

NCKU Programming Contest Training Course

2013/08/09

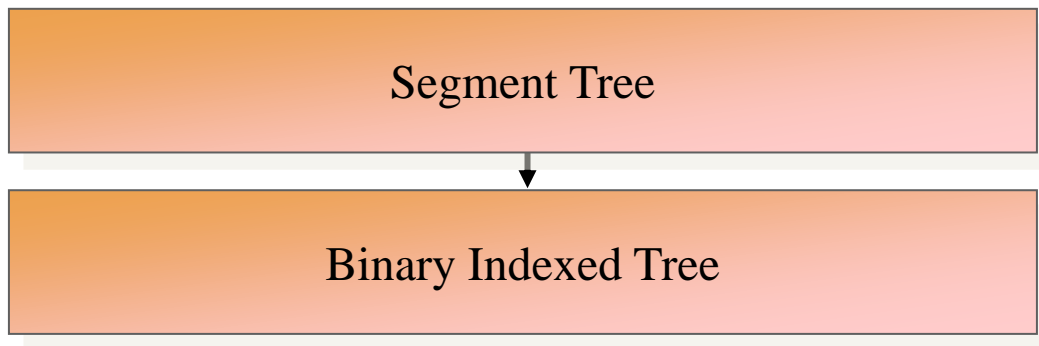
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<http://myweb.ncku.edu.tw/~p76014143/20130809.rar>

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National Cheng Kung University
Tainan, Taiwan



Outline



Segment Tree

Problem

求區間內數字最大

陣列長度 = N

Query 數 = Q

時間複雜度 = $O(?)$

3	8	9	2	5	6	7	9	1	4
---	---	---	---	---	---	---	---	---	---



Segment Tree

- Segment Tree
 - A tree-based data structure
 - Construct tree with $O(N)$
 - RMQ (range minimum/maximum query problem) in $O(\log N)$



