

Transaction Processing

Agenda



- What is Transaction
- Transaction Model
- Isolation
 - Classification
- Savepoint and Release point
- Locking
 - Locking Levels
- Views and its types
 - Horizontal
 - Vertical
 - Group
- CHECK OPTION

What is Transaction?



- Transaction is a process of executing multiple DML statements sequentially to achieve a final task.
- A transaction which is performing DML operations including Insertion, deletion, update of records ensures *data integrity*.
- In an online application, the transactions happening on the database concurrently ensures the data is accurate and reliable without any loss.
- Usually , in big organizations like banks perform millions of transactions per minute in database environment. All these transactions are processed through DML statements including INSERT, DELETE, UPDATE and SELECT statements.
- All of these transactions ensure the database integrity and accuracy of data for users.

What is Transaction?



- A transaction with DML statements follows the ACID properties.
- *Atomicity*: It ensures the transaction is *fulfilled*; otherwise *fail* the transaction.
- E.g: Ordering a Pizza is fulfilled with series of transactions like -
 - Choose an item - check the inventory
 - Order the dish - insert order details
 - Complete the order - includes payment transaction.
- If the customer recalls the order at any stage, all the DML operations are rolled back including changes to inventory, orders and payments.

However, if the transaction is completed, you will enjoy your Pizza.

This file is meant for personal use by lokesh.jejappa@gmail.com only.

Sharing or publishing the contents in part or full is liable for legal action.

Consistency :

Any failure of transactions due to database crash or network failures, the original state of data is retained.

During crash recovery, SQL engine replay the transactions recorded in doublewrite buffer. Hence the *data is consistent* in the database without loss of transaction status.

It means it will have a record of the user SQL statements on the database, which are used for recovery process .

Isolation:

Multiple DML statements concurrently hit the database and isolating one transaction from other transaction is the responsibility of database.

So one individual transaction is isolated from other transactions. It is achieved only by forming the Queue of transactions.

E.g: Assume Withdrawal of amount from ATM & internet transfer are two different transactions occurred at the same second. Both transactions are transferring money from the account.

The two individual transactions are isolated , it means the database keeps the transactions in Queue to perform one transaction after the other to ensure the balance is sufficient to transfer the money.

SQL Transaction Model

- E.g: A transaction is identified with series of DML operations.

Transaction

- SELECT Operation -1 : Check bank balance
- INSERT Operation -2 : Withdraw money from ATM - new entry into transaction table.
- UPDATE Operation -3 : Update the balance by deducting with drawn money.
- INSERT Operation - 4 : Calculate the service charge on withdrawn amount and

This file is meant for personal use by lokesh.jejappa@gmail.com only.

Sharing or publishing the contents in part or full is liable for legal action.

insert a new entry in Transaction table



Start of Transaction



START TRANSACTION;

/ Check bank balance */*

```
Select Balance from ACCOUNT where Acct_Num = '4000-1956-2001' ;
```

```
Select Balance from ACCOUNT where Acct_Num = '4000-1956-2001' ;
```

Filter Rows:		Export:		Wrap Cell Content:	IA
Balance					
400000					

/ Bank charges 0.2% on withdrawn money */*

```
Insert Transaction values ('4000-1956-2001' , -2300.00 * 0.02, 'ATM  
Withdrawal' , 'CA' , now());
```

/ Update the old balance with new balance. */*

```
Update ACCOUNT set balance = balance - 2300 * 0.02 where Acct_Num =  
'4000-1956-2001' ;
```

COMMIT ;

This file is meant for personal use by lokesh.jejappa@gmail.com only.
Sharing or publishing the contents in part or full is liable for legal action.

Transaction



START TRANSACTION;

/ Check bank balance */*

Select Balance from ACCOUNT where Acct_Num = '4000-1956-2001' ;

```
Select Balance from ACCOUNT where Acct_Num = '4000-1956-2001' ;
```

Filter Rows: <input type="text"/>	Export:	Wrap Cell Content:
Balance		
400000		

/ Withdraw money from ATM */*

Insert Transaction values ('4000-1956-2001' , -2300.00, 'ATM Withdrawal' , 'CA' , now()) ;

```
Insert Transaction values ('4000-1956-2001' , -2300.00 , 'ATM Withdrawal' , 'CA' , current_date() );
```

```
Select * from Transaction where acct_num = '4000-1956-2001' and tran_amount = -2300.00
```

Filter Rows: <input type="text"/>					Export:	Wrap Cell Content:
Acct_Num	Tran_Amount	Channel	Area	Tran_Date		
4000-1956-2001	-2300.00	ATM Withdrawal	CA	2025-05-07		

This file is meant for personal use by lokesh.jejappa@gmail.com only.

Sharing or publishing the contents in part or full is liable for legal action.

Transaction



/ Bank charges 0.2% on withdrawn money */*

Insert Transaction values

```
('4000-1956-2001', -2300.00 * 0.02, 'ATM Withdrawal', 'CA' ,current_date());
```

Insert Transaction values

```
('4000-1956-2001' , -2300.00 * 0.02, 'ATM Withdrawal' , 'CA' ,current_date());
```

```
Select * from Transaction where acct_num = '4000-1956-2001' and tran_amount = -46
```

Filter Rows: <input type="text"/> Export: Wrap Cell Content:				
Acct_Num	Tran_Amount	Channel	Area	Tran_Date
4000-1956-2001	-46.00	ATM Withdrawal	CA	2020-05-07

This file is meant for personal use by lokesh.jejappa@gmail.com only.
Sharing or publishing the contents in part or full is liable for legal action.

End of Transaction



```
/* Update the old balance with new balance. */
```

```
Update ACCOUNT  
Set      balance = balance - 2300 * 0.02  
WHERE    Acct_Num = '4000-1956-2001' ;
```

```
/* End of transaction */  
COMMIT;
```

```
Update ACCOUNT set  balance = balance - 2300 * 0.02  where Acct_Num = '4000-1956-2001' ;  
Select * from  ACCOUNT where acct_num  =  '4000-1956-2001'  ;
```

Cust_Id	Acct_Num	Acct_Type	Balance	Acct_status	Relation
123002	4000-1956-2001	SAVINGS	397654	ACTIVE	P

This file is meant for personal use by lokesh.jejappa@gmail.com only.
Sharing or publishing the contents in part or full is liable for legal action.

Transaction Summary



In the above Start and End of transaction,

the transaction is atomic and completes the transaction of ATM withdrawal of money with some charges

- *COMMIT* ensures that transaction is fulfilled

- *ROLLBACK* the transaction, if customer ignores to withdraw the money from ATM, all of the above transactions are reversed from database .

The database will retain your original balance in the account.

Set Transaction



SET transaction controls performs the below tasks

- Encapsulate the DML statements
- Separates the execution of DML statements from DML statements of other transactions.
- Inter- lock the underlying affected records.

Note:

To ensure data integrity, each transaction should be consistent .

It means each DML statement should know the latest value of the underlying record though the record is affected by multiple transactions at same time.

Finally, the data Integrity is achieved only by below:

-Restrict other transactions performing the tasks on same records.

-Maintain the isolation of transactions with other transactions. So that other transactions wait in Queue to perform .

This file is meant for personal use by lokesh.jejappa@gmail.com only.

Sharing or publishing the contents in part or full is liable for legal action.

