# ECH4905 ChemE Optimization HW 4

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# 1 Problem 1

Consider the following integer programming problem:

maximize 
$$1.2y_1 + y_2$$
  
subject to  $y_1 + y_2 \le 1$   
 $1.2y_1 + 0.5y_2 \le 1$   
 $y_1, y_2 \in \{0, 1\}$ 

#### 1.1 Part a

Solve the first relaxed LP subproblem by hand using the simplex method and derive Gomory cuts based on the LP relaxation.

**Solution:** To tackle this problem we first relax the problem and then turn the problem above into the standard form so we can create a simplex tableau from it.

$$\begin{array}{ll} \text{maximize} & 1.2y_1+y_2\\ \text{subject to} & y_1+y_2+s_1=1\\ & 1.2y_1+0.5y_2+s_2=1\\ & y_1+s_3=1\\ & y_2+s_4=1\\ & y_1,y_2,s_1,s_2,s_3,s_4\geq 0 \end{array}$$

In matrix notation,

minimize 
$$\mathbf{c}^{\top}\mathbf{x}$$
  
subject to  $\mathbf{A}\mathbf{x} = \mathbf{b}$   
 $\mathbf{x} \succeq 0$ 

where

$$\mathbf{c} = \begin{bmatrix} -1.2 \\ -1 \\ 0 \\ 0 \\ 0 \\ 0 \end{bmatrix}, \quad \mathbf{A} = \begin{bmatrix} 1 & 1 & 1 & 0 & 0 & 0 \\ 1.2 & 0.5 & 0 & 1 & 0 & 0 \\ 1 & 0 & 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 & 0 & 1 \end{bmatrix}, \quad \mathbf{b} = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}, \quad \mathbf{x} = \begin{bmatrix} y_1 \\ y_2 \\ s_1 \\ s_2 \\ s_3 \\ s_4 \end{bmatrix}$$

(with a flipped objective component)

The initial simplex tableau for the problem is as follows:

Basic Var	$y_1$	$y_2$	$s_1$	$s_2$	$s_3$	$s_4$	RHS	$\alpha$
$s_1$	1	1	1	0	0	0	1	
$s_2$	1.2	0.5	0	1	0	0	1	
$s_3$	1	0	0	0	1	0	1	
$s_4$	0	1	0	0	0	1	1	
obj	-1.2	-1	0	0	0	0	_	_

we can define the slack variables equal to the right hand side, and this is in turn a basic feasible solution, so we can jump into phase 2.

We select the  $y_1$  as the entering variable and calculate the alpha value for each basic variable

Basic Var	$y_1$	$y_2$	$s_1$	$s_2$	$s_3$	$s_4$	RHS	$\alpha$
$s_1$	1	1	1	0	0	0	1	$\frac{1}{1}$
$s_2$	1.2	0.5	0	1	0	0	1	$\frac{1}{12}$
$s_3$	1	0	0	0	1	0	1	$\frac{1}{1}$
$s_4$	0	1	0	0	0	1	1	$\frac{1}{0}$
obj	-1.2	-1	0	0	0	0	_	_

We pivot this on the 1st column  $(y_1)$  and the 2nd row  $(s_2)$ 

Basic Var	$y_1$	$y_2$	$s_1$	$s_2$	$s_3$	$s_4$	RHS	$\alpha$
$\overline{s_1}$	0	$\frac{7}{12}$	1	$-\frac{5}{6}$	0	0	$\frac{1}{6}$	
$y_1$	1	$\frac{5}{12}$	0	$\frac{5}{6}$	0	0	$\frac{5}{6}$	
$s_3$	0	$-\frac{5}{12}$	0	$-\frac{5}{6}$	1	0	$\frac{1}{6}$	
$s_4$	0	1	0	0	0	1	ľ	
obj	0	-0.5	_	_	_	_	_	_

With blands rule, we pick  $y_2$  and calculate the alpha value for each basic variable.

