

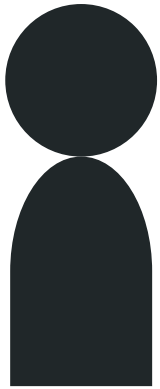
Classical Verification of Quantum Depth

Nai-Hui Chia (Rice University)

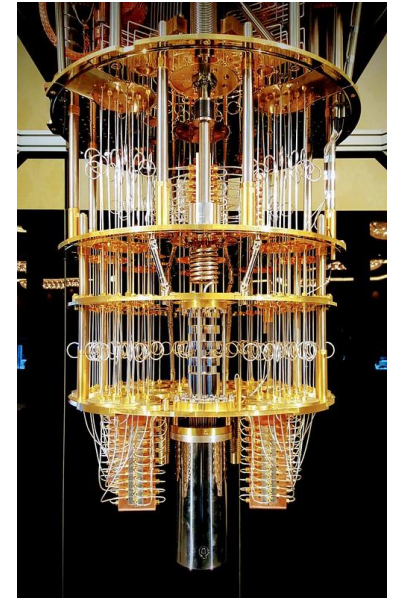
Based on joint work with Shih-Hang Hung

Certifying Quantum Resources

Can classical clients verify that a remote server has claimed quantum resources?



Classical client

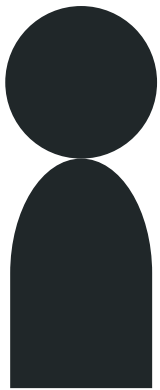


Remote server

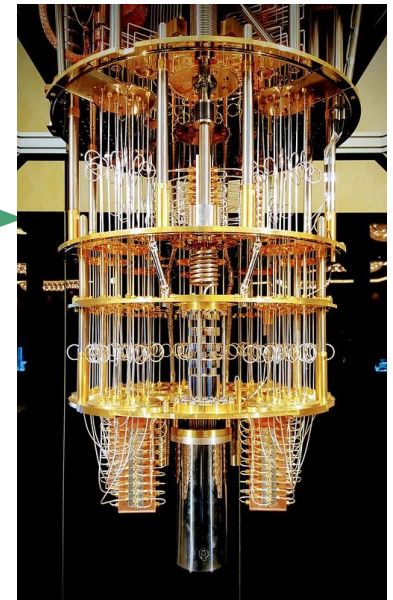
Certifying Quantum Resources

Can classical clients verify that a remote server has claimed quantum resources?

Do you have 500 qubits and 100 depths?



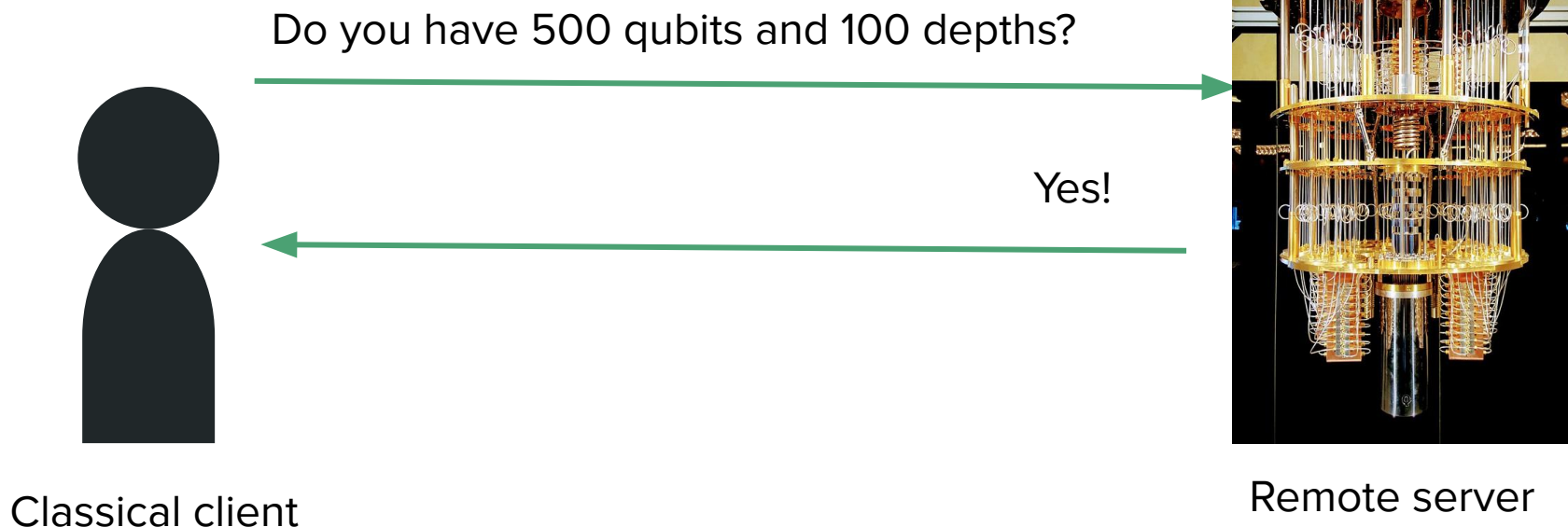
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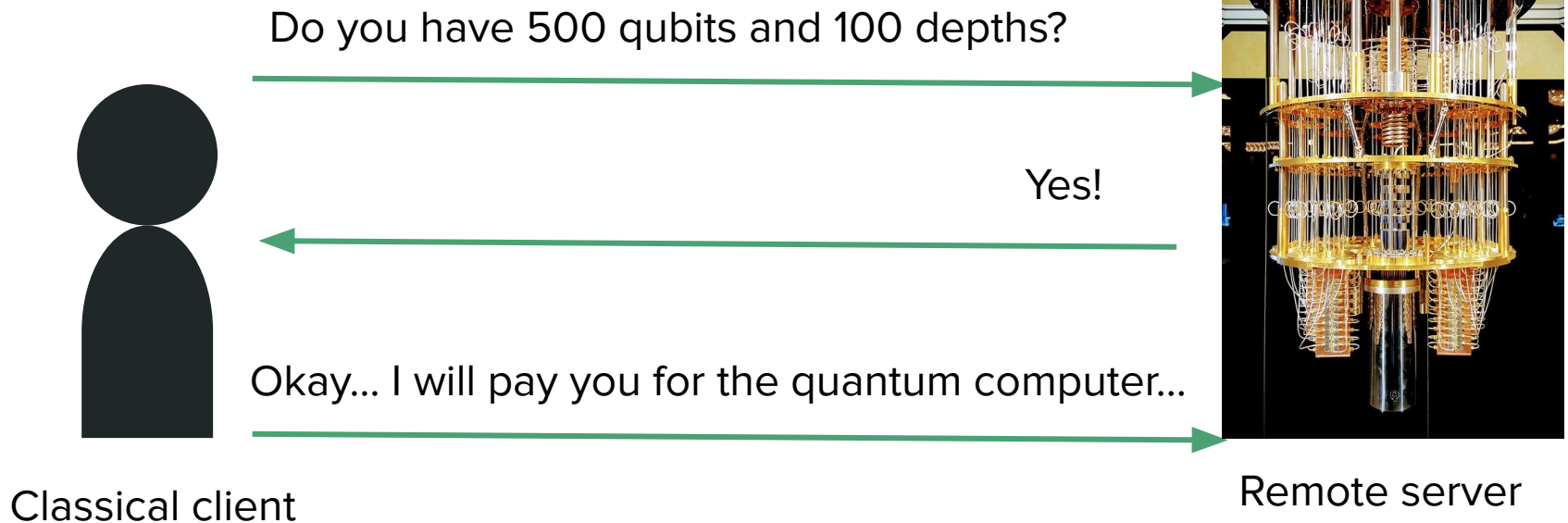
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Certifying Quantum Resources

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Certifying C

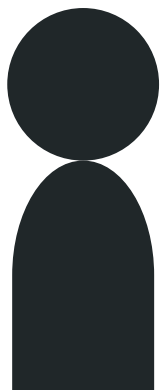
Actually, I only have 50 qubits and 10 depths. But you never know -- I can use some supercomputer to complete your task!

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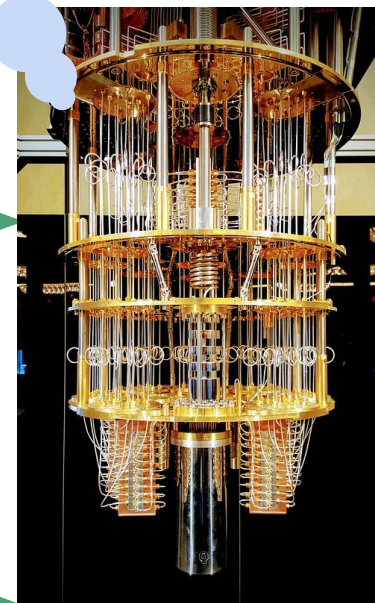
Do you have 500 qubits and 100 depths?

Yes!

Okay... I will pay you for the quantum computer...



Classical client



Remote server

Certifying C

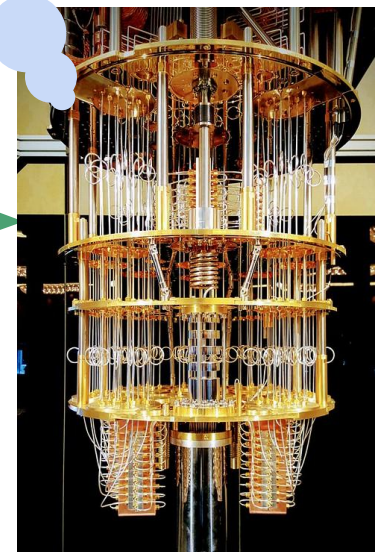
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We are especially interested in protocols for certifying
quantum depth

Why do we care about quantum depth?

