

Flink Checkpoint-轻量级分布式快照

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Checkpoint的执行机制

Checkpoint与State

二者的关系



Checkpoint与State

Checkpoint,在Flink中是一个执行操作,最终产生的结果作为分布式快照提供容错机制。

| Subtasks Task Metrics Wat | termarks Accumulators | Checkpoints | Back Pressure | | |
|-----------------------------|------------------------|-----------------|--------------------------|---------------------|----------------|
| Overview History Summa | ary Configuration | | | | |
| Checkpoint Counts | Triggered: 569027 In P | rogress: 0 Comp | leted: 569027 Failed: 0 | Restored: 0 | |
| Latest Completed Checkpoint | ID: 569027 Completion | Time: 18:19:02 | End to End Duration: 3ms | State Size: 9.32 KB | ▶ More details |
| Latest Failed Checkpoint | None | | | | |
| _atest Savepoint | None | | | | |
| _atest Restore | None | | | | |
| | | | | | |



Checkpoint与State

State是构成checkpoint的数据构成

| Overvie | w History | / Summary | Configura | ntion Details for Ch | eckpoint 569116 | | | | |
|---------|-------------|--------------|-----------------|---------------------------|------------------------|---------------|------------------------------|-----------|---|
| ID | Status | Acknowledged | Trigger Time | Latest Acknowledgement | End to End Duration | State Size | Buffered During Alignment | Discarded | Path |
| 569116 | ✓ Completed | 2/2 (100%) | 18:22:00 | 18:22:00 | 2ms | 9.17 KB | 0 B | Yes | <pre><checkpoint-not-externally- addressable=""></checkpoint-not-externally-></pre> |

Operators

| Name | Acknowledged | Latest Acknowledgment | End to End Duration | State Size | Buffered During Alignment | |
|-------------------------------------|--------------|--------------------------|------------------------|---------------|------------------------------|-----------------|
| Source: Custom Source | 1/1 (100%) | 18:22:00 | 2ms | 0 B | 0 B | Show Subtasks ✔ |
| Flat Map -> Sink: Print to Std. Out | 1/1 (100%) | 18:22:00 | 2ms | 9.17 KB | 0 B | Show Subtasks ✔ |



```
env.socketTextStream("localhost",9000)
    // split up the lines in pairs (2-tuples) containing: (word,1)
    .flatMap(new Tokenizer())
    // group by the tuple field "0" and sum up tuple field "1"
    .keyBy(0).sum(1)
    .print();
```

- 1. 执行上述代码
- 2. 本地启动netcat
- 3. 键入

hello world

\$ nc -lk 9000

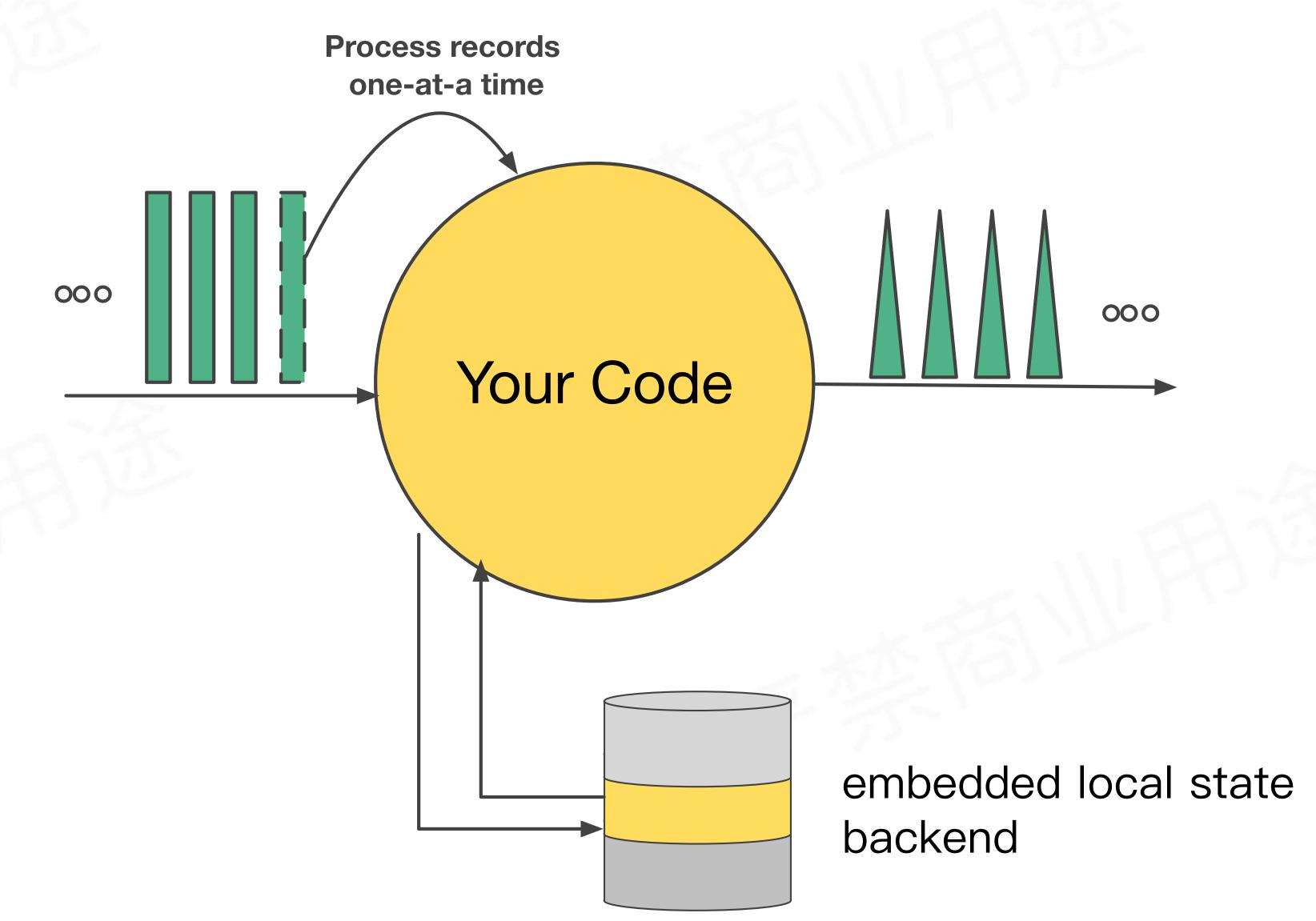
执行程序会输出什么?

