

Database Systems

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Triết lý Giáo dục - Giá trị cốt lõi

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1.1. Learning Objective

- Define Database and Database Management System (DBMS).
- Study structure of Relation model in Database.
- Query Data by Relational Algebra language and SQL.
- Analyses Data Integrity Constrains.
- Python GUI for database.

1.2. Databases Are Everywhere

- Database = a large (?) collection of related data.
- Classically, a DB models a real-world organization (e.g., enterprise, university).
 - Entities (e.g., students, courses)
 - Relationships (e.g., “Martin is teaching DBs in 2018/09”)
- Changes in the organization = changes in the database.
- Examples:
 - Personnel records
 - Students management
 - Banking
 - Airline reservations

1.3. Scientific Databases (Examples)

- Biology:
e.g., DNA sequences of genes, amino-acid sequences of proteins, genes expressed in tissues (up to several Gigabytes)
- Astronomy:
e.g., location and spectra of astronomic objects (up to several Terabytes)
- Physics:
e.g., sensor measurements in particle physics experiments (up to several Petabytes)

1.4. Operations with Databases

- Design
Define structure and types of data.
- Construction
Create data structures of DB, populate DB with data.
- Manipulation of Data
 - Insert, delete, update
 - Query: “Which department pays the highest salary?”
 - Create reports:
“List monthly salaries of employees, organized by department, with average salary and total sum of salaries for each dept”

1.5. Database Management System (DBMS)

- DBMS is a software package designed to store and manage databases.
- Several brands, e.g.
 - Oracle (Oracle)
 - DB2 (IBM)
 - SQL Server, Access (Microsoft),
 - MySQL, PostgreSQL, HSQLDB, SQLite (open source)

Use a DBMS?

- Separation of the Data definition and the Program.
- Abstraction into the simple Model.
- Data independence and efficient access.
- Reduced application development time.
- Data integrity and Security.
- Uniform data administration.
- Concurrent access, recovery from crashes.
- Support for multiple different views.

1.6. Study Database?

- Databases used to be specialized applications; now they are a central component in most applications.
- Knowledge of database concepts is essential for computer scientists.
- Databases are everywhere, even when you don't see them.
- Data is valuable, e.g. include tax records, student records, bank account records, photos, ...
- Data is typically structured. e.g. tax records follow the same structure, bank records follow the same structure, ...
- Database field has made a number of contributions to the field of computer science.

2.1. Data Model

Data model is a notation for describing data or information. The description generally consists of three parts:

- **Structure of the data.** Familiar with programming languages such as C: arrays, structures, objects. In the database world, data models referred to as a conceptual model to emphasize the difference in level.
- **Operations on the data.**
 - Operations that retrieve information: Set of queries.
 - Operations that change the database: Set of modifications.
- **Constraints on the data.** The way to describe limitations on what the data can be.

Important Data Models

Today, there are three data models of importance for database systems:

- **Relational data model.** Including object-relational extensions
- **Semi-structured data model.** Including XML and related standards.
- **NoSQL data model.** Four types of NoSQL databases are Document-oriented, Key-Value Pairs, Column oriented and Graph.

2.2. Relational Data Model

Relational model gives a single way to represent data: as a two-dimensional table called a relation.

e.g.

ST-ID	ST-NAME	CLASS-ID
S01	Nguyen Van Nam	L01
S02	Nguyen Van Nam	L01
S03	Tran Quoc Tuan	L01
S04	Le Van Cuong	L02
S05	Nguyen Van Cuong	L02

2.3. Attributes

- The columns of a relation are named by attributes.
- An attribute describes the meaning of entries in the column below.
- Each attribute belong to one **data type**, so that have **Domain**.
- **Domain**: Set of possible **atomic** values, including **Null** value.
- in **Domain** not permitted record structure, set, list, array, or any other type that can broken into smaller components.

2.4. Schemes

- Name of a relation and the set of attributes for a relation is called the **relation scheme** for that relation.
- Show scheme for the relation with the relation name followed by a parenthesized list of its attributes.
- e.g. **STUDENT(ST-ID, ST-NAME, CLASS-ID)**
- The attributes in a relation scheme are a set, not a list.
- A database consists of one or more relation schemes. That is called a relational database scheme or just database scheme.
- In general: $R(A_1 , A_2 , ..., A_n)$

2.5. Tuples

- The rows of a relation are called tuples.
- A tuple has one component for each attribute of the relation.
- For instance, the first of the five tuples of **STUDENT** has the three components:
S01, Nguyen Van Nam, L01
- Write a tuple in isolation, use commas to separate components, and use parentheses to surround the tuple.
e.g. *(S01, Nguyen Van Nam, L01)*
- **t** is a tuple. **t.Attribute** is a value of attribute in tuple **t**.
e.g. *t = (S01, Nguyen Van Nam, L01)*
t.(ST-ID) = 'S01'

2.6. Relation Instances

- Tuples in relation change over time.
- A set of tuples for a given relation in time is an instance of that relation.
- Denote: $r(R)$
- e.g. **STUDENT(ST-ID, ST-NAME, CLASS-ID)**
in time we see we have one *relation instance*

ST-ID	ST-NAME	CLASS-ID
S01	Nguyen Van Nam	L01
S02	Nguyen Van Nam	L01
S03	Tran Quoc Tuan	L01
S04	Le Van Cuong	L02
S05	Nguyen Van Cuong	L02

2.7. Keys

- A key of a relation R is a subset K of R such that for any distinct tuples t_1 and t_2 in $r(R)$ then $t_1(K) \neq t_2(K)$ and no proper subset K' of K shares this property.
- X is a **superkey** of R if X contains a key of R .
- In R maybe more than one key.
- Each key denote one underline.
- e.g. **STUDENT** have one key **ST-ID**.

3.1. Overview of Relational Algebra

The operations of the traditional relational algebra fall into four broad classes:

1. Set operations: union, intersection, and difference. Applied to relations.
2. Remove parts of a relation: Selection and Projection.
3. Combine the tuples of two relations: Cartesian product, Join.
4. Renaming: Changes the relation scheme.
5. Group.
6. ...

3.2. Conditions

Let \mathcal{R}, \mathcal{S} with conditions:

1. Identical sets of attributes, domains for each attribute must be the same.
2. Attributes must be ordered to the same for both relations.

3.3. Set Operators

Relation \mathcal{R} and \mathcal{S} : set of tuples.