Near-term QC:

- > 100 qubits
- Noisy gates
(not fault-tolerant)
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General QC:

- Many qubits
- Small noise (fault-tolerant)
+ long coherence time
=> **Poly(n) circuit depth**
- Can do all poly-time (quantum) algorithms

Building machine with large quantum depth is challenging

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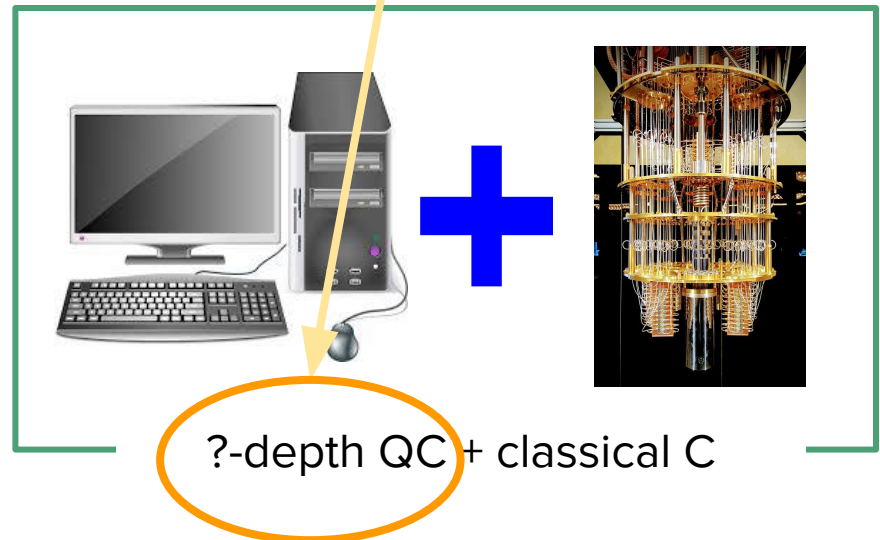
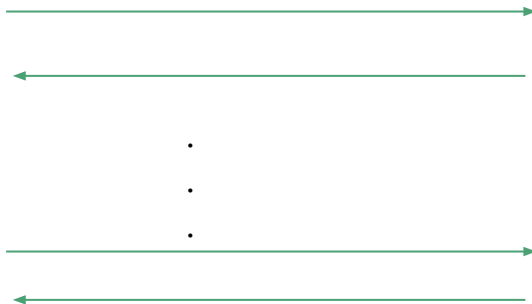
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Goal: A method to check quantum depth of a remote server

Classical Verification of Quantum Depth (CVQD)

A protocol which lets a classical verifier check the quantum depth of a remote server

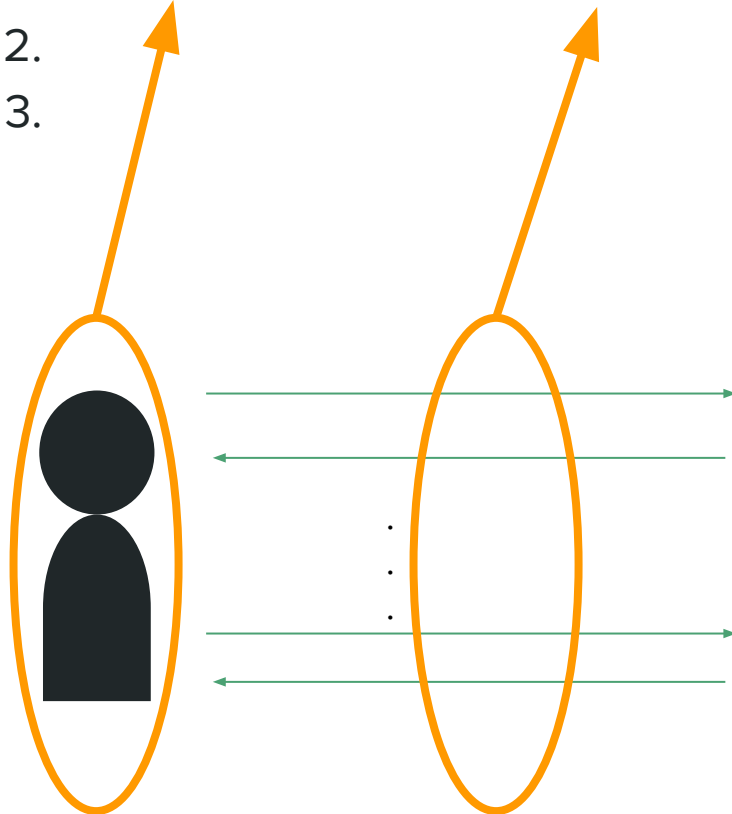


Classical Verification of Quantum Depth (CVQD)

A protocol which lets a classical verifier check the quantum depth of a remote server

Three requirements for CVQD

1. **Classical** verifier and **classical** protocol
- 2.
- 3.



Classical Verification of Quantum Depth (CVQD)

Three requirements for CVQD

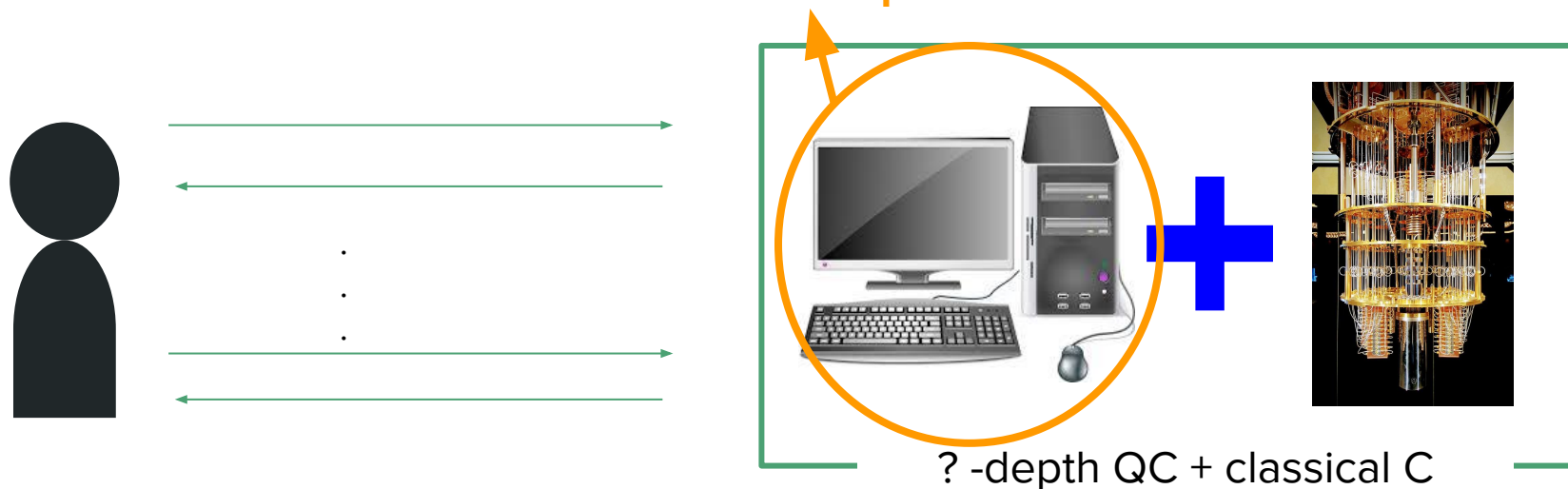
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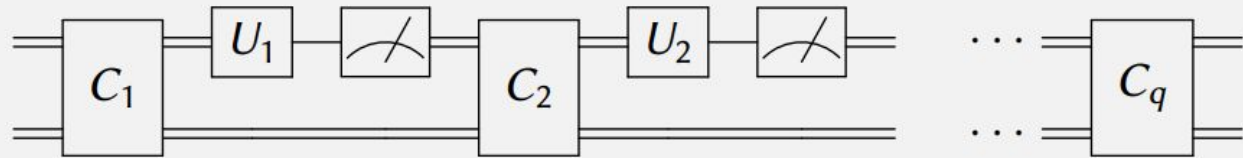
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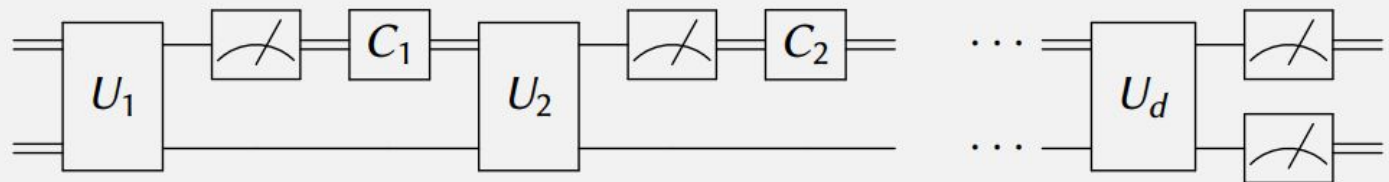


Two Models of Hybrid Quantum-Classical Comp.

d -CQ scheme



d -QC scheme



- d -CQ scheme: classical computer can access a d -depth quantum circuit
- d -QC scheme: d -depth quantum circuits can access classical computer after each layer of gates.

First Attempt for CVQD protocols

First Attempt for CVQD Protocols

Theorem [CCL20, CH22, HG22]: \exists a problem (called d-SSP) such that

- $[(d+3)$ -depth QC + Classical C] solves the problem
- $[d$ -depth QC+Classical C] cannot solve the problem

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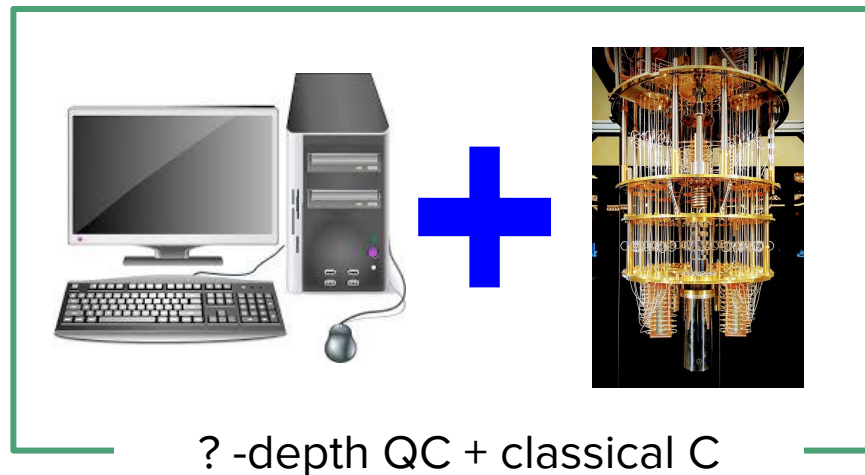
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Can you solve d-SSP?



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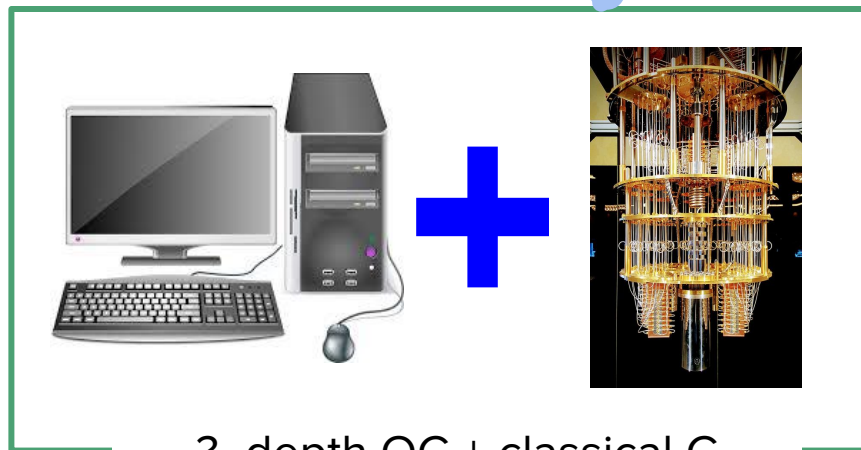
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Oops!

