

$$P_{ij}(s) \triangleq \sum_{n=0}^{\infty} \underbrace{P_{ij}(n)}_{\text{IP}(X_n=j | X_0=i)} \cdot s^n, \quad \begin{cases} P_{ii}(s) = 1 + F_{ii}(s) P_{ii}(s) \\ P_{ij}(s) = F_{ij}(s) \cdot P_{jj}(s) \end{cases} \quad (\forall i \neq j)$$

$$F_{ij}(s) \triangleq \sum_{n=0}^{\infty} \underbrace{f_{ij}(n)}_{\text{IP}(T_j=n | X_0=i)} \cdot s^n$$

state y

recurrent	$\text{IP}_y(T_y < \infty) = 1$
transient	$\text{IP}_y(T_y < \infty) < 1$

with $\text{IP}_y(T_y < \infty) = F_{yy}(1)$

So: y is recurrent iff $F_{yy}(1) = 1$

iff $\frac{P_{yy}(1) - 1}{P_{yy}(1)} = 1$

iff $P_{yy}(1) = \infty$

iff $\sum_{n=0}^{\infty} P_{yy}(n) = \infty$

(6.2.3)

e.g.: State y is recurrent for Markov chain $\{X_n\}$ iff the expected # of visits to y is ∞ if starts from $X_0=y$.

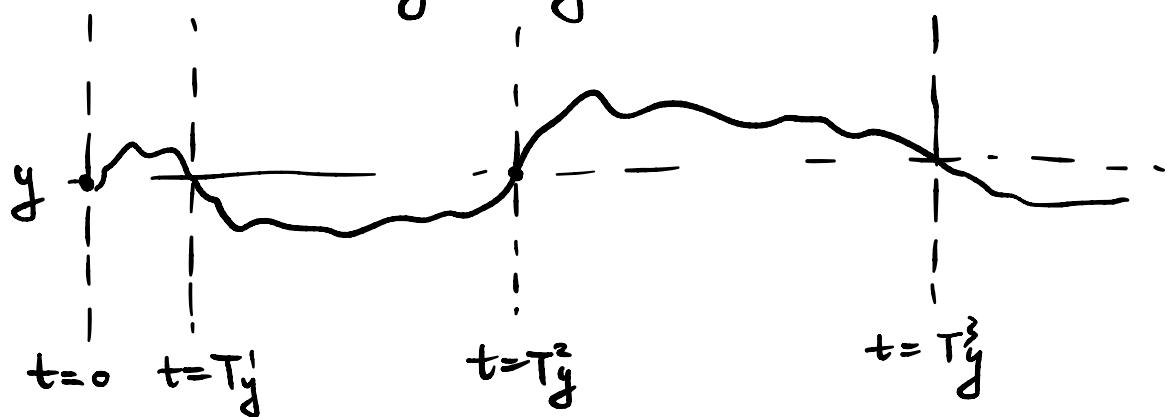
Def: $N_y \triangleq \sum_{n=0}^{\infty} I_{\{X_n=y\}}$ is the # of visits to y

$$E_N = E \sum_{n=0}^{\infty} I_{\{X_n=y\}} \stackrel{\text{Fubini}}{=} \sum_{n=0}^{\infty} P_y(X_n=y) \\ = \sum_{n=0}^{\infty} P_{yy}(n) = \infty$$

iff y is recurrent.

Intuition: if $P_y(T_y < \infty) = 1$, then

restart MC once it hits y , expect to see infinitely many hits.



Understanding recurrent/transient intuitively:

e.g.: If RW $\{S_n\}$ has i.i.d. increments $\{X_n\}$,
 $E|X_1| < \infty$, $EX_1 \neq 0$, then state 0 is transient.

Pf: SLLN: $\frac{S_n}{n} \xrightarrow{\text{a.s.}} EX_1 \quad (n \rightarrow \infty)$

WLOG, assume $EX_1 > 0$, so $S_n \xrightarrow{\text{a.s.}} +\infty \quad (n \rightarrow \infty)$

With prob 1, $\forall K > 0$, $\forall n \geq K$, $S_n \geq 1$.

So: $IP(N_0 \leq K) = 1$, $EN_0 \leq K < \infty$,

State 0 must be transient.

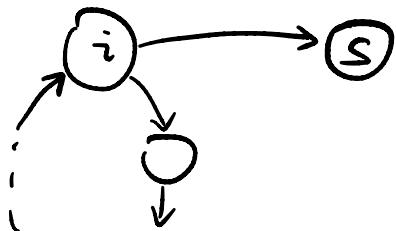
obvious "trend" \rightarrow transience

(e.g. asymmetric SRW, absorbing states, etc.)

e.g. (6.2.2) $\{X_n\}$ has absorbing state s and s communicate with any other states,

$\forall i \in S, \exists n = n(i), P_{is}(n) > 0$. Show that all states other than s are transient.

