# **Explanations**

#### Part-1

These are timetaken to complete using 1,2,3,4..8 threads. 8.619860 4.393746 3.956057 3.042890 2.971401 3.111116 2.968731 2.706698

## **Execution of parallel\_hashtable.c:**

```
PROBLEMS
              OUTPUT
                      DEBUG CONSOLE
þ
     $ ./parallel_hashtable 1
[main] Inserted 100000 keys in 0.032566 seconds
(\Omega)
```

What circumstances cause an entry to get lost? Analyze the initial code and write a short answer to describe what it means for an entry to be "lost," and which parts of the program are causing this unintended behavior when run with multiple threads.

In the initial code, an entry is considered "lost" when a thread attempts to retrieve a key from the hash table and fails to find it, even though it was previously inserted. This can happen due to a race condition between multiple threads accessing and modifying the hash table concurrently. The primary cause of this unintended behavior is the lack of proper synchronization mechanisms to ensure the atomicity of critical operations, such as inserting and retrieving elements from the hash table.

When multiple threads concurrently execute the 'insert' function without proper synchronization, they might interfere with each other, leading to race conditions. For instance, if one thread is in the process of updating the hash table, another thread might concurrently read from or modify the same location, resulting in the loss of data.

Similarly, during the `retrieve` operation, if one thread is retrieving an entry while another thread is concurrently modifying the hash table, inconsistencies may arise, causing the entry to appear as if it were never inserted.

If I can summarize then the absence of synchronization mechanisms, such as mutexes or locks, allows multiple threads to interfere with each other's access to the hash table, leading to race conditions and, consequently, the loss of entries.

Time Overhead = time taken for parallel hashtable – time taken for parallel mutex

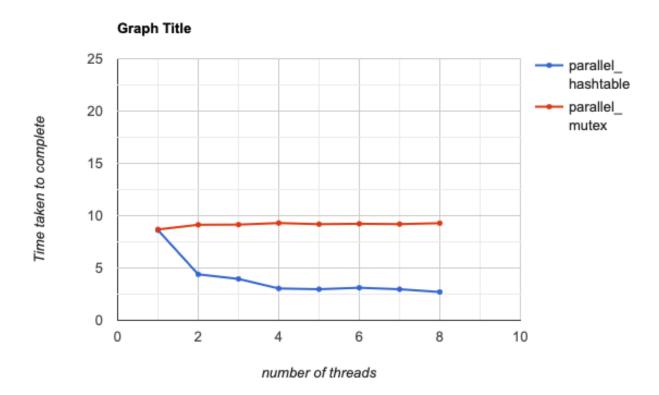
These are timetaken to complete using 1,2,3,4..8 threads. 8.681130 9.131634 9.147779 9.295456 9.189071 9.226427 9.194912 9.283103

- 1. Overhead for 1 thread: 9.131634 8.619860 = 0.511774 seconds
- 2. Overhead for 2 threads: 9.131634 4.393746 = 4.737888 seconds
- 3. Overhead for 3 threads: 9.147779 3.956057 = 5.191722 seconds
- 4. Overhead for 4 threads: 9.295456 3.042890 = 6.252566 seconds
- 5. Overhead for 5 threads: 9.189071 2.971401 = 6.217670 seconds
- 6. Overhead for 6 threads: 9.226427 3.111116 = 6.115311 seconds
- 7. Overhead for 7 threads: 9.194912 2.968731 = 6.226181 seconds
- 8. Overhead for 8 threads: 9.283103 2.706698 = 6.576405 seconds