Set Transaction - Example



Example:

At 9:00 AM , One transaction tries to update the balance of an Account - 123456 and expected to complete by 9:05 AM.

At 9:02 AM : Meanwhile, another transaction updates the balance of same Account - 123456

Here, the functionality of transactions is:

The second transaction occurred at 9:02 AM will not have access to the Account - 123456 until 9:05 AM.

Hence, it will keep on waiting until the first transaction is completed and execute after 9:05 AM.

This functionality is said to be in ISOLATION which is set by MySQL, and will study more.

ISOLATION



MySQL - ISOLATION plays an important role in keep waiting of the transactions. This is done internally and users need not aware of it.

However, these ISOLATION levels can be manually altered by user and can modify such constraints and still achieve data integrity and performance.

Classification of Isolation levels



Classification of ISOLATION levels:

ISOLATION levels are classified as

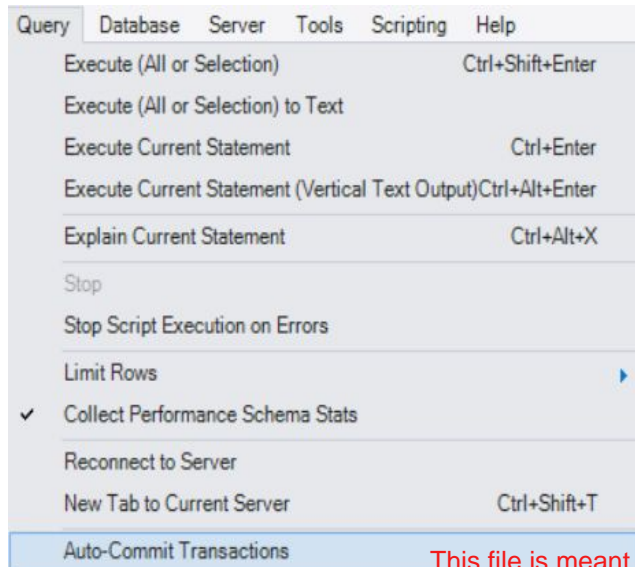
- REPEATABLE - READ
- READ COMMITTED
- SERIALIZABLE

Note:

Methods including READ , WRITE on transactions can also achieve Isolation.

Pre-requisite:

- Disable Auto- commit transaction



ISOLATION - Repeatable Read



- REPEATABLE READ: Lock the table in first Transaction affecting other Instances.

Perform below steps in "Instance - 1" and see result of each SQL statement.
(Ensure to *disable* Auto-commit)

```
/* MYSQL Workbench Instance - 1 */  
  
SET TRANSACTION ISOLATION LEVEL REPEATABLE READ;  
  
/*      Check bank balance      */  
Select Balance from ACCOUNT where Acct_Num =  
'4000-1956-2001' ;  
  
/*      Update balance      */  
Update ACCOUNT set      balance = balance - 2300  
where Acct_Num = '4000-1956-2001' ;  
  
/*      Check bank balance      */  
Select Balance from ACCOUNT where Acct_Num =  
'4000-1956-2001' ;
```

Result of Instance - 1

```
/*      Check bank balance      */  
Select Balance from ACCOUNT where Acct_Num = '4000-1956-2001' ;  
  
/*      Update balance      */  
Update ACCOUNT set      balance = balance - 2300 where Acct_Num = '4000-1956-2001' ;  
  
/*      Check bank balance      */  
Select Balance from ACCOUNT where Acct_Num = '4000-1956-2001' ;
```

A screenshot of the MySQL Workbench interface. At the top, there is a toolbar with icons for Filter Rows, Export, and Wrap Cell Content. Below the toolbar, a table is displayed with a single column labeled "Balance" and a single row containing the value "395354".

Balance
395354

ISOLATION - Repeatable Read



- REPEATABLE READ: Experience table level lock using UPDATE statement

```
/* MYSQL workbench , Instance - 2 */
```

```
/*      Check bank balance of different account */  
Select Balance  
from ACCOUNT  
where Acct_Num = '4000-1956-5698' ;
```

```
/*      Update balance */  
Update ACCOUNT  
set    balance = balance - 1500  
where Acct_Num = '4000-1956-5698' ;
```

Result of Instance - 2

```
Select Balance  
from ACCOUNT where Acct_Num = '4000-1956-5698' ;
```

Balance
397654

```
Update ACCOUNT  
set    balance = balance - 1500 where Acct_Num = '4000-1956-5698' ;
```

Time	Action	Message
0:18:18	INTERRUPT	OK - Query cancelled
0:18:30	Update ACCOUNT set balance = balance - 1500 where Acct_Num = '4000-1956-5698'	Running...

- In Instance - 2, the status of UPDATE keeps RUNNING.
- No additional records can be inserted. Because of session-1 acquired lock at table level.

This file is meant for personal use only. Sharing or publishing the contents in part or full is liable for legal action.

ISOLATION - Repeatable Read



- REPEATABLE READ: Experience table level lock using *INSERT statement*

Result of Instance - 2

```
/* MYSQL workbench , Instance - 2 */
```

```
/* Insert new record */
```

```
Insert into ACCOUNT  
Values ( 123001 , '4000-1956-9999' , 'SAVINGS',  
69000, 'ACTIVE' , 'P' ) ;
```

```
/* insert new record */
```

```
Insert into ACCOUNT
```

```
Values ( 123001 , '4000-1956-9999' , 'SAVINGS', 69000, 'ACTIVE' , 'P' ) ;
```

Output		
Time	Action	Message
10:23:30	Insert into ACCOUNT Values (123001 , '4000-1956-9999' , 'SAVINGS', 69000, 'ACTIVE' , 'P')	Error Code: 2013. Lost connec
10:25:33	Insert into ACCOUNT Values (123001 , '4000-1956-9999' , 'SAVINGS', 69000, 'ACTIVE' , 'P')	Running...

- In Instance - 2, the status of INSERT keeps RUNNING.
Because of session 1 has acquired the lock on the ACCOUNT table.

This file is result for personal use by Jakesh Jeyapalan@gmail.com only.
Sharing or publishing the contents in part or full is liable for legal action.

ISOLATION - Repeatable Read



- REPEATABLE READ: Experience table level lock with below practice, impacts other records.

```
/* MYSQL workbench , Instance - 2 */
```

```
/*      Check bank balance of different account */  
Select Balance  
from ACCOUNT  
where Acct_Num = '4000-1956-5698' ;
```

```
/*      Update balance */  
Update ACCOUNT  
set    balance = balance - 1500  
where Acct_Num = '4000-1956-5698' ;
```

Result of Instance - 2

```
Select Balance  
from ACCOUNT where Acct_Num = '4000-1956-5698' ;
```

Balance
397654

```
Update ACCOUNT  
set    balance = balance - 1500 where Acct_Num = '4000-1956-5698' ;
```

Time	Action	Message
0:18:18	INTERRUPT	OK - Query cancelled
0:18:30	Update ACCOUNT set balance = balance - 1500 where Acct_Num = '4000-1956-5698'	Running...

- In Instance - 2, the status of UPDATE keeps RUNNING.
- No additional records can be inserted. Because of session-1 acquired lock at table level.

