# RBE550 Motion Planning Bugs and Potential Fields



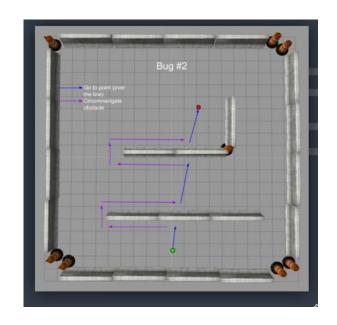
Constantinos Chamzas www.cchamzas.com www.elpislab.org

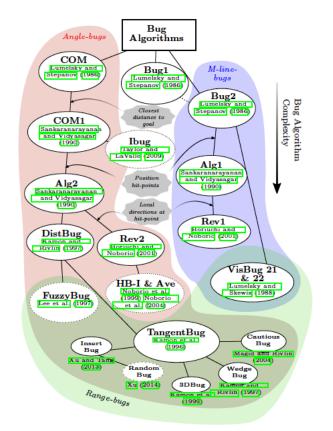
## Acknowledgements

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#### **Overview**

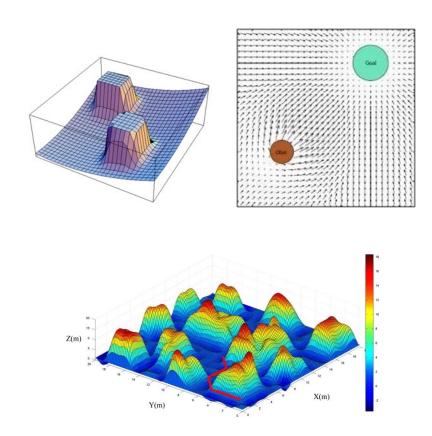
#### **Bug Algorithms**





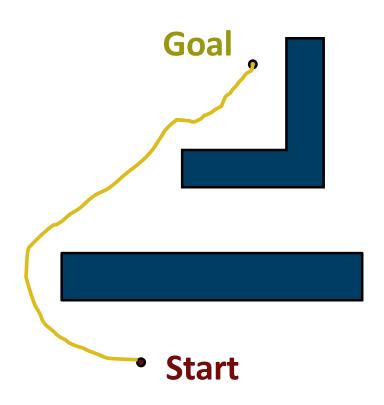
From McGuire at al, 2018

#### **Potential Fields**



## **Motion Planning Definition (Point Robot)**

- Workspace: The environment in which the robot operates
  - Notation: W
  - Generally, either 2- or 3-Dimensional Euclidean space,
  - For 2-D point robot  $W = \mathbb{R}^2$ 
    - Obstacle i in workspace W is denoted by WOi
    - Free workspace:  $W_{\text{free}} = W \setminus \bigcup_i W O_i$
- Configuration Space: a complete specification of the robot's state (informal definition)
  - Notation: Q or X or C
  - Dimension generally depends on the robot
  - For point robot  $Q = W = \mathbb{R}^2$  and
  - Collision-free space  $Q_{\text{free}} = W_{\text{free}}$ , and  $q_{start}$ ,  $q_{goal} \in Q_{free}$
- Robot path from  $q_{start}$  to  $q_{goal}$ : A continuous curve  $\sigma(t)$  such that  $\{\sigma(t) \in Q_{free} \mid \sigma(0) = q_{start}, \sigma(1) = q_{goal}\}$



## Think Like a Bug

- How would you navigate like a bug?
  - You can smell the food
  - You can't see the walls but you can feel them

Can you come up with an algorithm?





## **Assumptions**

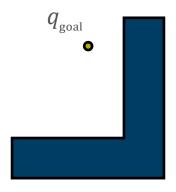
- Known direction to goal:
  - Robot can measure distance d(x, y) between points  $x, y \in W$
- Local sensing :
  - Sense walls/obstacles
  - Know position all the time (perfect encoders)
- Reasonable world:
  - Finitely many obstacles in finite area
  - A line will intersect an obstacle finitely many times
  - Workspace is bounded





## **Bugginner Strategy**

#### "Bug 0" Algorithm





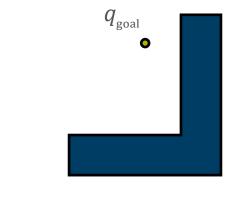
- Known direction to goal
- Local sensing walls/obstacles & encoders
- Reasonable world





## **Bugginner Strategy**

#### "Bug 0" Algorithm





#### Assumptions:

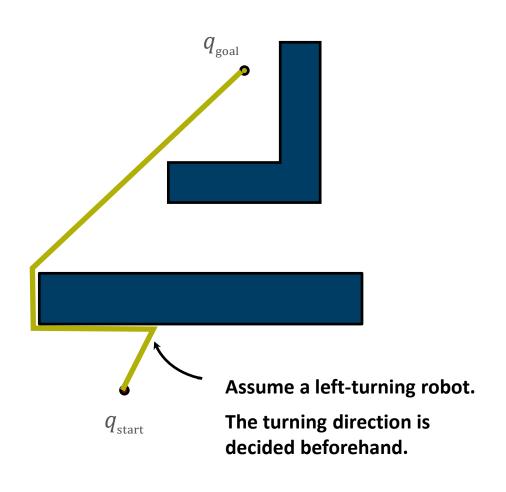
- Known direction to goal
- Local sensing walls/obstacles & encoders
- Reasonable world

#### Strategy:

- 1. Head toward goal
- 2. Follow obstacles until you can follow the goal again
- 3. Continue

Path?





#### "Bug 0" Strategy:

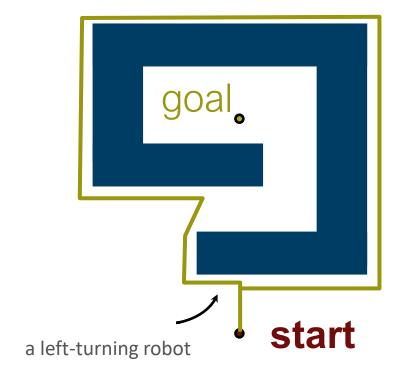
- 1. Head toward goal
- 2. Follow obstacles until you can follow the goal again
- 3. Continue

Done?



## **Bug Zapper**

#### Will "Bug 0" Strategy: work here?

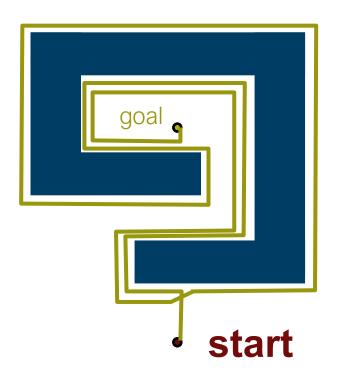


#### "Bug 0" Strategy:

- 1. Head toward goal
- 2. Follow obstacles until you can follow the goal again
- 3. Continue

How can we make it smarter?

- Add memory!

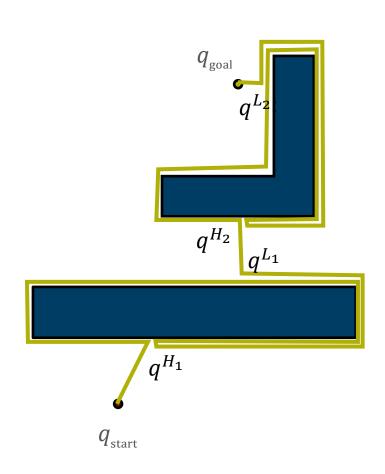


#### Assumptions:

- Known direction to goal
- Local sensing walls/obstacles & encoders
- Reasonable world
- Has Memory

#### Strategy:

- 1. Head toward goal
- 2. If an obstacle is encountered, circumnavigate it and remember how close you get to the goal
- 3. Then return to that closest point (by wall-following) and continue



#### "Bug 1" Strategy:

- 1. Head toward goal
- 2. If an obstacle is encountered, circumnavigate it and remember how close you get to the goal
- 3. Then return to that closest point (by wall-following) and continue

#### Some notation:

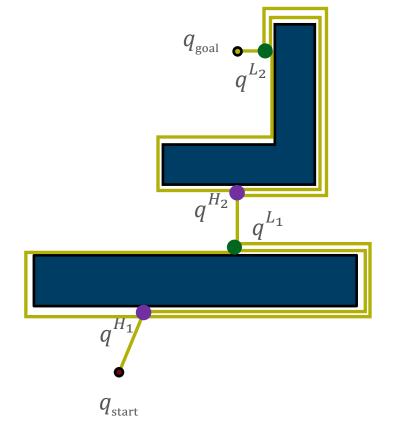
- Start and goal positions:  $q_{\text{start}}$  and  $q_{\text{goal}}$
- "hit point":  $q^{H_i}$
- "leave point":  $q^{L_i}$
- Path: a sequence of hit/leave pairs bounded by  $q_{\rm start}$  and  $q_{\rm goal}$



Algorithm: Bug 1

## "Bug 1" Algorithm pseudocode

```
Input: A point robot with a tactile sensor
Output: A path to the q_{goal} or a conclusion no such path exists
     Let q^{L_0} = q_{\text{start}}; i = 1
     Repeat
         Repeat
             from q^{L_{i-1}} move toward q_{goal}
         Until goal is reached or obstacle encountered at q^{H_i}
         If goal is reached,
              exit
         Repeat
              follow boundary recording point q^{L_i} with shortest distance to goal
         Until q_{goal} is reached or q^{H_i} is re-encountered
         If goal is reached
              exit
         Go to a^{L_i}
         If move toward q_{goal} moves into obstacle
              exit with failure
         Else
             i=i+1
              continue
```



 $q^{H_i}$ : hit point

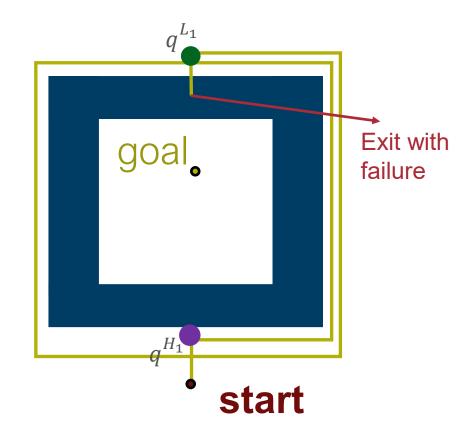
 $q^{L_i}$ : leave point



Algorithm: Bug 1

## Failure of "Bug 1" Algorithm

```
Input: A point robot with a tactile sensor
Output: A path to the q_{goal} or a conclusion no such path exists
     Let q^{L_0} = q_{\text{start}}; i = 1
     Repeat
         Repeat
             from q^{L_{i-1}} move toward q_{goal}
         Until goal is reached or obstacle encountered at q^{H_i}
         If goal is reached,
              exit
         Repeat
              follow boundary recording point q^{L_i} with shortest distance to goal
         Until q_{goal} is reached or q^{H_i} is re-encountered
         If goal is reached
              exit
         Go to a^{L_i}
         If move toward q_{goal} moves into obstacle
              exit with failure
         Else
             i=i+1
              continue
```

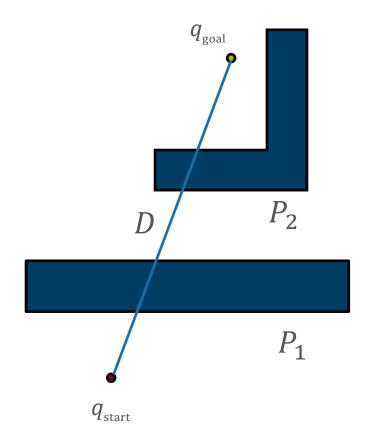


 $q^{H_i}$ : hit point

 $q^{L_i}$ : leave point

## "Bug 1" Analysis

Bug 1: Path Bounds



# What are upper/lower bounds on the path length that the robot takes?

D = straight-line distance from start to goal  $P_i$  = perimeter of the i<sup>th</sup> obstacle

