Petri Net Simulator for a Simple Processor

Petri Net Simulator for a Simple Processor: In this project, you will create a Petri Net simulator for a simple processor. The model will use colored tokens (token with values) rather than the default Petri net. Your simulator should generate step-by-step simulation of the Petri net model of the processor described below. You can use **C**, **C++**, **Java** or **Python** to implement this project. Please go through this document first, and then view the sample input/output files in the project assignment.
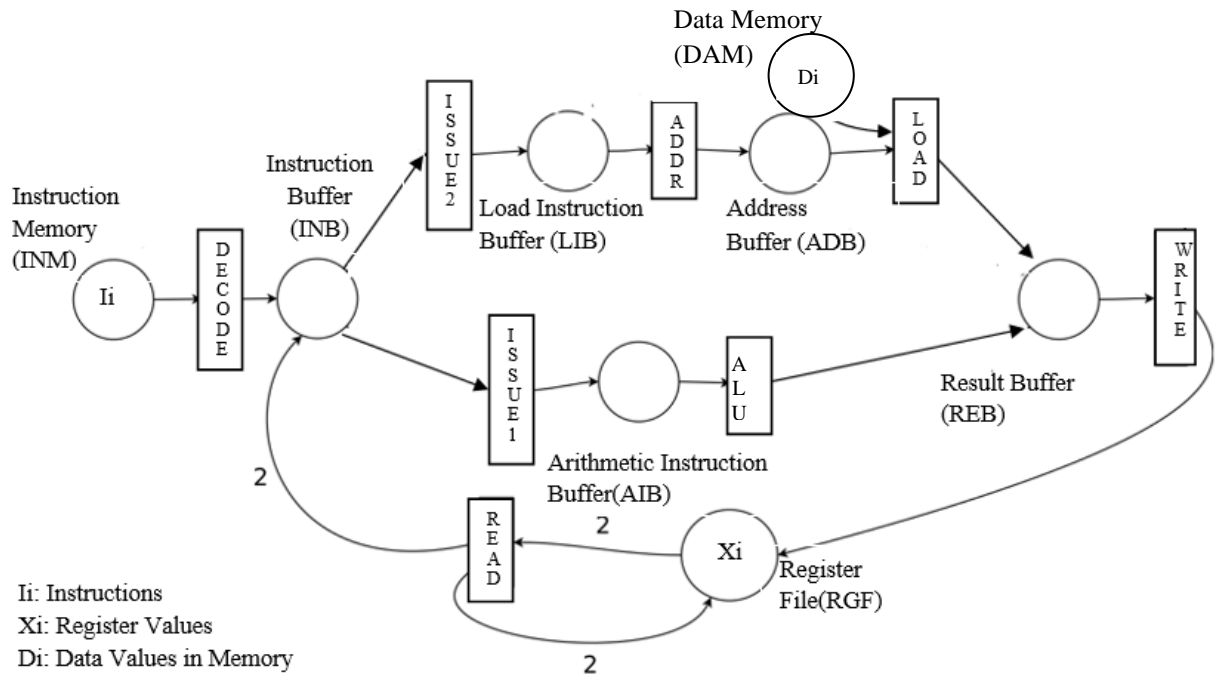

Figure 1. Petri Net Model of a MIPS Processor

We first describe three important places (instruction memory, register file, and data memory) of the Petri net model. Next we describe eight transitions. The remaining places can be viewed as buffers. Each arc carries (consumes) 1 token unless marked otherwise. **You do not have to worry about any hazards or exceptional scenarios. We will not provide any testcase that can produce any hazards or exceptions.**

THREE IMPORTANT PLACES

1. Instruction Memory (INM):

The processor to be simulated only supports five types of instructions: *add* (ADD), *subtract* (SUB), *logical and* (AND), *logical or* (OR), and *load* (LD). At a time step, the place denoted as Instruction Memory (INM) can have up to 16 instruction tokens. This is shown as **Ii** in Figure 1. We will provide an input file (instructions.txt) with up to 16 instruction tokens. It supports the following instruction format. Please note that both source operands are always registers.

<Opcode>, <Destination Register>, <First Source Operand>, <Second Source Operand>

Sample instruction tokens and equivalent functionality are shown below:
<ADD, R1, R2, R3>   ➔ R1 = R2 + R3
<SUB, R1, R2, R3>   ➔ R1 = R2 – R3
<AND, R1, R2, R3>   ➔ R1 = R2 & R3
<OR, R1, R2, R3>   ➔ R1 = R2 | R3
<LD, R1, R2, R3>   ➔ R1 = DataMemory[R2+R3]

2. Register File (RGF):

This processor supports up to 8 registers (R0 through R7). At a time step it can have up to 8 tokens. The token format is <registername, registervalue>, e.g., <R1, 5>. This is shown as **Xi** in Figure 1. We will provide an input file (registers.txt) with 8 register tokens that you can use to initialize the registers. You can assume that the content of a register can vary between 0 - 63.

