#### ECE30030/ITP30010 Database Systems

# Structured Query Language

Reading: Chapter 3

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## Agenda

- Structured query language (SQL)
- SQL data manipulation language (DML)
  - SELECT, FROM, WHERE
  - NULL values
  - Set operations
  - String operations, ordering
  - Aggregate functions, aggregation
- SQL data definition language (DDL) -- NEXT CLASS

## Structured Query Language (SQL)

- SQL: Structured Query Language
  - The principal language used to describe and manipulate relational databases
  - Very high-level
    - Say "what to do" rather than "how to do it"
    - SQL is not specifying data-manipulation details
    - DBMSs figure out the "best" way to execute queries
      - Called "query optimization"
  - Two aspects to SQL
    - Data definition: for declaring database schemas (DDL)
    - Data manipulation: for querying (asking questions about) databases and for modifying the database (DML)

#### **SQL Parts**

- DML provides the ability to query information from the database and to insert tuples into, delete tuples from, and modify tuples in the database
- Integrity the DDL includes commands for specifying integrity constraints
- View definition the DDL includes commands for defining views
- Transaction control includes commands for specifying the beginning and ending of transactions
- Embedded SQL and dynamic SQL define how SQL statements can be embedded within general-purpose programming language
- Authorization includes commands for specifying access rights to relations and views

## A Brief History

- IBM SEQUEL (Structured English Query Language) was developed as a part of the System R project (Chamberlin and Boyce, early 1970s)
  - Later on, SEQUEL was renamed SQL (structured query language)
  - System R → System/38 (1979), SQL/DS (1981), DB2 (1983)
- Relational Software, Inc released the first commercial implementation of SQL, Oracle V2 for VAX computers
  - Relational Software, Inc is now Oracle Corporation
- ANSI and ISO standardized SQL:
  - SQL-86, SQL-89, SQL-92, SQL:1999, ..., SQL:2011, SQL:2016 (current)
  - SQL-92 is supported by the most of database systems

## **Basic Query Structure**

A typical SQL query has the form:

**SELECT** 
$$A_1, A_2, ..., A_n$$
 **FROM**  $r_1, r_2, ..., r_m$  **WHERE**  $P$ 

- A<sub>i</sub> represents an attribute
- $R_i$  represents a relation
- *P* is a predicate
- The result of an SQL query is a relation

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- SQL data definition language (DDL)

## SQL Data Manipulation Language

• The SQL data-manipulation language (DML) allows querying (ask questions about) and modifying the databases

## Running Examples

• Relations (tables): instructor, teaches

#### *Instructor* relation

ID	<b>‡</b>	,⊞ name ÷	dept_name :	≣ salary :
10101		Srinivasan	Comp. Sci.	65000.00
12121		Wu	Finance	90000.00
15151		Mozart	Music	40000.00
22222		Einstein	Physics	95000.00
32343		El Said	History	60000.00
33456		Gold	Physics	87000.00
45565		Katz	Comp. Sci.	75000.00
58583		Califieri	History	62000.00
76543		Singh	Finance	80000.00
76766		Crick	Biology	72000.00
83821		Brandt	Comp. Sci.	92000.00
98345		Kim	Elec. Eng.	80000.00

#### teaches relation

₽ ID ÷	course_id ÷	sec_id +	<b>ş</b> semester ‡	📭 year 🛊
76766	BIO-101	1	Summer	2017
76766	BIO-301	1	Summer	2018
10101	CS-101	1	Fall	2017
45565	CS-101	1	Spring	2018
83821	CS-190	1	Spring	2017
83821	CS-190	2	Spring	2017
10101	CS-315	1	Spring	2018
45565	CS-319	1	Spring	2018
83821	CS-319	2	Spring	2018
10101	CS-347	1	Fall	2017
98345	EE-181	1	Spring	2017
12121	FIN-201	1	Spring	2018
32343	HIS-351	1	Spring	2018
15151	MU-199	1	Spring	2018
22222	PHY-101	1	Fall	2017

## **Basic Query Structure**

A typical SQL query has the form:

**SELECT** 
$$A_1, A_2, ..., A_n$$
 **FROM**  $r_1, r_2, ..., r_m$  **WHERE**  $P$ 

- A<sub>i</sub> represents an attribute
- $R_i$  represents a relation
- *P* is a predicate
- The result of an SQL query is a relation