再次键入 hello world

执行程序会输出什么?



State:流式计算中持久化了的状态



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Keyed State

- 只能应用于 KeyedStream 的函数与操作中,例如Keyed UDF, window state
- Keyed State是已经分区/划分好的,每一个key只能属于某一个keyed state

什么是state Keyed State



再来看这 段 word count代码

```
env.socketTextStream("localhost",9000)
    // split up the lines in pairs (2-tuples) containing: (word,1)
    .flatMap(new Tokenizer())
    // group by the tuple field "0" and sum up tuple field "1"
    .keyBy(0) sum(1)
    .print();
```

创建 KeyedStream (对key进行了划分, 不同task上不会出 现相同的key) 调用内置的 StreamGroupedReduce UDF

什么是state Keyed State



再来看这段 word count代码

```
env.socketTextStream("localhost",9000)
   // split up the lines in pairs (2-tuples) containing: (word,1)
    .flatMap(new Tokenizer())
   // aroup by the tuple field "0" and sum up tuple field "1"
   .keyBy(0).sum(1)
    .print();
 hello, 1
              world, 1
                                       apache, 1
              flink, 1
                                                     China, 1
              China, 1
                                                    world, 1
                                       hello,
```

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Operator State

- 又称为non-keyed state,每一个operator state都仅与一个operator的实例绑定。
- 常见的operator state是source state,例如记录当前source的offset



什么是state Operator State

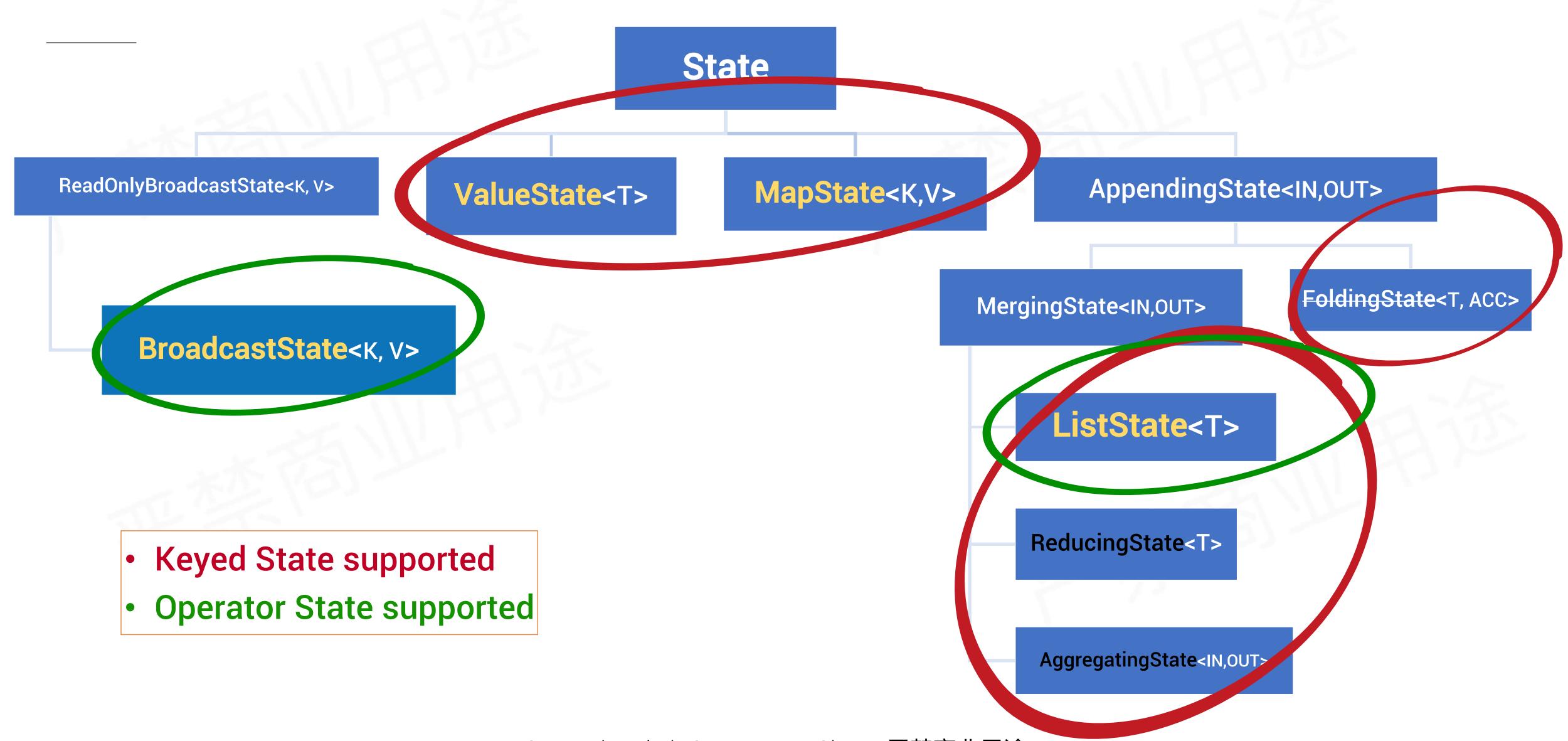
再来看一段 word count代码

```
env.fromElements(/ordCountData.WORDS)

// split up the lines in pairs (2-tuples) containing: (word,1)
.flatMap(new Tokenizer())

// group by the tuple field "0" and sum up tuple field "1"
.keyBy(0).sum(1)
.print();
```





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- Managed State:由Flink管理的state,刚才举例的所有state均是managed
- Raw State: Flink仅提供stream可以进行存储数据,对其而言 只是一些bytes

如何在Flink中使用state

使用指南

如何在Flink中使用state: Keyed State



再来看这段 word count代码

```
env.fromElements(WordCountData.WORDS)

// split up the lines in pairs (2-tuples) containing: (word,1)
    .flatMap(new Tokenizer())

// group by the tuple field "0" and sum up tuple field "1"
    .keyBy(0)
    .sum(1)
    .print();
```

创建KeyedStream (对key进行了划分, 不同task上不会出现相 同的key)

