

Directed Graphs

Directed Graphs (Digraphs)

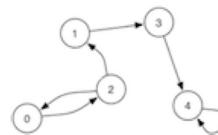
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In our previous discussion of graphs:

- an edge indicates a relationship between two vertices
- an edge indicates nothing more than a relationship

In many real-world applications of graphs:

- edges are directional ($v \rightarrow w \neq w \rightarrow v$)

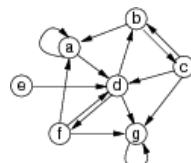


- edges have a *weight* (cost to go from $v \rightarrow w$)

... Directed Graphs (Digraphs)

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Example digraph and adjacency matrix representation:



	a	b	c	d	e	f	g
a	1	0	0	1	0	0	0
b	1	0	1	0	0	0	0
c	0	1	0	1	0	0	1
d	0	1	0	0	0	1	1
e	0	0	0	1	0	0	0
f	1	0	0	1	0	0	1
g	0	0	0	0	0	0	1

Undirectional \Rightarrow symmetric matrix

Directional \Rightarrow non-symmetric matrix

Maximum #edges in a digraph with V vertices: V^2

... Directed Graphs (Digraphs)

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Terminology for digraphs ...

Directed path: sequence of $n \geq 2$ vertices $v_1 \rightarrow v_2 \rightarrow \dots \rightarrow v_n$

- where $(v_i, v_{i+1}) \in \text{edges}(G)$ for all v_i, v_{i+1} in sequence
- if $v_1 = v_n$, we have a *directed cycle*

Digraph Applications

Potential application areas:

Domain	Vertex	Edge
Web	web page	hyperlink
scheduling	task	precedence
chess	board position	legal move
science	journal article	citation
dynamic data	malloc'd object	pointer
programs	function	function call
make	file	dependency

... Digraph Applications

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Problems to solve on digraphs:

- is there a directed path from s to t ? (transitive closure)
- what is the shortest path from s to t ? (shortest path)
- are all vertices mutually reachable? (strong connectivity)
- how to organise a set of tasks? (topological sort)
- which web pages are "important"? (PageRank)
- how to build a web crawler? (graph traversal)

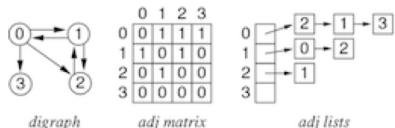
Digraph Representation

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Similar set of choices as for undirectional graphs:

- array of edges (directed)
- vertex-indexed adjacency matrix (non-symmetric)
- vertex-indexed adjacency lists

V vertices identified by $0 \dots V-1$



Reachability

Transitive Closure

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Given a digraph G it is potentially useful to know

- is vertex t reachable from vertex s ?

Example applications:

- can I complete a schedule from the current state?
- is a malloc'd object being referenced by any pointer?

How to compute transitive closure?

... Transitive Closure

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One possibility:

- implement it via `hasPath(G, s, t)` (itself implemented by DFS or BFS algorithm)
- feasible if $\text{reachable}(G, s, t)$ is infrequent operation

What if we have an algorithm that frequently needs to check reachability?

Would be very convenient/efficient to have:

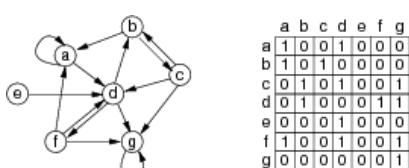
```
reachable(G, s, t):
| return G.tc[s][t] // transitive closure matrix
```

Of course, if V is very large, then this is not feasible.

Exercise #1: Transitive Closure Matrix

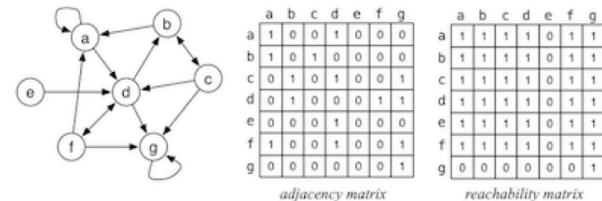
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Which reachable $s \dots t$ exist in the following graph?



	a	b	c	d	e	f	g
a	1	0	0	1	0	0	0
b	1	0	1	0	0	0	0
c	0	1	0	1	0	0	1
d	0	1	0	0	0	1	1
e	0	0	0	1	0	0	0
f	1	0	0	1	0	0	1
g	0	0	0	0	0	0	1

Transitive closure of example graph:



... Transitive Closure

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Goal: produce a matrix of reachability values

- if $tc[s][t]$ is 1, then t is reachable from s
- if $tc[s][t]$ is 0, then t is not reachable from s

So, how to create this matrix?

Observation:

$\forall i, s, t \in \text{vertices}(G):$

$$(s, i) \in \text{edges}(G) \text{ and } (i, t) \in \text{edges}(G) \Rightarrow tc[s][t] = 1$$

$tc[s][t] = 1$ if there is a path from s to t of length 2 ($s \rightarrow i \rightarrow t$)

... Transitive Closure

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If we implement the above as:

```
make tc[][] a copy of edges[][]
for all i in vertices(G) do
    for all s in vertices(G) do
        for all t in vertices(G) do
            if tc[s][i]=1 and tc[i][t]=1 then
                tc[s][t]=1
            end if
        end for
    end for
end for
```

then we get an algorithm to convert `edges` into a `tc`

This is known as *Warshall's algorithm*

... Transitive Closure

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How it works ...