## **Execution of parallel\_mutex.c:**

```
PROBLEMS
                                                      DEBUG CONSOLE
                                                                                                 TERMINAL
 $ ./parallel_mutex 1
 [main] Inserted 100000 keys in 0.038213 seconds
 [thread 0] 0 keys lost!
[main] Retrieved 100000/100000 keys in 8.681130 seconds
$ ./parallel_mutex 2
[main] Inserted 100000 keys in 0.040001 seconds
 [thread 0] 0 keys lost!
[thread 1] 0 keys lost!
[main] Inserted 100000 keys in 9.131634 seconds
$ ./parallel_mutex 3
[main] Inserted 100000 keys in 0.047560 seconds
[thread 1] 0 keys lost!
[thread 0] 0 keys lost!
[thread 2] 0 keys lost!
[main] Retrieved 100000/100000 keys in 9.147779 seconds
$ ./parallel_mutex 4
[main] Inserted 100000 keys in 0.048247 seconds
[thread 0] 0 keys lost!
[thread 3] 0 keys lost!
[thread 1] 0 keys lost!
[thread 2] 0 keys lost!
 [main] Retrieved 100000/100000 keys in 9.295456 seconds
[main] Retrieved 100000/100000 keys in 9.295456
$ ./parallel_mutex 5
[main] Inserted 1000000 keys in 0.050977 seconds
[thread 4] 0 keys lost!
[thread 2] 0 keys lost!
[thread 3] 0 keys lost!
[thread 0] 0 keys lost!
[thread 1] 0 keys lost!
[main] Retrieved 100000/100000 keys in 9.189071
 [main] Retrieved 100000/100000 keys in 9.189071 seconds
[main] Retrieved 100000/100000 keys in 9.1890/1
$ ./parallel_mutex 6
[main] Inserted 100000 keys in 0.044059 seconds
[thread 1] 0 keys lost!
[thread 0] 0 keys lost!
[thread 2] 0 keys lost!
[thread 5] 0 keys lost!
[thread 5] 0 keys lost!
[thread 4] 0 keys lost!
[thread 4] 0 keys lost!
[main] Retrieved 100000/100000 keys in 9.226427
 [main] Retrieved 100000/100000 keys in 9.226427 seconds
 $ ./parallel_mutex 7
[main] Inserted 100000 keys in 0.045579 seconds
[main] Inserted 100000 keys in 0.0455/9 seconds
[thread 3] 0 keys lost!
[thread 1] 0 keys lost!
[thread 2] 0 keys lost!
[thread 5] 0 keys lost!
[thread 4] 0 keys lost!
[thread 6] 0 keys lost!
$ ./parallel_mutex 8
[main] Inserted 100000 keys in 0.042317 seconds
[main] Inserted 100000 [thread 2] 0 keys lost! [thread 0] 0 keys lost! [thread 1] 0 keys lost! [thread 5] 0 keys lost! [thread 3] 0 keys lost! [thread 4] 0 keys lost! [thread 6] 0 keys lost!
```

# **Graph of time taken for parallel\_hashtable vs parallel\_mutex**



We can see that parallel\_mutex execution has remained almost the same regardless of increase in number of threads.

Whereas the time execuetion in parallel\_hashtable has decreased significantly by increasing the number of threads.

If you were to replace all mutexes with spinlocks, what do you think will happen to the running time? Write a short answer describing what you expect to happen, and why the differences in mutex vs. spinlock implementations lead you to that conclusion

When replacing mutexes with spinlocks, I would expect the running time to decrease compared to using mutexes. Spinlocks are generally considered more lightweight and have lower overhead than mutexes.

Mutexes involve blocking the thread when it tries to acquire a lock, and the thread is only unblocked when the lock is acquired. This blocking and unblocking operation introduces additional overhead, especially in situations with high contention.

On the other hand, spinlocks work by repeatedly "spinning" in a busy-wait loop until the lock is acquired. This approach is more efficient when lock contention is brief because it avoids the context-switching overhead associated with blocking and unblocking threads.

In summary, I expect that using spinlocks will result in reduced running time compared to mutexes, especially in scenarios with short critical sections and low contention but not sure about more critical or introducing more no of threads.

Time Overhead = time taken for parallel\_hashtable – time taken for parallel\_spin

These are timetaken to complete using 1,2,3,4..8 threads. 8.578704 8.665593 10.545755 10.754553 11.163455 16.094853 20.432355 22.897329

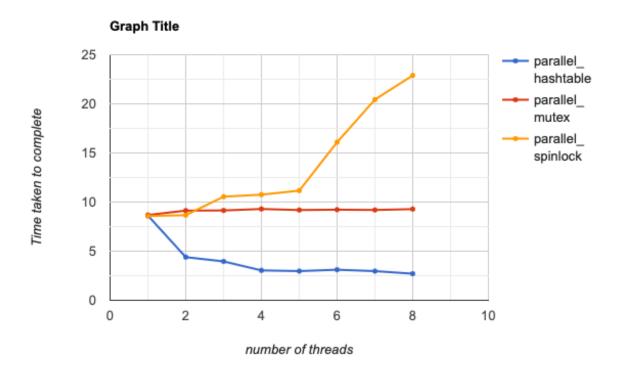
```
8.578704 - 8.614970 = -0.036266 seconds
8.665593 - 8.619860 = 0.045733 seconds
10.545755 - 8.614970 = 1.930785 seconds
10.754553 - 9.131634 = 1.622919 seconds
11.163455 - 9.189071 = 1.974384 seconds
16.094853 - 9.226427 = 6.868426 seconds
20.432355 - 9.194912 = 11.237443 seconds
22.897329 - 9.283103 = 13.614226 seconds
```