调用内置的StreamGroupedReduce



如何在Flink中使用state: Keyed State

```
public class StreamGroupedReduce<IN> extends AbstractUdfStreamOperator<IN, ReduceFunction<IN>>
       implements OneInputStreamOperator<IN, IN> {
   private transient ValueState<IN> values;
   @Override
   public void open() throws Exception {
       super.open();
       ValueStateDescriptor<IN> stateId = new ValueStateDescriptor<>(STATE_NAME, serializer);
       values = getRuntimeContext().getState(stateId);
                                                                 通过 RuntimeContext 访问state
   @Override
   public void processElement(StreamRecord<IN> element) throws Exception {
       IN value = element.getValue();
       IN currentValue = values.value();
                                                     访问和修改当前key对应的state数值
       if (currentValue != null) {
           IN reduced = userFunction.reduce(currentValue, value);
           values.update(reduced);
           output.collect(element.replace(reduced));
       } else {
           values.update(value);
           output.collect(element.replace(value));
                                e Abache i illik community china / <del>Mahith</del>is
```

如何在Flink中使用state: Operator State



再来看这段 word count代码

```
env.fromElements WordCountData.WORDS)

// split up the lines in pairs (2-tuples) containing: (word,1)
.flatMap(new Tokenizer())

// group by the tuple field "0" and sum up tuple field "1"
.keyBy(0).sum(1)
.print();
```

调用内置的FromElementsFunction



如何在Flink中使用state: Operator State

```
class FromElementsFunction< T > implements SourceFunction <math>< T > I CheckpointedFunction < T > I
private transient ListState<Integer> checkpointedState;
                                                                     通过 FunctionInitializationContext 访问state
@Override
public void initializeState(FunctionInitializationContext context) throws Exception {
   Preconditions.checkState(this.checkpointedState == null,
        "The " + getClass().getSimpleName() + " has already been initialized.");
    this.checkpointedState = context.getOperatorStateStore().getListState(
        new ListStateDescriptor<>("from-elements-state", IntSerializer.INSTANCE)
   );
   if (context.isRestored()) {
        List<Integer> retrievedStates = new ArrayList<>();
        for (Integer entry : this.checkpointedState.get()) {
            retrievedStates.add(entry);
        // given that the parallelism of the function is 1, we can only have 1 state
        Preconditions.checkArgument(retrievedStates.size() == 1, getClass().getSimpleName() + " retrieved invalid state.");
        this.numElementsToSkip = retrievedStates.get(0);
                                                                在snapshotState时将状态数据存储到state中
@0verride
public void snapshotState(FunctionSnapshotContext context) throws Exception {
    Preconditions.checkState(this.checkpointedState != null,
        "The " + getClass().getSimpleName() + " has not been properly initialized.");
   this.checkpointedState.clear();
    this.checkpointedState.add(this.numElementsEmitted);
```


Checkpoint的执行机制

Checkpoint internal



state-backend 分类:

MemoryStateBackend

FsStateBackend

RocksDBStateBackend

Operator state-backend 分类:

Keyed state-backend 分类:

DefaultOperatorStateBackend

HeapKeyedStateBackend

RocksDBKeyedStateBacken d

StreamExecutionEnvironment env = StreamExecutionEnvironment.getExecutionEnvironment();
env.setStateBackend(new FsStateBackend("hdfs://namenode:40010/flink/checkpoints"));



Checkpoint数据直接返回给 master节点

Checkpoint数据写入文件中,将 文件路径传递给master

Checkpoint数据写入文件中,将 文件路径传递给master

state-backend 分类:

MemoryStateBackend

FsStateBackend

RocksDBStateBackend

Operator state-backend 分类:

Keyed state-backend 分类:

数据存储在内存中

DefaultOperatorStateBackend

HeapKeyedStateBackend

RocksDBKeyedStateBacken d

数据存储在RocksDB中



HeapKeyedStateBackend 存储格式

- 支持异步checkpoint (默认): CopyOnWriteStateTable<K, N, S>[], 整体相当于一个map
- 仅支持同步checkpoint: Map<N, Map<K, S>>[] , 由嵌套map的数组构成