1. $\mathcal{R} \cup \mathcal{S}$: **Union** of \mathcal{R} and \mathcal{S} .

$$\mathcal{R} \cup \mathcal{S} = \{t \mid t \in r(\mathcal{R}) \vee t \in s(\mathcal{S})\}$$

2. $\mathcal{R} \cap \mathcal{S}$: **Intersection** of \mathcal{R} and \mathcal{S} .

$$\mathcal{R} \cap \mathcal{S} = \{t \mid t \in r(\mathcal{R}) \wedge t \in s(\mathcal{S})\}$$

3. $\mathcal{R} - \mathcal{S}$: **Difference** of \mathcal{R} and \mathcal{S} .

$$\mathcal{R} - \mathcal{S} = \{t \mid t \in r(\mathcal{R}) \wedge t \notin s(\mathcal{S})\}$$

3.4. Selection σ

Let relation \mathcal{R} and C is a conditional logic expression

- C made from logic operator: $=, \neq, <, \leq, >, \geq, \wedge, \vee, \text{not}$
- Selection operator applied to \mathcal{R} , produces a new relation with a subset of tuples of $r(\mathcal{R})$
- The tuples in the resulting relation are those that satisfy some condition C that involves the attributes of \mathcal{R}
- Notation: $\sigma_C(\mathcal{R})$
- $\sigma_C(\mathcal{R}) = \{t \mid t \in r(\mathcal{R}) , t \text{ satisfy } C\}$

e.g. σ

Suppose $\mathcal{R}(ABC)$ with instance $r(\mathcal{R})$:

A	B	C
1	3	1
2	3	2
3	4	3

$\sigma_{B=3 \wedge C>1}(\mathcal{R})$

A	B	C
2	3	2

3.5. Projection π

Let relation \mathcal{R} and $X \subseteq \mathcal{R}$. Projection of \mathcal{R} on to X :

- Notation: $\pi_X(\mathcal{R})$
- $\pi_X(\mathcal{R})$ only have X attributes
- $\pi_X(\mathcal{R}) = \{t.X \mid t \in r(\mathcal{R})\}$
- Remove duplicate tuples in $\pi_X(\mathcal{R})$

If $X_1 \subseteq X_2 \subseteq \dots \subseteq X_m$, then

$$\pi_{X_1}(\pi_{X_2}(\dots(\pi_{X_m}(\mathcal{R}))\dots)) = \pi_{X_1}(\mathcal{R})$$

e.g. π

Suppose $\mathcal{R}(ABC)$ with instance $r(\mathcal{R})$:

A	B	C
1	3	1
2	3	2
3	4	3

$\pi_B(\mathcal{R})$

B
3
4

3.6. Cartesian product \times

Seeing relation \mathcal{R} and \mathcal{S} to be a set of tuples.

- *Cartesian product* (cross-product, just product) of two sets \mathcal{R} and \mathcal{S} is the set of pairs of (element of \mathcal{R}) and (element of \mathcal{S}).
- Denote: $\mathcal{R} \times \mathcal{S}$
- Relation $\mathcal{R} \times \mathcal{S}$ have **all attributes** of \mathcal{R} and \mathcal{S}
- $\mathcal{R} \times \mathcal{S} = \{(t_1, t_2) \mid t_1 \in r(\mathcal{R}), t_2 \in s(\mathcal{S})\}$

e.g. \times

Let $\mathcal{R}(AB)$, $\mathcal{S}(CD)$; $\mathcal{R} \times \mathcal{S}$

\mathcal{R}	
A	B
1	2
3	4

\mathcal{S}	
C	D
5	6
7	8

$\mathcal{R} \times \mathcal{S}$			
A	B	C	D
1	2	5	6
1	2	7	8
3	4	5	6
3	4	7	8

3.7. Natural Join \bowtie

Let relation \mathcal{R} and \mathcal{S} with $X = \mathcal{R} \cap \mathcal{S}$

Natural Join of \mathcal{R} and \mathcal{S} which we pair only those tuples from $r(\mathcal{R})$ and $s(\mathcal{S})$ that agree in X attribute.

- X is called attribute common.
- Notation: $\mathcal{R} \bowtie \mathcal{S}$
- $\mathcal{R} \bowtie \mathcal{S} = \{(t_1, t_2) \mid t_1 \in r(\mathcal{R}), t_2 \in s(\mathcal{S}), t_1.X = t_2.X\}$

$$L = \mathcal{R} \cup (\mathcal{S} - \mathcal{R})$$

$$\mathcal{R} \bowtie \mathcal{S} = \pi_L(\sigma_{\mathcal{R}.X = \mathcal{S}.X}(\mathcal{R} \times \mathcal{S}))$$

If $X = \emptyset$ then *Natural Join* become *Cartesian Product*

e.g. \bowtie

Let $\mathcal{R}(ABC)$, $\mathcal{S}(CD)$; $\mathcal{R} \bowtie \mathcal{S}$

\mathcal{R}		
A	B	C
1	1	2
5	5	3
7	6	4
1	5	7

\mathcal{S}	
C	D
2	6
5	8
2	4
4	8

$\mathcal{R} \bowtie \mathcal{S}$			
A	B	C	D
1	1	2	6
1	1	2	4
7	6	4	8

3.8. Theta-join \bowtie_{θ}

Theta-join of relations \mathcal{R} and \mathcal{S} based on *condition* θ :

The result of this operation is constructed as follows:

1. Take the product of \mathcal{R} and \mathcal{S} : $\mathcal{R} \times \mathcal{S}$
2. Select from $\mathcal{R} \times \mathcal{S}$ so that satisfy the *condition* θ

$$\mathcal{R} \bowtie_{\theta} \mathcal{S} = \sigma_{\theta}(\mathcal{R} \times \mathcal{S})$$

e.g. \bowtie_{θ}

Let $\mathcal{R}(ABC)$, $\mathcal{S}(DE)$; $\mathcal{R} \bowtie_{C > D} \mathcal{S}$

\mathcal{R}		
A	B	C
1	1	2
5	5	3

\mathcal{S}	
D	E
2	6
5	8

$\mathcal{R} \bowtie_{C > D} \mathcal{S}$				
A	B	C	D	E
5	5	3	2	6

3.9. Left Semi-join \ltimes

Let \mathcal{R} and \mathcal{S} . The **Left Semi-join** is similar to the natural join.

The result is the set of **all tuples in $r(\mathcal{R})$** for which there is a tuple in $s(\mathcal{S})$ that is **equal on their common attribute**.

The difference from a natural join is that **other attributes of \mathcal{S} do not appear**.

- Notation: $\mathcal{R} \ltimes \mathcal{S}$
- $\mathcal{R} \ltimes \mathcal{S} = \{t \mid t \in r(\mathcal{R}) \wedge \exists s \in (\mathcal{R} \bowtie \mathcal{S})\}$, or
- $\mathcal{R} \ltimes \mathcal{S} = \pi_{\mathcal{R}}(\mathcal{R} \bowtie \mathcal{S})$

e.g. \times

Let $\mathcal{R}(ABC)$, $\mathcal{S}(CD)$; $\mathcal{R} \times \mathcal{S}$

\mathcal{R}		
A	B	C
1	1	2
5	5	3

\mathcal{S}	
C	D
2	6
5	8

$\mathcal{R} \times \mathcal{S}$		
A	B	C
1	1	2

3.10. Right Semi-join \bowtie

Let \mathcal{R} and \mathcal{S} . The **Right Semi-join** is similar to the natural join.