After iteration 1, $\text{tc}[s][t]$ is 1 if

- either $s \rightarrow t$ exists or $s \rightarrow 0 \rightarrow t$ exists

After iteration 2, $\text{tc}[s][t]$ is 1 if any of the following exist

- $s \rightarrow t$ or $s \rightarrow 0 \rightarrow t$ or $s \rightarrow 1 \rightarrow t$ or $s \rightarrow 0 \rightarrow 1 \rightarrow t$ or $s \rightarrow 1 \rightarrow 0 \rightarrow t$

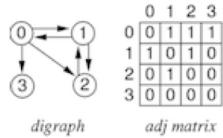
Etc. ... so after the V^{th} iteration, $\text{tc}[s][t]$ is 1 if

- there is any directed path in the graph from s to t

Exercise #2: Transitive Closure

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Trace Warshall's algorithm on the following graph:



1st iteration $i=0$:

tc	[0]	[1]	[2]	[3]
[0]	0	1	1	1
[1]	1	1	1	1
[2]	0	1	0	0
[3]	0	0	0	0

2nd iteration $i=1$:

tc	[0]	[1]	[2]	[3]
[0]	1	1	1	1
[1]	1	1	1	1
[2]	1	1	1	1
[3]	0	0	0	0

3rd iteration $i=2$: unchanged

4th iteration $i=3$: unchanged

... Transitive Closure

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Cost analysis:

- storage: additional V^2 items (each item may be 1 bit)

- computation of transitive closure: $O(V^3)$

- computation of `reachable()`: $O(1)$ after having generated `tc[][]`

Amortisation: would need many calls to `reachable()` to justify other costs

Alternative: use DFS in each call to `reachable()`

Cost analysis:

- storage: cost of queue and set during reachable

- computation of `reachable()`: cost of DFS = $O(V^2)$ (for adjacency matrix)

Digraph Traversal

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Same algorithms as for undirected graphs:

`depthFirst(v) :`

1. mark v as visited
2. for each $(v, w) \in edges(G)$ do
 - if w has not been visited then
`depthFirst(w)`

`breadth-first(v) :`

1. enqueue v
2. while queue not empty do
 - dequeue v
 - if v not already visited then
 - mark v as visited
 - enqueue each vertex w adjacent to v

Example: Web Crawling

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Goal: visit every page on the web

Solution: breadth-first search with "implicit" graph

```
webCrawl(startingURL):  
    mark startingURL as alreadySeen  
    enqueue(Q, startingURL)  
    while Q is not empty do  
        nextPage=dequeue(Q)  
        visit nextPage  
        for each hyperlink on nextPage do  
            if hyperlink not alreadySeen then  
                mark hyperlink as alreadySeen  
                enqueue(Q, hyperlink)  
            end if  
        end for
```

| end while

visit scans page and collects e.g. keywords and links

Weighted Graphs

Weighted Graphs

Graphs so far have considered

- edge = an association between two vertices/nodes
- may be a precedence in the association (directed)

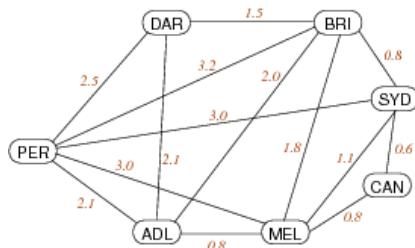
Some applications require us to consider

- a *cost* or *weight* of an association
- modelled by assigning values to edges (e.g. positive reals)

Weights can be used in both directed and undirected graphs.

... Weighted Graphs

Example: major airline flight routes in Australia



Representation: edge = direct flight; weight = approx flying time (hours)

... Weighted Graphs

Weights lead to minimisation-type questions, e.g.

1. Cheapest way to connect all vertices?

- a.k.a. *minimum spanning tree* problem
- assumes: edges are weighted and undirected

2. Cheapest way to get from A to B?

- a.k.a *shortest path* problem

- assumes: edge weights positive, directed or undirected

Exercise #3: Implementing a Route Finder

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If we represent a street map as a graph

- what are the vertices?
- what are the edges?
- are edges directional?
- what are the weights?
- are the weights fixed?

Weighted Graph Representation

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Weights can easily be added to:

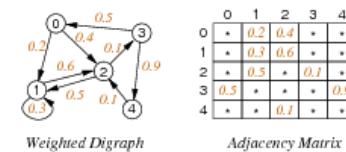
- adjacency matrix representation ($0/1 \rightarrow$ int or float)
- adjacency lists representation (add int/float to list node)

Both representations work whether edges are directed or not.

... Weighted Graph Representation

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Adjacency matrix representation with weights:

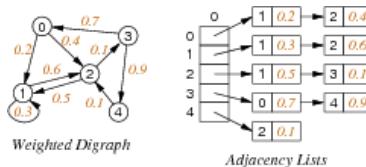


Note: need distinguished value to indicate "no edge".