Basic Var	$y_1$	$y_2$	$s_1$	$s_2$	$s_3$	$s_4$	RHS	$\alpha$
$\overline{s_1}$	0	$\frac{7}{12}$	1	$-\frac{5}{6}$	0	0	$\frac{1}{6}$	$\frac{2}{7}$
$y_1$	1	$\frac{5}{12}$	0	$\frac{5}{6}$	0	0	$\frac{5}{6}$	$\frac{2}{1}$
$s_3$	0	$-\frac{5}{12}$	0	$-\frac{5}{6}$	1	0	$\frac{1}{6}$	$-\frac{1}{5}$
$s_4$	0	1	0	0	0	1	ĺ	$\frac{1}{1}$
obj	0	-0.5	_	_	_	_	_	_

We pivot on the 2nd column  $(y_2)$  and the 1st row  $(s_1)$ .

Basic Var	$y_1$	$y_2$	$s_1$	$s_2$	$s_3$	$s_4$	RHS	$\alpha$
$y_2$	0	1	$\frac{12}{7}$	$-\frac{10}{7}$	0	0	$\frac{2}{7}$	
$y_1$	1	0	$-\frac{5}{7}$	$\frac{60}{42}$	0	0	$\frac{30}{42}$	
$s_3$	0	0	$\frac{5}{7}$	$-\frac{60}{42}$	1	0	12	
$s_4$	0	0	$-\frac{12}{7}$	$\frac{60}{42}^{2}$	0	1	$\frac{30}{42}$	
obj	0	0	0.857	0.286	0	0	0	-1.143

This is the optimal solution to the LP relaxed problem. Now we will derive Gomory cuts from this LP relaxed problem. Since each constraint has a non-integer solution, we can generate a Gomory cut on each constraint.

$$y_{2} + \operatorname{floor}(\frac{12}{7})s_{1} + \operatorname{floor}(\frac{-10}{7})s_{2} \leq \operatorname{floor}(\frac{2}{7})$$

$$y_{1} + \operatorname{floor}(\frac{-5}{7})s_{1} + \operatorname{floor}(\frac{10}{7})s_{2} \leq \operatorname{floor}(\frac{5}{7})$$

$$s_{3} + \operatorname{floor}(\frac{5}{7})s_{1} + \operatorname{floor}(\frac{-10}{7})s_{2} \leq \operatorname{floor}(\frac{2}{7})$$

$$s_{4} + \operatorname{floor}(\frac{-12}{7})s_{1} + \operatorname{floor}(\frac{10}{7})s_{2} \leq \operatorname{floor}(\frac{5}{7})$$

These turn into the cuts

$$y_2 + s_1 - 2s_2 \le 0$$

$$y_1 - s_1 + s_2 \le 0$$

$$s_3 - 2s_2 \le 0$$

$$s_4 - 2s_1 + s_2 \le 0$$

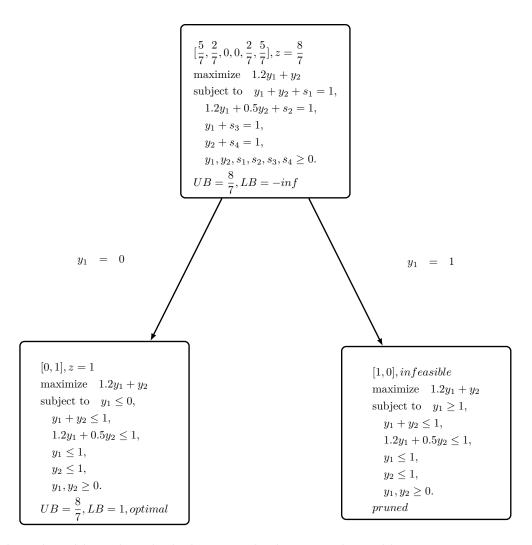
#### 1.2 Part b

Solve the above problem with the branch and bound method by enumerating nodes in the tree and solving the LP subproblems using GAMS.

**Solution:** The initial LP relaxed problem is solved in 1.1, so we can start with the parent node. An important note for this question, I will be solving the LPs in my custom gatorpy LP solver so that I can use them as verification tests. The code used will be available in section 4.1

$$\begin{bmatrix} \frac{5}{7}, \frac{2}{7}, 0, 0, \frac{2}{7}, \frac{5}{7} \end{bmatrix}, z = \frac{8}{7}$$
maximize  $1.2y_1 + y_2$ 
subject to  $y_1 + y_2 + s_1 = 1$ ,
$$1.2y_1 + 0.5y_2 + s_2 = 1$$
,
$$y_1 + s_3 = 1$$
,
$$y_2 + s_4 = 1$$
,
$$y_1, y_2, s_1, s_2, s_3, s_4 \ge 0$$
.
$$UB = \frac{8}{7}, LB = -inf$$

Since all variables are fractional, we can pick the first one  $y_1$  to branch on



Via the branch and bound method, the optimal solution to the problem

$$\begin{array}{ll} \text{maximize} & 1.2y_1 + y_2 \\ \text{subject to} & y_1 + y_2 \leq 1 \\ & 1.2y_1 + 0.5y_2 \leq 1 \\ & y_1, y_2 \in \{0, 1\} \end{array}$$

is  $y_1 = 0, y_2 = 1, z = 1$ .

