

Sorting Algorithms

Sorting

- *Sorting* is a process that organizes a collection of data into either ascending or descending order.
- Formally
 - Input: A sequence of n numbers $\langle a_1, a_2, \dots, a_n \rangle$
 - Output: A reordering $\langle a'_1, a'_2, \dots, a'_n \rangle$ of the sequence such that $a'_1 \leq a'_2 \leq \dots \leq a'_n$
 - Given the input $\langle 6, 3, 1, 7 \rangle$, the algorithm should produce $\langle 1, 3, 6, 7 \rangle$
 - Called an instance of the problem

Sorting Algorithms

- Some of sorting algorithms:
 - Bubble Sort
 - Selection Sort
 - Insertion Sort
 - Merge Sort
 - Quick Sort
 - Heap Sort
- These are among the most fundamental sorting algorithms

Bubble Sort

- Elements of the array are stored in $a[0 \dots n-1]$
- Repeatedly
 - Compare every pair of adjacent elements.
 - swap adjacent elements that are out of order.
- After one iteration, maximum/minimum element moves to the last position of the list like a bubble moving up to the top in boiling water.
- Repeat the process $n-1$ times to sort entire list.
- In each iteration $n-1$ pairs of adjacent elements are to be compared and swapped whenever necessary.
- Can we reduce the number of pairs elements to be compared?

Bubble Sort

23	78	45	8	32	56
23	45	78	8	32	56
23	45	8	78	32	56
23	45	8	32	78	56
23	45	8	32	56	78
23	8	45	32	56	78
23	8	32	45	56	78
8	23	32	45	56	78
8	23	32	45	56	78

Bubble Sort Algorithm -1

```
template <class Item>
void bubbleSort(Item a[], int n)
{
    for (int i = 0; i < n; i++)
    {
        for (int j = 0; j < n; j++)
            if (a[j] > a[j+1])
                swap(a[j], a[j+1]);
    }
}
```

Sorting Algorithms

- Some of sorting algorithms:

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- Merge Sort

- Quick Sort

- Heap Sort

6 elements

a_1

$a[0]$ → After 1st traversal of my entire process

a_1 a_2

$a[i]$ → After 2nd traversal

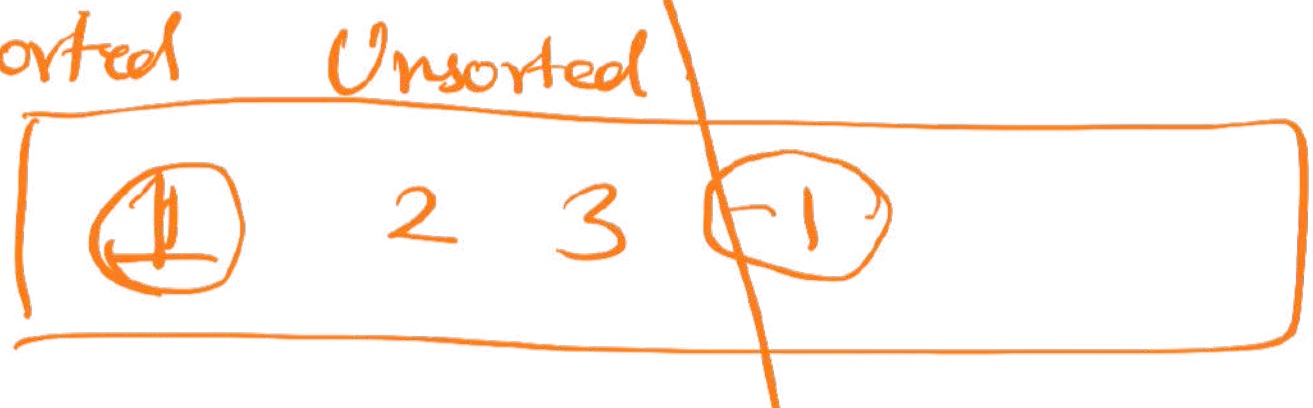
Given n elements, $n!$ permutations are possible.