... Weighted Graph Representation

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Adjacency lists representation with weights:



Note: if undirected, each edge appears twice with same weight

... Weighted Graph Representation

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Sample adjacency matrix implementation in C requires minimal changes to previous Graph ADT:

WGraph.h

```
// edges are pairs of vertices (end-points) plus positive weight
typedef struct Edge {
    Vertex v;
    Vertex w;
    int weight;
} Edge;

// returns weight, or 0 if vertices not adjacent
int adjacent(Graph, Vertex, Vertex);
```

... Weighted Graph Representation

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WGraph.c

```
typedef struct GraphRep {
    int **edges; // adjacency matrix storing positive weights
                  // 0 if nodes not adjacent
    int nV;      // #vertices
    int nE;      // #edges
} GraphRep;

void insertEdge(Graph g, Edge e) {
    assert(g != NULL && validV(g,e.v) && validV(g,e.w));
    if (g->edges[e.v][e.w] == 0) { // edge e not in graph
        g->edges[e.v][e.w] = e.weight;
        g->nE++;
    }
}

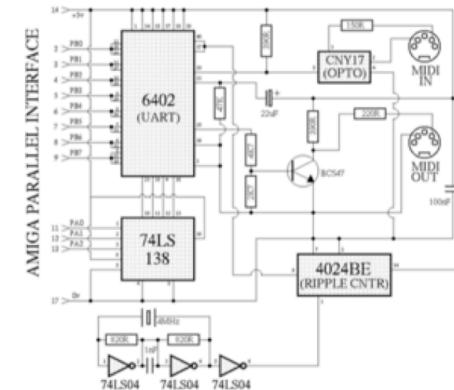
int adjacent(Graph g, Vertex v, Vertex w) {
    assert(g != NULL && validV(g,v) && validV(g,w));
    return g->edges[v][w];
}
```

Minimum Spanning Trees

Exercise #4: Minimising Wires in Circuits

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Electronic circuit designs often need to make the pins of several components electrically equivalent by wiring them together.



To interconnect a set of n pins we can use an arrangement of $n-1$ wires each connecting two pins.

What kind of algorithm would ...

- help us find the arrangement with the least amount of wire?

Minimum Spanning Trees

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Reminder: *Spanning tree ST* of graph $G=(V,E)$

- *spanning* = all vertices, *tree* = no cycles
 - ST is a subgraph of G ($G'=(V,E')$ where $E' \subseteq E$)
 - ST is *connected* and *acyclic*

Minimum spanning tree MST of graph G

- MST is a spanning tree of G
- sum of edge weights is no larger than any other ST

Applications: Computer networks, Electrical grids, Transportation networks ...

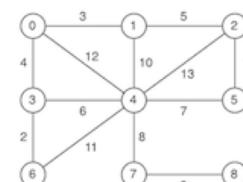
Problem: how to (efficiently) find MST for graph G ?

NB: MST may not be unique (e.g. all edges have same weight \Rightarrow every ST is MST)

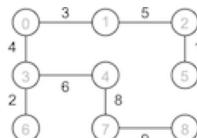
... Minimum Spanning Trees

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Example:



An MST ...



... Minimum Spanning Trees

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Brute force solution:

```
findMST(G):
    Input  graph G
    Output a minimum spanning tree of G

    bestCost=∞
    for all spanning trees t of G do
        if cost(t)<bestCost then
            bestTree=t
            bestCost=cost(t)
        end if
    end for
    return bestTree
```

Example of *generate-and-test* algorithm.

Not useful because #spanning trees is potentially large (e.g. n^{n-2} for a complete graph with n vertices)

... Minimum Spanning Trees

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Simplifying assumption:

- edges in G are not directed (MST for digraphs is harder)

Kruskal's Algorithm

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One approach to computing MST for graph G with V nodes:

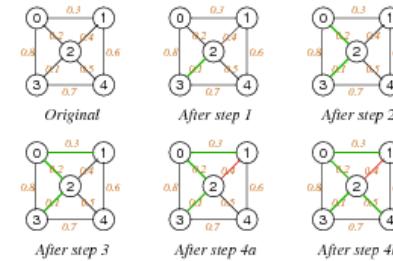
1. start with empty MST
2. consider edges in increasing weight order
 - add edge if it does not form a cycle in MST
3. repeat until $V-1$ edges are added

Critical operations:

- iterating over edges in weight order
- checking for cycles in a graph

... Kruskal's Algorithm

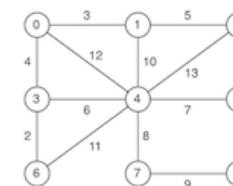
Execution trace of Kruskal's algorithm:



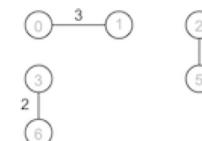
Exercise #5: Kruskal's Algorithm

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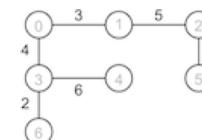
Show how Kruskal's algorithm produces an MST on:



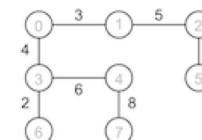
After 3rd iteration:



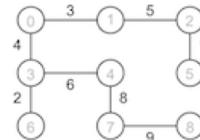
After 6th iteration:



After 7th iteration:



After 8th iteration ($V=8$ edges added):



... Kruskal's Algorithm

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Pseudocode:

```
KruskalMST(G):
    Input  graph G with n nodes
    Output a minimum spanning tree of G

    MST=empty graph
    sort edges(G) by weight
    for each e $\in$ sortedEdgeList do
        MST = MST U {e}
        if MST has a cycle then
            MST = MST \ {e}
        end if
        if MST has n-1 edges then
            return MST
        end if
    end for
```

... Kruskal's Algorithm

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Rough time complexity analysis ...