- The SELECT clause lists the attributes desired in the result of a query
  - Corresponds to the projection operation of the relational algebra
- Example: Find the names of all instructors
  - SQL: **SELECT** name **FROM** instructor;

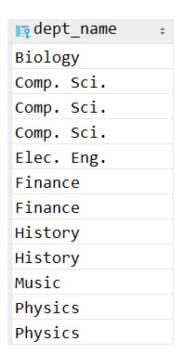




#### Note

- Note: SQL names are case insensitive
  - *E.g.*, *Name* ≡ *NAME* ≡ *name*
  - SQL commands (SELECT, FROM, WHERE, ...) are written in upper case (just a convention)
  - MySQL has an option flag, lower\_case\_table\_names
    - Link: <a href="https://dev.mysql.com/doc/refman/8.0/en/identifier-case-sensitivity.html">https://dev.mysql.com/doc/refman/8.0/en/identifier-case-sensitivity.html</a>

- SQL allows duplicates in relations as well as in query results
  - The keyword ALL specifies that duplicates should not be removed SELECT ALL dept\_name FROM instructor

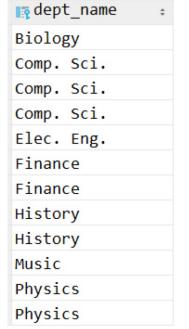


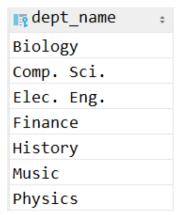


- SQL allows duplicates in relations as well as in query results
  - The keyword ALL specifies that duplicates should not be removed **SELECT ALL** *dept\_name* **FROM** instructor
  - To force the elimination of duplicates, insert the keyword DISTINCT after **SELECT** 
    - Find the department names of all instructor, removing duplicates:

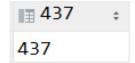
**SELECT DISTINCT** *dept\_name* 

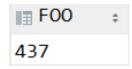
**FROM** *instructor*:



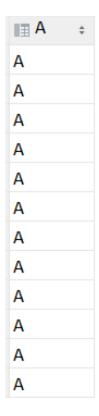


- An asterisk in the select clause denotes "all attributes"
   SELECT \* FROM instructor;
- An attribute can be a literal with no FROM clause SELECT '437';
  - Result is a table with one column and a single row with value "437"
  - Can give the column a name using AS:
     SELECT '437' AS FOO





- An attribute can be a literal with FROM clause
   SELECT 'A' FROM instructor
  - Result is a table with one column and N rows (number of tuples in the instructor table), each row with value "A"



- The SELECT clause can contain arithmetic expressions involving the operation, +, -, \*, and /, and operating on constants or attributes of tuples
  - The query: **SELECT** *ID, name, salary/12* **FROM** *instructor*

would return a relation that is the same as the *instructor* relation, except that the value of the attribute *salary* is divided by 12

III ID ;	: ∏aname ÷	¡ `salary/12` ÷
10101	Srinivasan	5416.666667
12121	Wu	7500.000000
15151	Mozart	3333.333333
22222	Einstein	7916.666667
32343	El Said	5000.0000000
33456	Gold	7250.000000
45565	Katz	6250.000000
58583	Califieri	5166.666667
76543	Singh	6666.666667
76766	Crick	6000.0000000
83821	Brandt	7666.666667
98345	Kim	6666.666667



- The SELECT clause can contain arithmetic expressions involving the operation, +, -, \*, and /, and operating on constants or attributes of tuples
  - Can rename "salary/12" using the AS clause:
     SELECT ID, name, salary/12 AS monthly\_salary
     FROM instructor

II ID	<b>‡</b>	I≣ name	<b>‡</b>	<pre>monthly_salary ;</pre>
10101		Srinivasan		5416.666667
12121		Wu		7500.000000
15151		Mozart		3333.333333
22222		Einstein		7916.666667
32343		El Said		5000.000000
33456		Gold		7250.000000
45565		Katz		6250.000000
58583		Califieri		5166.666667
76543		Singh		6666.666667
76766		Crick		6000.000000
83821		Brandt		7666.666667
98345		Kim		6666.666667