Total cumulative overhead = 37.25765

#### **Execution of parallel spin:**

```
PROBLEMS
                                DEBUG CONSOLE
                 OUTPUT
                                                          TERMINAL
                                                                           PORTS
$ ./parallel_spin 1
[main] Inserted 100000 keys in 0.038270 seconds [thread 0] 0 keys lost!
[main] Retrieved 100000/100000 keys in 8.578704 seconds
$ ./parallel_spin 2
[main] Inserted 100000 keys in 0.033884 seconds
[thread 1] 0 keys lost!
[thread 0] 0 keys lost!
[main] Retrieved 100000/100000 keys in 8.665593 seconds
$ ./parallel_spin 3
[main] Inserted 100000 keys in 0.034144 seconds
[thread 1] 0 keys lost!
[thread 2] 0 keys lost!
[thread 0] 0 keys lost!
[main] Retrieved 100000/100000 keys in 10.545755 seconds
$ ./parallel_spin 4
[main] Inserted 100000 keys in 0.033933 seconds
[thread 2] 0 keys lost!
[thread 0] 0 keys lost!
[thread 2] 0 keys lost!
[thread 1] 0 keys lost!
[main] Retrieved 100000/100000 keys in 10.754553 seconds
$ ./parallel_spin 5
[main] Inserted 100000 keys in 0.040241 seconds
[thread 2] 0 keys lost!
[thread 0] 0 keys lost!
[thread 1] 0 keys lost!
[thread 4] 0 keys lost!
[thread 3] 0 keys lost!
[main] Retrieved 100000/100000 keys in 11.163455 seconds
$ ./parallel_spin 6
[main] Inserted 100000 keys in 0.045086 seconds
[thread 4] 0 keys lost!
[thread 3] 0 keys lost!
[thread 5] 0 keys lost!
[thread 0] 0 keys lost!
[thread 1] 0 keys lost!
[thread 2] 0 keys lost!
[main] Retrieved 100000/100000 keys in 16.094853 seconds
$ ./parallel_spin 7
[main] Inserted 100000 keys in 0.034635 seconds
[thread 0] 0 keys lost!
[thread 2] 0 keys lost!
[thread 6] 0 keys lost!
[thread 5] 0 keys lost!
[thread 3] 0 keys lost!
[thread 4] 0 keys lost!
[thread 1] 0 keys lost!
[main] Retrieved 100000/100000 keys in 20.432355 seconds
$ ./parallel_spin 8
[main] Inserted 100000 keys in 0.039997 seconds [thread 2] 0 keys lost! [thread 0] 0 keys lost!
[thread 3] 0 keys lost!
[thread 6] 0 keys lost!
[thread 7] 0 keys lost!
[thread 1] 0 keys lost!
[thread 5] 0 keys lost!
```

# Graph:



We can see that spinlock execution has increased significantly when we have increased the number of threads but has consumed less time with 1 thread when compared to others.

My hypothesis was correct for one thread but not when more number of threads are introduced.

#### Part-4

Let's revisit your mutex-based code. When we retrieve an item from the hash table, do we need a lock? Write a short answer and explain why or why not.

When considering whether to use a lock during the retrieval of items from the hash table, it's crucial to assess the nature of the concurrent operations happening in the program.

In situations where multiple threads are exclusively performing read operations (retrievals) without any concurrent write operations (such as insertions or deletions), omitting a lock during retrieval might seem feasible. This is grounded in the understanding that reading from shared data structures is generally considered safe when there are no concurrent writes.

However, it's essential to recognize that the decision to forego a lock during retrieval is contingent on the absence of any concurrent write operations. If there's a possibility of simultaneous write operations occurring, such as insertions or deletions, it becomes imperative to introduce a lock. This precautionary measure is crucial for ensuring data consistency and preventing race conditions that could lead to inaccurate or partially updated results.

The necessity of a lock during retrieval hinges on the specific requirements of the application and the potential for concurrent write operations. When there's any chance of concurrent writes, opting for a lock becomes a prudent choice to uphold data integrity and prevent unintended inconsistencies.

#### Changes I made:

I just left the insert function same as in the parallel\_mutex.c and have removed mutex locks in retrive function.