The result is the set of **all tuples in $s(\mathcal{S})$** for which there is a tuple in $r(\mathcal{R})$ that is **equal on their common attribute**.

The difference from a natural join is that **other attributes of \mathcal{R} do not appear**.

- Notation: $\mathcal{R} \bowtie \mathcal{S}$
- $\mathcal{R} \bowtie \mathcal{S} = \{t \mid t \in s(\mathcal{S}) \wedge \exists r \in (\mathcal{R} \bowtie \mathcal{S})\}$, or
- $\mathcal{R} \bowtie \mathcal{S} = \pi_{\mathcal{S}}(\mathcal{R} \bowtie \mathcal{S})$

e.g. \times

Let $\mathcal{R}(ABC)$, $\mathcal{S}(CD)$; $\mathcal{R} \times \mathcal{S}$

\mathcal{R}		
A	B	C
1	1	2
5	5	3

\mathcal{S}	
C	D
2	6
5	8

$\mathcal{R} \times \mathcal{S}$	
C	D
2	6

3.11. Anti-join \triangleright

Let \mathcal{R} and \mathcal{S} . The **Anti-join** is similar to the semijoin, but the result of an antijoin is only those tuples in \mathcal{R} for which there is no tuple in \mathcal{S} that is equal on their common attribute names.

- Notation: $\mathcal{R} \triangleright \mathcal{S}$
- $\mathcal{R} \triangleright \mathcal{S} = \{t \mid t \in r(\mathcal{R}) \wedge \neg \exists s \in (\mathcal{R} \bowtie \mathcal{S})\}$, or
- $\mathcal{R} \triangleright \mathcal{S} = \{t \mid t \in r(\mathcal{R}), \text{ there is no tuple } s(\mathcal{S}) \text{ that satisfies } (\mathcal{R} \bowtie \mathcal{S})\}$, or
- $\mathcal{R} \triangleright \mathcal{S} = \mathcal{R} - \mathcal{R} \bowtie \mathcal{S}$

e.g. \triangleright

Let $\mathcal{R}(ABC)$, $\mathcal{S}(CD)$; $\mathcal{R} \triangleright \mathcal{S}$

\mathcal{R}		
A	B	C
1	1	2
5	5	3

\mathcal{S}	
C	D
2	6
5	8

$\mathcal{R} \triangleright \mathcal{S}$		
A	B	C
5	5	3

3.12. Complex Queries

- *Complex queries* refer to data in many relations.
- In relational algebra, like all algebras, allows us to *form expressions* by applying operations to the result of other operations.
- Possible to represent expressions as *expression trees*, so that machine program can do it easily.

e.g.

Let $\mathcal{R}(ABCDE)$

A	B	C	D	E
1	2	5	6	2
1	2	7	8	3
3	4	5	6	4

Query: Find **A** when **B** > **3** and **A** when **E** < **4**

Answer: $\pi_A(\sigma_{B>3}(\mathcal{R}) \cup \sigma_{E<4}(\mathcal{R}))$

A
1
3

Another answer: $\pi_A(\sigma_{B>3 \vee E<4}(\mathcal{R}))$

3.13. Rename ρ

- To control the names of the attributes or relation scheme, using *rename* operator. All data (tuples) not change in this operator.

- Notation: $\rho_{S(A_1, A_2, \dots, A_n)}(\mathcal{R})$

rename \mathcal{R} to S . n attributes of \mathcal{R} rename to A_1, A_2, \dots, A_n

- eg. Let $\mathcal{R}(ABC)$.

$\rho_S(\mathcal{R})$ give us the relation $S(ABC)$

$\rho_{ADE}(S)$ give us the relation $S(ADE)$

3.14. Assignment :=

- Definition: Assign an relation expression to relation.
- Notation: :=
- eg. Suppose

StudentMark(StudentID, SubjectID, Mark)

$\mathcal{R}(T, U, M) := \sigma_{SubjectID='S01' \wedge Mark > 7}(StudentMark)$

$\mathcal{S}(T, U, M) := \sigma_{SubjectID='S02' \wedge Mark > 9}(StudentMark)$

$Answer(StudentID) := \pi_T(\mathcal{R} \cap \mathcal{S})$

3.15. Div \div

Divide operator has a rather complex definition, but it does have some applications in natural situations.

Let \mathcal{R} and \mathcal{S} , with $\mathcal{S} \subseteq \mathcal{R}$

- Let $\mathcal{R}' = \mathcal{R} - \mathcal{S}$
- \mathcal{R} *divided by* \mathcal{S} , written $\mathcal{R} \div \mathcal{S}$
- $r'(\mathcal{R}') = \{t \mid \text{for every tuple } t_s \in s(\mathcal{S})$
there is a } t_r \in r(\mathcal{R}) \text{ with } t_r(\mathcal{R}') = t \text{ and } t_r(\mathcal{S}) = t_s\}

e.g. \div

Suppose $\mathcal{R}(AB)$ and $\mathcal{S}(B)$, with $\mathcal{S} \subseteq \mathcal{R}$

$$\mathcal{R}' = \mathcal{R} - \mathcal{S}$$

$$\mathcal{R}'(A) = \mathcal{R}(AB) \div \mathcal{S}(B)$$

\mathcal{R}' A
5
7

\mathcal{R}	
A	B
7	1
7	2
5	1
5	2
8	1

\mathcal{S} B
1
2

3.16. Outer Joins

We know:

- *Null* value is the value unknown or does not exist.
- All comparisons involving null are false by definition.

So that:

- Extension of the join operation that avoids loss of information.
- Computes the join and then adds tuples from one relation that does not match tuples in the other relation to the result of the join.

3.17. Left Outer Joins \bowtie

Let \mathcal{R} and \mathcal{S} . *Left outer join* of that two relations is

- Denote: $\mathcal{R} \bowtie \mathcal{S}$
- The result is the set of:
 - Tuples in \mathcal{R} and \mathcal{S} that are equal on their common attribute names, and
 - Addition to tuples in \mathcal{R} that have no matching tuples in \mathcal{S} .
- Let (\perp, \dots, \perp) : Singleton relation with attributes that are unique to the \mathcal{S} (are not attributes of \mathcal{R}):

$$(\mathcal{R} \bowtie \mathcal{S}) \cup ((\mathcal{R} - \pi_{\mathcal{R}}(\mathcal{R} \bowtie \mathcal{S})) \times \{(\perp, \dots, \perp)\})$$

e.g. \bowtie

Suppose $\mathcal{R}(AB)$ and $\mathcal{S}(BC)$.

