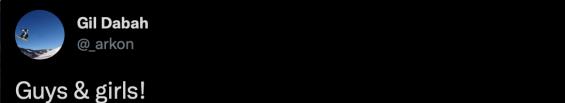
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Win32k: Smash The Ref

Gil Dabah CEO, Piiano



Exactly a year ago I promised over 15 bugs in win32k. You're welcome to read and find out about my biggest research so far: #win32k #SmashTheRef bug class - github.com/gdabah/win32k-...

Check out the paper and the POCs, there are some crazy stuff going on. Promise!

gdabah/win32kbugs



Dump of win32k POCs for bugs I've found

Issues

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Contributor

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339 Stars

Forks

github.com

win32k-bugs/SmashTheRef at master \cdot gdabah/win32k-bugs

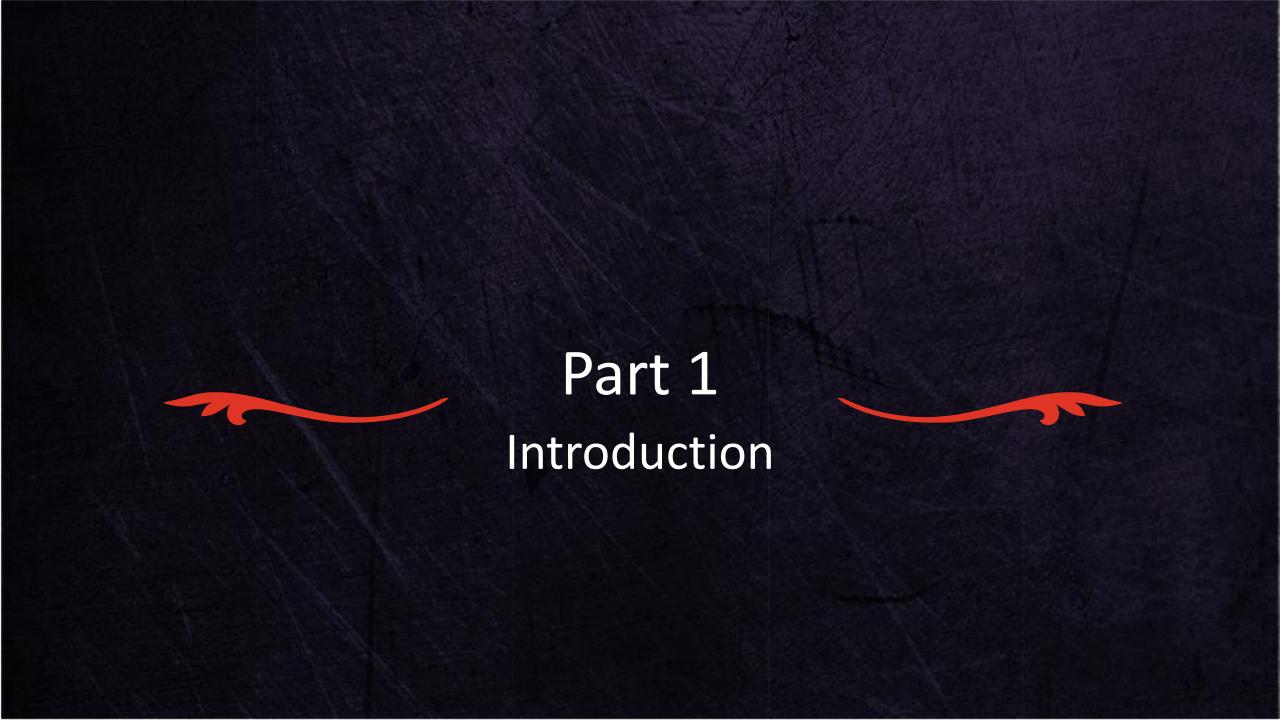
Dump of win32k POCs for bugs I've found. Contribute to gdabah/win32k-bugs development by creating an account on GitHub.