• 在MemoryStateBackend内使用时,checkpoint序列化数据阶段默认有最大5MB数据的限制

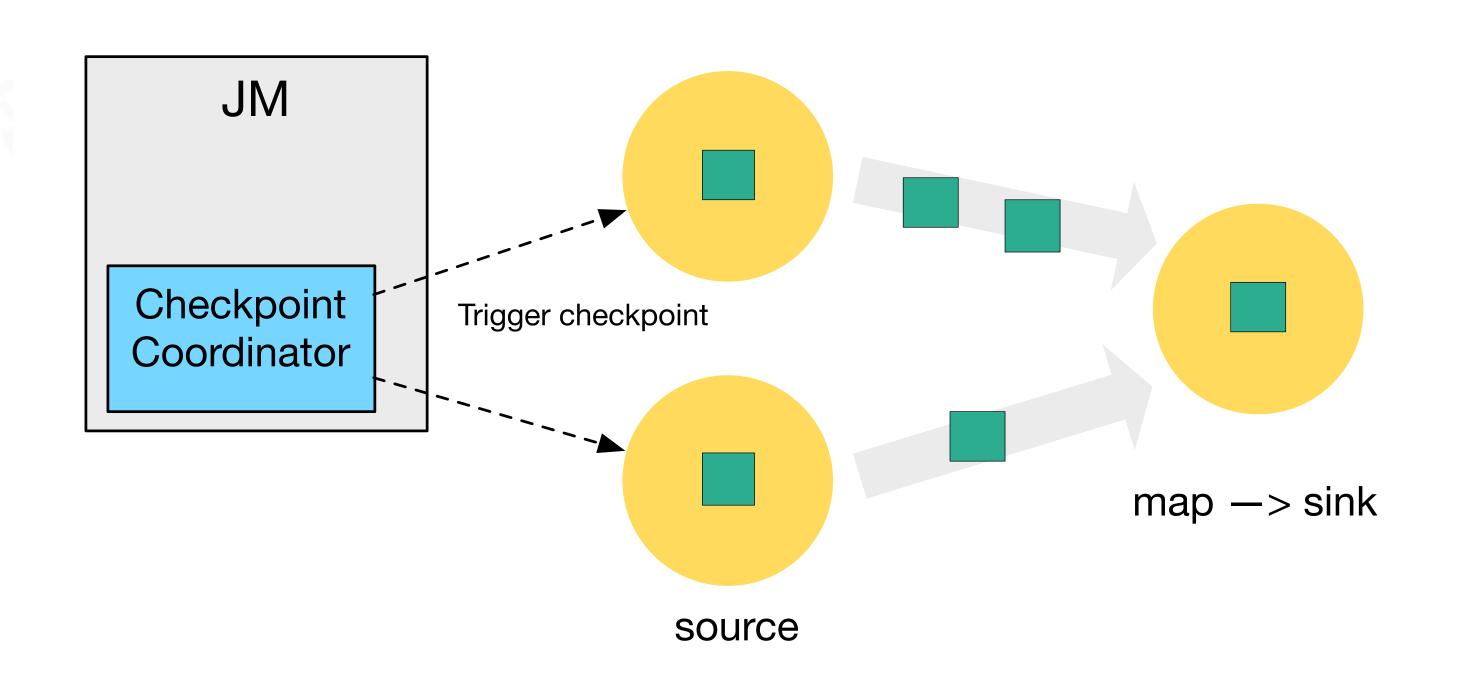


RocksDBKeyedStateBackend 存储格式

每个state都存储在一个单独的column family内 keyGroup, Key和Namespace进行序列化存储在DB作为key

| State1 | | State2 | | | |
|----------------------------|-------|----------------------------|-------|--|--|
| KeyGroup + Key + Namespace | value | KeyGroup + Key + Namespace | value | | |
| (1, K1, Window(10, 20)) | v1 | (2, K2, Window(10, 20)) | v2 | | |
| (1, K3, Window(10, 20)) | v3 | (2, K4, Window(10, 25)) | v4 | | |
| === | | | | | |