Can you solve d-SSP?



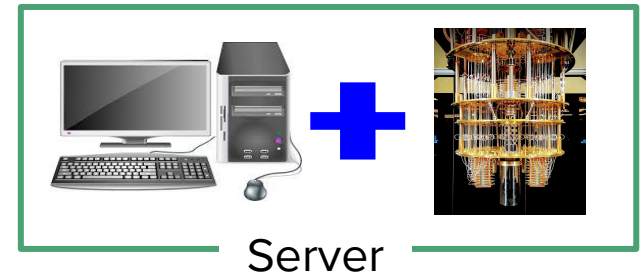
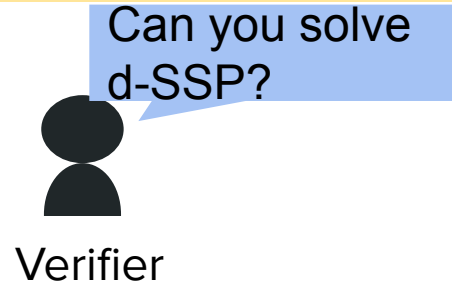
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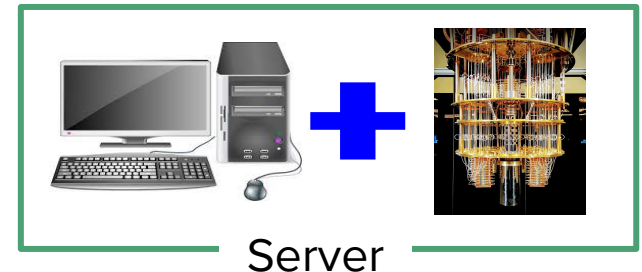
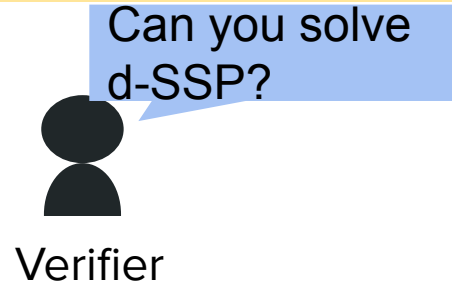


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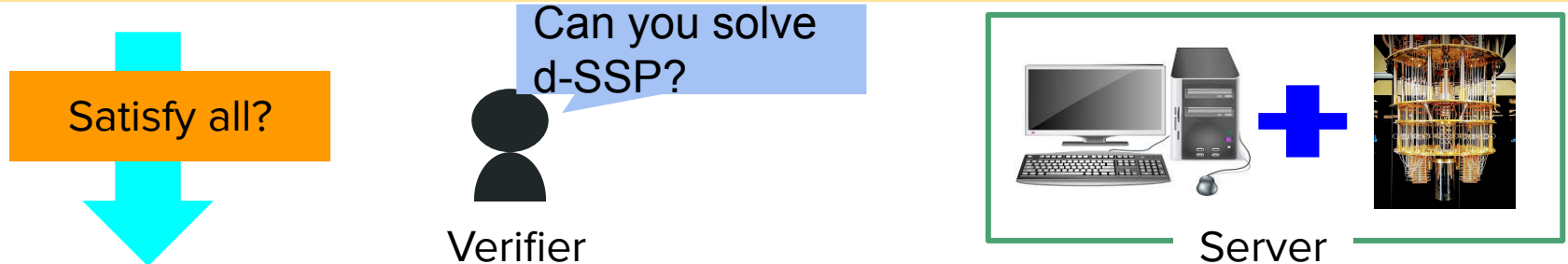
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Satisfy all?

Separate $d+3$ quantum depth from d .

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Satisfy all?

Server without sufficient quantum depth cannot convince the verifier even using classical computer

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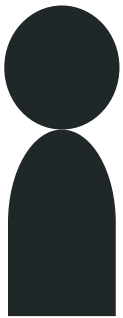
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Algorithm



? -depth QC + classical C

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Queries



Algorithm



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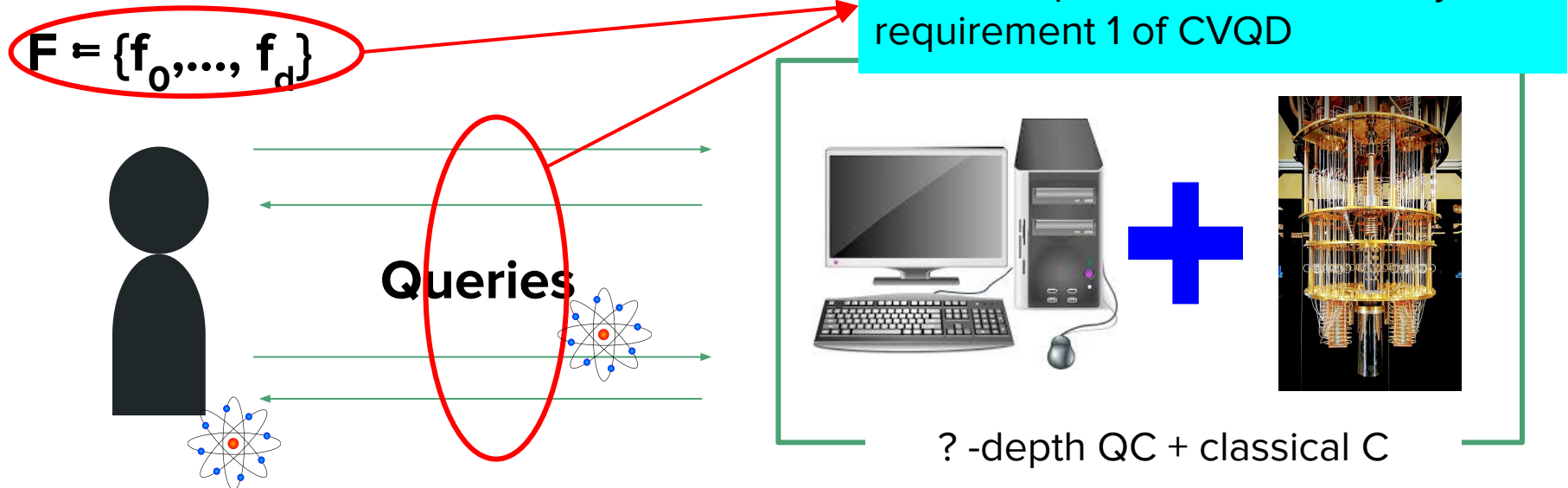
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Our CVQD Protocols

Two CVQD protocols:

Theorem 1 [Chia-Hung 22]: \exists two CVQD protocols where

- **Protocol A** distinguishes **d -depth quantum circuit** from **$d+c^*$** for any d

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Remarks of Protocol A

- The separation **(d v.s. $d+c^*$) is not optimal**
- Honest server **requires ($>c^*$)-depth** to implement the protocol
- Other improvements that we will discuss later

Two CVQD protocols:

Theorem 1 [Chia-Hung 22]: \exists two CVQD protocols where

- **Protocol A** distinguishes **d -depth quantum circuit** from **$d+c^*$** for any d
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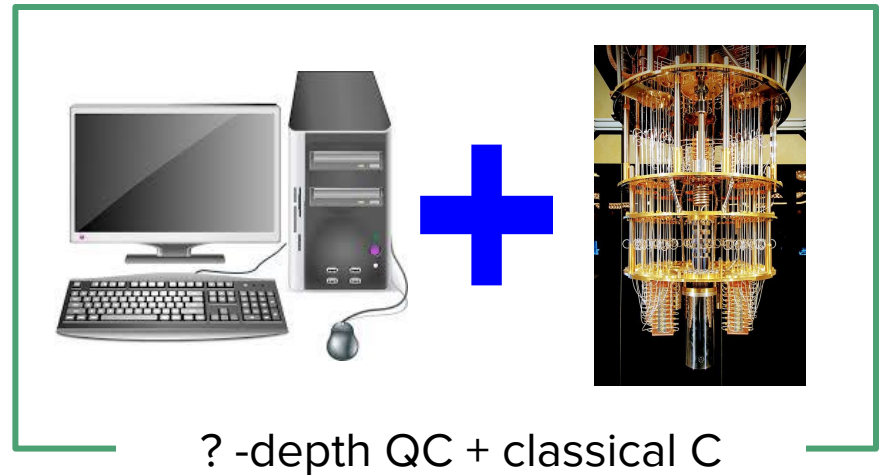
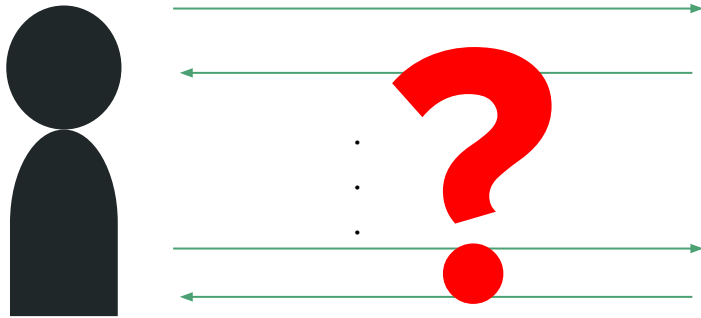
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Remarks of Protocol B

Better separation (d versus $d+3$), but protocol B **requires another dishonest quantum prover** to help.

What is our protocol?