I/p → $\{A\}$
1 permutation → $\{a_1 < a_2 < a_3 < \dots < a_n\}$
O/P

Selection Sort

- Partition the input list into a sorted and unsorted part (initially sorted part is empty)
- **Select** the smallest element in the unsorted list and put it to the end of the sorted part
- Increase the size of the sorted part by one
- Repeat this $n-1$ times to sort a list of n elements

I/P
↓



Sorted

Unsorted

	23	78	45	8	32	56	Original List
--	----	----	----	---	----	----	---------------

8		78	45	23	32	56	After pass 1
---	--	----	----	----	----	----	--------------

8	23		45	78	32	56	After pass 2
---	----	--	----	----	----	----	--------------

8	23	32		78	45	56	After pass 3
---	----	----	--	----	----	----	--------------

8	23	32	45		78	56	After pass 4
---	----	----	----	--	----	----	--------------

8	23	32	45	56		78	After pass 5
---	----	----	----	----	--	----	--------------

Selection Sort Algorithm

```
void selectionSort(Item a[], int n) {  
    for (int i = 0; i < n-1; i++) {  
        int min = i;  
        for (int j = i+1; j < n; j++)  
        {  
            if (a[j] < a[min])  
                min = j;  
        }  
        if (min != i)  
        {  
            swap(a[i], a[min]);  
        }  
    }  
}
```

Selection Sort -- Complexity

- What is the complexity of selection sort?
- What are best, average, and worst case complexities?
- It is $O(n^2)$ for all cases!

min = 4

$a[3]$
 \updownarrow
 $a[0]$
= index of minimum

Sorted
Unsorted

n
 $n-1$

Original List

2

$n-2$

After pass 1

2nd minimum

After pass 2

3rd minimum

After pass 3

After pass 4

After pass 5

n times

$O(n^2)$

$n-1$

Sorted

Unsorted

23 78 45 8 32 56
0 2 4 5

8 78 45 23 32 56
1 4 5

8 23 45 78 32 56
2 5

8 23 32 78 45 56

8 23 32 45 78 56

8 23 32 45 56 78

Selection Sort Algorithm

```
void selectionSort(Item a[], int n) {
```

```
  for (int i = 0; i < n-1; i++) {
```

```
    int min = i;
```

```
    for (int j = i+1; j < n; j++)
```

```
    { if (a[j] < a[min])
      min = j;
```

```
    } if (min != i)
```

```
    { swap(a[i], a[min]);
```

```
    }
```

```
}
```

