

## Graph Traversal Algorithms:

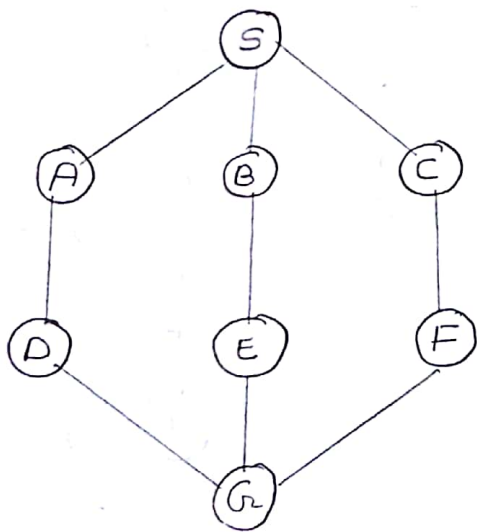
### Depth First Search:

is an algorithm for traversing or searching tree or graph data structures.

One starts at root (selecting some arbitrary node as the root in case of a graph) and explores as far as possible along each branch before backtracking.

Examples: Perform DFS on:

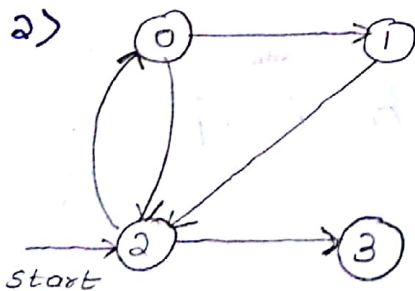
1)



Traversal:

S A D G E B F C

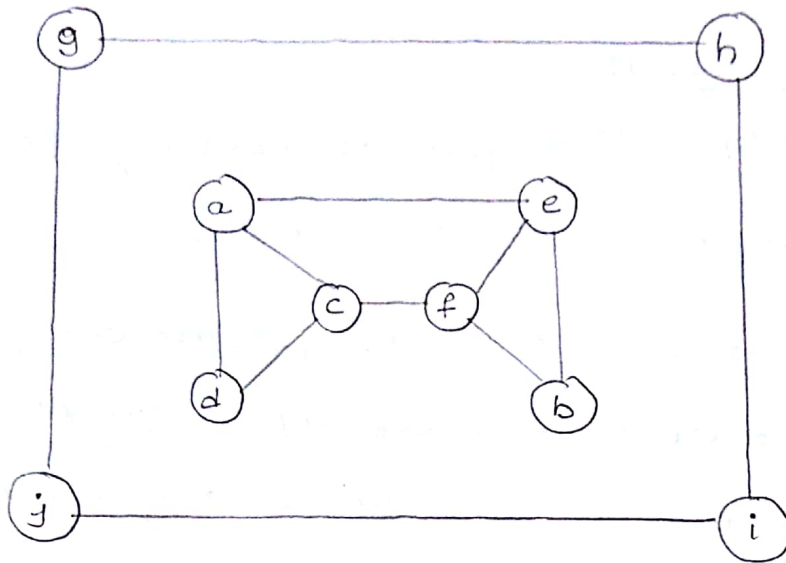
2)



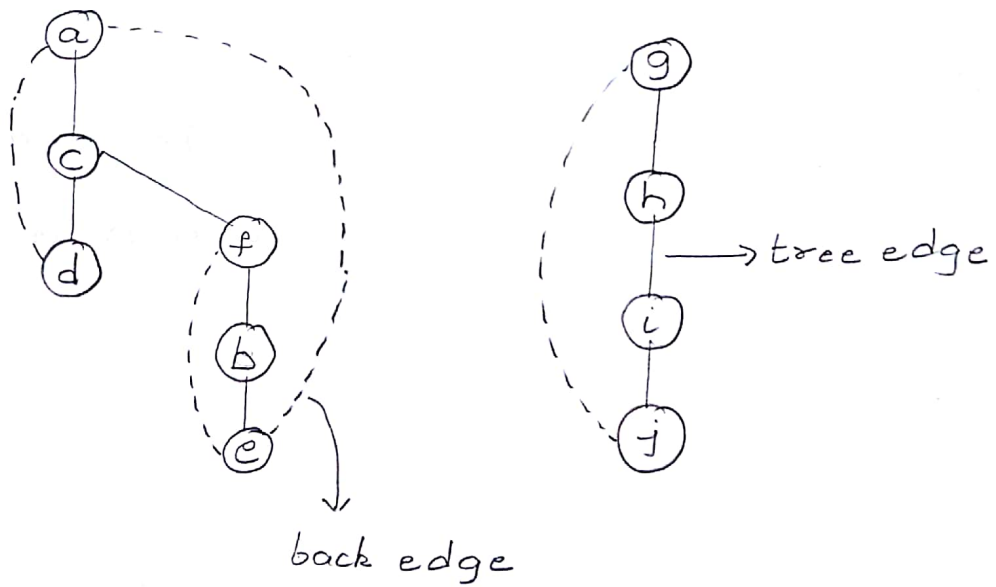
Traversal:

2 0 1 3

3)



DFS forest:



Traversal result:

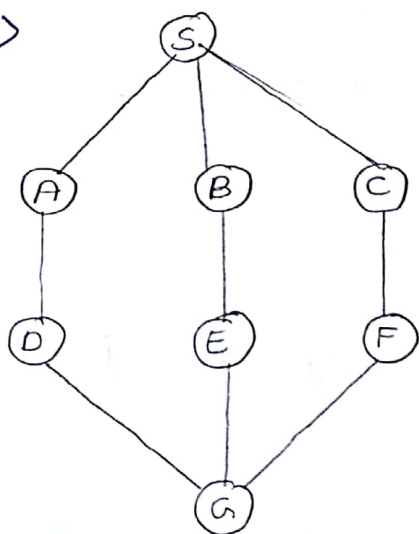
a c d f b e g h i j

## Breadth First Search:

is an algorithm for traversing or searching tree or graph data structures. It starts at the tree root (or some arbitrary node of a graph, sometimes referred to as a 'search key') and explores the neighbor nodes first, before moving to the next level neighbors.

Examples: Perform BFS on:

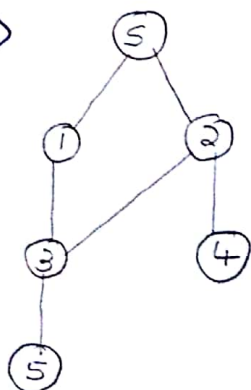
1)



Traversal:

S A B C D E F G

2)

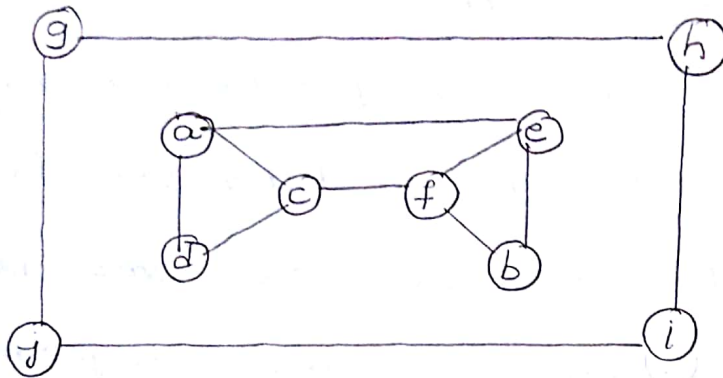
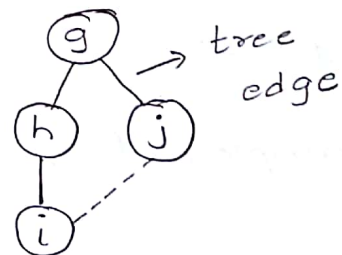
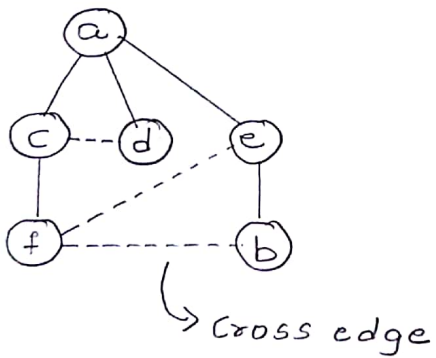


Traversal:

S 1 2 3 4 5

-PH

3)

BFS Forest:Traversal result:

