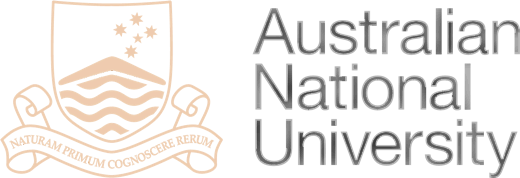
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***Overview***

In this report, it shows our answer of Part 1 fundamentals and the report part of Part 2 Reaction Timer. The implementation of Part 2 is using the Digilent Nexys4 DDR. All the codes are written via Vivado 2017.4. The version we handled are fully tested and it could be synthesis, implemented and generating bitstream with no bugs.

This assignment gets special page extension from Lyle, check it on Appendix.

Acknowledgement for the following people who gave us a help during the assignment:

Tutors:

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* Lyle Roberts
* Jonghyuk Kim

Laboratory:

* Xianjun Zheng

***Assignment One***

***Fundamentals & Reaction Timer***

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# Part 1: Fundamentals

1. Block statements must be executed in the order as they managed in the block. In contrast, non-blocking statements allows executing concurrently within the same time stamp. The order of non-blocking statements won’t affect the result.

For example, a 2-bit shift register could be implemented with 2 D-type flip-flops. If the shift register implemented with

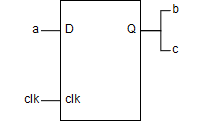
always @(posedge clk) begin

b = a;

c = b;

end

The implementation of this code should be



which makes the register flip-flops (b and c) have exactly the same value. This is not what we want. This is the blocking statements. However, if it is implemented with non-blocking assignment

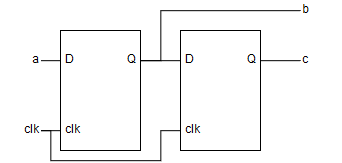
always @(posedge clk) begin

b <= a;

c <= b;

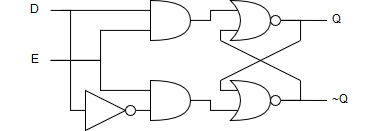
end

The implementation should be



And this is exactly what a shift register works as. Because when the code is running, b still keeps its previous value thanks for the clock. When the code is executing, these two statements runs at the same time, so c would get the value of b, and b would be updated to a.

1. D-type flip-flop (D-FF) is developed based on a RS-type flip-flop (RS-FF). The following diagram shows the internal of a D-FF.



The right part (two NOR gates) is the RS-FF. When the inputs of RS-FF are two 0, it would hold current input. However, when the inputs are two 1, the value would be undetermined.

D-FF introduced a NOT gates to ensured that the input from D could be guaranteed different. In the usage of the D-FF, pin E (stands for enable) is connected to a clock. When the E is 0, both input would be 0, which let the RS-FF to be stay at hold state. When E is 1, the input of RS-FF would be fully determined by D input. When D is 1, the inputs (S, R) of RS-FF are (1, 0). The RS-FF output (Q) would be 1 (work as Reset). When D is 0, the inputs of RS-FF are (0, 1). The RS-FF output (Q) would be 0. The output of the RS-FF is exactly the same as the D input. Hence, the D-FF could be used to store the data.

Sequential Logic is defined as that the output of the circuit depends on both the current and the *past inputs*. So a part of the circuit needs to be used to store the past inputs. D-FF could be used as the storage unit, this is first reason.

The second reason is D-FF would hold the past input when the circuit is synchronized with the clock rising edge. This feature allows the data to be delayed. These are the reasons that D-FF is important in Sequential Logic designs.

1. The minimum sizes of the output signals are
2. 36 bits (Multiply need sum up the value’s bits)
3. 4 bits (Consider 2’s complement, -4 is the smallest number that could be expressed in 3-bit signed number and 3 is the biggest number that could be expressed. If -4 subtract by 3, it should be -7. To save this number, it needs at least 4 bits which 3 bits to store its value and 1 bit for the sign as 1001.)
4. 15 bits (Binary expression of 17,000 is 100,0010,0110,1000. This needs 15 bits.)
5. 5 bits (The maximum 10-bit number is 1023 which is 210-1. A number divided by 32 means right shift for 5 bits, which means that it only saves the left-most 5 bits. Hence it needs 5 bits)
6. The synchronous reset means that the reset action could only be checked and triggered at the edge of the clock. When the settled clock edge is not coming, the reset signal would be delayed until the next settled edge of the clock. The asynchronous reset means that the reset action would be triggered as soon as possible which is not be synchronized with the clock edges. Asynchronous reset could give the maximum priority to the reset signal.

In a procedural block, synchronous reset could be expressed as

always @(posedge clock) begin

if (reset) begin

// Reset action here.

end else begin

……

end

end

and asynchronous reset could be expressed as

always @(posedge clock or posedge reset) begin

if (reset) begin

// Reset action here.

end else begin

……

end

end

The different of these two codes is the sensitivity list. For asynchronous reset, it would be triggered as the not only at the time of rising edge of the clock, but also the rising edge of the reset.

When a physical push button is used as an asynchronous reset input, the physical bouncing could be very bouncing when it is pressed. These bouncing could keep for several clock cycles (like the 100MHz on Nexys4-DDR). This could cause a problem when the reset signal releases at the settled edge of the system clock. The reset instructions would be executed at the same time as the normal instructions. This usually caused the register flip-flops are assigned at the same time by two blocks of codes at the same, which lead the circuit to metastable state. Compare to synchronous reset, it won’t have this problem as the reset signal would be treated as one input to the component and check as the same time with the other input wires.