A	B
1	2
2	4
3	6

B	C
2	7
3	8
5	9

$\mathcal{R} \bowtie \mathcal{S}$

A	B	C
1	2	7
2	4	\perp
3	6	\perp

3.18. Right Outer Joins $\bowtie\llcorner$

Let \mathcal{R} and \mathcal{S} . *Right outer join* of that two relations is

- Denote: $\mathcal{R} \bowtie\llcorner \mathcal{S}$
- The result is the set of:
 - Tuples in \mathcal{R} and \mathcal{S} that are equal on their common attribute names, and
 - Addition to tuples in \mathcal{S} that have no matching tuples in \mathcal{R} .
- Let (\perp, \dots, \perp) : Singleton relation with attributes that are unique to the \mathcal{R} (are not attributes of \mathcal{S}):

$$(\mathcal{R} \bowtie\llcorner \mathcal{S}) \cup (\{(\perp, \dots, \perp)\} \times (\mathcal{S} - \pi_{\mathcal{S}}(\mathcal{R} \bowtie\llcorner \mathcal{S})))$$

e.g. \bowtie

Suppose $\mathcal{R}(AB)$ and $\mathcal{S}(BC)$.

A	B
1	2
2	4
3	6

B	C
2	7
3	8
5	9

$\mathcal{R} \bowtie \mathcal{S}$

A	B	C
1	2	7
\perp	3	8
\perp	5	9

3.19. Full Outer Joins \bowtie

Let \mathcal{R} and \mathcal{S} . *Full outer join* of that two relations is

- Denote: $\mathcal{R} \bowtie \mathcal{S}$
- The result is the set of:
 - Tuples in \mathcal{R} and \mathcal{S} that are equal on their common attribute names, and
 - Addition to tuples in \mathcal{R} that have no matching tuples in \mathcal{S} and to tuples in \mathcal{S} that have no matching tuples in \mathcal{R} .
- $\mathcal{R} \bowtie \mathcal{S} = (\mathcal{R} \bowtie \mathcal{S}) \cup (\mathcal{R} \bowtie \mathcal{S})$

e.g. \bowtie

Suppose $\mathcal{R}(AB)$ and $\mathcal{S}(BC)$.

A	B
1	2
2	4
3	6

B	C
2	7
3	8
5	9

$\mathcal{R} \bowtie \mathcal{S}$

A	B	C
1	2	7
2	4	\perp
3	6	\perp
\perp	3	8
\perp	5	9

3.20. Duplicate Elimination Operators δ

Let \mathcal{R}

- Denote: $\delta(\mathcal{R})$
- Return the set consisting of *one copy of every tuple that appears one or more times* in relation \mathcal{R} .

e.g. $\mathcal{R}(ABC); \delta(\mathcal{R})$

\mathcal{R}		
A	B	C
1	2	3
4	5	6
1	2	3

$\delta(\mathcal{R})$		
A	B	C
1	2	3
4	5	6

3.21. Aggregate functions

Aggregate function takes a collection of values and returns a single value as a result.

- *sum*(·): Sum of values.
- *avg*(·): Average value.
- *min*(·): Minimum value.
- *max*(·): Maximum value.
- *count*(·): Number of values.

3.22. Group \mathcal{G}

Let \mathcal{U} is any relational-algebra expression

g_1, g_2, \dots, g_n : Attributes on which to group (can be empty)

Each f_i is an aggregate function

Each A_i is an attribute name

- Denote: $g_1, g_2, \dots, g_n \mathcal{G}_{f_1(A_1), f_2(A_2), \dots, f_m(A_m)}(\mathcal{U})$
- Partition the tuples of \mathcal{U} into groups g_1, g_2, \dots, g_n .
- For each group:
 - The **grouping attributes' values** for that group and
 - Can use aggregate functions f_i for value in group.

e.g. \mathcal{G}

Suppose $\mathcal{R}(ABC)$

A	B	C
1	4	3
1	3	3
4	5	6

$\mathcal{G}_{sum(A), min(C), count(C)}(\mathcal{R})$

SumA	MinB	CountC
6	3	3

e.g. \mathcal{G}

Suppose $\mathcal{R}(ABC)$

A	B	C
1	4	3
1	3	3
4	5	6
4	6	7
4	8	9

$\mathcal{G}_{\min(B), \max(B), \text{Count}(A)}(\mathcal{R})$

A	MinB	MaxB	CountA
1	3	4	2
4	5	8	3

3.23. Sorting Operator τ

Let \mathcal{R} , with $L \subseteq \mathcal{R}$

Sorting data in $r(\mathcal{R})$ over L attributes.

Denote: $\tau_L(\mathcal{R})$

Suppose $\mathcal{R}(ABC)$, $\tau_{A,B}(\mathcal{R})$

\mathcal{R}		
A	B	C
4	2	3
4	5	6
1	3	3
1	2	5

$\tau_{A,B}(\mathcal{R})$		
A	B	C
1	2	5
1	3	3
4	2	3
4	5	6

4.1. SQL

- * SQL (pronounced “sequel”) - “Structured Query Language” is the principal language used to describe and manipulate relational databases.
- * There are two aspects to SQL:
 1. The Data-Definition sublanguage for declaring database schemas (DDL).
 2. The Data-Manipulation sublanguage for querying databases and for modifying the database (DML).

4.2. Relations in SQL

- * SQL makes a distinction between three kinds of relations:
 1. **Table**: Relation that exists in the database and that can be modified by changing its tuples, as well as queried.
 2. **View**: Relations defined by a computation. These relations are not stored, but are constructed, in whole or in part, when needed.
 3. **Temporary table**: Constructed by the SQL language processor when it performs its job of executing queries and data modifications. These relations are then thrown away and not stored.

4.3. Create Table

Create table in current database:

```
CREATE TABLE tableName(  
    column1 datatype,  
    .....  
    columnN datatype,  
    PRIMARY KEY(one or more columns)  
);
```

e.g.

```
CREATE TABLE Company(  
    ID Int Primary Key,  
    Name Text Not Null,  
    Age Int Not Null,  
    Address Char(50),  
    Salary Real  
);
```

4.4. Drop, Alter Table

Drop table in current database:

```
DROP TABLE TableName;
```

- * Alter table statement is used to add, delete, or modify columns in an existing table.
- * Alter table statement is also used to add and drop various constraints on an existing table.

Alter Table - ADD Column:

```
ALTER TABLE tableName  
ADD columnName datatype;
```

Alter Table - DROP Column:

```
ALTER TABLE tableName  
DROP COLUMN columnName;
```

e.g.

- `CREATE TABLE contacts`
 (`id INTEGER PRIMARY KEY`,
 `name TEXT NOT NULL COLLATE NOCASE`,
 `phone TEXT NOT NULL DEFAULT 'Unknown'`,
 `UNIQUE (name,phone)`
);
- `DROP TABLE Shippers;`
- `ALTER TABLE Customers`
 `ADD Email varchar(255);`
- `ALTER TABLE Customers`
 `DROP COLUMN Email;`
- `ALTER TABLE Persons`
 `ALTER COLUMN DateOfBirth year;`

4.5. Insert

- **INSERT a single row into a table:**

```
INSERT INTO tableName (column1, column2 ,..)  
VALUES  
    (value1, value2 ,...);
```

- **INSERT multiple rows into a table:**

```
INSERT INTO tableName (column1, column2 ,..)  
VALUES  
    (value1, value2 ,...),  
    (value1, value2 ,...),  
    ...  
    (value1, value2 ,...);
```

e.g.