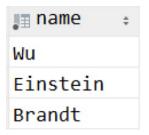
- The WHERE clause specifies conditions that the result must satisfy
  - Corresponds to the selection predicate of the relational algebra
- E.g., To find all instructors in Comp. Sci. dept:
   SELECT name FROM instructor
   WHERE dept\_name = 'Comp. Sci.';



- SQL allows the use of the logical connectives AND, OR, and NOT
- The operands of the logical connectives can be expressions involving the comparison operators <, <=, >, >=, =, and <>
  - <> means not equal (there is no != in SQL)
- Comparisons can be applied to results of arithmetic expressions
- E.g., To find all instructors in Comp. Sci. with salary > 70,000:
   SELECT name FROM instructor
   WHERE dept\_name = 'Comp. Sci.' AND salary > 70000;



- SQL includes a **BETWEEN** comparison operator
- Example: Find the names of all instructors with salary between \$90,000 and \$100,000 (that is,  $\geq$  \$90,000 and  $\leq$  \$100,000)
  - SELECT name
    FROM instructor
    WHERE salary BETWEEN 90000 AND 100000



- Tuple comparison: makes comparisons per tuple
  - SELECT name, course\_id
     FROM instructor, teaches
     WHERE (instructor.ID, dept\_name) = (teaches.ID, 'Biology');

<b>∏</b> name	<b>‡</b>	⊞ course_id	<b>‡</b>
Crick		BIO-101	
Crick		BIO-301	

#### The FROM Clause

- The FROM clause lists the relations involved in the query
  - Corresponds to the Cartesian-product operation of the relational algebra
- Find the Cartesian-product *instructor* × *teaches* **SELECT** \* **FROM** *instructor*, *teaches*;
  - Generates every possible instructor-teaches pairs, with all attributes from both relations
  - For common attributes (e.g., ID), the attributes in the resulting table are renamed using the relation name (e.g., instructor.ID)

### The FROM Clause

# • Find the Cartesian-product *instructor X teaches* **SELECT** \* **FROM** *instructor, teaches*;

instructor.ID	name	dept_name	salary	teaches.ID	course_id	sec_id	semester	year	
10101	Srinivasan	Comp. Sci.	65000	76766	BIO-101	1	Summer		2017
12121	Wu	Finance	90000	76766	BIO-101	1	Summer		2017
15151	Mozart	Music	40000	76766	BIO-101	1	Summer		2017
22222	Einstein	Physics	95000	76766	BIO-101	1	Summer		2017
32343	El Said	History	60000	76766	BIO-101	1	Summer		2017
•••	•••		•••	•••			•••		
10101	Srinivasan	Comp. Sci.	65000	10101	CS-101	1	Fall		2017
12121	Wu	Finance	90000	10101	CS-101	1	Fall		2017
15151	Mozart	Music	40000	10101	CS-101	1	Fall		2017
22222	Einstein	Physics	95000	10101	CS-101	1	Fall		2017
32343	El Said	History	60000	10101	CS-101	1	Fall		2017
	•••	•••	•••	•••	•••		•••		
	•••	•••	•••	•••	•••		•••		
10101	Srinivasan	Comp. Sci.	65000	83821	CS-190	2	Spring		2017
12121	Wu	Finance	90000	83821	CS-190	2	Spring		2017
15151	Mozart	Music	40000	83821	CS-190	2	Spring		2017
	•••		•••	•••			•••		
10101	Srinivasan	Comp. Sci.	65000	10101	CS-315	1	Spring		2018
12121	Wu	Finance	90000	10101	CS-315	1	Spring		2018
15151	Mozart	Music	40000	10101	CS-315	1	Spring		2018
	•••		•••	•••			•••		
•••				•••					



## Implementing JOIN

- Cartesian-product is not very useful directly; but useful combined with WHERE-clause condition (selection operation in relational algebra)
  - Cartesian-product + selection = join

• E.g., Find the names of all instructors who have taught some course and

the course\_id

**SELECT** name, course\_id

**FROM** *instructor* , *teaches* 

**WHERE** instructor.ID = teaches.ID

iii name	course_id :				
Srinivasan	CS-101				
Srinivasan	CS-315				
Srinivasan	CS-347				
Wu	FIN-201				
Mozart	MU-199				
Einstein	PHY-101				
El Said	HIS-351				
Katz	CS-101				
Katz	CS-319				
Crick	BIO-101				
Crick	BIO-301				
Brandt	CS-190				
Brandt	CS-190				
Brandt	CS-319				
Kim	EE-181				

