

ECS 170 Assignment 1 Part 4

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For the part 4, climbing Mount Saint Helens During the Eruption, the biggest challenge is the efficiency to calculate the optimal path. I used modified Dijkstra's method in part 3 using arrays like Visited, Distance etc.... This approach has poor efficiency. First, the way to calculate the minimum node(minimum $f(x)$) require $O(n)$ because we are searching in the array storing all point's $f(x)$. We used a priority queue with customized compare function and Hash maps to make the cost of searching and traversing cheaper. The searching in hash map only takes $O(1)$ and the priority queue will be sorted when we add Point to the queue but not $O(n)$ approach to find the minimum $f(x)$. Also, the array approach I used calls Point to index and index to Point calculation extensively, we fixed this problem in the new version using hash key.

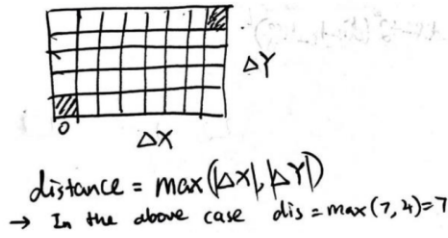
1. Exponential of Height Difference

$$D = \text{Height Difference}(\text{End} - \text{Current})$$

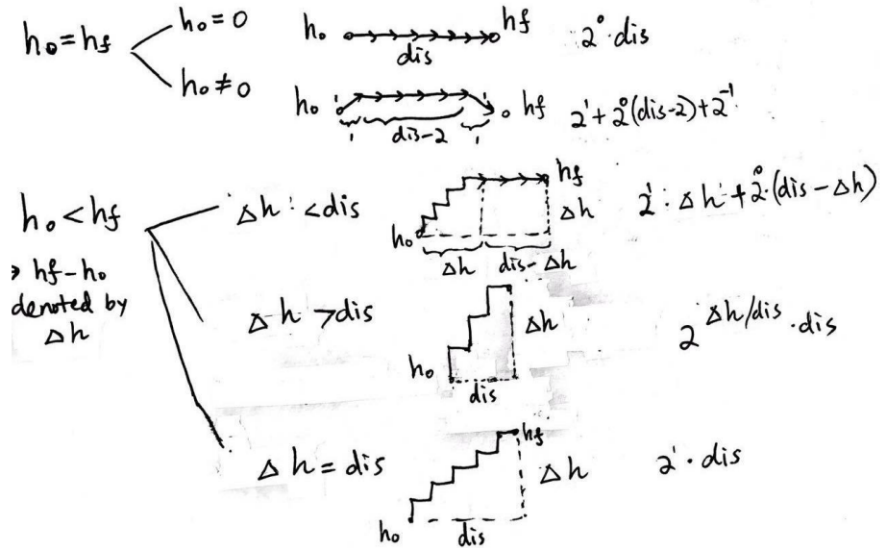
$$S = \text{Minimum Steps to reach the EndPoint}(\text{chebyshev distance})$$

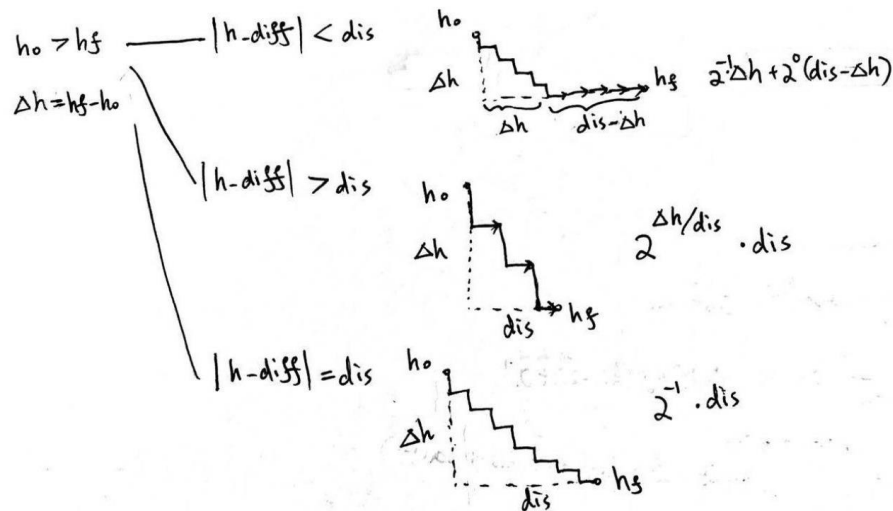
$$h(n) = \begin{cases} 2^D, & S = 1 & (1.1) \\ S, & D = 0 & (1.2) \\ S * 2^{\frac{D}{S}}, & D \neq 0, |D| > S & (1.3) \\ (S - |D|) + 2 * |D|, & |D| < S, D > 0 & (1.4) \\ (S - |D|) + 0.5 * |D|, & |D| < S, D \leq 0 & (1.5) \end{cases}$$

Graph illustration of D:



Graph illustration of Heuristic:





Prove of Consistency:

- (1.1) It is the exact form of the true cost to the goal state when we are one move away from EndPoint.
- (1.2) When the current point and end point are at the same height, we claim the Minimum moves underestimate the true cost from current to EndPoint.

