Project Proposal

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1 Project Description

Python LP Solver: This project will focus on building an LP solver in **python** from scratch. The project will use **numpy** as the linear algebra package behind the solver, but everything else will be implemented from scratch.

2 Methodology

For me to complete this project, I think this is the best way to approach the problem.

- 1. First, create a simplex_phase_2 method which will take a feasible initial \mathbf{x}_0 as well as $\mathbf{A}, \mathbf{b}, \mathbf{c}$ and return the optimal solution to the problem x^* . Some different variations on the entering variable algorithm are listed below. I will also investigate how these algorithms perform on different problem sizes.
 - Steepest descent picking the value with the greatest negative value to enter the basis.
 - Bland's Rule picking the variable with the first negative value as the entering variable.
 - Secretary's rule I want to try using the Secretary's rule, where you do the first $\frac{1}{e}$ proportion of variables, and pick the one with first value greater than that. I expect this won't work super well but I read that it is supposed to be the most efficient way to find the optimal sequential choice.
- 2. Second, create a simplex_phase_1 method which will take in any of the parameters $\mathbf{A}, \mathbf{b}, \mathbf{c}$ and return a feasible start (or an output stating that the problem is infeasible). I expect that this phase 1 simplex method will create the arbitrary variables \mathbf{h} and then call the simplex phase 2 to solve it. It will identify infeasibility as a solution to the auxiliary problem that is not $\mathbf{1}^{\top}\mathbf{h} = 0$.
- 3. Third, and likely the most complex part of this project, I will implement a python module called lp_reductions.py that will take any LP and turn it into standard form. In order to accomplish this, I plan to do the following:
 - (a) Implement a class of Variable and Expression. Variables will track the variables of a problem. It will probably have some methods like intermediate, non-negative, etc that will be helpful for the below things. Expression will track the equalities and objective function of a problem. Each variable will be assumed to be a vector \mathbb{R}^n and each expression be an affine matrix inequality or equality.

- (b) Accept an arbitrary number of $\mathbf{A}_i \mathbf{x}_i = \mathbf{b}_i$, $\mathbf{x}_i \succeq 0$ equations. Implement a condense_standard_forms function that will take the arbitrary number of affine matrix equalities and concatenate into one $\mathbf{A}\mathbf{x} = \mathbf{b}$, $\mathbf{x} \succeq 0$ problem.
- (c) Implement a function lower_ineq_to_eq. This function should take in the expression $\mathbf{A}\mathbf{x} \leq \mathbf{b}$ and add slack variables to turn it into $\mathbf{A}_s\mathbf{x}_s = \mathbf{b}_s$
- (d) Implement a function greater_ineq_to_lower_ineq. This function will turn $\mathbf{A}\mathbf{x} \succeq \mathbf{b}$ into $-\mathbf{A}\mathbf{x} \preceq -\mathbf{b}$
- (e) Implement a function convert_objective_to_standard_form. This function will take a maximization problem max $\mathbf{c}^{\top}\mathbf{x}$ and convert it into a minimization problem min $-\mathbf{c}^{\top}\mathbf{x}$, which is the standard form for LP solvers.
- (f) Implement a function combine_multivar_equality that will combine any linear combinations of matrix vector multiplications. $\mathbf{A}_0\mathbf{x}_0 + \mathbf{A}_1\mathbf{x}_1 + \dots \mathbf{A}_n\mathbf{x}_n = \mathbf{A}_{tot}\mathbf{x}_{tot}$ where

$$\mathbf{A}_{tot} = egin{bmatrix} \mathbf{A}_0 & \mathbf{A}_1 & \dots & \mathbf{A}_n \end{bmatrix}, \mathbf{x}_{tot} = egin{bmatrix} \mathbf{x}_0 \\ \mathbf{x}_1 \\ \vdots \\ \mathbf{x}_n \end{bmatrix}$$

- (g) Implement a function bound_all_vars which will accept a $\mathbf{A}\mathbf{x} = \mathbf{b}$ and if it is not already bounded by non-negativity, it will split it into x^+, x^- and make them non-negative.
- 4. Finally, put this altogether by testing it with some available toy and real-world LP problems with accessible data.
- 5. If I have time, it would be nice to make it so that variables can appear on any side of the equation and still work. Otherwise I think it is still sufficient to have to enter variables on the left side of the equation.

