

Project Report Omega

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1 Description of the game

2 Choice of programming language

I chose the programming language Julia. Since Mark told us that Python is quite slow, the programming language I'm most comfortable with fell through and I had to choose another one. Julia is quite new (version 1.0 came out on 8th of August 2018) [1] and has a similar syntax like Python but was created for high performance computations. With Julia it is possible to have C-like speeds [2]. Although this is mainly useful for data analysis and machine learning I wanted to try it out for this project. Furthermore, it is possible to import Python and C packages which can be useful if there is no native Julia package released yet.

3 Usage of the program

In order to use the implementation Julia needs to be installed. To do this go to <https://julialang.org/downloads/> and download it. For easier use you can add Julia to your Path variable. Since my implementation uses packages that are not installed natively you need to run Julia and open the so called REPL (Read-eval-print-loop). Type 'using Pkg' (this is the package installer of Julia) followed by the following statements: 'Pkg.add("Printf")', 'Pkg.add("IterTools")' and finally 'Pkg.add("DataStructures")'.

With Julia correctly set up you can go to the directory of my implementation (<https://github.com/HansBambel/Omega>). In order to play a game of Omega you have to type in bash 'julia main.jl <boardsize> <players>'. Boardsize is the dimension of the Omega board, e.g. 5 or 7. Players is a composite of two letters, h (Human), a (AI), r (random) or d (dumber AI), e.g. 'ah' would be AI as white playing against a human player as black. The sequence the players are entered also defines who is white and who is black. A complete call would be for example 'julia main.jl 5 ah'. Note that you will be asked for the parameters if you did not enter any, i.e. you just entered 'julia main.jl'.

When playing the game the current player is displayed, when the current player is a human he is asked for the position to put a stone. The position is a 2-tuple, consisting of the row and the column. After entering the first position the is asked for the next stone coordinates. If a mistake was done in either of these, he will be asked again. At the end of his turn he is able to undo the entered moves and can restart his turn.

Furthermore, it is possible to vary the amount of time the AI has to search. This is done by changing the variable *totalTurnTime* in main.jl. By default, the total time the AI gets is at 15 minutes, resulting in about 87 seconds per turn. In order to have a more timely game time I recommend 2-4 minutes. This reduces the time per turn to about 14 seconds.

4 Details of implementation

I chose to implement Omega on the command-line because Julia has no native graphical user interface library and it seemed to me a bit too much work to learn to program a GUI in Julia using another library additionally to learning Julia in the first place. Using unicode symbols still makes the graphical board representation appealing and useful.

```
##### TURN 9: Best AI PLAYER x #####
Turns left: 55
Possible moves: 111
Time for Turn 9: 22.283114354760002
Depth 1 took 0.010746897s Best Value: -Inf Moves investigated: 111
Depth 2 took 1.230231863s Best Value: 12.779816513761467 Moves investigated: 12321
Depth 3 took 3.166897214s Best Value: 2.9166666666666668 Moves investigated: 22754
Depth 4 took 17.921998838s Best Value: -5.046728971962617 Moves investigated: 176000
x stone set on (1, 6)
Δ stone set on (2, 7)
 1 :      Δ x . . . . . x x
 2 :      Δ x . . . . . Δ Δ
 3 :      . . . . . x . . .
 4 :      . . . . . x . . .
 5 :      . . . . . . . . .
 6 :      . . . . . . . . .
 7 : Δ . . . . Δ Δ x Δ Δ . . x
 8 : . . . . . . . . . . .
 9 :      . . . . . . . . .
10 :      . . . . . . . . .
11 :      . . . . . . . . .
12 :      . . . . . . . . .
13 :      Δ . . . . . x
Best AI needed 45.26279076s of its given 600.0s
Current Score: [2.0, 8.0]
Current Heuristic: [1.0, -7.0]
##### TURN 10: HUMAN PLAYER Δ #####
Where should a stone of player x be put? (Format: row, col or row col)
12 1
x stone set on (12, 1)
 1 :      Δ x . . . . . x x
 2 :      Δ x . . . . . Δ Δ
 3 :      . . . . . x . . .
 4 :      . . . . . x . . .
 5 :      . . . . . . . . .
 6 :      . . . . . . . . .
 7 : Δ . . . . Δ Δ x Δ Δ . . x
 8 : . . . . . . . . . . .
 9 :      . . . . . . . . .
10 :      . . . . . . . . .
11 :      . . . . . . . . .
12 :      x . . . . . . . .
13 :      Δ . . . . . x
Where should a stone of player Δ be put? (Format: row, col or row col)
13 2
```

The grid has a hashing function for the later use of a transpositiontable. The hashing is done using zobrist hashing [4] to be computationally efficient. A random 64-bit Integer is created when the grid is initialized to represent the empty board. Then, for every hexfield three 64-bit Integer values are generated representing that this field can be either free or have a white or a black stone on it. When a stone is set on the grid, the corresponding Integer of the previous occupying stone is removed from the old hashvalue using the XOR operation. Then, the new stone is inserted to the hashcode using the XOR operation with the corresponding Integer of the new stone.

The hexgrid, additionally to having the current boardstate stored, has two lists of groups and group-sizes stored. This enables a fast calculation of the current score. When putting a player stone on the board a new hashcode is calculated and also a union-find algorithm [3] with path-splitting is used to update potential new groups. There is a difference to this when a stone of a player is removed (a field is freed up): additionally to storing the current board state, groups and group-sizes I also store the past groups and group-sizes. This makes it easy to revert to a previous state where the latest player stone was not set. This is done because a de-unification and thus updating all the groups is inefficient and requires more overhead.

When calculating the scores it is only required to multiply all values of the player's group-sizes together and not a scan of the whole board is needed.