Segment Tree

1

3	8	9	2	5	6	7	9	1	4
---	---	---	---	---	---	---	---	---	---



Segment Tree

1

3	8	9	2	5	6	7	9	1	4
---	---	---	---	---	---	---	---	---	---

2

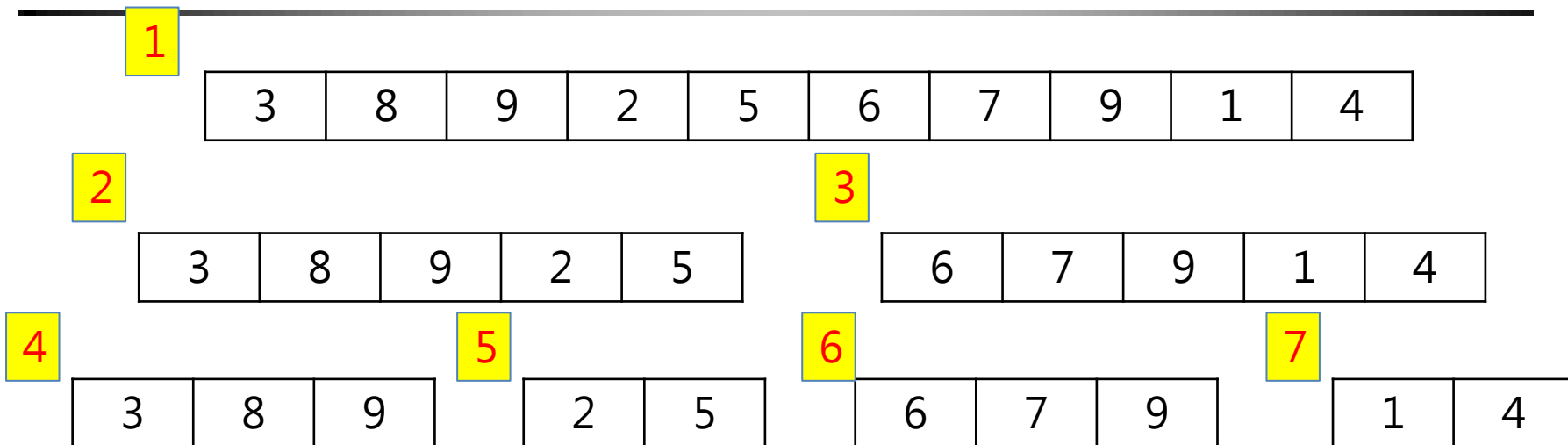
3	8	9	2	5
---	---	---	---	---

3

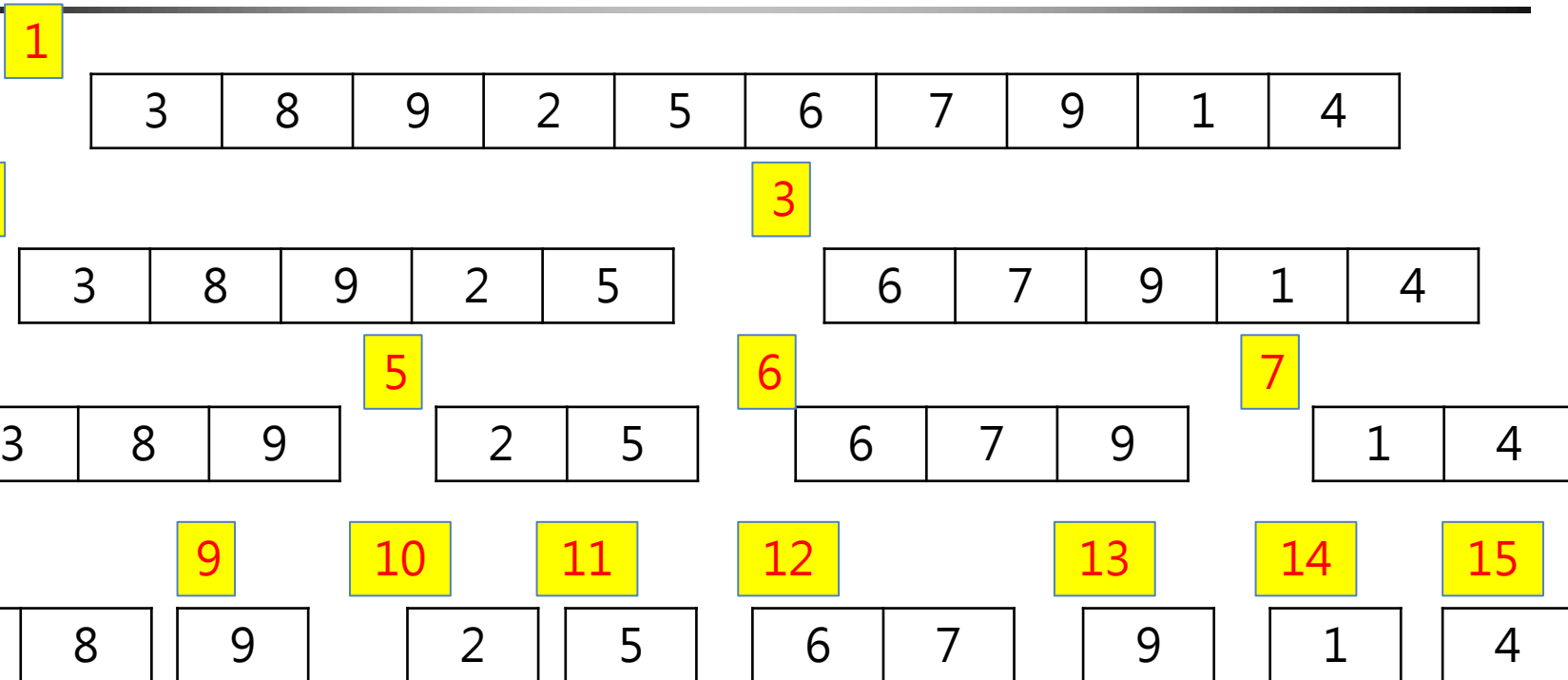
6	7	9	1	4
---	---	---	---	---



Segment Tree



Segment Tree



Segment Tree

1

3	8	9	2	5	6	7	9	1	4
---	---	---	---	---	---	---	---	---	---

2

3	8	9	2	5
---	---	---	---	---

3

6	7	9	1	4
---	---	---	---	---

4

3	8	9
---	---	---

5

2	5
---	---

6

6	7	9
---	---	---

7

1	4
---	---

8

3	8
---	---

9

9

10

2

11

5

12

6	7
---	---

13

9

14

1

15

4

16

3

17

8

24

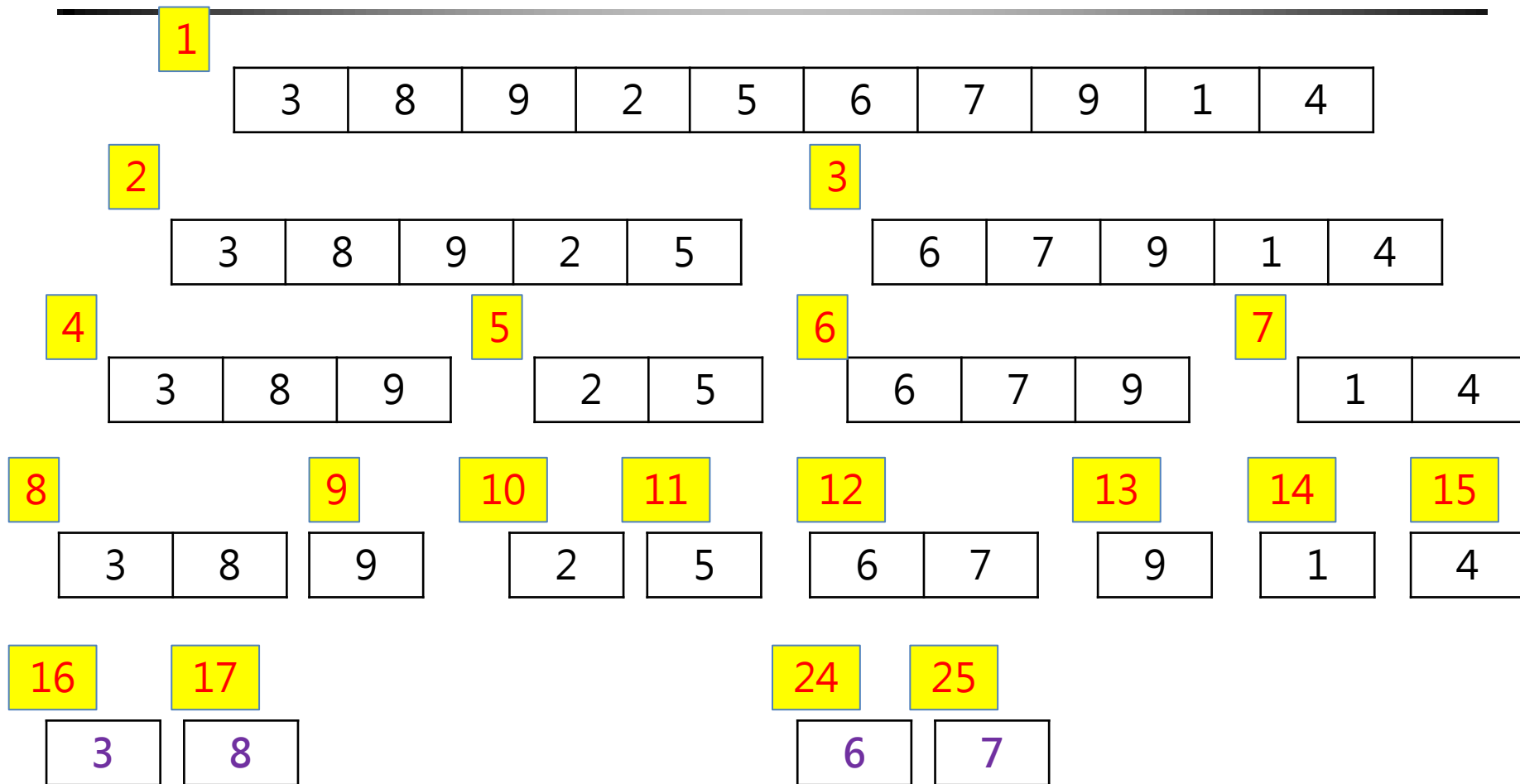
6

25

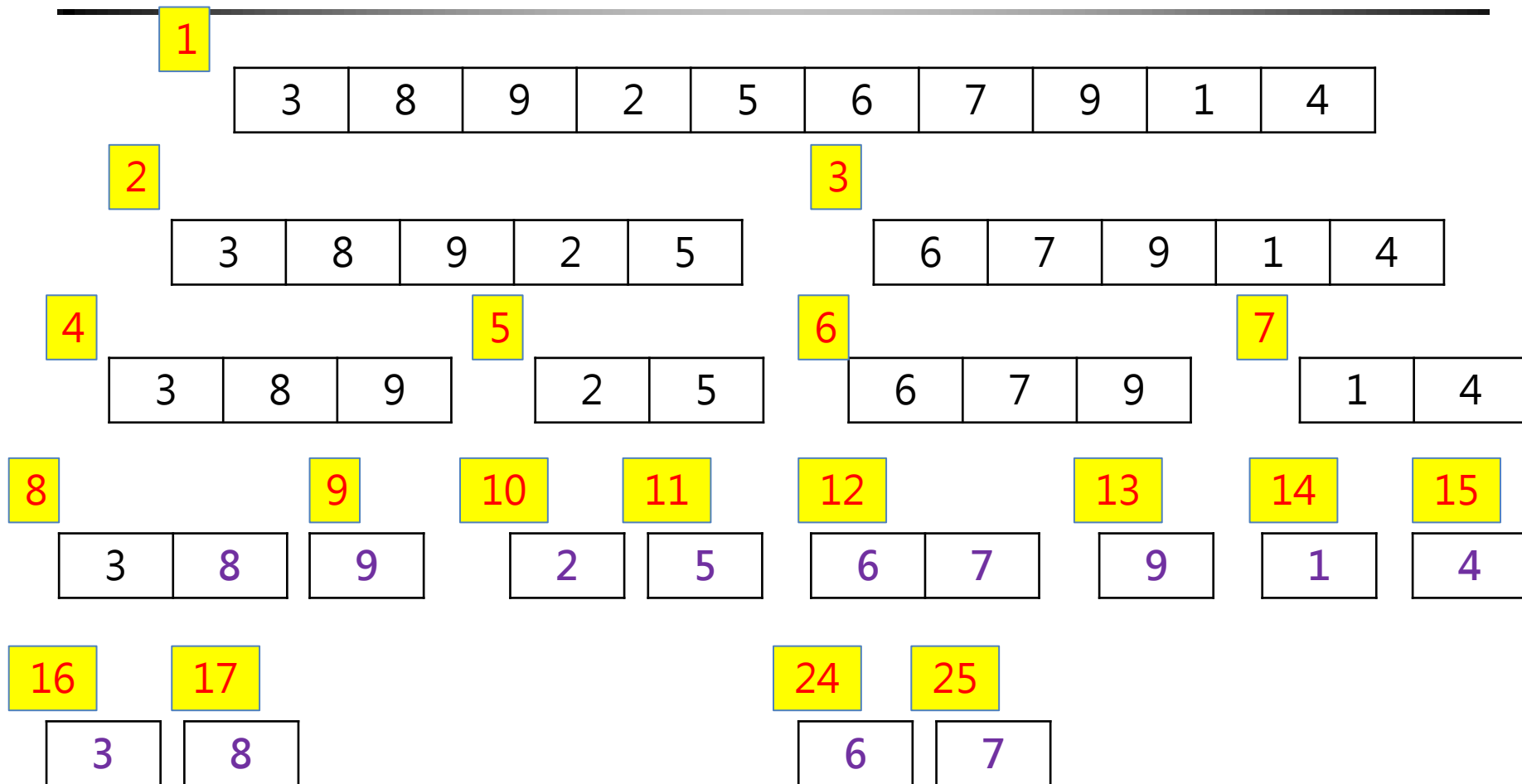
7



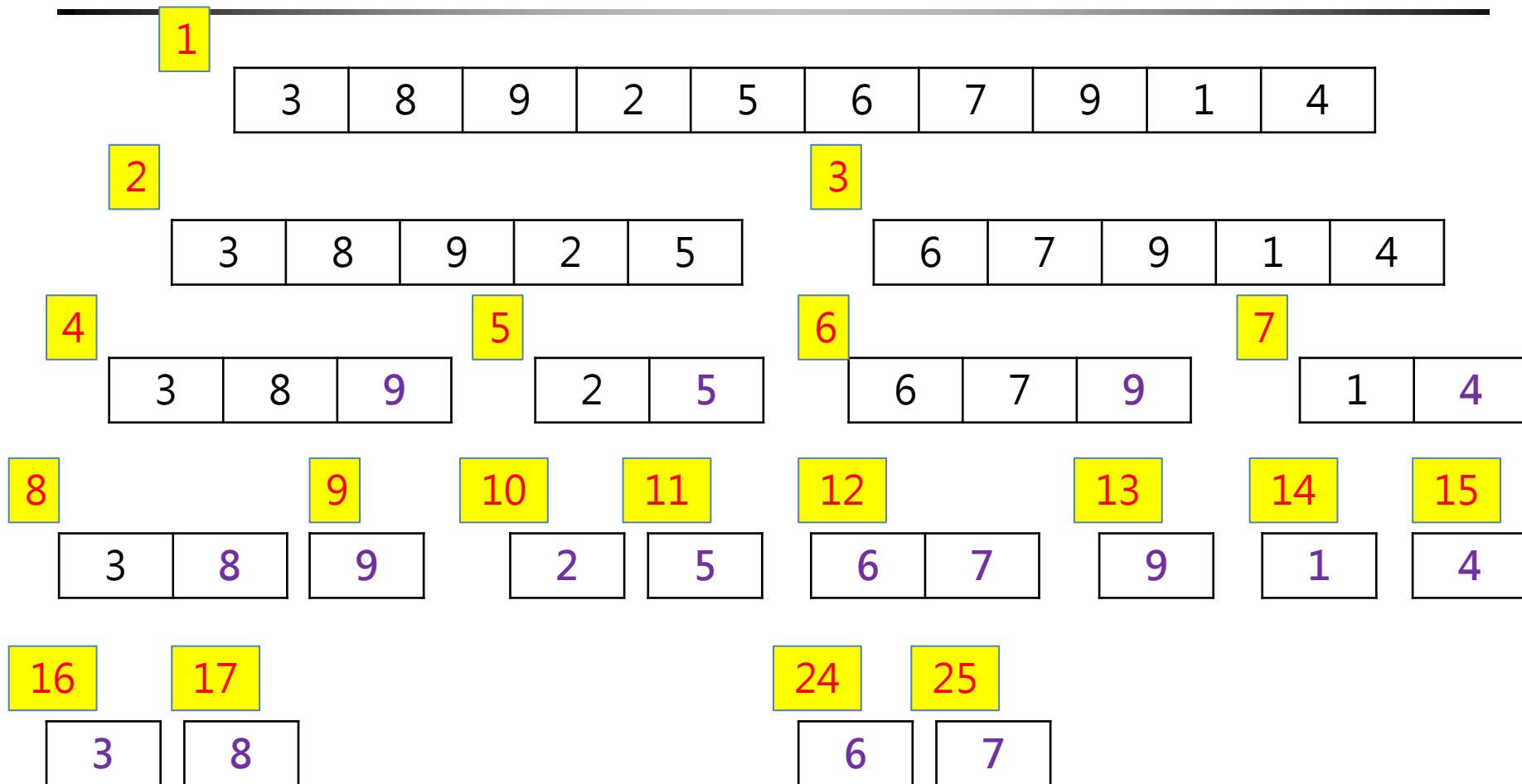
Segment Tree



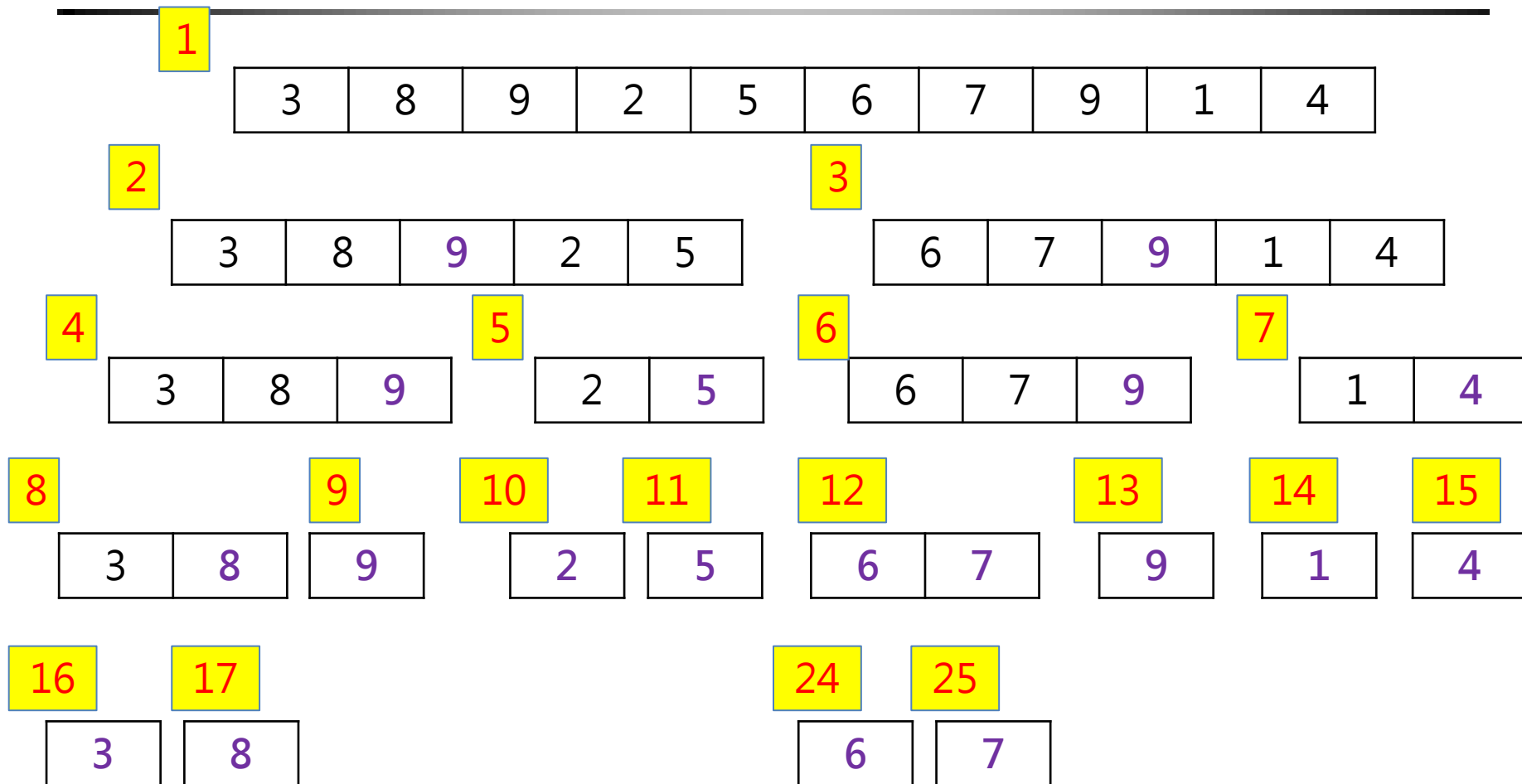
Segment Tree



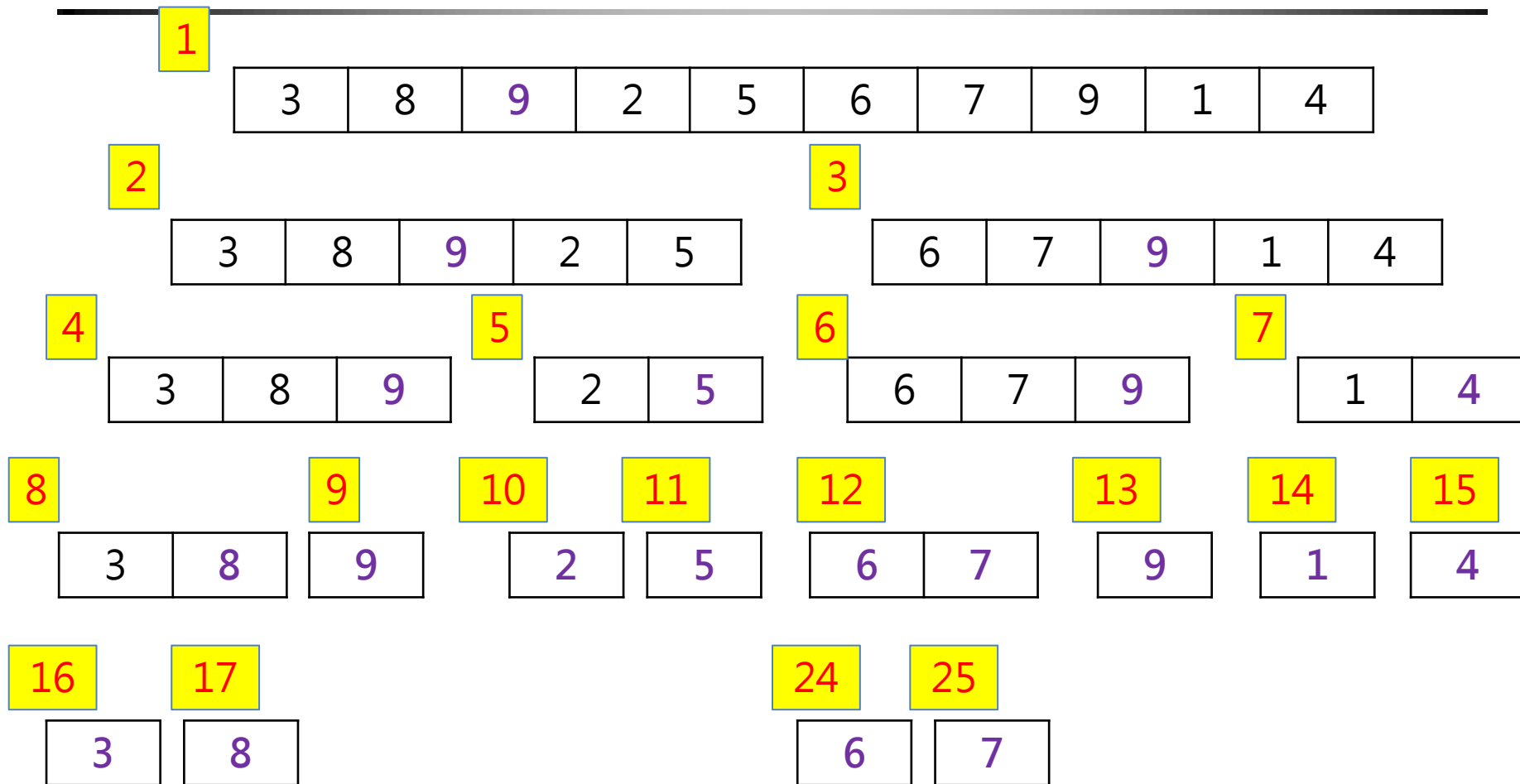
Segment Tree



