

# Introduction to Computer Science

## Assignment 08

Constructor University,  
Bremen, Germany

April 10, 2024

- Marks are only awarded for producing the correct solutions and justifying how you arrived at them.
- Please write legibly. Illegible answers will NOT be marked.

### Question 1

Using the Definition of the syntax of Boolean formulas, show that the expression:  $\neg(x \vee (y \wedge 1))$  is a well-formed formula.

[1 mark]

### Question 2

Consider the set of Boolean variables  $D = \{x_1, x_2, x_3\}$  and the interpretation  $\mathcal{I}(x_1) = 1$ ,  $\mathcal{I}(x_2) = 0$ , and  $\mathcal{I}(x_3) = 1$ . Determine the semantics of the formula:  $x_1 \vee \neg(x_2 \vee 0) \wedge x_3$  using the Definition of the extended interpretation for formulas  $\mathcal{I}^*$ .

[1 mark]

### Question 3

By means of a truth table, demonstrate the validity of the De Morgan's laws for three variables, i.e.:

$$\neg(\varphi \wedge \psi \wedge \chi) \equiv \neg\varphi \vee \neg\psi \vee \neg\chi$$

[1 mark]

### Question 4

By means of algebraic manipulation (Boolean equivalence laws), show that the following Boolean identity holds:

$$(\neg x \wedge y) \vee (y \wedge \neg z) \vee (y \wedge z) \vee (x \wedge \neg y \wedge \neg z) \equiv y \vee (x \wedge \neg z)$$

[1 mark]

**Question 5**

For the formula:  $\neg x \vee \neg y \vee (x \wedge y \wedge \neg z)$  obtain the *Disjunctive Normal Form* (DNF).

[1 mark]

**Question 6**

A binary T has the following traversal sequences:

preorder: J, G, A, E, V, L, M, W, Z

inorder: A, E, G, J, L, M, V, W, Z

Present the binary tree T.

[2 marks]

**Question 7**

Natural numbers are an inductive data type. The set of natural numbers is infinite, yet it can be defined in Haskell by an Algebraic Data Type:

```
data Natural = Zero | Succ Natural

instance Show Natural where
  -- show :: Natural -> String
  show n = show (iterate n) where
    iterate Zero      = 0
    iterate (Succ m) = 1 + (iterate m)

-- some constant
zero, one, two, three, four :: Natural
zero = Zero
one  = Succ Zero
two  = Succ one
three = Succ two

-- addition
infixl 6 <+>
(<+>) :: Natural -> Natural -> Natural
n <+> Zero      = n
n <+> (Succ m) = Succ (n <+> m)
```

In this representation, implement the multiplication (represented by  $\langle . \rangle$ ) as an infix left-associative operator with higher precedence than addition.

[3 marks]