5 Implementation of AI

The AI is implemented using a negamax search with iterative deepening and using a transposition table. To reduce the branching factor negamax is looking at a single move at each search depth and not a whole turn, which consists of two moves. This reduces the branching factor significantly: On a 6-by-6 boardgame it reduces the branching factor from 8190 to only 91 on the first move. Furthermore, with this division into two separate moves it allows negamax to prune even more resulting in less subtrees that are searched.

5.1 Transposition table

Using a transposition table enables iterative deepening to work, since without it the previous search is not used at all and only extends the search space. Furthermore, using a transposition table applies move ordering thus making iterative deepening more effective. As previously already mentioned, the transposition table is implemented using Zobrist-hashing.

5.2 Move ordering

Move ordering aims at investigating promising moves first as they may lead to additional pruning, making the tree search more effective. In my implementation, the AI investigates first the move that was saved in the transposition table, then it will look into up to two killer-moves. A killer-move is a move that resulted in a cut-off in a previous search at that search depth.

5.3 Heuristic

The heuristic the AI uses consists of two scores. The first only takes groups of two and three into consideration. Those groups are then checked for their free neighbouring hexagons. Every free hexagon of a group-hex serves as a punishment, encouraging *safe* groups (groups that cannot be extended). All other, bigger groups, are disregarded until a terminal state is reached. The second score is the current board score. The final heuristic score consists of a weighting of these scores. The further the game progresses the more the current score will be taken into account and less the first score.

When a terminal state is reached, the complete score will be calculated, because this score can not be an overestimate since nothing changes any more. This weighting of both scores has the effect that even later in the game the AI tries to form groups of two and three.

5.4 Early Moves

When the board is empty the possible moves are very high, but as a human good opening moves are quite easy to spot. A good idea is to put your own stone in the corner of the board while trying to put the other player's stones in the middle. This results in safer groups of your own and unsafe and easily extendible groups of the opponent. That is the reason that for the first five turns the AI puts their own stones in the corner of the board and groups the opponent's stones in the middle of the board.

5.5 Null move

Using a null-move should improve the search depth even further. It assumes that when passing your turn, you are still in a good spot later on, you can do even better when actually making a move. Since the search depth when doing a null-move is more shallow, it requires less time to return a value. A null-move should not follow a previous null-move, otherwise people would just skip all the time. In order to avoid this the null-move flag is disabled after executing a null-move.

5.6 Multi-cut

Another enhancement which the AI uses is multi-cut. Multi-cut assumes that when there are a certain amount of prunings already at the beginning of a (shallower) search then there will be even more later on. This is a form of forward pruning and may be more aggressive resulting in incorrect or bad prunings. The AI prunes where there are at least three children out of ten where a cut-off occurred.

5.7 Time management

What the AI does for time management is the following: it gets the remaining time and calculates how much time it maximally is allowed to use when every turn receives the same amount of search time. Disregarded are the last three turns, which require in total about 15 seconds to find a terminal state. The AI then uses the calculated time window to do a search. At the end of such a search the elapsed time of the search is subtracted from the remaining time resulting in new search windows for the next search.

6 Problems

6.1 Heuristic

Finding a good heuristic was a key problem. As my old AI-teacher always said: "An admissible heuristics never overestimates." Therefore, using the current

score is not suitable because having a high score now can change quickly in the next move when multiple groups are combined to a single one. In the earlier state of the game the score is not that important, making the heuristic useful for that part of the game. Due to the fact that the heuristic does not take into account bigger groups of the players, which could occur over the mid and late-game it will become less suitable for later stages of the game. In order to compensate for that I decided to also take the current score into account. The current score is multiplied with a discount factor that decreases the further the game progresses.

6.2 Side-effects

For efficiently calculating scores and evaluating the heuristic I use a union-find algorithm [3]. Every time a stone of a player is set on the field the union-find algorithm will unify it with its neighbouring groups if there are any. Unfortunately, separating a group is complicated / computationally inefficient and thus, I save the current groups in a stack. If a separation is necessary, I pop the last entry (where the latest unification occurred) of the stack and use the one before. This way I circumvent a delete function of the union-find algorithm. A problem I had was that the scores were not calculated correctly when the AI-player was searching. Even trying reproducing the exact same moves by hand did not reproduce the same bug. After 10 mind-boggling hours I found out that I needed to create a deepcopy of the latest entry to not have side-effects later on.

7 Experiments

In this section I will compare the different enhancements and how much of a difference they make. The time for the AI was 15 minutes across all experiments on a board size 5. For this I let the fully enhanced AI play two games against the lesser enhanced AI. Once as white and another time as black. In order to see which enhancement was the most effective I let different versions of the AI play against itself. After every run I added another enhancement. Since there is (almost) no randomization (the killermoves get randomly exchanged) we can consider all results deterministic and thus do not need multiple runs.

The first comparison was against an AI that just uses a transposition table. The fully enhanced AI won both games and was able to search at almost every stage of the game 2-ply deeper than the transposition table-AI. Next, I compared against an AI with transposition table and a null-move heuristic. Here the fully enhanced AI won again both games. The difference in search depth here was about 1-ply in favour of the fully enhanced AI. For the next comparison I added a multi-cut heuristic. After two games the AI was able to win both of them. Although in the mid-game the search-depth was similar, in later stages of the game the fully enhanced AI was able to search two and sometimes even three ply deeper than the less enhanced AI. This may be a reason for it to win both of the games. Lastly, adding killer-moves results in both AI being equally potent, thus resulting in a draw, i.e. each should win once. Combining killer-moves with multi-cut resulted in the full-AI winning only one of the games. This means that these two enhancements have a good effect on the tree search.

References

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