# 2 Problem 2

Consider the following superstructure for the separation of four chemical components using sharp distillation columns. The total cost of a distillation column is calculated as follows:

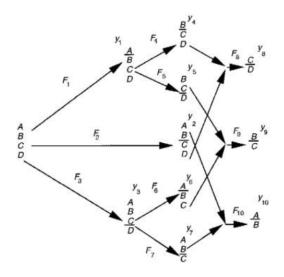


Figure 1: Problem 2 superstructure

$$cost_k = \alpha_k + \beta_k F_k + \gamma_{Hot} Q_k^{Hot} + \gamma_{Cold} Q_k^{Cold}$$

where:

- $\alpha_k$  represents a fixed capital cost,
- $\beta_k$  represents the variable investment cost,
- $\gamma_{\text{Hot/Cold}}$  is the cost of hot/cold utilities, and
- $Q_k^{
  m Hot}/Q_k^{
  m Cold}$  is the total demand of hot and cold utilities (assumed to be equal).

Given:

- Initial feed: 1000 Kmol/h,
- Feed composition (mole fraction): A = 0.15, B = 0.3, C = 0.35, D = 0.2.

and the following data:

#### 2.1 Part A: MILP Formulation

Formulate a Mixed-Integer Linear Programming (MILP) problem to find the optimal sequence of distillation columns. Identify at least 2 logicx based equations that can be formulated to tighten the problem formulation.

# 2.2 Part B: Solve Using GAMS

Solve the problem using GAMS.

		Investi	ment cost	Heat duty
k	Separator	$\alpha_k$ , fixed	$\beta_k$ , variable	coefficients, Kk,
		$(10^3 \text{/yr})$	(103\$hr/kmol yr)	(106kJ/kgmol)
1	A/BCD	145	0.42	0.028
2	AB/CD	52	0.12	0.042
3	ABC/D	76	0.25	0.054
6	A/BC	125	0.78	0.024
7	AB/C	44	0.11	0.039
4	B/CD	38	0.14	0.040
5	BC/D	66	0.21	0.047
10	A/B	112	0.39	0.022
9	B/C	37	0.08	0.036
8	C/D	58	0.19	0.044
Cost	of utilities:			
Co	ooling water	$C_C = 1.3 (10^3 \$/10^6 \text{kJ})$	yr)	
St	eam	$C_H = 34 (10^3 \$/10^6 \text{kJy})$	r)	

Figure 2: Problem 2 Data

# 2.3 Part C: Integer Cut

Once the solution is found:

- $\bullet\,$  Identify the active binary variables in the optimal solution.
- Formulate an integer cut to exclude this solution from the feasible space.
- $\bullet\,$  Solve the problem again to find the next best solution.

# 3 Problem 3

Given are three candidate reactors for the reaction  $A \to B$ , where we would like to produce 10 kmol/h of B. Up to 15 kmol/hr of reactant A are available at a price of \$2/kmol. The data on the three reactors is as follows:

Reactor	Conversion	Cost
1	0.8	$8 + 1.5 \cdot \text{Feed}$
2	0.667	5.4 + Feed
3	0.555	$2.7 + 0.5 \cdot \text{Feed}$

Table 1: Reactor Data

# 3.1 Part A: Superstructure Design

Design a superstructure to represent this problem.

Solution: Not sure if I am oversimplifying this. This superstructure assumes that no A is recycled

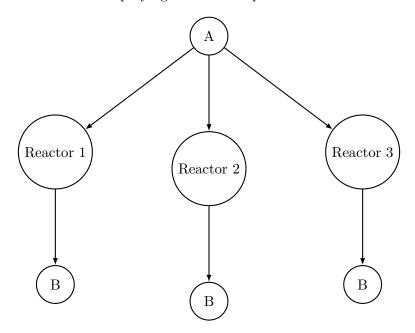


Figure 3: Superstructure for the reaction  $A \to B$  with three reactors.

back into the reactor, and we only choose one reactor. This is generally a simple problem, we are effectively picking which reactor to use to minimize the cost.

