: $\hat{y}^{(t)} = \mathbf{sign}(\frac{\sum_{n=1}^N f_i^{(t)}}{N})$;

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$w_n^{(t)} \leftarrow 0$ if expert E_n makes a mistake at time $t - 1$, i.e. kick E_n out of the committee;

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Cummulative loss:

$$\hat{L}^{(t)} = \sum_{i=t}^t l(\hat{y}^{(t)}, y^{(t)}) \leq \log_2 N; \quad (1)$$

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$$R_n^{(t)} = \hat{L}^{(t)} - L_n^{(t)} = \sum_{i=t}^t l(\hat{y}^{(t)}, y^{(t)}) - \sum_{i=t}^t l(f_n^{(t)}, y^{(t)}) \quad (2)$$

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the upper bound of regret

$$R^{(t)*} = \max_{n \in [1, N]} R_n^{(t)} = \hat{L}^{(t)} - \min_{n \in [1, N]} L_n^{(t)}$$

Weighted majority algorithm

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$$w_n^{(t+1)} = (1 - \eta) w_n^{(t)} \quad (3)$$

where $\eta \leq \frac{1}{2}$.

Analysis on weighted majority algorithm

Theorem (cummulative loss bound using weighted majority)

After T steps, $\hat{L}^{(T)} \leq 2(1 + \eta) \min_{1 \in [1, N]} L_n^{(T)} + \frac{2 \ln N}{\eta}$

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Proof: Let $\Gamma^{(t)} = \sum_{n \in [1, N]} w_n^{(t)}$, then $\Gamma^{(1)} = N$. Also, if the forecaster makes a mistake, $\hat{y}^{(t)} \neq y^{(t)}$

$$\Gamma^{(t+1)} \leq \Gamma^{(t)} \left(\frac{1}{2} + \frac{1}{2}(1 - \eta) \right) = \Gamma^{(t)} \left(1 - \frac{\eta}{2} \right) \quad (4)$$

therefore:

$$\Gamma^{(T+1)} \leq N \left(1 - \frac{\eta}{2} \right)^{\hat{L}^{(T)}} \quad (5)$$

for any individual expert n

$$w_n^{(T+1)} = (1 - \eta)^{L_n^{(T)}} \quad (6)$$

since $w_n^{(T+1)} \leq \Gamma^{(T+1)} \implies (1 - \eta)^{L_n^{(T)}} \leq N \left(1 - \frac{\eta}{2} \right)^{\hat{L}^{(T)}}$

Analysis on weighted majority algorithm, cont.

$$\begin{aligned}
 (1 - \eta)^{L_n^{(T)}} &\leq N(1 - \frac{\eta}{2})^{\hat{L}^{(T)}} \\
 \Leftrightarrow L_n^{(T)} \ln(1 - \eta) &\leq \ln N + \hat{L}^{(T)} \ln(1 - \frac{\eta}{2}) \\
 \Leftrightarrow -\ln(1 - \frac{\eta}{2}) &\leq -L_n^{(T)} \ln(1 - \eta) + \ln N \\
 \xRightarrow{x \leq -\ln(1-x)} \frac{\eta}{2} \hat{L}^{(T)} &\leq -L_n^{(T)} \ln(1 - \eta) + \ln N \\
 \xRightarrow{-\ln(1-x) \leq x + x^2, \text{ when } x \leq 1/2} \frac{\eta}{2} \hat{L}^{(T)} &\leq L_n^{(T)} \eta(1 + \eta) + \ln N \\
 \Leftrightarrow \hat{L}^{(T)} &\leq 2(1 + \eta) L_n^{(T)} + \frac{2 \ln N}{\eta}
 \end{aligned} \tag{7}$$

Randomized weighted majority algorithm

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Proof: Let $\Gamma^{(t)} = \sum_{n \in [1, N]} w_n^{(t)}$, then $\Gamma^{(1)} = N$.