Protocol A

Idea:

Proof of Quantumness Protocol + “Pointer chasing”

Proof of Quantumness (PoQ)

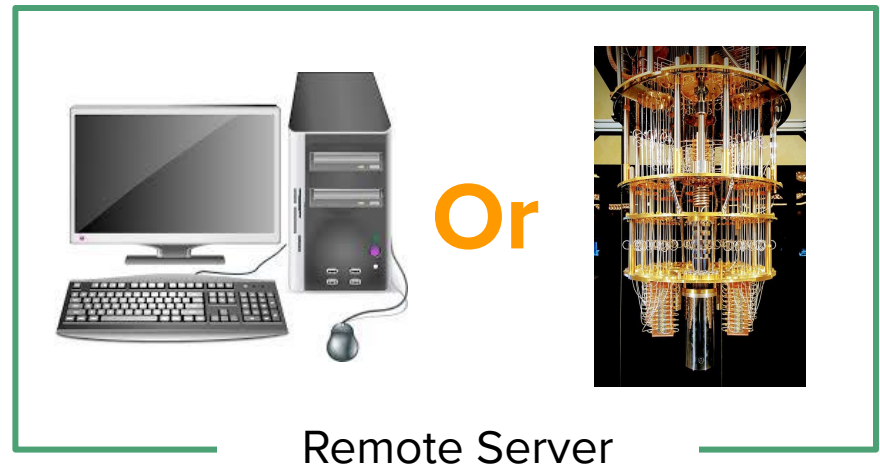
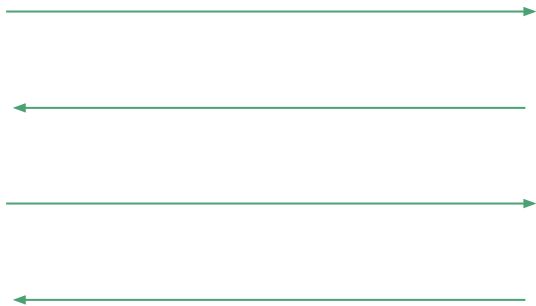
BCM²VV's proof of quantumness protocol (PoQ)

- Goal: Let a classical client to distinguish the following two cases
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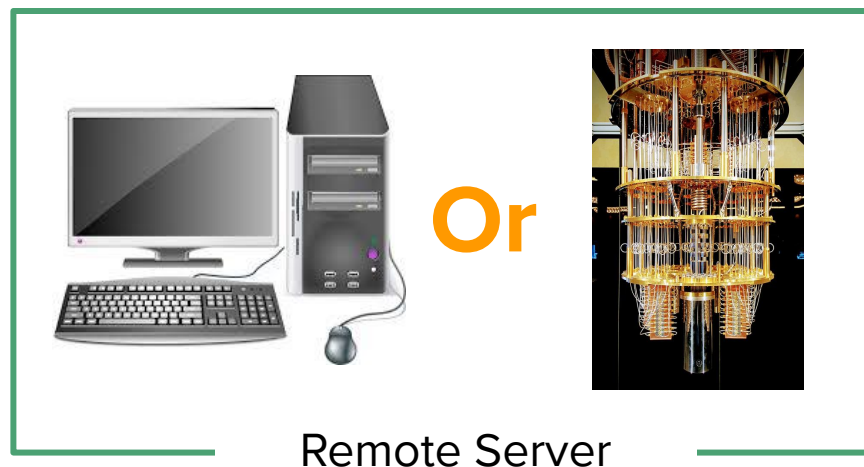
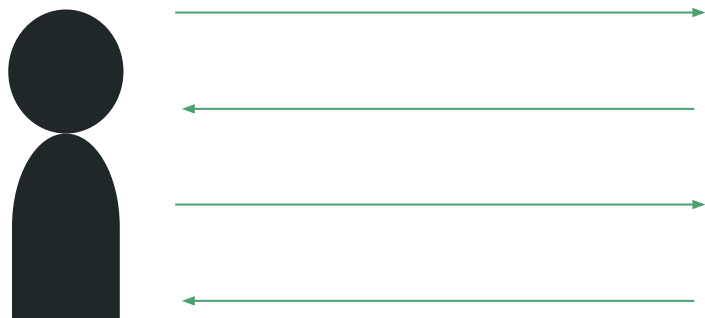
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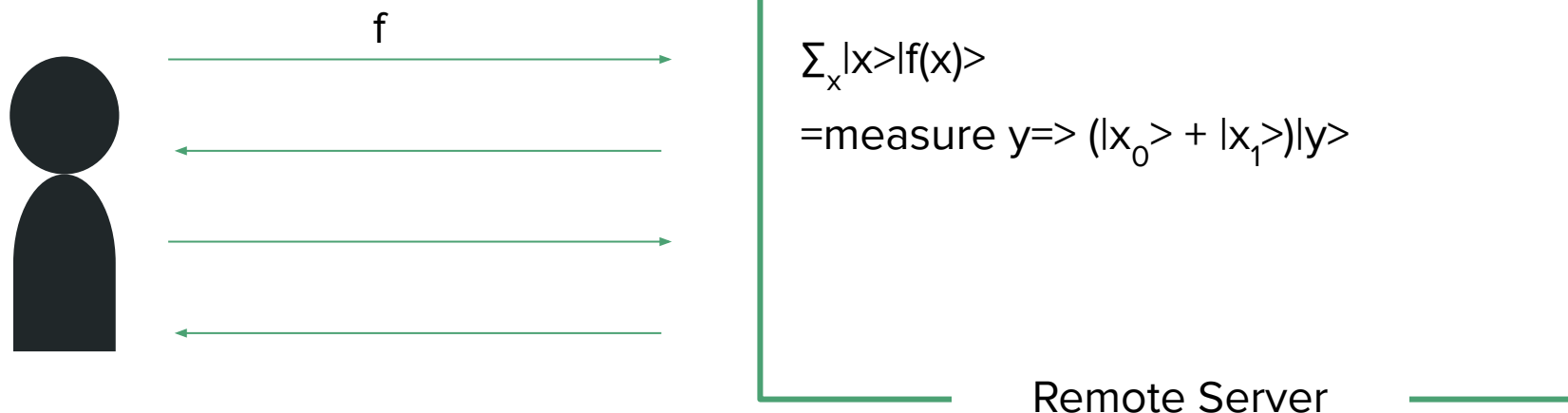


[Brakerski, Christiano, Mahadev, Vazirani, Vidick '18]

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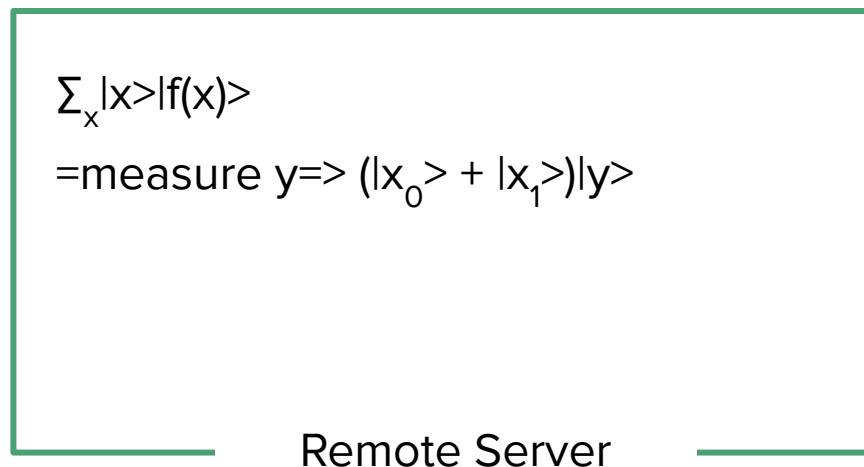
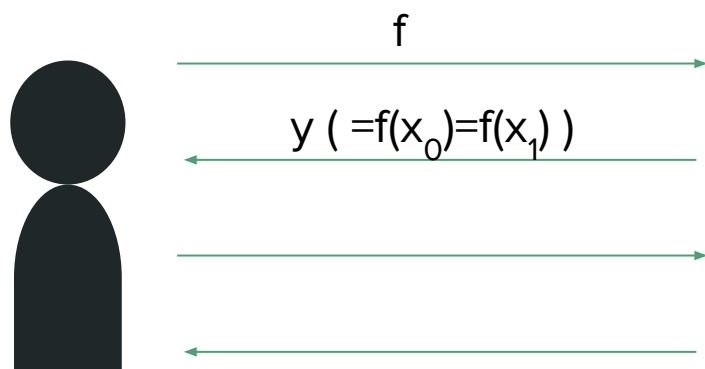
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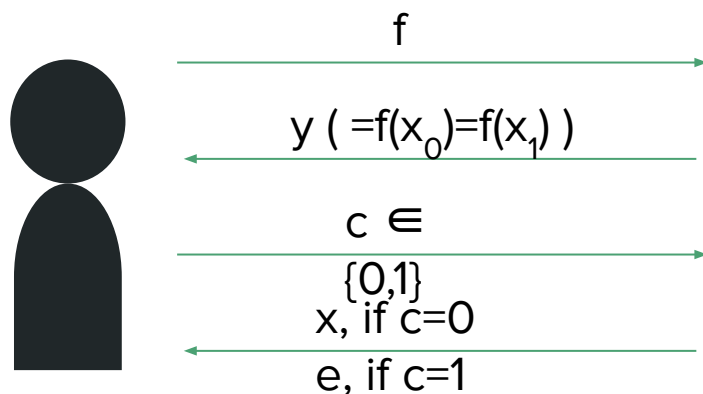
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$$\sum_x |x\rangle |f(x)\rangle$$

=measure $y \Rightarrow (|x_0\rangle + |x_1\rangle) |y\rangle$

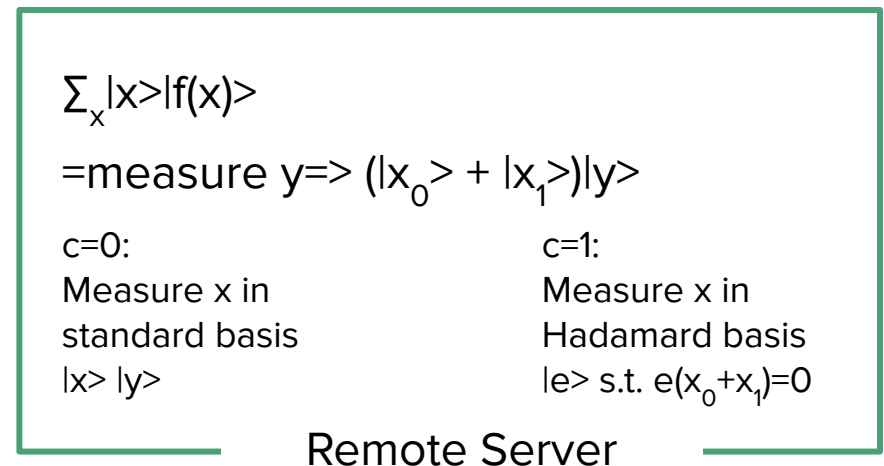
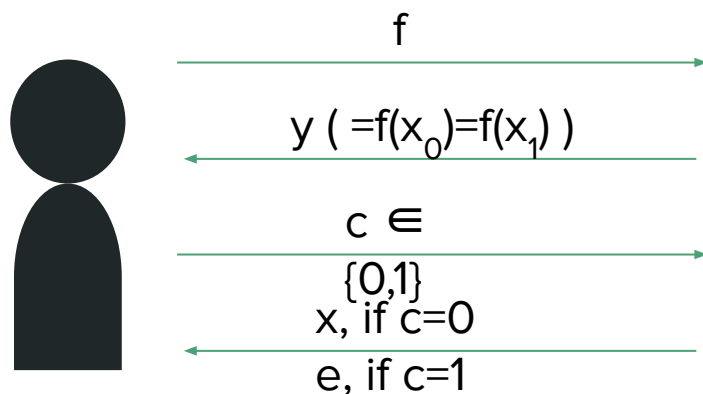
$c=0:$ Measure x in standard basis $ x\rangle y\rangle$	$c=1:$ Measure x in Hadamard basis $ e\rangle$ s.t. $e(x_0 + x_1) = 0$
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Remote Server

- A classical server who can win the game can break AHB by rewinding!
- [Hirahara-Le Gall '21] and [Liu-Gheorghiu '21] showed that f can be evaluated in depth- $O(1)$ (i.e., c^*) \Rightarrow separates depth- $O(1)$ from depth-0 (i.e., classical) device.