IP

Unsorted

1	2	3	4	5
---	---	---	---	---

$O(n)^2$

min = 2

Sorted

1	3	5	6	8	9
---	---	---	---	---	---

Unsorted

2	3	4	5
---	---	---	---

3	5	1	6	8	9
---	---	---	---	---	---

min = 2

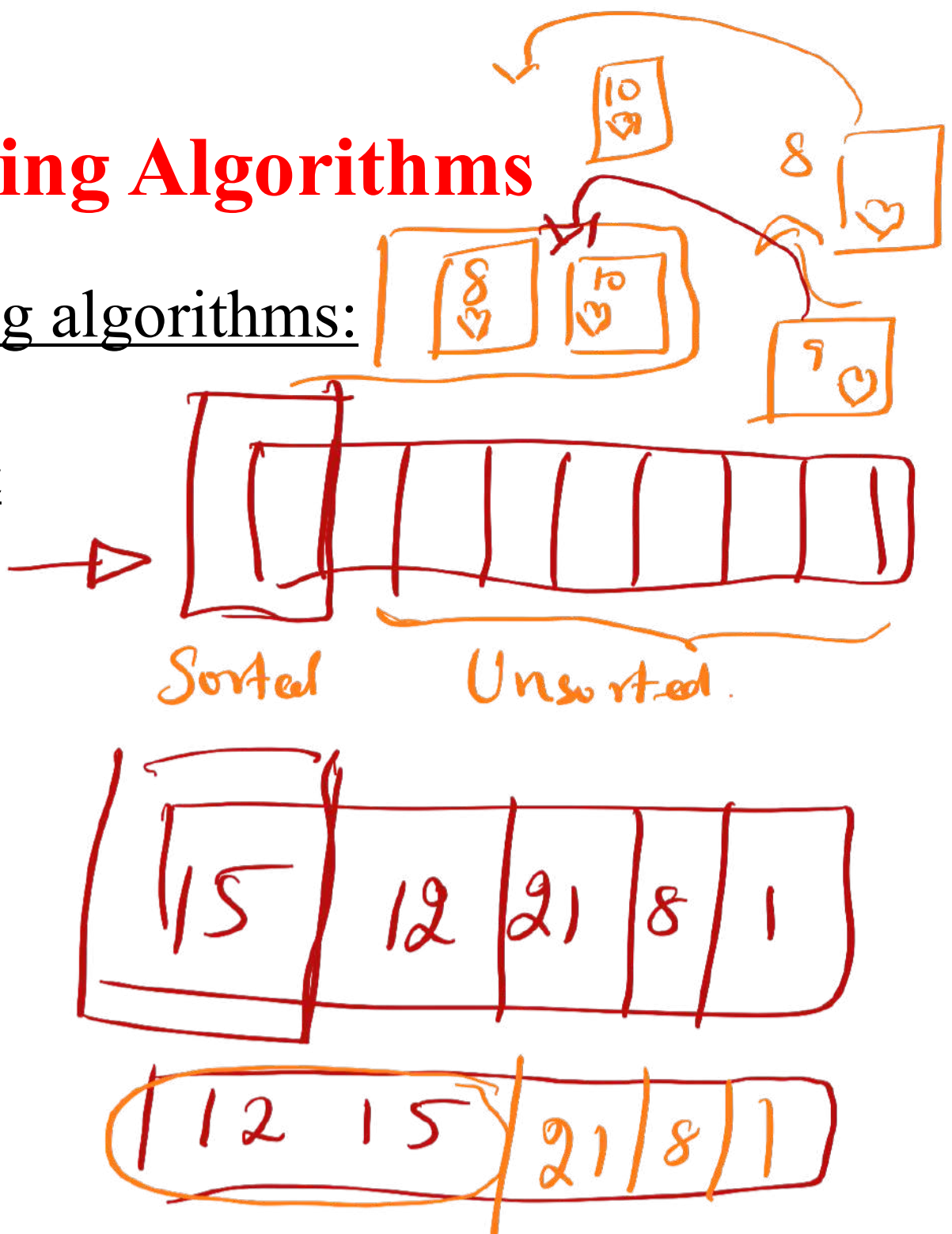
5	3	6	8	9
---	---	---	---	---

Unsorted

Sorting Algorithms

- Some of sorting algorithms:

- Bubble Sort
- Selection Sort
- Insertion Sort
- Merge Sort
- Quick Sort
- Heap Sort





- Insertion sort is a simple sorting algorithm that is appropriate for small size inputs.
- Again, the list is divided into two parts: sorted and unsorted.
- In each pass, the first element of the unsorted part is picked up, transferred to the sorted sublist, and **inserted** at the appropriate place.
- A list of n elements will take at most $n-1$ passes to sort the data.