1. The output of combinational logic circuit is a direct function of its current input. The output of sequential logic circuit is not only based on the current input, but also depends on past inputs. In sequential logic circuit, it needs circuit components to store the history values.

Sequential Logic circuit needs a clock to get synchronized. This allows the circuit to execute each operation step by step.

For the following circuits,

1. A 32-bit adder **doesn’t need** sequential logic elements. The result of the 32-bit adder only depends on the current input of two 32-bit numbers.

The Verilog code is

always @(\*) begin

result = a + b;

end

1. A 6-element shift-register **need** sequential logic elements. It needs to save the last result of the shift operation to do the current shift.

The Verilog code is

always @(posedge clock) begin

result <= {result[5:0], 1’b0};

end

1. A 10-bit overflow counter **need** sequential logic elements. It needs to save the last number of the counter to do the current counting.

The Verilog code is

always @(posedge clock) begin

counter <= counter + 10’d1;

end

1. A 16-bit signed multiplier **doesn’t need** sequential logic elements. The output value only depends on two input 16-bit signed values, which has no relationship with its last result.

The Verilog code is

always @(\*) begin

result = a \* b;

end

1. A 7-input multiplexer **doesn’t need** sequential logic elements. The output value only depends on the current input of the index selection wires and the current input of the values.

The Verilog code is

always begin

case(path)

3’d0: result <= route[0];

3’d1: result <= route[1];

3’d2: result <= route[2];

3’d3: result <= route[3];

3’d4: result <= route[4];

3’d5: result <= route[5];

3’d6: result <= route[6];

endcase

end

# Part 2: Reaction Timer Summary

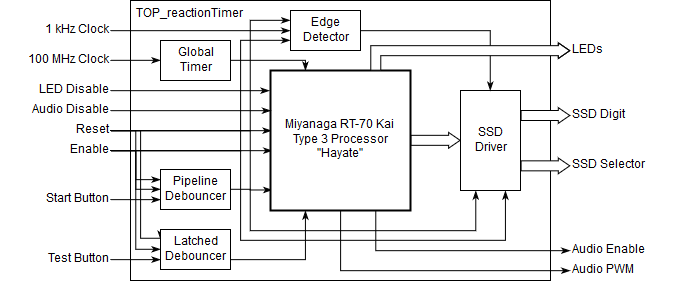
In this part, it shows the main feature list and high-level design of our reaction timer implementation. It also explains the details of how we implement these features. At the last of this report, we discuss the mean we applied testing during our development.

# Features

* Implement the four period of the Reaction Timer specification on Nexys4 DDR
* [Advanced] Shows the best record time and LED animation at the idle state
* [Advanced] Automatically return to idle state after display the reaction time for 10 seconds
* [Advanced] Display the reaction time as a 7-digit decimal floating point number in second
* [Advanced] Using Mersenne Twister MT19937 as one of the random number generators to set random event time
* [Advanced] Using a Linear Congruential Generator (LCG) manager as the other random number generator. This generator could manage maximum 4 LCG equations and automatically select between those equations for avoiding the periodic issues of the LCG method
* [Advanced] Introducing Advanced Random-number Election System (ARES) to manage all the pseudorandom number generators. The ARES allows to manage multiple random number generator with the same interface
* [Advanced] The recent reaction time would be to flash if it is the best time
* [Advanced] If tester pressed the button before the reaction signal, the reaction timer will mark this trial as a failure
* [Advanced] Support mono audio output as the optional reaction signal
* [Advanced] Support switches to enable and disable audio output and LED output
* [Advanced] Support using tri-color LED to display the reaction time level. At RESULT state, it would display the best recorded time level and current test level. At IDLE state, it would display the best recorded time level only
* [Advanced] Support disable the tri-color LED display
* [Advanced] Support clear the best recorded time at IDLE state
* [Advanced] Support VGA video signal output for visualization the testing configuration and operation reference manual.
* [Advanced] Using microphone to gather the environment wave to change the seed of the pseudorandom number generator.

# Architecture

The following diagram shows the top level design of the reaction module.



The input wires of the global module are 4 switches, 2 buttons, microphone wires and system 100MHz clock. One switch uses as reset, one uses as enable, one uses as audio output disable, and the other one uses as LED output disable. One button uses as the test start button and the other one uses as test button. The output wires are 16 LEDs, 7-segement display, VGA port and mono audio outputs.

The top module architecture is designed for processing the input and output signal directly to the FPGA board. All the input and output pins are chosen to be used CMOS 3.3V standard. The clock used in this design is the 100MHz clock provided by the crystal on the board. An 1kHz clock generated from the 100MHz clock is used as the clock for the Seven-Segments Display to select the digit.

The edge detector has the same logic as the one we used in Lab 4.

## Modules at Top Level

### Input Signal Debouncing

For all the switches used in this reaction timer, it doesn’t need any debouncer. But for those two buttons, they need different debouncing methods.

For the start button, it is only used as a signal for starting the test. So the debouncer for start button is just the sample pipeline debouncer as we implemented in the lab. The pipeline is described as a single register. The different from the lab version is the pipeline level could be modified by a parameter. If we still use the way as what we did in the Lab, the value used to detect the button pressed would be 2n+1-1 in decimal. To reduce the work of FPGA, we store the inverse of the button instead of its original value. When the button pressed for a while, the value of the pipeline should be 0. The debouncer module also uses logic equal to avoid the high impedance value comparison. Due to the limitation of the integer is 32-bit as default, hence this module could support up to 32-level pipeline. In the implementation of reaction timer, this level has been set to 5.

For the test button, it needs for lower latency to get accurate result. Here, we use a latched debouncer for the test button, which allows us to reduce latency within 10ns. The edge detector would add another 10ns delay to the signal. Hence, the total latency of the test button pressed is 20ns for these two modules. However, this won’t affect the precise of our final result. It will explain further later.

### Global Timer

The global timer module is a simple counter start running at the very beginning. It works like the UNIX timestamp. There is a 32-bit unsigned register and initializes as 0. This register increase 1 for each 100MHz rising edge. Hence, it would warp around every 42,949,672,950 nanoseconds, which is around 43 seconds. This value doesn’t reset neither reset signal is high nor enable signal is low.

This global timer would be used as the seed of the randomized number generator. This makes the possibility of the same interval as its previous trial reduce to every 42 seconds.

### Seven-Segments Display (SSD) Driver

The SSD driver have 13 inputs: eight 4-digit Enhanced Binary Coded Decimal (EBCD) number for each digit, one 8-bit dots display expression number, 1kHz clock rising wire, 100MHz system clock, reset and enable. Only 2 outputs are need for this module: one 8-bit output for the digit display with the point and one 8-bit for the digit selector. In the module implementation, the dot output wire separates from the digit expression.

There is a 3-bit register used as the digit selector. This number increases at each 1kHz clock rising edge. The digit selected integer could be calculated with index (displayIndex) with expression ~(8'd1 << displayIndex).

The EBCD is based on the standard Binary Coded Decimal (BCD) and defines some characters for the reaction timer. It expresses a character with a 4-bit unsigned integer. The following table shows the definition of the EBCD.

|  |  |  |  |  |  |  |  |  |  |  |  |
| --- | --- | --- | --- | --- | --- | --- | --- | --- | --- | --- | --- |
| EBCD | | Glyph | EBCD | | Glyph | EBCD | | Glyph | EBCD | | Glyph |
| 0000 | 0 | 0 | 0100 | 4 | 4 | 1000 | 8 | 8 | 1100 | 12 | A |
| 0001 | 1 | 1 | 0101 | 5 | 5 | 1001 | 9 | 9 | 1101 | 13 | F |
| 0010 | 2 | 2 | 0110 | 6 | 6 | 1010 | 10 | - | 1110 | 14 | I |
| 0011 | 3 | 3 | 0111 | 7 | 7 | 1011 | 11 | blank | 1111 | 15 | L |

The first 10 characters are the same as the normal BCD. Letters ‘A’, ‘F’, ‘I’, ‘L’ are defined for showing the word ‘FAIL’ when user failed to pressed the button within 9.9 seconds. The minus character is used to display the best record number when there is no valid record. Eight minus characters would be shown on SSD.

The EBCD code mapping is implemented in the SSD driver module as a single case statement. When the driver is at reset or disable state, it would output blank character with no dot displayed.

### VGA Driver and VRAM

The architecture only shows the major component of the entire system. It didn’t show the VGA output modules. The VGA output are used for displaying the reference manual of the reaction timer and display the configuration of the system. Now it could display the enable and disable state of audio output, LED output and Tri-color LED level output. The VGA display resolution has been fixed to 640480 running at 60Hz. But the display would detect it as 59Hz. It is now using the 8030 text mode output which is exactly the same as MS-DOS. It could display the full ASCII code with the IBM extension. The font library is ASC16 from the UC-DOS system, which means for each character it would be 8-bit width and 16-bit height.

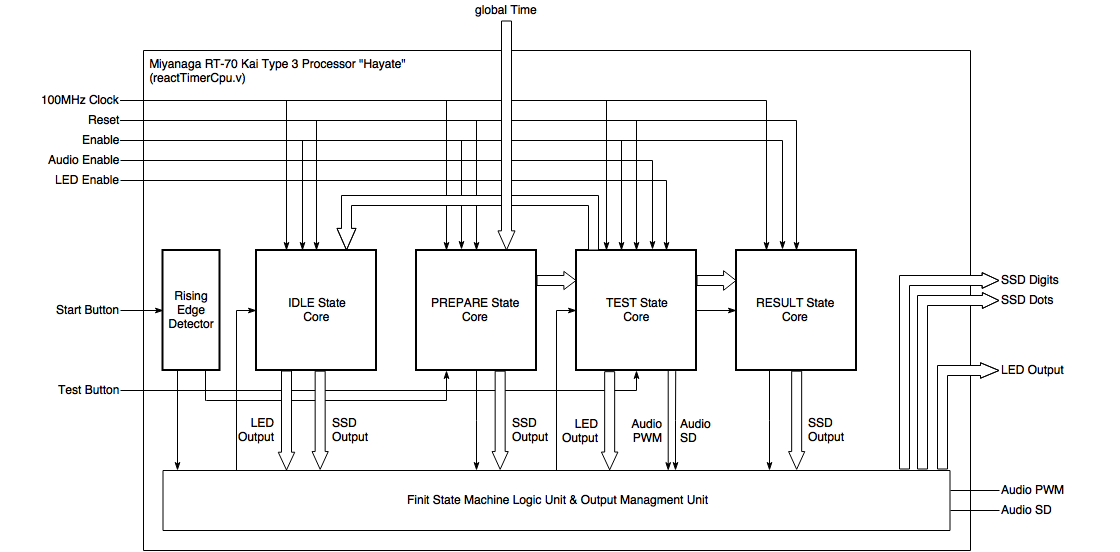
There is a video RAM module introduced to store the entire display data. It is an 8030 8-bit integer array which stores the ASCII code of each character. There is no double-buffering in this system.