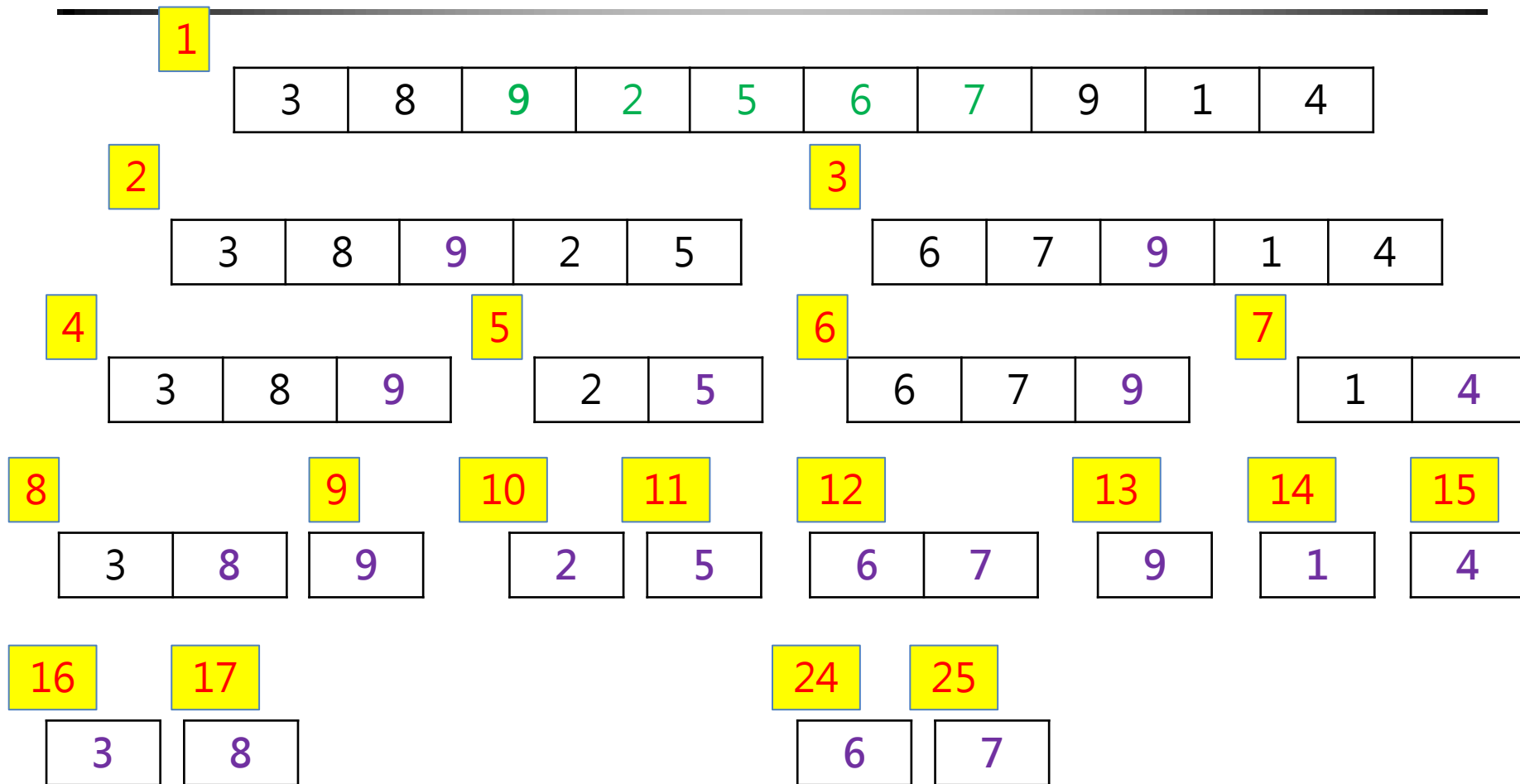
Segment Tree



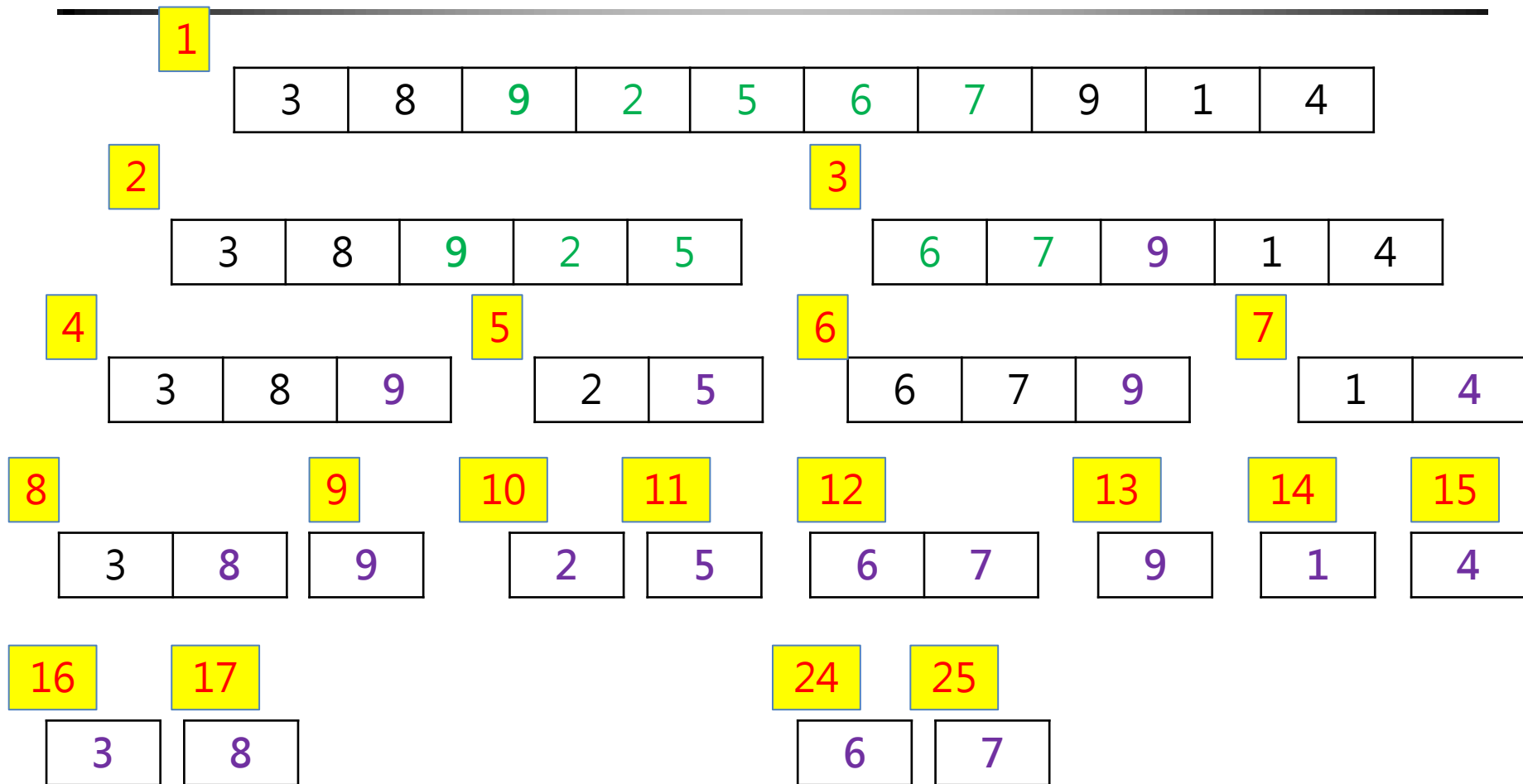
Segment Tree



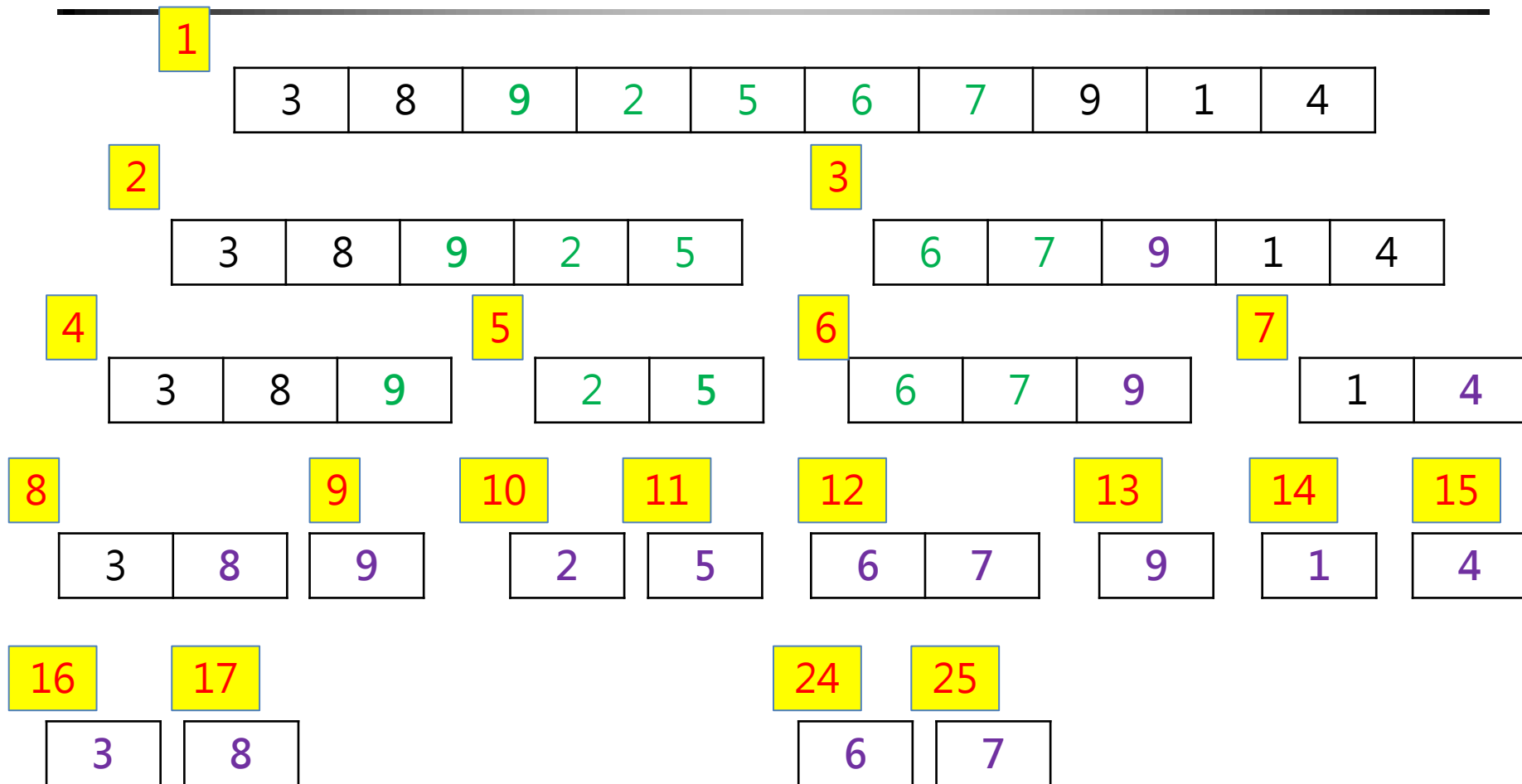
Segment Tree



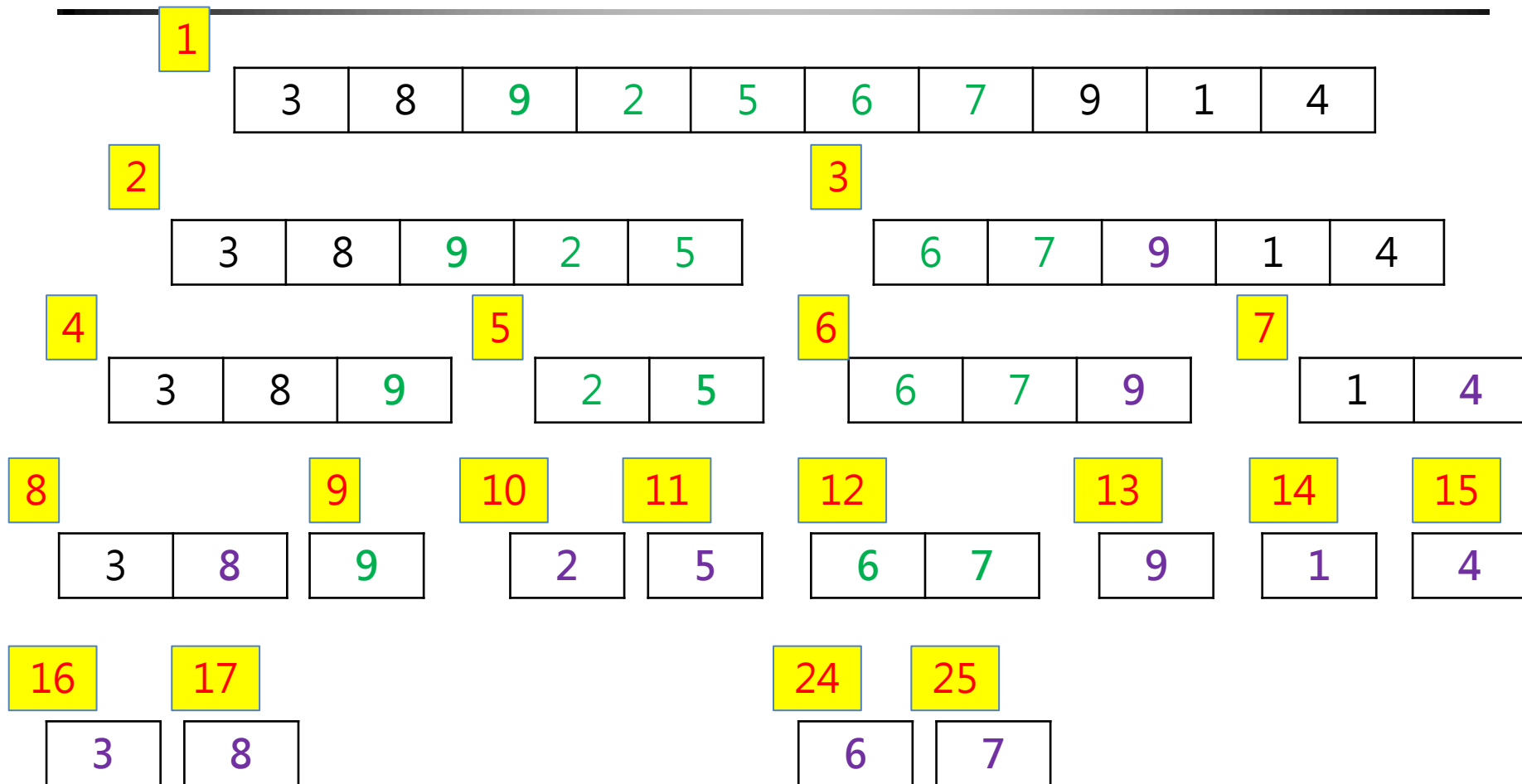
Segment Tree



Segment Tree



Segment Tree



Example

- POJ-3264

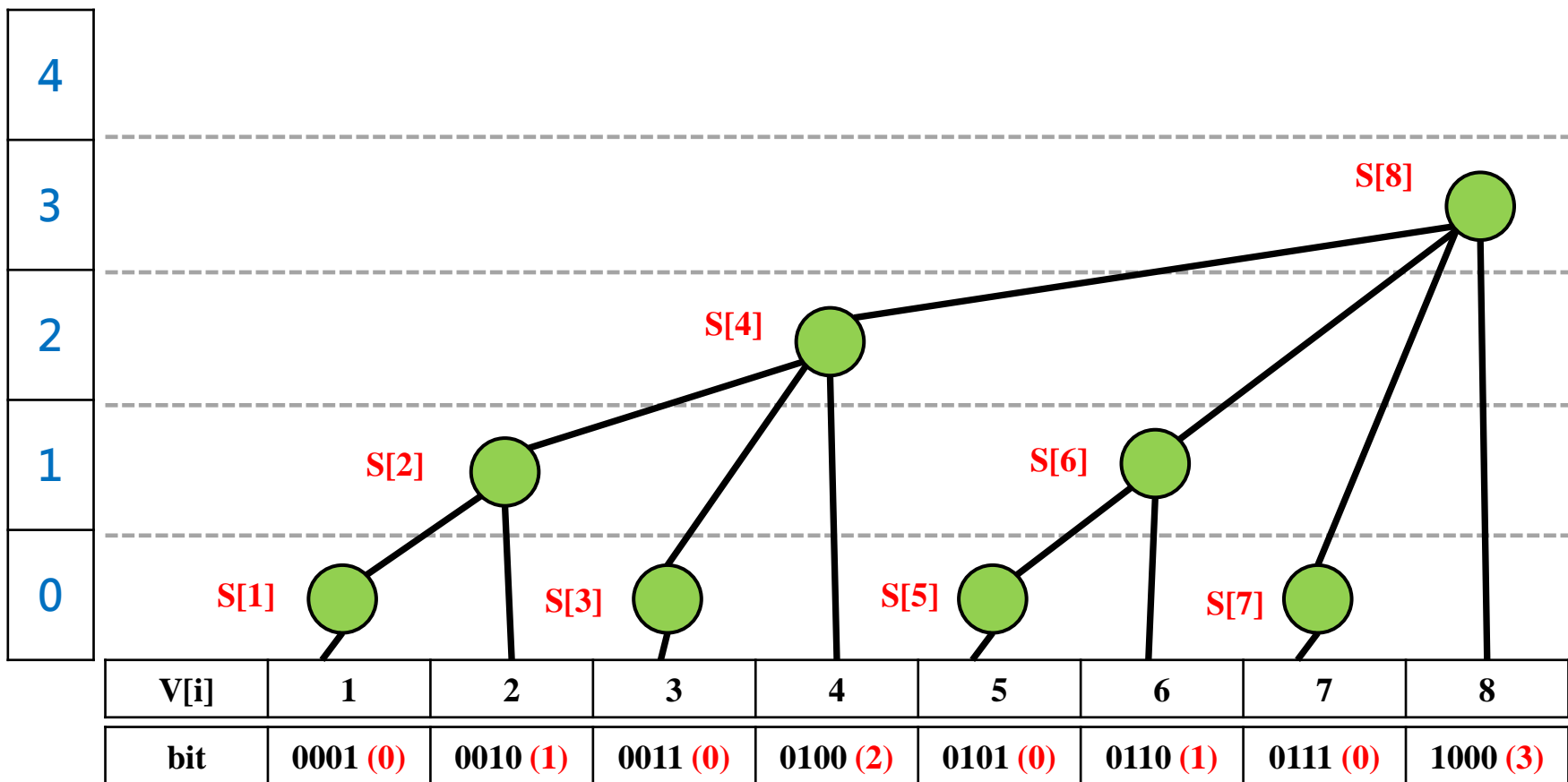


Binary Indexed Tree

- Operation
 - Increase/decrease the value of an element
 - Query the summation within the interval $[i...j]$
 - How to solve this problem efficiently?
- Naïve Solution
 - Each query requires $O(N)$
 - Total **$O(QN)$** ...

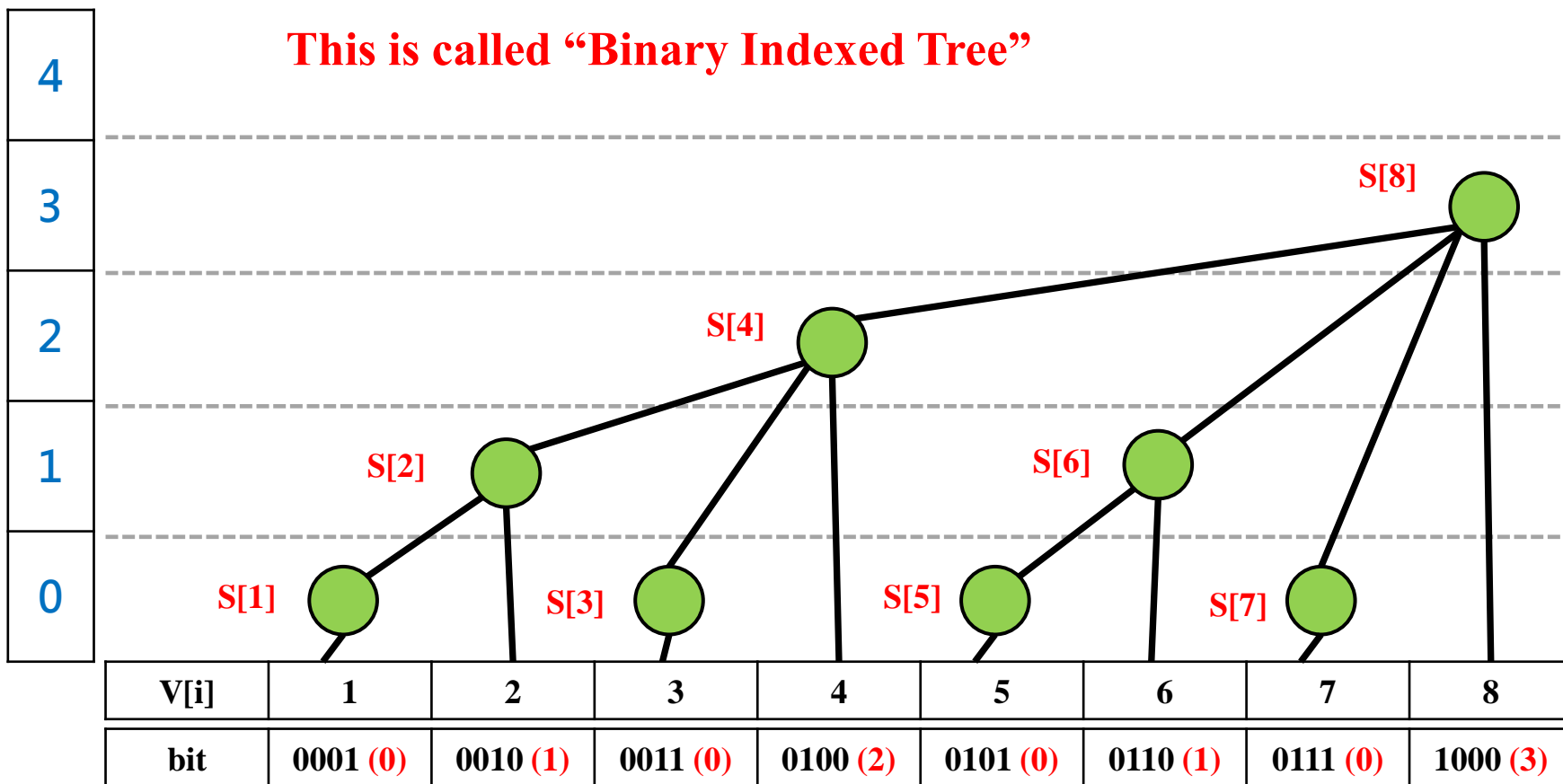


Binary Indexed Tree



Binary Indexed Tree

This is called “Binary Indexed Tree”



Binary Indexed Tree

$$s[1] = v[1]$$

$$s[2] = v[2] + s[1]$$

$$s[3] = v[3]$$

$$s[4] = v[4] + s[3] + s[2]$$

$$s[5] = v[5]$$

$$s[6] = v[6] + s[5]$$

$$s[7] = v[7]$$

$$s[8] = v[8] + s[7] + s[6] + s[4]$$

...

What's the regularity?