#### Execution of parallel mutex opt 1.c:

```
PROBLEMS
                                      DEBUG CONSOLE
                                                                   TERMINAL
                                                                                       PORTS
$ gcc -pthread parallel_mutex_opt_1.c -o parallel_mutex_opt_1
$ ./parallel_mutex_opt_1 1
[main] Inserted 100000 keys in 0.038280 seconds
[thread 0] 0 keys lost!
[main] Retrieved 100000/100000 keys in 8.665619 seconds
$ ./parallel_mutex_opt_1 2
[main] Inserted 100000 keys in 0.038347 seconds
[thread 1] 0 keys lost!
[thread 0] 0 keys lost!
[main] Retrieved 100000/100000 keys in 4.410507 seconds
$ ./parallel_mutex_opt_1 3
[main] Inserted 100000 keys in 0.046224 seconds
[thread 2] 0 keys lost!
[thread 1] 0 keys lost!
[thread 0] 0 keys lost!
[main] Retrieved 100000/100000 keys in 3.430454 seconds
$ ./parallel_mutex_opt_1 4
[main] Inserted 100000 keys in 0.051823 seconds
[thread 3] 0 keys lost!
[thread 1] 0 keys lost!
[thread 0] 0 keys lost!
[thread 2] 0 keys lost!
[main] Retrieved 100000/100000 keys in 2.484234 seconds
$ ./parallel_mutex_opt_1 5
[main] Inserted 100000 keys in 0.048700 seconds
[thread 4] 0 keys lost!
[thread 2] 0 keys lost!
[thread 2] 0 keys lost!
[thread 0] 0 keys lost!
[thread 3] 0 keys lost!
[main] Retrieved 100000/100000 keys in 3.200293 seconds
$ ./parallel_mutex_opt_1 6
[main] Inserted 100000 keys in 0.049645 seconds
[thread 2] 0 keys lost!
[thread 5] 0 keys lost!
[thread 1] 0 keys lost!
[thread 0] 0 keys lost!
[thread 3] 0 keys lost!
[thread 4] 0 keys lost!
[main] Retrieved 100000/100000 keys in 2.695721 seconds
$ ./parallel_mutex_opt_1 7
[main] Inserted 100000 keys in 0.048095 seconds
[thread 6] 0 keys lost!
[thread 0] 0 keys lost!
[thread 2] 0 keys lost!
[thread 3] 0 keys lost!
[thread 1] 0 keys lost!
[thread 4] 0 keys lost!
[thread 0] 0 keys lost!
[main] Retrieved 100000/100000 keys in 2.491166 seconds
$ ./parallel_mutex_opt_1 8
[main] Inserted 100000 keys in 0.056539 seconds
[thread 5] 0 keys lost!
[thread 3] 0 keys lost!
[thread 2] 0 keys lost!
[thread 6] 0 keys lost!
[thread 4] 0 keys lost!
[thread 1] 0 keys lost!
```

#### Part-5

# Last, let's consider insertions. Describe a situation in which multiple insertions could happen safely (hint: what's a bucket?).

In the context of a hash table, a bucket is a slot or container that holds key-value pairs. Each bucket corresponds to a specific hash value, and multiple keys can map to the same bucket due to hash collisions. The objective is to design the hash table in a way that allows multiple insertions to different buckets to happen safely and concurrently.

A situation in which multiple insertions could happen safely involves ensuring that each insertion operation targets a distinct bucket. Since each bucket is independent and doesn't share data with other buckets, concurrent insertions to different buckets won't result in data races or conflicts.

To achieve this, the hash function used to determine the bucket index for a given key should be designed such that different keys are more likely to hash to different buckets. This reduces the likelihood of collisions and ensures that threads inserting different keys can safely operate on their respective buckets without interfering with each other.

In summary, for safe concurrent insertions, the key is to design the hash function to distribute keys evenly across buckets, minimizing the chance of collisions and allowing threads to independently insert into different buckets without the need for locks or synchronization.

Multiple insertions can happen safely in parallel when they target different buckets in the hash table. In a hash table, a bucket is a container that holds key-value pairs. Each bucket is identified by a unique index, typically determined by the hash value of the keys.

Since different keys can hash to different indices, threads inserting key-value pairs with different keys (hashing to different buckets) can safely operate concurrently without conflicting with each other. This is because they are modifying different parts of the hash table (different buckets), and thus, there is no contention for the same memory location.

In other words, as long as the threads are inserting key-value pairs into distinct buckets, there won't be any data dependencies or conflicts, and the insertions can happen safely in parallel. This scenario allows for better utilization of multiple threads, promoting parallelism and potentially improving the overall performance of the hash table operations.

#### **Explanations on changes I made:**

The modifications made to the code involve introducing individual mutexes for each bucket in the hash table, thereby replacing the global mutex. Specifically, for each bucket, a mutex is initialized, and before performing any operations on the bucket (insertion or retrieval), the corresponding mutex is locked. After the operation is completed, the mutex is unlocked.