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 - b. A remote server is purely classical
- **AHB:** From LWE, \exists distribution over 2-to-1 functions f such that
 - a. No efficient Q. alg. can output (y, x, e) s.t. $f(x) = y$ and $(x_0 + x_1) \cdot e = 0$ w.p. $> 1/2$
 - b. After revealing y , (y, x) or (y, e) can be efficiently computed (by Q. alg.)



PoQ + “Pointer Chasing” = CVQD!

- Goal: Distinguish a depth- $(d + c^*)$ from a depth- d QC scheme

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f_1, f_2, \dots, f_d



$$\sum_x |x\rangle |f_1(x)\rangle, \sum_x |x\rangle |f_2(x)\rangle, \dots, \sum_x |x\rangle |f_d(x)\rangle \\ \Rightarrow (|x_0^1\rangle + |x_1^1\rangle) |y_1\rangle, \dots, (|x_0^d\rangle + |x_1^d\rangle) |y_d\rangle$$

Remote Server

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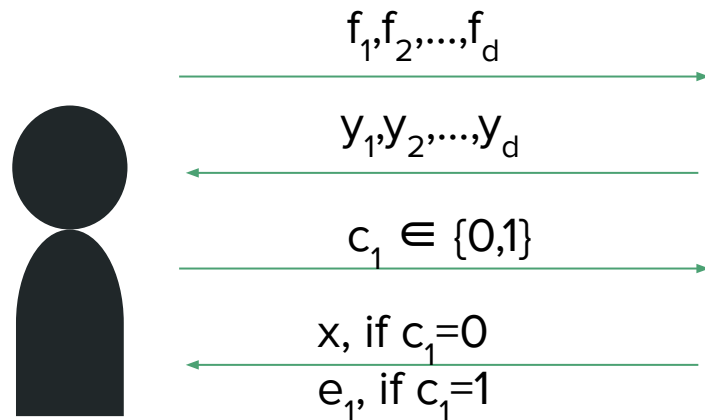
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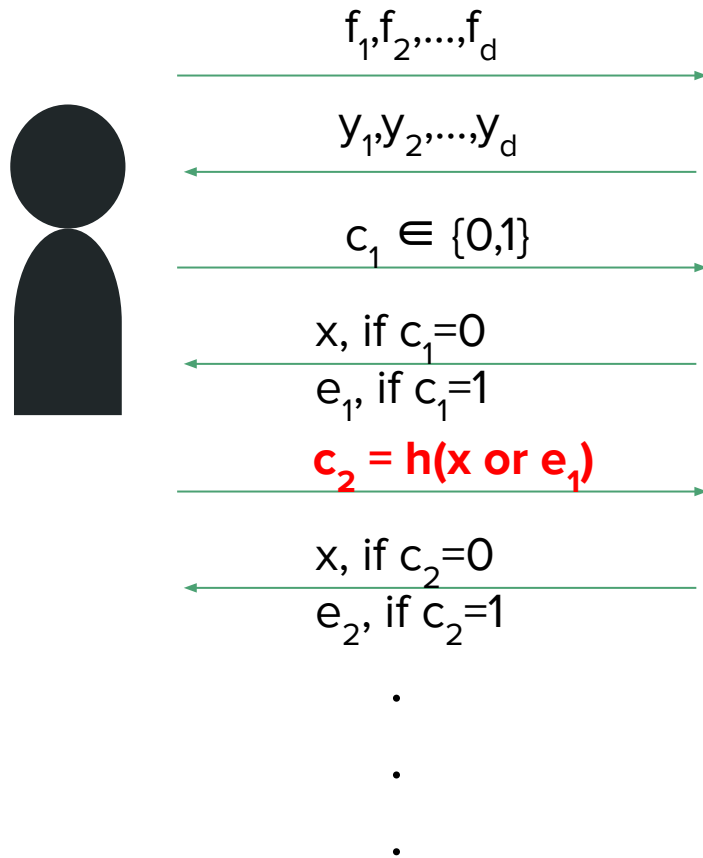
$c_1=0$:
Measure x in
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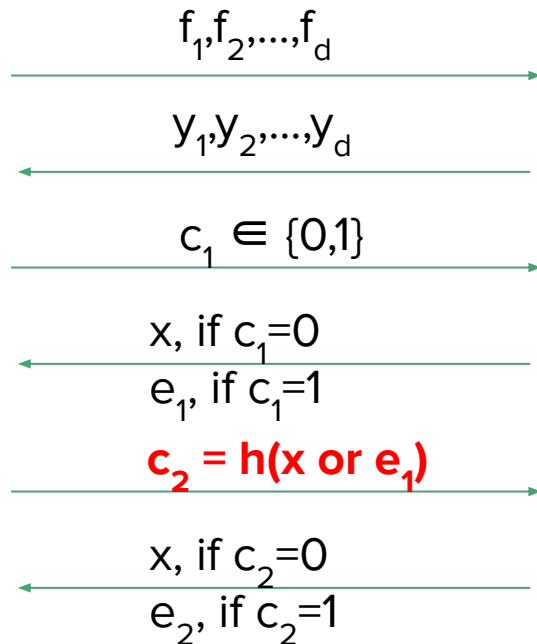
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c_i is determined by the answer of the $(i-1)$ -th round

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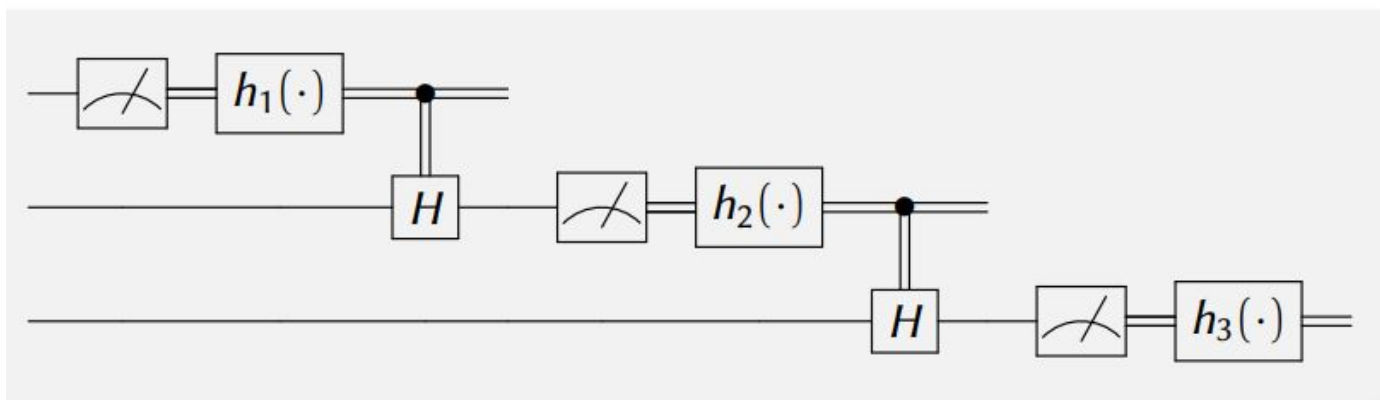
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Remote Server

Completeness and Soundness

Completeness: $\exists (d+c^*)$ -QC that can win the game!

- Need c^* depth to evaluate $\sum_x |x\rangle |f(x)\rangle$
- $(i-1)$ -th answer determines i -th challenge
 - i -th qubit keeps coherence until all previous $i-1$ rounds complete



Soundness: $\forall d'$ -QC with $d' < d$ can't win the game breaks AHB

- Pigeon holes:
 - **1** quantum depth can only be used in **1** round
 - **d'** depth vs **d** rounds $\Rightarrow \exists$ rounds that are executed classically
 \Rightarrow one can use these rounds to break AHB

First CVQD Protocol

Theorem 1 [Chia-Hung 22]: \exists CVQD protocols which distinguishes **d-depth quantum circuit** from **$d+c^*$** for any d

Protocol A

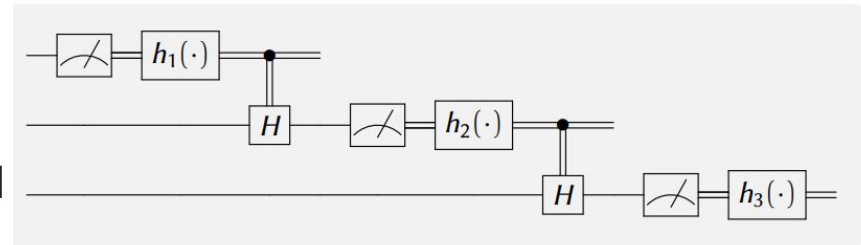
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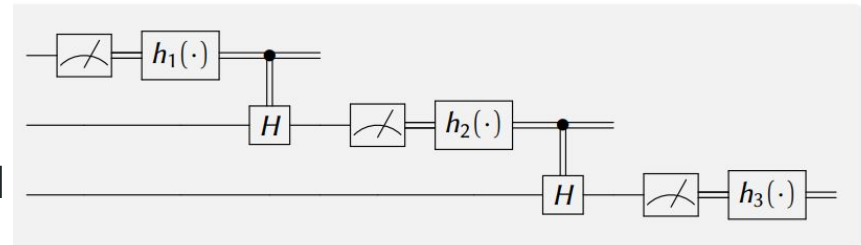
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Question: Can we solve these issues?