ISOLATION - Read Committed



- READ COMMITTED : Experience record level lock with below practice.

```
/* MYSQL workbench - Instance - 1 */
```

```
SET TRANSACTION ISOLATION LEVEL READ  
COMMITTED
```

```
/*    Check bank balance    */  
Select Balance from ACCOUNT where  
Acct_Num = '4000-1956-2001' ;
```

```
/*    Update balance    */  
Update ACCOUNT set  
balance = balance - 2300 where  
Acct_Num = '4000-1956-2001' ;
```

```
/*    Check bank balance    */  
Select Balance from ACCOUNT where  
Acct_Num = '4000-1956-2001' ;
```

Result of Instance -1

```
/*    Check bank balance    */  
Select Balance from ACCOUNT where Acct_Num = '4000-1956-2001' ;  
  
/*    Update balance    */  
Update ACCOUNT set balance = balance - 2300 where Acct_Num = '4000-1956-2001' ;  
  
/*    Check bank balance    */  
Select Balance from ACCOUNT where Acct_Num = '4000-1956-2001' ;
```

Filter Rows:	Export:	Wrap Cell Content:
Balance		
393054		

READ Committed



- READ COMMITTED : Experience *record level lock only* with below practice.

Result of Instance - 2

```
/* MYSQL workbench Instance - 2 */  
  
/* Check bank balance of different account */  
Select Balance  
from ACCOUNT  
where Acct_Num = '4000-1956-5698' ;  
  
/* insert a new record */  
Insert into ACCOUNT  
Values ( '4000-1956-9999' , 'SAVINGS',  
69000, 'ACTIVE' , 'P' ) ;
```

The screenshot shows two queries executed in MySQL Workbench. The first query selects the balance for account '4000-1956-5698', returning 397654. The second query inserts a new record and then selects all records for account '4000-1956-9999', showing a single row with details: Cust_Id 123001, Acct_Num 4000-1956-9999, Acct_Type SAVINGS, Balance 69000, Acct_status ACTIVE, and Relation P.

Balance
397654

Cust_Id	Acct_Num	Acct_Type	Balance	Acct_status	Relation
123001	4000-1956-9999	SAVINGS	69000	ACTIVE	P

- In Instance - 2, the status of INSERT happens successfully.
However , the Instance -1 acquired lock at table level but restricted to existing records only.
Hence it allowed to insert new records from session -2.

Transaction with READ and WRITE



Transactions Modifiers:

READ and WRITE can also perform the isolation of transactions in the same Instance .

Set Transaction with default - Read Write



READ WRITE : It is a default transaction,
So any statement is executed in current Instance.

Result of READ WRITE transaction

```
/* Instance - 1 */  
  
SET TRANSACTION READ WRITE ;  
  
/*    Check bank balance    */  
Select Balance  
from ACCOUNT  
where Acct_Num = '4000-1956-2001' ;  
  
/*    Update balance    */  
Update ACCOUNT  
set    balance = balance - 2300  
where Acct_Num = '4000-1956-2001' ;
```

```
SET TRANSACTION READ WRITE ;  
/*    Check bank balance    */  
Select Balance  
from ACCOUNT where Acct_Num = '4000-1956-2001' ;  
/*    Update balance    */  
Update ACCOUNT  
set    balance = balance - 2300 where Acct_Num = '4000-1956-2001' ;
```

Output	
Time	Action
08:17:32	SET TRANSACTION READ WRITE
08:17:36	Select Balance from ACCOUNT where Acct_Num = '4000-1956-2001' LIMIT 0, 2000
08:17:42	Update ACCOUNT set balance = balance - 2300 where Acct_Num = '4000-1956-2001'

Message
0 row(s) affected
1 row(s) returned
1 row(s) affected Rows matched: 1 Changed: 1 Warnings: 0

UPDATE statement is executed. Since Read Write is a default transaction modifier.

This file is meant for personal use by lokesh.jejappa@gmail.com only.

Sharing or publishing the contents in part or full is liable for legal action.

Set Transaction with READ Only



In the user Instance, transactions are isolated with READ / WRITE / both modifiers.

Result of READ only transaction

READ Only :

```
SET TRANSACTION READ ONLY ;
```

```
/* Check bank balance */
```

```
Select Balance  
from ACCOUNT  
where Acct_Num = '4000-1956-2001' ;
```

```
/* Update balance */
```

```
Update ACCOUNT  
set balance = balance - 2300  
where Acct_Num = '4000-1956-2001' ;
```

```
SET TRANSACTION READ ONLY ;  
  
/* Check bank balance */  
Select Balance from ACCOUNT where Acct_Num = '4000-1956-2001' ;  
  
/* Update balance */  
Update ACCOUNT set balance = balance - 2300 where Acct_Num = '4000-1956-2001' ;
```

In Output		
Time	Action	Message
08:04:37	SET TRANSACTION READ ONLY	0 row(s) affected
08:04:42	Select Balance from ACCOUNT where Acct_Num = '4000-1956-2001' LIMIT 0, 2000	1 row(s) returned
08:04:50	Update ACCOUNT set balance = balance - 2300 where Acct_Num = '4000-1956-2001'	Error Code: 1792. Cannot execute statement in a READ ONLY transaction.

UPDATE statement throws an error :

Cannot execute statement in a READ ONLY transaction.

This file is meant for personal use by laksh.jeppa@gmail.com only.
Sharing or publishing the contents in part or full is liable for legal action.

Altering the Modifiers : Read and Write

Modifiers when set at Instance level can be altered, but it doesn't effect the current transaction.

However it will be effective only from next transaction.

```
SET SESSION TRANSACTION READ WRITE ;
```

```
/* In read and write mode */
```

```
START Transaction;
```

```
SET SESSION TRANSACTION READ ONLY ; /* Effective from next transaction */
```

```
Update ACCOUNT  
set  
balance = balance - 2300  
where Acct_Num = '4000-1956-2001';
```

```
/* DML is executed */
```

```
COMMIT ; /* End the current transaction */
```

This file is meant for personal use by lokeshe.jejappa@gmail.com only.
Sharing or publishing the contents in part or full is liable for legal action.

Altering the Modifiers : Read and Write

Result

```

SET SESSION TRANSACTION READ WRITE ;
/* In read and write mode */
START Transaction;
SET SESSION TRANSACTION READ ONLY ;      /* Effective only in second transaction */
/*      Accepting DML */
Update ACCOUNT set balance = balance - 2300 where Acct_Num = '4000-1956-2001';
COMMIT; /* End the transaction */

```

Output

Time	Action	Message
09:05:33	commit	0 row(s) affected
09:05:47	SET SESSION TRANSACTION READ WRITE	0 row(s) affected
09:06:23	START Transaction	0 row(s) affected
09:06:25	SET SESSION TRANSACTION READ ONLY	0 row(s) affected
09:06:30	Update ACCOUNT set balance = balance - 2300 where Acct_Num = '4000-1956-2001'	1 row(s) affected Rows matched: 1 Changed: 1
09:06:38	COMMIT	0 row(s) affected

This file is meant for personal use by lakesh.jejappa@gmail.com only.

Sharing or publishing the contents in part or full is liable for legal action.

Set Transaction



Alter Access modifiers : READ/ WRITE in user SESSION.