3 Example Logic

To test how this logic will work, I will set up a problem below and show the outline of the future algorithm I will use to solve it.

Consider a problem with variables $\mathbf{w} \in \mathbb{R}^{20}$, $\mathbf{y} \in \mathbb{R}^{10}$, $z \in \mathbb{R}$ and the parameters $\mathbf{A}_0 \in \mathbb{R}^{5 \times 20}$, $\mathbf{A}_1 \in \mathbb{R}^{5 \times 10}$, $\mathbf{b}_0 \in \mathbb{R}^5$, $\mathbf{b}_1 \in \mathbb{R}^5$, $\mathbf{c}_0 \in \mathbb{R}^{20}$, $\mathbf{c}_1 \in \mathbb{R}^{10}$, $c_2 \in \mathbb{R}$

maximize
$$\mathbf{c}_0^{\mathsf{T}} \mathbf{w} - \mathbf{c}_1^{\mathsf{T}} \mathbf{y} + c_2 z$$
 (1)

subject to
$$\mathbf{A}_0 \mathbf{w} - \mathbf{A}_1 \mathbf{y} = \mathbf{b}_0$$
 (2)

$$\mathbf{A}_0 \mathbf{w} \le \mathbf{b}_0 \tag{3}$$

$$\mathbf{A}_1 \mathbf{y} \succeq \mathbf{b}_1 \tag{4}$$

$$z \le 2 \tag{5}$$

$$\mathbf{w} \succeq 0 \tag{6}$$

We can handle this line by line.

Objective function: call convert_objective_to_standard_form

maximize
$$\mathbf{c}_0^{\mathsf{T}} \mathbf{w} - \mathbf{c}_1^{\mathsf{T}} \mathbf{y} + c_2 z \rightarrow \text{minimize} - \mathbf{c}_0^{\mathsf{T}} \mathbf{w} + \mathbf{c}_1^{\mathsf{T}} \mathbf{y} - c_2 z$$
 (7)

Equality 1: call combine_multivar_equality

$$\mathbf{A}_0 \mathbf{w} - \mathbf{A}_1 \mathbf{y} = \mathbf{b}_0 + c_2 z \quad \to \quad \mathbf{A}_{01} \mathbf{x}_{wy} = \mathbf{b}_0 \tag{8}$$

where \mathbf{A}_{01} is a column concatenation of the parameters and \mathbf{x}_{wy} is a row contenation of the variables. Inequality 1: call lower_ineq_to_eq

$$\mathbf{A}_0 \mathbf{w} \leq \mathbf{b}_0 \quad \rightarrow \quad \mathbf{A}_0 \mathbf{w} + \mathbf{s}_0 = \mathbf{b}_0, \quad \mathbf{s}_0 \succeq 0$$
 (9)

Inequality 2: call greater_ineq_to_lower_ineq and then lower_ineq_to_eq

$$\mathbf{A}_1 \mathbf{y} \succeq \mathbf{b}_1 \quad \rightarrow \quad -\mathbf{A}_1 \mathbf{y} \preceq -\mathbf{b}_1 \quad \rightarrow \quad -\mathbf{A}_1 \mathbf{y} + \mathbf{s}_1 = -\mathbf{b}_1, \quad \mathbf{s}_1 \succeq 0$$
 (10)

Inequality 3: call lower_ineq_to_eq

$$z \le 2 \quad \to \quad z + s_2 = 2, \quad s_2 \succeq 0 \tag{11}$$

Final combined system:

minimize
$$-\mathbf{c}_0^{\mathsf{T}}\mathbf{w} + \mathbf{c}_1^{\mathsf{T}}\mathbf{y} - c_2 z$$
 (12)

subject to
$$\begin{bmatrix} \mathbf{A}_0 & -\mathbf{A}_1 & 0 \\ \mathbf{A}_0 & 0 & 0 \\ 0 & -\mathbf{A}_1 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} \mathbf{w} \\ \mathbf{y} \\ z \end{bmatrix} + \begin{bmatrix} 0 \\ \mathbf{s}_0 \\ \mathbf{s}_1 \\ s_2 \end{bmatrix} = \begin{bmatrix} \mathbf{b}_0 \\ \mathbf{b}_0 \\ -\mathbf{b}_1 \\ 2 \end{bmatrix}, \quad \begin{bmatrix} \mathbf{s}_0 \\ \mathbf{s}_1 \\ s_2 \end{bmatrix} \succeq 0$$
 (13)