#### Lower bound:

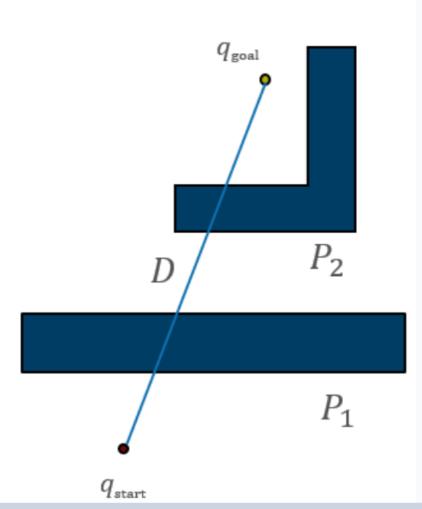
What's the shortest distance it might travel?

## Upper bound:

What's the longest distance it might travel?

#### What is the lower bound for the path that bug1 can take

### Bug 1: Path Bounds



D

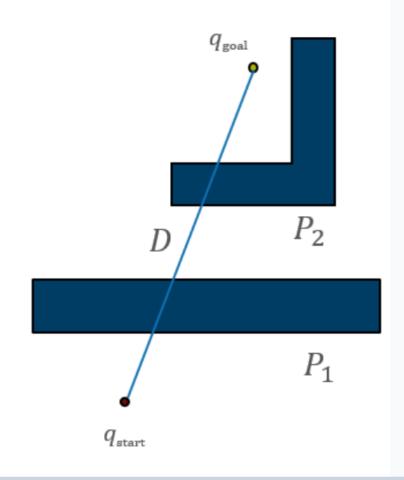
 $P_1 + P_2$ 

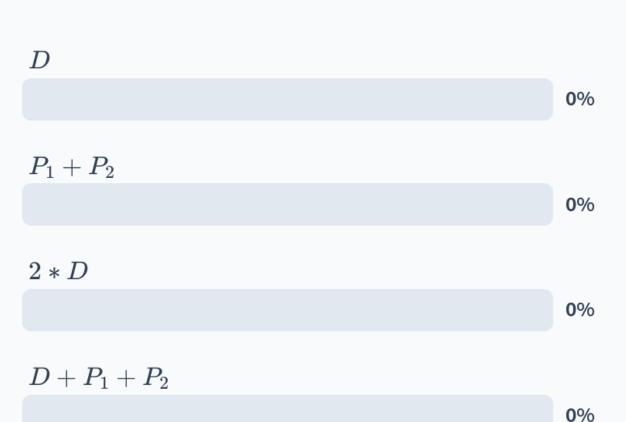
2\*D

$$D + P_1 + P_2$$

#### What is the lower bound for the path that bug1 can take

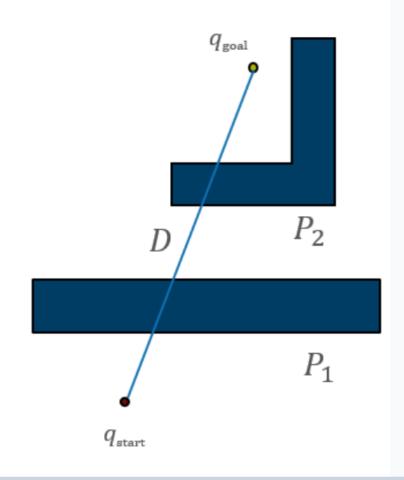
#### Bug 1: Path Bounds

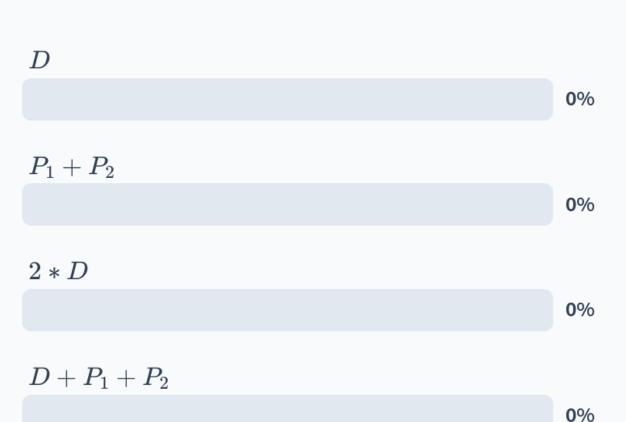




#### What is the lower bound for the path that bug1 can take

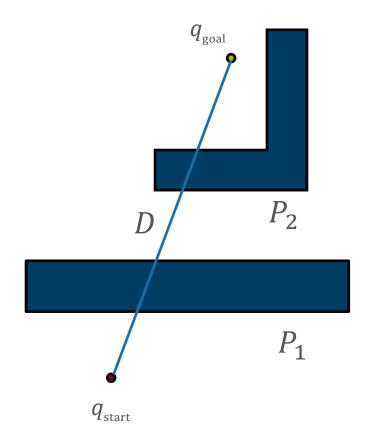
#### Bug 1: Path Bounds





## "Bug 1" Analysis

Bug 1: Path Bounds



# What are upper/lower bounds on the path length that the robot takes?

D = straight-line distance from start to goal  $P_i$  = perimeter of the i<sup>th</sup> obstacle

#### Lower bound:

What's the shortest distance it might travel?