/ Second transaction: */*

START transaction

/ Update will not execute */*

Update ACCOUNT

set

balance = balance - 2300

where Acct_Num = '4000-1956-2001' ;

```
SET SESSION TRANSACTION READ WRITE ;
/* In read and write mode */
START Transaction;
SET SESSION TRANSACTION READ ONLY ; /* Effective only in second transaction */
/* Accepting DML */
Update ACCOUNT set balance = balance - 2300 where Acct_Num = '4000-1956-2001';
COMMIT; /* End the transaction */
```

```
START transaction ;
/* Update balance but 1 record effected */
Update ACCOUNT set balance = balance - 2300 where Acct_Num = '4000-1956-2001' ;
```

In Output		
Time	Action	Message
09:06:25	SET SESSION TRANSACTION READ ONLY	0 row(s) affected
09:06:30	Update ACCOUNT set balance = balance - 2300 where Acct_Num = '4000-1956-2001'	1 row(s) affected Rows matched: 1 Changed: 1 Warnings: 0
09:06:38	COMMIT	0 row(s) affected
09:27:15	START transaction	0 row(s) affected
09:27:31	Update ACCOUNT set balance = balance - 2300 where Acct_Num = '4000-1956-2001'	Error Code: 1792. Cannot execute statement in a READ ONLY transaction.

This file is meant for personal use by lokesh.jejappa@gmail.com only.
Sharing or publishing the contents in part or full is liable for legal action.

SAVEPOINT and RELEASE SAVEPOINT statement

- So far, ROLLBACK is issued on entire transaction. SQL engine facilitates to do partially rollback and do final commit.
- SAVEPOINT statement is used to mark a portion of transaction with an identifier in a full length of transaction.
- These checkpoints are later referred by ROLLBACK to undo partial transactions.

SAVEPOINT



- Savepoints are so useful for recovery of errors in database applications.
- Each Savepoint is labelled with a unique name to identify the location of intermediate portion of transaction.

E.g:

```
SAVEPOINT Label_1;  
SAVEPOINT Label_2;
```

- Using these Labels, user can undo the portion of the transaction.

E.g:

```
ROLLBACK TO Label_1;
```

SAVEPOINT



```
SELECT *  
FROM   CUSTOMER  
WHERE  CUST_ID in  
       (123002, 123003 , 123004 )
```

Cust_Id	Name	Address	State	Phone
123002	George	194-6,New brighton	MN	189761700
123003	Harry	2909-5,walnut creek	CA	1897617866
123004	Jack	229-5, Concord	CA	1897627999

Try:

```
START TRANSACTION;  
SAVEPOINT customer_1;  
    DELETE FROM CUSTOMER WHERE  
CUST_ID=123004;  
SAVEPOINT customer_2;  
    DELETE FROM CUSTOMER WHERE  
CUST_ID=123003;  
SAVEPOINT customer_3;  
    DELETE FROM CUSTOMER WHERE  
CUST_ID=123002;
```

<pre>START TRANSACTION; SAVEPOINT customer_1; DELETE FROM CUSTOMER WHERE CUST_ID=123004; SAVEPOINT customer_2; DELETE FROM CUSTOMER WHERE CUST_ID=123003; SAVEPOINT customer_3; DELETE FROM CUSTOMER WHERE CUST_ID=123002;</pre>		
Output		
Time	Action	Message
09:54:00	SAVEPOINT customer_1	0 row(s) affected
09:54:00	DELETE FROM CUSTOMER WHERE CUST_ID=123004	1 row(s) affected
09:54:00	SAVEPOINT customer_2	0 row(s) affected
09:54:00	DELETE FROM CUSTOMER WHERE CUST_ID=123003	1 row(s) affected
09:54:00	SAVEPOINT customer_3	0 row(s) affected
09:54:00	DELETE FROM CUSTOMER WHERE CUST_ID=123002	1 row(s) affected

This file is meant for personal use by lokesh.jejappa@gmail.com only.
Sharing or publishing the contents in part or full is liable for legal action.

SAVEPOINT



```
/* Continue */
```

```
ROLLBACK TO customer_2;
```

```
COMMIT;
```

```
SELECT *  
FROM   CUSTOMER  
WHERE  CUST_ID in  
       (123002, 123003 , 123004 )
```

Cust_Id	Name	Address	State	Phone
123002	George	194-6,New brighton	MN	189761700
123003	Harry	2909-5,walnut creek	CA	1897617866

- *Note : Upon commit the transaction, savepoints defined in the transaction are lost.*

This file is meant for personal use by lokesh.jejappa@gmail.com only.

Sharing or publishing the contents in part or full is liable for legal action.

RELEASE SAVEPOINT



- Savepoints that are created can be deleted in a transaction without affecting the transaction.
- This is called releasing the individual savepoint from the total set of savepoints.
- If SAVEPOINT has been released, you can no longer use the ROLLBACK command to undo transactions performed since the earlier SAVEPOINT.

Let us see an example of below:

```
SELECT * FROM CUSTOMER WHERE CUST_ID in ( 123007, 123006, 123005 )
```

Cust_Id	Name	Address	State	Phone
123005	Jacob	325-7, Mission Dist	SFO	1897637000
123006	Noah	275-9, saint-paul	MN	1897613200
123007	Charlie	125-1, Richfield	MN	1897617666

RELEASE SAVEPOINT

Example:



```
START TRANSACTION;
```

```
SAVEPOINT customer_1;  
DELETE  
FROM CUSTOMER  
WHERE CUST_ID =123007;
```

```
SAVEPOINT customer_2;  
DELETE  
FROM CUSTOMER  
WHERE CUST_ID =123006;
```

```
SAVEPOINT customer_3;  
DELETE  
FROM CUSTOMER  
WHERE Cust_Id=123005;
```

```
RELEASE SAVEPOINT customer_2;
```

```
ROLLBACK to customer_2;  
ROLLBACK to customer_1;
```

Result

```
START TRANSACTION;  
SAVEPOINT customer_1;  
DELETE FROM CUSTOMER WHERE CUST_ID =123007;  
SAVEPOINT customer_2;  
DELETE FROM CUSTOMER WHERE CUST_ID =123006;  
SAVEPOINT customer_3;  
DELETE FROM CUSTOMER WHERE Cust_Id =123005;  
RELEASE SAVEPOINT customer_2;  
ROLLBACK TO customer_2;  
ROLLBACK TO customer_3;
```

Output		
Time	Action	Message
14:46:28	DELETE FROM CUSTOMER WHERE CUST_ID =123007	1 row(s) affected
14:46:28	SAVEPOINT customer_2	0 row(s) affected
14:46:28	DELETE FROM CUSTOMER WHERE CUST_ID =123006	1 row(s) affected
14:46:28	SAVEPOINT customer_3	0 row(s) affected
14:46:28	DELETE FROM CUSTOMER WHERE Cust_Id =123005	1 row(s) affected
14:46:28	RELEASE SAVEPOINT customer_2	0 row(s) affected
14:46:28	ROLLBACK TO customer_2	Error Code: 1305. SAVEPOINT customer_2 does not exist
14:46:42	ROLLBACK TO customer_3	Error Code: 1305. SAVEPOINT customer_3 does not exist

This file is meant for personal use by lokesh.jeyappa@gmail.com only.
Sharing or publishing the contents in part or full is liable for legal action.

RELEASE SAVEPOINT



Rollback issued before RELEASE will be successful.