## Central Logic Unit (Miyanaga RT-70 Kai Type 3 “Hayate”, M7T3)

This is the module (**M7T3**) which implements the finite state machine (FSM) of the main reaction timer logic. It contains quad cores inside the module for each state of the FSM and a unified FSM logic switching logic and output management unit. The entire module works as a multiplexer which output the correct signal of the state core module according to the current state of the FSM.

The input signal of **M7T3** are quite simple. Besides the standard 100MHz system clock, reset, enable wires, only four other wires are needed: start test button, test button, audio hint enables and LED hint enables. Outputs are just LED, SSD and audio outputs.

The following diagram shows the inner structure of the **M7T3** processor.



The FSM logic unit inside **M7T3** manages the state switching of the state. There are four states as the specification describes. The responsibility of each state and the policy that transfer between these states are

* IDLE: display the best score and waiting for the start pulse signal.
  + Start button pressed: move to PREPARE state.
  + Or else: play the LED animation.
* PREPARE: generate the random number and show the countdown animation on SSD.
  + Wait both the countdown animation and random number generation complete.
  + Or else: output the random number to TEST state and move to TEST state.
* TEST: according to the random number generate by the PREPARE state, give out the hint signal, wait for user pressing the test button.
  + Play the audio sound hint and/or lighten the LEDs hint after the random number time given by PREPARE state.
  + Test button pressed:
    - If the test button pressed before the hint occurs or not being pressed in 10 seconds, the result is FAIL and move to RESULT state.
    - If the test button pressed in 10 seconds, the result would be recorded and move to RESULT state.
* RESULT: show the test result of the TEST state.
  + After 10 seconds, go to the IDLE state.

If reset signal is on, the state would be back to IDLE state. A 2-bit register is used to store the current state. This register is checked at each 100MHz system clock rising edge by the FSM logic unit. The FSM logic unit is implemented as an *always* block. Within this block, when the state is not changed, the output signal would be refreshed from the output of the state core. If there is no output for one type signal, for instance, there is no audio output for IDLE state, the audio output data would be reset to default value.

### IDLE state core

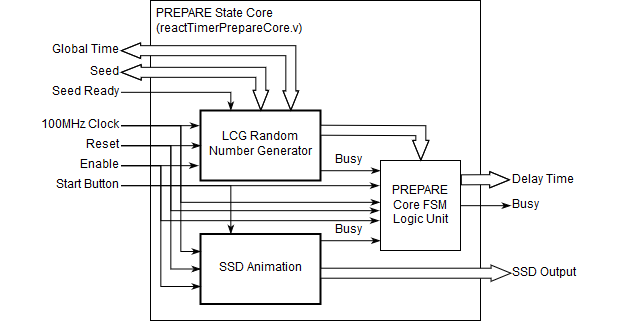
The logic of this core is pretty simple. It stores the binary of the best time. The test result comes directly from the TEST state core. There is also an LED animation playing module embedded in this core.

The outputs of this core are LED animation result and SSD signals.

### PREPARE state core

PREPARE state core implements in a small two state FSM. One state is idle and the other one is busy. For the idle state, this core just waits and detects the start signal rising. The busy state is waiting for the random number generator to generate the random number and SSD animation module to finish playing the countdown animation.

The following diagram shows the PREPARE core architecture.



Once the start button is at rising edge, PREPARE core FSM logic unit would change the state from idle to busy. LCG random number generator and SSD animation module would start to work concurrently. The LCG and SSD animation module both have a busy wire output which shows whether their mission is finished. Once all of these wires changed to 0, this module would change the busy output from 1 to 0.

The output of SSD animation module directly connects to the output of PREPARE core. The busy state of LCG generator is checked by the FSM Logic Unit at every 100MHz clock cycle rising edge. Once the falling edge of the LCG wire is detected, the output 32-bit random number would be modular within 700,000,000 that could guarantee that the waiting seconds are limited to maximum 7 seconds.

#### Random Number Generator

In M7T3, it contains two random number generator: Liner Congruential Generator and Mersenne Twister MT19937 standard implementation. All the pseudorandom number generators used in **M7T3** needs 32-bit seeds and could generate 32-bit output. The time deadline for generating a random number is 4 seconds because the animation duration of the countdown is 4 seconds. Even if the random number generator cannot meet this requirement, it is still okay to be used. However, the duration of the PREPARE state would be longer and this is transparent to tester.

These two pseudorandom number generators are managed under an Advanced Elective Generator Integrated System (AEGIS) unit. This system is embedded in the always block of the FSM logic unit in the PREPARE state core. All these generators are using the same port interfaces, which allows us to add and remove pseudorandom number generators easily on source code level. When output the random number result, AEGIS system would limit the output number less than 805,306,366 (32’h2FFFFFFE, calculated by ), which is around 8 seconds. This allows the tester won’t wait too long.

LCG generator is used as the random generator. There are four equations have been introduced in our design. Three of them are all used by previous version library and compiler: Apple Carbon, Microsoft Visual C/C++ 6.0 and C++11 official implementations. One is from the book *Numerical Recipes*. All these equations have been divided into two steps that could make sure that each step could be done within 10ns. According to the report, in **M7T3** implementation, it needs 10.568ns to complete the multiply and addition at the same time (7.336ns for logic and 3.232ns for route) for the Visual C/C++ 6.0 equation.