# 3.2 Part B: MILP Formulation

Determine a MILP formulation.

**Solution:** In order to solve this MILP formulation, I will start thinking about this in a GAMS/modeling language way starting with parameters.

Parameters: We have parameters  $k_i$  as the  $A \to B$  conversion of reactor i. We also have cost components  $d_i$  as the constant cost of using reactor i and  $c_i$  as the multiplicative cost of feeding A through reactor i. Our demand parameter D for the amount of B we would like to produce and

our supply S that we have available. We can also have the price of using a kmol of A as p. Below are the numerical representations of the problem parameters:

$$k_1 = \frac{4}{5}, k_2 = \frac{2}{3}, k_3 = \frac{5}{9}$$
  
$$d_1 = 8, d_2 = 5.4, d_3 = 2.7; c_1 = 1.5, c_2 = 1, c_3 = 0.5$$
  
$$D = 10, S = 15, p = 2$$

Variables: Our variables for this problem are as follows: We have  $x_i$  as the amount of A delivered to the i-th reactor. In order to model our decision to pick a reactor, we have variables  $y_i$  which are booleans that signify if reactor i has been chosen.

Constraints: We can model our constraints as follows:

$$x_i \leq S, \quad \forall i \in [1,2,3]$$
 Supply Constraint 
$$\sum_{i=1}^3 k_i x_i \geq D$$
 Demand Constraint 
$$\sum_{i=1}^3 y_i = 1$$
 One Reactor Constraint 
$$x_i \leq Sy_i, \quad \forall i \in [1,2,3]$$
 Big-M One Reactor Flow Constraint 
$$x_i \geq 0, \quad \forall i \in [1,2,3]$$
 Non-negativity 
$$y_i \in \{0,1\} \quad \forall i \in [1,2,3]$$
 Binaries

We can also get rid of the supply constraint since the Big-M naturally handles that. The GAMS solution below drops the supply constraint. *Objective:* Our objective in this problem is to minimize the cost while still meeting the demand. We can bunch these costs into different components

$$p \sum_{i=1}^{3} x_i$$
 Supply component 
$$\sum_{i=1}^{3} d_i y_i$$
 Reactor constant component 
$$\sum_{i=1}^{3} c_i y_i$$
 Reactor processing component

# 3.3 Part C: Solve Using GAMS

Solve in GAMS.

**Solution:** The optimal solution was found to be  $x_1 = 12.5, y_1 = 1, z = 34.5$ , and the rest of the variables equal to 0. GAMS code available in section 4.2