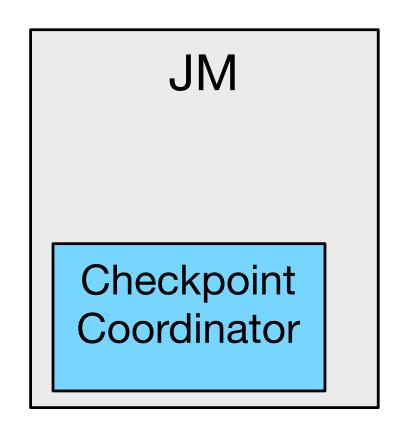


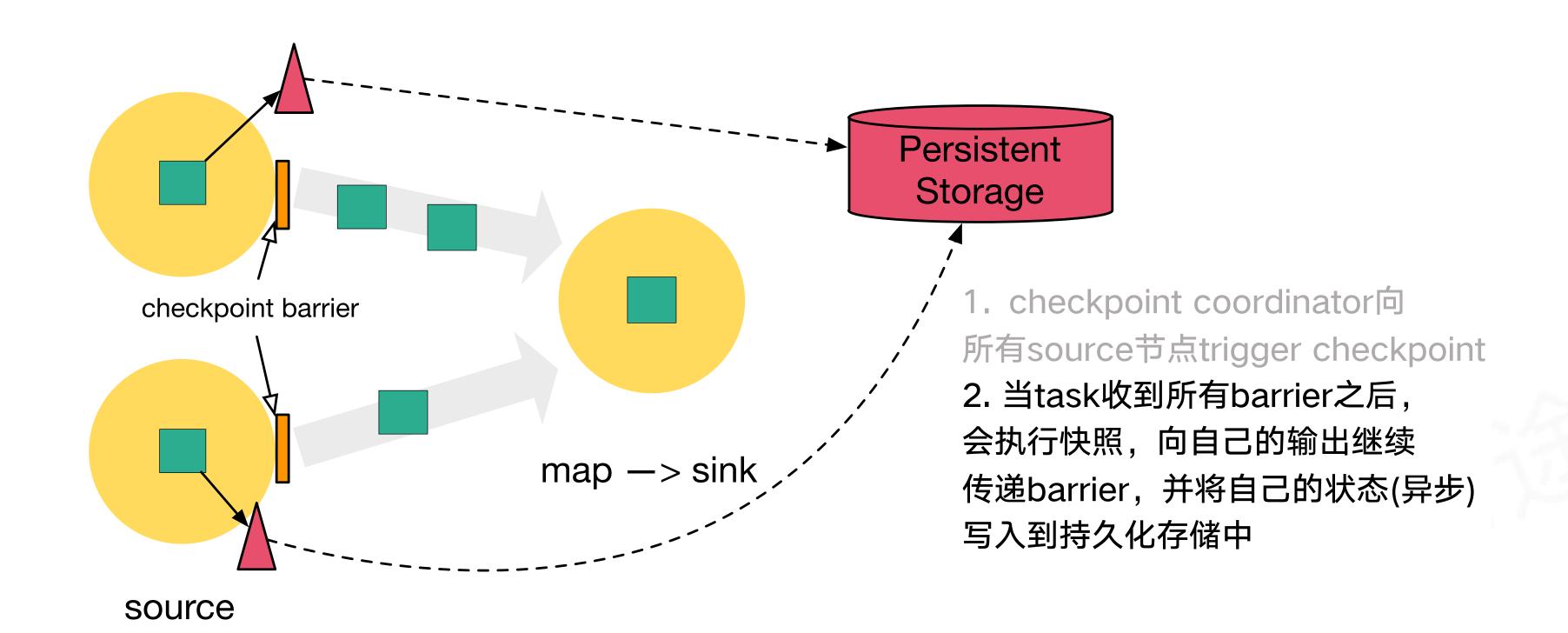


Persistent Storage

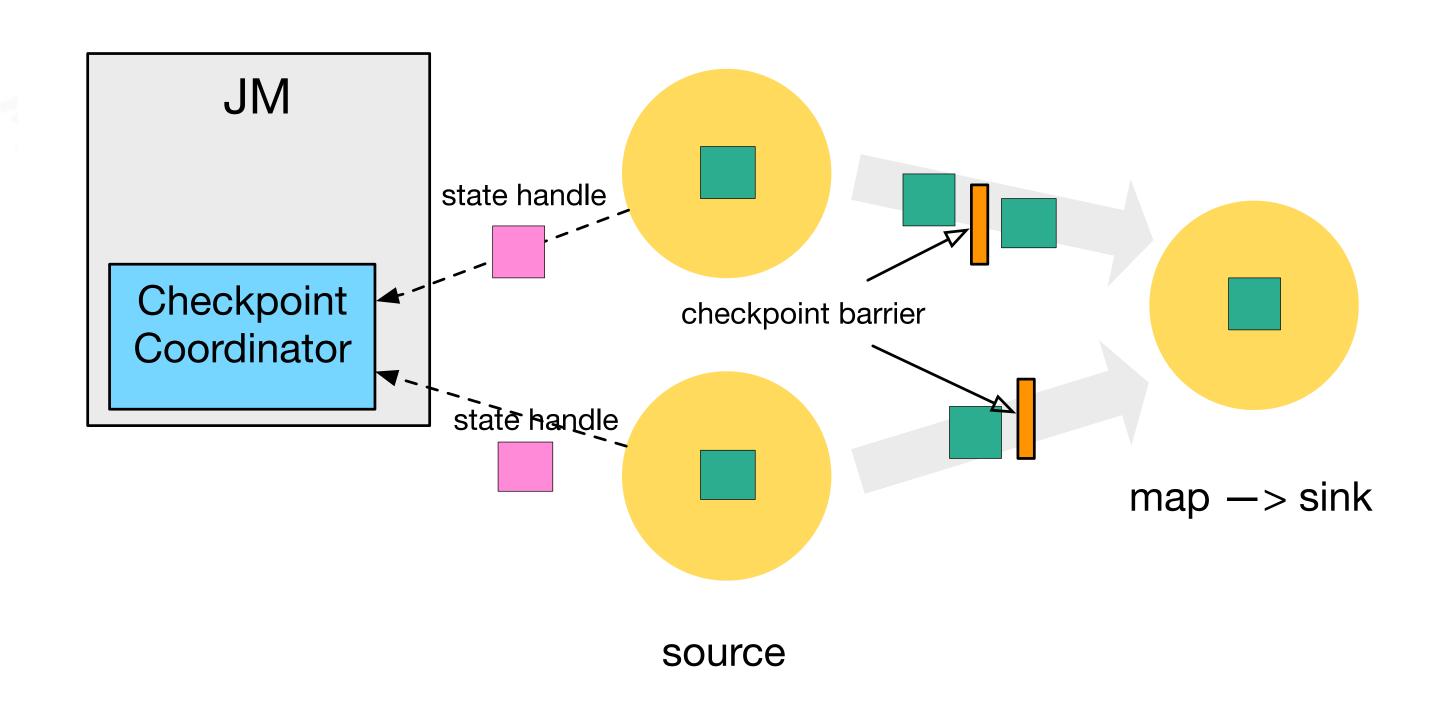
1. checkpoint coordinator 向所有source节点trigger checkpoint