Agenda

4 parts:

- Introduction classic win32k bugs
- Zombies new attack & techniques
- Exploitation doing the impossible
- Wrap up the real cool stuff





Win32K

- GUI done in kernel
- UI objects are managed in kernel
- What can go wrong?
- Good Performance

U->K instead of U->U

Bad - Security

From K->U for helpers

Ugly – the code quality

Kernel -> User Callbacks

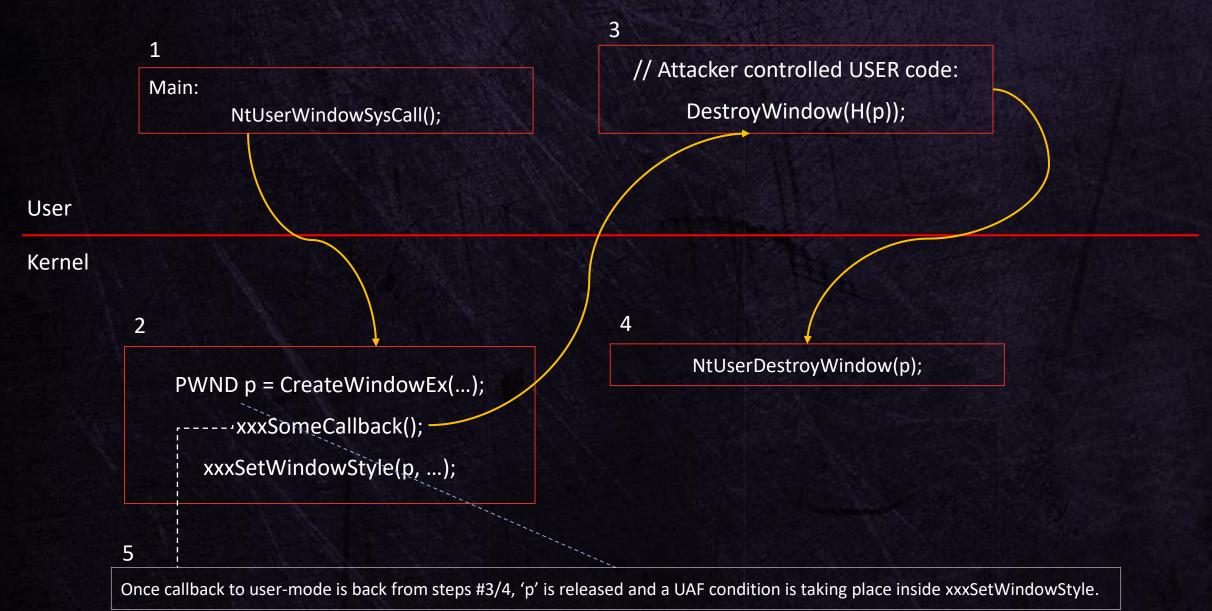
- Win32k calls back to user for
 - window events
 - hooks
 - notification events
 - helper functions
- Note 'xxx' prefix function names for leaving kernel->user
 - (aka 'usercrit')
 - xxxSendMessage
 - xxxCreateWindow
 - xxxDestroyWindow
 - & more

Abusing Callbacks — Introducing use-after-free

Kernel pseudo code snippet:

```
NtUserWindowSysCall()
    PWND p = xxxCreateWindowEx(...);
    xxxSomeCallback();
    xxxSetWindowStyle(p);
Where's the bug?
```