## Implementing JOIN

- Cartesian-product is not very useful directly; but useful combined with WHERE-clause condition (selection operation in relational algebra)
  - Cartesian-product + selection = join
  - Find the names of all instructors in the Music department who have taught

some course and the course\_id

**SELECT** name, course\_id

**FROM** *instructor* , *teaches* 

**WHERE** instructor.ID = teaches.ID

**AND** *instructor. dept\_name = 'Music'* 

course id
 course

MU-199

■ name

Mozart

## The Rename Operation

The SQL allows renaming relations and attributes using the AS clause:

old-name AS new-name

- Find the names of all instructors who have a higher salary than some instructor in 'Comp. Sci.'
  - SELECT DISTINCT T.name
     FROM instructor AS T, instructor AS S
     WHERE T.salary > S.salary AND S.dept\_name = 'Comp. Sci.'





## The Rename Operation

The SQL allows renaming relations and attributes using the AS clause:

old-name AS new-name

- Find the names of all instructors who have a higher salary than some instructor in 'Comp. Sci.'
  - SELECT DISTINCT T.name
     FROM instructor AS T, instructor AS S
     WHERE T.salary > S.salary AND S.dept\_name = 'Comp. Sci.'
- Keyword AS is optional and may be omitted instructor AS T ≡ instructor T

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- SQL data definition language (DDL)

#### **NULL Values**

- It is possible for tuples to have a NULL value for some of their attributes
  - NULL signifies an unknown value or that a value does not exist
- The result of any arithmetic expression involving NULL is NULL
  - E.g., 5 + NULL returns NULL

## IS NULL / IS NOT NULL

- The predicate IS NULL can be used to check for NULL values
  - E.g., Find all instructors whose salary is null SELECT name FROM instructor WHERE salary IS NULL
- The predicate IS NOT NULL succeeds if the value on which it is applied is not null

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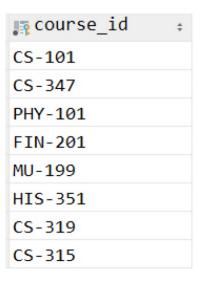
## **Set Operations**

- Set operations UNION, INTERSECT, and EXCEPT
  - Each of the above operations automatically eliminates duplicates
- To retain all duplicates, use ALL:
  - UNION ALL
  - INTERSECT ALL
  - EXCEPT ALL
- C.f., SELECT retains all duplicates by default

## **Set Operations: UNION**

- Find courses that ran in Fall 2017 or in Spring 2018
  - (SELECT course\_id FROM teaches WHERE semester = 'Fall' AND year = 2017) UNION

(**SELECT** *course\_id* **FROM** *teaches* **WHERE** *semester* = 'Spring' **AND** *year* = 2018)



## Set Operations: INTERSECT

- Find courses that ran in Fall 2017 and in Spring 2018
  - (SELECT course\_id FROM teaches WHERE semester = 'Fall' AND year = 2017)
     INTERSECT
     (SELECT course\_id FROM teaches WHERE semester = 'Spring' AND year = 2018)
  - C.f., MySQL does NOT support INTERSECT
    - One can emulate INTERSECT using JOIN (we'll study JOIN later)
    - SELECT LT.course\_id
       FROM (SELECT course\_id FROM teaches WHERE semester = 'Fall' AND year = 2017)
       AS LT
       JOIN (SELECT course\_id FROM teaches WHERE semester = 'Spring' AND year = 2018) AS RT
       ON LT.course id=RT.course id;

## **Set Operations: EXCEPT**

- Find courses that ran in Fall 2017 but not in Spring 2018
  - (SELECT course\_id FROM teaches WHERE semester = 'Fall' AND year = 2017)
     EXCEPT
     (SELECT course\_id FROM teaches WHERE semester = 'Spring' AND year = 2018)
  - *C.f.*, MySQL does NOT support EXCEPT
    - One can emulate EXCEPT using NOT IN
    - SELECT course\_id FROM teaches WHERE semester = 'Fall' AND year = 2017
       AND course\_id NOT IN
       (SELECT course\_id FROM teaches
       WHERE semester = 'Spring' AND year = 2018);

```
course_id :
CS-347
PHY-101
```