- sorting edge list is $O(E \cdot \log E)$
- at least V iterations over sorted edges
- on each iteration ...
 - getting next lowest cost edge is $O(1)$
 - checking whether adding it forms a cycle: cost = ??

Possibilities for cycle checking:

- use DFS ... too expensive?
- could use *Union-Find data structure* (see Sedgewick Ch.1)

Prim's Algorithm

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Another approach to computing MST for graph $G=(V,E)$:

1. start from any vertex v and empty MST
2. choose edge not already in MST to add to MST

- must be incident on a vertex s already connected to v in MST
- must be incident on a vertex t not already connected to v in MST
- must have minimal weight of all such edges

3. repeat until MST covers all vertices

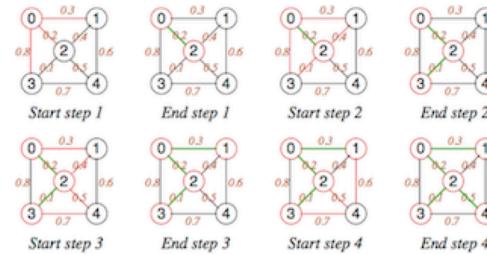
Critical operations:

- checking for vertex being connected in a graph
- finding min weight edge in a set of edges

... Prim's Algorithm

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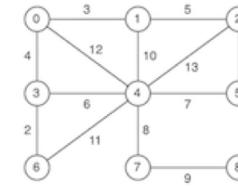
Execution trace of Prim's algorithm (starting at $s=0$):



Exercise #6: Prim's Algorithm

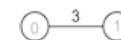
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Show how Prim's algorithm produces an MST on:

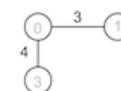


Start from vertex 0

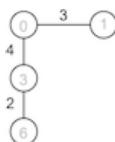
After 1st iteration:



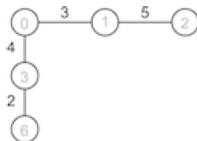
After 2nd iteration:



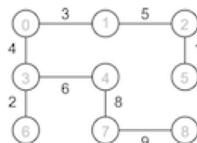
After 3rd iteration:



After 4th iteration:



After 8th iteration (all vertices covered):



... Prim's Algorithm

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Pseudocode:

```

PrimMST(G):
    Input graph G with n nodes
    Output a minimum spanning tree of G

    MST=empty graph
    usedV={0}
    unusedE=edges(g)
    while |usedV|<n do
        find e=(s,t,w)∈unusedE such that {
            s∈usedV, t∉usedV and w is min weight of all such edges
        }
        MST = MST ∪ {e}
        usedV = usedV ∪ {t}
        unusedE = unusedE \ {e}
    end while
    return MST

```

Critical operation: finding best edge

... Prim's Algorithm

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Rough time complexity analysis ...

- V iterations of outer loop

- in each iteration ...
 - find min edge with set of edges is $O(E) \Rightarrow O(V \cdot E)$ overall
 - find min edge with priority queue is $O(\log E) \Rightarrow O(V \cdot \log E)$ overall

Sidetrack: Priority Queues

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Some applications of queues require

- items processed in order of "priority"
- rather than in order of entry (FIFO — first in, first out)

Priority Queues (PQueues) provide this via:

- **join**: insert item into PQueue with an associated priority (replacing enqueue)
- **leave**: remove item with highest priority (replacing dequeue)

Time complexity for naive implementation of a PQueue containing N items ...

- $O(1)$ for join $O(N)$ for leave

Most efficient implementation ("heap") ...

- $O(\log N)$ for join, leave

Other MST Algorithms

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Boruvka's algorithm ... complexity $O(E \cdot \log V)$

- the oldest MST algorithm
- start with V separate components
- join components using min cost links
- continue until only one single component

Karger, Klein, and Tarjan ... complexity $O(E)$

- based on Boruvka, but non-deterministic
- randomly selects subset of edges to consider
- for the keen, here's [the paper](#) describing the algorithm

Shortest Path

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Shortest Path

$Path =$ sequence of edges in graph $G \quad p = (v_0, v_1), (v_1, v_2), \dots, (v_{m-1}, v_m)$

$cost(path) =$ sum of edge weights along path

$Shortest\ path$ between vertices s and t

- a simple path $p(s,t)$ where $s = \text{first}(p)$, $t = \text{last}(p)$
- no other simple path $q(s,t)$ has $\text{cost}(q) < \text{cost}(p)$

Assumptions: weighted digraph, no negative weights.

Finding shortest path between two given nodes known as *source-target* SP problem

Variations: *single-source* SP, *all-pairs* SP

Applications: navigation, routing in data networks, ...