Binary Indexed Tree

Define:

```
int lowbit (int in)
{
    return in&(-in);
}
```

ex:

lowbit(1) = 1

lowbit(2) = 2

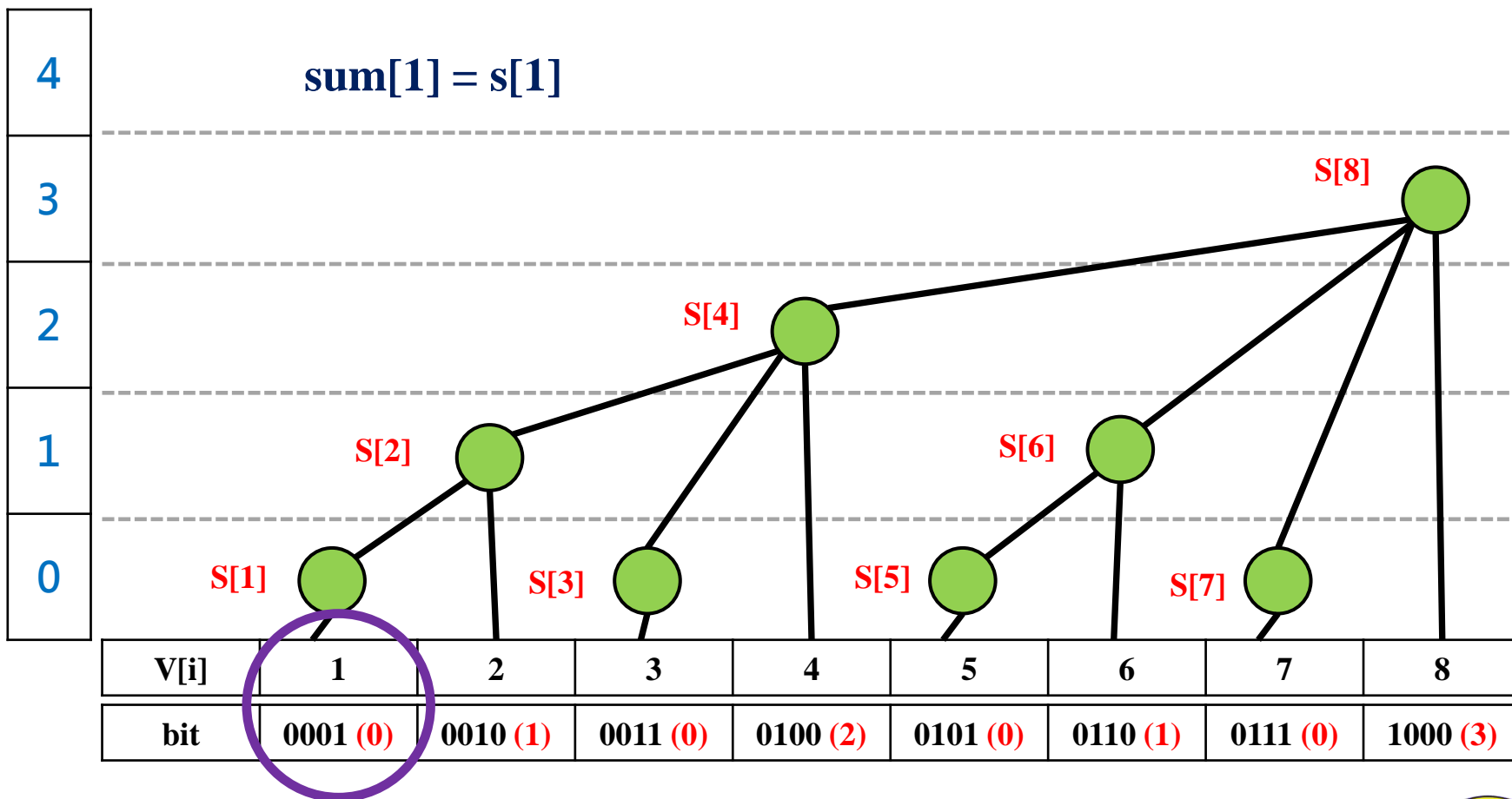
lowbit(3) = 1

lowbit(4) = 4

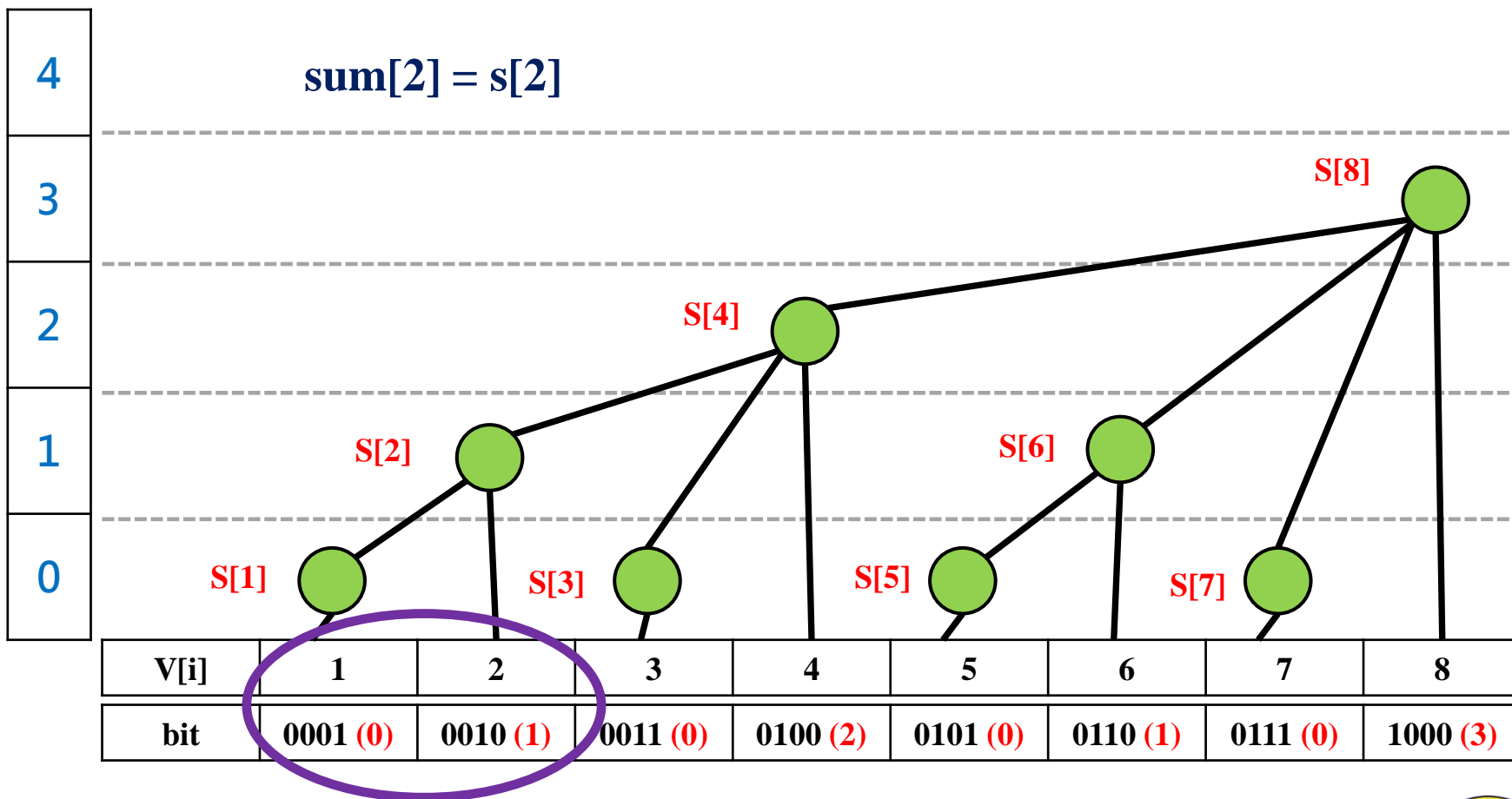
...



