Date..... ques find pivot element in an array after this element order break, it sund means 7 is the pivot element. If we apply Binary Search on above array, we have 2 different elements which break) the rule of Binary Search, i.e., 7&1 start end Let suppose, if our mid comes of 7 then the condition would be, if (arr[mid] > arr[mid tol] return mid:

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If our mid comes on 1, then our
If our mid comes or
condition would so
ViflarrEmid-1] > arrEmid])
it larremia-13.
return mid-1;
3 1187 1717/
xapro triangle
Above are the 2 conditions from which
we can handle the 2 oad elements
which break Binary Search rule.
Now, we have to handle the left
array & the right array.
1 september 19 to 19 19 19 19 15 C 9 VON BURGETOND
if (arr [start] > arr [mid])
end 2 mid-1; Vif start eleme
3 is a mid elem
else 1 1 d 2 - ent, it means
E rotation is at
start = mid +1: left side.
3
(T) 10 1 19 10 1 1 1 1 1 1 1 1 1 1 1 1 1 1
right if start element
array array means rotation
is at right side

```
code:
 Hinclude (iostream)
Lusing namespace std;
 int main ()
    int arr [] 2 £ 3, 4, 5, 6, 7, 1, 2 3;
    int start 20 end 2 n-1; 14
     int mid:
    while (start < 2 end)
       mid 2 (start + end)
       if ((mid+1) < n de avr Emid:
                    > arr [mid +1
       cout << avv Emid J;
        abreak; of months
      if ((mid-1) > 20 && arr [mid]
          Karr [mid-1])
         cout << arr [mid] -1;
         break!
      if (arr [start] )= arr [mid])
        end 2 mld -1,
         start = mld +
```

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------	---	-----	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--	--

aues search in rotated & sortedo.

array.

in this problem, we have a sorted array,

0 1 2 3 1 4 5 6 7

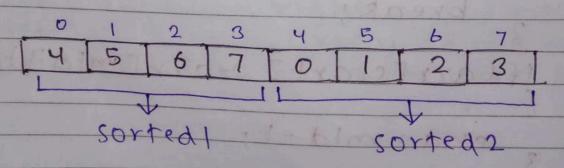
Now this array is rotated,

 -4
 5
 6
 7
 0
 1
 1
 2
 3

this rotated sorted array

In the problem, we have given a rotated sorted array & we have to find the given key in the array.

The condition of Binary Search is that it applies on sorted array only, in our above array we are getting 2 sorted arrays, i.e.,



we've find out our pivot element in the previous problem. we can use that approach in this problem as well.

2

9

3

3

>

3

5

5

0

•

2

3

Approach: If we find our pivot element using the previous approach then we can check that, if our given key is present in the sorted array, then apply Binary search on sorted 1 array. If our given key is present in sorted 2 array, then apply Binary search on sorted 2 array, then apply Binary search on sorted 2 array.

0 1 2 3 4 5 6 7 4 5 6 7 0 1 2 3 1 Pivot index, i.e., 3

if (key > = arr[0] && key <= arr[piv])

is present in sorted I array, then we apply Binary Search on sorted I array & search the key.

1	ate							
~	uu	••	• •	•••	• • •		8	•

if (key) = arrEpiv + 1] && key <=
arrEn-1])

if this condition is true,
it means key is present
in sorted 2 array, then
we apply Binary Search
on Sorted 2 array & search
the key.

code:tol [Tryp tous tous tous tous tous # include (iostream) using namespace std; int binary search (int arrE], int n, int target, int start, int end) int mid; while (start <= end) { mid = (start + end) 12; if (target == arr[mid]) { return mid; else if (target < arr [mid]) { end = mid - 1; else & sold start = mid + 1: retyrn -1: Binary Search is applied on the parts i. of array to search the given target element.

```
int findPivot (int arr[], int n)
  int start = 0, end = n-1, mid;
 while (start < end) {
  mid = (start + end) 12;
  if ((mid + 1) < n && arr[mid] >
                 3([It pim] ryp
         return mid;
         break;
     if((mid-1))= 0 && arr[mid] (
          HIM arrEmid-13){
          return mid -1;
        break;
     if (arr[start] > arr[mid]) {
         end = mid - 1;
        start = mid:
  return start;
        this function will find the
        index of pivot element &
        return that index to search ()
```

function.

Date..... int search (int arr [], int n, int key) - int pivot Element = find Pivot (arr, n); if (target > = arr[0] && target <= arr[pivot Element]){ 11 search in I sorted array int ans = binary search larr, n, target, O, pivot Eleme -nt); return ans; if (pivot Element + 1 < n && target >= arr [pivot Element + 1] && target de arren-1]/2 11 search in 2 sorted array int ans = binary search (arr, n, target, pivotElement +1, n-1); return ans; return -1; > search () will first find the index of our pivot element, the according to the pivot element we check that if our target lies between oth element to pivot elem

right part.

