Edge E nodes

non negative length

Graph Primitives

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Dijkstra's Algorithm: The Basics

Design and Analysis of Algorithms I

Single-Source Shortest Paths

Input: directed graph $G=(V, E)$. ($m=|E|$, $n=|V|$)

- each edge has non negative length l_e
- source vertex s

Length of path
= sum of edge lengths

Output: for each $v \in V$, compute
  $L(v)$:= length of a shortest s-v path in G

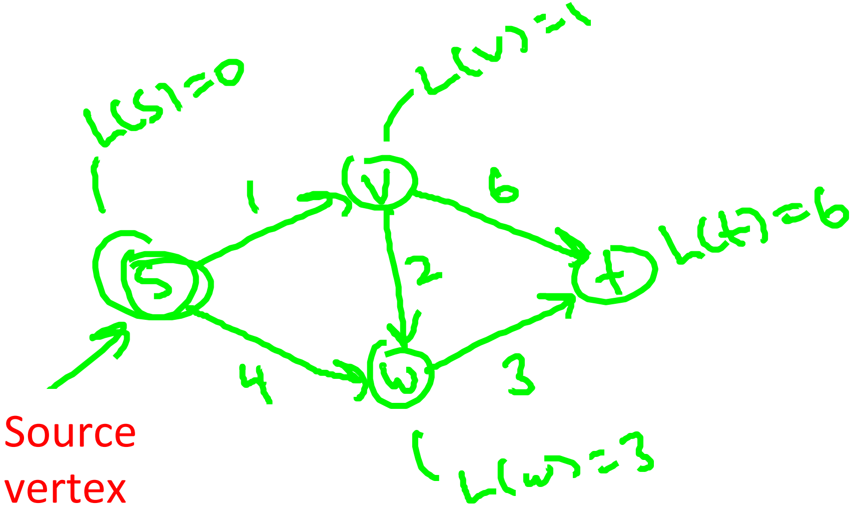
Path length = 6

Assumption:

1. [for convenience] $\forall v \in V, \exists s \Rightarrow v$ path
2. [important] $l_e \geq 0 \;\; \forall e \in E$

One of the following is the list of shortest-path distances for the nodes s,v,w,t, respectively. Which is it?

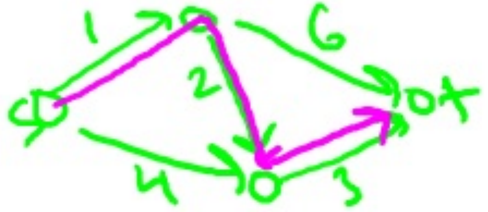
- 0,1,2,3
- 0,1,4,7
- 0,1,4,6
- 0,1,3,6



Why Another Shortest-Path Algorithm?

Question: doesn't BFS already compute shortest paths in linear time?

Answer: yes, IF $l_e = 1$ for every edge e.



Question: why not just replace each edge e by directed path of l_e unit length edges:

Answer: blows up graph too much

Solution: Dijkstra's shortest path algorithm.

Dijkstra's Algorithm

This array only to help explanation!

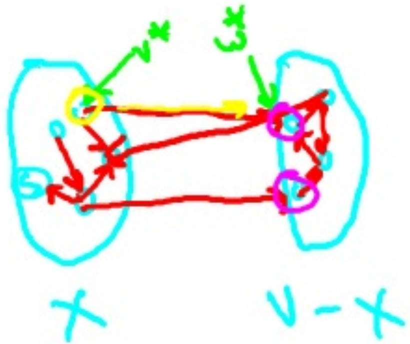
Initialize:

- X = [s]   [vertices processed so far]
- A[s] = 0 [computed shortest path distances]
- B[s] = empty path [computed shortest paths]

Main Loop

- while $X \neq V$:

-need to grow x by one node



Main Loop cont'd:

- among all edges $(v, w) \in E$ with $v \in X, w \notin X$, pick the one that minimizes $A[v] + l_{vw}$

