

Big-oh Notation

'O'

 $f(n), g(n)$ Non-negative \downarrow
Step Count

$$f(n) = O(g(n)) \text{ iff } \forall n \geq n_0 \exists c \geq 0 (f(n) \leq c \cdot g(n))$$

$$f(n) = O(g(n)) \text{ iff } \exists n_0 \geq 0 \exists c \geq 0 \forall n \geq n_0 (f(n) \leq c \cdot g(n))$$

$$f(n) = 3n + 6 \quad 3n + 6 \leq 5n \quad c = 5, \forall n \geq 3, n_0 = 3 \Rightarrow 3n + 6 = O(n)$$

$$3n + 6 \leq 15n \quad c = 15, n_0 = 1, \forall n \geq 1 \Rightarrow 3n + 6 = O(n)$$

$$f(n) = n^5 + 5n - 2 \quad n^5 + 5n - 2 \leq 10 n^5 \quad c = 10, n_0 = 1, \forall n \geq 1 \Rightarrow n^5 + 5n - 2 = O(n^5)$$

$$n^5 + 5n - 2 \leq 3 n^5 \quad c = 3, n_0 = 2, \forall n \geq 2 \Rightarrow n^5 + 5n - 2 = O(n^5)$$

$$n^5 + 5n - 2 \leq 3 n^3 \quad c = 3, n_0 = 2, \forall n \geq 2 \Rightarrow O(n^3)$$

$$n^5 + 5n - 2 \leq 2 \cdot n^{10} \quad c = 2, n_0 = 2, \forall n \geq 2 \Rightarrow O(n^{10})$$

$$n^5 + 5n - 2 \leq 2^n \quad c = 1, n_0 = 7, \forall n \geq 7 \Rightarrow O(2^n)$$

Big-oh Notation

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$f(n), g(n)$

$f(n) = O(g(n))$ iff $\exists c > 0 \exists n_0 > 0 \forall n > n_0 (f(n) \leq c \cdot g(n))$

$$3n^2 + 4n - 2 \leq 6n^2 \quad c=6, n_0=3, \forall n \geq 3$$

$$\leq 25n^2 \quad c=25, n_0=6, \forall n \geq 6$$

$$\underline{n_0=1 \quad c=1}$$

$$3n^2 + 4n - 2 \leq 1 \cdot n^2$$

$$c=1, n_0=2 \Rightarrow 3 \cdot 4 + 4 \cdot 2 - 2 \leq 4 \quad \times$$

$$n_0=3 \Rightarrow 3 \cdot 9 + 4 \cdot 3 - 2 \leq 9 \quad \times$$

$$c=4, n_0=1 \Rightarrow 3n^2 + 4n - 2 \leq 4 \cdot n^2$$

$$\hookrightarrow 3 + 4 - 2 \leq 4 \quad \times$$

$$n_0=2 \Rightarrow 3 \cdot 4 + 4 \cdot 2 - 2 \leq 4 \cdot 4 \quad \checkmark$$

$$3n^2 + 4n - 2 \not\leq 0.5 n^2$$

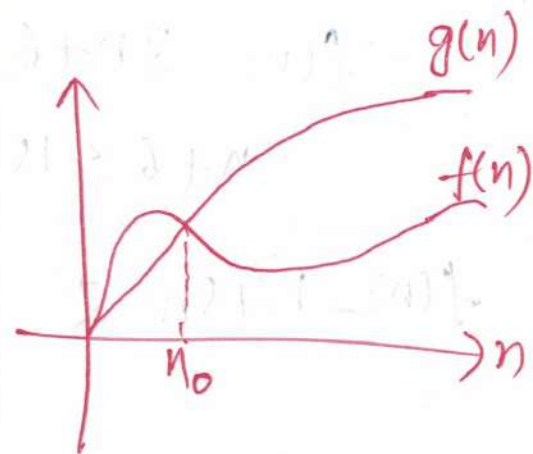
for any n_0

$$3n^2 + 4n - 2 \not\leq 2 n^2$$

$$3n^2 + 4n - 2 \leq (3+\epsilon)n^2$$

for some

$\epsilon > 0 \ \& \ n_0 > 0$



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Big-oh Notation

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$$3n^v + 4n - 2 \leq 3 \cdot n^3 \\ \leq 10 \cdot 2^n$$

