Slide 1:

Hallo, everyone, I am Han Zhao, thank you for waiting me for a long time. Now, it’s my turn to introduce my topic “Post-quantum secure PUF authentication using LPN”.

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I will introduce my topic in these parts. Firstly, I describe the introduction of this topic, it contains the concept of LPN and the motivation of the topic. And then, I will introduce the construction of this authentication system in detail, it refers to two stages, the first stage is the enrollment phase and the second is the authentication phase. After that, I will give a summary of the whole process and show you my schedule in the following time. Finally, we can discuss the problems or questions about this topic together.

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Firstly, the concept of LPN, Learning Parity with Noise. It can be described as mathematical problem, selecting a secret s in Galois filed 2 randomly with the length of n then selecting a vector A in Galois filed 2 randomly, the length of A is equal to n, after that, selecting a bit noise e, and e follows the Bernoulli distribution with the parameter epsilon. That means, the probability is epsilon when the noise e is equal to one. The output b is equal to the product of A and s plus noise. This is only one sample, if we collect m samples, then A is a matrix with the dimension of m multiple n and normally m is larger than n, the length of error is m, but the length of secret is always n. So, at last, the goal of this problem is to extract s given public known matrix A and the outputs b.

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Ok, then, the motivation of my topic, there are two reasons, on the theoretical side, LPN based schemes offer a very strong security guarantee. The LPN problem is equivalent to the problem of decoding random linear codes. Until now, there is no algorithm that can solve s in polynomial time, so this problem can be considered as a very hard problem.

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I selected 3 famous algorithms for solving LPN problem, the result shows that these algorithms solve s in exponential time with a number of samples and the computational complexity of LPN problem is related to the length of the secret n and the parameter of noise distribution. The larger n is, the harder the problem is.

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On the practical side, LPN based schemes are often extremely simple and efficient in terms of code-size as well as time and space requirements, so it’s very suitable for hardware implementation. In practice, there are many cryptographic applications based on LPN, such as authentication or encryption protocols, in the past, strong PUFs were used for authentication, the inputs are very long, and the corresponded outputs are very short. Its security is based on information theoretical complexity, however, nowadays it’s not secure anymore, because some of strong PUFS’ outputs are repetitive, it can be analyzed out by machine learning. However, compared with strong PUF-based authentication, LPN-based authentication is more secure, so far there is no quantum computer that can solve the LPN problem in polynomial time, that means, LPN problem has no known quantum attacks, so it can be used for post-quantum authentication-system.

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After the introduction of the background, now I go into the main part of the topic, the construction of the authentication system, it refers to two phases, enrollment phase and authentication phase. In the first stage, the Challenge-Response-Pairs are collected and stored in database. In the second stage, the extracted Challenge-Response-Pairs are matched with the reference in the database. If they are identical, then the authentication succeeds, otherwise, it fails. I introduce the enrollment phase firstly.

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In this stage, the aim is to collect challenge and response pairs, there are mainly 3 parts, POK, encoding-module and Hash module.

POK: namely the weak PUF, Due to the internal manufacturing variability the PUF output is unclonable and unique. Here the weak PUF is used for key storage. E is the POK output, it’s sign of the differential count in frequency between two Ring Oscillators. r\_noise is the noise from the extra environment, the two are added into the encoding module as noise.

TRNG: stands for True Random Number Generator, it’s used for producing random number, namely the secret, binary value, 0 or 1.

Here is encoding module: matrix A is also random, the output of this module is b, b is equal to A multiple s and then plus e, so this is the LPN problem.

Then, Hash module: Due to the properties of hash function onewayness and non-malleability, so we regard the hash function of s as the response, b and the hash function of s,b is challenge. Even if you know the Hash value of s, you can’t calculate s, or you know both the Hash value of s,b and b, you can’t calculate s.

In the process, POK is used as one-time-pad and s is also random, so the message b is unique. This method is different from the traditional encryption method, there is no encryption-module, the traditional method, such as AES or RSA encryption algorithm, they adopt encryption algorithm to exchange information, the ciphertext is usually short. But here the text can be longer, namely the value of m. So, this structure is more secure than the traditional.

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In the authentication phase, the aim is to match the extracted challenge response pairs with the reference. There are mainly 4 parts, POK, decoding module, Hash module and a small matching module.

Here the POK output: is different from the former measured in the first stage, produces the noise for the decoding module.

In decoding module: given public known matrix A and challenge b, if the device can return the corresponded response, then the authentication succeeds. As for the matching module, there are two variables as challenge, b and Hash function of b,s, before returning the hash function of s, the hash function of s,b should be matched firstly, this step is necessary in case that one bit of b is changed by the adversary and this bit doesn’t influence the extraction of correct s, so it can improve the security of this system. If this step succeeds, then checking the hash value of s, if they are identical, this authentication process succeeds, otherwise it fails. In this module, solving s is equivalent to solving LPN problem. Because s is random and e prime is also random, it’s difficult for the adversary to compute s in polynomial time. So I think this system is secure. But for the device, the problem is how to extract s.