# 4 Code

#### 4.1 Problem 1 Code

```
# Parameters
A_{arr} = np.array([[1,1],[1.2,0.5]])
b_arr = np.array([1,1])
c_arr = np.array([1.2,1])
A = Parameter(A_arr)
b = Parameter(b_arr)
c = Parameter(c_arr)
# Variables
y = Variable(2)
# Problem
problem = Problem({
    'maximize': c.T @ y,
    'subject to': [
        A @ y \le b,
        y >= 0,
        y <= 1
    ]
})
# CVX Vars
y_cvx = cp.Variable(2)
# CVX Problem
objective = cp.Maximize(c_arr @ y_cvx)
constraints = [
    A_arr @ y_cvx <= b_arr,
    y_cvx >= 0,
    y_cvx <= 1
problem_cvx = cp.Problem(objective, constraints)
get_test_results(problem, problem_cvx, "HW5 Problem 1 initial LP relaxation")
# branch and bound arrays
I_y_1_arr = np.array([[1,0],[0,0]])
I_y_2_arr = np.array([[0,0],[0,1]])
I_y_1 = Parameter(I_y_1_arr)
I_y_2 = Parameter(I_y_2_arr)
# LEFT SPLIT
# Problem
problem = Problem({
    'maximize': c.T @ y,
    'subject to': [
        A @ y <= b,
        y >= 0,
        y <= 1,
```

```
I_y_1 @ y \le 0,
    ]
})
# CVX Vars
y_cvx = cp.Variable(2)
# CVX Problem
objective = cp.Maximize(c_arr @ y_cvx)
constraints = [
    A_arr @ y_cvx <= b_arr,
    y_cvx >= 0,
    y_cvx <= 1,
    I_y_1_arr @ y_cvx <= 0</pre>
problem_cvx = cp.Problem(objective, constraints)
get_test_results(problem, problem_cvx, "HW5 Problem 1 Branch and Bound Left Split y_1 <= 0")</pre>
# RIGHT SPLIT
# Problem
problem = Problem({
    'maximize': c.T @ y,
    'subject to': [
        A @ y <= b,
        y >= 0,
        y <= 1,
        I_y_1 @ y >= 1,
    ]
})
# CVX Vars
y_cvx = cp.Variable(2)
# CVX Problem
objective = cp.Maximize(c_arr @ y_cvx)
constraints = [
    A_arr @ y_cvx <= b_arr,
    y_cvx >= 0,
    y_cvx <= 1,
    I_y_1_arr @ y_cvx >= 1
problem_cvx = cp.Problem(objective, constraints)
get_test_results(problem, problem_cvx, "HW5 Problem 1 Branch and Bound Right Split y_1 >=1")
Test ID: HW5 Problem 1 initial LP relaxation
CVX: (np.float64(1.14), array([[0.71, 0.29]]), True)
GatORPy: (array(1.14285714), array([0.71428571, 0.28571429, 0.
                                                                        , 0.
                                                                                     , 0.285714
       0.71428571]), True)
```

```
Test passed: False
 Test ID: HW5 Problem 1 Branch and Bound Left Split y_1 <= 0
 CVX: (np.float64(1.0), array([[0., 1.]]), True)
 GatORPy: (array(1.), array([0., 1., 0., 0.5, 1., 0., 0., 0.]), True)
 Test passed: False
 Test ID: HW5 Problem 1 Branch and Bound Right Split y_1 >=1
 CVX: (None, None, False)
 GatORPy: (array(0.8), array([ 1. , -0.4, 0.4, 0. , 0. , 1.4, 0. , -1. ]), True)
 Test passed: False
4.2 Problem 3 Code
 GAMS 49.1.0 5c4d4ed6 Feb 15, 2025
                                         DAX-DAC arm 64bit/macOS - 04/14/25 18:53:07 Page
General Algebraic Modeling System
Compilation
  1 SETS
  2 i
               'set of reactors'
                                      /r1,r2,r3/
  3
  4;
  5 Parameters
  6 k(i)
            'conversion factors'
                                   /r1 0.8, r2 0.666667, r3 0.555555/
            'Reactor constant costs'/r1 8,r2 5.4,r3 2.7/
  7 d(i)
  8 c(i)
            'Reactor coeff costs'
                                   /r1 1.5,r2 1,r3 0.5/
  9;
 10 Scalars
 11 De 'Demand' /10/
 12 Su 'Supply' /15/
 13 pr 'Price' /2/
 14 ;
 15 Variables
 16 z
 17 ;
 18 Positive Variables
 19 x(i)
 20 ;
 21 Binary Variables
 22 y(i)
 23 ;
 24 Equations
 25
         demand_constraint
 26
         one_reactor_constraint
 27
         flow_constraints(i)
 28
         objective_eq
 29 ;
```

```
30 demand_constraint..
                               sum(i,k(i) * x(i)) = g = De;
 31 one_reactor_constraint.. sum(i,y(i)) =e= 1;
 32 flow_constraints(i)..
                              x(i) = 1 = Su*y(i);
 33 objective_eq..
                                z = e = sum(i,pr*x(i) + d(i) * y(i) + c(i) * y(i));
 34
 35 Model Superstructure / all /;
 36 solve Superstructure using MIP minimizing z;
COMPILATION TIME
                            0.000 SECONDS
                                              3 MB 49.1.0 5c4d4ed6 DAX-DAC
GAMS 49.1.0 5c4d4ed6 Feb 15, 2025
                                         DAX-DAC arm 64bit/macOS - 04/14/25 18:53:07 Page 2
General Algebraic Modeling System
Equation Listing SOLVE Superstructure Using MIP From line 36
---- demand_constraint =G=
demand_constraint.. 0.8*x(r1) + 0.666667*x(r2) + 0.555555*x(r3) = G = 10; (LHS = 0, INFES = 10)
---- one_reactor_constraint =E=
one_reactor_constraint.. y(r1) + y(r2) + y(r3) = E = 1; (LHS = 0, INFES = 1 ****)
---- flow_constraints =L=
flow_constraints(r1).. x(r1) - 15*y(r1) = L = 0; (LHS = 0)
flow_constraints(r2).. x(r2) - 15*y(r2) = L = 0; (LHS = 0)
flow_constraints(r3).. x(r3) - 15*y(r3) = L = 0; (LHS = 0)
---- objective_eq =E=
objective_eq.. z - 2*x(r1) - 2*x(r2) - 2*x(r3) - 9.5*y(r1) - 6.4*y(r2) - 3.2*y(r3) = E = 0; (L)
GAMS 49.1.0 5c4d4ed6 Feb 15, 2025
                                         DAX-DAC arm 64bit/macOS - 04/14/25 18:53:07 Page 3
General Algebraic Modeling
                                                  System
                SOLVE Superstructure Using MIP From line 36
Column Listing
---- z
z
               (.LO, .L, .UP, .M = -INF, O, +INF, O)
               objective_eq
       1
```

```
---- x
x(r1)
                (.LO, .L, .UP, .M = 0, 0, +INF, 0)
        0.8
                demand_constraint
                flow_constraints(r1)
        1
       -2
                objective_eq
x(r2)
                (.LO, .L, .UP, .M = 0, 0, +INF, 0)
                demand_constraint
        0.6667
                flow_constraints(r2)
       -2
                objective_eq
x(r3)
                (.LO, .L, .UP, .M = 0, 0, +INF, 0)
        0.5556 demand_constraint
                flow_constraints(r3)
        1
                objective_eq
       -2
---- y
y(r1)
                (.LO, .L, .UP, .M = 0, 0, 1, 0)
                one_reactor_constraint
        1
                flow_constraints(r1)
      -15
       -9.5
                objective_eq
y(r2)
                (.L0, .L, .UP, .M = 0, 0, 1, 0)
        1
                one_reactor_constraint
      -15
                flow_constraints(r2)
                objective_eq
       -6.4
y(r3)
                (.LO, .L, .UP, .M = 0, 0, 1, 0)
        1
                one_reactor_constraint
                flow_constraints(r3)
      -15
       -3.2
                objective_eq
Proven optimal solution
MIP Solution:
                        34.500000
                                      (0 iterations, 0 nodes)
Final Solve:
                                      (0 iterations)
                        34.500000
```