At each time t , let $F^{(t)} = \frac{\sum_{n: f_n^{(t)} \neq y^{(t)}} w_n^{(t)}}{\sum_n w_n^{(t)}}$, then

$$\Gamma^{(t+1)} = \Gamma^{(t)} (1 - F^{(t)} + F^{(t)}(1 - \eta)) = \Gamma^{(t)} (1 - F^{(t)}\eta) \quad (9)$$

therefore:

$$\Gamma^{(T+1)} = N \prod_{t=1}^T (1 - F^{(t)}\eta) \quad (10)$$

for any individual expert n : $w_n^{(T+1)} = (1 - \eta)^{L_n^{(T)}}$,

since $w_n^{(T+1)} \leq \Gamma^{(T+1)} \implies (1 - \eta)^{L_n^{(T)}} \leq N \prod_{t=1}^T (1 - F^{(t)}\eta)$

Go beyond binary sequence and 0-1 loss

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 - **Forecaster:** $\hat{y}^{(t)} = \pi(f_1^{(t)}, f_2^{(t)}, \dots, f_N^{(t)})$;
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- **Task: general online prediction**, i.e. sequentially forecast a value $y^{(t)} \in \mathbb{R}$ at time t based on historical values $\{y^{(i)}\}_{i=1}^{(t-1)}$;
 - **Loss:** a cost vector $[m_1^{(t)}, m_2^{(t)}, \dots, m_N^{(t)}]$ is incurred to experts' predictions at time t ;
 - **Weights update:** multiplicative manner

$$w_n^{(t+1)} = w_n^{(t)} g(m_n^{(t)}) \quad (11)$$

where $g(x)$ is a decreasing function w.r.t. x .

Revisit aggregation policy in binary case

Forecaster: $\hat{y}_t \sim \text{Bernoulli} \left(\frac{\sum_{n: f_{n,t}=1} w_n^{(t-1)}}{\sum_n w_n^{(t-1)}}, \frac{\sum_{n: f_{n,t}=-1} w_n^{(t-1)}}{\sum_n w_n^{(t-1)}} \right);$

Let $p_n^{(t-1)} = \frac{w_n^{(t-1)}}{\sum_{n=1}^N w_n^{(t-1)}}$, then

$$\begin{aligned} \mathbb{E} \{ \hat{y}^{(t)} \} &= (-1) \cdot \sum_{n: f_{n,t}=-1} p_n^{(t-1)} + 1 \cdot \sum_{n: f_{n,t}=1} p_n^{(t-1)} \\ &= \sum_{n=1}^N f_n^{(t)} p_n^{(t-1)} \\ &= \mathbb{E}_{p^{(t-1)}} \{ f^{(t)} \} \end{aligned} \tag{12}$$

A further randomized version:

$$\hat{y}^{(t)} = f_{n_{\dagger}}^{(t)} \text{ with } n_{\dagger} \sim [p_1^{(t-1)}, p_2^{(t-1)}, \dots, p_N^{(t-1)}] \tag{13}$$

Randomized weighted majority algorithm IN GENERAL

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$$w_n^{(t+1)} = w_n^{(t)}(1 - \eta m_n^{(t)}) \quad (14)$$

where $\eta \leq 1/2$

Analysis

Theorem (Regret bound using randomized weighted majority)

When $m_n^{(t)} \in [-1, 1], \forall n, t$, after T steps,
$$\hat{L}^{(T)} \leq \min_{1 \in [1, N]} L_n^{(T)} + \eta \sum_{t=1}^T |m_n^{(t)}|_1 + \frac{\ln N}{\eta}$$

Regret bound: $\mathbf{R}^* \leq \eta T + \frac{\ln N}{\eta}$

Hedge algorithm

Weights update:

$$w_n^{(t+1)} = w_n^{(t)} \exp(-\eta m_n^{(t)}) \quad (15)$$

where $\eta \in [0, 1]$.

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One-player game

- One player plays a game, at time t he takes one action $a^{(t)} = i \in [1, N]$, then the environment releases the cost for each action $\mathbf{m}^{(t)} = [m_1^{(t)}, m_2^{(t)}, \dots, m_N^{(t)}]^\top, m_n^{(t)} \in [-1, 1]$,
- **Note** that the loss function $\mathbf{m}^{(t)}$ can change over time, i.e. the environment is changing.