7

2

-ent then we apply Binary Search on left part else apply Binary search on

int main () the town town int arr [] = {5, 6, 7, 8, 9, 10, 1, 2, 33; intende 9; fint target = 3; int targetIndex 2 search (arr, n, target); if (targetIndex == -1) { Here, we call cout << "not exist"; search() to else {

cout << target << 'is at ' <<

target Index of

target Index << 'index'

} return o:

Ques find the square root of a son

If we want to find the square root of 100 then it means it lies between 0 to 100.

Similarly, if n = 50, then it means square root of 50 lies between

0 10 50

Approach: If a square root to so lies between 0 to 50 then we say that our start = 0 & end = 50. we simply calculate mid, i.e., (0 + 50) / 2 = 25.

Now, we check the square of mid, means 25 x 25 = 625.

No, we check if saware of mid is equal to n, then return mid.

If saware of mid is less than n,

then we store the answer first, i.e., store the mid & check at the right part of the array, that is there any

nearest number present to n.
But if square is greater than n
then we simply check at left

part, i.e., end = mid -1.

0 -> 50 -> search space

Date.....

code :a to took groups ant brill int squareRoot (int n) int start = 0, end = n, mid; int ans; on one of magnified while (start <= end) { mid = (start + end) 12; int square = mid * mid; if (sauare == n) 3 return mid; else if (square < n) ans = mid; start = mid + 1; Dim else to ment of 1811s 3 end = mid -1; return ans; From the above code we just only find out the integer part of the square root. Now we find the

floating part also.

```
code! - (For Floating numbers also)
    double squareRoot (int n)
  int start = 0, end = n, mid;
       int ans:
       while (start <= end)
          mid = (start + end) 12:
     int square = mid * mid;
if (square == n) &
          3 11946be tiol to 21 6113
        else if (square 4 n) {
          ans = mid:
       vote start = mid + 1;
     else E
      - bry end = mid - 1;
9
        int precision;
cout << "enter precision: ";
         cin >> precision;
         double step = O.1;
         double final Ans = ans;
      for (int i=0; i < precision; i+t) {
      for (double je final Ans;
      j*j <= n; j += step) 
              final Ans = j;
                         Teacher's Sign .....
```

Date.....

Ques Apply binary search on 20 9

100					rows(n)=5
1	416	2	3	4	column (m) = 4
	5	В	7	8	acro twin h
	49	10	11	12) of Nu
100	13	14	15	16	The triff section
	17	18	19	20	570

In linear array, start = 0 & end = n-1.

same in Matrix case, start is at 0

& end is at last address.

In 2D array, last address is (n x m)-1.

111111111	to 1 1/4		400		ex Start 20
start->	1	2	3	4	end = (5x4)-1
pile of Albert	5	6	7	8	end = 20 -1
	9	710	11	12	* end = 19
mid	13	14	15	16	
- Walter of	17	18	19	20	tend :

mid = (start + end) 12 mid = (0 + 19) /2 mid = 19/2 * mid = 9

> start = 0 end = 19 mid = 9

In 1D array, we directly print the mid element like arr Emid].

In 2D array, how do we print mid element, in 2D array we need to give row no. & column no. to print I that element, but we have only mid. so, now do we print mid element? Stuamala pina da

we have a formula that will find that at which index our s element is stored in the memory. when we have row no. (i) & column

no. (j), the formula to find thework

memory location is,

9

9

9

9

This is the formula to find memory location using row & column.

But to find row index & column index from a given memory location is.

row index = mid / columns

column index = mid % columns

	11 101	101	39 Y		
Starte	TE	2	3/	49	til trigment lit
0)((1)	5	6	7	8	
winny to	9	700	110	512	IN ID array, no.
axist des	13	14	15	16	element, in 20
twive of	17	18	19	20	-> end wor guio
UINO	29V	DO	944	40	V, transpla Tont

How to know row no. & column no. of mid element?

* row no. = mid = 9 = 2 columns 4

« column no. 2 mid % columns 2 9 % 4

So, row no. & column no. of middle element is 2,

2 rivulas 1 bim = xobri cubs

DAMNIBA de KINY = xabyi NMU/00

```
ocode:-
   Hindlyde (iostream)
   using namespace std;
   bool binary search (int arr [][4],
            int rows, int columns,
                          int target
      int start = 0; Manny das
      int end = (rows * columns)-1;
      int mid;
      while (start (= end) {
         mid = (start + end) /2;
         int rowlndex = mid / columns;
          int columnindex = mid %.
                             columns:
          if (target == arr [rowIndex]
                   [column index]
                    De Emplifool Fri
            cout << " element found at"
            << rowlindex << ", " <<
               columnindex;
            return true:
            if (target < arr [rowlnder])

[columninder])
               end = mid -1;
```

```
else
             start = mid +1;
  200 seturn false;
int main ()
                    £10, 15, 20, 253
   int arr[5][4] 2 {
   61 (bus + +00+2 {30,35, 40, 453
 11100 \ 611 \ 60,653
 1 bin x boule 20,75,80,85
                    590,92,96,983
 int rows 25;
 int columns 2 4;
 int target;
cout << "enter target: ";
   cin>> target;
   bool ans 2 binary search (arr,
               rows, columns, target);
if (ans) {

cout << "found";
else {

cout << 'not found';

Spiral 3
                          Teacher's Sign .....
```