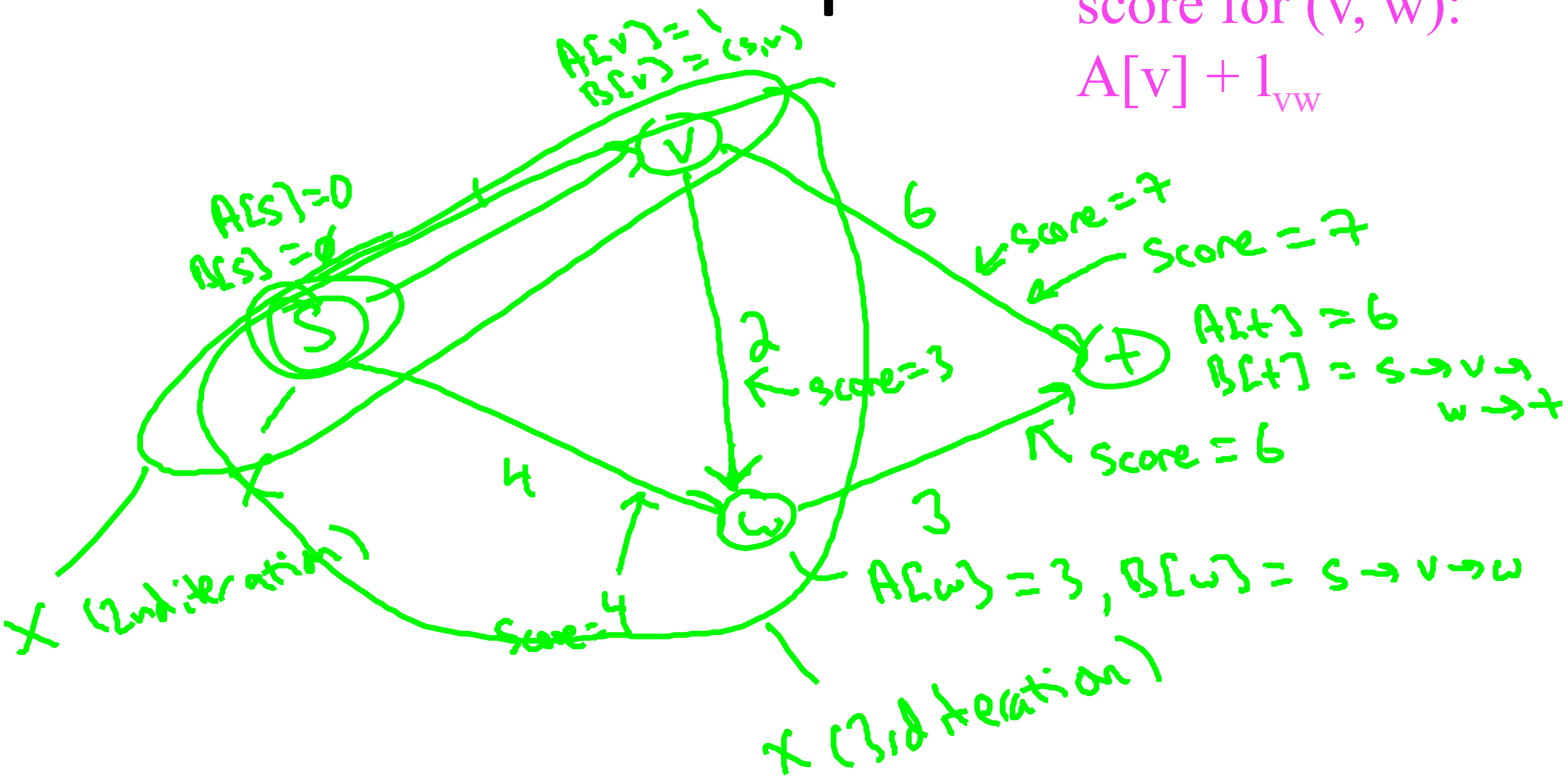
[call it (v*, w*)]

Already computed in earlier iteration

- add w* to X
- set $A[w^*] := A[v^*] + l_{v^*w^*}$
- set $B[w^*] := B[v^*]u(v^*, w^*)$

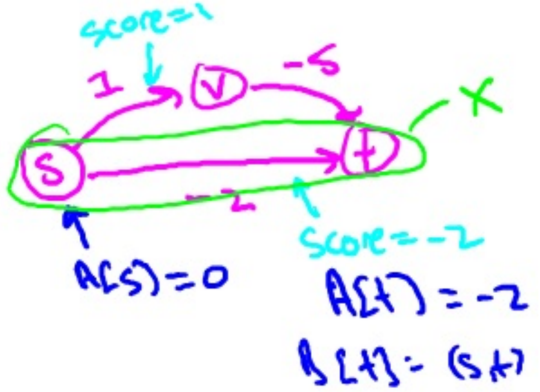
Example

Dijkstra's greedy score for (v, w): $A[v] + l_{vw}$



Non-Example

Question: why not reduce computing shortest paths with negative edge lengths to the same problem with non negative lengths? (by adding large constant to edge lengths)



Problem: doesn't preserve shortest paths !

Also: Dijkstra's algorithm incorrect on this graph !
(computes shortest s-t distance to be -2 rather than -4)