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- **Note** that the loss function $\mathbf{m}^{(t)}$ can change over time, i.e. the environment is changing.
- if the player selects actions by the probabilities $\mathbf{p}^{(t)}$ computed by randomized weighted majority and update them accordingly, then

$$\sum_{t=1}^T \langle \mathbf{m}^{(t)}, \mathbf{p}^{(t)} \rangle \leq \sum_{t=1}^T m_i^{(t)} + \eta T + \frac{\ln n}{\eta} \quad (16)$$

Two-player game

Two-person zero-sum: two players ($K = 2$) play a game which is defined by a $N \times N$ cost matrix \mathbf{C} (N is the number of possible actions for players), where each entry c_{ij} defines the loss to the row player when the row player takes the action $i \in [1, N]$ and the column player takes the action $j \in [1, N]$.

An example cost matrix:

$$\begin{bmatrix} & 1 & 2 & 3 & 4 \\ 1 & c_{11} & c_{12} & c_{13} & c_{14} \\ 2 & c_{21} & c_{22} & c_{23} & c_{24} \\ 3 & c_{31} & c_{32} & c_{33} & c_{34} \\ 4 & c_{41} & c_{42} & c_{43} & c_{44} \end{bmatrix}$$

The row player's goal is to minimize its loss while the objective of the column player is to maximize it

Nash Equilibrium

- if the row player chooses his action from a distribution \mathbf{p} , then the most adversary action the column player should take is

$$j := \arg \max_{j \in [1, N]} \mathbb{E}_{i \sim \mathbf{p}} \{c_{ij}\} \quad (17)$$

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- similarly, if column player use a distribution \mathbf{q} , then the row player should take:

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von Neumann's min-max theorem

$$\min_{\mathbf{p}} \max_j \mathbb{E}_{i \sim \mathbf{p}} \{c_{ij}\} = \max_{\mathbf{q}} \max_i \mathbb{E}_{j \sim \mathbf{q}} \{c_{ij}\} = \lambda^*$$

- no player has an incentive of changing his strategy (distribution) if the player does not change his, i.e. every player is happy about current status;

Hanan's algorithm for two-player zero-sum game

Follow the perturbed leading expert

Forecaster: select the action $i^{(t)} = \arg \min_{i \in [1, N]} \left\{ L_i^{(t-1)} + \tau_i \right\}$ where τ_i are randomly sampled from $[0, 1/\epsilon]$

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Equivalence to exponential weighted majority

In Hanan's algorithm, when $\tau_i = \frac{1}{\eta} \ln \ln \frac{1}{u_i}$, where $u_i \sim [0, 1]$,

$$Pr[i^{(t)} = j] = \frac{e^{-\eta L_j^{(t)}}}{\sum_{k=1}^N e^{-\eta L_k^{(t)}}}$$

after $T = \frac{4 \ln n}{\epsilon}$ iterations, the algorithm can converge to a $\tilde{\mathbf{p}}$ which yield $\lambda^* + \epsilon$

Arbitrary convex loss function

- a convex function family on $\mathbf{p} : g(\mathbf{p})$
- the whole objective function is presented sequentially:
$$G(\cdot) = g^{(1)}(\cdot) + g^{(2)}(\cdot) + \cdots g^{(T)}(\cdot)$$

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- re-design a loss $\mathbf{m}^{(t)} = \frac{1}{\rho} \Delta f^{(t)}(\mathbf{p})$
- $$\sum_{t=1}^T f^{(t)}(\mathbf{p}^{(t)}) - \min_{\mathbf{p}} G(\mathbf{p}) \leq 2\rho\sqrt{\ln n T}$$

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- $\sum_{t=1}^T f^{(t)}(\mathbf{p}^{(t)}) - \min_{\mathbf{p}} G(\mathbf{p}) \leq 2\rho\sqrt{\ln n T}$
- applications
 - 1 online portfolio management: $g^{(t)}(\mathbf{p}) = \log(-\mathbf{p}^\top \Delta \mathbf{v}^{(t)})$
 - 2 online learning algorithms: $g^{(t)}(\mathbf{p}) = \|\mathbf{y}^{(t)} - \mathbf{p}^\top \mathbf{x}^{(t)}\|_2^2$

Bandit Optimization in metric spaces

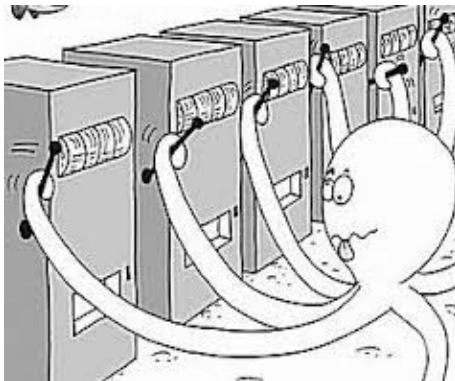
Observe only one cost

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- playing games in a Bandit setting falls into **classic exploitation v.s. exploration dilemma**