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There are two methods, Gaussian elimination and error correction algorithm. Gaussian elimination is only suitable for solving linear equations, obviously, in this system this is not linear, unless we can find the stable POK bits, but the workload is very large, in hardware using Gaussian elimination algorithm is also too complex and can’t guarantee the accuracy of s, so I adopted error correction algorithm. As for the security analysis, maybe error correction code can cause security reduction, this is possible, but there is POK output, it’s random and unpredicted, so I think it’s secure and a trade-off decision must be made between the security and the accuracy.

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Which kind of code can be used? Before I planed to use LDPC code, the parity check matrix of LDPC code is sparse, it can ‘Eigenschaft von LDPC, warum passt nicht?’ 2 minuten 1 seite PPT

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There are several candidates, such as LDPC, it’s useful for implementation in hardware, but its generator isn’t sparse, that means, the encoding process is very complex, as for Hamming code, the number of error correction is limited, there are about 10 to 15 percent errors, so I select Reed Muller code.

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Firstly, the generator matrix of Reed Muller code is very simple to construct and in decoding module, there can be no parity-check-matrix to correct error, that means, in hardware there is no need for the extra storage of this matrix. Besides, it has better error correction ability, especially for long code.

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The Plotkin-Construction method can be used to generate Reed Muller code, Reed Muller code can be made from two subcodes, namely u and v, u belongs to Reed Muller r, m-1, and v belongs to Reed Muller r-1, m-1. This is a relationship between the parameters r, m and the length of codeword n and the length of information bits. The Hamming distance of the codeword is equal to 2 square m minus r, so the long code has better error correction property.

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This is a tree of Plotkin-Construction.

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Trapdoor：

Trapdoor is defined as c, c is the confidence information, the actual differential count between two ring oscillators. e is the sign of c, the larger c is, the more stable e is.

Hash function:

Hash function is a function that can be used to map data of arbitrary size to data of fixed size. It has the properties: One-wayness and Non-malleability. That means, even though somebody knows H(s), he cannot calculate s or according to b and H(s,b) he cannot calculate s. It can improve the security of the system.

Strong PUF vs Weak PUF:

Strong PUF: normally, strong PUF are used for authentication, their inputs are very long, and the corresponded output is very short. That means, the input and the output of the authentication are not same. Its security is based on information theoretical complexity, but nowadays it is not secure any more, because some of its outputs are repetitive so it can be analyzed by machine learning.

However,

Weak PUF: it is used to generate the symmetric key, the length of the weak PUF’s input and output are the same, that means, the weak PUF outputs can be used as intended error to produce one-time-pad. Combined with LPN, the security of the authentication system can increase quitely. Because LPN is based on computational complexity and Hardware. Until now

LPN has no known quantum attacks, so this makes it a suitable choice for the post-quantum authentication-system. Many researchers have gone into the solution of LPN for many years, however, LPN is still a NP problem, NP is the set of problems solvable by a nondeterministic algorithm whose running time is bounded by a polynomial in the length of the input.

Cryptosystem VS Authentication-system:

The traditional cryptosystem, such as AES and RSA, they adopt encryption algorithm to exchange information. The ciphertext is usually short. However, in this authentication system there is no encryption algorithm, the ciphertext is very long and only for once use, the weak PUF outputs are random and can’t be predicted, as well as that the secret is random.

Post-quantum cryptography refers to cryptographic algorithms (usually public-key algorithms) that are thought to be secure against an attack by a quantum computer.

Post-quantum cryptography is distinct from [quantum cryptography](https://en.wikipedia.org/wiki/Quantum_cryptography), which refers to using quantum phenomena to achieve secrecy and detect eavesdropping.

Notation:

A is uniformly public-known random matrix\\

s is chosen uniformly at random with length n\\

e is the added noise \\

Trapdoor：

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Hash function:

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How to extract correct s according to the challenge, namely the helper data and the public known matrix A? Firstly, the Gaussian elimination algorithm is recommended, but this algorithm is only suitable for solving linear equations. In this system, the intended error contains the POK output and the noise from the extra environment, so if we use Gaussian elimination algorithm here, the probability of extracting incorrect s is very high, unless we can find the stable POK bits, but this workload is very large, it seems impractical. Besides, the computation complexity of Gaussian is also very high, so we adopt error correction algorithm in the decoding module to check the errors and correct them. Then the accuracy of extracting correct s can increase.

Error correction codes:

We considered several kinds of error correction codes, such as Hamming code, LDPC code and so on. Among them, we decided to select Reed Muller code. There are several reasons. The construction of Reed Muller code is easy and there is no parity check matrix, only generator matrix. That means, in the decoding module, if I use the other code, such as LDPC code, the parity check code is sparse, it is very useful for decoding, but its generator matrix is not sparse, it can cause that its encoding is very complex. And generally, it is used for thousand long code.

Finite-state machine, FSM