



Embedded Operating Systems

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Sockets : Handling multiple clients

server

```
signal(SIGCHLD, sigchld_handler);
```

```
sfid = socket();
```

```
bind(sfid, ...);
```

```
listen(sfid, ...);
```

```
while(1) {
```

```
  cfd = accept(sfid, ...);
```

3 *if(accept() failed. → Interrupted SubCall.*

```
    if(cfd < 0) {
```

```
      if(errno == EINTR)
```

```
        continue;
```

```
    if(cfd > 0) {
```

```
      pid = fork();
```

```
      if(pid == 0) {
```

```
        handle-client(...);
```

```
        exit(0);
```

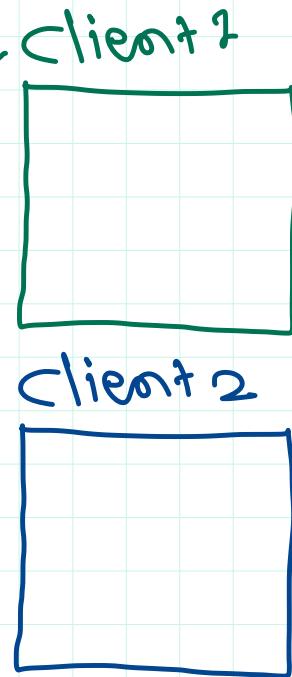
3 3

3 3

```
void sigchld_handler(..) {
```

```
  wait(&s);
```

3



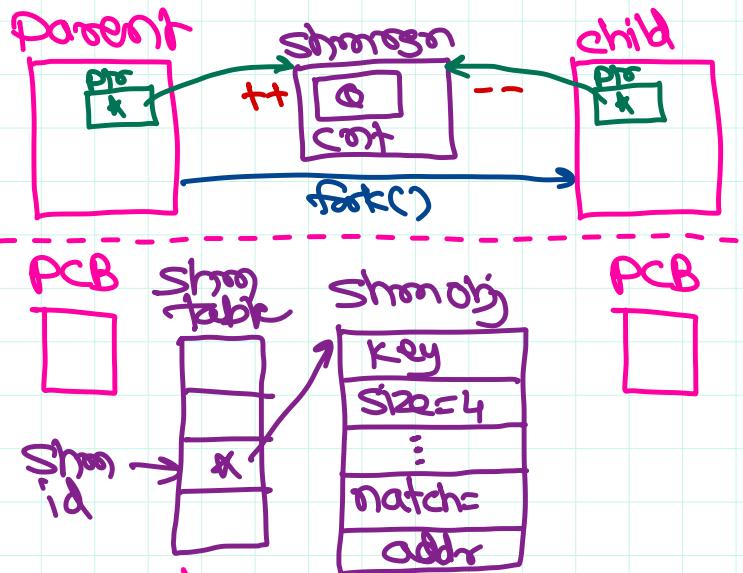
```
handle-client(cfd, ...){
```

3

```
  read(cfd, ...);
```

```
  write(cfd, ...);
```

```
  close(cfd);
```



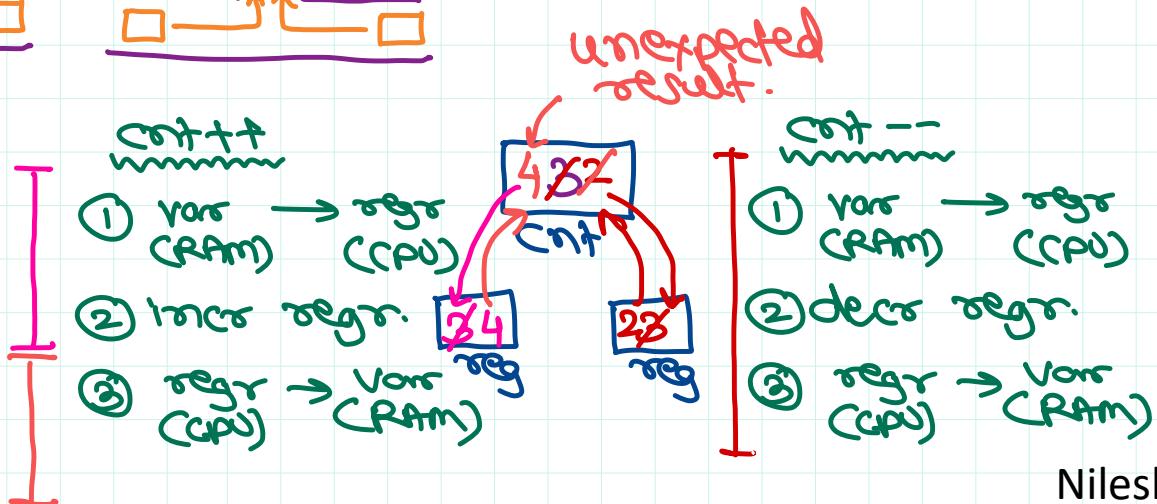
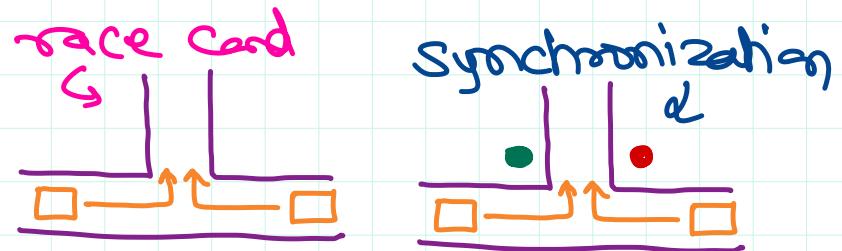
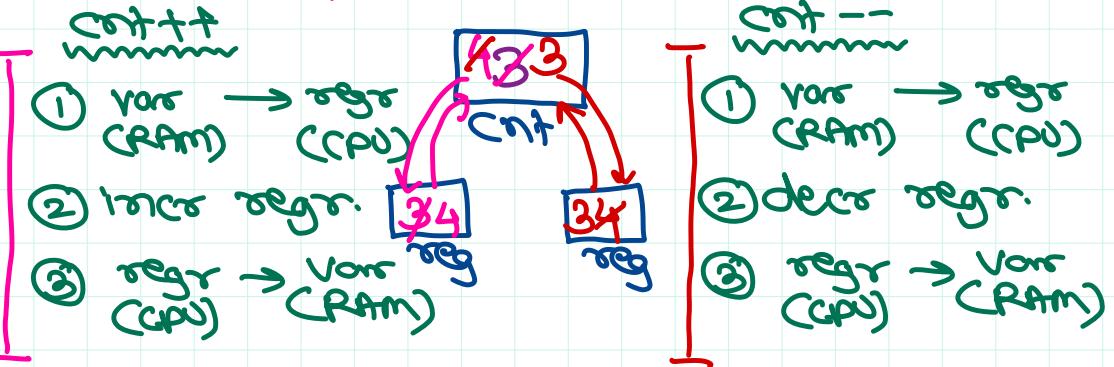
Synchronization: is a process of coordinating multiple concurrent processes to ensure that they can access shared resources without data inconsistency/corruption.