Trigger UAF Flow

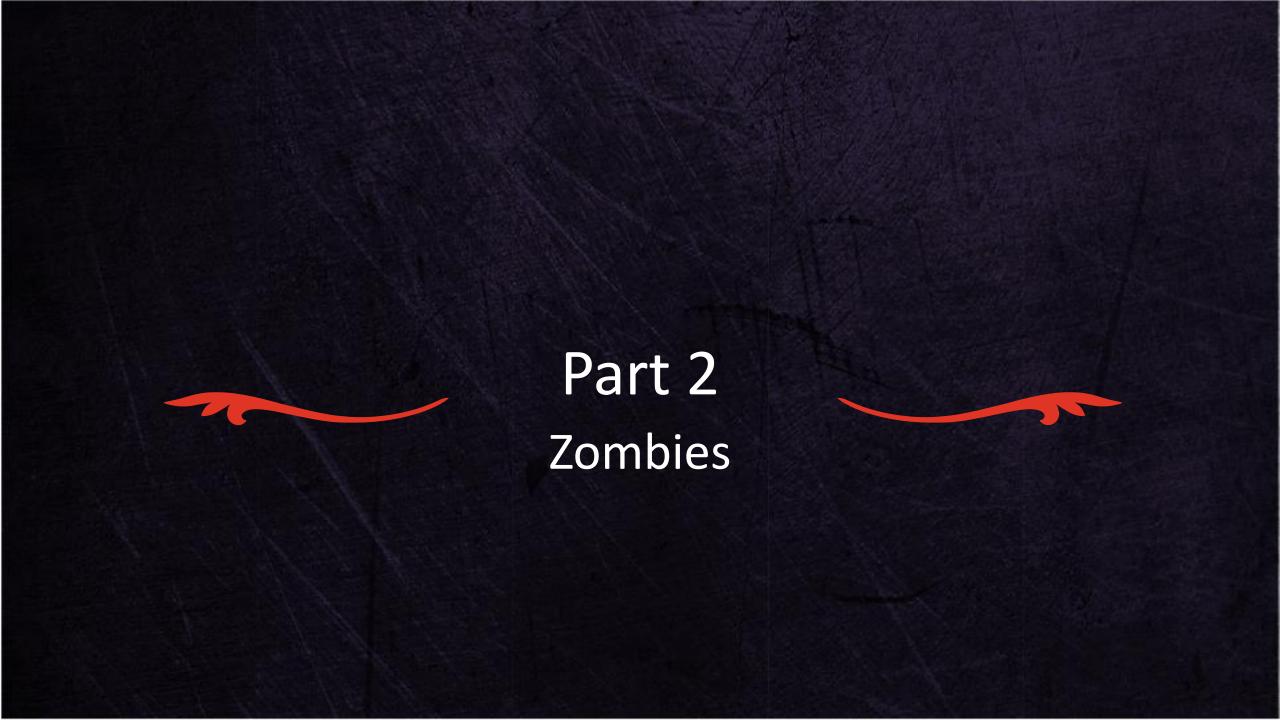


Kernel UI Objects Locking Mechanism

- Create/Destroy syscalls for objects
- Objects start with 0 ref count
- Destroy functions free objects only if 0 refs!
- ThreadLock/ThreadUnlock local scope for locking
 - Blocks usermode from killing objects (hmmm)

Introducing The Fix — Thread Scope Locking

```
PWND p = xxxCreateWindowEx(...);
ThreadLock(p);
xxxSomeCallback();
// Zombie window is used next, immune!
xxxSetWindowStyle(p);
ThreadUnlock();
```



Introducing Zombie Objects

Destroying an object with >0 refs???

```
PMENU p = CreateMenu();
ThreadLock(p);
xxxSomeCallback();
//'p' menu is now a zombie!
ThreadUnlock(p);
// Really frees 'p' inside
decrement ref - last ref goes from 1 to 0!
```

Zombie Disorders

Destroyed objects that are still referenced are called zombies! (otherwise, they're freed)

Zombies still exist (allocated) in kernel (otherwise bad things will happen)

Zombies cannot be accessed from usermode

Each object-destruction function behaves differently

Pretty much cannot be used in this state, only waiting to be freed

DestroyWindow API

- DestroyWindow releases all linked objects immediately
 - Menus, carets, child-windows, timers, etc

- xxxDestroyWindow is first called from user mode
- Then again (a second time!) when last ref is gone side effects?!

• Remember - like OOP:

When a destructor is called, it calls the destructors of its members!