- `INSERT INTO Students (name)`
`VALUES`
`('Bud Powell');`
- `INSERT INTO Students (name)`
`VALUES`
`("Buddy Rich"),`
`("Candido"),`
`("Charlie Byrd");`

4.6. Update

The search condition in the **WHERE** has the following form:

Left Expression **Comparicon** **Operator** *Right Expression*

```
UPDATE tableName  
SET column1 = newValue1,  
      column2 = newValue2  
WHERE  
      SearchCondition
```

e.g.

```
WHERE column1 = 100;
```

```
WHERE column2 IN (1,2,3);
```

```
WHERE column3 LIKE 'An%';
```

```
WHERE column4 BETWEEN 10 AND 20;
```

4.7. Comparison Operators

Operator	Meaning
=	Equal to
<> or !=	Not equal to
<	Less than
>	Greater than
<=	Less than or equal to
>=	Greater than or equal to

4.8. Logical Operators

Notice that 1 means **TRUE**, and 0 means **FALSE**.

Operator	Meaning
ALL	Returns 1 if all expressions are 1.
AND	Returns 1 if both expressions are 1, and 0 if one of the expressions is 0.
ANY	Returns 1 if any one of a set of comparisons is 1.
BETWEEN	Returns 1 if a value is within a range.
EXISTS	Returns 1 if a subquery contains any rows.
IN	Returns 1 if a value is in a list of values.
LIKE	Returns 1 if a value matches a pattern.
NOT	Reverses the value of other operators such as NOT EXISTS , NOT IN , NOT BETWEEN , etc.
OR	returns true if either expression is 1.

e.g.

```
UPDATE employees  
SET lastname = 'Smith'  
WHERE  
    employeeid = 3;
```

```
UPDATE employees  
SET city = 'Toronto',  
    state = 'ON',  
    postalcode = 'M5P 2N7'  
WHERE  
    employeeid = 4;
```

4.9. Delete

DELETE FROM tableName
WHERE searchCondition

e.g.

DELETE FROM employees
WHERE
name **LIKE** '%Santana%';

Database Scheme

Suppose *Order Management* Database scheme:

1. **Categories(CategoryID, CategoryName, Description)**
2. **Products(ProductID, ProductName, UnitPrice, CategoryID)**
3. **Customers(CustomerID, CustomerName, Address, Phone, Email, Country)**
4. **Orders(OrderID, OrderDate, RequiredDate, CustomerID)**
5. **OrderDetails(OrderID, ProductID, UnitPrice, Quantity, Discount)**

4.10. Products and Joins in SQL

e.g. Suppose we want to know the name of the customer in which he order with OrderID = 'D01'. To answer this question we need the following two relations:

Customers(CustomerID, CustomerName, Address,
Phone, Fax, Country)

Orders(OrderID, OrderDate, RequiredDate, CustomerID)

```
SELECT B.CustomerID, CustomerName
FROM Orders AS A, Customers AS B
WHERE A.CustomerID = B.CustomerID
      AND OrderID = 'D01';
```

Same as the relational algebra expression:

$$\pi_{CustomerID, CustomerName} \sigma_{OrderID = 'D01'} (Orders \bowtie Customers)$$

4.11. Union \cup

e.g. Suppose we wanted the ID and name of all product it is belong to Category = 'C01' or 'C02'.

```
( SELECT ProductID, ProductName  
  FROM Products  
 WHERE CategoryID = 'C01' )
```

UNION

```
( SELECT ProductID, ProductName  
  FROM Products  
 WHERE CategoryID = 'C02' );
```

Same as the relational algebra expression:

$$\pi_{ProductID, ProductName} \sigma_{CategoryID = 'C01'}(Products) \cup \pi_{ProductID, ProductName} \sigma_{CategoryID = 'C02'}(Products)$$

4.12. Intersection \cap

e.g. Suppose we wanted the ID and name of product so that order with customer ID = 'T01' and 'T02'.

```
( SELECT C.ProductID, ProductName
  FROM Orders A, OrderDetails B, Products C
 WHERE CustomerID = 'T01'
       AND A.OrderID = B.OrderID
       AND B.ProductID = C.ProductID )
```

INTERSECT

```
( SELECT C.ProductID, ProductName
  FROM Orders A, OrderDetails B, Products C
 WHERE CustomerID = 'T02'
       AND A.OrderID = B.OrderID
       AND B.ProductID = C.ProductID );
```

4.13. Difference —

e.g.

Suppose we wanted the ID and name of product that never order.

```
( SELECT ProductID, ProductName
  FROM Products
   EXCEPT
  ( SELECT B.ProductID, ProductName
    FROM OrderDetails A, Products B
   WHERE A.ProductID = B.ProductID );
```


4.14. Subqueries

- * **Subquery**: A query that is part of another query.
- * Subqueries can have subqueries, and so on, down as many levels as we wish.
- * There are a number of other ways that subqueries can be used:
 1. Subqueries can return a *single constant*, and this constant can be *compared* with another value in a **WHERE** clause.
 2. Subqueries can return *relations* that can be used in various ways in **WHERE** clauses.
 3. Subqueries can appear in **FROM** clauses.

e.g.

Finding information about customer that order in OderID = 'D01'.

```
SELECT *  
FROM Customers  
WHERE CustomerID = ( SELECT CustomerID  
                     FROM Orders  
                     WHERE OrderID = 'D01' );
```

Conditions Involving Relations

- * There are a number of SQL operators that we can apply to a relation \mathcal{R} and produce a boolean result.
- * However, the relation \mathcal{R} must be expressed as a subquery and s is a scalar value. In this situation, the subquery \mathcal{R} is required to produce a *one column relation*.
- * Here are the definitions of the operators:
 1. $\text{EXISTS } \mathcal{R}$ is *true* if and only if \mathcal{R} is not empty.
 2. $s \text{ IN } \mathcal{R}$ is true if and only if s is equal to one of the values in \mathcal{R} . Likewise, $s \text{ NOT IN } \mathcal{R}$ is true if and only if s is equal to no value in \mathcal{R} .