E.g: **ROLLBACK** TO *Customer_1* ;

```
START TRANSACTION;
SAVEPOINT customer_1;
  DELETE FROM CUSTOMER WHERE CUST_ID =123007;
SAVEPOINT customer_2;
  DELETE FROM CUSTOMER WHERE CUST_ID =123006;
SAVEPOINT customer_3;
  DELETE FROM CUSTOMER WHERE Cust_Id =123005;
RELEASE SAVEPOINT customer_2;
ROLLBACK TO customer_2;
ROLLBACK TO customer_3;
ROLLBACK TO customer_1 ;
```

Output		
Time	Action	Message
14:46:28	SAVEPOINT customer_2	0 row(s) affected
14:46:28	DELETE FROM CUSTOMER WHERE CUST_ID =123006	1 row(s) affected
14:46:28	SAVEPOINT customer_3	0 row(s) affected
14:46:28	DELETE FROM CUSTOMER WHERE Cust_Id =123005	1 row(s) affected
14:46:28	RELEASE SAVEPOINT customer_2	0 row(s) affected
14:46:28	ROLLBACK TO customer_2	Error Code: 1305. SAVEPOINT customer_2 does not exist
14:46:42	ROLLBACK TO customer_3	Error Code: 1305. SAVEPOINT customer_3 does not exist
14:50:33	ROLLBACK TO customer_1	0 row(s) affected

This file is meant for personal use by lokesh.jejappa@gmail.com only
Sharing or publishing the contents in part or full is liable for legal action.

COMMIT Statement

COMMIT



- COMMIT ends the current transaction.
- It permanently saves the transaction effected changes to the physical data files and make others visible to the data changes.
- After the COMMIT has been issued, the current transaction will exit and savepoint will disappear.
- There is no option of database that will "undo" the transaction after committing the data.

COMMIT



Example:

```
SELECT *  
FROM ACCOUNT  
WHERE Acct_status = 'INACTIVE'  
AND (balance = 0 or balance is NULL ) ;
```

Cust_Id	Acct_Num	Acct_Type	Balance	Acct_status	Relation
123007	9000-1700-7777-4321	CREDITCARD	0	INACTIVE	P
123008	5000-1700-7755	SAVINGS	NULL	INACTIVE	P

```
DELETE  
FROM ACCOUNT  
WHERE Acct_status = 'INACTIVE'  
AND (balance = 0 or balance is NULL);
```

COMMIT;

This file is meant for personal use by lokesh.jejappa@gmail.com only.
Sharing or publishing the contents in part or full is liable for legal action.

ROLLBACK



- In a current running transaction which is not yet physically committed to database, the transactional changes can be rolled back and retain the original status of the data.
- ROLLBACK plays a great roll during interactive operations when user tries to restrain from modifying the data.

ROLLBACK



Delete few rows and try to ROLLBACK to retain the original data.

```
DELETE FROM CUSTOMER WHERE CUST_ID =123006
```

```
select * from CUSTOMER;
```

CUST_ID	NAME	ADDRESS	STATE_CODE	TELEPHONE
123001	Oliver	225-5, Emeryville	452	1897614500
123002	George	194-6, New brighton	875	189761700
123003	Harry	2909-5, walnut creek	232	1897617866
123004	Jack	229-5, Concord	501	1897627999

ROLLBACK



Performing ROLLBACK

```
ROLLBACK;
```

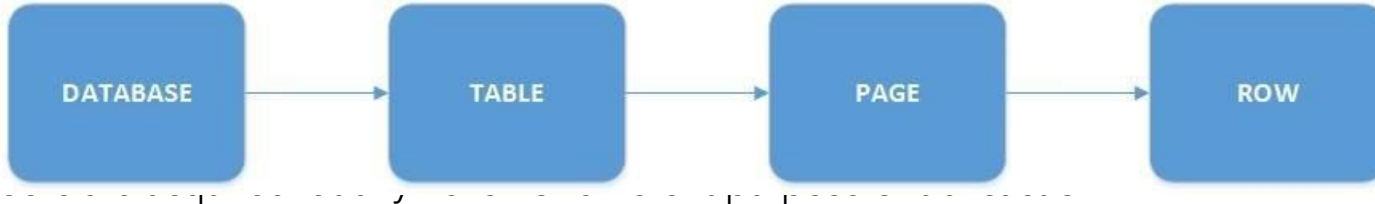
```
select * from CUSTOMER;
```

CUST_ID	NAME	ADDRESS	STATE_CODE	TELEPHONE
123006	Kiran		532	9542434422
123006	Kiran		532	9542434422
123001	Oliver	225-5, Emeryville	452	1897614500
123002	George	194-6, New brighton	875	189761700
123003	Harry	2909-5, walnut creek	232	1897617866
123004	Jack	229-5, Concord	501	1897627999

Locking

Lock

- LOCKS are acquired by user transaction in a session.
- SQL Engine locks objects when the transaction begins the DML statements on each record.
- To ensure data integrity, DML statements acquire explicit locks on data records or tables.
- Multiple users in an integrated environment can see changes applied by each other.
- After
- Lock The



Locking Levels

Locking Levels



- Several types of locks are used for concurrency that each lock is specific to a transaction and does not affect other transaction.
- Below are the system resources on which Locks are acquired.
- RID: (Row ID) : Row-id is a unique address of individual record in a table. DML statements when executed will try to isolate the records from other transactions using this rowid. Hence the row-ids are locked by DML statements, and then perform the operations.
- Table: A transaction with DML statements lock the entire table and corresponding indexes. This locks are more reliable in batch processing when performing huge volume of data, and restricts users to access the tables. So that the performance will not be degraded in terms of taking snapshots of data.
- Key: Indexes are generated by default when primary Key is created on a table. Hence the locks are acquired on these keys .

Locking Levels



Memory level locks : Database memory is allocated with pages and pages forms an extent. Locks do occur at these stages.

Page: A page consists of 8-kilobyte (KB) used by a data file or index . Acquire the locks at page level and extent level.

Extent: Similarly an extents can be locked which consists of contiguous 8 data pages.

Database: Entire Database can be locked for maintenance purpose .

Shared and Exclusive Locks

Shared and Exclusive Locks



- Every transaction executes DML statements needs to acquire *Row-level locking*.
- Here are the two types.
 - Shared lock or Read Lock
 - Exclusive lock or Write Lock

Shared Locks



- Shared lock is an initial lock acquired by a SELECT statement of a transaction to lock the rows for applying next DML statements.
- Shared lock is useful to avoid deadlocks of multiple transactions if they are using same DML statements.
- However, the other transactions can also acquire multiple shared locks (S) on the same records locked by first transaction.

Shared Locks between two SELECT Queries



For example: Consider the following two sessions

SHARED lock is acquired by two sessions without any error.

Instance - 1

```
START TRANSACTION;
```

```
SELECT * FROM ACCOUNT  
WHERE Acct_Num = '4000-1956-2001'  
LOCK IN SHARE MODE
```

```
START TRANSACTION;
```

```
SELECT * FROM ACCOUNT
```

```
WHERE Acct_Num = '4000-1956-2001' LOCK IN SHARE MODE
```

Filter Rows: <input type="text"/> Export: Wrap Cell Contents:					
Cust_Id	Acct_Num	Acct_Type	Balance	Acct_status	Relation
123002	4000-1956-2001	SAVINGS	383854	ACTIVE	P
Output					
Time	Action	Message			
15:28:54	START TRANSACTION	0 row(s) affected			
15:28:54	SELECT * FROM ACCOUNT WHERE Acct_Num = '4000-1956-2001' LIMIT 0, 2000 LOCK IN SHARE MODE	1 row(s) returned			

This file is meant for personal use by lokesh.jeappa@gmail.com only.