$$c=3, n_0=4, \forall n \geq 4$$

$$c=10, n_0=3, \forall n \geq 3$$

$$O(n^{v+\epsilon}) = 3n^v + 4n - 2 = O(n^v) \\ \epsilon \geq 0$$

$$O(c^n), c > 1 = O(n^v) \\ = O(n^3)$$

$$O(1 \cdot n) = O(n^0)$$

$$O(1.001^n) = O(2^n)$$

$$IS \ 3n^v + 4n - 2 = O(n)??$$

$$3n^v + 4n - 2 \leq c \cdot n$$

$$3n + 4 - \frac{2}{n} \leq c$$

$$\exists c > 0 \ \exists n_0 > 0 \ \forall n \geq n_0$$

$$c = 10^9 \quad n_0 = 10^{12} \quad \forall n \geq 10^{12}$$

$$3n + 4 - \frac{2}{n} \neq 10^9$$

$$3n^v + 4n - 2 \neq O(n)$$

$$\neq O(n^{1.5})$$

$$\neq O(n^{2-\epsilon}) \\ \epsilon > 0$$

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Big-Oh Notation

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$$5n^3 + 6n \cdot 2^n - 100 \leq 10 \cdot n \cdot 2^n \quad c=10 \quad n_0=4$$

$$= O(n \cdot 2^n)$$

$$= O(n^3 \cdot 2^n)$$

$$= O(n \cdot 3^n)$$

$$= O(n^{2+\epsilon} \cdot (2+\delta)^n)$$

$$\epsilon \geq 0, \delta \geq 0$$

$$\neq O(2^n)$$

$$\neq O(n \cdot 2^n)$$

$$= O(3^n)$$

$$= O(2^n)$$

$$= O(2.001^n)$$

$$n \log n + 10 = O(n \log n)$$

$$= O(n^2)$$

$$= O(n^{1.1})$$

$$= O(2^n)$$

$$\neq O(n)$$

$$= O(n \log n \log n)$$

$$\neq O(n \log \log n)$$

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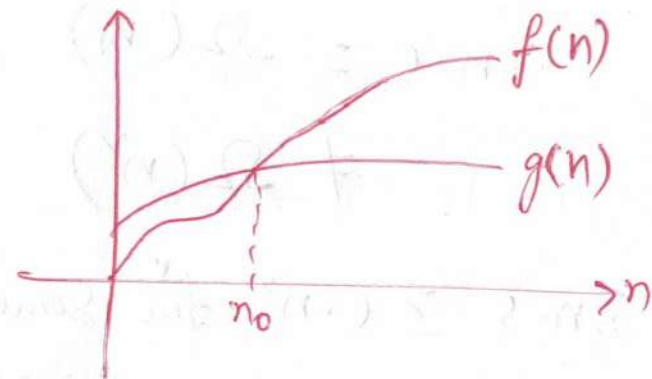
Big-Omega Notation

$$f(n) = \Omega(g(n)) \text{ iff } \exists c > 0 \exists n_0 > 0 \forall n \geq n_0 (f(n) \geq c \cdot g(n))$$

$$3n-5 \geq 2 \cdot n \quad c=2, n_0=5 \Rightarrow 3n-5 = \Omega(n)$$

$$3n-5 \geq 0.5n$$

$$3n-5 \geq 2 \log_2 n \Rightarrow \Omega(\log_2 n)$$



$$4n^{\sqrt{e}} - n + 2 \geq 2 \cdot n^{\sqrt{e}} \quad c=2, n_0=2$$

$$= \Omega(n^{2-\epsilon})$$

$$\epsilon \geq 0$$

$$= \Omega(n^{\sqrt{e}})$$

$$\geq 0.001 n^{\sqrt{e}} \Rightarrow \Omega(n^{\sqrt{e}})$$

$$\geq 10 \cdot n \Rightarrow \Omega(n)$$

$$\geq 5\sqrt{n} \Rightarrow \Omega(\sqrt{n})$$

$$\geq 6 \cdot n^{1.9} \Rightarrow \Omega(n^{1.9})$$

$$= \Omega(\log n)$$

$$= \Omega(n \log n)$$

Big-Omega Notation

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$$f(n) = \Omega(g(n)) \text{ iff } \exists c > 0 \exists n_0 > 0 \forall n \geq n_0 (f(n) \geq c \cdot g(n))$$

$$3n-5 = \Omega(n)$$

$$\neq \Omega(n^2)$$

$$3n-5 \geq c \cdot n^2 \text{ for some } c > 0 \\ \forall n \geq n_0$$

$$\frac{3}{n} - \frac{5}{n^2} \geq c \quad \text{Suppose } c = 10^{-6} \\ \text{choose } n_0 = 10^9 \\ n \geq n_0$$

The above inequality
is false

$$3n-5 \neq \Omega(2^n)$$

$$\neq \Omega(n \log n)$$

$$6n^2 - n + 100 = \Omega(n^2)$$

$$= \Omega(n)$$

$$= \Omega(\log n)$$

$$= \Omega(n^{\epsilon}) \quad \epsilon > 0$$

$$\neq \Omega(n^3)$$

$$\neq \Omega(n^{2+\delta}) \quad \delta > 0$$

$$\neq \Omega(c^n) \quad c > 1$$

Tight Bounds (Theta Notation)