4.15. Conditions Involving Relations

3. $s > \text{ALL } \mathcal{R}$ is true if and only if s is greater than every value in unary relation \mathcal{R} . Similarly, the $>$ operator could be replaced by any of the other five comparison operators, with the analogous meaning: s stands in the stated relationship to every tuple in \mathcal{R} . For instance, $s <> \text{ALL } \mathcal{R}$ is the same as $s \text{ NOT IN } \mathcal{R}$.
4. $s > \text{ANY } \mathcal{R}$ is true if and only if s is greater than at least one value in unary relation \mathcal{R} . Similarly, any of the other five comparisons could be used in place of $>$, with the meaning that s stands in the stated relationship to at least one tuple of \mathcal{R} . For instance, $s = \text{ANY } \mathcal{R}$ is the same as $s \text{ IN } \mathcal{R}$.

Conditions Involving Relations

The **EXISTS**, **ALL**, and **ANY** operators can be negated by putting **NOT** in front of the entire expression.

- **NOT EXISTS** \mathcal{R} is true if and only if \mathcal{R} is empty.
- **NOT** $s \geq \text{ALL } \mathcal{R}$ is true if and only if s is not the maximum value in \mathcal{R} .
- **NOT** $s > \text{ANY } \mathcal{R}$ is true if and only if s is the minimum value in \mathcal{R} .

4.16. Conditions Involving Tuples

- * If a tuple t has the same number of components as a relation \mathcal{R} , then it makes sense to compare t and \mathcal{R} in expressions.
- * Examples are $t \text{ IN } \mathcal{R}$ or $t \text{ <> ANY } \mathcal{R}$.

e.g. Support *Warehouse Management* database scheme:

1. **Categories**(CategoryID, CategoryName, Description)
2. **Products**(ProductID, ProductName, UnitPrice, CategoryID)
3. **Warehouses**(WarehouseID, Name, Address, CategoryID)
4. **InStocks**(WarehouseID, ProductID, Quantity)

e.g.

Finding information about product, that it can be store in warehouse 'W01'.

```
SELECT *  
FROM Products  
WHERE (ProductID, CategoryID) IN  
    ( SELECT ProductID, B.CategoryID  
      FROM Warehouses A, Products B  
      WHERE WarehouseID = 'W01' AND  
        A.CategoryID = B.CategoryID  
    );
```

4.17. Correlated Subqueries

- * The simplest subqueries can be evaluated *once and for all*, and the result used in a higher-level query.
- * *Nested subqueries* requires the subquery to be evaluated many times. This type is called a *correlated subquery*.

e.g. Finding information about warehouse and product that instock with quantity > 30.

```
SELECT A.*, B.*  
FROM Warehouses A, Products B  
WHERE A.CategoryID = B.CategoryID  
      AND 30 < ( SELECT Quantity  
                  FROM Instocks  
                  WHERE WarehouseID = A.WarehouseID  
                    AND ProductID = B.ProductID )  
ORDER BY WarehouseID, ProductID;
```


4.18. Subqueries in FROM Clauses

- * Subqueries is as relations in a **FROM** clause.
- * In a **FROM** list, use a parenthesized subquery and give it a alias.

e.g. Finding information about product that store in warehouse with WarehouseID = 'W01'.

```
SELECT A.*  
FROM Products A, ( SELECT ProductID  
                   FROM Instocks  
                   WHERE WarehouseID = 'W01'  
                 ) B  
WHERE A.ProductID = B.ProductID;
```

4.19. SQL Join Expressions

Let $\mathcal{R}_1, \mathcal{R}_2$ are relations. \mathcal{P} is a condition expression.

1. Cross Join: $\mathcal{R}_1 \times \mathcal{R}_2$. FROM $\mathcal{R}_1, \mathcal{R}_2$
2. Inner Join:
 - Join: FROM \mathcal{R}_1 [Inner] Join \mathcal{R}_2 On \mathcal{P}
 - Theta Join: FROM \mathcal{R}_1 Join \mathcal{R}_2 On \mathcal{P}
 - Natural Join: FROM \mathcal{R}_1 Natural Join \mathcal{R}_2
FROM \mathcal{R}_1 Join \mathcal{R}_2 On $\mathcal{R}_1.X = \mathcal{R}_2.X$

SQL Join Expressions

Outer Join:

- Left Outer Join:

FROM \mathcal{R}_1 Left Outer Join \mathcal{R}_2 On \mathcal{P}

FROM \mathcal{R}_1 Natural Left Outer Join \mathcal{R}_2

- Right Outer Join:

FROM \mathcal{R}_1 Right Outer Join \mathcal{R}_2 On \mathcal{P}

FROM \mathcal{R}_1 Natural Right Outer Join \mathcal{R}_2

- Full Outer Join:

FROM \mathcal{R}_1 Full Join \mathcal{R}_2 On \mathcal{P}

FROM \mathcal{R}_1 Natural Full Join \mathcal{R}_2

e.g.

Finding information about customer have never order.

```
SELECT A.*
```

```
FROM Customers A Natural Left Outer Join Orders B
```

```
WHERE B.OrderID is Null
```

```
ORDER BY CustomerName;
```

4.20. Eliminating Duplicates

- * SQL system does not eliminate duplicates. Thus, the SQL response to a query may list the same tuple several times.
- * If we do not wish duplicates in the result, then we may follow the keyword **SELECT** by the keyword **DISTINCT**.

SELECT DISTINCT X

- * One might be tempted to place DISTINCT after every SELECT, on the theory that it is harmless.
- * So the relation must be sorted or partitioned.

4.21. Grouping and Aggregation Operators

In algebra: Aggregation functions and Group.

$$g_1, g_2, \dots, g_n \mathcal{G}_{f_1(A_1), f_2(A_2), \dots, f_m(A_m)}(\mathcal{U})$$

In SQL:

SELECT < One or more Attributes in Group clause >

[, Aggregation Functions]

FROM < Relations >

[**WHERE** < Tuple-Condition >]

GROUP BY < Attribute want to Group >

[**HAVING** < Group-Condition >]

4.22. Aggregation Operators

- * Five aggregation operators: `Sum(·)`, `Avg(·)`, `Min(·)`, `Max(·)`, and `Count(·)`.
- * Augument is a scalar valued expression, typically a column name, in a `SELECT` clause.
- * `Count(*)`: Counts all the tuples in the relation that agree with condition on `WHERE` clause.
- * Another option of eliminating duplicates from the column before applying the aggregation operator by using the keyword `Distinct`.
- * `Count(Distinct A)` counts the number of distinct values in column A.
- * We could use any of the other operators in place of `Count` here.

4.23. Grouping

* To group tuples, using **GROUP BY** clause, following the **WHERE** clause.

Notice:

1. * **GROUP BY** are followed by a list of *grouping attributes*.
2. * *Whatever aggregation operators* are used in the **SELECT** clause are applied *only within groups*.

e.g. Finding total quantity for each product its order.

```
SELECT ProductID, Sum(Quantity)
FROM OrderDetails
GROUP BY ProductID;
```


4.24. Grouping, Aggregation, and Nulls

When tuples have nulls, there are a few rules we must remember:

- The value **Null** is ignored in any aggregation. It does not contribute to a **sum**, **average**, or **count of an attribute**, nor can it be the **minimum or maximum** in its column. For example, **Count(*)** is always a count of the number of tuples in a relation, but **Count(A)** is the number of tuples with **non-NULL** values for attribute A.
- On the other hand, **Null** is treated as an ordinary value when forming groups. That is, we can have a group in which one or more of the grouping attributes are assigned the value **Null**.
- When we perform any aggregation except count over an empty bag of values, the result is **Null**. The count of an empty bag is 0.

e.g.