Sharing or publishing the contents in part or full is liable for legal action.

Shared Locks between two SELECT queries




Shared lock is also acquired by Instance -2 .

Instance - 2


```
START TRANSACTION;  
  
SELECT *  
FROM ACCOUNT  
WHERE Acct_Num = '4000-1956-2001'  
LOCK IN SHARE MODE
```

```
START TRANSACTION;  
SELECT * FROM ACCOUNT  
WHERE Acct_Num = '4000-1956-2001' LOCK IN SHARE MODE
```




Filter Rows:

Export:



Wrap Cell Content:



Cust_Id	Acct_Num	Acct_Type	Balance	Acct_status	Relation
123002	4000-1956-2001	SAVINGS	383854	ACTIVE	P

Shared Locks with Update statement



Session - 1

```
START TRANSACTION;
```

```
SELECT * FROM ACCOUNT  
WHERE Acct_Num = '4000-1956-2001'  
LOCK IN SHARE MODE;
```

```
UPDATE ACCOUNT  
Set balance =  
balance - balance * 0.04  
WHERE Acct_Num = '4000-1956-2001'
```

```
START TRANSACTION;  
SELECT * FROM ACCOUNT  
WHERE Acct_Num = '4000-1956-2001' LOCK IN SHARE MODE;  
/* Update after */  
UPDATE ACCOUNT  
Set balance = balance - balance * 0.04  
WHERE Acct_Num = '4000-1956-2001'
```

Output

Time	Action	Message
15:37:21	START TRANSACTION	0 row(s) affected
15:37:21	SELECT * FROM ACCOUNT WHERE Acct_Num = '4000-1956-2001' LIMIT 0, 2000 LOCK IN SHARE MODE	1 row(s) returned
15:37:33	UPDATE ACCOUNT Set balance = balance - balance * 0.04 WHERE Acct_Num = '4000-1956-2001'	1 row(s) affected Rows matched: 1 Changed: 1 Warnings: 0

Shared Locks



Session - 2

```
START TRANSACTION;
```

```
SELECT *  
FROM ACCOUNT  
WHERE  
Acct_Num = '4000-1956-2001'  
LOCK IN SHARE MODE
```

```
START TRANSACTION;
```

```
SELECT * FROM ACCOUNT
```

```
WHERE Acct_Num = '4000-1956-2001' LOCK IN SHARE MODE
```

Output

Time	Action	Message
15:39:52	START TRANSACTION	0 row(s) affected
15:39:52	SELECT * FROM ACCOUNT WHERE Acct_Num = '4000-1956-2001' LIMIT 0, 2000 LOCK IN SHARE MODE	Running...

Running / Awaiting :

Since the DML statement is executed in Instance -1 for same record, a shared lock is occurred in Instance - 2 .

Instance - 2 is in Queue and gains the lock on the record after instance - 1 activity is done.

This file is meant for personal use by lokesh.jeppa@gmail.com only.

Sharing or publishing the contents in part or full is liable for legal action.

Exclusive (X) Locks



- Exclusive locks are acquired by DML statements: DELETE, INSERT and UPDATE.
- Data cannot be modified at a time by two DML statements. If they then its a deadlock or one transaction keeps waiting for other transaction.
- Hence this lock is mandatorily acquired by any DML statement by default when executed to modify the records.
- This lock cannot be shared with any other transactions as they have to keep waiting until the first DML transaction completes and commits its changes to database.
- Note: Once the Exclusive lock (X) is acquired , other transactions gain neither of the Shared (S) nor Exclusive (X) Locks

Exclusive (X) Locks



Session - 1 /* acquired shared lock */

```
START TRANSACTION
SELECT * FROM ACCOUNT WHERE Acct_Num = '4000-1956-2001' LOCK IN SHARE MODE
```

Session - 2 /* acquired shared lock */

```
START TRANSACTION
SELECT * FROM ACCOUNT WHERE Acct_Num = '4000-1956-2001' LOCK IN SHARE MODE
```

Session - 1 /* acquired shared lock and exclusive lock */

```
START TRANSACTION
SELECT * FROM ACCOUNT WHERE Acct_Num = '4000-1956-2001' LOCK IN SHARE MODE
UPDATE ACCOUNT Set balance = balance - balance * 0.04 WHERE Acct_Num = '4000-1956-2001'
```

Session - 2 /* acquired shared lock but did not gain exclusive lock (waiting) */

```
START TRANSACTION
SELECT * FROM ACCOUNT WHERE Acct_Num = '4000-1956-2001' LOCK IN SHARE MODE
UPDATE ACCOUNT Set balance = balance - balance * 0.04 WHERE Acct_Num = '4000-1956-2001'
```

Session - 3 /* Neither acquires Shared lock(S) nor Exclusive (X) lock */

```
START TRANSACTION
SELECT * FROM ACCOUNT WHERE Acct_Num = '4000-1956-2001' LOCK IN SHARE MODE
UPDATE ACCOUNT Set balance = balance - balance * 0.04 WHERE Acct_Num = '4000-1956-2001'
```

This file is meant for personal use by lokesh.jejappa@gmail.com only.

Sharing or publishing the contents in part or full is liable for legal action.

Views

Virtual Tables



- A virtual table is a *derived form of data* from a physical database tables.
- Virtual tables are logically represented using physical data but *do not store* any data.
- A virtual table is a construct of record set that are not actually tables. so that these record sets are referenced in SQL statements like normal tables.
- Virtual tables otherwise called as views and are stored as database objects in MYSQL Information_schema, and can be retrieved without re-writing it again.

What is a view?

What is view?



- A view is a database object that is created using a Select Query with complex logic.
- So views are said to be a logical representation of physical data
- Views behave like a physical table and users can use them as database objects in any part of SQL queries
- DML operations can be performed on base tables via views.
- Since they are stored in the data dictionary , they can be retrieved easily

Types of views: Simple View



- Simple views are created with a select query written using a single table
- DML operations are easily performed just like performing on tables

```
Create VIEW Simple_view as  
Select * from CUSTOMER ;
```


Complex View:



- Complex view is created using a Select query written with multiple tables using JOINS, sub-queries, with clauses and conditionally filtered

```
Create VIEW Complex_view as
SELECT  c.Cust_Id ,
        a.ACCOUNT
From CUSTOMER c
JOIN ACCOUNT a
Where c.Cust_Id = a.Cust_Id
And    a.balance > 300000
```

Inline View:



- A subquery is also called as an inline view if and only if it is called in FROM clause of a SELECT query
- It is not stored in any data dictionary like other views. It is dynamically defined in queries

```
SELECT * FROM  
( SELECT  c.Cust_Id , ba.ACCOUNT  
  From    CUSTOMER bc  
JOIN ACCOUNT  ba  
Where c.Cust_Id = ba.Cust_Id  
And ba.balance > 300000)
```

Materialized View:



- Unlike views, Materialized views stores the results of a complex SQL queries
- These are so useful when the SQL engine takes time in execution of the SQL query logic and retrieve the results
- Materialized views are also called as Snapshots of database tables

Materialized View:



- These views are refreshed with new data with scheduled intervals
- Materialized views are mostly used in data marts where there is a need for handling large volume of data that refers to history data

Data marts is a repositories of summarized data. This data is collected for analysis on a specific section or unit within an organization, for example, the sales department.