Usually this is done by blocking one or more processes while a process is using resource. Sync is done using OS sync obj/ primitives:

- ① Semaphore
- ② mutex
- ③ cond variable

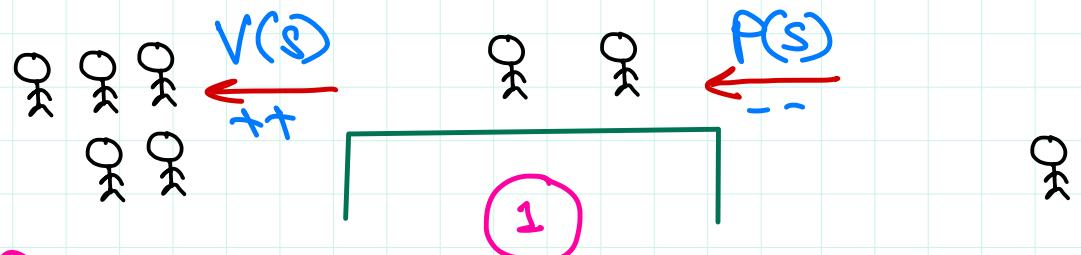
race cond: when a common resource accessed by multiple processes at the same time, it may yield unexpected results due to resource corruption.

Peterson's problem:



Semaphore

- sync object
- Dijkstra designed.
- sema is a counter.
 - decr op - wait op - P op.
 - decr cnt by 1.
 - if cnt < 0, block the current process.
 - incr op - signal op - V op
 - incr cnt by 1.
 - if one/more processes are blocked, one of the process is woken up.
- sema types
 - counting sema - count resources, availability, or processes.
 - binary sema - two states of sema (locked/unlocked)



OS scheduling → multiple queues

- ① job queue a.k.a. process table/list → all processes
- ② ready queue a.k.a. run list → processes ready for execution on CPU.
- ③ wait queue(s) → per IO device, IPC mech, Sync objs, ...

process block/sleep/wait:

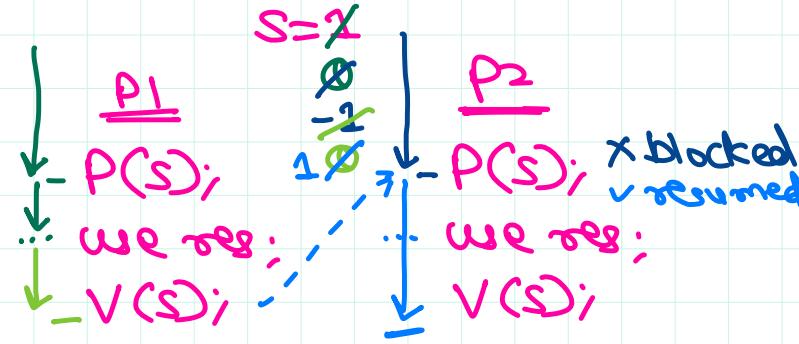
- PCB removed from ready queue
- PCB.state = waiting
- PCB is now added to wait queue.

process wake up

- PCB removed from wait queue.
- PCB.state = ready
- PCB is added to ready queue.

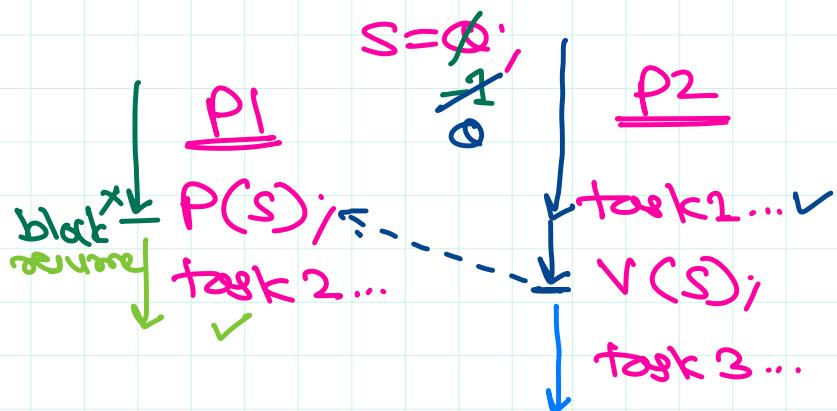
Semaphore Usage

① Mutual Exclusion



mutual exclusion \rightarrow only one process can access resource at a time.

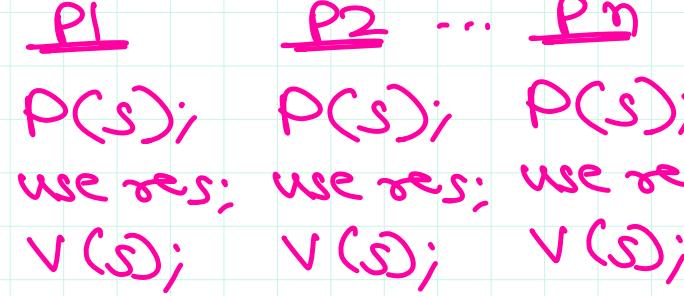
③ Flag/Event



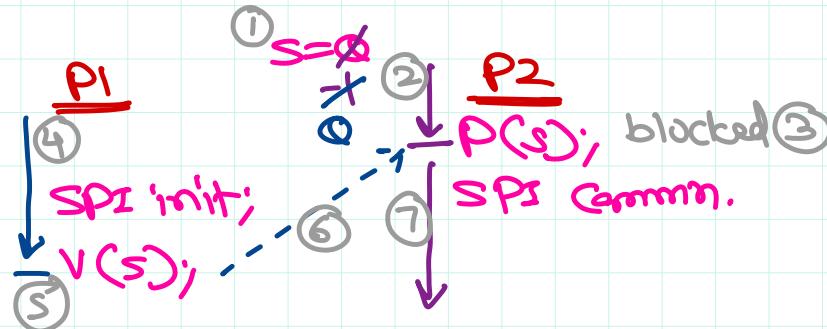
A process (P1) waits for another process (P2) to complete certain task (task1) before P1 continues its work (task2).

② Counting

$$S=n$$



task 1 (P1) \rightarrow SPI init
 task 2 (P2) \rightarrow SPI send/recv



resource \rightarrow Vart
 tasks (P1, P2, ...) \rightarrow use Vart (tx)

Px $S=1$

P(S);
 tx on Vart;
 V(S);

5 Varts to share b/w
 20 tasks.

array

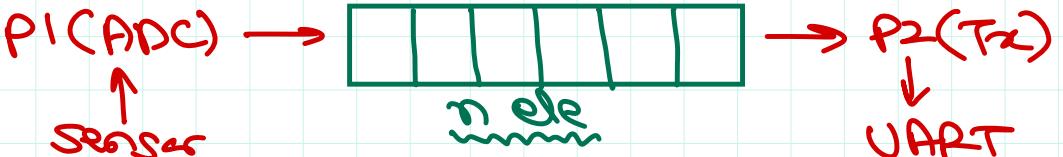
U0	U1	U2	U3	U4
g1	g1	g1	g1	g1

 Varts.

g1	g1	g1	g1	g1
inuse				

sem_avail = 5
 sem_mnt = 1

ADC task (P1) \rightarrow read ADC sensor \rightarrow producer
 UART task (P2) \rightarrow send ADC read on UART \rightarrow consumer.



