## Exercise 1 (Open a Bank Account)

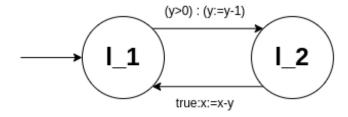
## **Exercise 2 (Transition Systems)**

- a)  $TS_1 = \{ \{s_0, s_1, s_2, s_3\}, \{\alpha, \beta, \gamma\}, \{(s_0, \alpha, s_2), (s_0, \gamma, s_1), (s_1, \gamma, s_1), (s_1, \alpha, s_3), (s_1, \beta, s_4), (s_2, \alpha, s_0), (s_2, \beta, s_4), (s_4, \alpha, s_2), (s_4, \gamma, s_3)\}, \{s_0\}, \{\{a\}, \{b\}, \{a, b\}\}, L_1\}$ with  $L_1 : \{\{s_0 \mapsto \{a\}, s_1 \mapsto \{a\}, s_2 \mapsto \{a, b\}, s_3 \mapsto \{b\}, s_4 \mapsto \{a, b\}\}.$
- b) Here is an example for a finite execution:  $\rho_{finite} = s_0 \gamma s_1 \alpha s_3$  and an example for an infinite execution:  $\rho_{infinite} = s_0 \gamma s_1 \gamma s_1 \gamma s_1 \dots$
- c) (i)  $TS_1$  is AP-deterministic, because  $|I| = |\{s_0\}| = 1 \le 1$  and there are only at most 2 states s and s' for which L(s) = L(s') holds: For these pairs  $(s_0, s_1)$  with  $L(s_0) = L(s_1) = \{a\}$  and  $(s_2, s_4)$  which  $L(s_2) = L(s_4) = \{a, b\}$  are never both in Post(s'') for all  $s'' \in S$ .

(ii)

- d)
- e)
- f)

## Aufgabe 3 (Program Graphs)



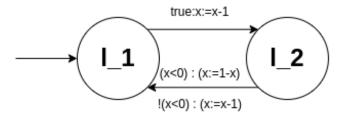


Abbildung 1:

- a)
- b)
- c)

## Aufgabe 4 (Handshaking)

- a)
- b)
- c)

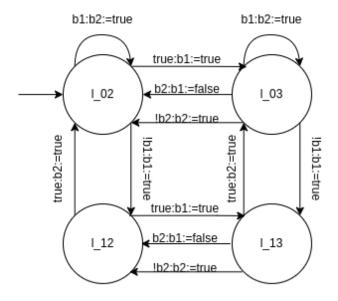


Abbildung 2:

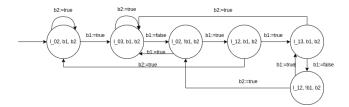


Abbildung 3: