

Floyd-Warshall Algorithm

The Floyd-Warshall algorithm is a method for finding the shortest paths between all pairs of vertices in a weighted graph. It's a type of dynamic programming algorithm that systematically finds the shortest paths by checking every possible intermediate vertex.

→ works on both directed and undirected graph.

Why Use it:-

The Floyd-Warshall algorithm is useful when you need to find the shortest path between every single pair of nodes in a graph.



① In network routing, it can be used to find the most efficient path for data to travel between any two routers.

② It's also a great tool to detecting negative cycles in a graph.

How It Works

① Initialization- Start with a distance matrix that represents the graph. If there's a

direct edge between two vertices, use its weight.

- * If not, use infinity to show there's no direct path. The distance from a vertex to itself is always 0.

(2) Iterative Improvements- The algorithm uses three nested loops to go through every possible path. The outer loop, controlled by a variable k , represents an intermediate vertex.

For each k , the algorithm checks every pair of vertices (i, j) to see-

- * If going from i to j via k is shorter than the currently known path.

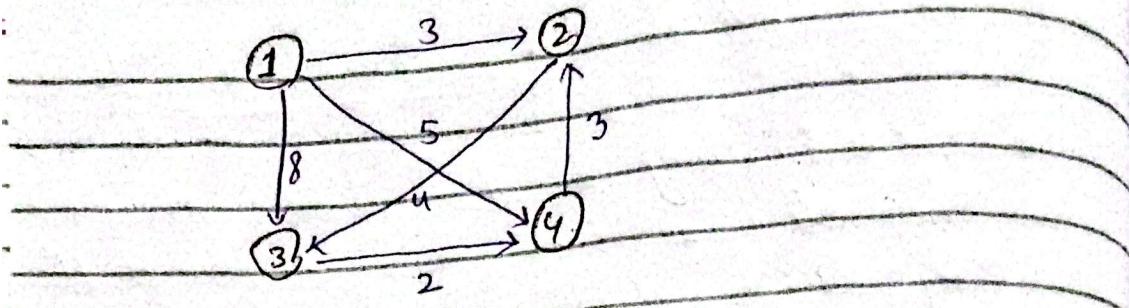
(3) Updated Rule- The core of the algorithm is a simple formula:

$$D[i][j] = \min(D[i][j], (D[i][j] + D[k][j])]$$

This means the new shortest distance from i to j is either the existing distance or the distance from i to k plus the distance from k to j , whichever is smaller.

(4) Final Result- After the loops are complete, the matrix will contain the shortest path distances between all pairs of vertices.

Example:-



Step 1:- Distance Matrix.

$$d^0 = \begin{bmatrix} 0 & 3 & 8 & 5 \\ \infty & 0 & 4 & \infty \\ \infty & \infty & 0 & 2 \\ \infty & 3 & \infty & 0 \end{bmatrix}$$

Step 2:- Iteration with intermediate vertex

$k=1$ (vertex 1).

$$d^{(1)} = \begin{bmatrix} 0 & 3 & 8 & 5 \\ \infty & 0 & 4 & \infty \\ \infty & \infty & 0 & 2 \\ \infty & 3 & \infty & 0 \end{bmatrix}$$

No any value update

→ diagonal remains same and row one and column 1 also same because 1-anynumber and anynumber to 1. distance remain same in this iteration.

Step 3:- Iteration, $k=2$, using (vertex 2)

$$d^{(2)} = \begin{bmatrix} 0 & 3 & 7 & 5 \\ \infty & 0 & 4 & \infty \\ \infty & \infty & 0 & 2 \\ \infty & 3 & 7 & 0 \end{bmatrix}$$

$0 \rightarrow 1 \rightarrow 2 \rightarrow 3 = 7$
 $7 < 8$ update
 for $4 \rightarrow 3$
 $0 \rightarrow 4 \rightarrow 2 \rightarrow 3 = 7$
 $7 < \infty$ update

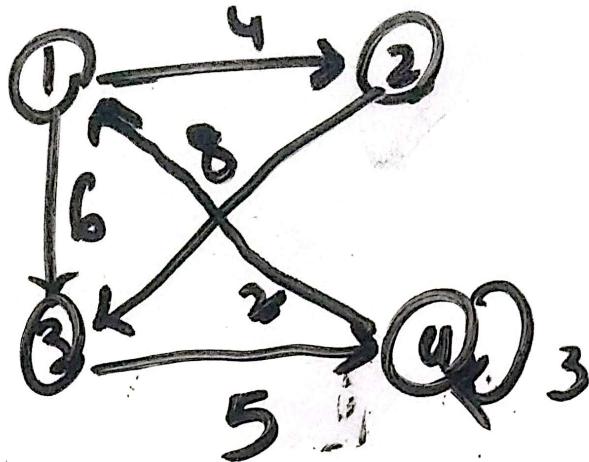
Step 4: Iteration, $k=3$, using (vertex 3)

	1	2	3	4	
$d^{(3)}$	1	0 3 7 5	① for $2 \rightarrow 4$		
	2	∞ 0 4 6	$2 \rightarrow 3 \rightarrow 4 = 6$		
	3	∞ ∞ 0 2	$6 < \infty$ update		
	4	∞ 3 7 0			

Step 5: Iteration, $k=4$, using (vertex 4)

