

# \* Sorting Algorithms \*

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## Bubble Sort

Trace for 89 45 68 90 29 34 17

### Iteration 01:

89<sup>?</sup>↔45 68 90 29 34 17

45 89<sup>?</sup>↔68 90 29 34 17

45 68 89<sup>?</sup>↔90<sup>?</sup>↔29 34 17

45 68 89 29 90<sup>?</sup>↔34 17

45 68 89 29 34 90<sup>?</sup>↔17

45 68 89 29 34 17 | 90 → 90 has found its home

### Iteration 02:

45<sup>?</sup>↔68<sup>?</sup>↔89<sup>?</sup>↔29 34 17 90

45 68 29 89<sup>?</sup>↔34 17 90

45 68 29 34 89<sup>?</sup>↔17 90

45 68 29 34 17 | 89 90

### Iteration 03:

45<sup>?</sup>↔68<sup>?</sup>↔~~29~~ 34 17 89 90

45 29 68<sup>?</sup>↔34 17 89 90

45 29 34 68<sup>?</sup>↔17 89 90

45 29 34 17 | 68 89 90

Ituation 04:

45  $\leftrightarrow$  29 34 17 68 89 90  
29 45  $\leftrightarrow$  34 17 68 89 90  
29 34 45  $\leftrightarrow$  17 68 89 90  
29 34 17 | 45 68 89 90

Ituation 05:

29  $\leftrightarrow$  34  $\leftrightarrow$  17 45 68 89 90  
29 17 | 34 45 68 89 90

Ituation 06:

29  $\leftrightarrow$  17 34 45 68 89 90  
17 | 29 34 45 68 89 90

Ituation 07:

| 17 29 34 45 68 89 90



## Algorithm

ALGORITHM Bubble Sort ( $A[0..n-1]$ )

// sorts given array by bubble sort

// Input: An array  $A[0..n-1]$  of orderable  
// elements

// Output: Array  $A[0..n-1]$  sorted in ascending  
// order

for  $i \leftarrow 0$  to  $n-2$  do

  for  $j \leftarrow 0$  to  $n-2-i$  do

    if  $A[j+1] < A[j]$

      Swap  $A[j]$  and  $A[j+1]$

## Basic Operation:

Comparison.

- Is the same, for any kind of input

$$C(n) = \sum_{i=0}^{n-2} \sum_{j=0}^{n-2-i} 1$$

$$= \sum_{i=0}^{n-2} n-2-i-0+1$$

$$= \sum_{i=0}^{n-2} n-i-1$$

$$= \sum_{i=1}^{n-1} n-i - \sum_{i=0}^{n-2} 1$$

$$= n-1 \sum_{i=0}^{n-2} 1 - \sum_{i=0}^{n-2} i$$

$$= (n-1)(n-2-0+1) - \frac{(n-2)(n-1)}{2}$$

$$= (n-1)(n-1) - \frac{(n-1)(n-2)}{2}$$

$$= \frac{2(n-1)(n-1) - (n-1)(n-2)}{2}$$

$$= \frac{(n-1)}{2} [2n-2-n+2]$$

$$= \frac{n-1}{2} \times n$$

$$= \frac{n^2-n}{2} \approx \frac{n^2}{2}$$

$$C(n) \in \Theta(n^2)$$

Note:

- For a decreasing array input, being the worst case, the number of key comparisons & swaps are the same.
- Bubble Sort is in-place

Improvement:

In an iteration pass through, if no swaps were made then the list is sorted. We can ~~stop~~ stop the algorithm.

Variants:

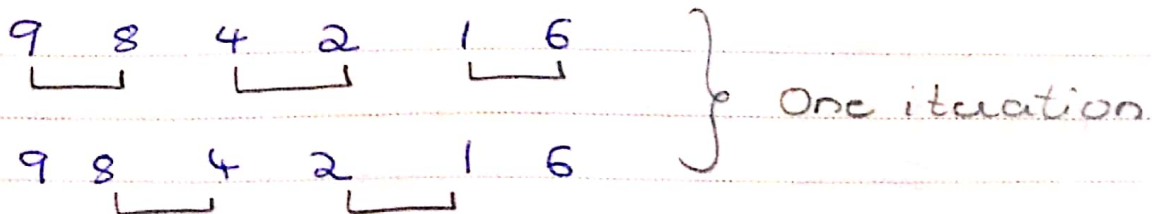
1. Recursive Bubble sort

2. Cocktail Sort

- First stage loops through the array from left to right
- Second stage loops through the array from right to left

3. Odd-Even Sort / Brick Sort

- Each iteration has two phases
- In odd phase, we perform a bubble sort on odd indexed elements & in even phase we perform bubble sort on even indexed elements.