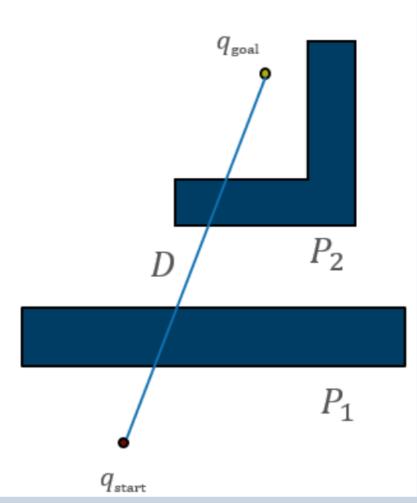
D

## Upper bound:

What's the longest distance it might travel?

#### What is the upper bound for the path that bug1 can take

### Bug 1: Path Bounds



3D

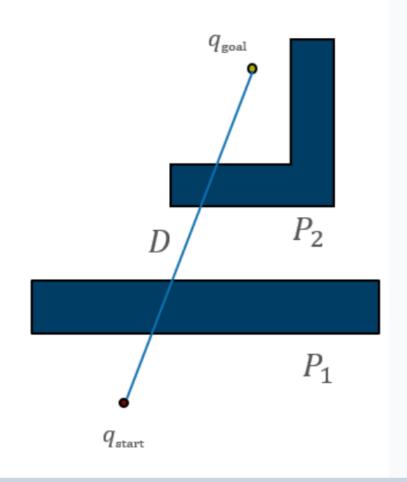
$$P_1 + P_2$$

$$1.5(P_1 + P_2) + D$$

$$2.5(P_1 + P_2) + D$$

#### What is the upper bound for the path that bug1 can take

#### Bug 1: Path Bounds



$$P_1 + P_2$$

$$1.5(P_1 + P_2) + D$$

$$2.5(P_1+P_2)+D$$

0%

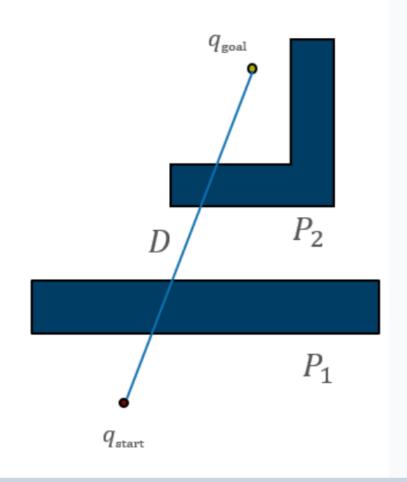
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#### What is the upper bound for the path that bug1 can take

#### Bug 1: Path Bounds



$$P_1 + P_2$$

$$1.5(P_1 + P_2) + D$$

$$2.5(P_1+P_2)+D$$

0%

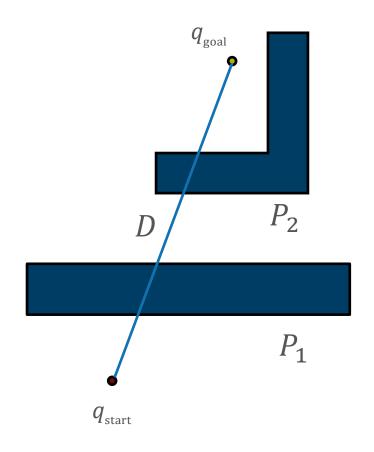
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## "Bug 1" Analysis

Bug 1: Path Bounds



#### What are upper/lower bounds on the path length that the robot takes?

*D* = straight-line distance from start to goal  $P_i$  = perimeter of the  $i^{th}$  obstacle

#### Lower bound:

What's the shortest distance it might travel?

# Upper bound: $D + \frac{3}{2}\sum P_i$ What's the longest

distance it might travel?

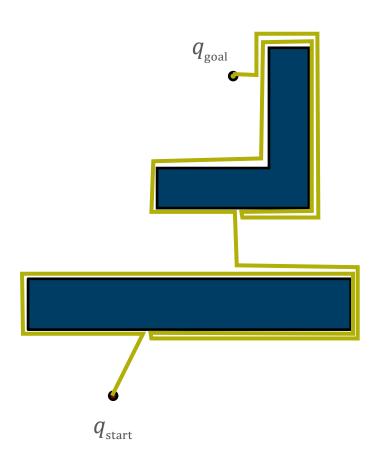
$$D + \frac{3}{2}\sum P_i$$



## Is "Bug 1" Complete?

- Definition (Completeness):
  - an algorithm is *complete* if, in finite time, it finds a path if such a path exists or terminates with failure if it does not
- Prove Bug 1 is complete (Proof Sketch):
  - Suppose BUG1 was incomplete, and there is a path from start to goal
  - BUG1 does not find it
    - Either it never terminates case1, or, it spends an infinite amount of time
    - Suppose it never terminates
      - But each leave point is closer to the goal than corresponding hit point
      - Each hit point is closer than the last leave point
      - Thus, there are a finite number of hit/leave pairs; after exhausting them, the robot will proceed to the goal and terminate
    - Suppose it terminates (incorrectly). Then, the closest point after a hit must be a leave where it would have to move into the obstacle
      - But, then line from robot to goal must intersect object even number of times (Jordan curve theorem)
      - But then there is another intersection point on the boundary closer to object. Since we assumed there is a path, we must have crossed this point on boundary which contradicts the definition of a leave point.