Single-source Shortest Path (SSSP)

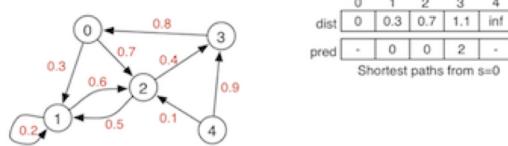
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Given: weighted digraph G , source vertex s

Result: shortest paths from s to all other vertices

- $\text{dist}[]$ V -indexed array of cost of shortest path from s
- $\text{pred}[]$ V -indexed array of predecessor in shortest path from s

Example:



Edge Relaxation

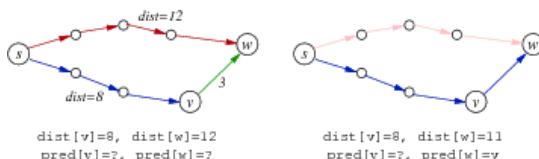
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Assume: $\text{dist}[]$ and $\text{pred}[]$ as above (but containing data for shortest paths *discovered so far*)

$\text{dist}[v]$ is length of shortest known path from s to v

$\text{dist}[w]$ is length of shortest known path from s to w

Relaxation updates data for w if we find a shorter path from s to w :



Relaxation along edge $e=(v,w,\text{weight})$:

- if $\text{dist}[v]+\text{weight} < \text{dist}[w]$ then update $\text{dist}[w]:=\text{dist}[v]+\text{weight}$ and $\text{pred}[w]:=v$

Dijkstra's Algorithm

One approach to solving single-source shortest path problem ...

Data: $G, s, \text{dist}[], \text{pred}[]$ and

- $vSet$: set of vertices whose shortest path from s is unknown

Algorithm:

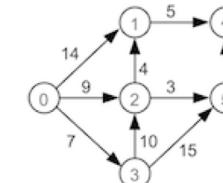
```
dist[] // array of cost of shortest path from s
pred[] // array of predecessor in shortest path from s
```

```
dijkstraSSSP(G,source):
    Input graph G, source node
    initialise dist[] to all ∞, except dist[source]=0
    initialise pred[] to all -1
    vSet=all vertices of G
    while vSet≠∅ do
        | find s∈vSet with minimum dist[s]
        | for each (s,t,w)∈edges(G) do
            |   relax along (s,t,w)
        | end for
        vSet=vSet\{s}
    end while
```

Exercise #7: Dijkstra's Algorithm

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Show how Dijkstra's algorithm runs on (source node = 0):



	[0]	[1]	[2]	[3]	[4]	[5]
dist	0	∞	∞	∞	∞	∞
pred	-	-	-	-	-	-

	[0]	[1]	[2]	[3]	[4]	[5]
dist	0	∞	∞	∞	∞	∞

pred	-	-	-	-	-	-
------	---	---	---	---	---	---

dist	0	14	9	7	∞	∞
pred	-	0	0	0	-	-

dist	0	14	9	7	∞	22
pred	-	0	0	0	-	3

dist	0	13	9	7	∞	12
pred	-	2	0	0	-	2

dist	0	13	9	7	20	12
pred	-	2	0	0	5	2

dist	0	13	9	7	18	12
pred	-	2	0	0	1	2

... Dijkstra's Algorithm

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Why Dijkstra's algorithm is correct:

Hypothesis.

- (a) For visited $s \dots dist[s]$ is shortest distance from source
- (b) For unvisited $t \dots dist[t]$ is shortest distance from source via visited nodes

Proof.

Base case: no visited nodes, $dist[\text{source}] = 0$, $dist[s] = \infty$ for all other nodes

Induction step:

1. If s is unvisited node with minimum $dist[s]$, then $dist[s]$ is shortest distance from source to s :
 - o if \exists shorter path via only visited nodes, then $dist[s]$ would have been updated when processing the predecessor of s on this path
 - o if \exists shorter path via an unvisited node u , then $dist[u] < dist[s]$, which is impossible if s has min distance of all unvisited nodes
2. This implies that (a) holds for s after processing s
3. (b) still holds for all unvisited nodes t after processing s :
 - o if \exists shorter path via s we would have just updated $dist[t]$
 - o if \exists shorter path without s we would have found it previously

... Dijkstra's Algorithm

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Time complexity analysis ...

Each edge needs to be considered once $\Rightarrow O(E)$.

Outer loop has $O(V)$ iterations.

Implementing "find sevSet with minimum $dist[s]$ "

1. try all $s \in \text{vSet} \Rightarrow \text{cost} = O(V) \Rightarrow \text{overall cost} = O(E + V^2) = O(V^2)$
2. using a PQqueue to implement extracting minimum
 - o can improve overall cost to $O(E + V \cdot \log V)$ (for best-known implementation)

All-pair Shortest Path (APSP)

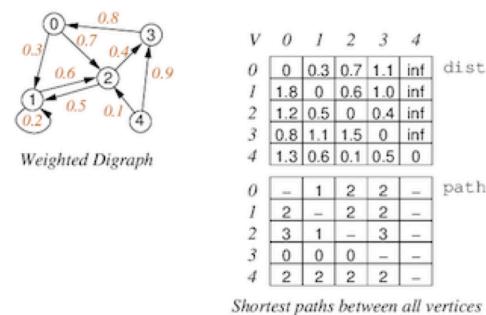
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Given: weighted digraph G

Result: shortest paths between all pairs of vertices

- $\text{dist}[\][]$ $V \times V$ -indexed matrix of cost of shortest path from v_{row} to v_{col}
- $\text{path}[\][]$ $V \times V$ -indexed matrix of next node in shortest path from v_{row} to v_{col}

Example:



Floyd's Algorithm

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One approach to solving all-pair shortest path problem...