Smashable Site - Recipe

- 1. A dumb pointer that's assigned to a UI object
- 2. A window object's last decref functionality (ThreadUnlock & friends)
- 3. Control where the last decref happens
- 4. Use-after-free of #1's pointer
- 5. A side effect behavior that can be reloaded unto a zombie window

xxxDestroyWindow – Releasing Sub-resources

```
void xxxDestroyWindow(PWND pwnd)
 xxxFW DestroyAllChildren(); // Destroy child windows, if exist!
 if (NULL != pwnd->spmenu) // If there's a menu, remove and destroy it
      DestroyMenu (pwnd->spmenu);
  if (pwnd == ptiCurrent->pq->capturedPwnd)
    xxxReleaseCapture();
 DereferenceClass(pwnd);
  if (HMMarkObjectDestroy(pwnd)) // Check for zero refs!
    HmFreeObject(pwnd); // Only now free the object
```

A Real-Life Example: xxxMnOpenHierarchy

```
ThreadLock(pwndParent);
// Notice there's a parent-child relation :)
pwnd = xxxCreateWindow(..., pwndParent, ...);
ThreadUnlock();
if (NULL == pwnd) return ...
// From now on the window exists and can be used.
xxxSendMessage(pwnd...)
```

We got a UAF, give me a yeahhh!

Killing in The Stack Of — The Secret Sauce

xxxDestroyWindow

xxxDestroyWindowIfSupported HMDestroyUnlockedObject

HMUnlockObjectInternal

HMAssignmentUnlock ←

xxxFreeWindow

xxxDestroyWindow

xxxDestroyWindowIfSupported

HMDestroyUnlockedObject

ThreadUnlock1

xxxMNOpenHierarchy *

Chain-effect:

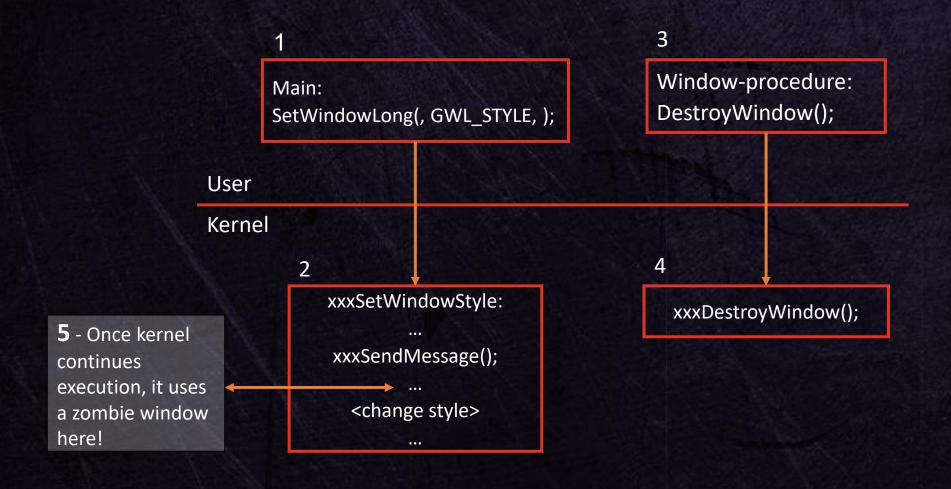
ThreadUnlock kills parent window first and then kills the child window too!