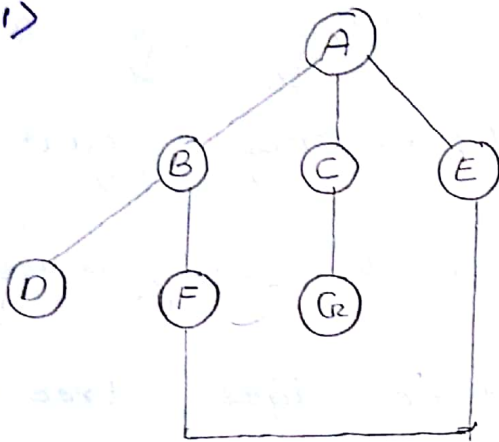
a c d e f b g h j i

## DFS and BFS Comparison:

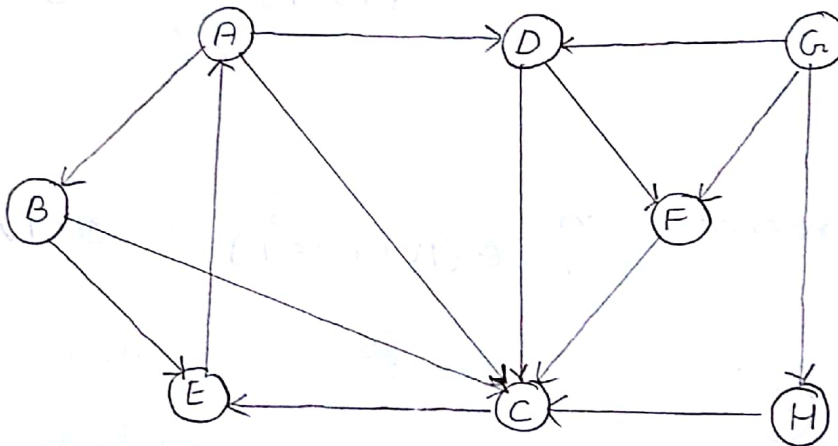
<u>Feature</u>	<u>DFS</u>	<u>BFS</u>
1. For	Brave ones	cautious ones
2. Data Structure	Stack	Queue
3. Edge types	tree edges back edges	tree edges cross edges
4. Matrix representation efficiency	$\Theta(V^2)$	$\Theta(V^2)$
5. List representation efficiency	$\Theta(V + E)$	$\Theta(V + E)$
6. Vertex Orderings number	2 orderings	1 ordering
7. Applications	<ul style="list-style-type: none"> <li>- Detecting cycles</li> <li>- Path finding</li> <li>- Bipartite graph</li> <li>- maze with one solution</li> <li>- etc</li> </ul>	<ul style="list-style-type: none"> <li>- Crawlers</li> <li>- Social n/w websites</li> <li>- GPS Systems</li> <li>- cycle detection</li> <li>- etc</li> </ul>

Exercise: Perform DFS & BFS on:

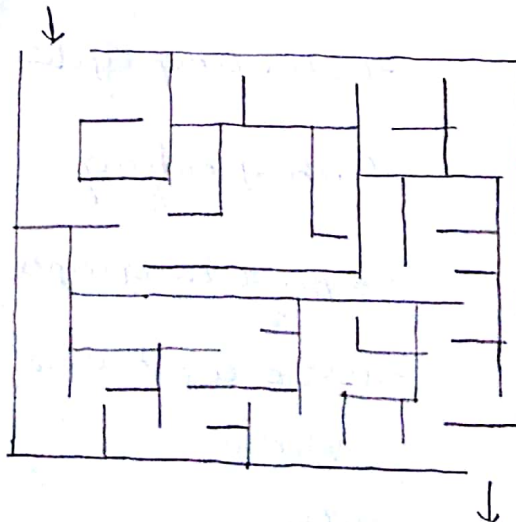
1)



2)



3) Is the maze for brave ones or cautious ones?





## Generic Graph Search:

### Goals:

- a) Find everything reachable from given start vertex
- b) Don't explore anything twice

Generic Algorithm: Given graph  $G$ , vertex  $s$

- initially  $s$ -explored, all other vertices unexplored
  - while possible:
    - choose an edge  $(u, v)$  with  $u$  explored &  $v$  unexplored
    - mark  $v$  explored
- $\left\{ \begin{array}{l} \text{if none,} \\ \text{halt} \end{array} \right.$

### Discovering all the Nodes:

Input Graph  $(s)$

DiscoveredNodes  $\leftarrow \{s\}$

while there is an edge  $e$  leaving DiscoveredNodes that has not been explored:

add vertex at other end of  $e$  to the

DiscoveredNodes

return DiscoveredNodes

### Book-keeping:

- How to keep track of which edges/vertices we have dealt with?
- What order do we explore new edges in?
- What if edges have weights?
- Can we sort the edges?
- Is the graph directed or undirected?
- How do we represent the graph?
- Does the graph have more than one component?
- Is the graph sparse or dense?
- Do we need to keep track of discovered edges or discovered vertices?
- How do we make sure that every node or ~~vis~~ is visited only once?
- What other Supporting Data Structures would be needed for traversal?