Mersenne Twister is a pseudorandom number generator widely used in multiple language compilers. It was introduced by Dr. Matsumoto and Dr. Nishimura in 1997. The original paper could be found at <http://www.math.sci.hiroshima-u.ac.jp/~m-mat/MT/ARTICLES/mt.pdf>. It has a standard implementation called MT19937. The Mersenne Twister module of our implementation in **M7T3** is compatible with MT19937 standard.

To implement a Mersenne Twister generator, the following parameters are needed in the following description

* – the length of the vector. This number is 624 in **M7T3** implementation.
* – the bits of the final output random number generator. This number is 32 in **M7T3** implementation.
* – an -length long row vector construct with 0 and 1. Hence a 32-bit unsigned binary.
* – the number a bit from the rightmost of .
* – integer in range .
* – integer in range .
* – a matrix.

* – a concatenation of higher bits of and filling 0 with lower bits. In Verilog, it could be expressed as {x[w:r], r{0}}.
* – a concatenation of lower bits and filling 0 with higher bits. In Verilog, it could be expressed as {(w-r){0}, x[r:0]}.
* – a concatenation with as higher bits and as lower bits.

To initial the MT19937 generator, first prepare a -length long row vector with the following equations:

According to the paper, the result of is

In the implementation, this is exactly using the Linear Feedback Shifting Register (LFSR) as the assignment paper mentioned. The description above is actually a level LFSR: take the higher bits of the  and the lower bits of the , then goes through a linear transformation . However, the equation above is actually take rounds of LFSR. It won’t affect the LFSR period if  and are relatively prime.

With this equation, this allows us to extend the period to , which is reason of the name contains 19937. To extract the value from the result of the data, we only need to post multiply an invertible matrix. With integer parameters (, , and ) and two -width bit mask integers ( and ), this multiplication could be expressed by

And the result is the output of one round. Summarized the MT19937 implementation, there are three states: seed initialization, twister and extraction.

In M6T3, we tried to implement this MT19937 with a multiple level FSM. However, it would take 20,118 look up tables and 21,276 flip-flops which takes a very long time to synthesis and implementation (and even gets an unknown error from Vivado). So a new pipelined optimization has been introduced to **M7T3** implementation. With this implementation, the number of look up tables has been reduced to only 304. Flip-flops usages has been reduced to 165, which makes MT19937 could be fit in the implementation.

In the pipeline implementation, there is still one FSM has been used to distinguish the idle and busy state. In the busy state, it would go through the entire three states of the MT19937 standard and update all the values of the 624 vectors to the final result. This would take 624 clock cycles to complete the entire computation. Then update all the values every 624 times access for asking the next random number.

The seed using by LCG and MT19937 modules are generated by the Advanced Seed Management (ASM) system. ASM would use the data from the microphone which continuously gathered the environment data and XOR to the global timer counter to make the seed to be more unpredictable.

In each clock cycle of busy state (or the SEED state in the implementation), it would complete one round calculation within 10ns (100MHz system clock period) and meet the time constraints without providing a negative WNS.

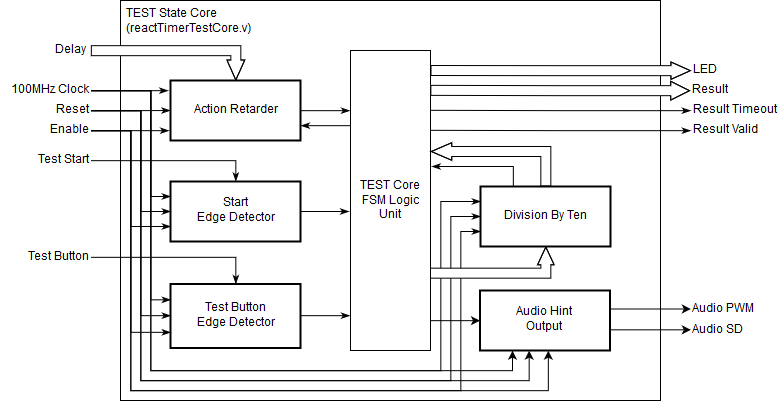
#### SSD Animation

The animation of SSD is implemented by an increase counter which increase 1 at when the timer clock rising. The clock is set to be 1Hz and increase the frame index and update the output frame according to the frame index at the rising edge of system 100MHz clock. The frames are stored as a 32-bit integer array inside the module.

### TEST state core

After the PREPARE core provides the delay duration, the TEST state core start working. This core is responsible for showing the hint after a specific period of time, this is implemented by an action retarder. Then it would measure the gap of the hint provided to the test button pressed. Last, it calculates the data to be output.

The following diagram shows the inner structures of the TEST state core.



The inner framework of the TEST state core is a FSM as well. There are 3 states for this core:

* Idle state: waiting for the test start signal. When the start signal pulse comes, the state switch to wait state. At the same time, the FSM Logic Unit would give action retarder a pulse. The action retarder would give another pulse after the time provided by the PREPARE state core.
* Wait state: At this state, the action retarder would hold the pulse for a while, and output the pulse after several clock cycles, then move to test state. FSM Unit detects this pulse for each 100MHz clock cycle rising edge, and then move to test state. If tester hits the test button at this state, this test would be marked as ‘FAIL’ and move to finish state.
* Test state: According to the configuration, the LED would be shown and/or the audio would be played. At the same time, counter starts to count from 0 and increases 1 for each 100MHz clock cycle. Once the user hits the test button, the counter would stop to count and then move to the finish state. If the counter reaches the limitation (by the default, is 999,999,999, which is actually 9.999 seconds), this test would also be marked as ‘FAIL’, then move to finish state.
* Finish state: the result of the counter is a 9-digit integer in decimal. The number goes into a division module which could divide a number by 10. If the result is not ‘FAIL’, it would wait until the divider finish and then mark the result as valid. Or else set the result to be 0, set the timeout signal as 1, then mark the result as valid.