## Can we make something better than "Bug1"?

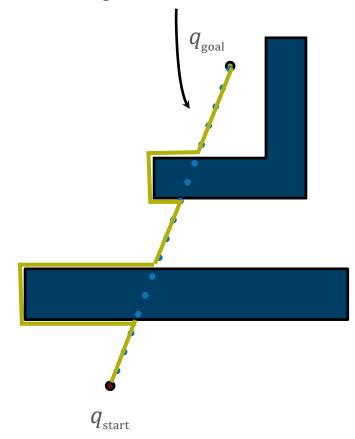


#### "Bug 1" Strategy:

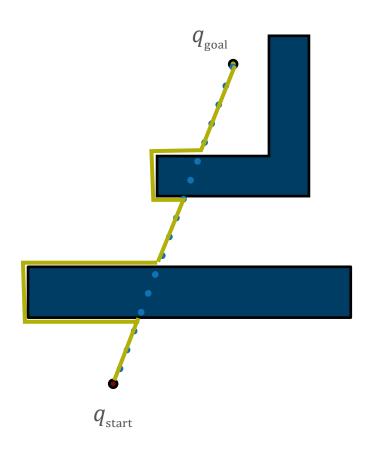
- 1. Head toward goal
- 2. If an obstacle is encountered, circumnavigate it and remember how close you get to the goal
- 3. Then return to that closest point (by wall-following) and continue

## A better bug?

Call the line from the starting point to the goal the *m-line* 



- 1. head toward goal on the m-line
- 2. if an obstacle is in the way, follow it until you encounter the m-line again
- 3. leave the obstacle and continue toward the goal

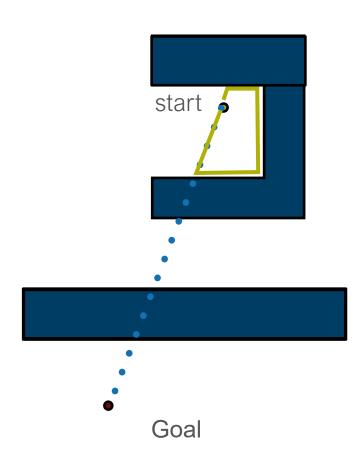


#### "Bug 2" Strategy:

- 1. Head toward goal on the m-line
- 2. If an obstacle is in the way, follow it until you encounter the m-line again
- 3. Leave the obstacle and continue toward the goal

Done?





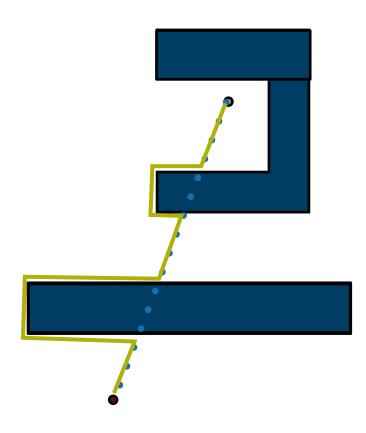
#### "Bug 2" Strategy:

- 1. Head toward goal on the m-line
- 2. If an obstacle is in the way, follow it until you encounter the m-line again
- 3. Leave the obstacle and continue toward the goal

## Done?

NO! How do we fix this?



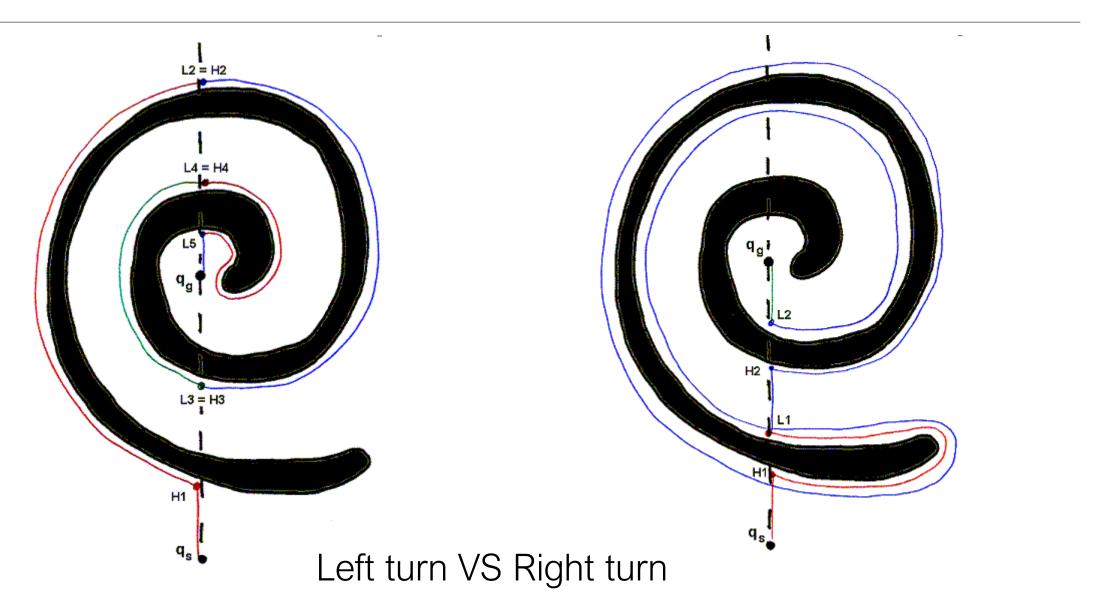


#### "Bug 2" Strategy:

- 1. Head toward goal on the m-line
- 2. If an obstacle is in the way, follow it until you encounter the m-line again closer to the goal
- 3. Leave the obstacle and continue toward the goal

Better or worse than Bug1?







Algorithm: Bug 2

```
Input: A point robot with a tactile sensor
Output: A path to the q_{goal} or a conclusion no such path exists
     Let q^{L_0} = q_{\text{start}}; i = 1
     Repeat
         Repeat
              from q^{L_{i-1}} move toward q_{goal} along the m-line
         Until goal is reached or obstacle encountered at q^{H_i}
         If goal is reached
              exit
         Repeat
              follow boundary
         Until q_{\text{goal}} is reached or q^{H_i} is re-encountered or m-line is re-encountered, x is not q^{H_i},
          d(x, q_{\text{goal}}) < d(q^{H_i}, q_{\text{goal}}) and way to goal is unimpeded
         If goal is reached
              exit
         If q^{H_i} is reached
              exit with failure
         Else
              i=i+1
              continue
```