34.500000

Best possible:

Absolute gap: 0.000000 Relative gap: 0.000000

		LOWER	LEVEL	UPPER	MARGINAL				
	J demand_co~ J one_react~	10.0000	10.0000 1.0000	+INF 1.0000	2.5000				
EQU	J flow_constrain	ts							
	LOWER	LEVEL	UPPER	MARGINAL					
r1 r2 r3	-INF -INF -INF	-2.5000							
		LOWER	LEVEL	UPPER	MARGINAL				
EQU	Jobjective~				1.0000				
		LOWER	LEVEL	UPPER	MARGINAL				
VAR	lΖ	-INF	34.5000	+INF					
VAR	ιx								
	LOWER	LEVEL	UPPER	MARGINAL					
r1 r2 r3		12.5000	+INF +INF +INF	0.3333 0.6111					
VAR									
	LOWER	LEVEL	UPPER	MARGINAL					
r1 r2 r3		1.0000	1.0000 1.0000 1.0000	9.5000 6.4000 3.2000					
**** REPORT SUMMARY:  O NONOPT  O INFEASIBLE  O UNBOUNDED									
EXECUTIO	EXECUTION TIME = 0.050 SECONDS 4 MB 49.1.0 5c4d4ed6 DAX-DAC								

USER: GAMS Demo, for EULA and demo limitations see G250131/0001CB-GEN https://www.gams.com/latest/docs/UG%5FLicense.html DC0000

# \*\*\*\* FILE SUMMARY

Input /Users/andresespinosa/Documents/GAMS/Studio/workspace/hw5.gms Output /Users/andresespinosa/Documents/GAMS/Studio/workspace/hw5.lst