

Chapter 5

The Linker

Objectives

This chapter introduces the operation of the linker and gives some specifics about the CodeWarrior® Smart Linker.

5.1 Introduction

A *linker program* combines *object* files to produce an *absolute* file to be loaded into the microcontroller's memory.

In this chapter we investigate the operation of a program used to link together separately assembled and compiled source files into one file. The *linker* produces an output file to be loaded into the microcontroller using the correct types of memory for each portion of the program.

5.2 Assumptions

In this chapter we assume you are using the CodeWarrior® tool set and that you will be running the linker from within the CodeWarrior Integrated Development Environment (IDE). Further, at least until we get to Chapter 8 where we consider C programs, we assume you are developing assembly language programs. In the following sections we describe how the linker locates each of the various parts of an assembly language program into the proper memory.

5.3 Why Use a Relocatable Assembler/Compiler and a Linker?

The relocatable assembler and compiler allow efficient and effective program development using modular programming methodologies much like you use for complex C++ or other high level language programs. You may develop functional modules separately, and independently, from one another and then to *link* them to *build* the final application to be loaded into the microcontroller's memory. Thus a large project can be allocated to different members of the development team. The linker also allows you to have local labels and to expose only a smaller subset of labels to the global program and other modules. Imagine the problems you could have with duplicate labels if you took many modules built by separate groups and tried to assemble them in one giant assembly language file of many thousand lines.

5.4 Memory Types, Sections, and Section Types

In Chapter 2 we described the operation of the relocatable assembler and compiler. We also described the different types of memory, RAM and ROM, and showed that each type is used for different purposes in our embedded system programs. When using the relocatable assembler, the placement of the various parts of our programs in the different memory is called *locating*. That is, we must locate the different parts of the program in the proper memory.

Another task to be done is linking together the separate modules. The need for this arises when code in one module calls or refers to code or data elements in another. This process is also achieved by the linker.

Every program, whether written in assembly language or C, will have various *segments* or *sections*.¹ These are the *code section*, *constant data section*, *variable data section*, and *stack section*.

Code Section

The code section contains all of the program code. The program **MUST** be located in

All parts of the program that must remain when power is turned off, including the *code* and *constant data*, are placed in the code section which is then located in ROM memory.

read-only memory in an embedded system. In a development environment, where you may have a *development system* or an *evaluation board (EVB)*, you may place your program and constants into RAM. This makes it easier to change and modify as you are developing and debugging your software. Ultimately, however, you will change the physical location of the code

section to ROM for the embedded application.

Constant Data Section

Constants are to be located in ROM too.

Constant data elements are located also in ROM memory. Examples of constant data include variable data initialization values and messages to be printed.

Variable Data Section

The variable data section contains the storage locations for all variable data used in the program. There are several points to make about variable data storage used in a program:

The *variable data* section is located in RAM memory.

- All variable data storage locations must be allocated in the program by the assembler or compiler using a *define storage (DS)* directive.
- All variable data storage must be located in RAM memory.
- The program code must initialize all variable data must be initialized, if needed, when the processor is running. You **MUST** not assume any data value will be in a RAM data storage location when the system is power on.

Stack Section

As we will see in later chapters, the stack is an allocation of RAM used for temporary

The *stack* section is located also in RAM memory.

information storage. This could be data saved temporarily or return addresses for subroutines. The stack pointer register must be initialized pointing to RAM before the stack can be used.

¹ We will generally use the term section, although you will see segment used in some of the CodeWarrior files. The two terms are synonymous.

5.5 Linker Operation

Locating in Proper Memory

Before we show how to set up the linker to create your program to be loaded into the microcontroller memory, let us have a look at what it is supposed to do.

Example 5-1 and Example 5-2 show the *main* and *subroutine* programs for two separately assembled modules. This program does not do anything sensible so we will not worry particularly about the assembly language code.

Main Module

In the main program, Example 5-1, we see there are four distinct uses of memory. The order in which they appear is not important although we will see some programming style suggestions in Chapter 7. The first is a variable data section shown between lines 15 – 17. A memory allocation for a single byte of data storage is made at line 17 by the assembler directive

```
main_data:  DS.B 1.
```

The label `main_data` marks a single byte of storage and will be assigned an address during the linking process. Any other *variable* data storage you might need must be defined in this section or other named sections that can be located in RAM. The section continues until another `SECTION` directive or something other than the define storage (`DS`) directive is found. Everything in this section must be located in RAM memory.

The next memory use is found starting at line 20. This is the *code section* and it starts with the directive

```
MainCode:  SECTION.
```

Everything in this section, from line 20 to the end of the program is to be located in ROM memory. You should choose the label `MainCode` to be something meaningful so you can see where it is located when you view the linker map file.

A *constant data section*, `MainConst: SECTION`, is defined at line 30 and a constant byte is defined by

```
main_const: DC.B $44
```

at line 33. Because constants go into the same type of memory as the code, you could simply define the bytes following the program code. However, it is far better to define a constant section to allow all constants to be grouped together in the linking process or to locate them in a different part of the ROM than the code.

A fourth memory use is implied in line 12 where the label `__SEG_END_SSTACK`² is defined to be external. The value of this label will be defined by the linker and is used with the stack pointer initialization instruction at line 22 in the main program. As we will learn later, the stack pointer register must be initialized before calling any subroutines. Also, the stack itself must be located in RAM memory.

² This label starts with two underbar (__) characters.