 $q^{H_i}$ : hit point

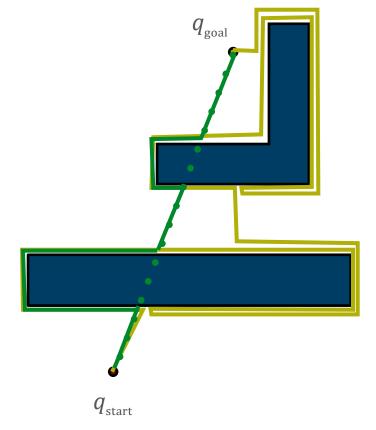
 $q^{L_i}$ : leave point



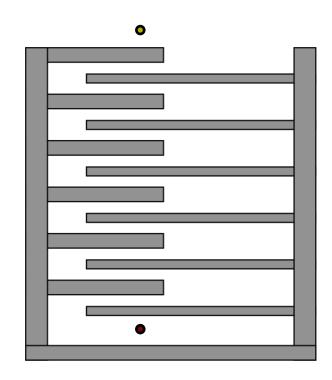
## "Bug 1" vs "Bug 2"

Worlds in which Bug 2 does better than Bug 1 (and vice versa).

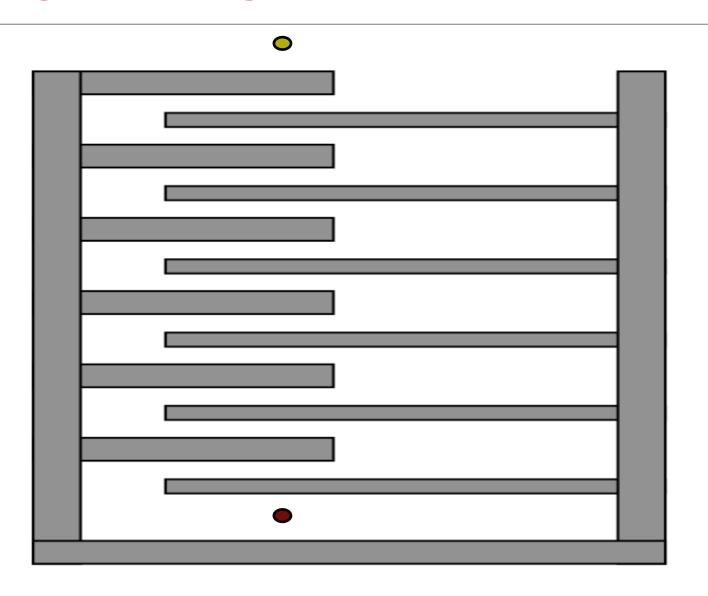
Bug 2 beats Bug 1



Bug 1 beats Bug 2

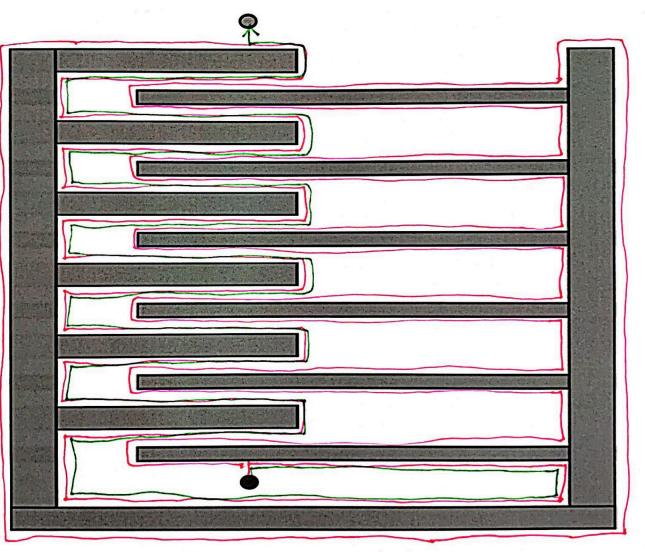


## Bug1 Vs Bug 2





## Bug 1

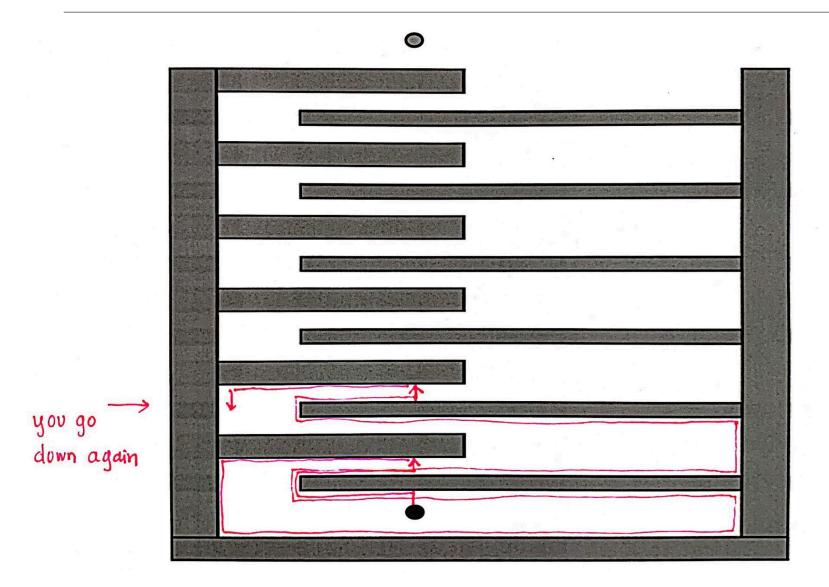


Red Line is the circumnavigation

Green Line is going to the closest recorded point

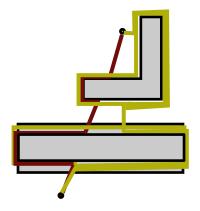


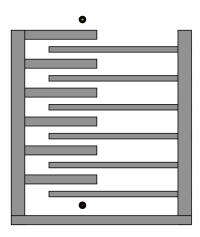
## **Bug2 – partial strategy**



## "Bug 1" VS "Bug 2"

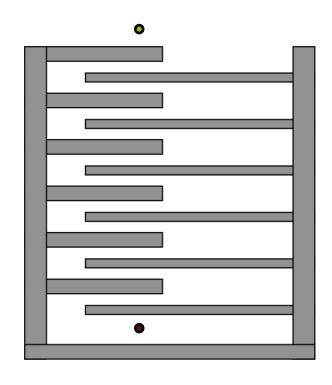
- "Bug 1" is an exhaustive search algorithm
  - it looks at all choices before committing
- "Bug 2" is a greedy algorithm
  - it takes the first thing that looks better
- In many cases, "Bug 2" will outperform "Bug 1", but
- "Bug 1" has a more predictable performance overall





