



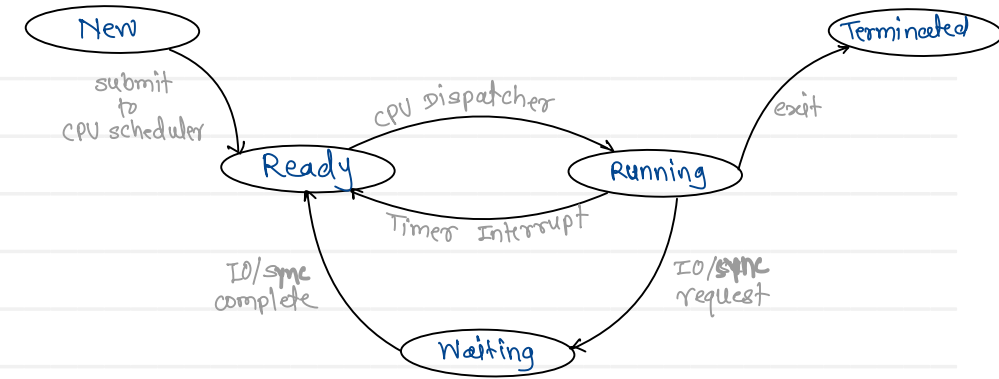
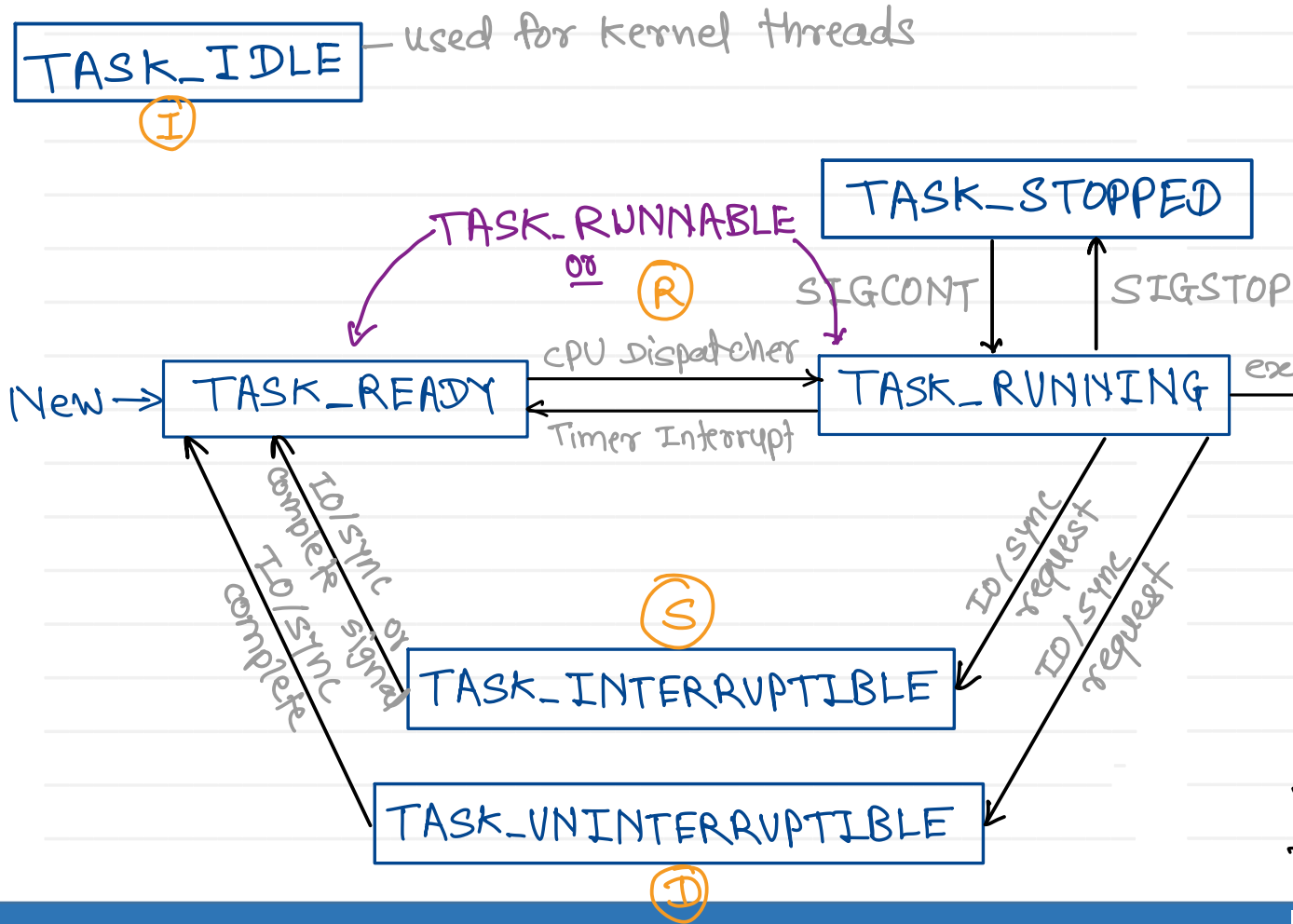
**Sunbeam Institute of Information Technology**  
**Pune and Karad**