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## **String Operations**

- SQL includes a string-matching operator for comparisons on character strings
- The operator LIKE uses patterns that are described using two special characters:
  - percent (%) The % character matches any substring
  - underscore (\_) The \_ character matches any character
- Find the names of all instructors whose name includes the substring "ri"

**SELECT** name **FROM** instructor **WHERE** name **LIKE** '%ri%'



## **String Operations**

- Escape character: Use backslash (\) as the escape character
  - *E.g.*, Match the string "100%" LIKE '100 \%' ESCAPE '\'

## **String Operations**

- Patterns are case sensitive
- Pattern matching examples:
  - 'Intro%' matches any string beginning with "Intro"
  - '%Comp%' matches any string containing "Comp" as a substring
  - '\_\_\_' matches any string of exactly three characters
  - '\_\_\_ %' matches any string of at least three characters
- SQL supports a variety of string operations such as
  - concatenation (using "||")
  - converting from upper to lower case (and vice versa)
  - finding string length, extracting substrings, etc.

# Ordering the Display of Tuples

- List in alphabetic order the names of all instructors
  - SELECT DISTINCT name FROM instructor ORDER BY name







# Ordering the Display of Tuples

- Can sort on multiple attributes
  - E.g., SELECT dept\_name, name FROM instructor ORDER BY dept\_name, name

dept_name	. name ÷
Biology	Crick
Comp. Sci.	Brandt
Comp. Sci.	Katz
Comp. Sci.	Srinivasan
Elec. Eng.	Kim
Finance	Singh
Finance	Wu
History	Califieri
History	El Said
Music	Mozart
Physics	Einstein
Physics	Gold

# Ordering the Display of Tuples

- We may specify DESC for descending order or ASC for ascending order, for each attribute; ascending order is the default
  - E.g., ORDER BY name DESC



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## **Aggregate Functions**

- These functions operate on the multiset of values of a column of a relation, and return a value
  - AVG: average value
  - MIN: minimum value
  - MAX: maximum value
  - **SUM:** sum of values
  - **COUNT:** number of values

# Aggregate Functions Examples

- Find the average salary of instructors in the Computer Science department
  - SELECT AVG(salary)
     FROM instructor
     WHERE dept\_name= 'Comp. Sci.';

- Find the total number of instructors who teach a course in the Spring 2018 semester
  - SELECT COUNT(DISTINCT ID)

    FROM teaches

    WHERE semester = 'Spring' AND year = 2018;
- Find the number of tuples in the teaches relation
  - SELECT COUNT (\*) FROM teaches;



## Aggregate Functions: Group By

- Find the average salary of instructors in each department
  - SELECT dept\_name, AVG(salary) AS avg\_salary
     FROM instructor
     GROUP BY dept\_name;

ID	пате	dept_name	salary
76766	Crick	Biology	72000
45565	Katz	Comp. Sci.	75000
10101	Srinivasan	Comp. Sci.	65000
83821	Brandt	Comp. Sci.	92000
98345	Kim	Elec. Eng.	80000
12121	Wu	Finance	90000
76543	Singh	Finance	80000
32343	El Said	History	60000
58583	Califieri	History	62000
15151	Mozart	Music	40000
33456	Gold	Physics	87000
22222	Einstein	Physics	95000

dept_name	<b>‡</b>	⊞ avg_salary :
Biology		72000.000000
Comp. Sci.		77333.333333
Elec. Eng.		80000.000000
Finance		85000.000000
History		61000.000000
Music		40000.000000
Physics		91000.000000

### Aggregation

- Attributes in SELECT clause outside of aggregate functions must appear in GROUP BY list
  - /\* erroneous query \*/
    SELECT dept\_name, ID, AVG(salary)
    FROM instructor
    GROUP BY dept\_name;

dept_name	‡ <b>∥≣ ID</b>	<b>‡</b>	<pre>Tava (salary) *</pre>
Biology	76766		72000.000000
Comp. Sci.	10101		77333.333333
Elec. Eng.	98345		80000.000000
Finance	12121		85000.000000
History	32343		61000.000000
Music	15151		40000.000000
Physics	22222		91000.000000