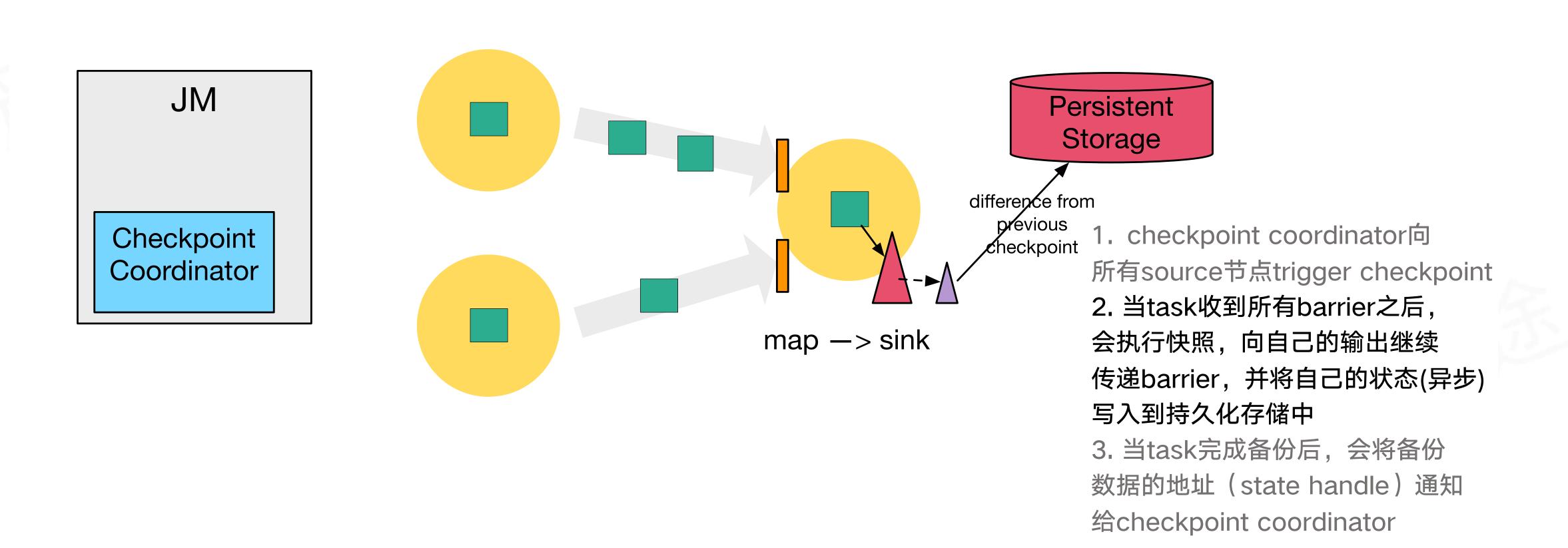




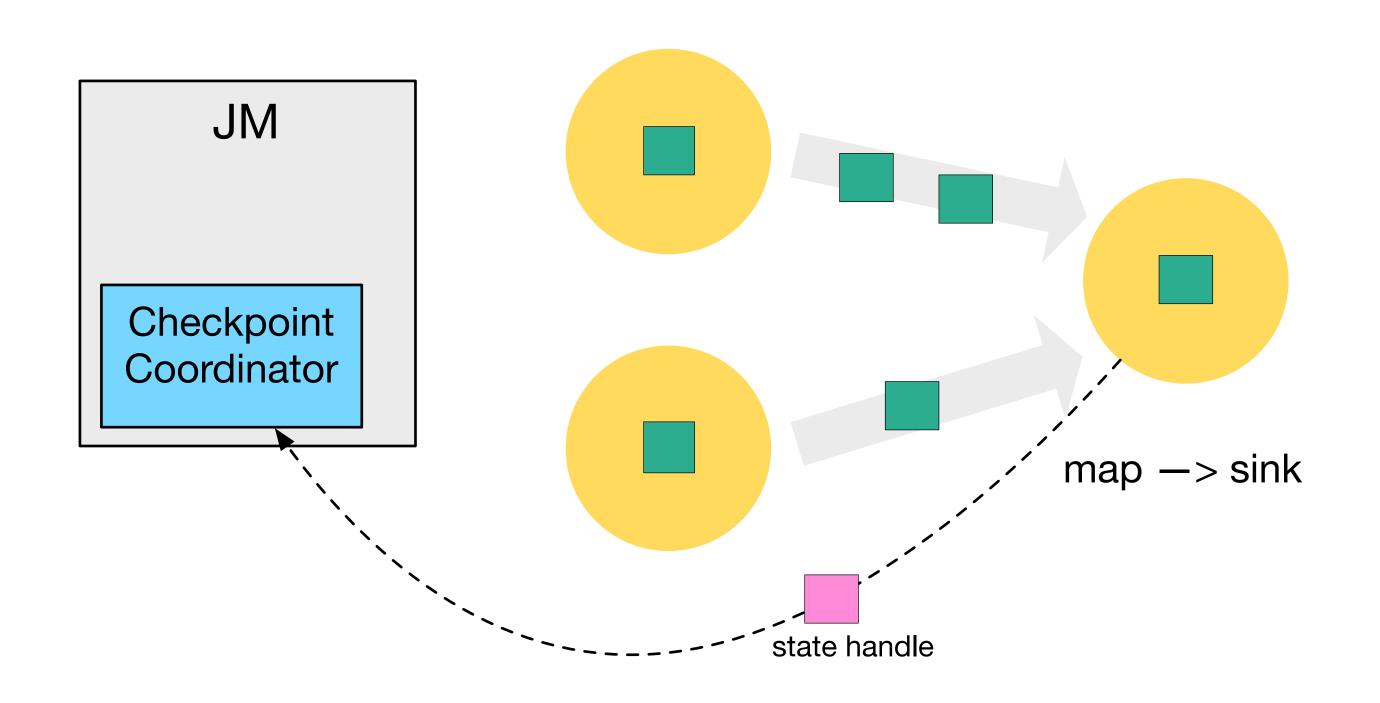
Persistent Storage

- 1. checkpoint coordinator向 所有source节点trigger checkpoint
- 2. 当task收到所有barrier之后, 会执行快照,向自己的输出继续 传递barrier,并将自己的状态(异步) 写入到持久化存储中
- 3. 当task完成备份后,会将备份数据的 地址(state handle)通知 给checkpoint coordinator