## **Module - Embedded Operating System**

Trainer - Devendra Dhande

Email – [devendra.dhande@sunbeaminfo.com](mailto:devendra.dhande@sunbeaminfo.com)

# Linux process life cycle



```
terminal> ps          - process of current terminal
terminal> ps -e /-A    - all system processes
terminal> ps -A -o pid,ppid,state,cmd
```

## exit() vs \_exit()

### exit()

- C library function
- stdio.h
- internally calls \_exit()
- & perform C runtime cleanup

### \_exit()

- system call
- unistd.h
- release process resources (memory, file....) and writes exit status in its PCB

## wait() vs waitpid()

### wait()

- 1> block execution of current process until one of its child is terminated.
- 2> after child is terminated, get its exit status from its PCB & put it into out parameter.
- 3> release PCB of child.

### waitpid()

- similar to wait(), but wait for a given child (pid - arg1) & given flags (default=0)

$\text{waitpid}(-1, \&s, 0) \Leftrightarrow \text{wait}(\&s)$

↑                      ↑  
any child          no flags

main.out

Exe headers
T
D
sym tbl

exec()

exec0/.out

S
P
D
T

fork()

exec0/.out

<del>S</del> S'
<del>P</del> P'
<del>D</del> D'
<del>T</del> T'

exec0/.out

Exe headers
T
D
sym tbl

Loaders

pid=x
EC

PCB

Parent process

pid=y
<del>EC</del> EC

PCB

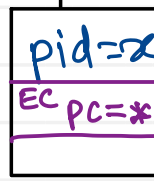
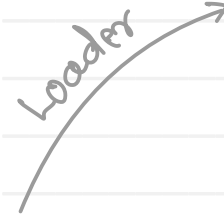
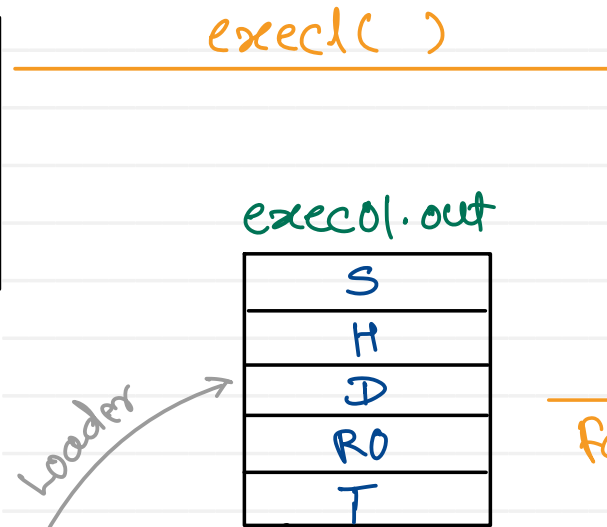
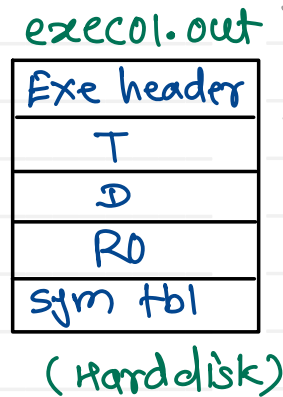
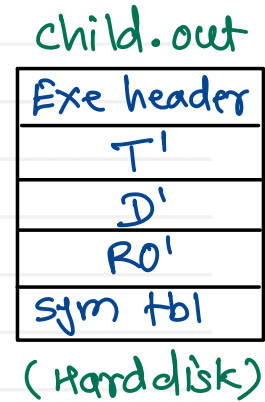
child process