Materialized View:



```
create MATERIALIZED VIEW M_VW as SELECT  c.Cust_Id ,  
ba.ACCOUNT  
From CUSTOMER c  
JOIN ACCOUNT    ba  
Where bc.Cust_Id = ba.Cust_Id  
And  ba.balance > 99999
```

Refresh Methods



- Fast & Complete
- A fast refresh uses a log that compares the history and recent changes to its database tables and captures only the changes
- Complete refresh truncates the whole materialized view data and refreshes with new changed data
- Complete refresh is slow when compared with fast refresh method

How the RDBMS handles views

How RDBMS handles views



- Views are stored in data dictionary. It means once they are created , they are used by several users at a time in different sessions

```
Select * from ALL_VIEWS where view_name in ('SIMPLE_VIEW'  
, 'COMPLX_VIEW')
```


How RDBMS handles views



- The underlying data of base tables dynamically reflects in views since it logically represents the physical data
- However, any updates to data is visible to views only when it is committed
- E.g. Open two different SQL sessions and test Simple_View
- Step1: In one session, issue the below:

```
Select * from Simple_view where Cust_Id = 123001
```

How RDBMS handles views



- Step2: In another session, issue the below (no auto commit/ commit)
- Update CUSTOMER set telephone = 9701499878 where Cust_Id = 123001
- Step3: Come back to first session , and again check the view to see the data doesn't reflects the updated telephone

```
Select * from Simple_view where Cust_Id = 123001
```

How RDBMS handles views



- Step4: Now go back to second session , and issue commit on Update DML statement

Commit;

- Step5: Once again come back to first session, and check the updated telephone number.

```
Select * from Simple_view where Cust_Id = 123001
```

How RDBMS handles views



- Views are invoked in FROM clauses of SELECT statements and avoids the complexity of debugging and writing large code

E.g:

```
Select  CA_vw.Cust_Id,  
        bat.Acct_Num,  
        bat.Tran_Amount,  
        CA_vw.balance
```

```
from
```

```
Transaction bat
```

```
JOIN
```

```
Complex_view  CA_vw
```

```
ON bat.Acct_Num = CA_vw.Acct_Num
```

This file is meant for personal use by lokesh.jejappa@gmail.com only.

Sharing or publishing the contents in part or full is liable for legal action.

How RDBMS handles views



- In the previous example, the query looks so simple as *Complex_view* has taken care of additional logic of conditions and joining CUSTOMER table, and avoided too much of parent Query logic.

How RDBMS handles views



- Views avoids direct access to physical data and protects from unauthorized users
- Ex: If any user fires a SELECT query directly on physical table while the data is being loaded into base tables at the same time
- Here RDBMS takes a before and after images of table data each time when updating is happening to data

How RDBMS handles views



- The after image data is then stored to physical disk
- During this time, if any user tries to do a SELECT on physical data , it will show the data from before image since the data is not committed
- Before images are however deleted once the after image data is committed
- Hence there is a lot of overhead to system to maintain before images for multiple users to ensure data integrity

Advantages and Disadvantages of views

Advantages:



- Views are stored queries and re-used
- They are accessible by multiple users in different sessions
- Scalability is high
- They hide the confidential columns data like SSN, date of birth, Address , telephone

Advantages:



- Avoids direct access to physical data
- Views protect the data when views are invoked by downstream users to build web applications that are prone to hackers
- Materialized views are useful in a centralized environment where multiple users access it in batches of data rather than streaming data

Disadvantages:



- When a base table of an associated view is dropped, it becomes obsolete
- Queries using views are bit slow compared to accessing data directly from physical tables
- Because , SQL engine will take an additional processing of semantics checking of view each time when queries are running on it

Disadvantages:



- Views created using aggregate functions are restricted to use in WHERE clauses
- Materialized views stores physical data hence it holds redundant data and results in consumption of additional space for its Mviews and Logs
- Views are in-efficient when joined with remote database tables as they consume lot of I/O transactions between local and remote database tables

Horizontal View

Horizontal View



- A Horizontal view dynamically represents the data without aware of the columns
- The view does not necessary to know about the number of columns and its data types of the respective base tables
- However the base tables are explicitly defined in the view

Example



```
Create view Horizontal_view as  
Select ba.*  
From ACCOUNT  where balance > 99999
```

Example



- In this previous example, the wildcard - "*" references all of the columns belongs to base tables mentioned in the view
- Any changes to the column names , data types doesn't affect the view creation

Vertical View

Vertical View



- A vertical view built with predefined columns , however the changes to column data types does not affect the view unless they are type-casted or applied with any system functions.

```
CREATE VIEW Vertical_view AS
SELECT Cust_Id,
       Acct_Num,
       Balance,
       Acct_Type,
       Acct_status,
       Relation
FROM ACCOUNT
```

Vertical View



- In this example, the columns that are explicitly mentioned while creating the view and those are same as base tables
- Otherwise it throws an error if any changes applied to columns

customer_id	account_number	balance_amount	account_type	Account_status	Relation_ship
123002	4000-1956-2001	400000	SAVINGS	ACTIVE	P
123003	4000-1956-2900	750000	SAVINGS	INACTIVE	P
123004	4000-1956-3401	655000	SAVINGS	ACTIVE	P
123001	4000-1956-3456	200000	SAVINGS	ACTIVE	P

Horizontal view Vs Vertical view



Horizontal views

Columns are represented with “*” card

Query execution is dynamic and re-calculate the execution plan each time because columns are unknown.

Maintenance is easy as the view is not worried about changes to columns

Vertical views

Columns are explicitly defined

Underlying query execution plan is not repeated because view columns are in sync with base tables.

View maintenance is required and need to be adjusted according to base table changes.

Vertical View



Horizontal views

Users have access to all view columns like accessing all base table columns.

Users can apply any data type conversions while using the views.

Views can be updated easily.

Vertical views

Many of the columns are hidden.

Any data type conversions can be applied unless the other functions inside view definition are not compatible.

These views can also be updated unless any functions are used on columns in view.

Row Column subset view

Row column subset View



- A View can be created to represent the subset of records and lesser columns in order to filter the records and hide the columns

```
CREATE VIEW Subset_view AS
SELECT
    Cust_Id,
    Acct_Num,
    Balance
FROM ACCOUNT
Where Balance > 700000 ;
```

Row column subset View



- In this example, only few columns are displayed that are defined in the subset view. The other columns are hidden
- Also, the records are filtered for which the balance is less than or equal to 700000

customer_id	account_number	balance_amount	account_type	Account_status	Relation_ship
123002	4000-1956-2001	400000	SAVINGS	ACTIVE	P
123003	4000-1956-2900	750000	SAVINGS	INACTIVE	P
123004	4000-1956-3401	655000	SAVINGS	ACTIVE	P
123001	4000-1956-3456	200000	SAVINGS	ACTIVE	P

This file is meant for personal use by lokesh.jejappa@gmail.com only.