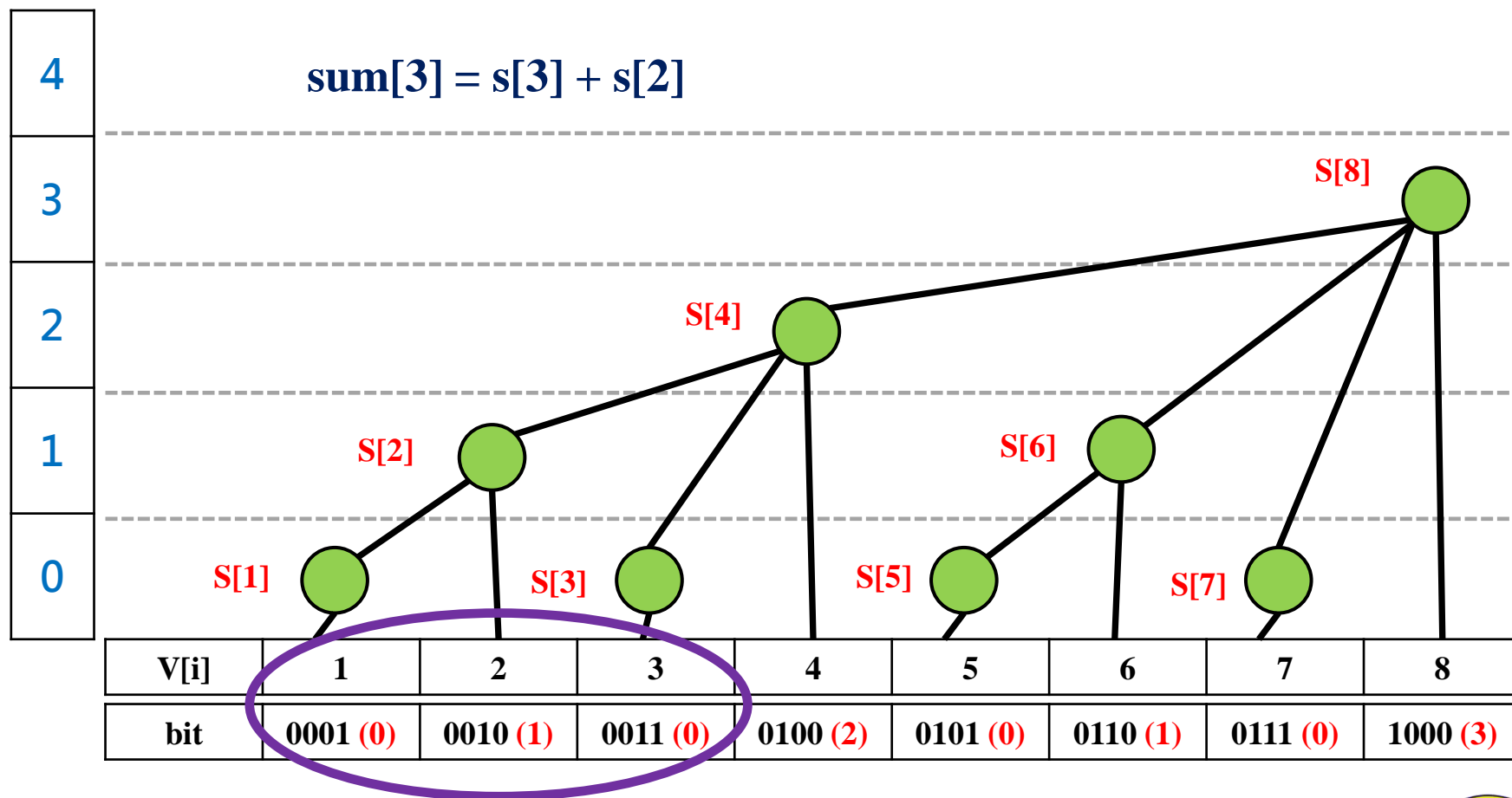
Binary Indexed Tree



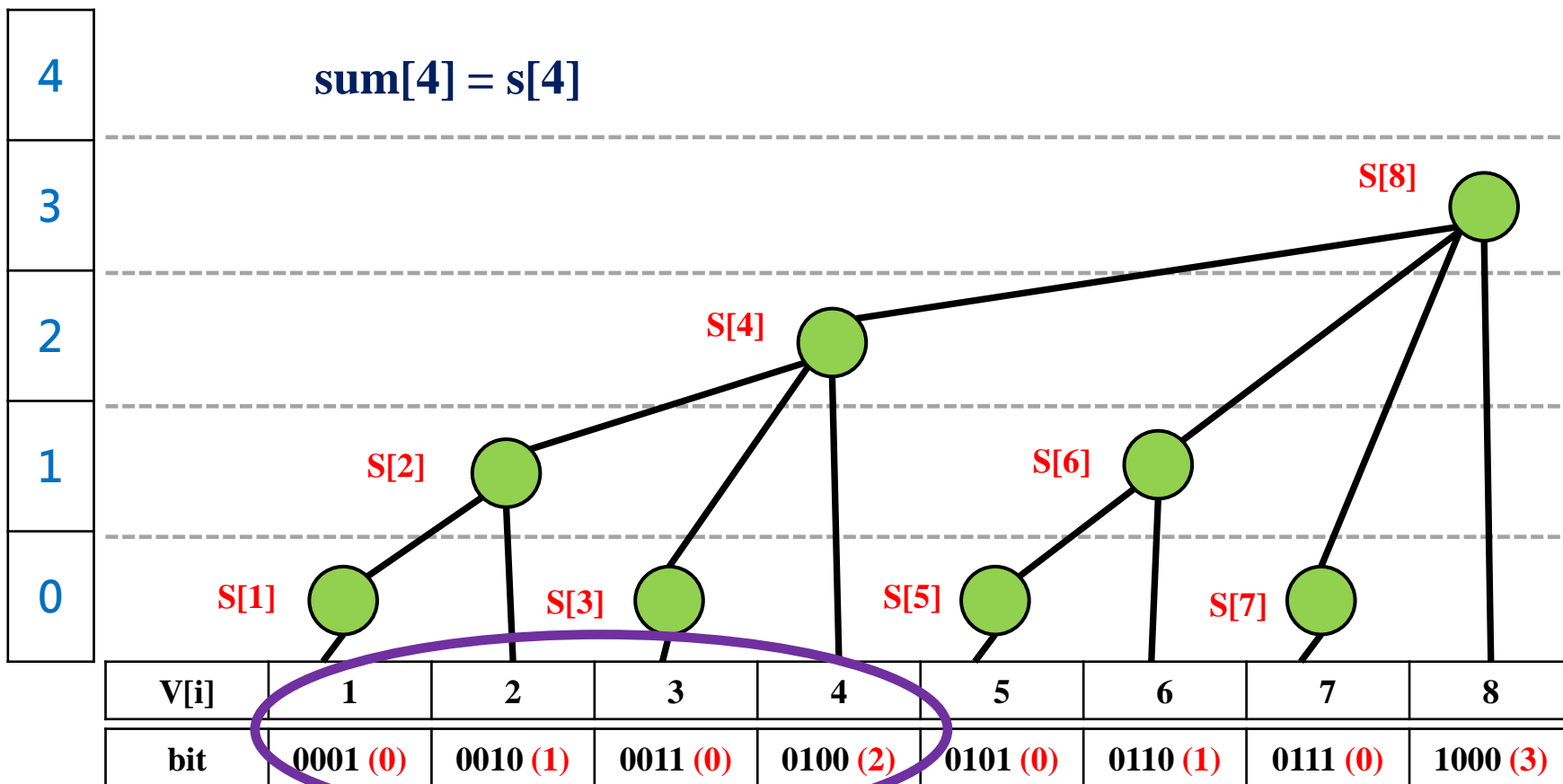
Binary Indexed Tree



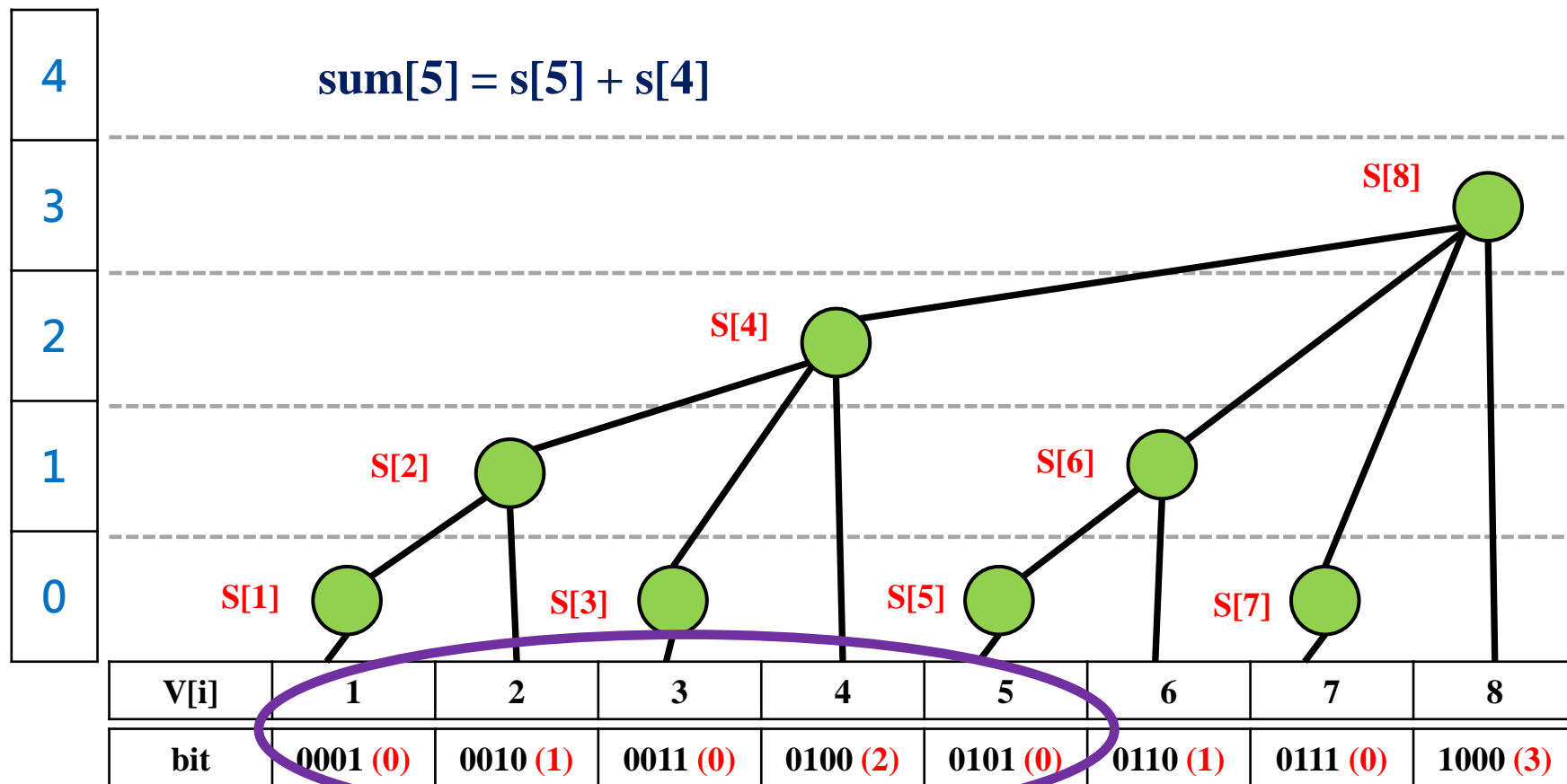
Binary Indexed Tree



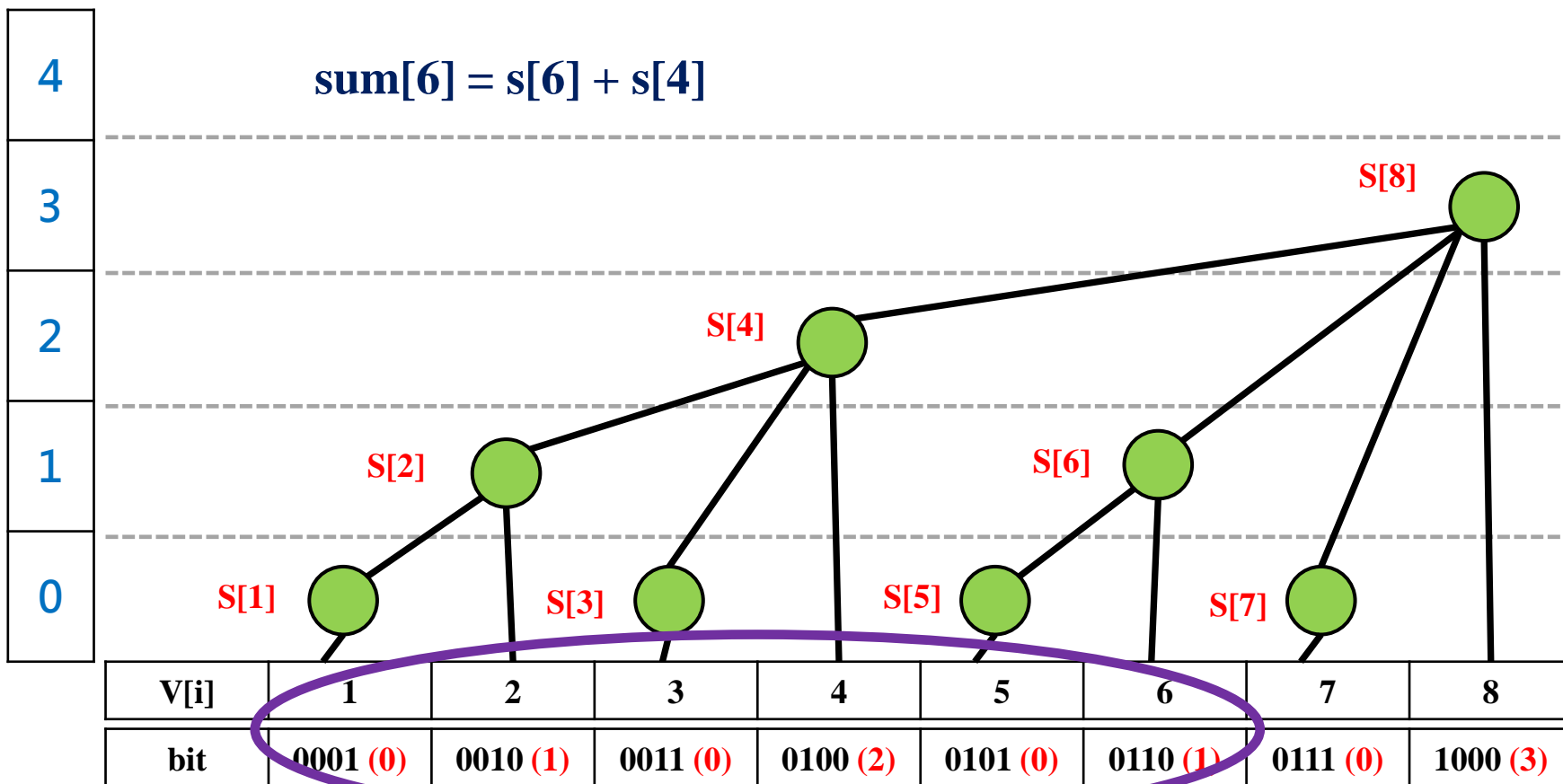
Binary Indexed Tree



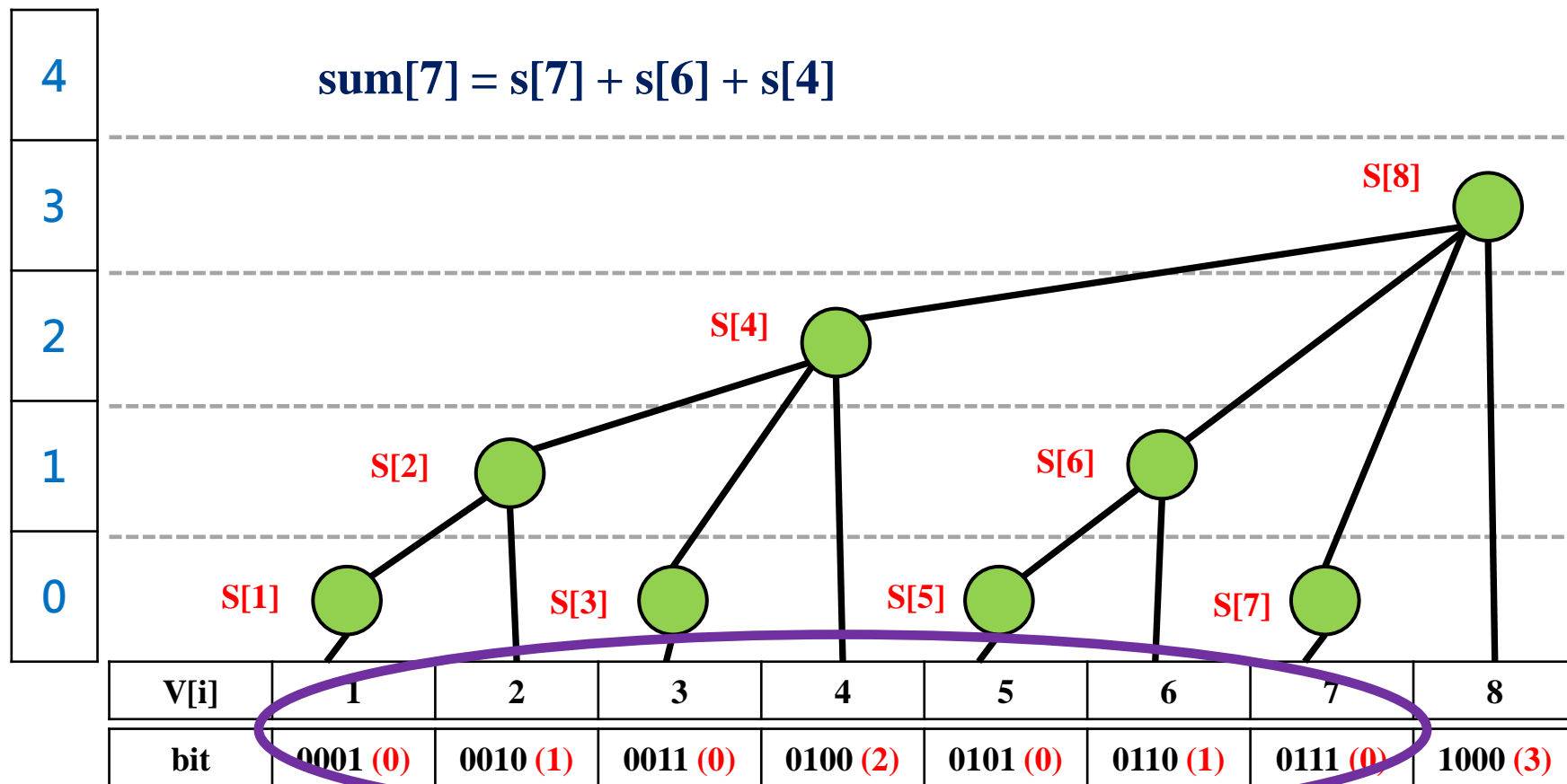
Binary Indexed Tree



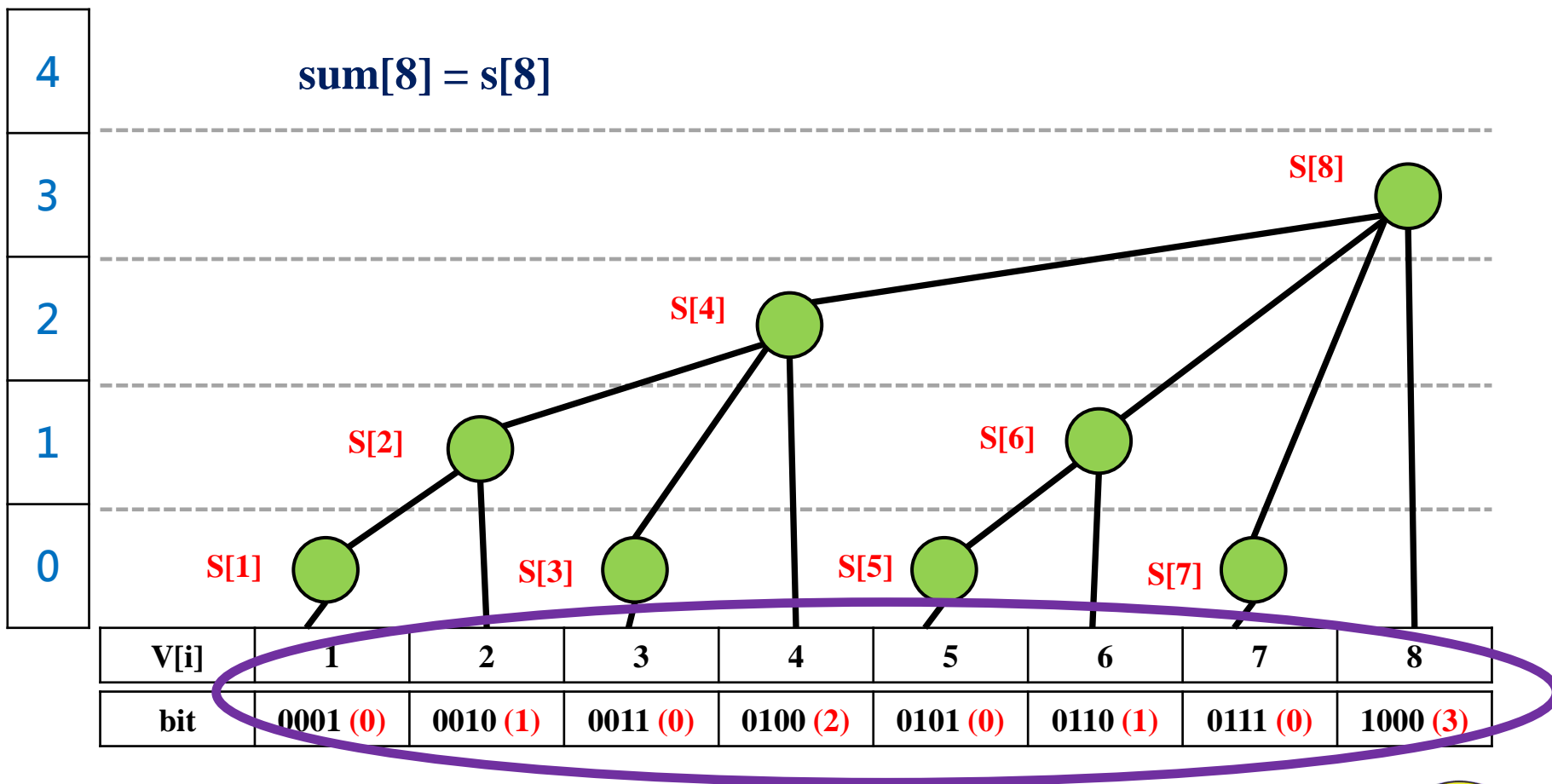
Binary Indexed Tree



Binary Indexed Tree



Binary Indexed Tree



Binary Indexed Tree

Define:

```
int lowbit (int in)
{
    return in&(-in);
}
```

ex:

lowbit(1) = 1

lowbit(2) = 2

lowbit(3) = 1

lowbit(4) = 4

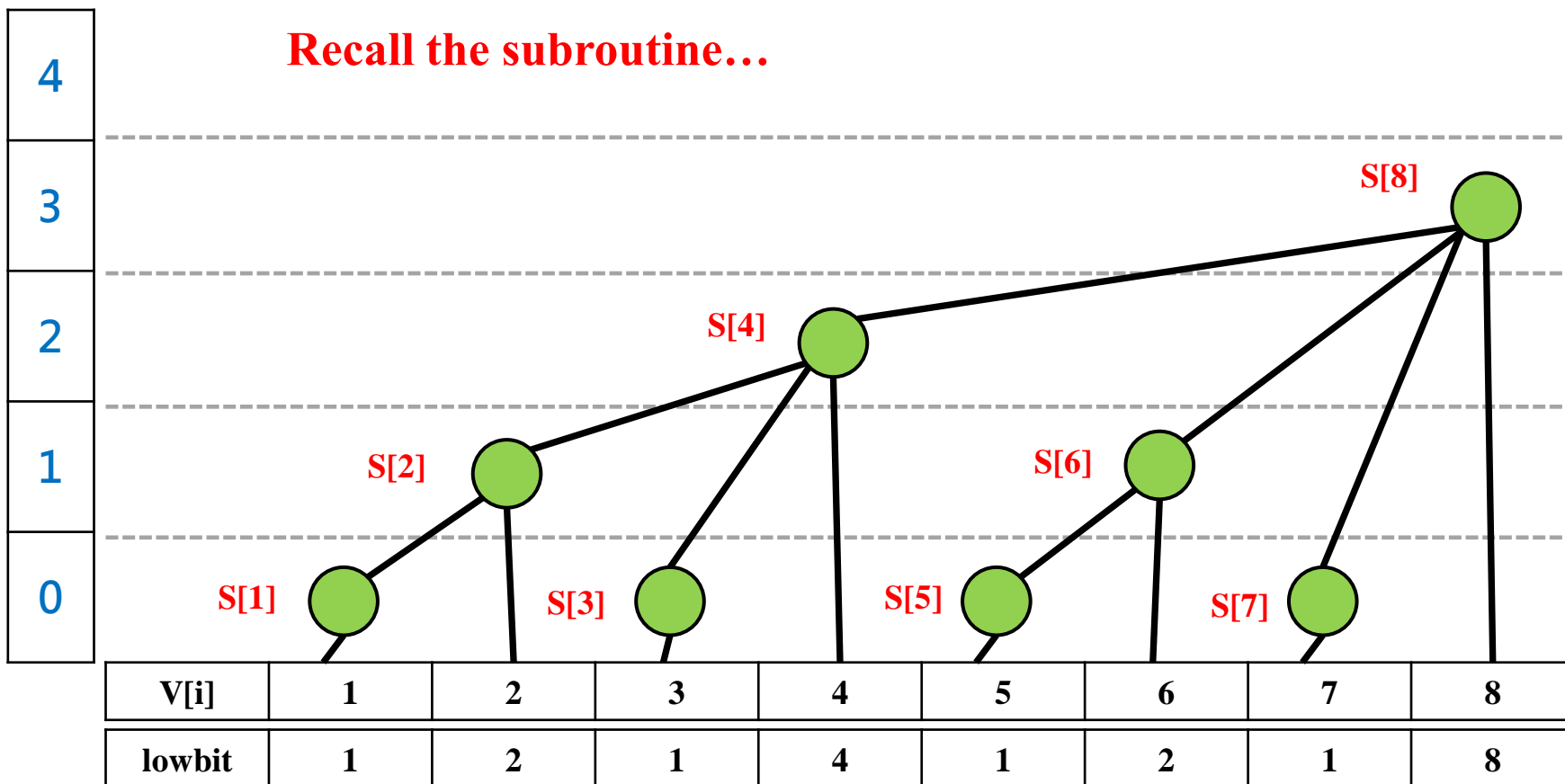
...