# exec() syscall

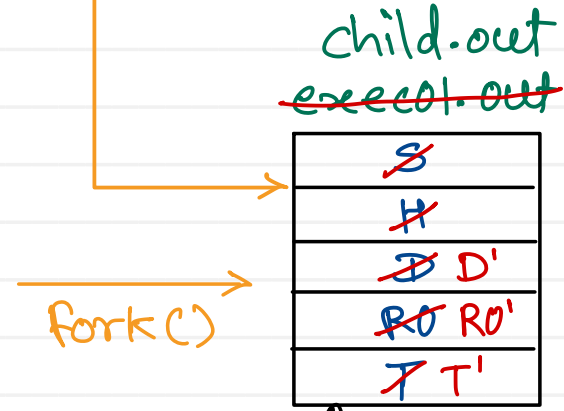
- exec() loads a new program in calling process memory by replacing old program image.

```
#include <stdio.h>
#include <unistd.h>
#include <sys/wait.h>
```

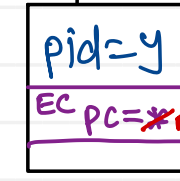
```
int main() {
    int pid, ret, s;
    pid = fork();
    if (pid == 0) {
        ret = exec("/path/of/child", args);
        if (ret < 0) {
            printf("exec() failed");
            _exit(-1);
        }
    }
    else
        waitpid(ret, &s, 0);
    return 0;
}
```



PCB  
Parent process



fork()



PCB  
child process

addr of entry point function of child program

shell\$ ls -l -i -a -h

program : ls

cmd line args : ls, -l, -i, -a, -h, NULL

argv	
0	ls
1	-l
2	-i
3	-a
4	-h
5	NULL

```
int execve (pathname, argv[], envp[])
```

pathname : path of program  
argv : array of cmd line args  
envp : array of env vars.

## exec() lib functions(3)

- 1) `execl()`
- 2) `execv()`
- 3) `execlp()`
- 4) `execvp()`
- 5) `execle()`
- 6) `execvpe()`

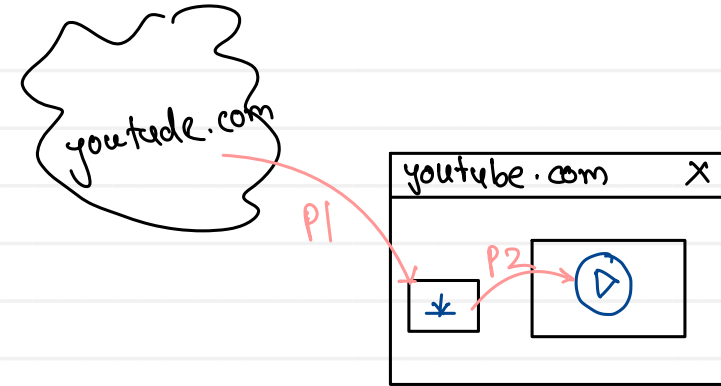
## exec() sys call(2)