Protocol B

Disadvantages of Protocol A

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Protocol B

Protocol B

- The separation $(d \text{ v.s. } d+3)$ is not optimal
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d -depth circuit with classical post processing



Protocol B

Protocol B

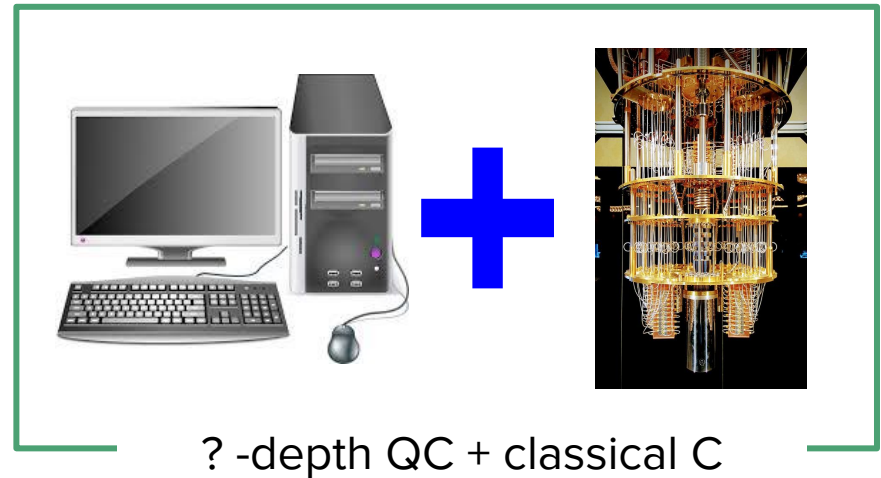
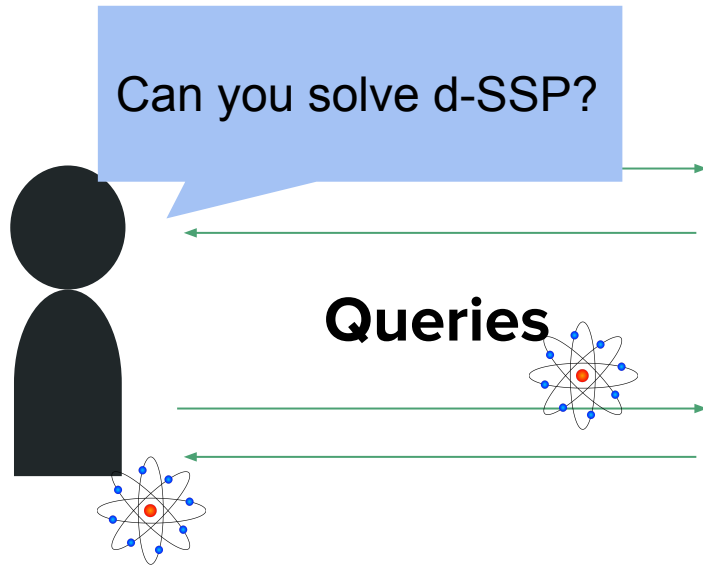
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Idea:

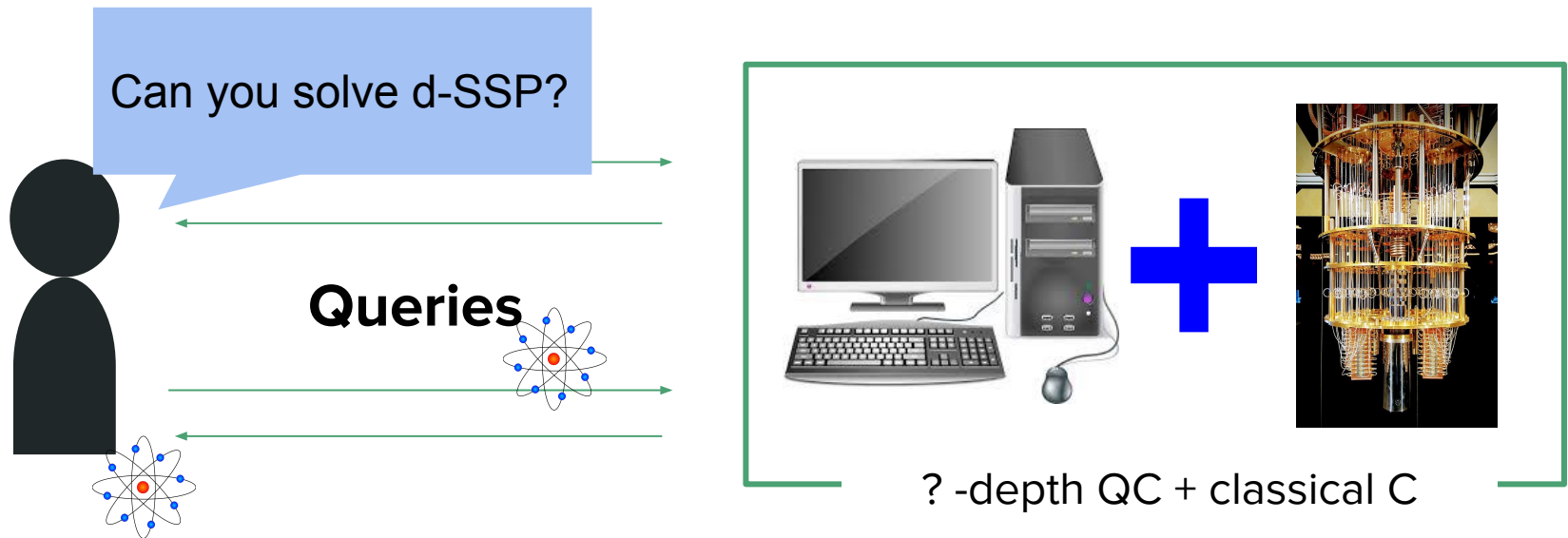
Revisit the first failed attempt (the protocol from d -SSP problem) and dequantize the protocol via **non-local game**

Need to “Dequantize” the Protocol



How to make the verifier and the protocol classical?

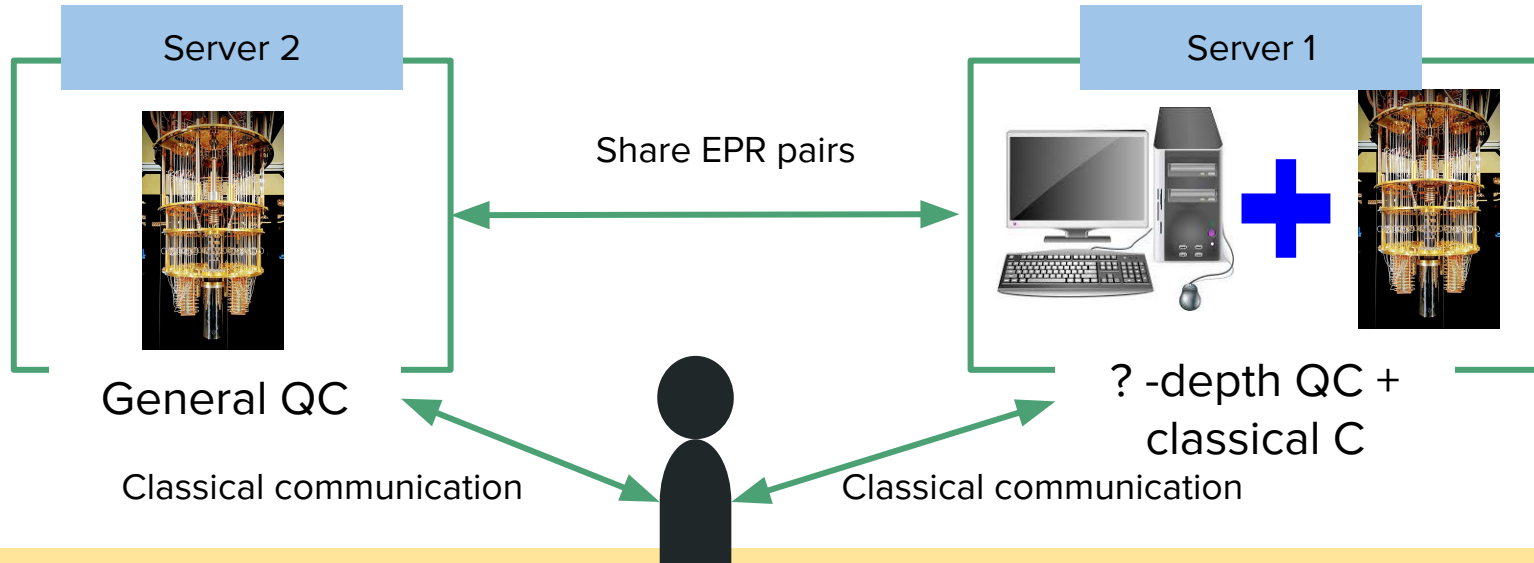
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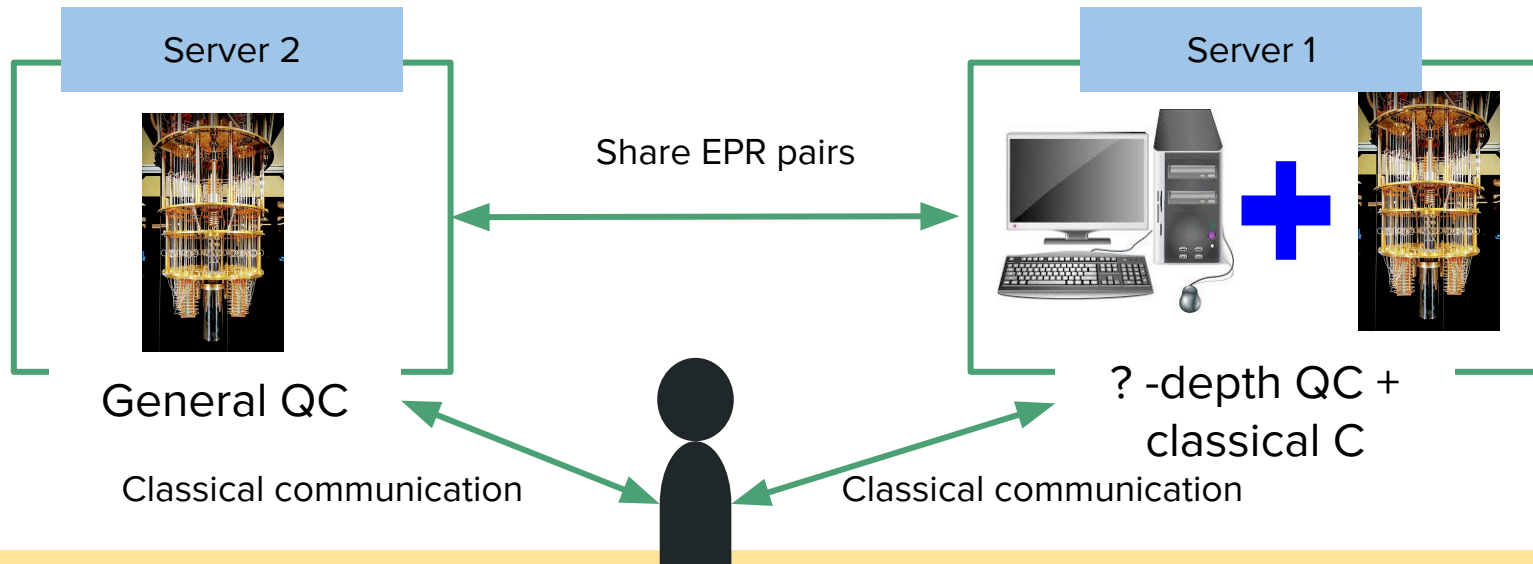
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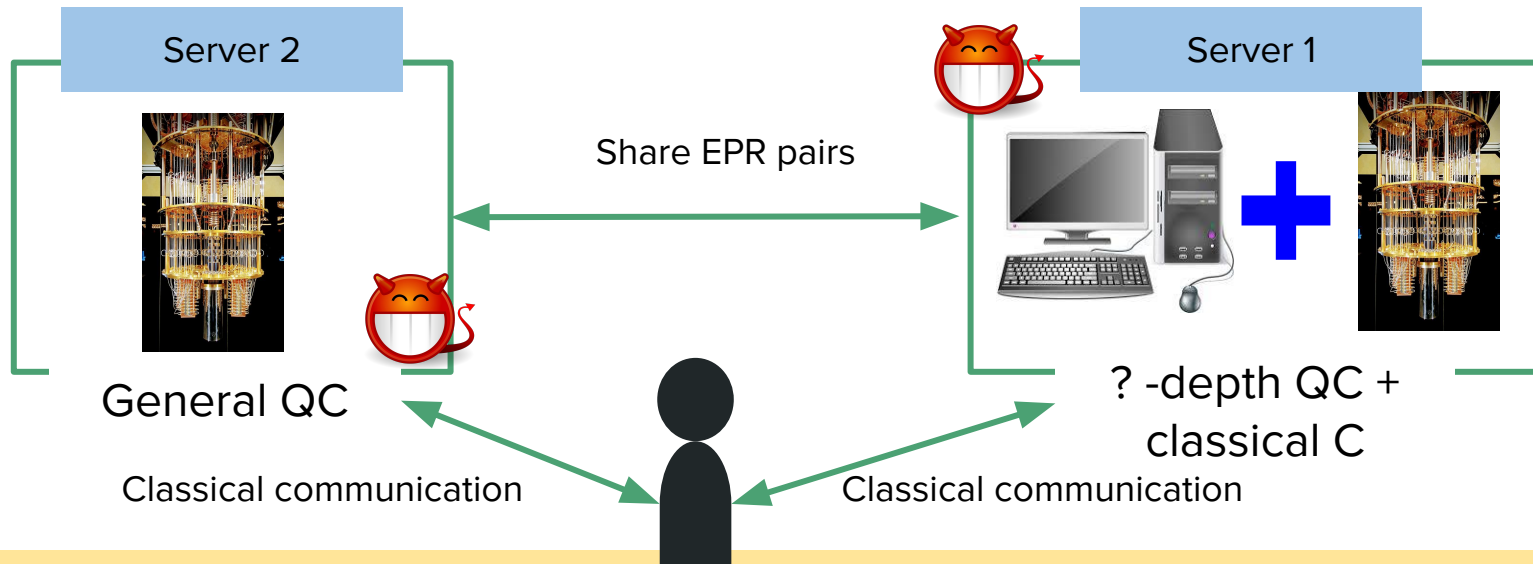


