

CSCI 570 HOMEWORK 2

Spring 2024

Due: Jan 24 , 11:59PM PST

Q1. What is the tight bound on worst-case runtime performance of the procedure below? Give an explanation for your answer. (10 points)

```
int c = 0;
for(int k = 0; k <= log2n; k++)
    for(int j = 1; j <= 2^k; j++)
        c=c+1
return c
```

Q2. Given an undirected graph G with n nodes and m edges, design an $O(m+n)$ algorithm to detect whether G contains a cycle. Your algorithm should output a cycle if G contains one. (10 points)

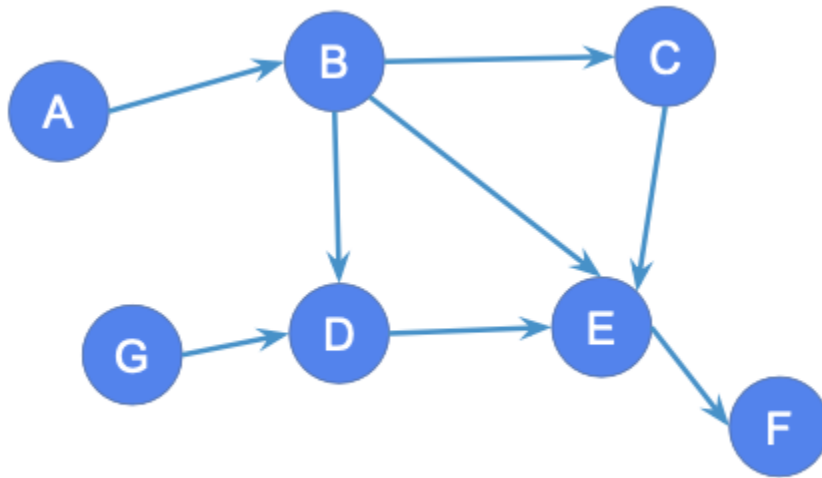
Q3. For each of the following indicate if $f = O(g)$ or $f = \Theta(g)$ or $f = \Omega(g)$ (10 points)

	$f(n)$	$g(n)$
1	$n \log(n)$	$n^2 \log(n^2)$
2	$\log(n)$	$\log(\log(5^n))$
3	$n^{1/3}$	$(\log(n))^3$
4	2^n	2^{3n}
5	$n^4 / \log(n)$	$n(\log(n))^4$

Q4. Indicate for each pair of expressions (A,B) in the table below, whether A is O, Ω , or Θ of B (in other words, whether $A=O(B)$, $A= \Omega(B)$, or $A= \Theta(B)$). Assume that k and C are positive constants. You can mark each box with Yes or No. No justification needed. (9 points)
(Note: log is base 2)

A	B	O	Ω	Θ
$n^3 + \log(n) + n^2$	$C \cdot n^3$			
n^2	$C \cdot n \cdot 2^{\log(n)}$			
$(2^n) \cdot (2^k)$	n^{2k}			

Q5. Find the total number of possible topological orderings in the following graph and list all of them (15 points)



Q6. Given a directed graph with m edges and n nodes where every edge has weight as either 1 or 2, find the shortest path from a given source vertex 's' to a given destination vertex 't'. Expected time complexity is $O(m+n)$. (8 points)

Q7. Given functions f_1, f_2, g_1, g_2 such that $f_1(n) = O(g_1(n))$ and $f_2(n) = O(g_2(n))$. For each of the following statements, decide whether you think it is true or false and give a proof or counterexample. (12 points)

- (a) $f_1(n) \cdot f_2(n) = O(g_1(n) \cdot g_2(n))$
- (b) $f_1(n) + f_2(n) = O(\max(g_1(n), g_2(n)))$
- (c) $f_1(n)^2 = O(g_1(n)^2)$
- (d) $\log_2 f_1(n) = O(\log_2 g_1(n))$

Q8. Design an algorithm which, given a directed graph $G = (V, E)$ and a particular edge $e \in E$, going from node u to node v determines whether G has a cycle containing e . The running time should be bounded by $O(|V| + |E|)$. Explain why your algorithm runs in $O(|V| + |E|)$ time. (8 points)

Q9. Solve Kleinberg and Tardos, **Chapter 3, Exercise 6**. (8 points)

Q10. Solve Kleinberg and Tardos, **Chapter 3, Exercise 9**. (10 points)