Action retarder is just a 2-state FSM which detect the input pulse and then wait by increasing a counter as it reaches the number provided by the TEST state core FSM logic unit. It would ignore all the other pulses while holding the first signal it detects.

#### Division by Ten Module

The precise of the counter is 10ns. However, we have only 8 digits which allows us to display the number at 100ns. So we have to divide the counter result by 10. We originally use counter / 10 here to calculate this value. However, it fails to reach the time constraints (provides a negative WNS). Instead of using the division calculation directly, we are using an 8 clock cycles delay divider to do this calculation.

In this algorithm, we need two parameters: and . Suppose the number we want to divide by 10 is . First step, we can get () by

next, get () by

after that, get () by

then, get () by

afterwards, get () by

This number is actually very closed to the data as what we want. To make the data more precise, we need to calculate and check it is greater than 9 or not by

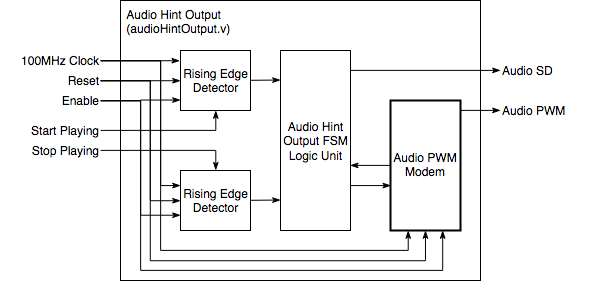
and finally get the result

which is correct in the 32-bit integer value for division by 10. This module is implemented by a 9-state FSM, which means that it adds a 9-cycle delay to the system at 100MHz.

In the implementation, there is a busy wire which marked whether the output data is ready. The TEST core unit would detect the falling edge of the divider. Once it gets the result, it will move on to the next state.

#### Audio Hint Module

The following diagram shows the architecture of the audio output module.



Inside this module, it stores the hardcoded audio samples. The PWM modem would provide a rising edge to update for the sample. The whole module is implemented by a 2-state FSM: idle state and playing state. When the start playing input wire occurs a rising edge, the state machine would move from idle to playing state, and sending audio samples to audio PWM modem. The state would be kept until the stop playing signal occurs a rising edge.

For the PWM modem using in the implementation, we decided to use 8-bit output for the sample. In the original plan, this should be 16-bit. However, the 100MHz system clock would limit the signal. To express a 16-bit number using PWM, we need 65,536 clock cycle for one number. As the system could provide a 100MHz clock, it allows us to express 1,525 samples only. To express a sine wave, we need at least 4 samples to express a cycle (which is 0, 1, 0 and -1), hence the limitation of the 16-bit PWM module with 100MHz as the input clock could only generate a 381Hz sine wave. However, the Stuttgart pitch (A440) is 440Hz, this could even fail to meet this requirement. If we are using an 8-bit PWM module with 100MHz system clock, it allows 390,625 samples per second, and this limitation would be 97,656Hz for the 4-sample sine wave. It allows use to generate A11 which is a 56,320Hz sine wave. Hence, the resolution for the PWM is decided as 8-bit.

The PWM modem is implemented by an 8-bit counter, which counts from 0 to 255 and then warp back to 0. The PWM output is a simple comparator of the counter and the sample number. When this counter start counting, it would ignore the change of the input during the counter count from 0 to 255. It would update the current sample at each time the counter become 0. It would give signal back to FSM logic unit to switch to next sample as a rising edge when the counter reach 128.

In **M7T3**, it uses a counter to control the frequency of the output signal. There is another counter which counts from 0 to a pre-set limitation parameter. This counter is called as frequency counter. Once the frequency counter reaches the limitation, the signal counter increase 1. For example, when the limitations set as 1, the signal counter increase 1 only when the frequency counter reach 1. The frequency of output PWM has been reduced to its half.

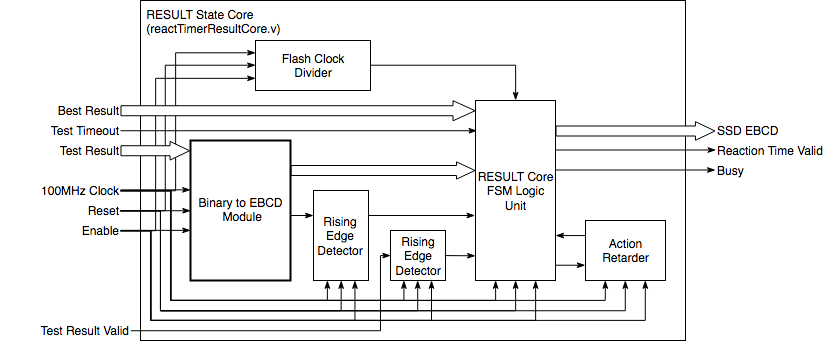
According to the Nexys4 DDR reference manual, the PWM is using the time of 1 divided by the entire period the get the value. As the document mentioned, we should output the data as PWM to express the data as what we want. However, we output the sine wave data to the audio pin, the sound is very strange. We finally choose to output a constant value which is 8’h80 which is exactly half of the duty cycle. And performs very well. Using Moku, we could actually find out that the PWM wave output directly without transfer the PWM to analog signal as the DAC module output. This is one of the thing that not mentioned in the reference manual. **M7T3** could output sine wave, but for better usage, we decide to output square wave.