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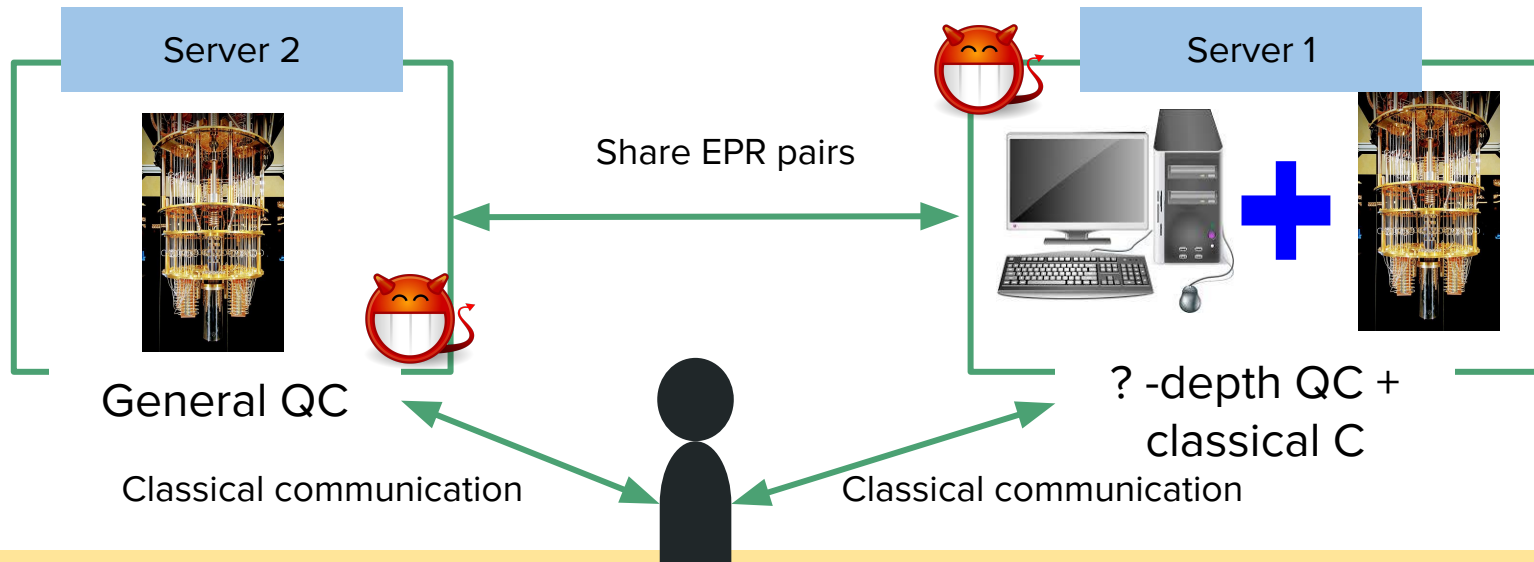


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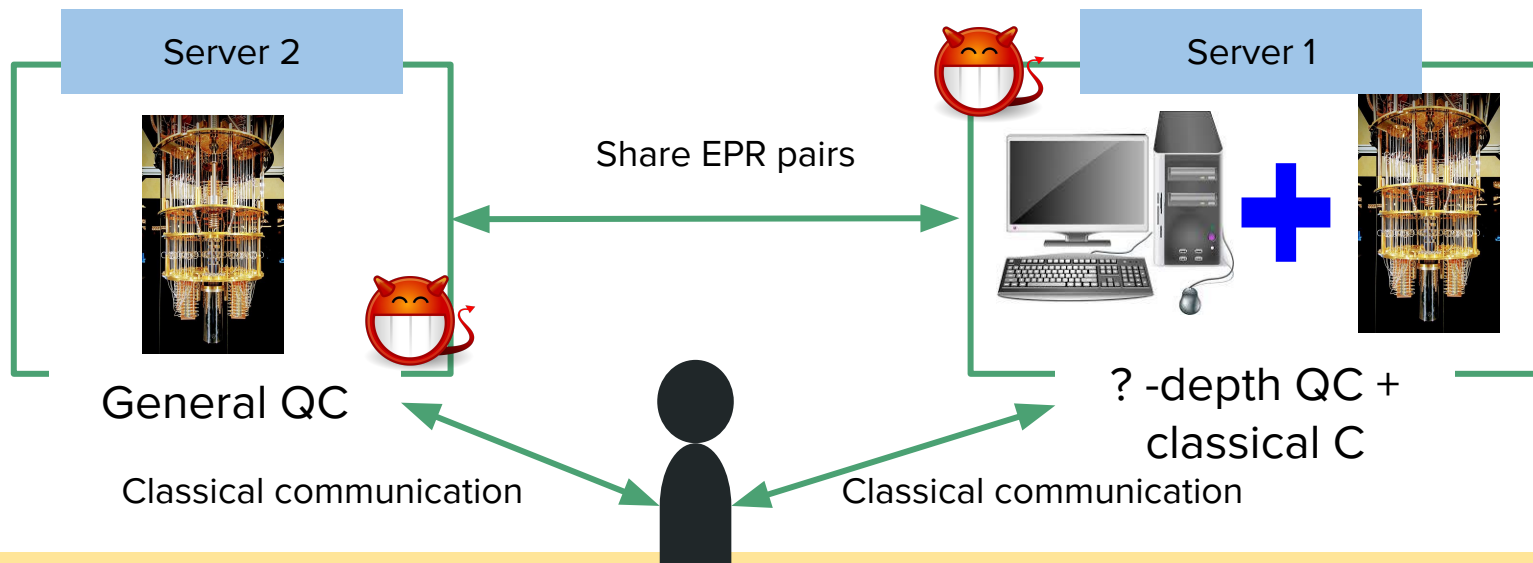