N: Current Node on Optimal Path N': Node successor of N

$$h(N) \leq h(N') + c(N, N')$$

In order to return to same height, there must be

N'' that counteract the change in height

$$S \leq (S - 2) + 2^{\Delta D} + 2^{-\Delta D}$$

Because $2^{\Delta D} + 2^{-\Delta D} \geq 2$, So it is consistent and admissible

(1.3) Let $\frac{D}{S} = a$, $|D| > S$, then we need to prove $2^{aS} \geq S * 2^a$

$$\ln(2^{aS}) \geq a \ln(2S)$$

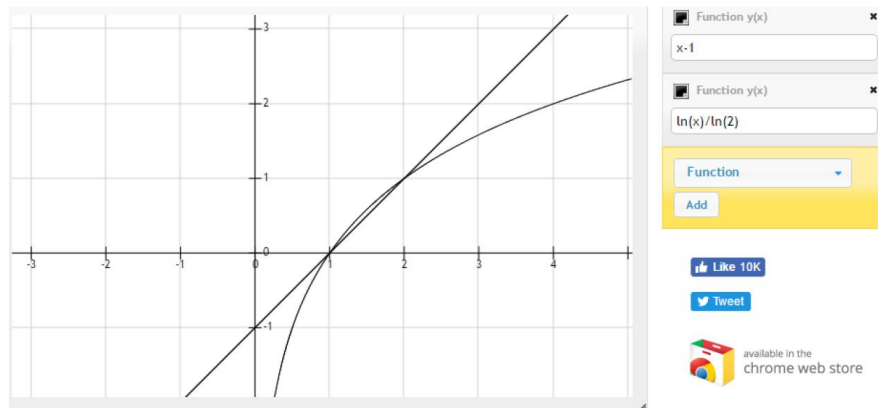
$$aS \ln(2) \geq a \ln(2) + a \ln(S)$$

$$S \geq \frac{\ln(2) + \ln(S)}{\ln(2)}$$

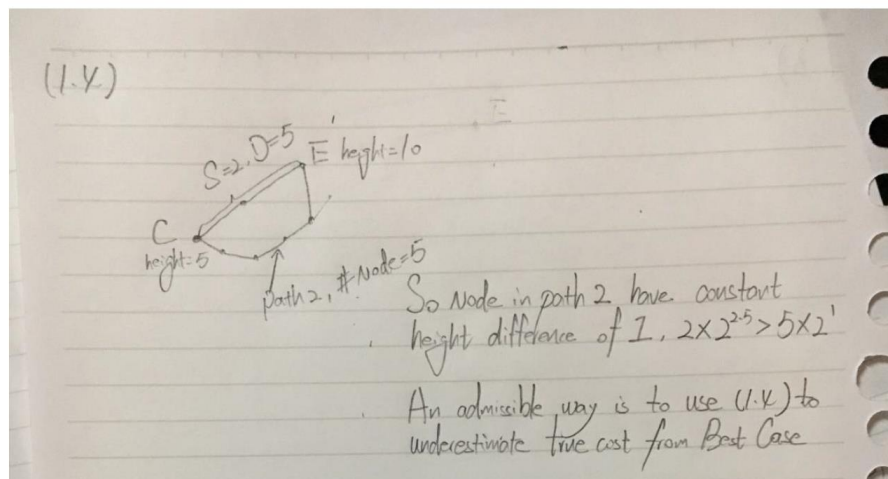
$$S - 1 \geq \frac{\ln(S)}{\ln(2)}$$

From the plot below, S-1 is always larger or equal to RHS when S-1 is Integer value.

So we proved that the (1.3) is admissible to divide height into different steps as much as possible (When $|D|$ is larger than S)



- (1.4) When $|D| < S$, so we cannot move by divide difference equally into integer steps. We give an example of underestimate the true cost.



- (1.5) Similar to (1.4) except we multiply by $2^{-1} = 0.5$