Multi-armed bandit problem



Upper confidence bound

- at each time t , play the arm $i^{(t)}$ by

$$i^{(t)} = \arg \max_{i \in [1, N]} \left\{ \mathbf{avg}.[m_i] + \sqrt{\frac{2 \ln t}{T_i}} \right\} \quad (19)$$

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- Regret bound:

$$\mathbf{R}^{(t)*} \leq \left[8 \sum_{i: \mu_i < \mu^*} \frac{\ln t}{\mu^* - \mu_i} \right] + (1 + \frac{\pi^2}{3}) (\sum_{i=1}^N (\mu^* - \mu_i)) \quad (20)$$

where μ_i denotes the true expected reward for arm i , and μ^* denotes the best one.

Contextual bandit optimization

- in classic multi-armed bandit problem, all arms are **independent**;

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- in classic multi-armed bandit problem, all arms are **independent**;
- in many practical applications, **there exists a context beneath arms**,
e.g. **representation of arms is in a metric space**;

Response surface optimization

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Response surface optimization

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- model the response of arms using a smooth surface function f
- a demo mimicking the style to understand surface-response optimization;
- a most general case, **Lipschit continuity**: if

$$d_x(\mathbf{x}_1, \mathbf{x}_2) \leq L \cdot d_f(f(\mathbf{x}_1), f(\mathbf{x}_2)) \quad (21)$$

then we say f is a L -Lipschit continuous function.

Gaussian process

Definition

A Gaussian process (GP) defines is a collection of random variables, any finite number of of which have joint Gaussian distribution.

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joint Gaussian distribution

$$\begin{bmatrix} f(\mathbf{X}_{train}) \\ f(\mathbf{x}_{test}) \end{bmatrix} \sim \left(\begin{bmatrix} \mu(\mathbf{X}_{train}) \\ \mu(\mathbf{x}_{test}) \end{bmatrix}, \begin{bmatrix} k(\mathbf{X}_{train}, \mathbf{X}_{train}), & k(\mathbf{X}_{train}, \mathbf{x}_{test}) \\ k(\mathbf{X}_{train}, \mathbf{x}_{test})^\top, & k(\mathbf{x}_{test}, \mathbf{x}_{test}) \end{bmatrix} \right) \quad (23)$$

Gaussian process, cont.

Conditional probability

$$f(\mathbf{x}_{test})|f(\mathbf{X}_{train}) \sim \mathcal{N}(\mu_{pos}(\mathbf{x}_{test}), k_{pos}(\mathbf{x}_{test}, \mathbf{x}')) \quad (24)$$

Gaussian process, cont.

Conditional probability

$$f(\mathbf{x}_{test})|f(\mathbf{X}_{train}) \sim \mathcal{N}(\mu_{pos}(\mathbf{x}_{test}), k_{pos}(\mathbf{x}_{test}, \mathbf{x}')) \quad (24)$$

Posterior Gaussian process

$$f(\mathbf{x})|\mathcal{D}_{train} \sim \mathcal{GP}(\mu_{pos}(\mathbf{x}), k_{pos}(\mathbf{x}, \mathbf{x}')) \quad (25)$$

where $\mu_{pos}(\mathbf{x}) = \mu(\mathbf{x}) + k(\mathbf{X}_{train}, \mathbf{x})^\top k(\mathbf{X}_{train}, \mathbf{X}_{train})^{-1}(f(\mathbf{X}_{train}) - \mathbf{X}_{train})$
 and $k_{pos}(\mathbf{x}, \mathbf{x}') = k(\mathbf{x}, \mathbf{x}') - k(\mathbf{X}_{train}, \mathbf{x})^\top k(\mathbf{X}_{train}, \mathbf{X}_{train})^{-1}k(\mathbf{X}_{train}, \mathbf{x}')$.