## Aggregate Functions – Having Clause

- Find the names and average salaries of all departments whose average salary is greater than 65000
  - SELECT dept\_name, AVG(salary) AS avg\_salary
     FROM instructor
     GROUP BY dept\_name
     HAVING AVG(salary) > 65000;

<pre>■ dept_name</pre>	■ avg_salary ‡
Biology	72000.000000
Comp. Sci.	77333.333333
Elec. Eng.	80000.000000
Finance	85000.000000
Physics	91000.000000

## Aggregate Functions – Having Clause

 Note: predicates in the HAVING clause are applied after the formation of groups whereas predicates in the WHERE clause are applied before forming groups

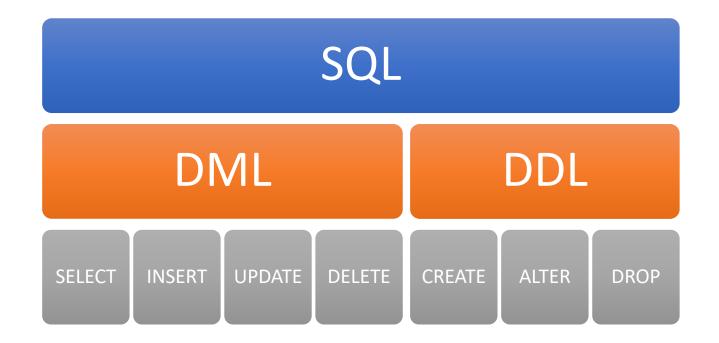
**SELECT** dept\_name, **AVG**(salary) **AS** avg\_salary **FROM** instructor **GROUP BY** dept\_name **HAVING AVG**(salary) > 65000;

**SELECT** dept\_name, **AVG**(salary) **AS** avg\_salary **FROM** instructor **WHERE** salary > 65000 **GROUP BY** dept\_name;

■ dept_name	<pre>avg_salary \$</pre>
Biology	72000.000000
Comp. Sci.	77333.333333
Elec. Eng.	80000.000000
Finance	85000.000000
Physics	91000.000000

■ dept_name ‡	■ avg_salary ‡
Biology	72000.000000
Comp. Sci.	83500.000000
Elec. Eng.	80000.000000
Finance	85000.000000
Physics	91000.000000

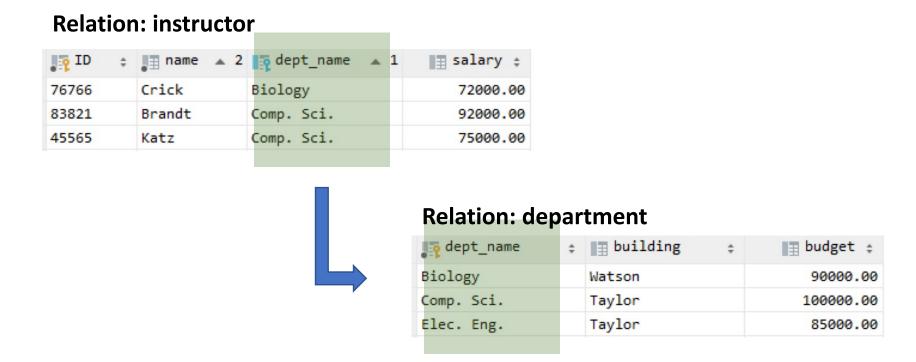
# **SQL Commands**



- Basic syntax
  - Insert data into every column:
    - INSERT INTO tablename VALUES (col1\_value, col2\_value, ...)
    - Must list values in the same order as in the table schema
    - If some data values are unknown, must type NULL
    - For character sequences, use quotation marks
      - Single quotation marks are preferred (but double quotation marks are allowed)
      - Value in quotations is case-sensitive
  - Insert data into selected columns
    - INSERT INTO tablename (col1\_name, col3\_name, col4\_name, ...) VALUES (col1\_value, col3\_value, col4\_value, ...)

