CS 211: High Performance Computing Project 3

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Parallel Sieve of Eratosthenes 10¹⁰

1.1 sieve0

The basic version is using Eratosthenes algorithm to find all prime within 10^{10} . This parallel algorithm shows that the performance has been improved by increasing the number of processing cores, and the running time is reduced to 50% when the number of processing cores is doubled.

1.2 sieve1

In this version, all even number are removed from original data, because all even number except 2 cannot be a prime. The result shows that all processing time (for different processor number) are decreased to a half of sieve0.

1.3 sieve2

In this version, all each process of the program finds its own sieving primes via local computations instead of broadcasts for saving communications time. The result shows that the processing time does not decrease so much (even cost more) when there is only few processing cores, but it shows relatively decreasing by increasing the process number.

For implementation, there is two way to calculate local primes. The first one is one-step calculation, which means the calculations are done in loops with marking those who is not prime. This implementation share the cache with marking operations, but it only calculates the primes for its own process and is limited by high_value, which means this implementation reduces the loop number.

The second one is all processes calculate all local primes first, then do loop to marking all non-prime numbers. Comparing with first implementation, this one make sure every process doing same amount of works as others for calculating local primes, saving loops for other processes may not help because the performance of a parallel program is decided by the slowest process (usually the last one).

Here is results for sieve2 with different problem size for both implementations.

Table 1: implementation 1

Table 2: implementation 2	2
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cores	10^{9}	10^{10}	$2*10^{10}$	cores	10^{9}	10^{10}	
32	1.345676	13.149903	27.425850	32	1.361698	25.021929	
64	0.684752	6.769014	25.098223	64	0.808290	6.944194	
128	0.419530	3.526987	6.987258	128	0.276708	3.448321	
256	0.078973	1.890770	4.127071	256	0.075725	1.944055	

There is no such huge difference between two implementations, for more cache available purpose, the implementation are used in sieve 3.

1.4 sieve3

To use the idea of cache, blocking is a good method to decrease the cache hit miss rate.

By checking the L2 cache size for tardis, it shows the L2 cache size is **2097152**, I chose **2000000** as final block size.

Here is the cache configuration and command I used to check:

getconf -a — grep LEVEL2_CACHE

LEVEL2_CACHE_SIZE 2097152 LEVEL2_CACHE_LINESIZE 64

The results shows that the performance has been improved almost 4 folds by increase the cache hit.

1.5 result

First, the results proving the correctness.

sieve0:

The total number of prime: 455052511, total time: 27.548977, total node 32 The total number of prime: 455052511, total time: 26.136182, total node 64 The total number of prime: 455052511, total time: 7.099779, total node 128 The total number of prime: 455052511, total time: 3.609159, total node 256 The total number of prime: 455052511, total time: 13.517305, total node 32 The total number of prime: 455052511, total time: 6.924514, total node 64 The total number of prime: 455052511, total time: 3.616797, total node 128 The total number of prime: 455052511, total time: 3.120928, total node 256 sieve2: The total number of prime: 455052511, total time: 13.592637, total node 32 The total number of prime: 455052511, total time: 6.769014, total node 64 The total number of prime: 455052511, total time: 3.418407, total node 128 The total number of prime: 455052511, total time: 1.965367, total node 256 sieve3: The total number of prime: 455052511, total time: 4.767698, total node 32 The total number of prime: 455052511, total time: 2.008644, total node 64 The total number of prime: 455052511, total time: 1.006978, total node 128 The total number of prime: 455052511, total time: 0.537084, total node 256

Table 3: sieve0					
cores	10^{9}	10^{10}	$2*10^{10}$	C	
32	2.837629	27.548977	57.065409	3	
64	2.053819	26.136182	54.029782	6	
128	0.808047	7.099779	14.476197	1	
256	0.326352	3.609159	7.389089	2	
Table 5: sieve2					
cores	10^{9}	10^{10}	$2*10^{10}$		

Table 4: sievel					
cores	10^{9}	10^{10}	$2*10^{10}$		
32	1.178136	13.517305	27.735075		
64	0.802593	6.924514	14.082999		
128	0.426171	3.616797	7.066504		
256	0.114588	3.120928	3.648910		

Table 5: sieve2						
cores	10^{9}	10^{10}	$2*10^{10}$			
32	1.361698	13.592637	27.831473			
64	0.808290	6.769014	26.397712			
128	0.276708	3.418407	7.020489			
256	0.075725	1.965367	4.120300			

Table 6: sieve3					
cores	10^{9}	10^{10}	$2*10^{10}$		
32	0.394069	4.767698	8.060845		
64	0.222897	2.008644	4.057865		
128	0.137524	1.006978	2.136247		
256	0.093506	0.537084	1.114526		

The result shows that, the performance has been improved 8 folds by increasing processes number from 32 - 256, and the performance has been improved 7 folds by optimizations from sieve0 - sieve3, total in almost 56 folds (27.548977 - 0.537084).