Data: $G, \text{dist}[\][], \text{path}[\][]$ Algorithm:

```
dist[][] // array of cost of shortest path from s to t
path[][] // array of next node after s on shortest path from s to t
```

floydAPSP(G):

Input graph G

```
initialise dist[s][t]=0 for each s=t
          =w for each (s,t,w) ∈ edges(G)
          =∞ otherwise
initialise path[s][t]=t for each (s,t,w) ∈ edges(G)
          =-1 otherwise
for all i ∈ vertices(G) do
  | for all s ∈ vertices(G) do
```

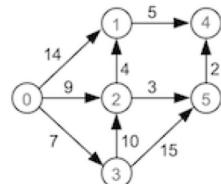
```

for all t in vertices(G) do
    if dist[s][i] + dist[i][t] < dist[s][t] then
        dist[s][t] = dist[s][i] + dist[i][t]
        path[s][t] = path[s][i]
    end if
end for
end for

```

Exercise #8: Floyd's Algorithm

Show how Floyd's algorithm runs on:



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dist	[0]	[1]	[2]	[3]	[4]	[5]	path	[0]	[1]	[2]	[3]	[4]	[5]
[0]	0	14	9	7			[0]		1	2	3		
[1]		0			5		[1]				4		
[2]		4	0			3	[2]		1			5	
[3]			10	0		15	[3]			2			5
[4]					0		[4]						
[5]					2	0	[5]					4	

After 1st iteration i=0: unchanged After 2nd iteration i=1:

dist	[0]	[1]	[2]	[3]	[4]	[5]	path	[0]	[1]	[2]	[3]	[4]	[5]
[0]	0	14	9	7	19	∞	[0]	-	1	2	3	1	-
[1]	∞	0	∞	∞	5	∞	[1]	-	-	-	-	4	-
[2]	∞	4	0	∞	9	3	[2]	-	1	-	-	1	5
[3]	∞	∞	10	0	∞	15	[3]	-	-	2	-	-	5
[4]	∞	∞	∞	∞	0	∞	[4]	-	-	-	-	-	-
[5]	∞	∞	∞	∞	2	0	[5]	-	-	-	-	4	-

After 3rd iteration i=2:

dist	[0]	[1]	[2]	[3]	[4]	[5]	path	[0]	[1]	[2]	[3]	[4]	[5]
[0]	0	13	9	7	18	12	[0]	-	2	2	3	2	2
[1]	∞	0	∞	∞	5	∞	[1]	-	-	-	-	4	-
[2]	∞	4	0	∞	9	3	[2]	-	1	-	-	1	5
[3]	∞	14	10	0	19	13	[3]	-	2	2	-	2	2
[4]	∞	∞	∞	∞	0	∞	[4]	-	-	-	-	-	-
[5]	∞	∞	∞	∞	2	0	[5]	-	-	-	-	4	-

After 4th iteration i=3: unchanged After 5th iteration i=4: unchanged After 6th iteration i=5:

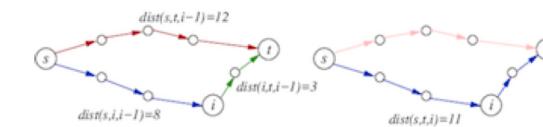
dist	[0]	[1]	[2]	[3]	[4]	[5]	path	[0]	[1]	[2]	[3]	[4]	[5]
[0]	0	13	9	7	14	12	[0]	-	2	2	3	2	2
[1]	∞	0	∞	∞	5	∞	[1]	-	-	-	-	4	-
[2]	∞	4	0	∞	5	3	[2]	-	1	-	-	5	5
[3]	∞	14	10	0	15	13	[3]	-	2	2	-	2	2
[4]	∞	∞	∞	∞	0	∞	[4]	-	-	-	-	-	-
[5]	∞	∞	∞	∞	2	0	[5]	-	-	-	-	4	-

... Floyd's Algorithm

Why Floyd's algorithm is correct:

A shortest path from s to t using only nodes from $\{0, \dots, i\}$ is the shorter of

- a shortest path from s to t using only nodes from $\{0, \dots, i-1\}$
- a shortest path from s to i using only nodes from $\{0, \dots, i-1\}$
plus a shortest path from i to t using only nodes from $\{0, \dots, i-1\}$



Also known as Floyd-Warshall algorithm (can you see why?)

... Floyd's Algorithm

Cost analysis ...

- initialising $\text{dist}[\cdot][\cdot]$, $\text{path}[\cdot][\cdot] \Rightarrow O(E)$

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- V iterations to update $\text{dist}[\cdot][\cdot]$, $\text{path}[\cdot][\cdot] \Rightarrow O(V^3)$

Time complexity of Floyd's algorithm: $O(V^3)$ (same as Warshall's algorithm for transitive closure)

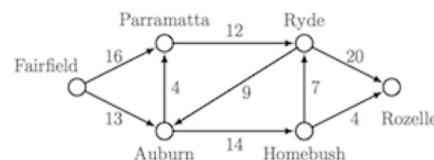
Network Flow

Exercise #9: Merchandise Distribution

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Lucky Cricket Company ...