Here are the key changes made to the code:

#### 1. Mutex Initialization for Each Bucket:

For each bucket in the hash table, a separate mutex is initialized. This is achieved by introducing an array of mutexes.

```
pthread_mutex_t mutex[NUM_BUCKETS];

Initialization is done in the main function:

for (i = 0; i < NUM_BUCKETS; i++) {
   if (pthread_mutex_init(&mutex[i], NULL) != 0) {
      perror("Mutex initialization failed");
      exit(1);
   }
}</pre>
```

# 2. Locking and Unlocking Mutexes in Insert and Retrieve Functions:

The critical sections in the 'insert' and 'retrieve' functions, where operations are performed on the hash table buckets, are protected by locking and unlocking the corresponding mutex.

```
// Lock the mutex before performing operations on the bucket
pthread_mutex_lock(&mutex[i]);
// ... Perform operations on the bucket ...
// Unlock the mutex after the operations are completed
pthread mutex unlock(&mutex[i]);
```

#### 3. Mutex Destruction:

After all the operations are completed, the mutexes are destroyed in the main function.

```
for (i = 0; i < NUM_BUCKETS; i++) {
   pthread_mutex_destroy(&mutex[i]);
}</pre>
```

These changes address the issues related to concurrent access to the hash table buckets. Each bucket is now protected by its own mutex, ensuring that insertions and retrievals are performed atomically and preventing race conditions that could lead to data corruption. The use of individual mutexes for each bucket enhances parallelism by allowing multiple threads to operate on different buckets simultaneously, reducing contention compared to a single global mutex.

## **Execution of parallel\_mutex\_opt.c:**

```
PROBLEMS
                           OUTPUT
                                                DEBUG CONSOLE
                                                                                        TERMINAL
$ gcc -pthread parallel_mutex_opt.c -o parallel_mutex_opt
$ ./parallel_mutex_opt 1
[main] Inserted 100000 keys in 0.038962 seconds
 [thread 0] 0 keys lost!
 [main] Retrieved 100000/100000 keys in 8.644975 seconds
final() Retrieved 100000 2
s./parallel_mutex_opt 2
[main] Inserted 100000 keys in 0.035725 seconds
[thread 1] 0 keys lost!
[thread 0] 0 keys lost!
[main] Retrieved 100000/100000 keys in 4.624151 seconds
$ ./parallel_mutex_opt 3
[main] Inserted 100000 keys in 0.036477 seconds
[thread 0] 0 keys lost!
[thread 2] 0 keys lost!
[thread 2] 0 keys lost!
[main] Retrieved 100000/100000 keys in 3.043021 seconds
$ ./parallel_mutex_opt 4
[main] Inserted 100000 keys in 0.035476 seconds
[thread 3] 0 keys lost!
[thread 2] 0 keys lost!
[thread 1] 0 keys lost!
[thread 0] 0 keys lost!
[main] Retrieved 100000/100000 keys in 3.300220 seconds
$ ./parallel_mutex_opt 5
[main] Inserted 100000 keys in 0.041079 seconds
[main] Inserted 100000 keys in 0.0410/9 seconds
[thread 2] 0 keys lost!
[thread 3] 0 keys lost!
[thread 1] 0 keys lost!
[thread 0] 0 keys lost!
[thread 0] 0 keys lost!
[main] Retrieved 100000/100000 keys in 3.794860 seconds

* (narallel mutey ont 6)
$ ./parallel_mutex_opt 6
[main] Inserted 100000 keys in 0.038144 seconds
[thread 5] 0 keys lost!
[thread 4] 0 keys lost!
[thread 0] 0 keys lost!
[thread 1] 0 keys lost!
[thread 3] 0 keys lost!
[thread 2] 0 keys lost!
 [main] Retrieved 100000/100000 keys in 2.487802 seconds
[main] Inserted 100000 keys in 0.039894 seconds
[main] Inserted 100000 keys in 0.039894 seconds
[thread 5] 0 keys lost!
[thread 6] 0 keys lost!
[thread 0] 0 keys lost!
[thread 3] 0 keys lost!
[thread 4] 0 keys lost!
[thread 2] 0 keys lost!
[thread 1] 0 keys lost!
[thread 1] 0 keys lost!
[main] Retrieved 100000/100000 keys in 2.598037 seconds
$ ./parallel_mutex_opt 8
[main] Inserted 1000000 keys in 0.038463 seconds
[thread 0] 0 keys lost!
[main] Inserted 100000 [thread 0] 0 keys lost! [thread 7] 0 keys lost! [thread 3] 0 keys lost! [thread 6] 0 keys lost! [thread 4] 0 keys lost! [thread 1] 0 keys lost!
```