V[i]	1	2	3	4	5	6	7	8
lowbit	1	2	1	4	1	2	1	8

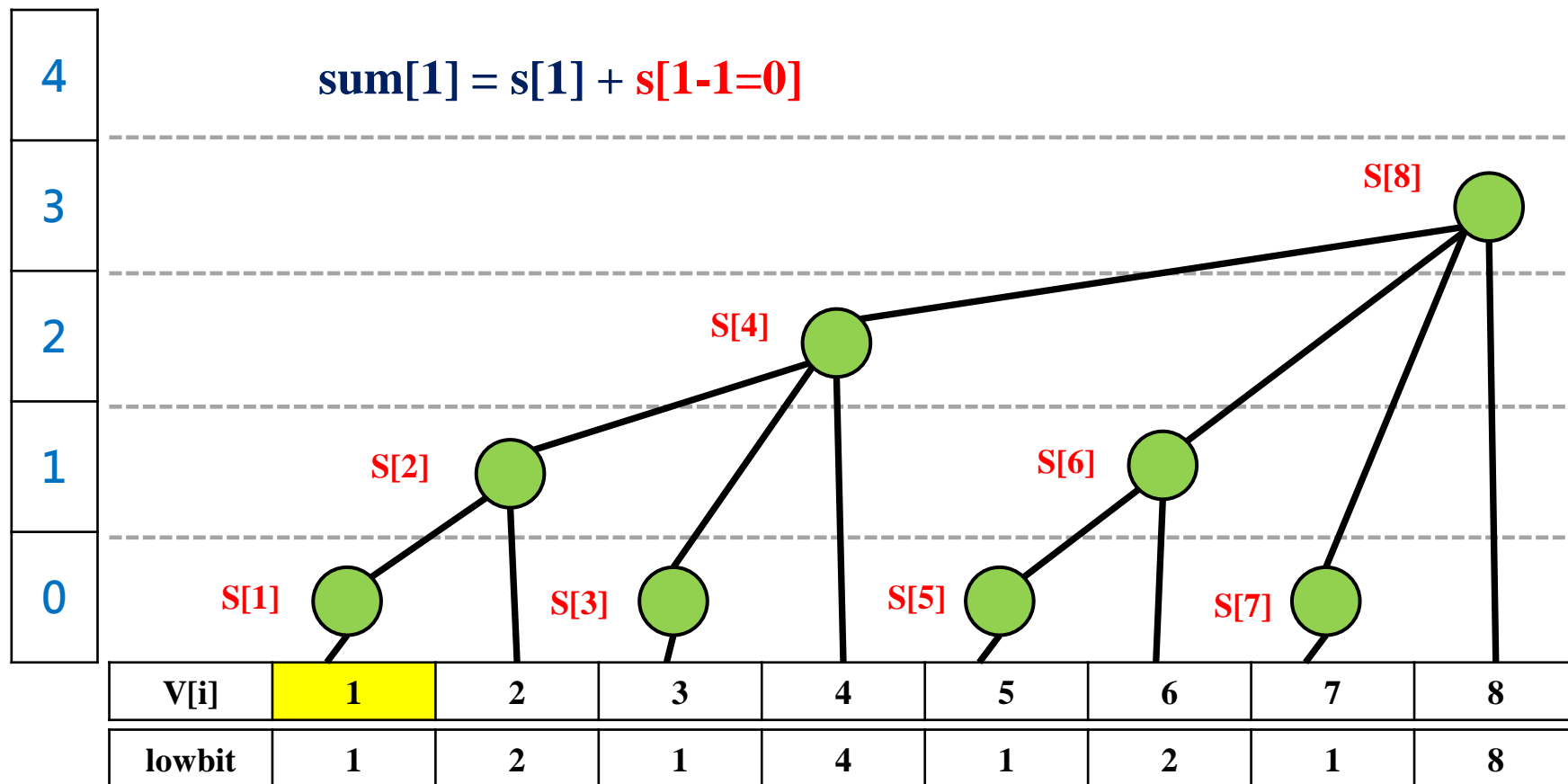


Binary Indexed Tree

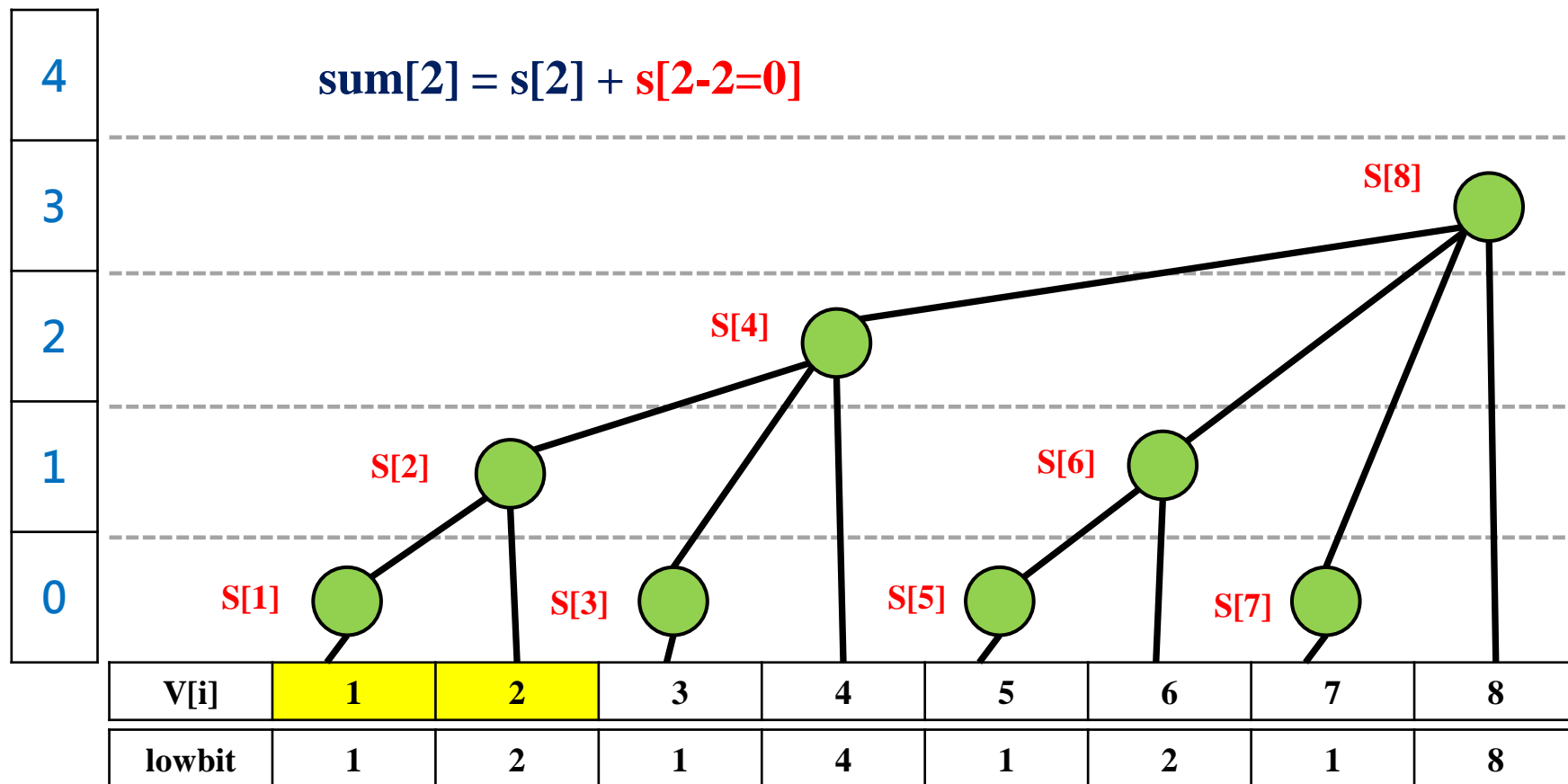
Recall the subroutine...