- produces cricket balls in Fairfield
- has a warehouse in Rozelle that stocks them
- ships them from factory to warehouse by leasing space on trucks with limited capacity:



What kind of algorithm would ...

- help us find the maximum number of crates that can be shipped from Fairfield to Rozelle per day?

Flow Networks

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Flow network ...

- weighted graph $G=(V,E)$
- distinct nodes $s \in V$ (source), $t \in V$ (sink)

Edge weights denote *capacities* Applications:

- Distribution networks, e.g.
 - source: oil field
 - sink: refinery
 - edges: pipes
- Traffic flow

... Flow Networks

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Flow in a network $G=(V,E)$... nonnegative $f(v,w)$ for all vertices $v,w \in V$ such that

- $f(v,w) \leq \text{capacity}$ for each edge $e=(v,w, \text{capacity}) \in E$
- $f(v,w)=0$ if no edge between v and w
- total flow *into* a vertex = total flow *out of* a vertex:

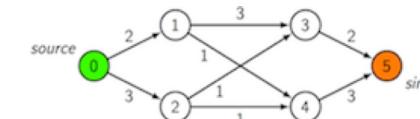
$$\sum_{x \in V} f(x, v) = \sum_{y \in V} f(v, y) \quad \text{for all } v \in V \setminus \{s,t\}$$

Maximum flow ... no other flow from s to t has larger value

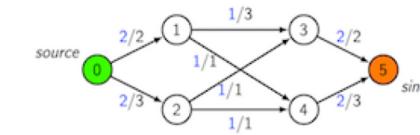
... Flow Networks

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Example:



A (maximum) flow ...



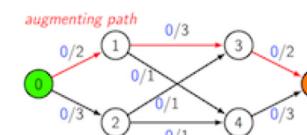
Augmenting Paths

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Assume ... $f(v,w)$ contains current flow

Augmenting path: any path from source s to sink t that can currently take more flow

Example:



Residual Network

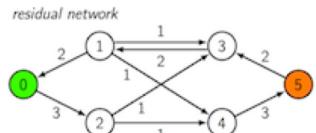
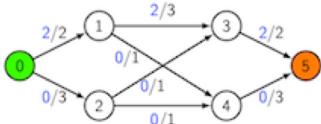
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Assume ... flow network $G=(V,E)$ and flow $f(v,w)$

Residual network (V,E') :

- same vertex set V
- for each edge $v \rightarrow^c w \in E$...
 - $f(v,w) < c \Rightarrow$ add edge $(v \rightarrow^{c-f(v,w)} w)$ to E'
 - $f(v,w) > 0 \Rightarrow$ add edge $(v \leftarrow^{f(v,w)} w)$ to E'

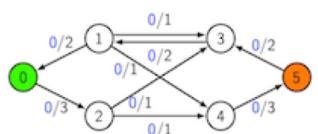
Example:



Exercise #10: Augmenting Paths and Residual Networks

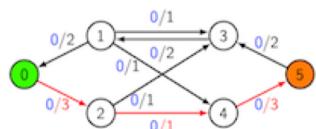
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Find an augmenting path in:



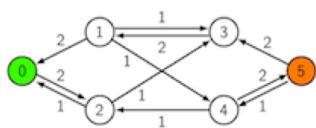
and show the residual network after augmenting the flow

1. Augmenting path:



maximum additional flow = 1

2. Residual network:



Can you find a further augmenting path in the new residual network?

Edmonds-Karp Algorithm

One approach to solving maximum flow problem ...

maxflow(G) :

1. Find a shortest augmenting path
2. Update $\text{flow}[][]$ so as to represent residual graph
3. Repeat until no augmenting path can be found

... Edmonds-Karp Algorithm

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Algorithm:

```
flow[][] // VxV array of current flow
visited[] /* array of predecessor nodes on shortest path
from source to sink in residual network */
```

maxflow(G) :

```
Input flow network G with source s and sink t
Output maximum flow value

initialise flow[v][w]=0 for all vertices v, w
maxflow=0
while  $\exists$  shortest augmenting path visited[] from s to t do
| df = maximum additional flow via visited[]
// adjust flow so as to represent residual graph
v=t
while v $\neq$ s do
| flow[visited[v]][v] = flow[visited[v]][v] + df;
| flow[v][visited[v]] = flow[v][visited[v]] - df;
| v=visited[v]
end while
maxflow=maxflow+df
end while
return maxflow
```

Shortest augmenting path can be found by standard BFS

... Edmonds-Karp Algorithm

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Time complexity analysis ...