- 7) `execve()`

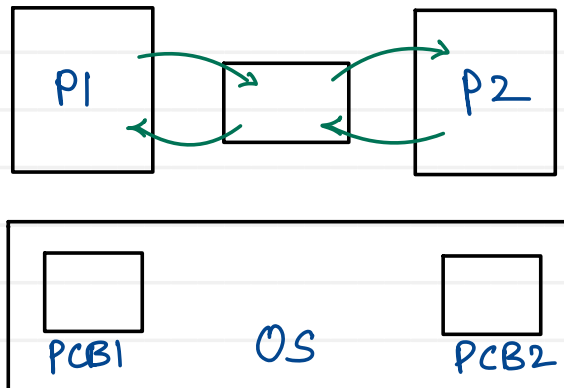
- l - variable arg list
- v - variable arg vector/array
- p - find exe in PATH var automatically.
- e - env parameters

1. `ret = execl("exePath", "exeName", "arg1", ..., NULL);`
2. `char *args[] = { "exeName", "arg1", ..., NULL };`  
`ret = execv("exePath", args);`
3. `ret = execlp("exeName", "exeName", "arg1", ..., NULL);`
4. `char *args[] = { "exeName", "arg1", ..., NULL };`  
`ret = execvp("exeName", args);`
5. `char *envp[] = { "K1=V1", "K2=V2", ..., NULL };`  
`ret = execl("exePath", "exeName", "arg1", ..., NULL, envp);`
6. `ret = execve("exePath", args, envp);`
7. `ret = execvpe("exeName", args, envp);`

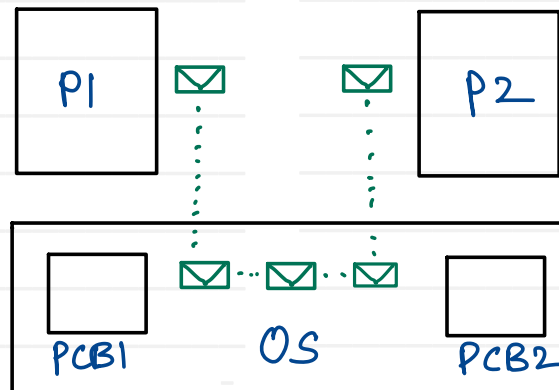
- single system can have multiple processes running together (multitasking)
- those processes can communicate / share / exchange information through IPC mechanisms.
- OS provides few ways to do comm<sup>n</sup>.



i) Shared Memory model



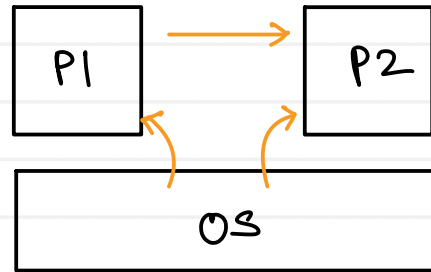
ii) Message Passing Model



- Linux IPC
- i) 1. Shared Memory
  2. Signal
  3. Message queue
  4. pipe
  5. socket



- set of predefined signals which processes need to follow.
  - approximate 64 signals are available
- terminal > kill -l → list of all signal