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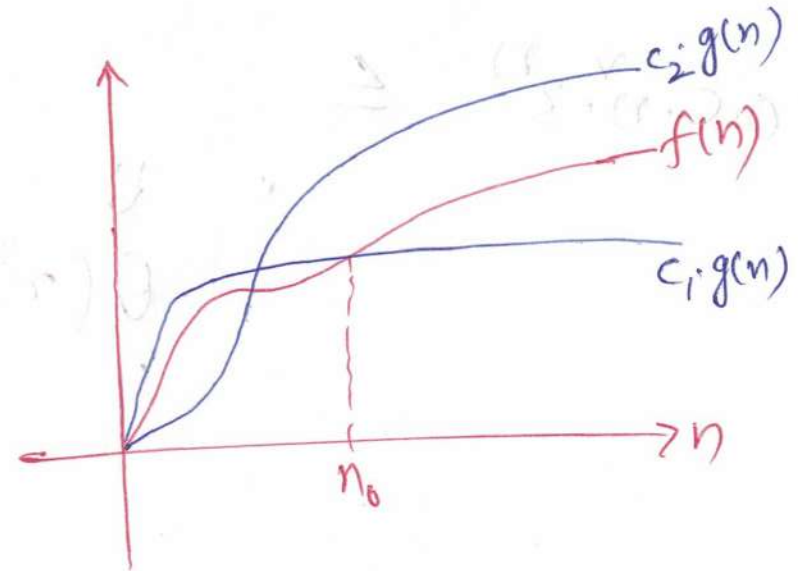
$f(n), g(n)$

$f(n) = \Theta(g(n))$ iff $\exists c_1 > 0 \exists c_2 > 0 \exists n_0 > 0 \forall n \geq n_0 (c_1 g(n) \leq f(n) \leq c_2 g(n))$

$f(n) = \Theta(g(n))$ iff $f(n) = \Omega(g(n))$
and
 $f(n) = O(g(n))$

$$2n \leq 3n+2 \leq 4n \quad \forall n \geq 2 \quad c_1=2 \quad c_2=4$$
$$\Rightarrow 3n+2 = \Theta(n)$$

$$\underbrace{0.5 n^2}_{\Omega(n^2)} \leq \underbrace{n^2 - 3n + 100}_{O(n^2)} \leq 4n^2 \quad \forall n \geq 6 \quad c_1=0.5 \quad c_2=4$$
$$\Rightarrow n^2 - 3n + 100 = \Theta(n^2)$$



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Tight Bounds (Theta Notation)

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$$2^n + n^{100} + n \cdot 3^n$$

$$0.5 \cdot n \cdot 3^n \leq$$

$$\leq 6 \cdot n \cdot 3^n$$

$$\Downarrow \\ \theta(n \cdot 3^n)$$

$$2^n = O(3^n)$$

$$\neq O(1.5^n)$$

$$2^n = O(n \cdot 3^n)$$

$$n^{100} = O(3^n)$$

$$n^{100} = O(n \cdot 3^n)$$

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Insight into Asymptotic notation

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$$3n+5 = O(n)$$

$$= O(n^2)$$

$$= O(n^3)$$

$$3n+5 = O(n^2)$$

$$3n+5 \leq c \cdot n^2 \quad \exists c > 0 \quad \exists n_0 > 0 \quad \forall n \geq n_0$$

$$3n+5 \leq c \cdot n \cdot n \quad \forall c > 0 \quad \exists n_0 > 0 \quad \forall n \geq n_0$$

$$3n+5 \leq 0.01 n^2 \quad c=0.01, n_0=1000$$

$$3n+5 \leq 0.0001 n^2 \quad n_0=10^6$$

work for $\rightarrow 3n+5 = O(n)$
some $c > 0$

work for $\rightarrow \begin{cases} 3n+5 = O(n^{1.1}) \\ \text{any } c > 0 \end{cases}$
 $= O(n^{1+\epsilon}) \quad \epsilon > 0$