	1	2	3	4	
$d^{(4)}$	1	0 3 7 5	① for $3 \rightarrow 2$		
	2	∞ 0 4 6	$3 \rightarrow 4 \rightarrow 2 = 5$		
	3	∞ 5 0 2	$5 < \infty$ update		
	4	∞ 3 7 0			

* $d^{(4)}$ contains the shortest path distances between all pairs of vertices.



graph having
loop.

$$d^{(0)} = \begin{bmatrix} 0 & 4 & 6 & \infty \\ \infty & 0 & 2 & \infty \\ \infty & \infty & 0 & 5 \\ 8 & \infty & \infty & 0 \end{bmatrix}$$

$4 \rightarrow 4 = 3$
first shortest
Path vertex to
itself is 0

$$k=1$$

$$d^{(1)} = \begin{bmatrix} 0 & 4 & 6 & \infty \\ \infty & 0 & 2 & \infty \\ \infty & \infty & 0 & 5 \\ 8 & 12 & 14 & 0 \end{bmatrix}$$

for $4 \rightarrow 2$
 $4 \rightarrow 1 \rightarrow 2 = 12$
 $12 < \infty$ update
for $4 \rightarrow 3$
 $4 \rightarrow 1 \rightarrow 3 = 14$

$$d^{(2)} = \begin{bmatrix} 0 & 4 & 6 & \infty \\ \infty & 0 & 2 & \infty \\ \infty & \infty & 0 & 5 \\ 8 & 12 & 14 & 0 \end{bmatrix}$$

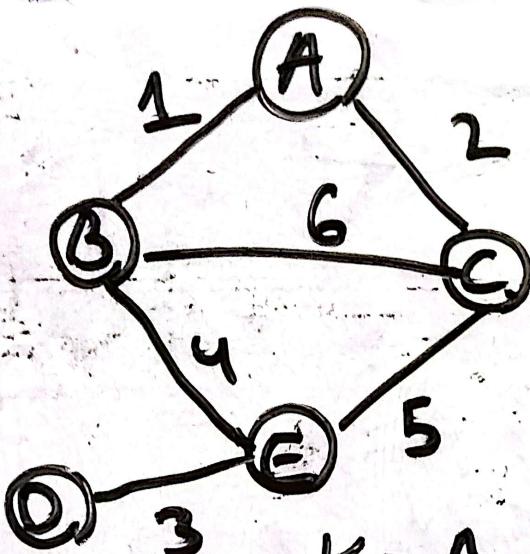
$14 < \infty$ update
No update.

$$K = 3$$

$d^{(3)}$	1	2	3	4	① for $1 \rightarrow 4$
$=$	1	0	4	6	② for $1 \rightarrow 3 \rightarrow 4 = 11$
	2	∞	0	2	$11 < \infty$ update
	3	∞	∞	0	③ for $2 \rightarrow 4$
	4	8	12	14	$2 \rightarrow 3 \rightarrow 4 = 7$
				0	$7 < \infty$ update

$d^{(4)}$	1	2	3	4	① for $2 \rightarrow 1$
	0	4	6	11	② for $2 \rightarrow 3 \rightarrow 4 \rightarrow 1 = 15$
	2	15	0	2	$15 < \infty$
	3	13	17	0	③ for $3 \rightarrow 1$
	4	8	12	14	$3 \rightarrow 4 \rightarrow 1 = 13$
				0	$13 < \infty$

③ for $3 \rightarrow 2$
 $3 \rightarrow 4 \rightarrow 1 \rightarrow 2 = 17 < \infty$ update



	A	B	C	D	E
A	0	1	2	∞	∞
B	1	0	6	4	∞
C	2	6	0	≈ 5	∞
D	∞	∞	∞	0	3
E	∞	4	5	3	0

$K = A$

	A	B	C	D	E
A	0	1	2	∞	∞
B	1	0	3	∞	4
C	2	3	0	∞	5
D	∞	∞	∞	0	3
E	∞	4	5	3	0

① for $B \rightarrow C$
 $B \rightarrow A \rightarrow C = 3 + 6 = 9$

② for $C \rightarrow B$
 $C \rightarrow A \rightarrow B = 3 + 6 = 9$

$K = B$

	A	B	C	D	E
A	0	1	2	∞	5
B	1	0	3	∞	4
C	2	3	0	∞	5
D	∞	∞	∞	0	3
E	5	4	5	3	0

① for $A \rightarrow E$
 $A \rightarrow B \rightarrow E = 5 + 4 = 9$

② for $E \rightarrow A$
 $E \rightarrow B \rightarrow A = 5$

$5 \leq 9$

	A	B	C	D	E
A	0	1	2	∞	5
B	1	0	3	∞	4
C	2	3	0	∞	5
D	∞	∞	∞	0	3
E	5	4	5	3	0

$$K = D$$

A B C D E

	A	B	C	D	E
A	0	1	2	∞	5
B	2	0	3	∞	4
C	∞	3	0	∞	4
D	∞	∞	∞	0	3
E	5	4	4	3	0

$$\overline{K = E}$$

D E OA \rightarrow D

A \rightarrow B \rightarrow E \rightarrow D
 $= 8 < \infty$

	A	B	C	D	E
A	0	1	2	8	5
B	1	0	3	7	4
C	2	3	0	8	4
D	8	7	8	0	3
E	5	4	4	3	0

② B \rightarrow D
B \rightarrow E \rightarrow D = 7
 $7 < \infty$

③ C \rightarrow D
C \rightarrow E \rightarrow D = 8
 $8 < \infty$