## Selection Sort

Trace for: 89 45 68 90 29 34 17

| 89 45 68 90 29 34 17

17 | 45 68 90 29 34 89

17 29 | 68 90 45 34 89

17 29 34 | 90 45 68 89

17 29 34 45 | 90 68 89

17 29 34 45 68 | 90 89

17 29 34 45 68 89 | 90

### Algorithm:

ALGORITHM SelectionSort ( $A[0 \dots n-1]$ )

// sorts a given array by Selection Sort

// Input: An array  $A[0 \dots n-1]$  of orderable elements

// Output: Array  $A[0 \dots n-1]$  sorted in ascending order

```

for i ← 0 to n-2 do
  min ← i
  for j ← i+1 to n-1 do
    if A[j] < A[min]
      min ← j
  swap A[i] and A[min]

```

Basic Operation:

Comparison.

$$c(n) = \sum_{i=0}^{n-2} \sum_{j=i+1}^{n-1} 1 = \sum_{i=0}^{n-2} (n-1-i-1+1)$$

$$= \sum_{i=0}^{n-2} (n-1-i) \approx \frac{n^2}{2}$$

$$c(n) \in \Theta(n^2)$$

The number of key swaps is  $\Theta(n)$ ,  
 $n-1$  to be precise

Improvement:

Also take maximum on every pass & place it in correct position. In every pass, we keep track of both maximum & minimum & array becomes sorted from both ends.

Variant:

## 1. Stable Selection Sort:

i/p: 4 5 3 2 4 1

o/p: 1 5 3 2 4 4  $\rightarrow$  X

Stable must produce two keys with same value appear in same order in sorted output as they appear in input.

Stable selection sort - instead of Swapping, the minimum element is placed in its position without swapping & by pushing every element one step forward.

## 2. Recursive Selection Sort.



# Insertion Sort

Trace for: 89 45 68 90 29 34 17

```

89 | 45 68 90 29 34 17
45 89 | 68 90 29 34 17
45 68 89 | 90 29 34 17
45 68 89 90 | 29 34 17
29 45 68 89 90 | 34 17
29 34 45 68 89 90 | 17
17 29 34 45 68 89 90
  
```

## Algorithm:

ALGORITHM Insertionsort( $A[0..n-1]$ )

// sorts the given array by insertion sort

// Input: An Array  $A[0..n-1]$  of orderable elements

// Output: An array  $A[0..n-1]$  sorted in increasing order

for  $i \leftarrow 1$  to  $n-1$  do

$U \leftarrow A[i]$

$j \leftarrow i-1$

    while  $j > 0$  and  $A[j] > U$  do

$A[j+1] \leftarrow A[j]$

$j \leftarrow j-1$

$A[j+1] \leftarrow U$

Basic Operation:  $A[j] > v$

Worst Case: Descending order sorted input.

$$C_{\text{worst}}(n) = \sum_{i=1}^{n-1} \sum_{j=0}^{i-1} 1$$

$$= \sum_{i=1}^{n-1} i - 1 - 0 + 1 = \sum_{i=1}^{n-1} i$$

$$= \frac{(n-1)n}{2} \in \Theta(n^2)$$

Best Case: Ascending order sorted input

$$C_{\text{best}}(n) = \sum_{i=1}^{n-1} 1 = n - 1 - 1 + 1$$

$$= n - 1$$

$$\in \Omega(n)$$

Average Case:

Consider an insertion of 3<sup>rd</sup> element, we have 3 combinations with same probability. No. of Comparisons = 1 + 2 + 3

$$\text{Average} = \frac{1+2+3}{3} = 2$$

element can be placed at any position.

Generalizing:

$$\sum_{j=1}^i j = \frac{i(i+1)}{2} = \frac{i+1}{2} \quad \text{no. of comparisons}$$

So,

$$C_{avg}(n) = \sum_{i=1}^{n-1} \frac{i+1}{2} = \frac{1}{2} \sum_{i=1}^{n-1} i+1$$

$$= \frac{1}{2} [2 + 3 + 4 + 5 + \dots + n]$$

$$= \frac{1}{2} \left[ \frac{n(n+1)}{2} - 1 \right]$$

$$= \frac{1}{2} \left[ \frac{n^2}{2} + \frac{n}{2} - \frac{1}{2} \right]$$

$$= \frac{n^2}{4} + \frac{n}{4} - \frac{1}{4}$$

$$\approx n^2$$

$$\in \Theta(n^2)$$



### Improvements & Variants:

1. Finding position of insert with left to right scan or right to left scan
2. Binary insertion sort
3. Recursive Insertion Sort

4.