Gaussian process bandit optimization

Kernel function

squared exponential kernel: $k(\mathbf{x}, \mathbf{x}') = \exp(-\frac{\|\mathbf{x} - \mathbf{x}'\|}{2l^2})$

Information gain

the informativeness of a set of points $A \in \mathcal{X}$ for learning f is defined as:

$$I(\mathbf{y}_A; \mathbf{f}_A) = H(\mathbf{y}_A) - H(\mathbf{y}_A | \mathbf{f}_A) \quad (26)$$

Maximum information gain after T iterations

$$\gamma^{(T)} = \max_{A \in \mathcal{X}: |A|=T} I(\mathbf{y}_A; \mathbf{f}_A)$$

- maximum information gain basically reflect the complexity of f ;
- for RBF kernel, $\gamma^{(T)} = \mathcal{O}((\log T)^{d+1})$

Upper confidence bound

- at each time t , chose:

$$\mathbf{x}^{(t)} = \arg \max_{\mathbf{x} \in \mathcal{X}} \mu^{(t-1)}(\mathbf{x}) + \kappa^{(t)} \sigma^{(t-1)}(\mathbf{x}) \quad (27)$$

where $\kappa^{(t)} = 2B + 300\gamma^{(t)} \log^3(t/\delta)$, $\|f\|^2 \leq B$

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where $\kappa^{(t)} = 2B + 300\gamma^{(t)} \log^3(t/\delta)$, $\|f\|^2 \leq B$

- Regret bound:

$$\Pr \left\{ \mathbf{R}^{(T)*} \leq \sqrt{\frac{8}{\log(1 + \sigma^{-1})} T \beta^{(T)} \gamma^{(T)}} \right\} \geq 1 - \delta \quad (28)$$

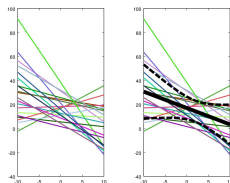
- Regret bound complexity: $\mathbf{R}^{(T)*} = \mathcal{O} \left(\sqrt{T(\log T)^{d+1}} \right)$.

Regression with ensemble models

- Gaussian process bandit optimization is an instance of Bayesian optimization;
- any function f with certain **smoothness assumption** and **uncertainty measurement** can be used as a pseudo Bayesian optimization;

Regression with ensemble models

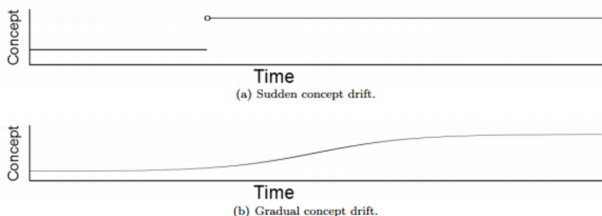
- Gaussian process bandit optimization is an instance of Bayesian optimization;
- any function f with certain **smoothness assumption** and **uncertainty measurement** can be used as a pseudo Bayesian optimization;
- **one example**: an ensemble of linear functions



- ① **smoothness assumption**: average of linear functions
- ② **uncertainty measurement**: std. of linear functions
- advantage: can better exploit task-relevant features instead of **isotropic length-scale** in kernel function

Concept drift in online decision

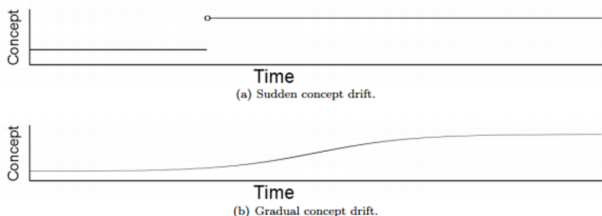
What do people do when the concept drifts ?



Assume that the concept drift can be successfully detected,

- forget the data before the new concept in learning algorithms;

What do people do when the concept drifts ?



Assume that the concept drift can be successfully detected,

- forget the data before the new concept in learning algorithms;
- change a new learning algorithm (strategy)

Stability and adaptivity when forgetting data

Definition

A learning algorithm \mathcal{A} has error stability β_n with respect to the loss function l if

$$\forall Z_n \in \mathcal{X}^n, \forall i \in \{1, \dots, n\} |\mathbb{E}\{l(\mathcal{A}_{Z_n})\} - \mathbb{E}\{l(\mathcal{A}_{Z_n}^{-i})\}|$$

Examples

for k-NN, SVM, support vector regression and ridge regression, $\beta_n = \mathcal{O}(\frac{1}{n})$

Detect concept drift

- Essentially, the detection has been conducted via hypothesis test sequentially.

$$I(\mathcal{A}_{\{\mathbf{x}_1, \dots, \mathbf{x}_{n-t}\}}, \underbrace{\{\mathbf{x}_{n-t+1}, \dots, \mathbf{x}_n\}}_{\text{time window}}) \implies$$

null hypothesis v.s. alternative hypothesis

- Can work well with good hyper-parameter settings.