Pf: $\forall i \in S, n_i \triangleq \min\{n : P_{is}(n) > 0\} < \infty$



Consider two procedures: $\begin{cases} i \rightarrow s \\ i \rightarrow i \end{cases}$,

let $T_i^k \triangleq$ the k -th time Markov chain hits i
(except time 0)

then $(0, T_i^1), (T_i^1, T_i^2), \dots, (T_i^{k-1}, T_i^k), \dots$ are visits to state i

If i is recurrent, by contradiction,

$$P_i(T_i^k < \infty) = P_i(T_i^{k-1} < \infty, T_i^k < \infty)$$

$$= P_i(T_i^{k-1} < \infty) \cdot P_i(T_i^k < \infty | T_i^{k-1} < \infty)$$

Strong Markov

$$\underline{\underline{IP_i(T_{i^+}^{k-1} < \infty) \cdot IP_i(T_i^1 < \infty)}}$$

As a result, $\underline{\underline{IP_i(T_i^k < \infty) = [IP_i(T_i^1 < \infty)]^k}}$.

By recurrence of i , $\underline{\underline{IP_i(T_i^k < \infty) = 1}}$

if a state is recurrent,
with prob 1, the k -th
hitting time is finite for
 $\forall k$



Not only expected # of hitting is ∞

$$E_i N_i = \infty$$

the # of hitting is almost surely ∞

$$IP_i(N_i = \infty) = 1.$$

However, $IP_i(X_{n_i} = s) > 0$ and whenever it hits s , it permanently stays at s , so

$$IP_i(T_i^{n_i} = \infty) \geq IP_i(X_{n_i} = s) > 0,$$

contradiction!

Classification of states:

Irreducible: one communication class

Closed: never transit out of the state subset

{ Communication class share recurrence/transience

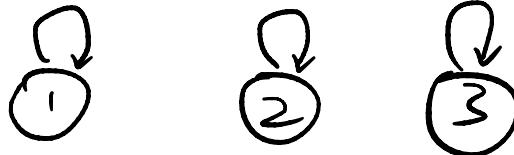
Finite closed set $C \subseteq S$ contains positive recurrent state

recurrence { null recurrent ($E_i T_i = \infty$)
positive recurrent ($E_i T_i < \infty$)

e.g: (6.3.3)

$$P = \begin{pmatrix} 1-2p & 2p & 0 \\ p & 1-2p & p \\ 0 & 2p & 1-2p \end{pmatrix}, p \in [0, \frac{1}{2}]$$

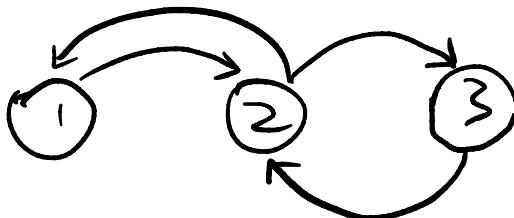
①: $p=0$



all positive recurrent

mean recurrence time = 1

②: $p=\frac{1}{2}$



Irreducible, finite state \Rightarrow all positive recurrent

③: $0 < p < \frac{1}{2}$



all positive recurrent

Calculate mean recurrence time?

first step decomposition

$$IE_1 T_1 = IP_1(X_1=1) \cdot IE_1(T_1 | X_1=1) + IP_1(X_1=2) \cdot IE_1(T_1 | X_1=2)$$

$$\stackrel{\text{Markov}}{=} (1-2p) \cdot 1 + 2p \cdot (1 + IE_2 T_1)$$
$$= 1 + 2p \cdot IE_2 T_1$$

$$IE_2 T_1 = IP_2(X_1=1) \cdot IE_2(T_1 | X_1=1) + IP_2(X_1=2) \cdot IE_2(T_1 | X_1=2)$$
$$+ IP_2(X_1=3) \cdot IE_2(T_1 | X_1=3)$$

$$= p \cdot 1 + (1-2p) \cdot (1 + IE_2 T_1) + p \cdot (1 + IE_3 T_1)$$
$$= 1 + (1-2p) \cdot IE_2 T_1 + p \cdot IE_3 T_1$$

$$IE_3 T_1 = 2p \cdot (1 + IE_2 T_1) + (1-2p) \cdot (1 + IE_3 T_1)$$

$$\Rightarrow \begin{cases} IE_1 T_1 = 1 + 2p IE_2 T_1 \\ 2p IE_2 T_1 = 1 + p IE_3 T_1 \\ 2p IE_3 T_1 = 1 + 2p IE_2 T_1 \end{cases} \Rightarrow \begin{cases} IE_1 T_1 = 4 \\ IE_2 T_1 = \frac{3}{2p} \\ IE_3 T_1 = \frac{2}{p} \end{cases}$$

mean rec
time for
state 1

Easier way to calculate mean rec time?

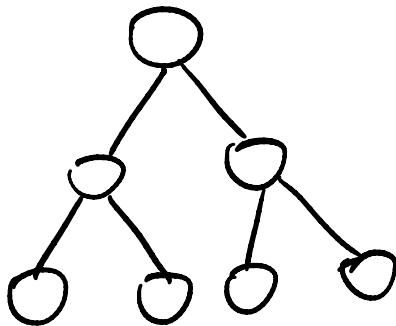
Relies on stationary dist π !

$$E_i T_i = \frac{1}{\pi_i}$$

Remarkable Result

Talk about next week.

e.g.: SRW on binary trees



{ finite tree — all positive recurrent
infinite tree — all transient