Sorted

Unsorted

23 78 45 8 32 56

23 78 45 8 32 56

23 45 78 8 32 56

8 23 45 78 32 56

8 23 32 45 78 56

8 23 32 45 56 78

2 3 4 5 7 8 8
2 3 4 5 ? 7 8

Original List

2 3 ? 4 5 7 8

After pass 1

8 2 3 4 5 7 8

After pass 2




After pass 3

After pass 4

After pass 5

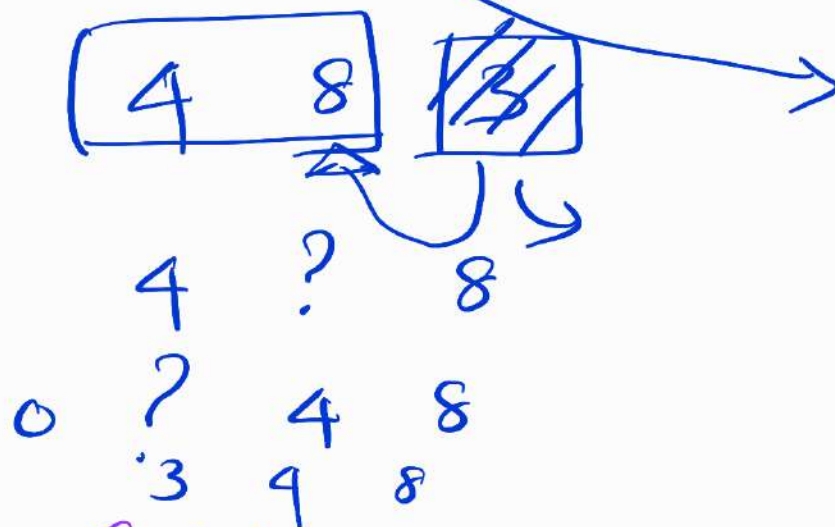
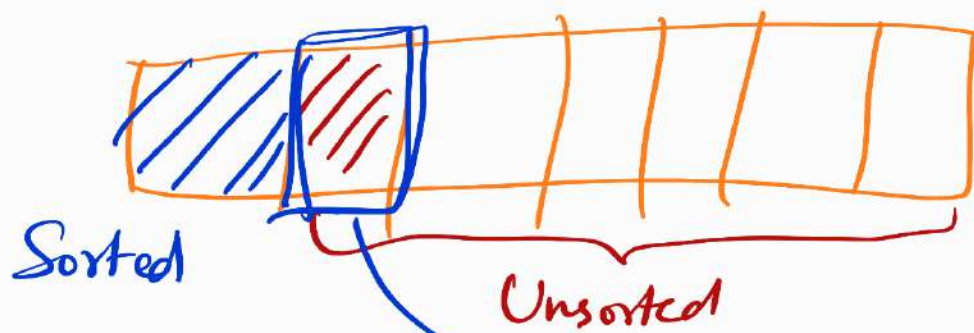
Final output

Insertion Sort – Complexity

- Running time depends on not only the size of the array but also the contents of the array.
- *Best-case:* $\rightarrow O(n)$
 - .
- *Worst-case:* $\rightarrow O(n^2)$
 - .
- *Average-case:* $\rightarrow O(n^2)$
 - .

Insertion Sort – Complexity

- Running time depends on not only the size of the array but also the contents of the array.
- ***Best-case:*** $\rightarrow O(n)$
 - Array is already sorted in ascending order.
- ***Worst-case:*** $\rightarrow O(n^2)$
 - Array is in reverse order:
- ***Average-case:*** $\rightarrow O(n^2)$
 - We have to look at all possible initial data organizations.



Insertion Sort:-

for ($i = 1$; $i < n$; $i++$)

{

hole = $a[i]$;

$j = i - 1$;

while ($j \geq 0$ && $a[j] > \text{hole}$)

{

$a[j+1] = a[j]$

$j = j - 1$

}

$a[i]$

$O(N)$

1, 2, 3, 4, 5

$a[j+1] = \text{hole}$

I/p (9) | (16 5 0 8 2) ←
 Ist Sorted / Unsorted. hole
↓
0

IInd hole = 16
 9 | (?) 5 0 8 2
 6 | 9 5 0 8 2
 (?) 9 5 0 8 2
 Reached j = 0; then j = -1, while loop fails.
 → 6 9 | (5) 0 8 2

IIIrd hole = 5
 6 ? 9 0 8 2
 ? 6 9 0 8 2
 Reached j = 0; then j = -1;
 5 6 9 | (0) 8 2

IVth Same process like IIIrd hole = 0
 0 5 6 ? | (8) 2

hole = 8
 0 5 6 ? 9 2
 0 5 6 8 9 2

V^{th}

When $6 > 8$; Condition fails.

$a[j] > \text{hole}$

So no need to check for preceding elements.

0 5 6 8 9 | 2

hole = 2

0 5 6 8 ? 9

0 5 6 ? 8 9

$V1^{th}$

0 5 ? 6 8 9

~~0~~ ? 5 6 8 9

No w $0 > 2$; Condition fails, So no need to

$a[j] > \text{hole}$

Check for preceding elements.