Show order total cost for each customer (ID and Name).

```
SELECT C.CustomerID, CustomerName,  
       Sum(Quantity × Price) As 'Total Cost'  
FROM Orders A, OrderDetails B, Customers C  
WHERE A.OrderID = B.OrderID  
       And A.CustomerID = C.CustomerID  
GROUP BY C.CustomerID, CustomerName;
```

e.g. Show order total cost for each order at year 2018.

```
SELECT A.OrderID, Sum(Quantity × Price) As 'Total Cost'  
FROM Orders A, OrderDetails B  
WHERE A.OrderID = B.OrderID  
       And strftime('%Y,OrderDate) = 2018  
GROUP BY A.OrderID;
```

4.25. HAVING Clauses

Condition after WHERE restrict tuples prior to grouping.

Condition after HAVING restrict group.

e.g. Suppose we want to print the order total cost for only those cost of one product in its order at least 100.

```
SELECT OrderID, Sum(Quantity × Price) As 'Total Cost'  
FROM OrderDetails  
GROUP BY OrderID  
HAVING Min(Quantity × Price) >= 100;
```

There are several rules we must remember about HAVING clauses:

- An aggregation in a **HAVING** clause applies only to the tuples of the group being tested.
- Any attribute of relations in the **FROM** clause may be aggregated in the **HAVING** clause.

4.26. Finding Maximum Group

e.g. Suppose finding the Order in which Total Cost is the maximum.

```
SELECT OrderID, Sum(Quantity × Price) As 'TotalFree'
FROM OrderDetails
GROUP BY OrderID
HAVING TotalFree >= ALL ( SELECT Sum(Quantity × Price)
                        FROM OrderDetails
                        GROUP BY OrderID
                        )
```

But SQLite is not support \geq ALL operator, so that we use this algorithm below:

4.27. Finding Maximum Group in SQLite

This algorithm find maximum group in SQLite:

```
SELECT OrderID, Sum(Quantity × Price) As 'TotalFree'
FROM OrderDetails
GROUP BY OrderID
HAVING TotalFree =
    ( SELECT Max(TotalFree)
      FROM ( SELECT Sum(Quantity × Price) AS 'TotalFree'
            FROM OrderDetails
            GROUP BY OrderID
          )
    )
```

5.1. Integrity Constraints

Parts of Integrity Constraints

Suppose $\mathcal{R}(ABC)$ with $B > C$

There is an *integrity constraint* of scheme R .

Integrity Constraints

- Context: R
- Condition:
 $\forall t \in \forall r$
 $t.B > t.C$
end.
- Influence table:

	Insert	Delete	Update
R	+	-	+(B/C)

5.2. Parts of Integrity Constraints

- In *influence table*, plus(+) sign if statement (operation) on relation maybe enforce business rule's wrong. Otherwise, minus(-) sign.
- With plus sign we write one *Trigger* to control that operation.
- Trigger is an event-driven action that is run automatically when a specified change operation (*INSERT*, *UPDATE*, and *DELETE* statement) is performed on a specified table.
- Triggers are useful for tasks such as enforcing business rules, validating input data, and keeping an audit trail.

Trigger (SQLite)

```
CREATE TRIGGER [IF NOT EXISTS] triggerName
    [BEFORE/AFTER/INSTEAD OF]
    [INSERT/UPDATE/DELETE]
    ON tableName
BEGIN
    [WHEN condition]
    BEGIN
        statements
    END;
END;

DROP TRIGGER [IF EXISTS] triggerName;
```


5.3. Trigger (SQLite)

We can access the data of the row being inserted, deleted, or updated using the **OLD** and **NEW** references in the form:

- **OLD.Attribute**
- **NEW.Attribute**

OLD and NEW references are available depending on the event that causes the trigger to be fired:

Action	Reference
INSERT	NEW is available
UPDATE	Both NEW and OLD are available
DELETE	OLD is available

e.g.

Suppose we want to validate email address before insert a new customer into the Customer relation.

```
CREATE TRIGGER validateEmailCustomer
    BEFORE INSERT ON Customer
BEGIN
    SELECT CASE
        WHEN NEW.email NOT LIKE '%_@__%.__%' THEN
            RAISE(ABORT, 'Invalid email address')
    END;
END;
```

e.g.

When insert a tuple, if the email is not valid, the **RAISE** function **aborts** the insert and issues an error message **'Invalid email address'**.

e.g.

```
INSERT INTO Customer  
VALUES('C01', 'John', '188, ...','0123',  
      'John.Nguyen@database.net','vietnam');
```

Because the email is valid, the insert statement executed successfully.

```
SELECT * FROM Customer
```

Result:

```
('C01', 'John', '188, ...','0123', 'John.Nguyen@database.net','vietnam')
```

5.4. SQL Common Integrity Constraints

```
CREATE TABLE tableName (  
    att1 datatype columnConstraint,  
    att2 datatype columnConstraint,  
    ...,  
    Constraint tableConstraintName < Logic expression >  
);
```

- Constraints are the rules enforced on a data columns on table.
- Constraints could be column level or table level.
- Column level constraints are applied only to one column, whereas table level constraints are applied to the whole table.
- Following are commonly used constraints available in SQLite.

SQL Common Integrity Constraints

1. **NOT NULL** constraint: Ensures that a column cannot have **NULL** value.
2. **DEFAULT** constraint: Provides a default value for a column when none is specified.
3. **UNIQUE** constraint: Ensures that all values in a column are different.
4. **PRIMARY Key**: Uniquely identifies each row/record in a database table.
5. **CHECK** constraint: Ensures that all values in a column satisfies certain conditions.

At **column level** constraint, each column followed by common constraint above.