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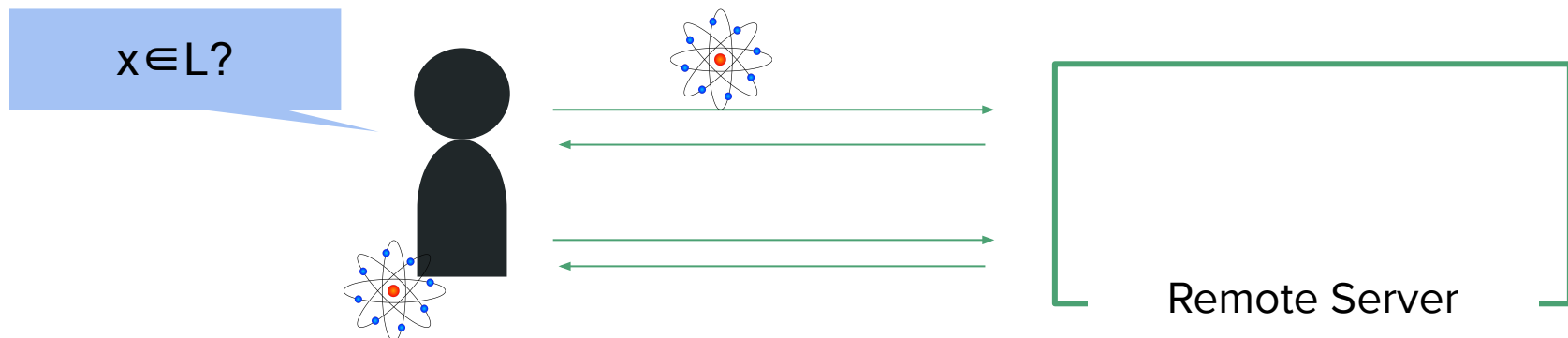
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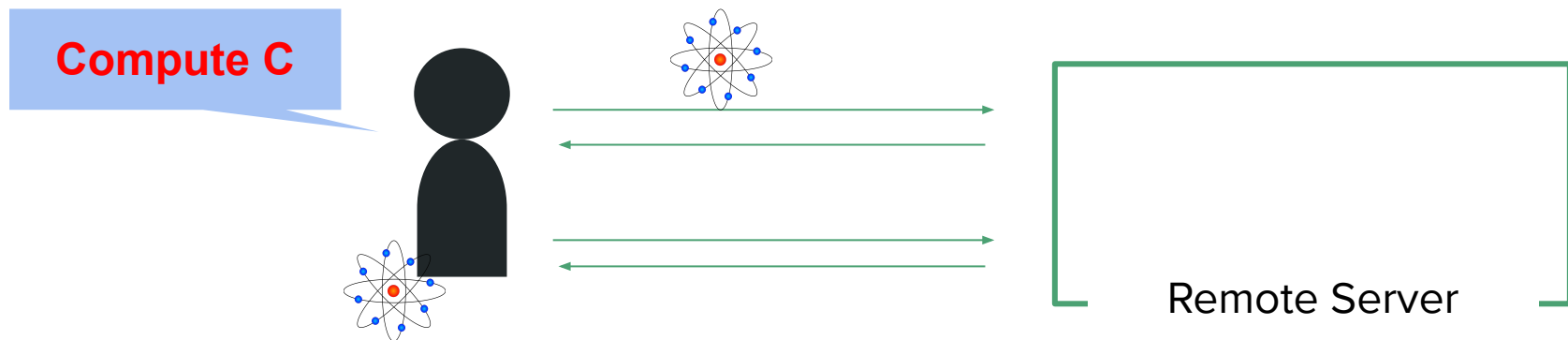
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- The Broadbent protocol for **BQP** [Broadbent '18]
 - Three types: Comp, X-, and Z-test (**indistinguishable from server's view**)
 - Tests for detecting bit-flip (X-test) and phase-flip (Z-test).



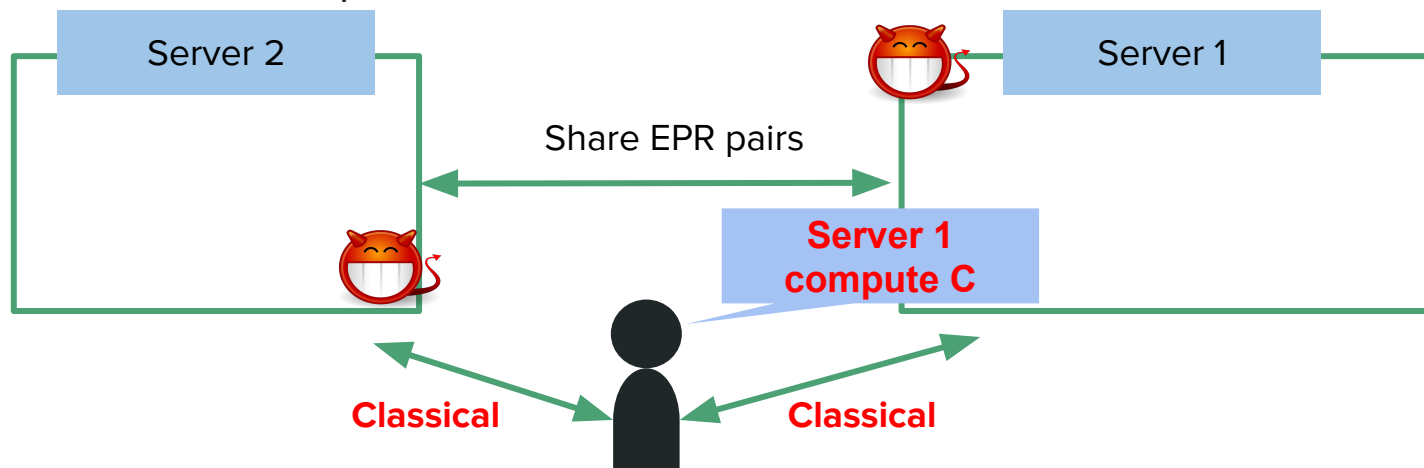
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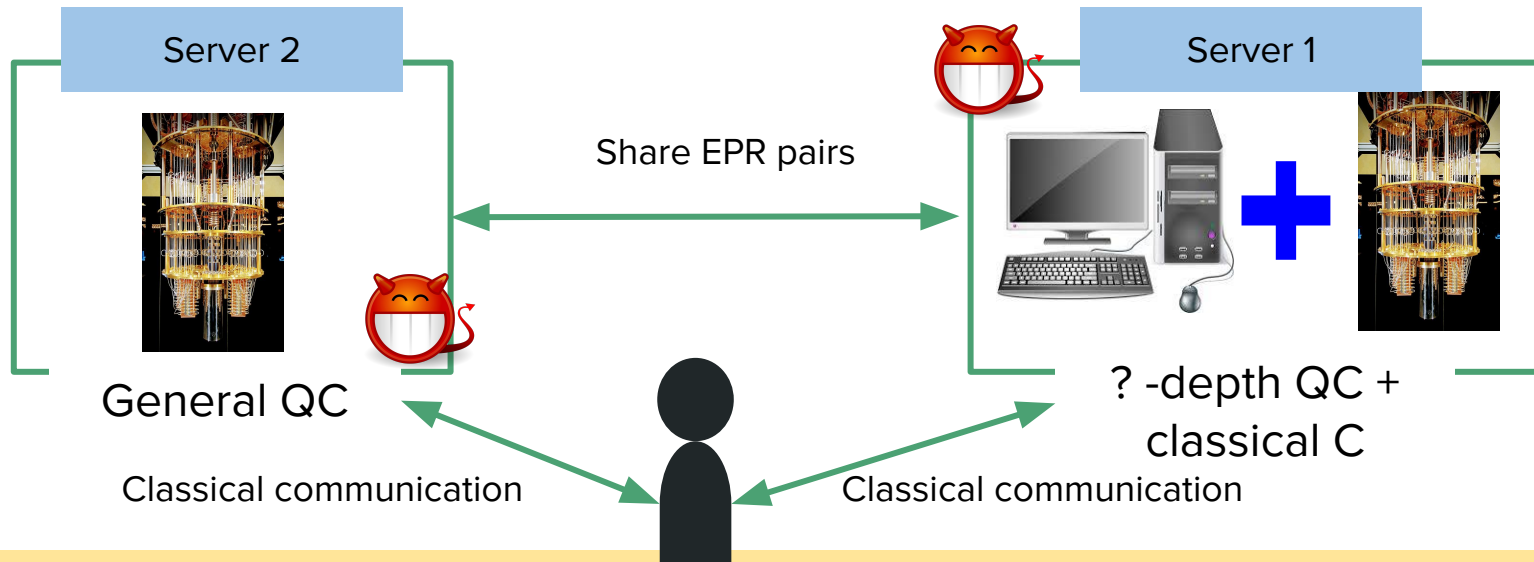


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- Apply the verifier-on-a-leash protocol [Coladangelo, Grilo, Jeery, Vidick '19]:
 - \exists two-player test which certifies Clifford measurements: if the player succeeds w.h.p., then the measurements is correct.



Need to “Dequantize” the Protocol



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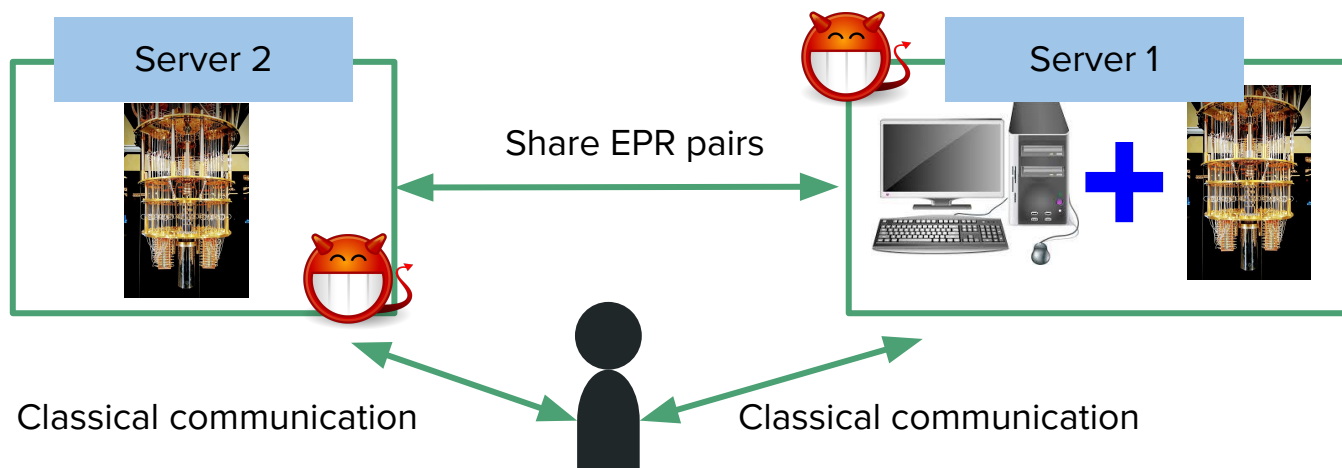
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Recap: Classical Verification of Quantum Depth

Theorem 2 [Chia-Hung 21]: There exists a two-prover classical protocol that satisfies the following:

1. The classical verifier accepts if the server has $\geq \underline{d+3}$ quantum depth
2. The classical verifier rejects if the server has $\leq \underline{d}$ quantum depth

- Unconditionally secure (inefficient) / Nearly optimal separation



Two-Prover CVQD

Theorem 1 [Chia-Hung 22]: \exists two CVQD protocols where

- **Protocol A** distinguishes **d-depth quantum circuit** from **$d+c^*$** for any d
- **Protocol B, a two-prover protocol**, distinguishes **d-depth quantum circuit** from **$d+3$** for any d

Protocol B

1. **Classical** verifier and **classical** protocol
2. Recognize server's quantum depth
 - a. Server's quantum depth $> d+3 \Rightarrow$ Verifier accepts
 - b. Server's quantum depth $\leq d \Rightarrow$ Verifier rejects
3. The server 1 only needs **d-depth circuit + classical post processing**

Unconditionally secure (inefficient) / Nearly optimal separation

Disadvantage of Protocol B

Protocol B **requires another dishonest quantum prover** to help...

Open Questions

Open question 1: Classical verification of quantum **resources**, e.g., quantum memory, quantum volume, etc.

Open question 2: Improved CVQD protocol that can be implemented **on near-term devices**

Open question 3: Separations between hybrid quantum-classical computation (d -CQ and d -QC schemes) **in the plain model**

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Thank you!