# "Bug 2" Analysis

Bug 2: Path Bounds



#### What are upper/lower bounds on the path length that the robot takes?

*D* = straight-line distance from start to goal  $P_i$  = perimeter of the i<sup>th</sup> obstacle

#### Lower bound:

What's the shortest distance it might travel?

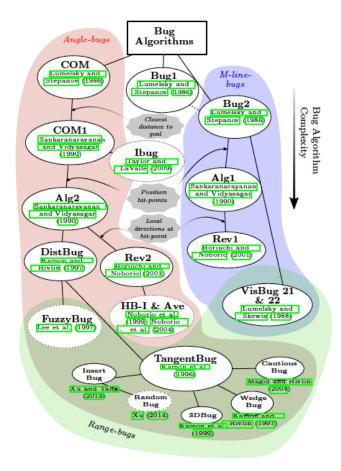
Upper bound:  $D + \frac{1}{2}\sum n_i P_i$ What's the longest

distance it might travel?

$$D + \frac{1}{2} \sum n_i P_i$$

 $n_i$ : # of m-line intersections with the *i*-th obstacle

# There is a whole family of Bugs out there

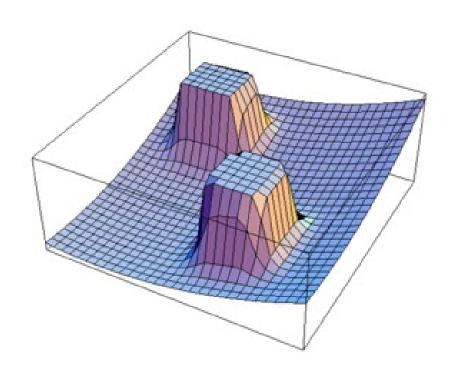


From McGuire at al, 2018

## **Summary**

- Bug 1: safe and reliable
- Bug 2: better in some cases; worse in others
- Overall there are some issues:
  - Knowing exactly where the boundary is
  - Being able to follow it safely
  - Non optimal solutions
  - Applies only to simple robots

### **Potentials Fields**



#### General Idea:

- Assume known obstacles
- Treat robot as a rolling ball
- Make mountains out of obstacles and a sink hole out of goal destination

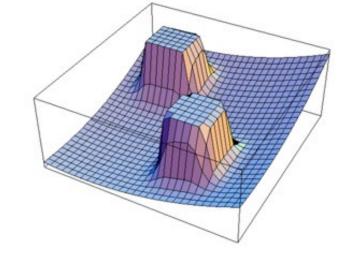
#### **Potential Function**

Potential function is a differentiable real-valued function.

$$U:\mathbb{R}^m\to\mathbb{R}$$

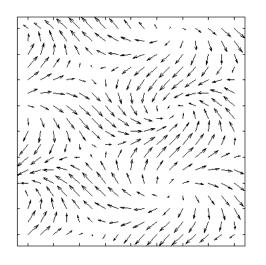
- Can be viewed as energy
- Gradient is a vector

$$\nabla U(q) = DU(q)^T = \left[ \frac{\partial U(q)}{\partial q_1}, \dots, \frac{\partial U(q)}{\partial q_m} \right]^T$$



which points in the direction that locally maximally increases U.

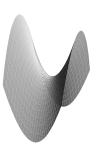
- Can be viewed as force. Here, assume 1<sup>st</sup> order dynamics, i.e., ∇ U(q) viewed as velocity.
- Can be used to define vector field.



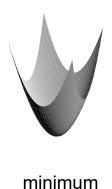
#### **Potential Function**

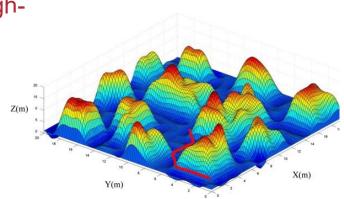
- Potential functions viewed as landscapes moving the robot from highvalue state to low-value state
- Robot moves downhill by following the negated gradient of U:
  - i.e., Gradient descent:  $\dot{c} = -\nabla U(c(t))$
- Robot terminates when it reaches at point  $q^*$  where  $\nabla U(q^*) = 0$
- Terminating point  $q^*$ , called critical point
  - Can be maximum, minimum, or saddle point





saddle



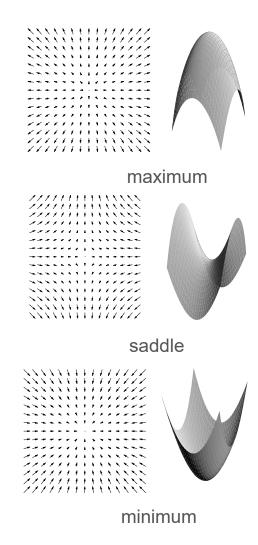


#### **Potential Function**

- Second derivative determines the type of critical point
- Hessian matrix

$$H = \begin{bmatrix} \frac{\partial^2 U}{\partial q_1^2} & \cdots & \frac{\partial^2 U}{\partial q_1 \partial q_n} \\ \vdots & \ddots & \vdots \\ \frac{\partial^2 U}{\partial q_1 \partial q_n} & \cdots & \frac{\partial^2 U}{\partial q_n^2} \end{bmatrix}$$