- Add a new tuple to course
  - INSERT INTO course VALUES ('CS-437', 'Database Systems', 'Comp. Sci.', 4);
- or equivalently
  - INSERT INTO course (course\_id, title, dept\_name, credits)
     VALUES ('CS-437', 'Database Systems', 'Comp. Sci.', 4);
- Add a new tuple to student with tot\_creds set to null
  - INSERT INTO student VALUES ('3003', 'Green', 'Finance', null);

- A foreign key specifies that an attribute from one relation has to map to a tuple in another relation
  - Value in one relation must appear in another relation



- A foreign key specifies that an attribute from one relation has to map to a tuple in another relation
  - Value in one relation must appear in another relation
- Make sure all foreign keys that new row references have already been added to database
  - One cannot insert a foreign key value unless the corresponding value exists in the referenced relation

- Inserting results of other SELECT query
  - Make each student in the Music department who has earned more than 144 credit hours an instructor in the Music department with a salary of \$18,000
    - INSERT INTO instructor
       SELECT ID, name, dept\_name, 18000
       FROM student
       WHERE dept\_name = 'Music' AND total\_cred > 144;
  - The **SELECT FROM WHERE** statement is evaluated fully before any of its results are inserted into the relation
    - Otherwise queries like
       INSERT INTO table1 SELECT \* FROM table1
       would cause problem

### **UPDATE**

- Basic syntax
  - Updating a table
    - UPDATE tablename SET col1\_name = new\_col1\_value, col2\_name = new\_col2\_value, ...;
  - Updating a table with conditions
    - UPDATE tablename SET col1\_name = new\_col1\_value, col2\_name = new\_col2\_value, ... WHERE predicate;

#### **UPDATE**

- Give a 5% salary raise to all instructors
  - **UPDATE** *instructor* **SET** *salary* = *salary* \* 1.05
- Give a 5% salary raise to those instructors who earn less than 70000
  - **UPDATE** *instructor*  **SET** *salary* = *salary* \* 1.05 **WHERE** *salary* < 70000;
- Give a 5% salary raise to instructors whose salary is less than average
  - UPDATE instructor
     SET salary = salary \* 1.05
     WHERE salary < (SELECT AVG(salary) FROM instructor);</li>

#### **UPDATE**

- Increase salaries of instructors whose salary is over \$100,000 by 3%, and all others by a 5%
  - Write two UPDATE statements:

```
UPDATE instructor
SET salary = salary * 1.03
WHERE salary > 100000;
```

```
UPDATE instructor
SET salary = salary * 1.05
WHERE salary <= 100000;
```

- The order is important
- Can be done better using the case statement (next slide)

## **CASE Statement for Conditional Update**

- The following query is equivalent to the previous UPDATE queries
  - UPDATE instructor
    SET salary = CASE
    WHEN salary <= 100000 THEN salary \* 1.05
    ELSE salary \* 1.03
    END

## **UPDATE** with Scalar Subqueries

- Recompute and update tot\_creds value for all students

#### DELETE

- Basic syntax
  - To remove specific rows
    - **DELETE FROM** *tablename* **WHERE** *predicate*;
  - To remove all rows
    - **DELETE FROM** *tablename*;
    - This is equivalent to TRUNCATE: TRUNCATE (TABLE) tablename;
    - One cannot truncate a table with foreign key constraints
      - Must disable the constraints first (we will cover ALTER when we study SQL DDL):
         ALTER TABLE tablename
         DISABLE CONSTRAINT constraint name;

#### DELETE

- Delete all instructors
  - **DELETE FROM** *instructor*;
- Delete all instructors from the Finance department
  - DELETE FROM instructor
     WHERE dept\_name= 'Finance';
- Delete all tuples in the instructor relation for those instructors associated with a department located in the Watson building
  - DELETE FROM instructor
     WHERE dept name IN (SELECT dept name
     FROM department
     WHERE building = 'Watson');

#### DELETE

- Delete all instructors whose salary is less than the average salary of instructors
  - Example: DELETE FROM instructor
     WHERE salary < (SELECT AVG (salary)</li>
     FROM instructor);
- Issue: as we delete tuples from *instructor*, the average salary changes
  - Solution used in SQL:
    - 1. First, compute **AVG**(salary) and find all tuples to delete
    - 2. Next, delete all tuples found above (without recomputing **AVG** or retesting the tuples)

### **EOF**

- Coming next:
  - More on Structured Query Language