## Preorder & Postorder Numbers for DFS:

DFS marks all vertices as visited. We can certainly keep track of other data that can be useful. Design Augment functions to store additional information.

Explore( $v$ )

visited( $v$ )  $\leftarrow$  true

previsit( $v$ )

for ( $v, w$ )  $\in E$ :

if not visited( $w$ ):

    explore( $w$ )

postvisit( $v$ )

what can be previsit( $v$ ) & postvisit( $v$ )?

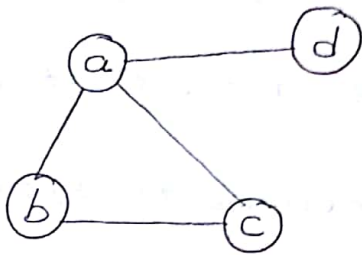
One possible option:

- keep track of order of visits
- record previsit & postvisit times for each  $v$ .

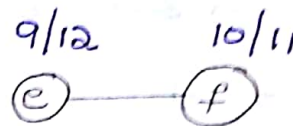
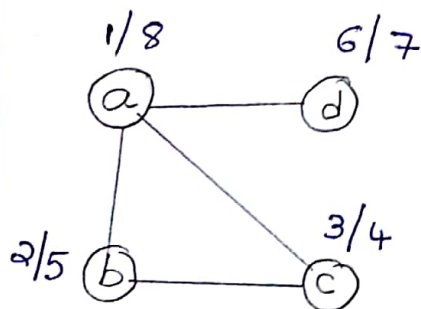
- Its like having sort of clock ticks

We can certainly design the functions any other way as the application demands.

Example:



Ordering:



$tick \leftarrow 1$

$previsit(u)$

$pre(u) \leftarrow tick$

$tick \leftarrow tick + 1$

$postvisit(u)$

$post(u) \leftarrow tick$

$tick \leftarrow tick + 1$

Pre visit and Postvisit numbers tell us about the execution of DFS.

Lemma: For any vertices  $u, v$  the intervals  $[pre(u), post(u)]$  and  $[pre(v), post(v)]$  are either nested or disjoint.

## Nested & Disjoint:

Nested :



Disjoint :



The interleaved case is not possible:



Which of the following tables is not a valid set of pre- & post- orders?

Table 01:

V	Pre	Post
A	1	8
B	9	10
C	3	4
D	2	7
E	5	6

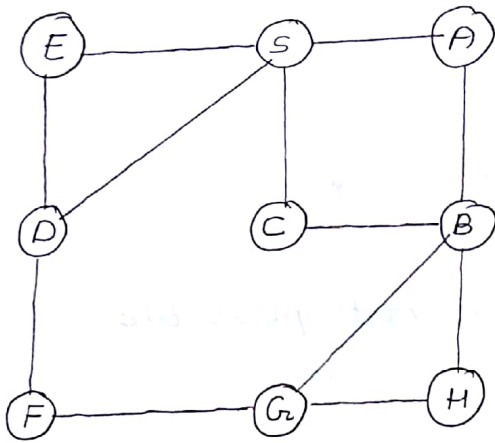
Table 02:

V	Pre	Post
A	1	9
B	8	10
C	2	7
D	3	6
E	4	5

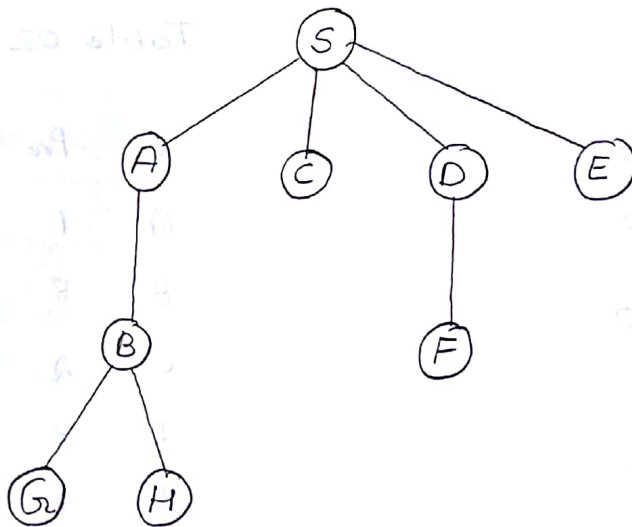
Answer: Table 02.

Can you justify why?

Question: Draw a shortest path tree for the given graph



We can apply BFS traversal to obtain the Shortest path tree.



Note how BFS creates a tree with shortest path from given source vertex to every other vertices.