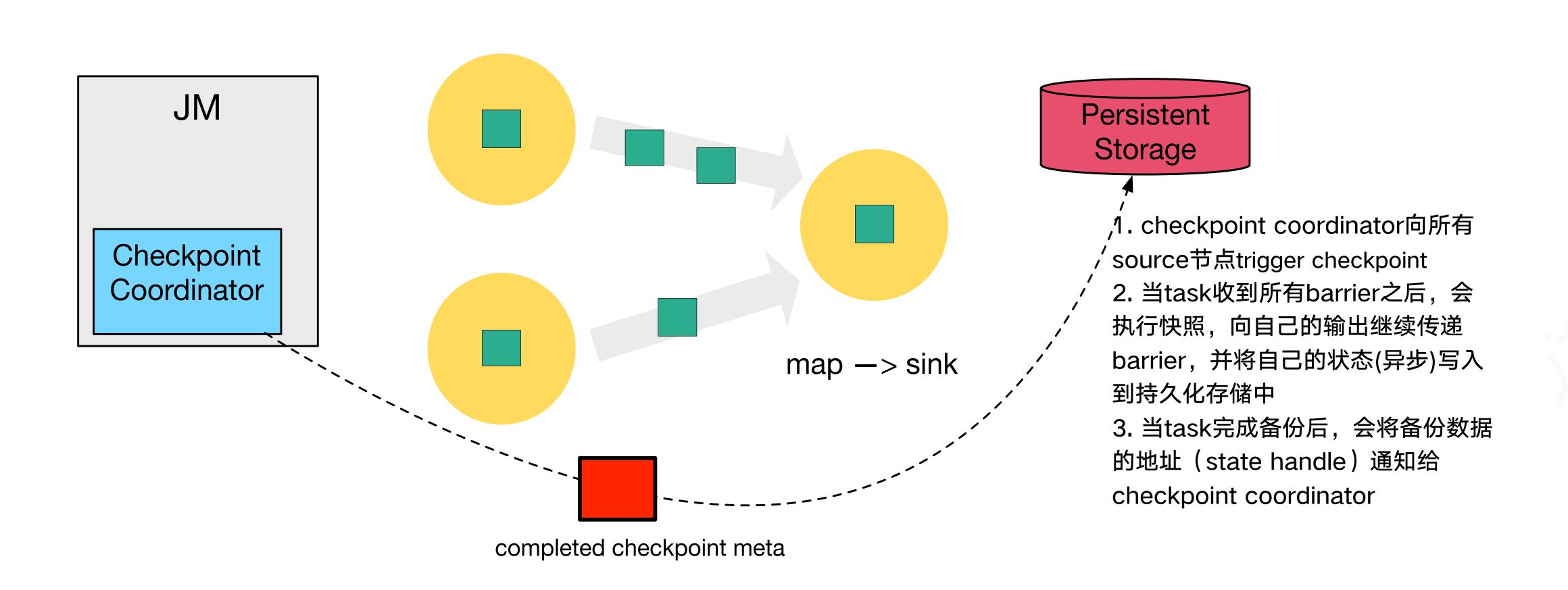




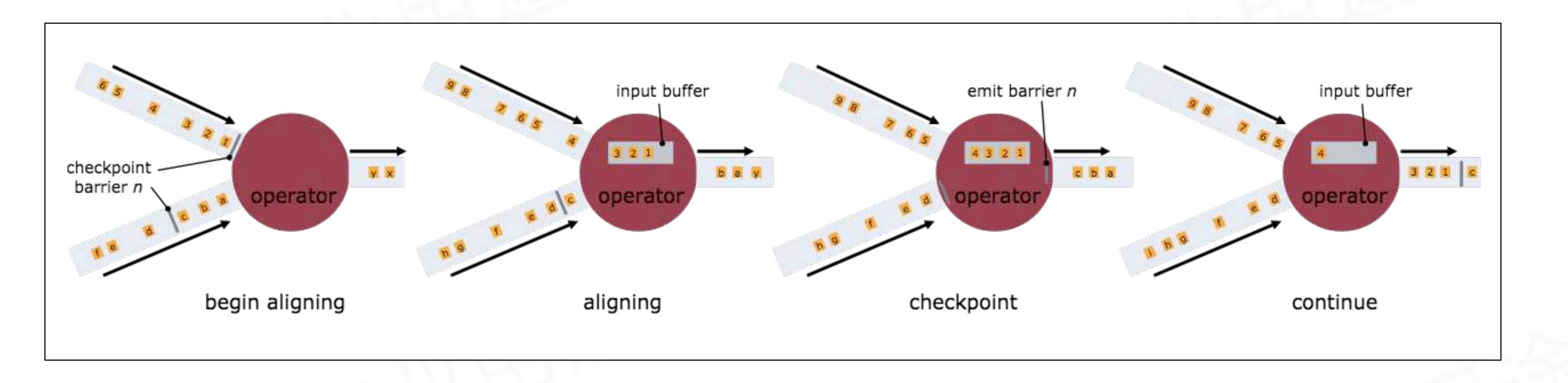
Persistent Storage

- 1. checkpoint coordinator向 所有source节点trigger checkpoint
- 2. 当task收到所有barrier之后, 会执行快照,向自己的输出继续 传递barrier,并将自己的状态(异步) 写入到持久化存储中
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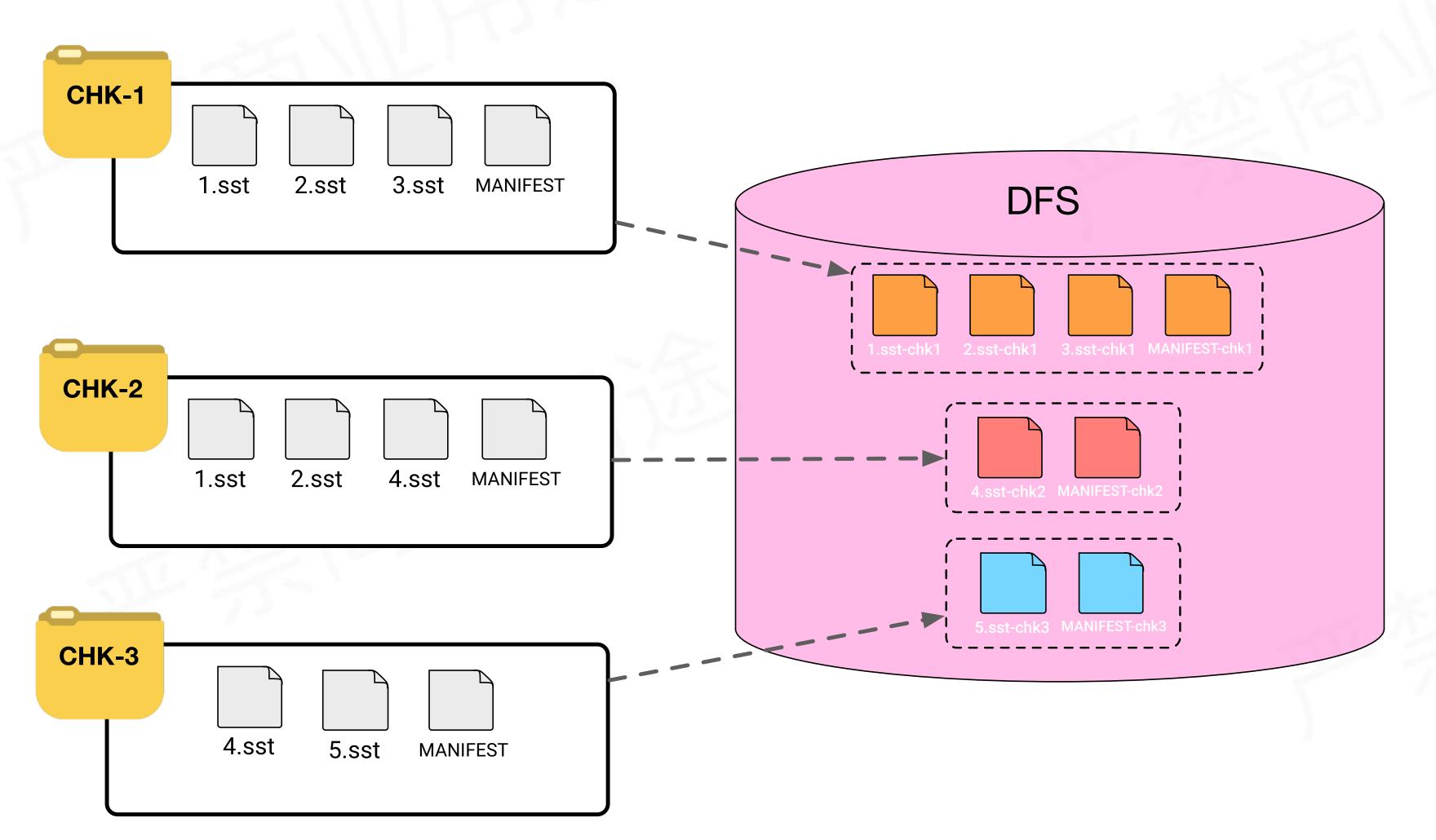




- 为了实现EXACTLY ONCE语义,Flink通过一个input buffer将在对齐阶段收到的数据缓存起来, 等对齐完成之后再进行处理。
- 对于AT LEAST ONCE语义,无需缓存收集到的数据,会对后续直接处理,所以导致restore时,数据可能会被多次处理。
- Flink的checkpoint机制只能保证Flink的计算过程可以做到EXACTLY ONCE, end-to-end的 EXACTLY ONCE需要source和sink支持

基于RocksDB的增量checkpoint





- 1. 本地snapshot目录创建当前DB 内容的备份
- 2. 与上一次成功的checkpoint本地sst文件列表比对,将不在其中的文件上传到DFS中
- 3. 所有文件都会重命名防止冲突
- 4. 包含了所有新旧文件的 handle返回给checkpoint coordinator

