

Facoltà di scienze informatiche

Operating Systems Spring - 2022

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Project 5 - User programs II

PintOS

In this project you are asked to implement two system calls: *wait* and *exec*. To implement *exec*, you need synchronization and also you need to check the pointers passed by the user which requires some understanding of memory paging.

The wait() system call

The calling process blocks until its child p, which has id pid, has finished. Here is the syscalls declaration:

• int wait (pid_t pid)

The return value must be:

- the exit status of p, if all was fine (even if p finished before wait was called!);
- -1 if p was killed;
- -1 if p is not a direct child of the caller (this makes it easier);
- -1 if the calling process already waited for that *pid* before.

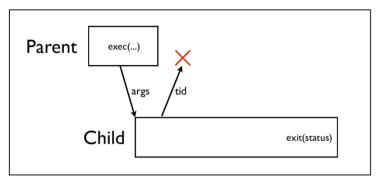
The exec() system call

A user program can use this system call to execute a command by creating a child process. Here is the systcall's declaration:

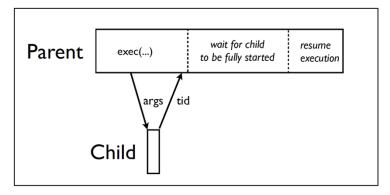
• pid_t exec (const char *cmd_line)

It tries to run the executable given in cmd_line . The return value is the pid (tid) of the child process (or -1 in case of error). A correct implementation synchronize between parent and child processes in order to prevent the following corner cases:

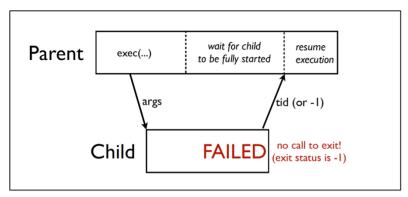
The parent thread must wait for the child creation:



The child thread may finish too fast:



The child thread may fail before telling some status:



Hints

The synchronization between parent and child process can be handled with *thread_block()*, or just a semaphore.

PintOS semaphore API

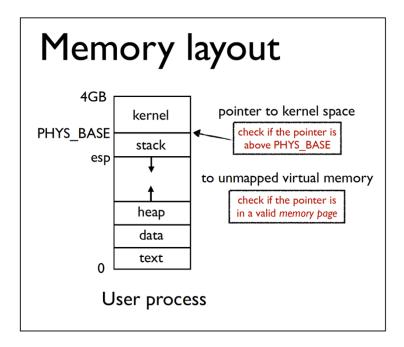
PintOS semaphore is a tool for synchronizing threads. It can solve the "wait for child" problem of the *exec* syscall. The API is available in "thread/synch.h".

- **struct semaphore** semaphore type; must be initialized with *sema_init*;
- sema_init (struct semaphore * s, unsigned val) initialize semaphore pointed by s with value val;
- sema_down (struct semaphore * s) if s is 0 (zero), block the calling thread and put it in a list waiting for s; otherwise, decrement s by 1;
- sema_up (struct semaphore * s) if s is 0 (zero) and there is some thread waiting for s, unblock one of the threads that are waiting for s; otherwise, increment s by 1;
- neither **sema_up** or **sema_down** can be interrupted (they're atomic)

Memory Access

The user may provide an invalid pointer in a syscall. It is important to check the pointer before using it. You should check if the pointer is:

- a. a null pointer;
- b. a pointer to kernel address space;
- c. a pointer to unmapped virtual memory; make sure user pointer is in a valid memory page (hints in **vaddr.h** and **pagedir.c**).



Tests

Your implementation should pass the tests:

- exec-once
- exec-arg
- $\bullet \ exec\text{-}multiple$
- $\bullet \ exec\text{-}missing$
- $\bullet \ exec\text{-}bad\text{-}ptr$
- $\bullet \ wait\text{-}simple$
- wait-twice
- wait-bad-pid

Readings

- PintOS documentation
 - Chapter 3 (Specially sections 3.1.4 and 3.3.4.)