Little oh [Loose upper bounds]

$$f(n) = o(g(n)) \text{ iff } \forall c > 0 \quad \exists n_0 > 0 \quad \forall n \geq n_0 (f(n) \leq c \cdot g(n))$$

$$3n+5 \neq o(n)$$

$$= O(n^{1.1})$$

$$= O(n^2)$$

$$= O(n^{1+\epsilon}) \quad \epsilon > 0$$

$$3n+5 \neq o(n)$$

for $c=0.001 \quad 3n+5 \not\leq 0.001 n$

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Little oh [Loose upper bounds]

$$\begin{aligned}n^{\sqrt{}} + 6n + 5 &= O(n^{\sqrt{}}) \\&= O(n^{2+\epsilon}), \epsilon > 0 \\&= O(c^n), c > 1\end{aligned}$$

$$\begin{aligned}n^{\sqrt{}} + 6n + 5 &\leq c \cdot n^{2.1}, \forall c > 0 \exists n_0 > 0 \forall n \geq n_0 \\&< 0.01 n^{2.1} \quad n_0 = 10^6\end{aligned}$$

$$\begin{aligned}n^{\sqrt{}} + 6n + 5 &\neq o(n^{\sqrt{}}) - \text{Little oh} \\&= o(n^{2+\epsilon}) \quad \epsilon > 0 \\&= o(2^n) \\&= o(n!) \\&= o(n^n)\end{aligned}$$

Little-Omega (Loose lower bounds)

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$$f(n) = \omega(g(n)) \text{ iff } \forall c > 0 \exists n_0 > 0 \forall n \geq n_0 (f(n) \geq c \cdot g(n))$$

$$n^v + 2^n - n = \Omega(n) \} = \omega(n)$$

$$= \Omega(n^v) \} = \omega(n^v)$$

$$= \Omega(n^3) \} = \omega(n^3)$$

$$= \Omega(2^n) \neq \omega(2^n)$$

$$\neq \omega(2^n)$$

$$c = 0.01 \checkmark$$

$$c = 10^4 \times$$

$$n^3 + 4n^v - n = \Omega(n^3)$$

$$= \Omega(n^{3-\epsilon}), \epsilon \geq 0$$

$$= \omega(n^{3-\epsilon}), \epsilon > 0$$

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Remarks

$$n^3 + 2n + 5 = O(n^3)$$

$$= O(n^3)$$

$$= O(n^{2+\epsilon}), \epsilon \geq 0$$

$$n^3 + 2n + 5 = \Omega(n^3)$$

$$= \Omega(n)$$

$$= \Omega(n^{2-\epsilon}), \epsilon \geq 0$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} \leq c, \quad c: \text{Constant}$$

$$\Rightarrow f(n) = O(g(n))$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} \geq c, \quad c: \text{Constant}$$

$$\Rightarrow f(n) = \Omega(g(n))$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = c, \quad c: \text{Constant}$$

$$\Rightarrow f(n) = \Theta(g(n))$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = 0 \Rightarrow f(n) = o(g(n)) \quad \text{small oh.}$$

$$\lim_{n \rightarrow \infty} \frac{f(n)}{g(n)} = \infty \Rightarrow f(n) = \omega(g(n))$$

① Reflexivity \downarrow reflexive property

$$f(n) = O(f(n))$$

$$f(n) = \Omega(f(n))$$

$$f(n) = \Theta(f(n))$$

$$n = O(n) \quad n \log n = O(n \log n)$$

$$n = O(n) \quad = O(n \log n)$$

$$n = \Omega(n) \quad = \Omega(n \log n)$$

② Symmetric

$$f(n) = O(g(n))$$

$$\nRightarrow g(n) = O(f(n))$$

$$f(n) = \Omega(g(n))$$

$$\nRightarrow g(n) = \Omega(f(n))$$

$$n = O(n^2)$$

$$n^2 \neq O(n)$$

$$n = \Omega(1)$$

$$1 \neq \Omega(n)$$

Θ is symmetric

$$f(n) = \Theta(g(n)) \Rightarrow g(n) = \Theta(f(n))$$

③ Transitive

$$f(n) = O(g(n)), g(n) = O(h(n))$$

$$\Rightarrow f(n) = O(h(n))$$

$$f(n) = \Omega(g(n)), g(n) = \Omega(h(n))$$

$$\Rightarrow f(n) = \Omega(h(n))$$

$$f(n) = \Theta(g(n)), g(n) = \Theta(h(n))$$

$$\Rightarrow f(n) = \Theta(h(n))$$

$$n^v = O(n^4), n^4 = O(2^n) \Rightarrow n^v = O(2^n)$$

$$n^v = \Omega(n \log n), n \log n = \Omega\left(\frac{1}{n}\right) \Rightarrow n^v = \Omega\left(\frac{1}{n}\right)$$

④ Transpose Symmetry

$$f(n) = O(g(n)) \Rightarrow g(n) = \Omega(f(n))$$

$$f(n) = \Omega(g(n)) \Rightarrow g(n) = O(f(n))$$

$$n^3 = O(2^n) \Rightarrow 2^n = \Omega(n^3)$$

$$n^3 = \Omega(n^v) \Rightarrow n^v = O(n^3)$$