0 2 5 6 8 9

0 2 5 6 8 9

Merge Sort

* External Sorting

Merge Sort (Arr[], int l, int r)

$$\{ \rightarrow \text{ if } (l \leq r)$$
$$\text{int } m = \left\lfloor \frac{l+r}{2} \right\rfloor$$

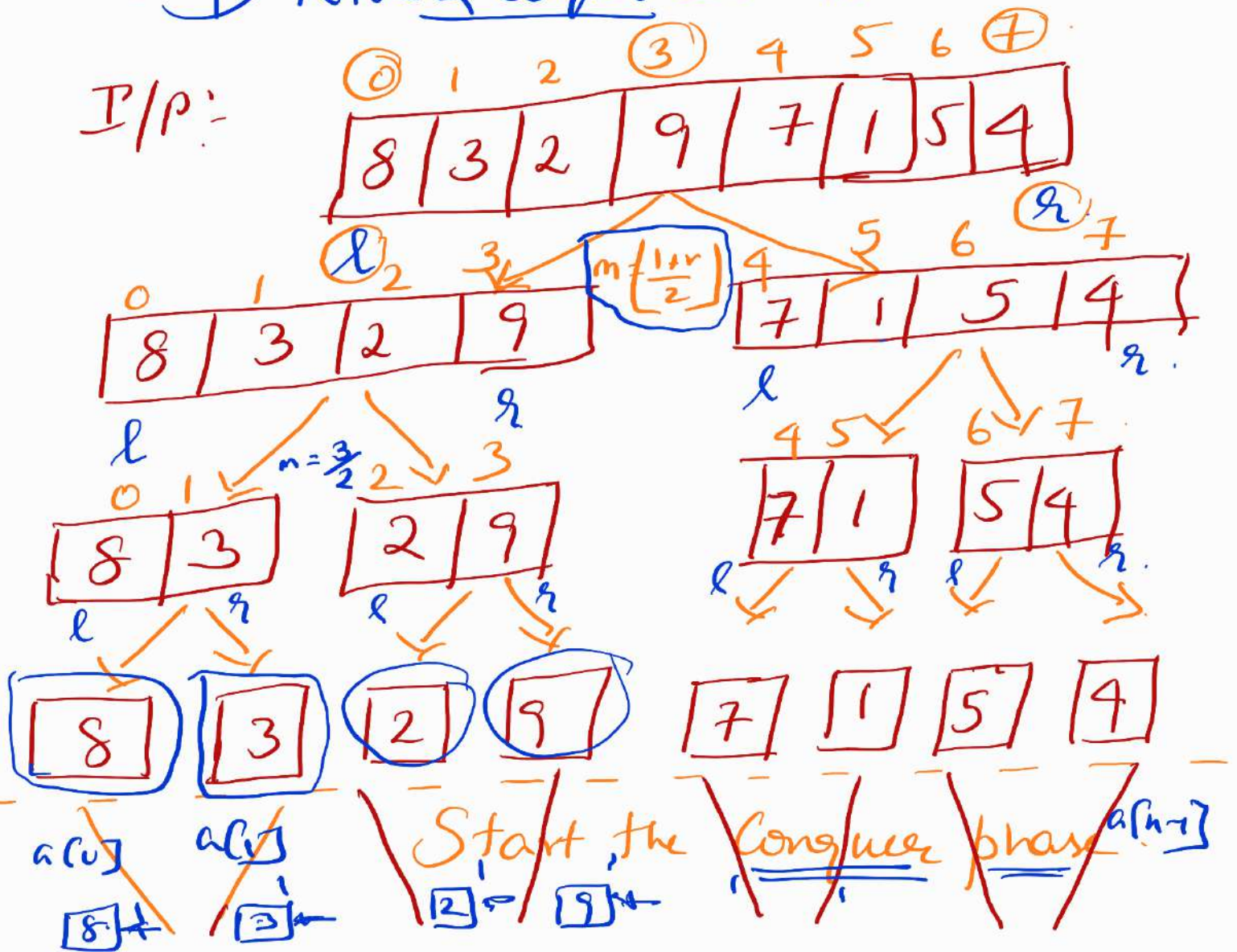
Divide \rightarrow $\left\{ \begin{array}{l} \text{MergeSort (Arr, l, m)} \leftarrow \text{left} \\ \text{MergeSort (Arr, m+1, r)} \leftarrow \text{right} \end{array} \right.$