Another thing which is not mentioned in the reference manual, is the existence of the AUD\_SD pin. In section 16, i.e. Mono Audio Output, it only mentioned about the PWM pin (A11), but not mentioned the AUD\_SD pin. AUD\_SD connects to the FPGA pin D12. This pin connects to the amplifiers of the low pass filter in the Figure 29 of the reference manual (Sollen-Key Butterworth Low-Pass 4th Order Filter), which controls the enable of the amplifiers. Hence, AUD\_SD is actually controls the enable of the low pass filter. In **M7T3** implementation, the AUD\_SD pin is hard coded to output high which means that it is always enable for output data.

### RESULT state core

RESULT state is implemented as a FSM. It contains 3 states: idle state, convert state and wait state. It starts at idle state, waiting for the test result from TEST state available. Then, RESULT state core would transfer the reaction time result from binary number into EBCD, and then according to the result recorded from the IDLE core, check whether the current result is the best or not. This stage is called convert state. After that, it switches to wait state. If the result is the best result, the result would be flashed for 10 seconds, or just display the time or ‘FAIL’ for 10 seconds. Then back to idle state.

The framework of the RESULT state core is shown in the following figure.



A clock divider is introduced to be used as the flash frequency. In **M7T3**, this clock has been set to be 1Hz. The SSD output signal is updated based on the divided clock from the clock divider at each rising edge of the 100MHz clock.

A busy wire is used as the signal for the core is at none-idle state. **M7T3** FSM logic unit checks this signal at every 100MHz rising clock. When the falling edge of busy is detected, **M7T3** state goes back to IDLE state.

An action retarder is introduced with a fixed parameter for waiting 10 seconds. It would start working right after the Binary to EBCD module finishes its work. Once the delayed signal pulse is detected by the FSM logic unit, the busy output would change from 1 to 0.

#### Binary to EBCD

To display the number on SSD, we have to separate the number from binary to each digit of decimal. Generally, we could divided by and modular by 10 to get the th digit of the number. However, this would give a very bad performance on WNS using division. Hence, we need another method to get each digit.

First, we need a way to express each digit of the number. There are 10 numbers of each decimal digit (0 to 9), which needs for 4 bits to store all of them. This is called Binary Coded Decimal (BCD), or actually 8421 coding. In **M7T3**, we are using an enhanced version called Enhanced Binary Coded Decimal (EBCD) which is compatible with BCD. The method of transferring binary to BCD could be used in our design.

The method we are using is called Shift-Add Binary to BCD transfer. In **M7T3**, we will finally display an 8-digit decimal. Hence, we need 32 bits to store the BCD result. The algorithm can be display as the following

1. Prepare two shift registers. One 28-bit for storing the reset of the original binary bits. The other one 32-bit for storing the transferred BCD result. Initial the 28-bit register as the original data. Initial the 32-bit register with a 32-bit decimal 0 (32’d0). Initial a counter to 28.
2. Check each decimal half-byte of the 32-bit register, if all the digit is greater than 4, increase this half-byte by 3.
3. Left shift the 32-bit register, and append the most-significant bit of the 28-bit register to the right-most bit. Left shift the 28-bit register. Decrease the counter by 1. If the counter is 0 now (left shift for 28 times), the transfer is finished, or else go to step 2.

Here is an example of transferring decimal 255 (binary 0000,0000,0000,0000,0000,1111,1111) to 8421 coding BCD with this algorithm. For the first 20 bits, there are no different because they are all 0. The diagram shows from step of 20 to 28. Step dot 2 means Shift operation from 28-bit register to 32-bit register and step dots 3 means the Add 3 operation for checking each decimal digit.

|  |  |  |
| --- | --- | --- |
| Steps | 28-bit register | 32-bit register |
| 0 | 0000 …… 0000 0000 1111 1111 | 0000 0000 0000 0000 0000 0000 0000 0000 |
| … | …… | …… |
| 20 | 1111 1111 0000 …… 0000 0000 | 0000 0000 0000 0000 0000 0000 0000 0000 |
| 21 | 1111 1110 0000 …… 0000 0000 | 0000 0000 0000 0000 0000 0000 0000 0001 |
| 22 | 1111 1100 0000 …… 0000 0000 | 0000 0000 0000 0000 0000 0000 0000 0011 |
| 23.2 | 1111 1000 0000 …… 0000 0000 | 0000 0000 0000 0000 0000 0000 0000 0111 |
| 23.3 | 1111 1000 0000 …… 0000 0000 | 0000 0000 0000 0000 0000 0000 0000 1010 |
| 24.2 | 1111 0000 0000 …… 0000 0000 | 0000 0000 0000 0000 0000 0000 0001 0101 |
| 24.3 | 1111 0000 0000 …… 0000 0000 | 0000 0000 0000 0000 0000 0000 0001 1000 |
| 25 | 1110 0000 0000 …… 0000 0000 | 0000 0000 0000 0000 0000 0000 0011 0001 |
| 26.2 | 1100 0000 0000 …… 0000 0000 | 0000 0000 0000 0000 0000 0000 0110 0011 |
| 26.3 | 1100 0000 0000 …… 0000 0000 | 0000 0000 0000 0000 0000 0000 1001 0011 |
| 27.2 | 1000 0000 0000 …… 0000 0000 | 0000 0000 0000 0000 0000 0001 0010 0111 |
| 27.3 | 1000 0000 0000 …… 0000 0000 | 0000 0000 0000 0000 0000 0001 0010 1010 |
| 28 | 0000 0000 0000 …… 0000 0000 | 0000 0000 0000 0000 0000 0010 0101 0101 |
| Result | 0000 0000 0000 …… 0000 0255 | 0000 0000 0000 0000 0000 0002 0005 0005 |

The reason for shift 28 times is we have to shift all the bits from the raw data to BCD expression. Using left shift and take the MSB would be simple for the implementation by taking the bit and shift for one bit.