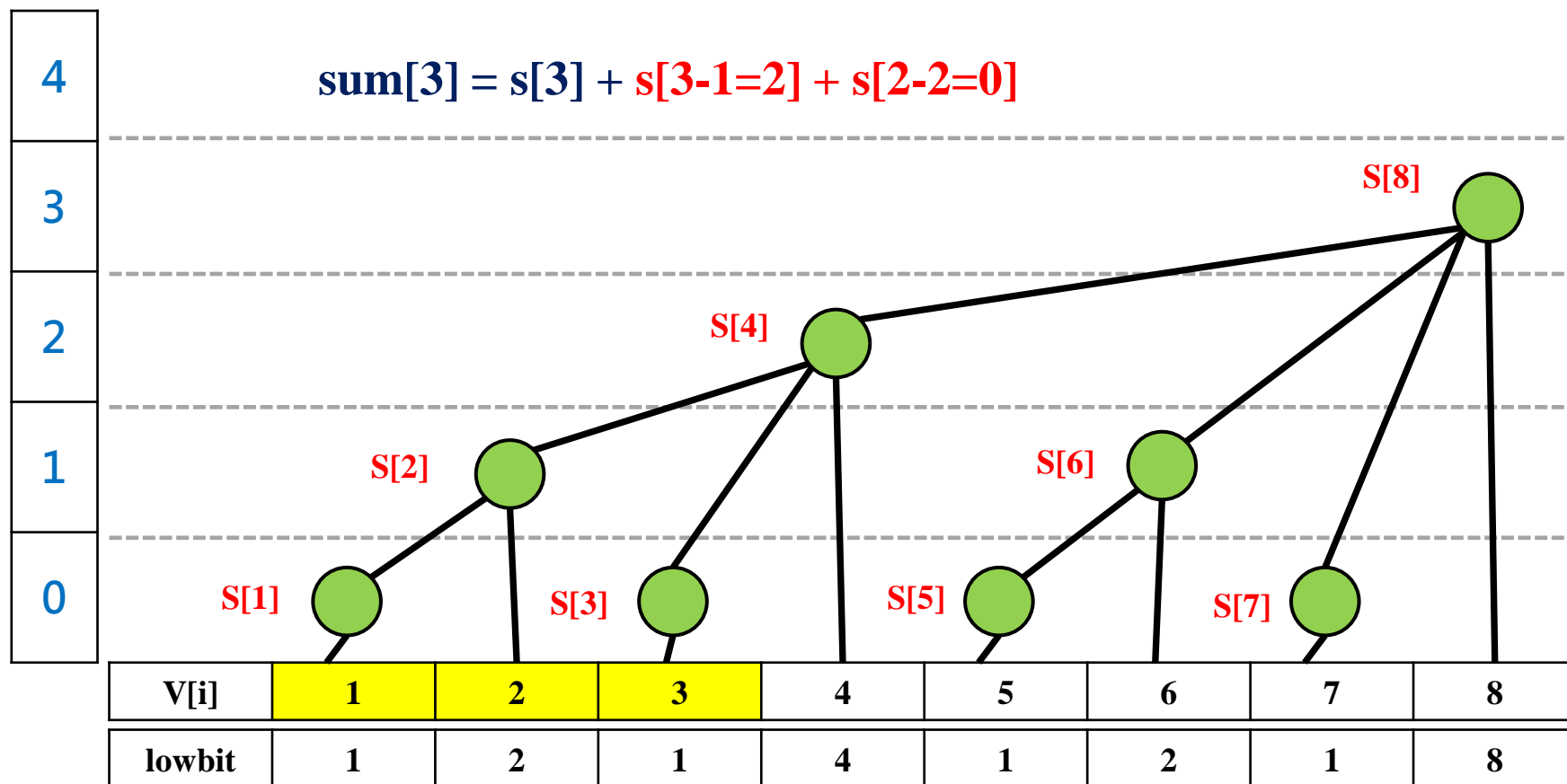
Binary Indexed Tree



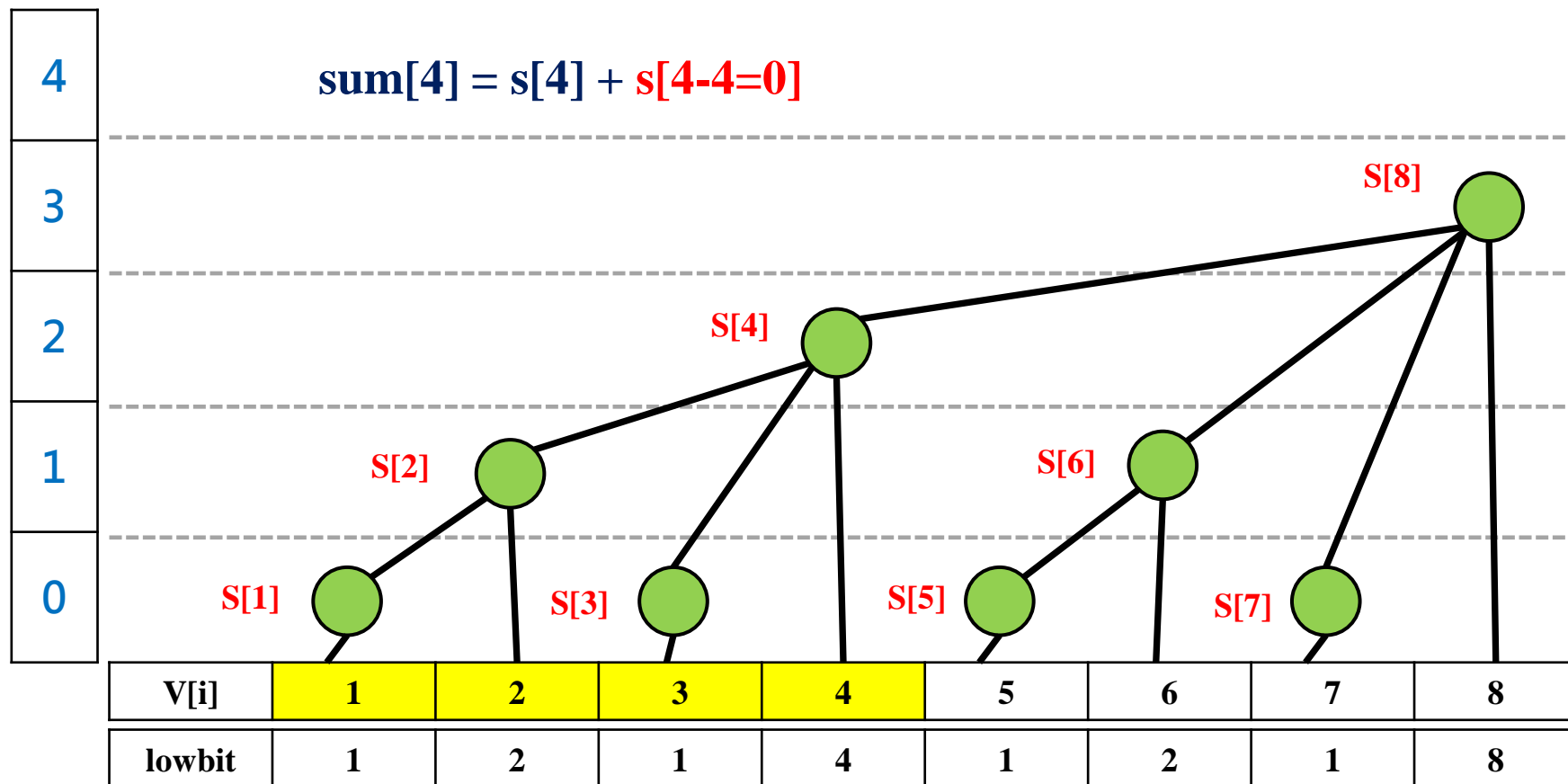
Binary Indexed Tree



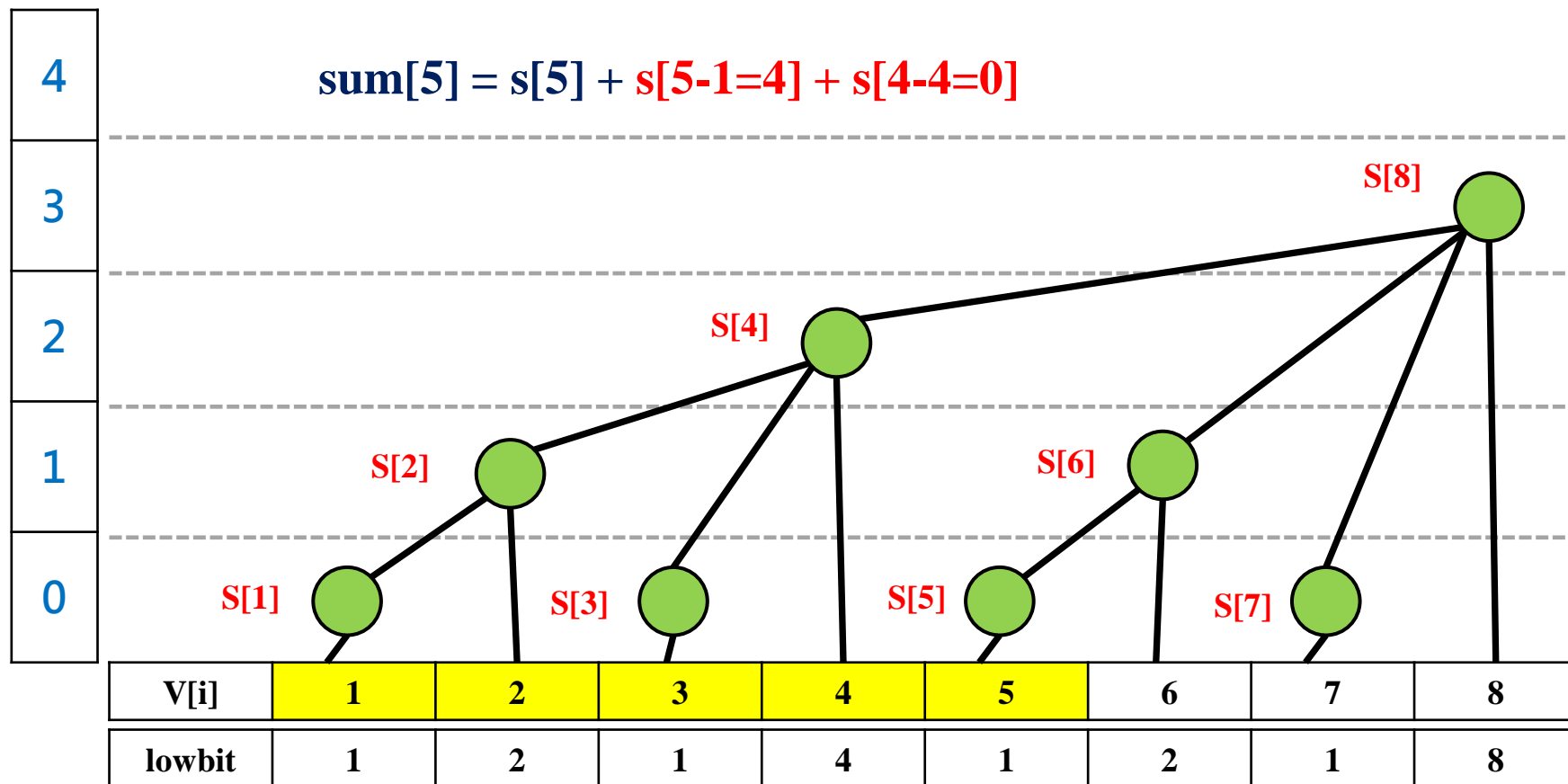
Binary Indexed Tree



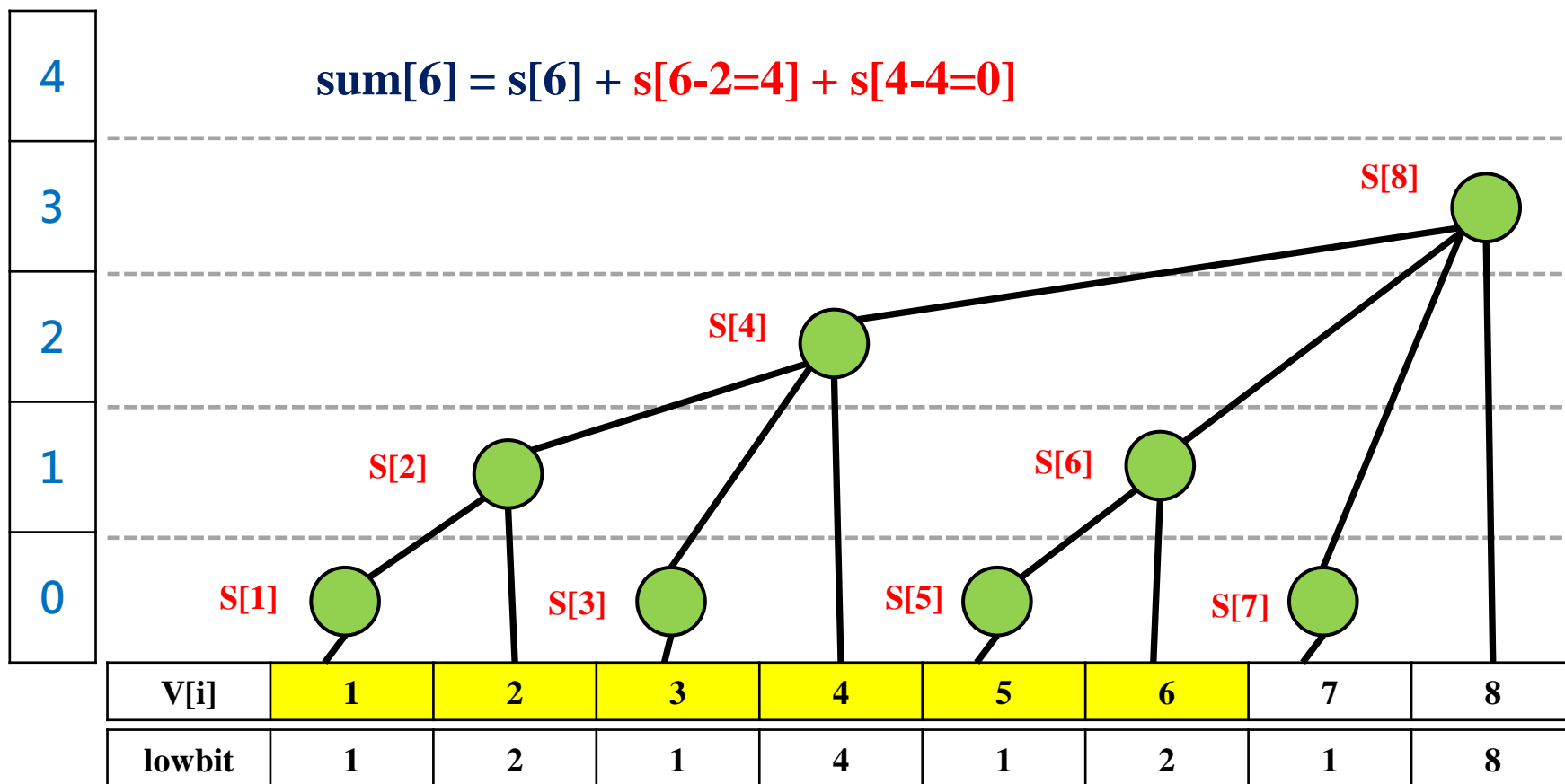
Binary Indexed Tree



Binary Indexed Tree

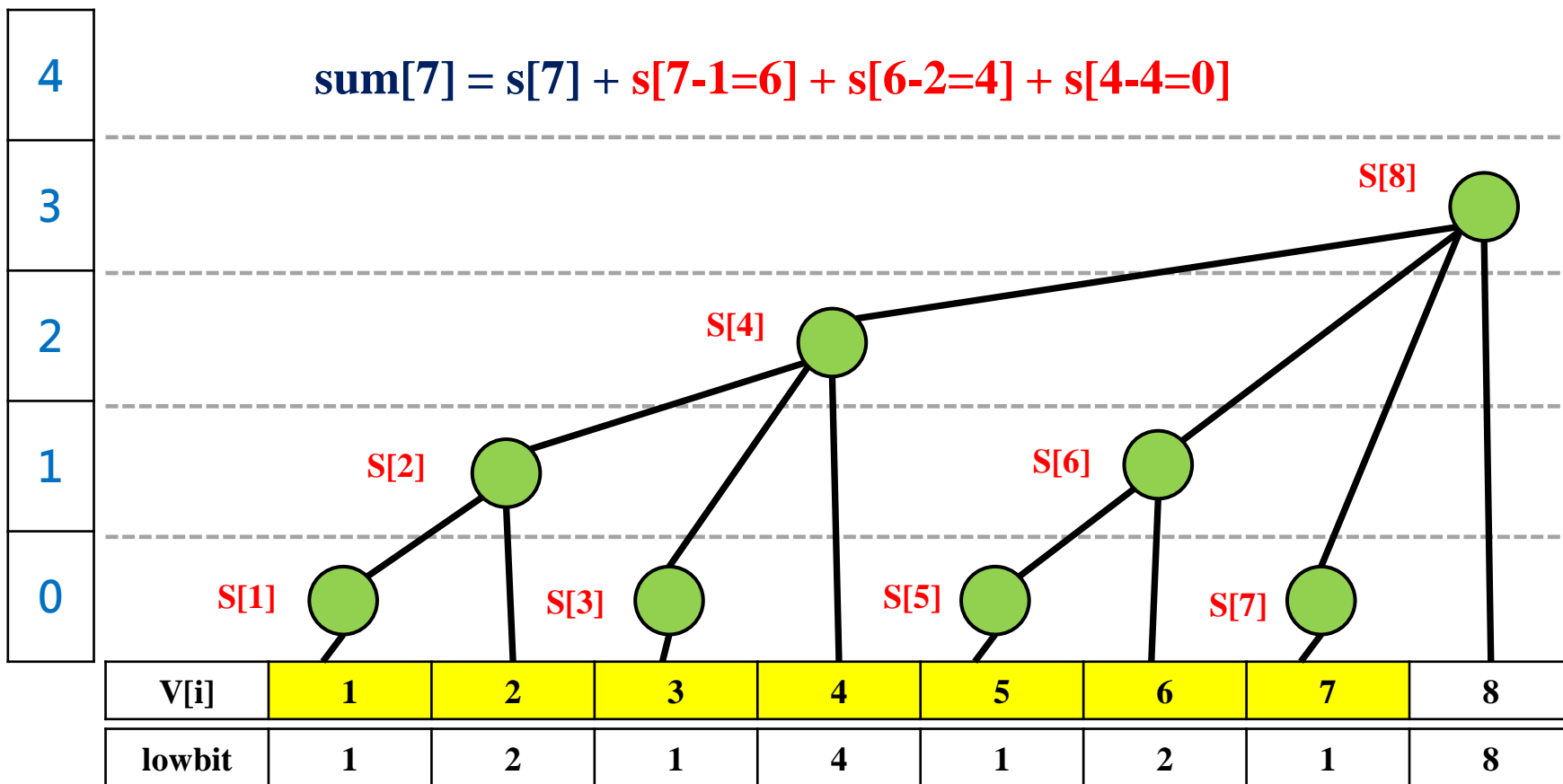


Binary Indexed Tree

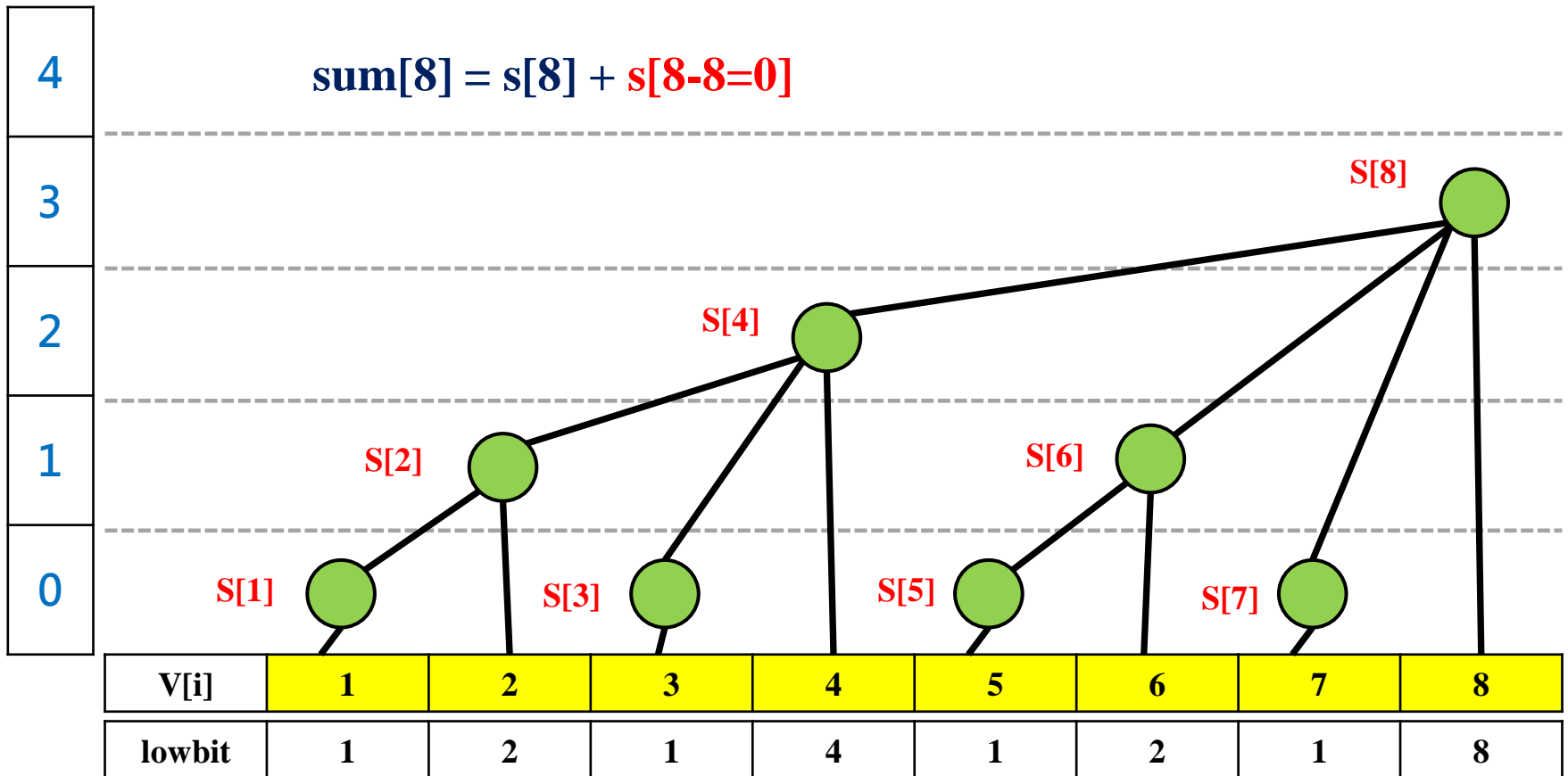


Binary Indexed Tree

$$\text{sum}[7] = s[7] + s[7-1=6] + s[6-2=4] + s[4-4=0]$$



Binary Indexed Tree



Binary Indexed Tree

Define:

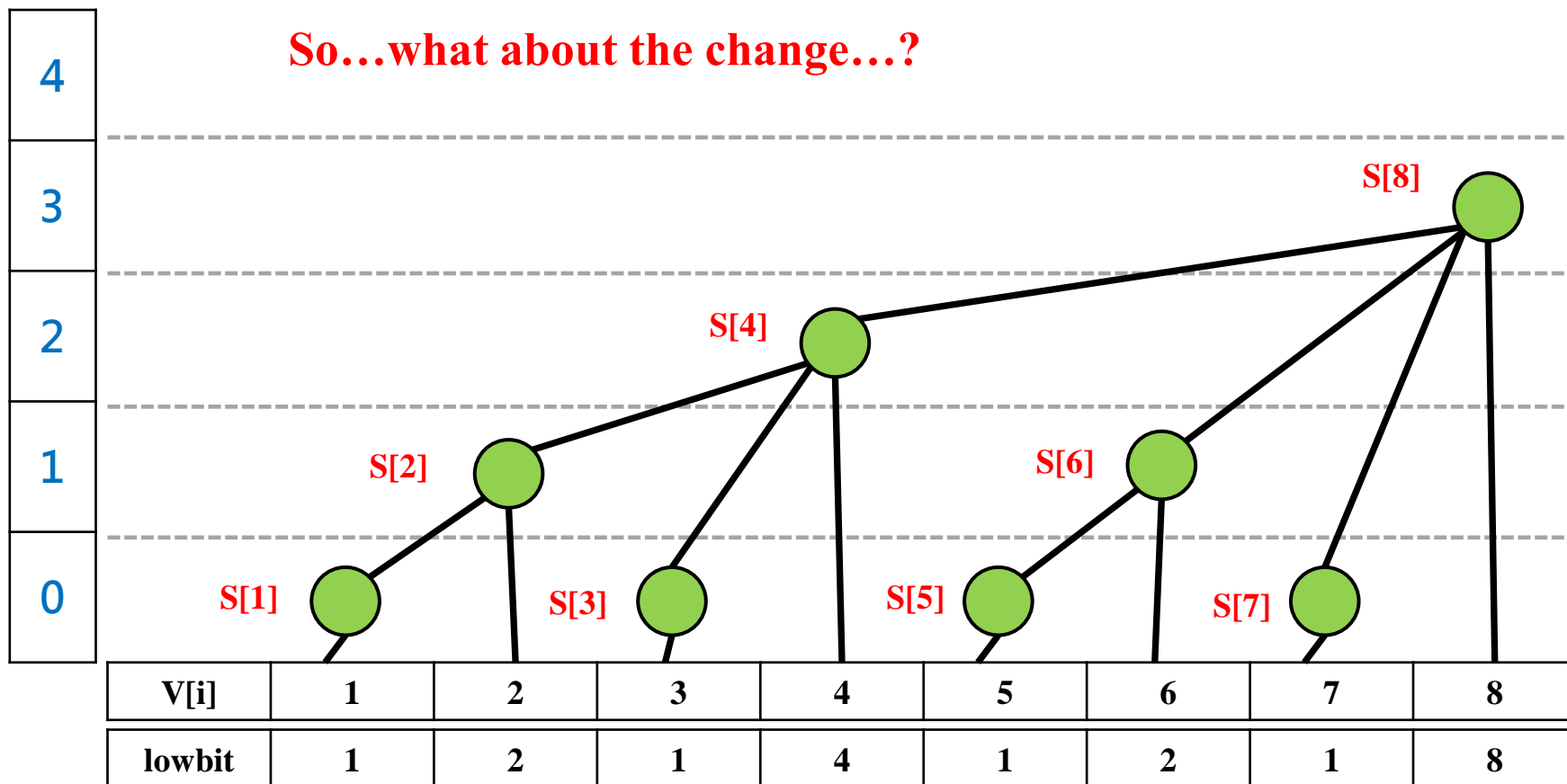
```
int lowbit (int in)
{
    return in&(-in);
}
```

```
int getsum (int end)
{
    int ans = 0;
    while(end>0)
    {
        ans += s[end];
        end -= lowbit(end);
    }
}
```