To fold the adding variables into the matrix equation, we can augment the matrix \mathbf{A} and the variable vector \mathbf{x} to include the slack variables. Here's how you can rewrite the final combined system:

minimize
$$\begin{bmatrix} -\mathbf{c}_0^{\top} & \mathbf{c}_1^{\top} & -c_2 & \mathbf{0} & \mathbf{0} \end{bmatrix} \begin{bmatrix} \mathbf{w} \\ \mathbf{y} \\ \mathbf{z} \\ \mathbf{s}_0 \\ \mathbf{s}_1 \\ s_2 \end{bmatrix}$$
 (14)

subject to
$$\begin{bmatrix} \mathbf{A}_{0} & -\mathbf{A}_{1} & 0 & 0 & 0 & 0 \\ \mathbf{A}_{0} & 0 & 0 & \mathbf{I} & 0 & 0 \\ 0 & -\mathbf{A}_{1} & 0 & 0 & \mathbf{I} & 0 \\ 0 & 0 & 1 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} \mathbf{w} \\ \mathbf{y} \\ z \\ \mathbf{s}_{0} \\ \mathbf{s}_{1} \\ s_{2} \end{bmatrix} = \begin{bmatrix} \mathbf{b}_{0} \\ \mathbf{b}_{0} \\ -\mathbf{b}_{1} \\ 2 \end{bmatrix}$$
(15)

$$\begin{bmatrix} \mathbf{w} \\ \mathbf{s}_0 \\ \mathbf{s}_1 \\ s_2 \end{bmatrix} \succeq 0 \tag{16}$$

Here, the slack variables $\mathbf{s}_0, \mathbf{s}_1, s_2$ are now part of the augmented variable vector, and the identity matrices \mathbf{I} are used to incorporate the slack variables into the constraints. This ensures that the system remains in the form $\mathbf{A}\mathbf{x} = \mathbf{b}$, where \mathbf{x} includes all original variables and slack variables. Finally, we call bound_all_vars to bound the variables that have not been bounded yet.

minimize
$$\begin{bmatrix} -\mathbf{c}_0^{\mathsf{T}} & \mathbf{c}_1^{\mathsf{T}} & -\mathbf{c}_1^{\mathsf{T}} & -c_2 & c_2 & \mathbf{0} & \mathbf{0} \end{bmatrix} \begin{bmatrix} \mathbf{w} \\ \mathbf{y}^+ \\ \mathbf{y}^- \\ \mathbf{z}^+ \\ z^- \\ \mathbf{s}_0 \\ \mathbf{s}_1 \\ s_2 \end{bmatrix}$$
 (17)

subject to
$$\begin{bmatrix} \mathbf{A}_{0} & -\mathbf{A}_{1} & \mathbf{A}_{1} & 0 & 0 & 0 & 0 & 0 \\ \mathbf{A}_{0} & 0 & 0 & 0 & 0 & \mathbf{I} & 0 & 0 \\ 0 & -\mathbf{A}_{1} & \mathbf{A}_{1} & 0 & 0 & 0 & \mathbf{I} & 0 \\ 0 & 0 & 0 & 1 & -1 & 0 & 0 & 1 \end{bmatrix} \begin{vmatrix} \mathbf{w} \\ \mathbf{y}^{+} \\ \mathbf{y}^{-} \\ z^{+} \\ z^{-} \\ \mathbf{s}_{0} \\ \mathbf{s}_{1} \\ s_{2} \end{bmatrix} = \begin{bmatrix} \mathbf{b}_{0} \\ \mathbf{b}_{0} \\ -\mathbf{b}_{1} \\ 2 \end{bmatrix}$$
(18)

$$\begin{bmatrix} \mathbf{w} \\ \mathbf{y}^+ \\ \mathbf{y}^- \\ \mathbf{z}^+ \\ z^- \\ \mathbf{s}_0 \\ \mathbf{s}_1 \\ \mathbf{s}_2 \end{bmatrix} \succeq 0 \tag{19}$$