To calculate the BCD of the binary, we first consider one decimal digit case. For 4-bit binary, it would increase 1 to the higher 4 bits when it reaches 16. For 4-bit 8421 coding BCD, it would increase 1 when it reaches 9. Hence, to transfer 4-bit binary to 4-bit 8412 coding could be processed by plus 6 to the binary.

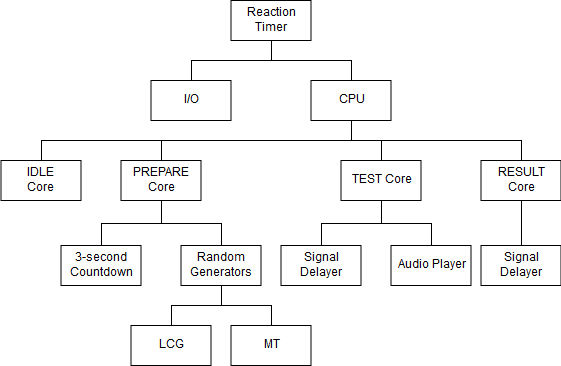
Then we consider the way to transfer multiple digit decimal into BCD. For example, decimal 30 (in binary 11110). It could be get by left shifting decimal 15 (in binary 1111). The 8421 coding BCD of 15 is 0001 0101. Shifting left 1 bit we could get 0010 1010. Check each half byte data of the BCD expression, 1010 is invalid, so we increase it for 9, the BCD is now becoming 0011 0000, which is exactly 30.

However, using increase 9 would affected the entire number. Instead of calculate after shift, we could actually check the BCD 8421 coding before we shift the bits. To increase one in binary, it needs the number to be 16. However, in 8421 coding BCD, this number should be 10. Left shift means multiply the number by 2. When the BCD digit is greater than 4, for example 5, shift left would let it to be 10, and need to increase 1 in the higher 4-bits and reset the current 4-bit to 0, and express this action is actually binary 1,0000, which is 16. In other words, the 5 here is actually 8 as what we want for the shift. Hence, before we shift the number, we check each digit is greater than 4 or not. If so, plus 3 to make this number could increase 1 in the higher 4 bits. Then do the shift and append the new bit. And this is the method that shown before and used in **M7T3**.

This algorithm needs 28 rounds for one transfer. Each round has 2 stages. Hence, 56 clock cycles to calculate the BCD from a 28-bit binary. The module is implemented by a 3-state FSM. A busy wire is present as an output for showing the signal of whether the transfer is completed. The FSM logic unit of the RESULT state core is going to check this wire for each rising edge of the 100MHz clock cycle.

# Design Verification and Validation

During the development, we are using test-driven development and go through the entire development period. We first determined the high level design. Next, we separated the function of each state, and drew the top level design diagram. The right diagram is the one that we draw before we started our implementation.



Then, we wrote each leaf-level modules, and write a test bunch for each of them as unit test. The All the unit tests are using strong robust equivalence classes black box testing. During this stage, we found the following two critical problems when we write the simulation test bunch for unit testing:

The first problem is some modules needs to execute very long time if we are using the real clock settings in the simulation. For example, the countdown animation module using in the PREPARE state core need a simulation of at least 4 seconds. Doing 1 second simulation would take a very long time on a Intel Core i7-3770 processor (4 cores 8 threads at 3.7GHz). To solve this problem, we introduced the dynamic parameter for each module. All these parameters are providing a default value which would be used as implementation configuration. But inside the test bunch, we are using a different parameter which allows us to do the simulation in a different condition, but see how the module works. This method is widely deployed in **M7T3** implementation.

The other problem is simulation cannot simulate the time constraints of the real board. In the simulation, if we use division and modular, it would still provide the same result and correct output. However, division cannot reach the requirement of the Nexys4-DDR board constraints. For 10ns period of time, it cannot complete these calculations. This data would only be available when the implementation stage completed.

The Worst Negative Slack (WNS) and Total Negative Slack (TNS) show whether our design is valid to the board and meet the time constraints. According to Lyle’s explanation, we found the information within the timing summary log in the reports (impl\_1\_route\_report\_timing\_summary\_0). Inside the report, it has a section named “Design Timing Summary”. If there is a negative WNS had been found in our design, it will list it here. In our design, the following parts are mention failed to meet the time constraints:

1. Separate the binary number into decimal digits with division and modular directly. (Solved by replacing the original method with Binary to EBCD module)
2. LCG result calculation. (Solved by introducing a FSM and separate the multiplication and addition into two steps)
3. Test result divided by 10. (Solved by introducing a new module which is especially designed to divide a number by 10)

Now, all the module designs specified timing constraints are met according to the implementation reports. The maximum delay path appears at the multiplication in LCG module. The data path delay is 9.579ns, which is very closed to the 10ns requirement. The final design takes 1,059 look up tables and 1,047 flip-flops after implementation optimization.

In the design, most of the function has been packaged into a single module. Some of the small common code pieces are gathered into a task. This is the light-weight cohesion code solution, which allows us to improve the robust of our implementation.

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# Appendix

This assignment got special page extension from Lecture Lyle.