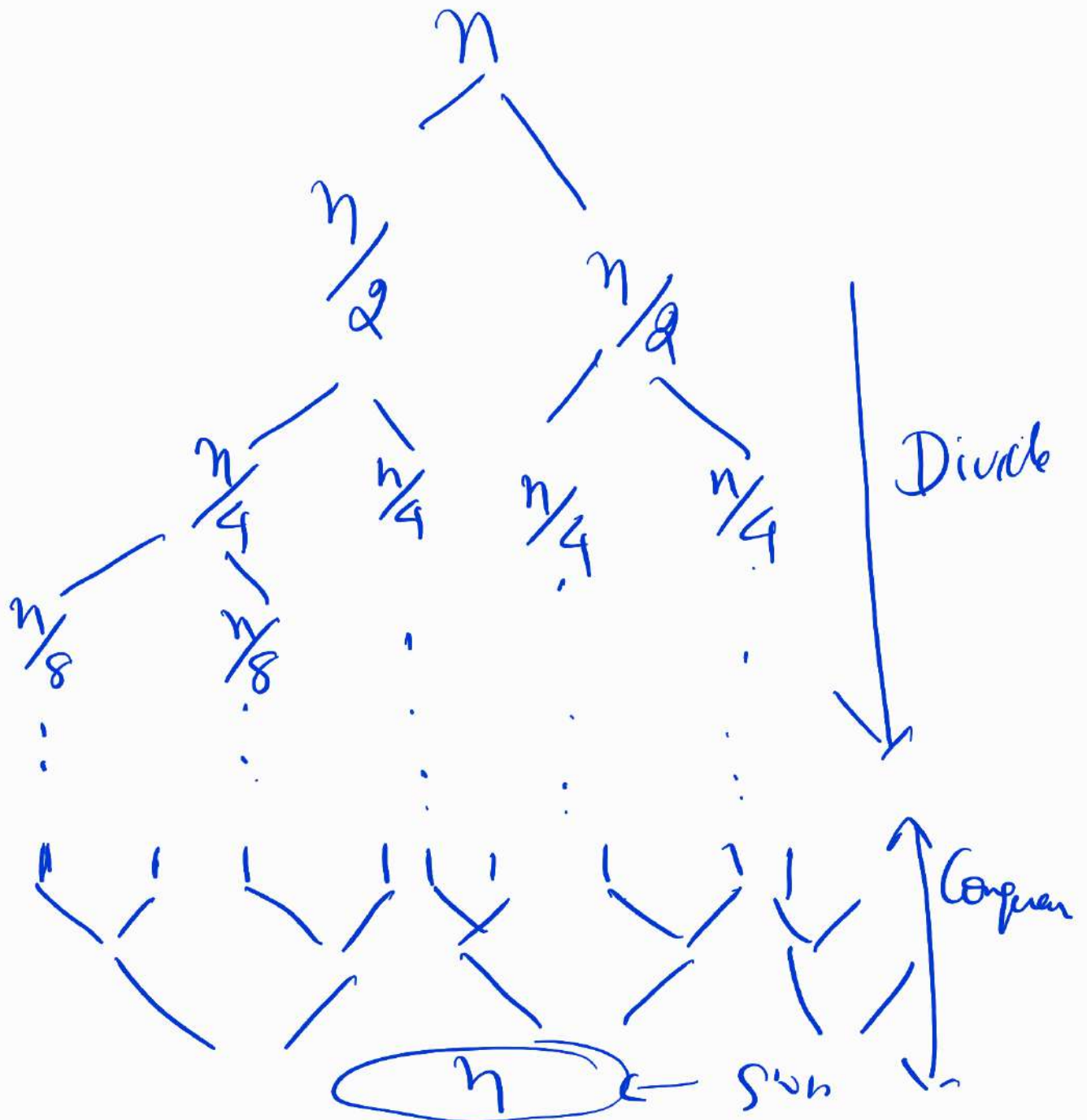
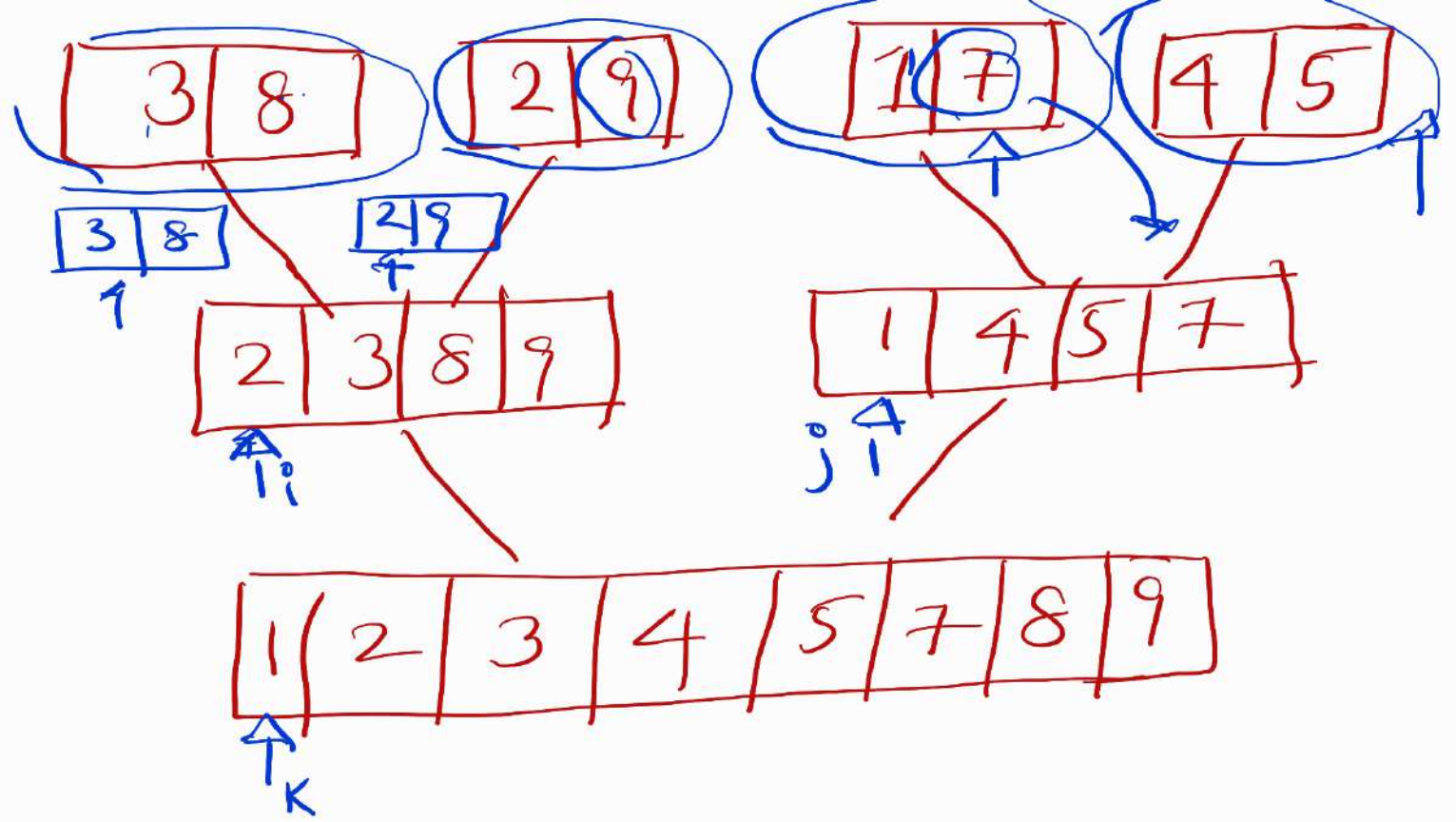
Conquer \leftarrow Merge (Arr, l, m, r) \leftarrow Conquer

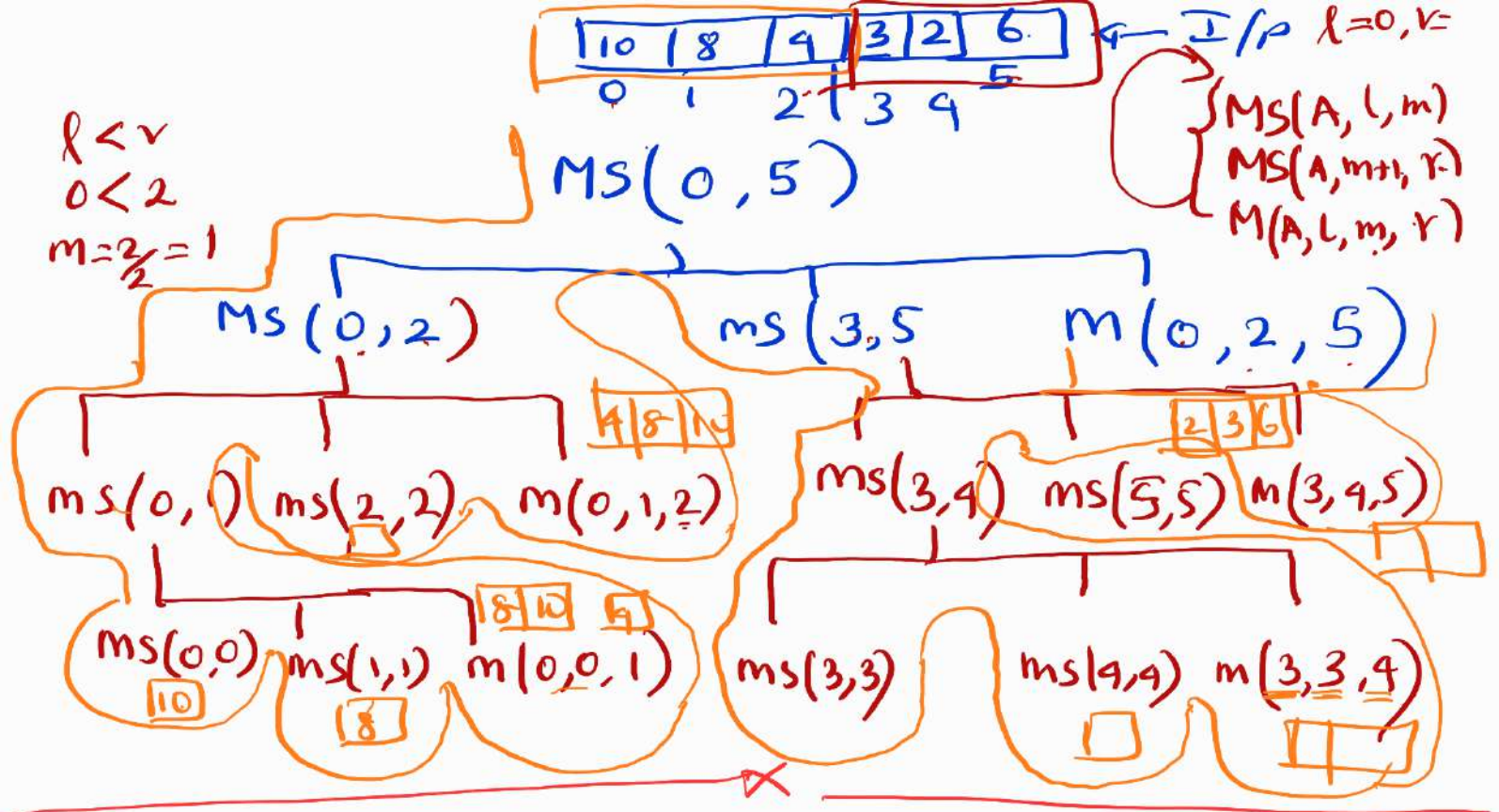
MergeSort (Arr, m+1, r) \leftarrow right

Conquer \leftarrow Merge (Arr, l, m, r) \leftarrow Conquer

3 Divide & Conquer Strategy

$$I/\rho:$$






void merge (int arr[], int l, int m, int r)

int i, j, k;

int n1 = m - l + 1; // Size of left temp array

int n2 = r - m; // Size of right temp array

From main array
put it into two
temporary array
 $L[n1], R[n2]$

for (i = 0; i < n1; i++)
 $L[i] = Arr[l + i];$
 for (j = 0; j < n2; j++)
 $R[j] = Arr[m + 1 + j];$

i = 0;

j = 0;

k = l;

while (i < n1 && j < n2) {

if (L[i] < R[j]) {

arr[k] = L[i];

i++;

k++;

}

else {

arr[k] = R[j];

j++;

k++;

}

while (i < n1) {

arr[k] = L[i];

i++;

k++;

while (j < n2) {

arr[k] = R[j];

j++;

k++;

}

void mergeSort(int arr[], int l, int r)

{

if (l < r)

{

int m = $\lfloor \frac{l+r}{2} \rfloor$

```
mergesort(arr, l, m);  
mergesort(arr, m+1, r);  
merge(arr, l, m, r);
```

}

}