Temporal-spatial kernel

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- Regret bound:

$$\Pr \left\{ \mathbf{R}^{(T)*} \leq \sqrt{\frac{8}{\log(1 + \sigma^{-1})}} T^{\beta(T)} \hat{\gamma}^{(T)} \right\} \geq 1 - \delta \quad (29)$$

Concept drift in expert committee

① A expert committee

- an expert which uses some side information;
- an expert which uses short-memory of instances;
- an expert which uses whole sequence of instances;
- an expert which can detect concept drifts and forget old data;

② concept drift == best expert shift

Adaptive regret and Fixed-share algorithm

- regular regret:

$$R^{(t)*} = \max_{n \in [1, N]} R_n^{(t)} = \hat{L}^{(t)} - \min_{n \in [1, N]} L_n^{(t)}$$

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- weight-share algorithm:

Fixed-share algorithm

Weight update if one expert E_n predicts wrongly, decrease its weight

$$\hat{w}_n^{(t+1)} = w_n^{(t)} \exp(-\eta m_n^{(t)}) \quad (30)$$

followed by a weight-share update:

$$w_n^{(t+1)} = (1 - \alpha) \hat{w}_n^{(t+1)} + \frac{\alpha}{N - 1} \sum_{m \neq n} \hat{w}_m^{(t+1)} \quad (31)$$

New item proposal plans

Expert committee interpretation

by considering each expert is a sku proposer,

- fully observed feedback for all experts;
- can flexibly define cost function ;
- can flexibly add diverse experts;
- regret bound complexity: $\mathcal{O}(\sqrt{T \ln N})$
- clear understanding of performance in non-stationary environments;
- can be computational expensive by maintaining many experts;
- there exist some gap between **sequential proposal** and **search**

Bandit interpretation within ensemble learning

- partially observed feedback only on the selected sku;
- no explicit loss function;
- non-trivial analysis of the information gains for different types of functions;
- regret bound complexity: $\mathcal{O}\left(\sqrt{T(\log T)^{d+1}}\right)$
- difficult to analyze information gain in non-stationary environments;
- GP complexity $\mathcal{O}(n^3)$; also can be computational expensive by maintaining many models;
- it stands closer to the nature of the *Search* task;

Connections

if $\kappa = 0$ for all t , then **Bandit setting == Expert committee setting** in the aggregation policy:

A reminder

$$\hat{y}^{(t)} = \sum_{n=1}^N p_n^{(t-1)} f_n^{(t)}$$

Plans and outlooks

Under expert committee interpretation framework:

- define more diverse proposal experts (strategies);
- design an informative loss function;
 - ① loss function for classifiers

$$m_n^{(t)} = \exp^{-p_n^{(t)}(y^{(t)})}$$

where $y^{(n)} = \{-1, 1\}$ and $p_n^{(t)}$ denotes the classification probability of E_n

- ② loss function for regressor (ranker)

??

Plans and outlooks

Under expert committee interpretation framework:

- define more diverse proposal experts (strategies);
- design an informative loss function;
 - 1 loss function for classifiers

$$m_n^{(t)} = \exp^{-p_n^{(t)}(y^{(t)})}$$

where $y^{(n)} = \{-1, 1\}$ and $p_n^{(t)}$ denotes the classification probability of E_n

- 2 loss function for regressor (ranker)

??

- enable some experts with concept drift detection and data forgetting mechanism;
- track the best expert using weight-share algorithm

Quantitative evaluation

- Expert committee interpretation:
cumulative loss of the forecaster or regret;
- Bandit interpretation:
use pseudo targets (wishlist);

[References 1](#): Nicolo Cesa-Bianchi and Gabor Lugosi. 2006. Prediction, Learning, and Games. Cambridge University Press, New York, NY, USA.

[References 2](#): Online Learning: A Comprehensive Survey

The End