- ① only producer/consumer should access buffer at a time.
- ② if buffer is full, producer should wait.
- ③ if buffer is empty, consumer should wait.

emp = 0
 mnt = 1
 sidra = 0
 sidrx = 1

fil = 0
 sidra = 2

producer

P(emp);
 P(mnt);
 get data
 write in buf
 V(mnt);
 V(fil);

consumer

P(fil);
 P(mnt);
 read the buf
 use data
 V(mnt);
 V(emp);

int acquire_vart() {
 P(sem_avail);
 P(sem_mnt);
 for(i=0; i<5; i++) {
 if(inuse[i]==0)
 inuse[i]=1;
 V(sem_mnt);
 return i;
}

void release_vart(i) { task

inuse[i]=0;
 V(sem_avail);

i = acquire_vart();
 use vart[i];
 release_vart(i);

UNIX Semaphore (Sys V)

One sema obj may have one/more sema counters.

① Create a sema `IPC_CREAT, 2`
`semid = semget(key, ncount, flag);`

② init sema
`1224` ③ `loop cont`
`count <=`
`fi1 cont`

`unsigned short initval = {1, 1, 0};`
`union semun su;`
`su.array = initval;`
`semctl(semid, 0, SETALL, su);`

③ destroy sema

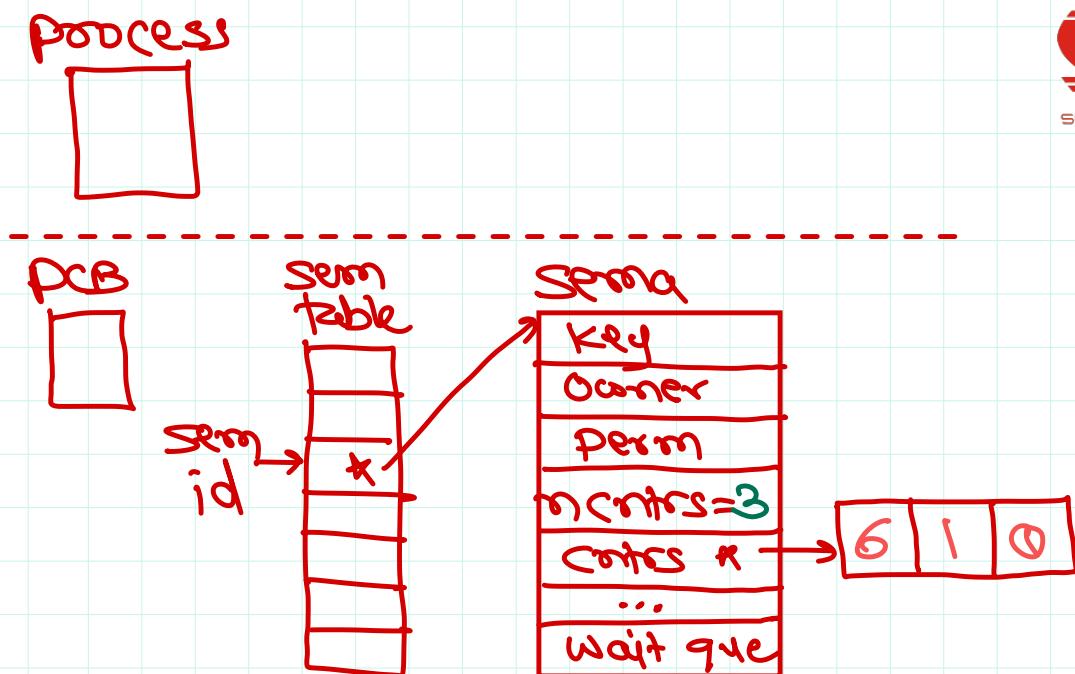
`semctl(semid, 0, IPC_RMID);`

④ Semop op - P or V

sem buf struct
- nrm = sem idx
- op = +1 or -1;
- flag = 0

V P

struct sembuf ops[2];
ops[0].sem_nrm = 0; } P(S[0])
ops[0].sem_op = -1;
ops[0].sem_flag = 0;
ops[1].sem_nrm = 1; } V(S[1])
ops[1].sem_op = +1;
ops[1].sem_flag = 0;
semop(semid, ops, 2);





Thank you!

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