3. Data Memory (DAM):

This processor supports up to 8 locations (0 – 7) in the data memory. At a time step it can have up to 8 tokens. The token format is <address, value>, e.g., <6, 5> implies that memory address 6 has value 5. This is shown as **Di** in Figure 1. We will provide an input file (datamemory.txt) with 8 data tokens that you can use to initialize the data memory locations. You can assume that the content of a data memory location can vary between 0 - 63.

EIGHT TRANSITIONS

1. READ:
The READ transition is a slight deviation from traditional Petri net semantics since it does not have any direct access to instruction tokens. Assume that it knows the top (in-order) instruction in the Instruction Memory (INM). It checks for the availability of the source operands in the Register File (RGF) for the top instruction token and passes them to Instruction Buffer (INB) by replacing the source operands with the respective values. For example, if the top instruction token in INM is <ADD, R1, R2, R3> and there are two tokens in RGF as <R2,5> and <R3,7>, then the instruction token in INB would be <ADD,R1,5,7> once both READ and DECODE transitions are activated. Both READ and DECODE transitions are executed together. Please note that when READ consumes two register tokens, it also returns them to RGF in the same time step (no change in RGF due to READ).

2. DECODE:

The DECODE transition consumes the top (in-order) instruction (one token) from INM and updates the values of the source registers with the values from RGF (with the help of READ transition, as described above), and places the modified instruction token in INB.

3. ISSUE1:

ISSUE1 transition consumes one arithmetic/logical (ADD, SUB, AND, OR) instruction token (if any) from INB and places it in the Arithmetic Instruction Buffer (AIB).

4. ISSUE2:

ISSUE2 transition consumes one load (LD) instruction token (if any) from INB and places it in the Load Instruction Buffer (LIB).

5. Arithmetic Logic Unit (ALU)

ALU transition performs arithmetic/logical computations as per the instruction token from AIB, and places the result in the result buffer (REB). The format of the token in result buffer is same as a token in RGF i.e., <destination-register-name, value>.

6. Address Calculation (ADDR)

ADDR transition performs effective (data memory) address calculation for the load instruction by adding the contents of two source registers. It produces a token as <destination-register-name, data memory address> and places it in the address buffer (ADB).

7. LOAD:

The LOAD transition consumes a token from ADB and gets the data from the data memory for the corresponding address. Assume that you will always have the data for the respective address in the data memory in the same time step. It places the data value (result of load) in the result buffer (REB). The format of the token in result buffer is same as a token in RGF i.e., <destination-register-name, data value>.

8. WRITE

Transfers the result (one token) from the Result Buffer (REB) to the register file (RGF). If there are more than one token in REB in a time step, the WRITE transition writes the token that belongs to the in-order first instruction.

<u>Command Line and Input/Output Formats:</u>

Command Line: The simulator should be executed with the following command line. Assume that the input files (with the following names) will be available in the same directory. You are also expected to produce the output file in the same directory.

./Psim or **java Psim** or **python3 Psim.py**

Please hardcode the input and output files as follows:
>   Instructions (input): instructions.txt
>   Registers (input): registers.txt
>   Data Memory (input): datamemory.txt
>   Simulation (output): simulation.txt

File Formats:

We will provide inputs in the specific format as listed below:

Input Register File Format: (see registers.txt for example)

```
<register name,value>
…
```

Input Data Memory File Format: (see datamemory.txt for example)

```
<memory address,data value>
…
```

Input Instruction Memory File Format (see instructions.txt for example):

```
<opcode,dest,src1,src2>
…
```

Step-by-step Snapshot Output File Format (see simulation.txt for example): **Please note the following comments are not part of the output format**.

```
STEP 0:
INM: I1,I2,I3,…         # Where Ii are comma separated instruction tokens.
INB:                    # Comma separated tokens with source values.
AIB:                    # Comma separated arithmetic instruction tokens
LIB:                    # Comma separated load instruction tokens
ADB:                    # Comma separated address tokens
REB:                    # Comma separated result buffer tokens
RGF:RF1,RF2,…           # Comma Separated register file tokens.
DAM:D1,D2,…             # Comma Separated data memory tokens.
<blank_line>
STEP 1:
…
```

Continue until the end of simulation. End of simulation is determined when none of the transitions can be fired in a time step.

Additional Notes: (refer to sample input and output files for ease of understanding)

- STEP 0 values represent initial states of all the places. STEP 1 represents the tokens of all the places at the end of first time step. In general, STEP **i** values should reflect the tokens of all the places at the end of time step **i**. In each time step, you are supposed to execute each transition exactly once (if it has required input tokens at the beginning of that cycle).

- When there are more than one tokens in a place, please print them in instruction order (**in-order**) except for DAM and RGF. The token for DAM and RGF should be printed in the sorted order based on memory address or register name (starting with the smallest), respectively.