**kill/pkill** - used to send signals to the processes

terminal > kill -signo PID > process ID

terminal > kill -signame PID

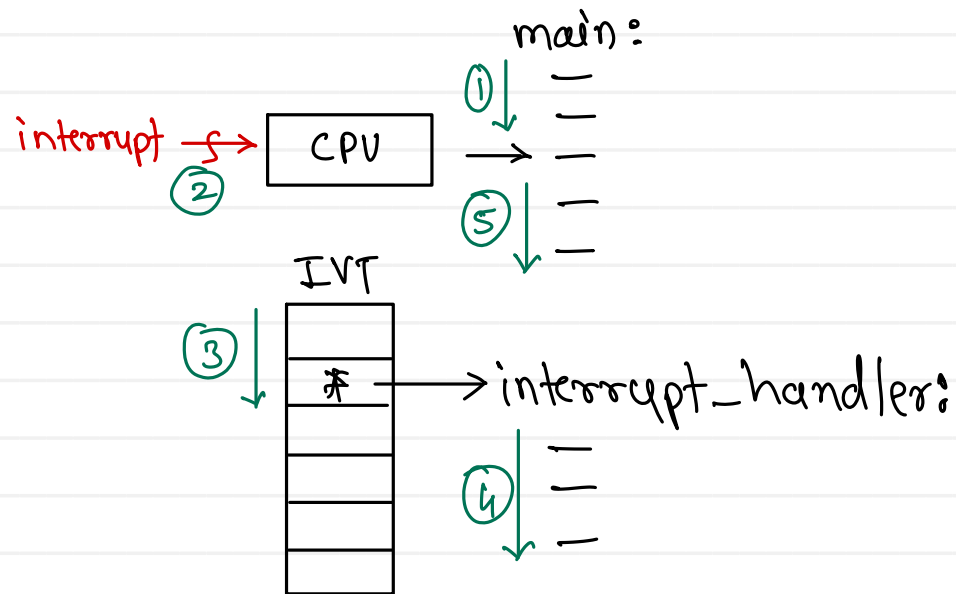
terminal > pkill -signo cmd > process name

terminal > pkill -signame cmd

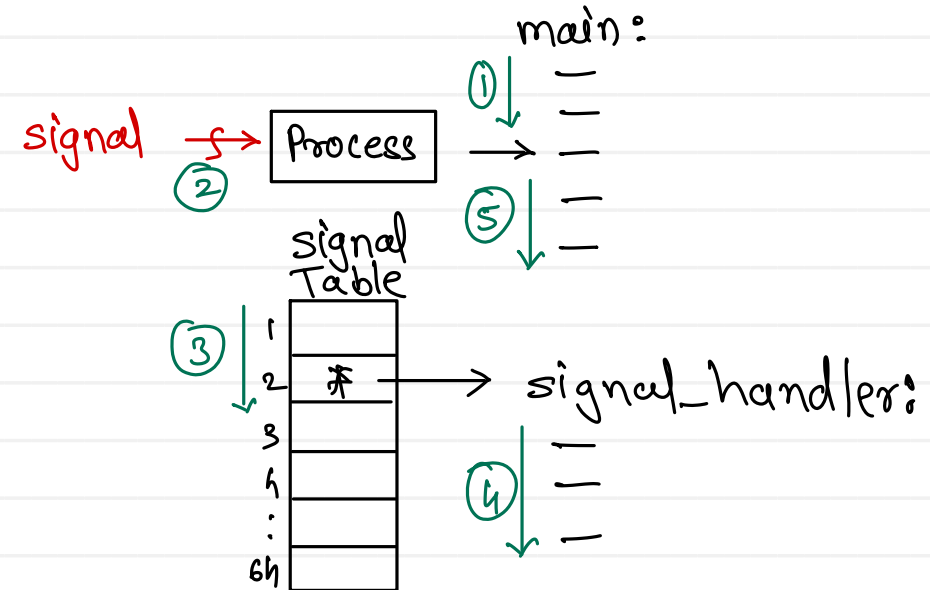
- both commands internally calls kill() system call
- kill(pid, sig)

- 1) SIGHUP - send to processes of terminal when we close it
- 2) SIGINT - used to interrupt process  
ctrl + c : terminate
- 9) SIGKILL - used to terminate forcefully.
- 11) SIGSEGV - on invalid memory access  
core dump
- 15) SIGTERM - used to terminate process  
(process gets chance to release resources)
- 17) SIGCHLD - when child terminates, it is given to parent (ignore)
- 19) SIGSTOP - pause/suspend the process  
ctrl + s : stopped
- 18) SIGCONT - resume the process  
ctrl + q : continued.

- signals are software counter part of hardware interrupts



step 1: implement interrupt handler  
step 2: register your interrupt handler



step 1: implement signal handler  
`void my_sigHandler(int);`  
 step 2: register your signal handler  
`signal( SigNum, my_sigHandler)`



Thank you!!!

Devendra Dhande

[devendra.dhande@sunbeaminfo.com](mailto:devendra.dhande@sunbeaminfo.com)