Sharing or publishing the contents in part or full is liable for legal action.

Grouped Views

Group View



- Grouping views are created using a SELECT query with grouping functions applied on multiple rows of data
- The grouping functions are dynamically applied on rows of data and displays aggregate results only

Group View



```
Create view Grouping_view
(Acct_Num, total_transaction)
as
select  Acct_Num ,
        sum(Tran_Amount )
from Transaction
group by Acct_Num ;
```

Group View



account_number	total_transaction
4000-1956-2001	8400.00
4000-1956-3401	8000.00
4000-1956-3456	-2000.00
4000-1956-5102	-6500.00
4000-1956-5698	9000.00
4000-1956-9977	229000.00
5000-1700-6091	40000.00
5000-1700-7791	40000.00
5800-1700-9800-7755	-18000.00
5890-1970-7706-8912	-26000.00
5900-1900-9877-5543	-20000.00
9000-1700-7777-4321	-15500.00

In this example, only the aggregate transaction amount is displayed

Joined Views

Joined View



- These views represent the data from more than one base tables joined using key columns
- The underlying query can have JOINS or subqueries on base tables

Joined View



```
Create View Joined_view
As
Select c.Cust_Id,
       c.Name,
       a.Acct_Num,
       a.Acct_Type,
       a.Balance,
       c.telephone
FROM ACCOUNT  a
JOIN CUSTOMER c
ON a.Cust_Id = c.Cust_Id ;
```

Joined View



customer_id	customer_name	account_number	account_type	balance_amount	telephone
123002	George	4000-1956-2001	SAVINGS	400000	1897617000
123003	Harry	4000-1956-2900	SAVINGS	750000	1897617866
123004	Jack	4000-1956-3401	SAVINGS	655000	1897627999
123001	Oliver	4000-1956-3456	SAVINGS	200000	1897614500
123005	Jacob	4000-1956-5102	SAVINGS	300000	1897637000

Joined View



- In this example, the underlying query is joined with two tables and selected their respective columns
- From the view point, the representation of columns are viewed as single entity without any complexity

CHECK OPTION

WITH CHECK OPTION



- Simple views can be queried and allow DML operations
- In order to restrict users to perform DML operations on sensitive columns, "WITH CHECK OPTION" is optionally used as a constraint

Example



```
Create view View_with_Check  
As Select * From ACCOUNT  
Where Balance > 5000  
WITH CHECK OPTION ;
```

```
Insert into View_with_Check values  
( 123001, '4000-7843-3002', 'SAVINGS', 4000, 'ACTIVE',  
'P')
```

Example



```
#1369 - CHECK OPTION failed 'tantu.view_with_check'
```

- In this example, you can insert any balance amount without the CHECK OPTION constraint irrespective of the condition in the view

WITH CHECK OPTION (LOCAL)



- LOCAL CHECK option pertains to only local conditions of the view and ignores the conditions of dependent views

```
Create view normal_view as  
Select * from ACCOUNT  
where Balance >= 80000;
```

```
create view local_view as  
select * from normal_view  
where Balance <= 90000  
with local check option;
```

WITH CHECK OPTION (LOCAL)



```
Insert into cascaded_view values  
( 123001, '6500-3823-4032', 'SAVINGS', 65000,  
'ACTIVE', 'P')
```

```
#1369 - CHECK OPTION failed 'tantu.view_with_check'
```

WITH CHECK OPTION (LOCAL)



- In this first example, you can bypass the condition and insert balance amount of 85000 in normal_view
- Local_view is created on normal_view will be local will validate its own conditions
- Hence the records with any balance less than 90k are accepted, and ignores the dependent view condition

WITH CHECK OPTION (CASCADED)



- A normal view conditions are bypassed when performing DML statements
- CASCADED CHECK option ensures to satisfy all of the conditions of dependent views

WITH CHECK OPTION (CASCADED)



```
Create view normal_view as  
Select * from ACCOUNT  
where Balance >= 80000;
```

```
create view cascaded_view as  
select * from normal_view  
where Balance =< 90000  
with cascaded check option;
```

```
Insert into cascaded_view values  
(123001, '5000-1253-4312', 'SAVINGS' , 70000,  
'ACTIVE','P')
```

WITH CHECK OPTION (CASCADED)



```
#1369 - CHECK OPTION failed 'tantu.view_with_check'
```

- In this first example, you can bypass the condition and insert balance amount of 85000 in normal_view
- Such cases, cascaded_view is created on normal_view to ensure that the conditions of both views are satisfied
- Hence the records with balance between 80k and 90k are only accepted

Drop view

Drop View



- Without affecting the base table data, any type of view can be dropped

E.g:

```
Drop view view_name;
```

Materialized view

Materialized views (Oracle)



- M-views are logical created using physical tables but it stores the data when they are executed
- Materialized views are executed frequently to refreshed with new data of base tables
- M-views uses mainly two refresh methods : Fast and Complete

Materialized views (Oracle)



```
Create materialized view mview_account  
As Select * From ACCOUNT  
Where Balance > 500000 ;
```

- In this example, materialized view is created with a SQL query logic but does not hold any physical data
- In order to feed the data to materialized views, materialized views must be refreshed

Materialized views (Oracle)



- Materialized views are refreshed with physical data. Let us look at how they are refreshed

E.g:

```
execute dbms_mview.refresh( 'mview_account' );
```

```
Select * from mview_account ;
```

Materialized views (Oracle)



Customer_id	Account_Number	Account_type	Balance_amount	Account_status	Relation_ship
123003	4000-1956-2900	SAVINGS	750000	INACTIVE	P
123004	4000-1956-3401	SAVINGS	655000	ACTIVE	P
123007	4000-1956-9977	RECURRING DEPOSITS	7025000	ACTIVE	S
123001	5000-1700-3456	RECURRING DEPOSITS	9400000	ACTIVE	S
123002	5000-1700-5001	RECURRING DEPOSITS	7500000	ACTIVE	S
123004	5000-1700-6091	RECURRING DEPOSITS	7500000	ACTIVE	S

- Mview_account is refreshed with data only after it is explicitly refreshed with system package

This file is meant for personal use by lokesh.jejappa@gmail.com only.
Sharing or publishing the contents in part or full is liable for legal action.

Complete Refresh:



- Materialized view uses complete refresh method to truncate and refresh with new data from Query it is defined with

Create materialized view

REFRESH COMPLETE

mview_account_RC

As Select * From ACCOUNT

Where Balance > 500000

execute dbms_mview.refresh(list => 'mview_account_RC' , method
=> 'C');

Complete Refresh:



- In this example, materialized view always uses complete method to reload the data every time it is refreshed explicitly
- Here the refresh time is equals to the time of query execution and loading of the data

FAST Refresh:



- Materialized view uses FAST refresh method with only changed data of the base query

Create materialized view fastrefresh_log on ACCOUNT
with sequence ;

Create materialized view

REFRESH FAST

mview_account_RFast

As Select * From ACCOUNT

Where Balance > 500000 ;

FAST Refresh:



```
execute dbms_mview.refresh( list => `
mview_account_RFast' , method => 'F');
```

- In this example, materialized view always uses FAST method to capture only updated or inserted data every time it is refreshed explicitly
- The changed data is always maintained in a log created on the base table which being used by fast refresh later

Thank You