

# Multiple Testing Project: Multiple Testing under independence assumptions

**MT under independence assumptions 2.1 Show that**  $P_H[\cap_{i=1}^m \{y_i > \alpha\}] = (1 - \alpha)^m$

As  $y_1, \dots, y_m$  are independent uniform random variables, the joint distribution of the probability that for every  $y_i > \alpha$  is nothing but the product of the probabilities,

$$P_H[\cap_{i=1}^m \{y_i > \alpha\}] = \prod_{i=1}^m P(y_i > \alpha)$$

We know from the previous exercises that  $P(y_i < \alpha) = \alpha$ , or, equivalently,  $P(y_i > \alpha) = 1 - \alpha$ . Therefore,

$$\prod_{i=1}^m P(y_i > \alpha) = (1 - \alpha)^m$$

**2.2 Show that the probability of rejecting at significance level  $\alpha$  at least one of the independent tests is  $1 - (1 - \alpha)^m$ .**

The probability of rejecting at least one of the tests could be expressed as the union,

$$P_H[\cup_{i=1}^m \{y_i < \alpha\}]$$

And we know that the probability of at least one rejection is the complement of the probability of having no rejections, which we have already determined as being equal to  $(1 - \alpha)^m$ . Thus,

$$P_H[\cup_{i=1}^m \{y_i < \alpha\}] = 1 - P_H[\cap_{i=1}^m \{y_i > \alpha\}]$$

$$P_H[\cup_{i=1}^m \{y_i < \alpha\}] = 1 - (1 - \alpha)^m$$

**2.3 Under the above assumption, show that if we wish that the overall type I error is  $\alpha$ , each independent test should be rejected at significance level  $1 - (1 - \alpha)^{1/m}$**

Type I error is rejecting the null when it is actually true. We know from 1.2 that the probability of type one error for  $m$  hypotheses each having been tested at the confidence level  $\alpha$  is:

$$1 - (1 - \alpha)^m$$

So if we want the overall type I error to be  $\alpha$  for  $m$  tests, each test should be rejected at a certain probability  $p_r$  that satisfies:

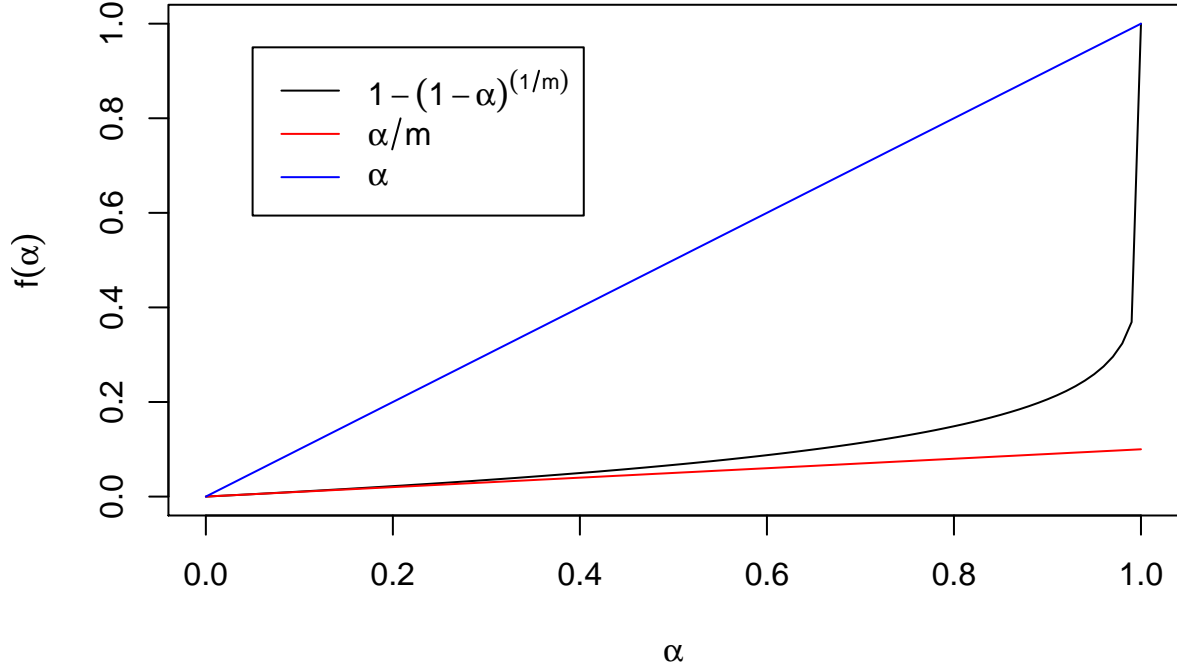
$$P_H[\cup_{i=1}^m \{y_i < p_r\}] = 1 - (1 - p_r)^m = \alpha$$

Solving for  $p_r$  to determine the level  $p_r$  at which each test should be evaluated:

$$1 - (1 - p_r)^m = \alpha$$

$$p_r = 1 - (1 - \alpha)^{\frac{1}{m}}$$

**2.4**  $f_1(\alpha) = 1 - (1 - \alpha)^{1/m}$ ,  $f_2(\alpha) = \alpha/m$ ,  $f_3(\alpha) = \alpha$ . for  $\alpha \in [0, 1]$ . Check that  $f_2 \leq f_1 \leq f_3$ .