5.5. Table Constraint

- * At the table level, SQLite supports the **UNIQUE**, **CHECK**, and **PRIMARY KEY** constraints.
- * The check constraint is very similar, requiring only an expression (**CHECK(expression)**).
- * Both the **UNIQUE** and **PRIMARY KEY** constraints, may be require a list of columns:
e.g.
UNIQUE (columnName, [...]),
PRIMARY KEY (columnName, [...])).

e.g. NOT NULL Constraint

```
CREATE TABLE Products (  
    ProductID      Text    Primary Key,  
    ProductName    Text    Not Null,  
    UnitPrice      Real    Not Null,  
    CategoryID     Text    Not Null  
);
```

e.g. DEFAULT Constraint

```
CREATE TABLE Orders (  
    OrderID      Text      Primary Key,  
    OrderDate    Text      Default Current_Date,  
    RequiredDate Text      Not Null,  
    CustomerID   Text  
);
```

When insert data into table, we can miss default column name to get default value:
e.g. suppose now is '2018-10-20'

```
insert into Orders (OrderID, RequiredDate, CustomerID)  
Values ('O01', '2018-10-22', 'C01');
```

```
Select * from Orders;  
( 'O01', '2018-10-20', '2018-10-22', 'C01')
```


e.g. UNIQUE Constraint

```
CREATE TABLE Categories (  
    CategoryID      Text    Primary Key,  
    CategoryName    Text    Not Null Unique,  
    Deacription     Text  
);
```

5.6. CHECK Constraint

* *CHECK constraint at the column level:*

```
CREATE TABLE tableName (  
    ...,  
    columnName dataType CHECK(expression),  
    ...  
);
```

* *CHECK constraint at the table level:*

```
CREATE TABLE tableName (  
    ...,  
    CHECK(expression)  
);
```

e.g. Table Constraint

```
CREATE TABLE Orders (  
  OrderID      Text      Primary Key,  
  OrderDate    Text      Default Current_Date,  
  RequireDate  Text      Not Null,  
  CustomerID   Text,  
  Constraint ValidateRequireDate Check (  
    strftime('%d', OrderDate) <=  
    strftime('%d', RequiredDate) )  
);
```

5.7. SQL Foreign Key Constraints

Foreign key: In a relation, attributes (one or many) are called foreign key if they are not a key in this relation, but a primary key in another relation.

Foreign key assert that a value appearing in one relation must also appear in the primary key component(s) of another relation.

e.g.

Class(ClassID, Description)

C01

C02

Student(StudentID, Name, Address, Email, *ClassID*)

C01

C02

C02

~~C03~~

5.8. Declaring Foreign Key Constraints

Let $\mathcal{R}_1(\underline{XY})$, $\mathcal{R}_2(Z, \overline{X})$

Declaring attributes of relation to be a foreign key, referencing some attribute(s) of a second relation (possibly the same relation) must be have twofold:

- Attribute(s) X of the \mathcal{R}_1 relation must be declared **UNIQUE** or the **PRIMARY KEY**. Otherwise, \mathcal{R}_2 cannot make the **foreign key** declaration.
- $\pi X \mathcal{R}_2 \subseteq \pi X \mathcal{R}_1$

Declaring Foreign Key Constraints

When foreign key constraint is declared, the system has to reject the violating modification.

There are two policy for Update and Delete tuples in foreign key:

1. *Default* policy: Reject Violating Modifications.
2. *Cascade* policy: When changes to the referenced attribute(s) in \mathcal{R}_1 . There are following change to the foreign key in \mathcal{R}_2 .
3. *Null* policy: When a modification to the referenced relation (\mathcal{R}_1) affects a foreign-key value (in \mathcal{R}_2) is changed to **NULL**.

Usually declare: Update Cascade; Delete Set Null.

e.g.

```
CREATE TABLE  $\mathcal{R}_2$  (  
  Z datatype,  
  X datatype,  
  Foreign Key (X) References  $\mathcal{R}_1(X)$   
    [On Update < Cascade / Set Null > ]  
    [On Delete < Cascade / Set Null > ] );
```

e.g.

```
Create Table ShipperDetail (  
  ShipperID Text, ProductID Text, Quantity int,  
  Primary Key (ShipperID, ProductID),  
  Foreign Key (ShipperID) References Shippers(ShipperID)  
    On Delete Set Null On Update Cascade,  
  Foreign Key (ProductID) References Products(ShipperID)  
    On Delete Set Null On Update Cascade );
```

6.9. Views

- `CREATE TABLE` statement create table in database, it stores in physical.
- There is another class of SQL relations, called *(virtual) views*, that do not exist physically. Rather, they are defined by an expression much like a query.
- Views, in turn, can be queried as if they existed physically, and in some cases, we can even modify views.
- Views in SQLite *do not allow* to perform any operations like `INSERT`, `UPDATE` and `DELETE` on views.

Views

```
CREATE VIEW viewName AS  
SELECT column1, column2, ...  
FROM tableName  
[ WHERE condition ];
```

```
DROP VIEW viewName;
```

e.g. Create a view for customer with total orders

```
CREATE VIEW CustomerTotalOrder AS  
SELECT CustomerID, Sum(UnitPrice * Quantity) AS Total  
FROM Orders A, OrderDetail B  
WHERE A.OrderID = B.OrderID  
GROUP BY CustomerID;
```

6.10. Querying Views

A view may be queried exactly as if it were a stored table.

e.g. Finding information about customer with total of orders:

```
SELECT A.*, Total  
FROM Customers A, CustomerTotalOrder B  
WHERE A.CustomerID = B.CustomerID
```

Querying Views

We can replace each view in a FROM clause by a subquery that is identical to the view definition.

```
SELECT A.*, Total
FROM Customers A, ( SELECT CustomerID,
                     Sum(UnitPrice * Quantity) AS Total
                     FROM Orders A, OrderDetail B
                     WHERE A.OrderID = B.OrderID
                     GROUP BY CustomerID ) B
WHERE A.CustomerID = B.CustomerID
```

6.11. Renaming Attributes

We can give a view's attributes in new names.

```
CREATE VIEW ProductInstock(Product, TotalQuantity) AS  
SELECT ProductID, Sum(Quantity)  
FROM Instocks  
GROUP BY ProductID
```

Notice: Views can modify in tuples by Insert, Delete and Update by SQL statements. But, it refer some *confuse conditions*. We do not present in here.

6.12. What is an index?

- Each row in table has a *rowID* sequence number used to identify the row.
- Therefore, we can consider a table as a list of pairs: *rowID* and *one or more column* we want to sort values; It is called *Index table*.
- Index is an additional data structure that helps speed up querying, join, and grouping operations.
- But, Index *slows down data input*, with the update and the insert statements.
- Indexes can be created or dropped with *no effect on the data*.

6.13. Create & Drop Indexes

CREATE INDEX indexName **ON** tableName column;

CREATE INDEX indexName **ON** tableName
(column1, column2, ...);

CREATE UNIQUE INDEX indexName **ON** tableName column;

DROP INDEX indexName;

e.g.

Categories:	rowID	CategoryID	CategoryName	Description
	1	c1	Rice	...
	2	c2	Electronic	...
	3	c3	Flower	...

```
CREATE INDEX CategoryNameIndex ON  
Categories(CategoryName ASC)
```

Index:	CategoryName	rowID
	Electronic	2
	Flower	3
	Rice	1

6.14. When should indexes be created

- A column contains a wide range of values.
- A column does not contain a large number of null values.
- One or more columns are frequently used together in a where clause or a join condition.

6.15. When should indexes be avoided

- The table is small.
- The columns are not often used as a condition in the query.
- The column is updated frequently.