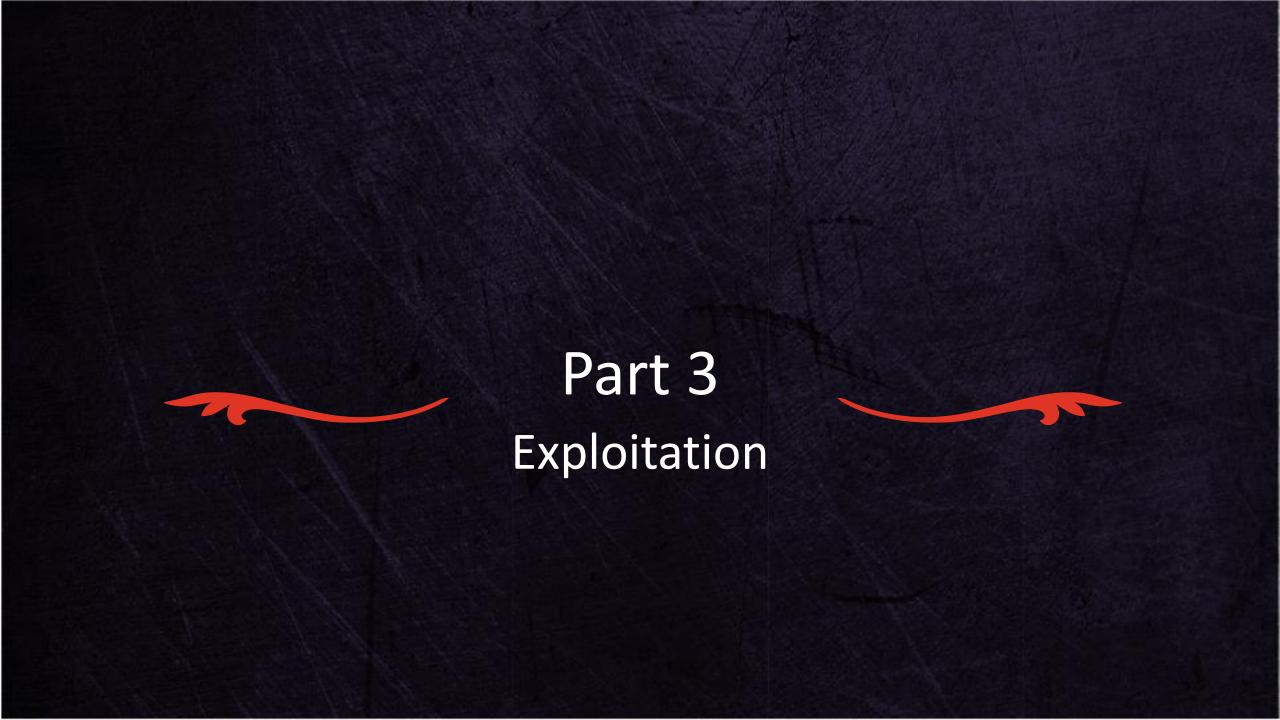
Reloading a Zombie-woOT??

But if an object is destroyed – how can we link sub-resources to it?

We **stack** the operation *before* it's destroyed, but (happening) *after* it's destroyed

Reloading a Zombie-A Simple Example





UAF Exploitation Problems

Race Condition

- Destruction doesn't leave (the usercrit) to umode!!!!
- NTGDI
- Pool type allocations
- kLFH

TypeIsolation Objects (allocator pools per object type)

SmartObjectStackRef OOP

Let's Go Back To The 90's

By calling back to usermode whenever WE want

Then we can exploit the freed object without a race

But how???

The Ultimate Reloading!

```
PWND xxxCreateWindowEx (...)
  PVOID ptr =
xxxClientAllocWindowClassExtraBytes();
  if (NULL == ptr) // error handling
  pwnd->extraWindowUMPtr = ptr;
  ...
  xxxCallHook(WH CBT, ..);
  if (windowIsDestroyed(pwnd)) return NULL;
```

Reload 'Back To User-Mode'

```
3
                                                        user32!
       WNDCLASS wc = \{0\};
                                             xxxClientAllocWindowClassExtraBytes
       wc.cbWndExtra = 100;
                                                       hook stub:
                                                                                          CBTHookProc(...):
       RegisterClass(&wc);
CreateWindow(wc.lpszClassName, ...);
                                                    DestroyWindow();
User
Kernel
                                             xxxCreateWindowEx:
                                                PVOID UMPtr =
                               xxxClientAllocWindowClassExtraBytes(sz);
                                     pwnd->extraWindowUMPtr = UMptr;
                                xxxCallHook(WH CBT, HCBT CREATEWND, ...);
```