## 2.5 (optional)

We can say  $f_1 \leq f_3$  from the following:

$$1 - (1 - \alpha)^{frac{1}{m}} \leq \alpha$$

$$1 - (1 - \alpha)^{frac{1}{m}} - \alpha \leq 0$$

We know the expression  $(1 - \alpha)^{frac{1}{m}}$  when  $m = 1$  because  $m \in [1, \inf]$  and  $1 - \alpha \in [0, 1]$ , and a fraction raised to a fraction is a greater fraction, increasing with the denominator of the exponent.

So the minimum value of  $1 - (1 - \alpha)^{frac{1}{m}} - \alpha$  is given by  $m = 1$  and

$$1 - 1 + \alpha - \alpha = 0 \leq 0, \text{ and } f_1 \leq f_3 \text{ Q.E.D.}$$

We can say  $f_2 \leq f_1$  from the following:

$$\frac{\alpha}{m} \leq 1 - (1 - \alpha)^{\frac{1}{m}}$$

$$0 \leq m - m(1 - \alpha)^{1/m} - \alpha$$

The expression  $m - m(1 - \alpha)^{1/m}$  will always be less than or greater than alpha because as described above, a fraction raised a fraction is a greater fraction so the difference in these two terms  $m$  and  $m - m(1 - \alpha)^{1/m}$  is always greater than  $\alpha$  and the difference increases with  $m$ . The minimum of  $m$  is 1 so for further proof, we can take the minimum to be:

$$0 \leq 1 - (1 - \alpha) - \alpha = 0 \text{ Q.E.D.}$$

## 2.6

In 2.3 it was shown the level needed for each test is  $p_r = 1 - (1 - \alpha)^{\frac{1}{m}}$  in order to achieve a level of  $\alpha$  for multiple tests. The level that makes no correction for multiple testing is assuming  $m = 1$ :

$$\text{No correction: } p_{uncorrected} = 1 - (1 - \alpha)$$

$$\text{Corrected: } p_{corrected} = 1 - (1 - \alpha)^{\frac{1}{m}}$$

We want to show that  $p_{corrected} \leq p_{uncorrected}$

$$1 - (1 - \alpha)^{\frac{1}{m}} \leq 1 - (1 - \alpha)$$

$$-(1 - \alpha)^{\frac{1}{m}} \leq -(1 - \alpha) \text{ Multiplying both sides by -1, we reverse the sign.}$$

$$(1 - \alpha)^{\frac{1}{m}} \geq (1 - \alpha)$$

Since  $m \in [1, \infty]$  and  $\alpha \in [0, 1]$ , this comparison is always true.

**A conservative but robust test: Bonferroni** 3. **Upper-bounding the probability of at least one rejection.** Show that  $\alpha \leq P_{C-H}[\cup_{i=1}^m \{p_i(Y_i) < \alpha\}] \leq m\alpha$

Using Boole's inequality we can say that:

$$P_{C-H}[\cup_{i=1}^m \{p_i(Y_i) < \alpha\}] \leq \sum_{i=1}^m P(p_i(Y_i) < \alpha)$$

From section 1 of the project we know that  $P(P(p_i(Y_i) < \alpha) = \alpha$ . Therefore,  $\sum_{i=1}^m P(p_i(Y_i) < \alpha) = m\alpha$ .

Hence, we have proved the upper bound  $P_{C-H}[\cup_{i=1}^m \{p_i(Y_i) < \alpha\}] \leq m\alpha$ .

The lower bound of  $\alpha$  turns into an equality at  $P_{C-H}[\cup_{i=1}^m \{p_i(Y_i) < \frac{\alpha}{m}\}]$ , since, solving for  $p_{size}$ ,

$$P_{C-H}[\cup_{i=1}^m \{p_i(Y_i) < p_{size}\}] \leq \sum_{i=1}^m P(p_i(Y_i) < p_{size}) = \alpha$$

$$mp_{size} = \alpha$$

$$p_{size} = \frac{\alpha}{m}$$

This means that using a threshold probability of  $\frac{\alpha}{m}$  for rejecting each individual test ensures that the overall type I error equals  $\alpha$ .

Now, if  $\alpha = P_{C-H}[\cup_{i=1}^m \{p_i(Y_i) < \frac{\alpha}{m}\}]$ , for any  $p_{size} \geq \frac{\alpha}{m}$  such as  $p_{size} = \alpha$ ,  $\alpha \leq P_{C-H}[\cup_{i=1}^m \{p_i(Y_i) < p_{size}\}]$ , proving the lower bound.

**Ordered p-values, family-wise error rate and a new MT correction** 4. If we reject test  $i$  when  $p_{(i)} < \frac{i\alpha}{m}$ , the probability of at least one false rejection is:

$$P_H[\cup_{i=1}^m \{y_i < \frac{i\alpha}{m}\}]$$

$$P_H[\cup_{i=1}^m \{y_i < \frac{i\alpha}{m}\}] = 1 - P_H[\cap_{i=1}^m \{y_i > \frac{i\alpha}{m}\}]$$

$$1 - P_H[\cap_{i=1}^m \{y_i > \frac{i\alpha}{m}\}] = 1 - (1 - \alpha)^m = 1 - (1 - \alpha) = \alpha$$