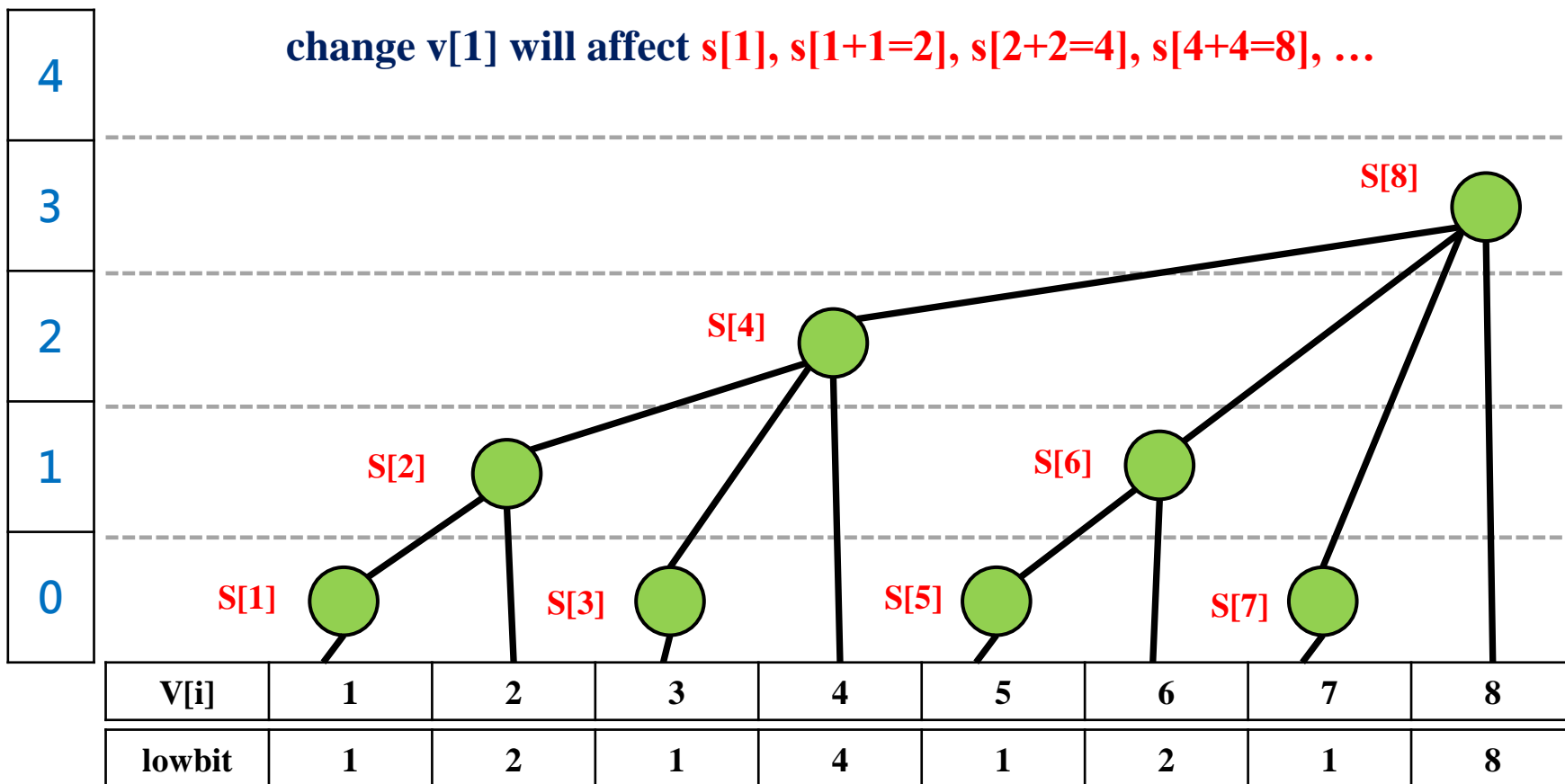
Binary Indexed Tree

So...what about the change...?



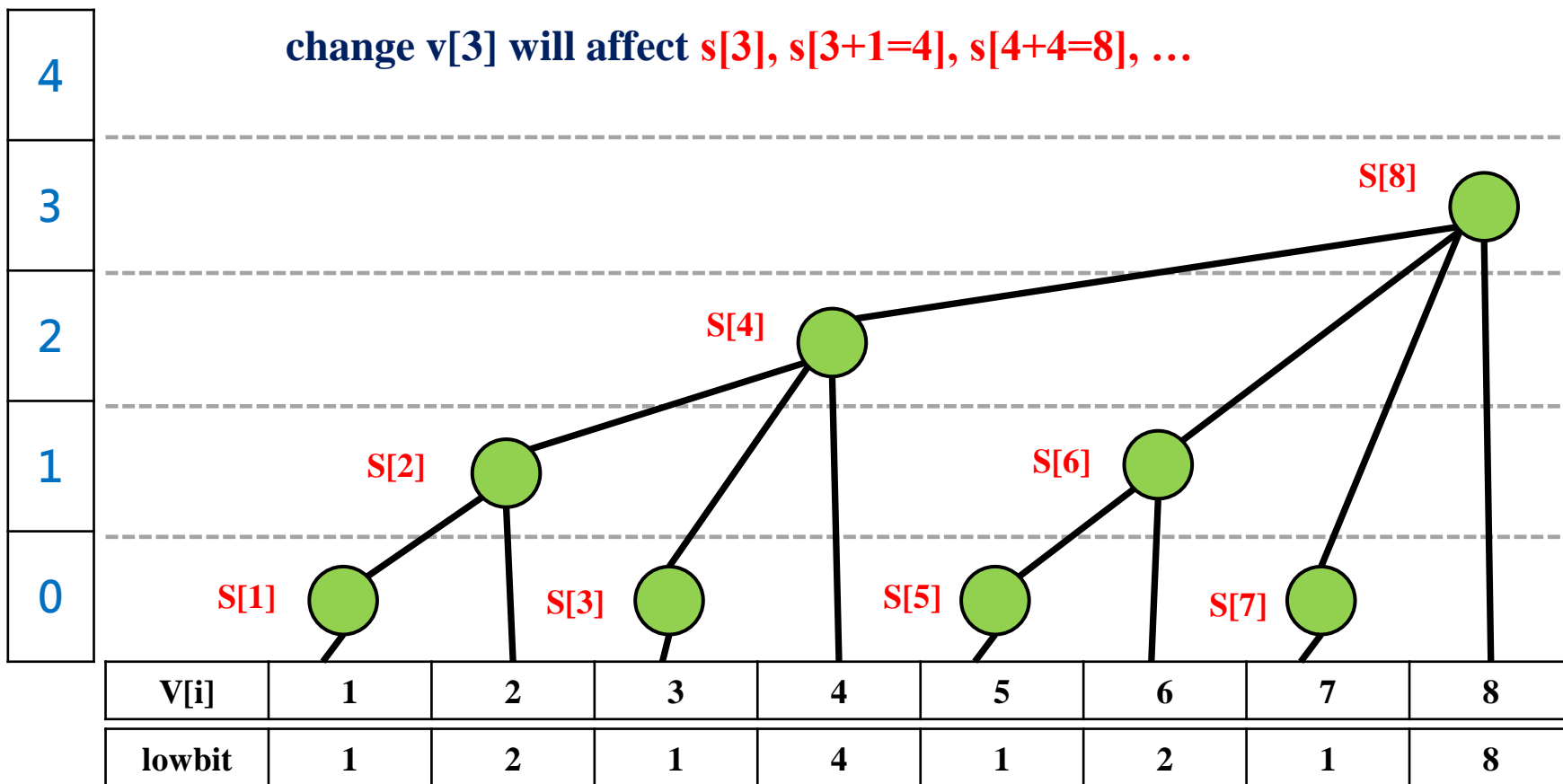
Binary Indexed Tree

change $v[1]$ will affect $s[1]$, $s[1+1=2]$, $s[2+2=4]$, $s[4+4=8]$, ...



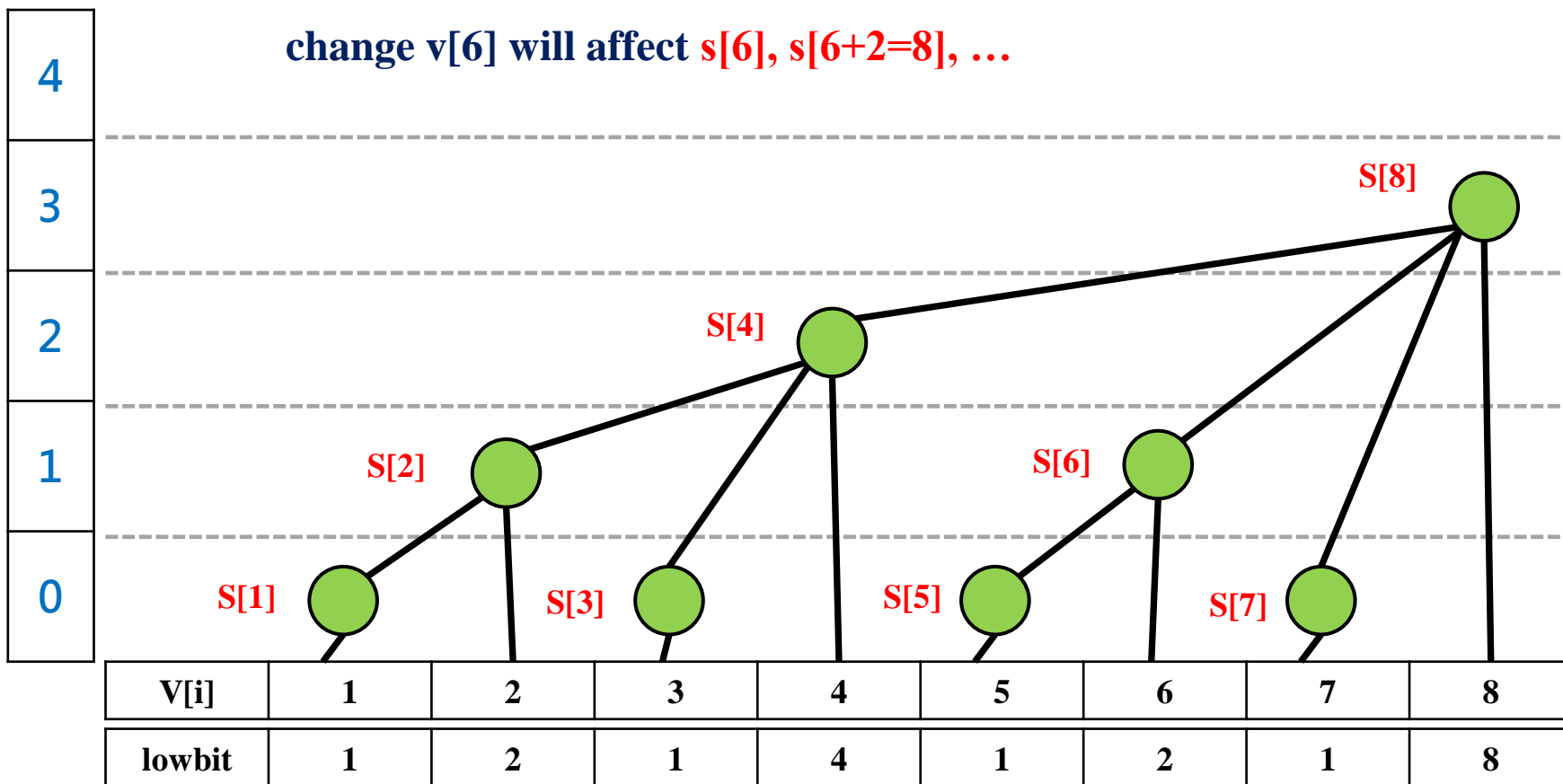
Binary Indexed Tree

change $v[3]$ will affect $s[3]$, $s[3+1=4]$, $s[4+4=8]$, ...



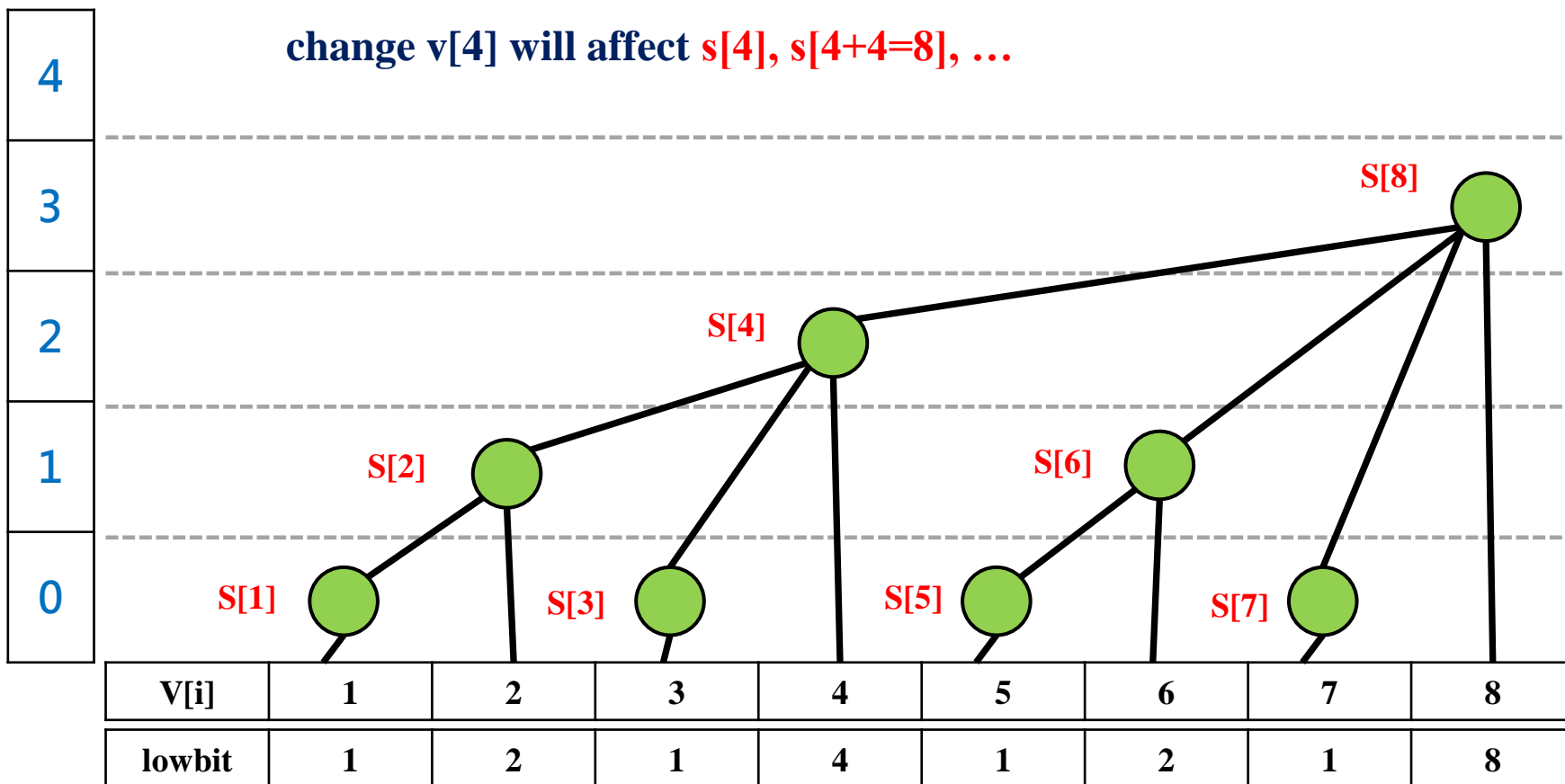
Binary Indexed Tree

change $v[6]$ will affect $s[6], s[6+2=8], \dots$



Binary Indexed Tree

change $v[4]$ will affect $s[4], s[4+4=8], \dots$



Binary Indexed Tree

Define:

```
int change (int end, int delta)
{
    for(int i=end; i<=maxsize; i+=lowbit(i))
        s[i] += delta;
}
```



Binary Indexed Tree

- How to find the summation between interval $[i..j]$?
 - call the subroutine “**getsum[j] – getsum[i-1]**”
- Expand the 1 dimension into 2 dimension by yourself
- Replace such routines with a segment tree by yourself



Example

- POJ 2352



HomeWork

pku-1195

PKU-3368

pku-3321

PKU-2528

pku-2155

PKU-2828

pku-2352

PKU-2777

pku-3067

PKU-2886

pku-2481

PKU-2750

pku-2299

PKU-2482

PKU-2352

