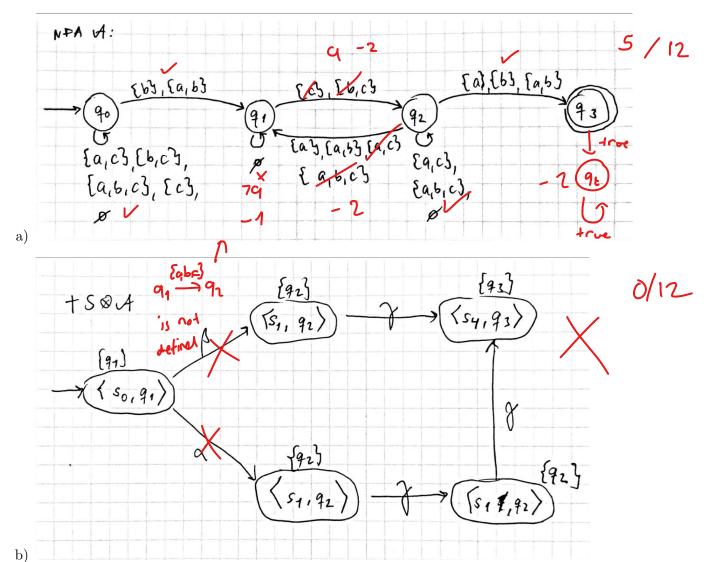
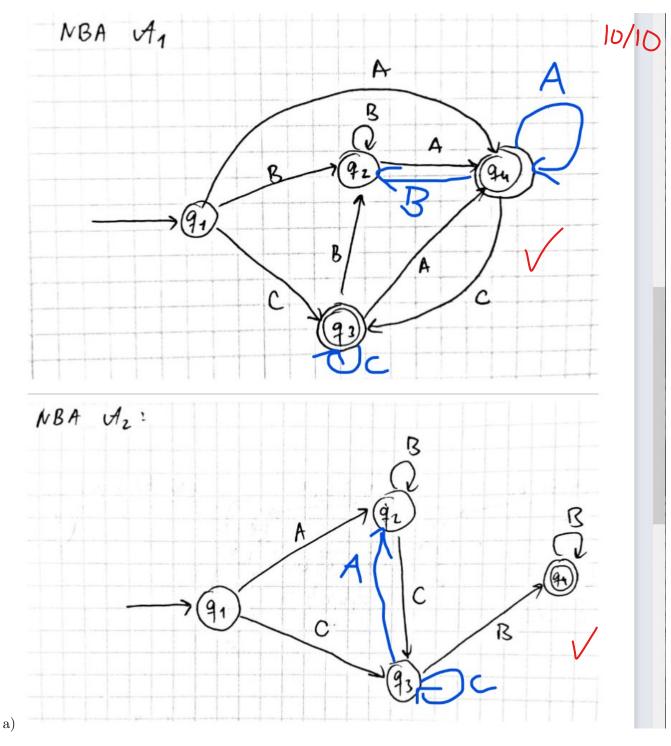
Exercise Sheet 3

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Exercise 1 (Regular Safety Property): 5/24



Exercise 2 (Closure of LT Properties): 10/32



c)

d)

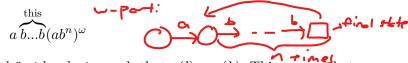
Aufgabe 3 (Finite Trace Equivalence and Safety Properties) 4/24

Consider the property: $P_{UP} := \{\{a,b\}^*(ab^k)^\omega | k \in \mathbb{N}\}$. Obviously $Pref_{fin}(P_{UP}) \supseteq \{a,b\}^*$ and thus is a safety property. However it is not a ω -regular property. Let $P_{UP} := P_{CEF}(P_{UP}) = P_{CEF}(P_{UP})$

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Let ρ be an accepting run of w on A. We know that there is a state repetition in ρ in





part of the word. In other words we can find $2 \le l < k \le n$ such that $\rho(l) = \rho(k)$. This means that

You skip here and Proceed into the

W-Port

$$\rho'(x) = \begin{cases} \rho(x) \\ \rho(x-l+k) \\ \rho'(x \mod (2n-k+l+2)) \end{cases}$$

$$\rho'(x) = \begin{cases} \rho(x) & \text{, iff. } x \leq l \\ \rho(x-l+k) & \text{, iff. } l < x \leq 2n+2-k+l \\ \rho'(x \mod (2n-k+l+2)) & \text{, otherwise } \text{, maybe confuse} \text{ tensor} \end{cases}$$

is an accepting run for the word: $(ab^{n-k+l}ab^n)^{\omega} \notin P_{UP}$ which is a contradiction.

Aufgabe 4 (Safety and Liveness Prperties) 2/20

a) Let $L_1 := (bb|c(aa)^*c)^*$ then A recognizes the language:

9_/10

 $L_1.b^{\omega} + L_1.a^{\omega}$

it ends with cu or au. b)

> however after words in La the NFA vill be in state 90 so a is not accepted directly after L1. L1.c.a is the right expression.