从已停止的作业进行状态恢复



Savepoint

用户通过命令触发,由用户管理其创建与删除

标准化格式存储,允许作业升级或者 配置变更

用户在恢复时需要提供用于恢复作业 状态的savepoint路径

Externalized Checkpoint

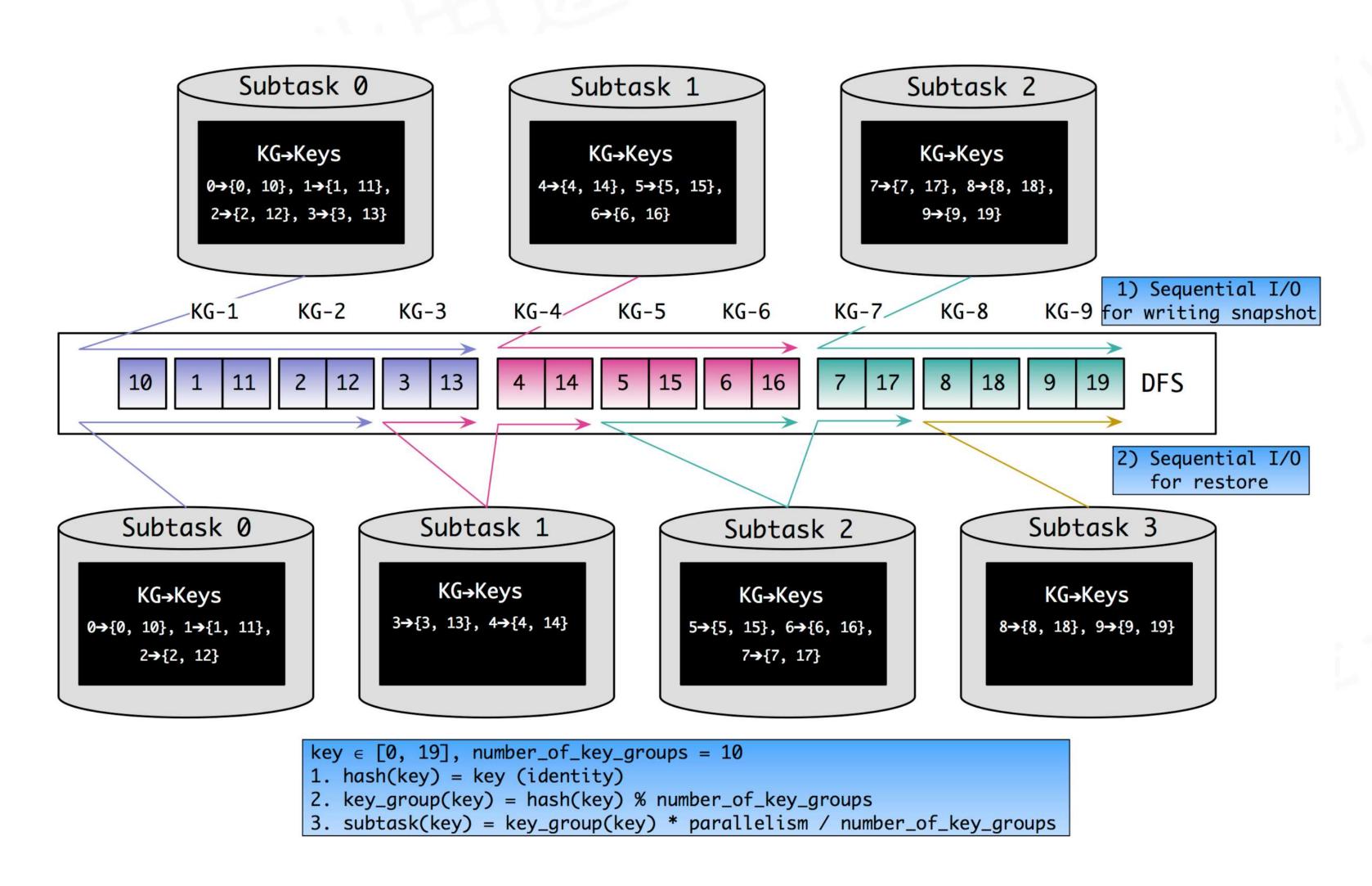
Checkpoint完成时,在用户给定的外部持久化存储保存

当作业FAILED(或者CANCELED)时,外部存储的checkpoint会保留下来

用户在恢复时需要提供用于恢复的作业状态的checkpoint路径

从已停止的作业进行状态恢复 Keyed State的改并发

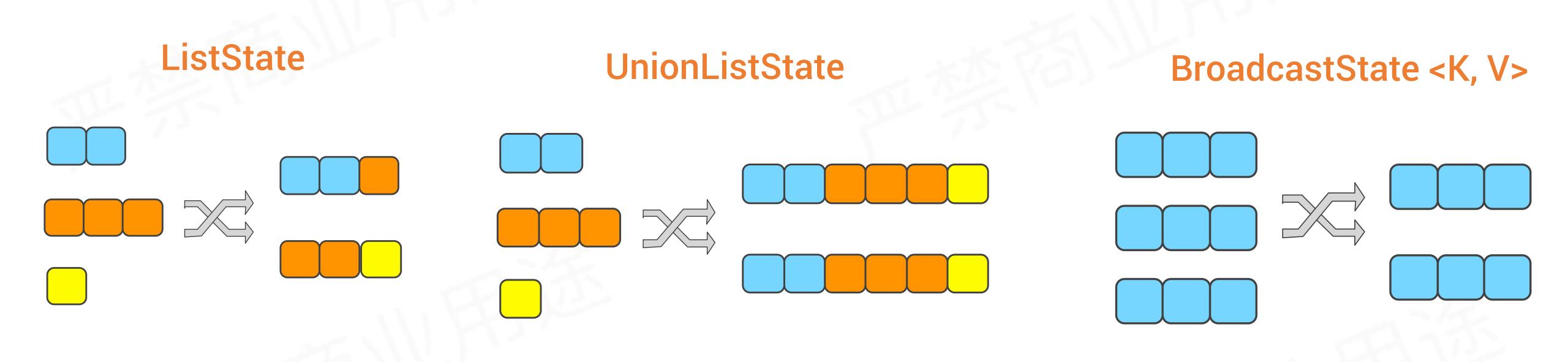




并发度 3 → 4 利用KeyGroup改并发

从已停止的作业进行状态恢复 Operator State的改并发





所有task的元素均匀划分 给新的task (Even-split redistribution) 所有task的元素全部划分给新的task (Union redistribution) 所有task的state相同, 改并发时,新的task获 得state的一个备份