- *Theorem.* The number of augmenting paths needed is at most $V \cdot E / 2$.
⇒ Outer loop has $O(V \cdot E)$ iterations.
- Finding augmenting path ⇒ $O(E)$ (consider only vertices connected to source and sink ⇒ $O(V+E)=O(E)$)

Overall cost of Edmonds-Karp algorithm: $O(V \cdot E^2)$

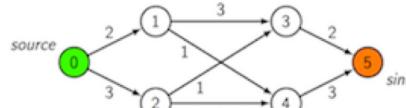
Note: Edmonds-Karp algorithm is an implementation of general *Ford-Fulkerson method*

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Exercise #11: Edmonds-Karp Algorithm

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Show how Edmonds-Karp algorithm runs on:



flow	[0]	[1]	[2]	[3]	[4]	[5]	c-f	[0]	[1]	[2]	[3]	[4]	[5]
[0]													
[1]													
[2]													
[3]													
[4]													
[5]													

flow	[0]	[1]	[2]	[3]	[4]	[5]	c-f	[0]	[1]	[2]	[3]	[4]	[5]
[0]	0	0	0	0	0	0							
[1]	0	0	0	0	0	0							
[2]	0	0	0	0	0	0							
[3]	0	0	0	0	0	0							
[4]	0	0	0	0	0	0							
[5]	0	0	0	0	0	0							

augmenting path: 0-1-3-5, df: 2

flow	[0]	[1]	[2]	[3]	[4]	[5]	c-f	[0]	[1]	[2]	[3]	[4]	[5]
[0]	0	2	0	0	0	0							
[1]	-2	0	0	2	0	0							
[2]	0	0	0	0	0	0							
[3]	0	-2	0	0	0	2							
[4]	0	0	0	0	0	0							
[5]	0	0	0	-2	0	0							

augmenting path: 0-2-4-5, df: 1

flow	[0]	[1]	[2]	[3]	[4]	[5]	c-f	[0]	[1]	[2]	[3]	[4]	[5]
[0]	0	2	1	0	0	0							
[1]	-2	0	0	2	0	0							
[2]	-1	0	0	0	1	0							
[3]	0	-2	0	0	0	2							
[4]	0	0	-1	0	0	1							
[5]	0	0	0	-2	-1	0							

augmenting path: 0-2-3-1-4-5, df: 1

flow	[0]	[1]	[2]	[3]	[4]	[5]	c-f	[0]	[1]	[2]	[3]	[4]	[5]
[0]	0	2	2	0	0	0							
[1]	-2	0	0	1	1	0							
[2]	-2	0	0	1	1	0							
[3]	0	-1	-1	0	0	2							
[4]	0	-1	-1	0	0	2							
[5]	0	0	0	-2	-2	0							

Digraph Applications

PageRank

81/86

Goal: determine which Web pages are "important"

Approach: ignore page contents; focus on hyperlinks

- treat Web as graph: page = vertex, hyperlink = directed edge
- pages with many incoming hyperlinks are important
- need to compute "incoming degree" for vertices

Problem: the Web is a *very* large graph

- approx. 10^{14} pages, 10^{15} hyperlinks

Assume for the moment that we could build a graph ...

Most frequent operation in algorithm "Does edge (v,w) exist?"

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... PageRank

Simple PageRank algorithm:

```
PageRank(myPage):  
    rank=0  
    for each page in the Web do  
        if linkExists(page,myPage) then  
            rank=rank+1  
        end if  
    end for
```

Note: requires *inbound* link check

... PageRank

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$V = \# \text{ pages in Web}$, $E = \# \text{ hyperlinks in Web}$

Costs for computing PageRank for each representation:

Representation	linkExists(v,w)	Cost
Adjacency matrix	edge[v][w]	I
Adjacency lists	inLL(list[v],w)	$\approx E/V$

Not feasible ...

- adjacency matrix ... $V \approx 10^{14} \Rightarrow$ matrix has 10^{28} cells
- adjacency list ... V lists, each with ≈ 10 hyperlinks $\Rightarrow 10^{15}$ list nodes

So how to really do it?

... PageRank

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Approach: the random web surfer

- if we randomly follow links in the web ...
- ... more likely to re-discover pages with many inbound links

```
curr=random page, prev=null  
for a long time do  
    if curr not in array ranked[] then  
        rank[curr]=0  
    end if  
    rank[curr]=rank[curr]+1  
    if random(0,100)<85 then // with 85% chance ...  
        prev=curr  
        curr=choose hyperlink from curr // ... crawl on  
    else
```

```
        curr=random page  
        prev=null  
    end if  
end for
```

// avoid getting stuck

Could be accomplished while we crawl web to build search index

Exercise #12: Implementing Facebook

85/86

Facebook could be considered as a giant "social graph"

- what are the vertices?
- what are the edges?
- are edges directional?

What kind of algorithm would ...

- help us find people that you might like to "befriend"?

Summary

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- Digraphs, weighted graphs: representations, applications
- Reachability
 - Warshall
- Minimum Spanning Tree (MST)
 - Kruskal, Prim
- Shortest path problems
 - Dijkstra (single source SPP)
 - Floyd (all-pair SSP)
- Flow networks
 - Edmonds-Karp (maximum flow)

- Suggested reading (Sedgewick):
 - digraphs ... Ch. 19.1-19.3
 - weighted graphs ... Ch. 20-20.1
 - MST ... Ch. 20.2-20.4
 - SSP ... Ch. 21-21.3
 - network flows ... Ch. 22.1-22.2

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