Reload 'Back To User-Mode'- OMG

```
user32!
        WNDCLASS wc = \{0\};
                                             xxxClientAllocWindowClassExtraBytes
       wc.cbWndExtra = 100;
                                                       hook stub:
       RegisterClass(&wc);
CreateWindow(wc.lpszClassName, ...);
                                                    DestroyWindow();
User
Kernel
                                             xxxCreateWindowEx:
                                                PVOID UMPtr =
                               xxxClientAllocWindowClassExtraBytes(sz);
                                     pwnd->extraWindowUMPtr = UMptr;
                                xxxCallHook (WH CBT, HCBT CREATEWND, ...);
```

CBTHookProc(...): ExitThread();



New Opportunities

Sometimes you cannot chain objects – but you can go back to usermode and then kill'em all

```
void UnlockNotifyWindow(pmenu)
 pitem = pmenu->items;
  for (; pitem != &pmenu->items[pmenu->cItems]; pitem++)
    if (NULL != pitem->psubmenu)
      UnlockNotifyWindow(pitem->psubmenu);
 HMAssignmentUnlock (&pmenu->notificationWnd);
```



Smash The Ref - Summary

Code that calls some destructors

• like DestroyMenu, FreeTimer, DestroySMWP and many more expects (primary) objects to be destroyed

However – it doesn't expect *other (secondary)* objects to be freed

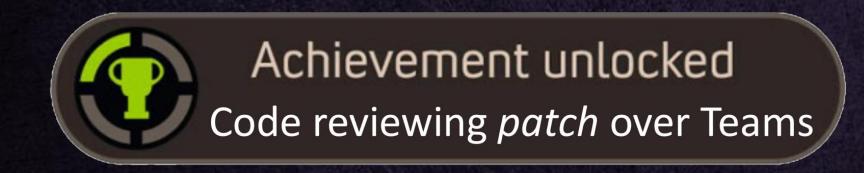
- with the 'zombie chain' and 'reloading a zombie'
- Primary vs secondary objects, no leaving usercrit
- 'Ultimate reloading' > xxx and sync race exploitation
- This attack works for multiple types of objects
 - Some don't even have refcount, oi oi oi

MSFT's Mitigation

Class IdentifyPrimaryDestroyTarget that is used in non-xxx destructors

Whenever a second object is destroyed during this scope,
 it is queued instead of being really destroyed

Global queue is emptied when it's safe - upon return to user-mode



He has made us all very busy, but we are enjoying the work tremendously, and he his findings are helping us to improve the security of win32k*.sys quite significantly

How It All Started — rare.c!xxxSendMinRectMessages

```
pdeskinfo = GETDESKINFO(pti);
   while (pwndShellHook = VWPLNext(pdeskinfo->pvwplShellHook)
       ThreadLock(pwndShellHook, &tlpwnd);
       xxxSendMessageTimeout(pwndShellHook, ...);
       ThreadUnlock(&tlpwnd);
```

POCs:

- 1. xxxMnOpenHierarchy window UAF
- 2. FreeTimer timer UAF
- 3. xxxCreateCaret PQ UAF
- 4. Ultimate reloading exit to user-mode from xxxFreeWindow of a zombie window
- 5. FreeSPB window UAF
- 6. xxxCapture window UAF
- 7. xxxCapture PQ UAF
- 8. zzzAttachThreadInput PQ UAF
- 9. xxxSendMinRectMessages DesktopInfo UAF (POC tries to catch free block before kLFH)
- 10. UnlockNotifyWindow menu UAF via user-mode callback
- 11. CSRSS Arbitrary user-mode heap-free

Spotted More Instances

- 1. xxxMnKeyDown
- 2. xxxQueryDropObject
- 3. xxxSetFocus
- 4. xxxMouseActivate
- 5. xxxDragObject
- 6. xxxCapture
- 7. SetMenultemInfo
- 8. xxxQueryDropObject

- 9. SendMsgCleanup
- 10. MnFreePopup
- 11. Un/LockPopupMenu
- 12. zzzDestroyQueue
- 13. LockMFMWFPWindow
- 14. MNFreeltem
- 15. xxxFreeWindow

MSRC top 30 for 2020 \$300K bounty award

Thanks to Tomer Schwartz and Saar Amar of MSRC-IL ©



Thank You

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