- When  $H(q^*)$  is non-singular,  $q^*$  is isolated
  - When  $H(q^*)$  is positive definite:  $q^*$  local minimum
  - When  $H(q^*)$  is negative definite:  $q^*$  local maximum
- · In gradient descent,
  - the robot never terminates in maximum or saddle points because, with a perturbation, it is freed.
  - Local minimum is a problem



## **Attractive Repulsive Field**

• Simplest potential function in  $Q_{\text{free}}$ 

$$U(q) = U_{att}(q) + U_{rep}(q)$$

 $U_{att}$  attracts the robot (goal) and  $U_{rep}$  repels the robot (obstacles)

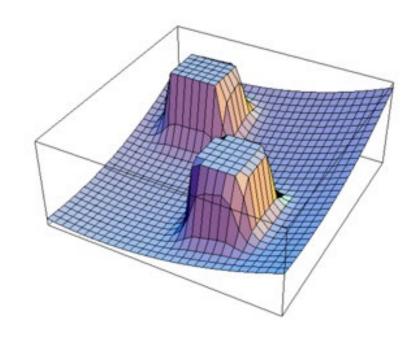
- Designing  $U_{att}$ 
  - should be monotonically increasing with distance from  $q_{
    m goal}$
  - Simplest: conic potential

$$U_{att}(q) = \xi d(q, q_{\text{goal}})$$

parameter to scale the effect of attractive force

## Repulsive Field

- Repulsive potential function  $U_{rep}$  keeps the robot from an obstacle
- Strength of  $U_{rep}$  depends on robot's proximity to an obstacle
  - The closer to an obstacle, the stronger the repulsive force
  - Define  $U_{rep}$  in terms of distance to the closest obstacle: D(q)



# **Gradient Descent Algorithms**

- Gradient descent is well-known approach to optimization problems
- Idea:
  - Start at the initial configuration
  - Small step in the direction opposite to the gradient
  - Now at a new (start) configuration
  - Repeat until gradient is zero

# **Gradient Descent Algorithm**

• Potential function  $U = U_{att} + U_{rep}$ :

```
Algorithm Gradient Descent

Input: A means to compute the gradient \nabla U(q) at a point q

Output: A sequence of points \{q(0), q(1), \dots, q(i)\}

q(0) = q_{\text{start}}
i = 0

while \nabla U(q(i)) \neq 0 do

q(i+1) = q(i) + \alpha(i) \nabla U(q(i))
i = i+1

end while
```

- q(i): value of q at iteration i
- $\alpha(i)$ : step size at the i iteration
  - Needs to be small enough not to allow "jump into obstacles"
  - Needs to be large enough not to require excessive computation time
  - The value is typically chosen ad hoc or empirically
    - e.g., based on distance to the nearest obstacle or to the goal

# **Gradient Descent Algorithm**

• Potential function  $U = U_{att} + U_{rep}$ :

# Algorithm Gradient Descent Input: A means to compute the gradient $\nabla U(q)$ at a point qOutput: A sequence of points $\{q(0), q(1), \dots, q(i)\}$ $q(0) = q_{\text{start}}$ i = 0while $\nabla U(q(i)) \neq 0$ do $q(i+1) = q(i) + \alpha(i) \nabla U(q(i))$ i = i+1end while

- Highly unlikely to exactly satisfy  $\nabla U(q(i)) = 0$
- More realistic condition:

$$\left\|\nabla U\big(q(i)\big)\right\| \leq \epsilon$$



# **Computing Distance**

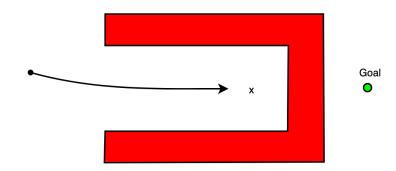
- Potential function  $U = U_{att} + U_{rep}$ 
  - Computing distance in  $U_{att}$  is simple because distance to a point
  - Computing distance in  $U_{rep}$  challenging because distance to obstacle

- For  $U_{rep}$ , can use various distance definitions, e.g.,
  - Range sensor distance
  - Discrete distance through discretization of continuous space

# Can we solve all problems now?

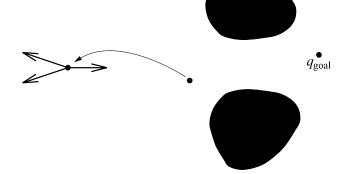
#### **Local Minima Problem**

- Problem with all gradient descent algorithms:
  - Getting stuck in local minima



- Gradient descent is guaranteed to to converge to a minimum in the field
  - No guarantee gradient descent will find a path to  $q_{\mathrm{goal}}$

 Gradient vanishes when the sum of the attractive gradient and the repulsive gradient is zero



# **Potential Fields Summary**

- Potential functions + gradient descent
  - Need careful design of these functions
    - Simplest form:  $U = U_{att} + U_{rep}$
  - Use gradient descent for motion planning
  - Challenge: computing distances
  - Problem: getting stuck in local minima

# **Bug Algorithms Potentials Fields Summary**

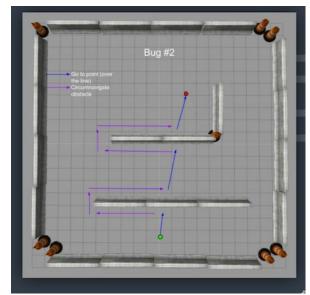
## **Bug Algorithms**

#### Pros:

- Simple to implement
- Complete
- Continuous

#### Cons:

- 2D robots
- Non Optimal Motions



#### Potential Fields

#### Pros:

- Beyond 2D
- Smooth Motions
- Continuous
- Efficient Computation

#### Cons:

- Stuck in Local Minima
- Not-Complete
- Non-Optimal Motions
- Challenging distance calculation