Example 5-1 Relocatable "Main" Program

Metrowerks HC12-Assembler
(c) COPYRIGHT METROWERKS 1987-2003

Abs.	Rel.	Loc	Obj. code	Source line
----	----	-----	-----	-----
1	1			
2	2			; Rel_ex_1a_main.asm
3	3			; The main module in a relocatable example.
4	4			; This program does nothing that is sensible.
5	5			; FM Cady 7/04
6	6			; *****
7	7			; Define entry point for the main program.
8	8			XDEF Entry
9	9			; Define symbols used in the subroutine.
10	10			XREF sub_1
11	11			; *****
12	12			; Stack section
13	13			XREF __SEG_END_SSTACK
14	14			; *****
15	15			; Variable Data Section
16	16			MainData: SECTION
17	17	000000		main_data: DS.B 1
18	18			; *****
19	19			; code section
20	20			MainCode: SECTION
21	21			Entry:
22	22	000000	CFxx xx	lds #__SEG_END_SSTACK
23	23			; Initialize the variable main_data.
24	24	000003	180C xxxx	movb main_const,main_data
		000007	xxxx	
25	25	000009	B6xx xx	loop: ldaa main_data ; Get the data.
26	26	00000C	16xx xx	jsr sub_1 ; Pass to sub
27	27			; Increment data for the next loop.
28	28	00000F	72xx xx	inc main_data
29	29	000012	06xx xx	jmp loop ; Continue forever.
30	30			; *****
31	31			MainConst: SECTION
32	32			; Define constants in the code section
33	33	000000	44	main_const: DC.B \$44

Subroutine Module

Let us now look at how memory is used in the subroutine shown in Example 5-2. There are three memory uses in this example. There is a section of code starting at line 16 and continuing through line 23. We see a constant data section at line 26 and a variable data section at line 31.

Example 5-2 Relocatable Subroutine

Metrowerks HC12-Assembler
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Abs.	Rel.	Loc	Obj. code	Source line
1	1			; rel_ex_1a_sub_1.asm
2	2			; Relocatable subroutine example with
3	3			; variable and constant data.
4	4			; FM Cady 7/04
5	5			; Input Registers:
6	6			; A = some input data
7	7			; Output Registers:
8	8			; A = some output data
9	9			; Registers modified:
10	10			; A, CCR
11	11			; *****
12	12			; Define symbols used in the subroutine.
13	13			XDEF sub_1
14	14			; *****
15	15			; Code Section
16	16			SubCode: SECTION
17	17			sub_1:
18	18			; On entry, the data is in Accumulator A.
19	19			; Store the data in the subroutine.
20	20	000000	7Axx xx	staa sub_data
21	21			; Get some data to return to the main.
22	22	000003	B6xx xx	ldaa sub_const
23	23	000006	3D	rts ; Return to main program.
24	24			; *****
25	25			; Define constant data.
26	26			SubConst: SECTION
27	27	000000	33	sub_const:DC.B \$33
28	28			; *****
29	29			; Variable Data Section.
30	30			SubData: SECTION
31	31	000000		sub_data: DS.B 1

Memory Needs

Our analysis of memory requirements for the main and subroutine is summarized in Table 5-1. Notice that we have included an estimate of the size of the stack needed for this program. We will defer how we find this out to a later chapter.

Table 5-1 Program Memory Requirements

Module		ROM		RAM
Main	Code	21 bytes	Data	1 byte
	Constant Data	1 byte	Stack	2 bytes
Subroutine	Code	7 bytes	Data	1 byte
	Constant Data	1 byte		

The Hardware Memory Map

As programmers, now we know how much ROM and RAM our program will require. As embedded system engineers, we must look at the hardware upon which our program is

Software engineers must know the system memory map to be able to load their code into the microcontroller's memory.

going to be installed. We do that by looking at the *memory map* of the system. In a single-chip microcontroller, such as we described in Chapter 2, the manufacturer determines the memory map when the chip is designed and manufactured.

Figure 5-1 shows a MC9S12C32 microcontroller operating in *single-chip* mode in an embedded system with no external RAM memory.

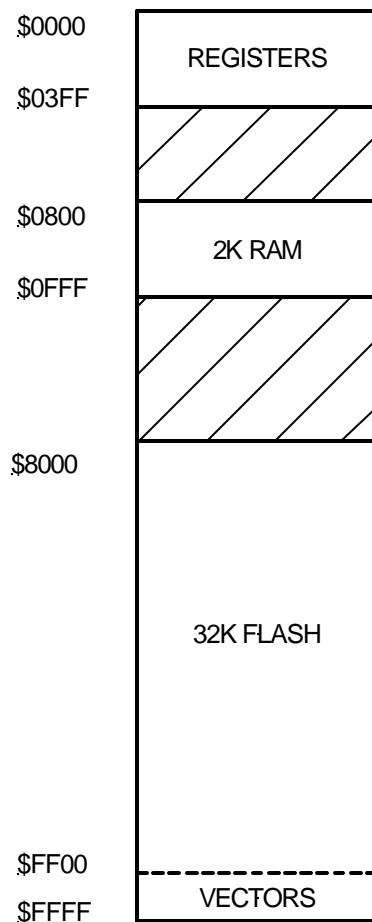


Figure 5-1 Memory Map for MC9S12C32 Microcontroller

Linking the Code

Before we show how to instruct the linker to locate our code correctly, we need to discuss the linking part of this job. Linking refers to establishing the final address for certain instructions, such as a *jsr – jump-to-subroutine*, when the starting address of the subroutine is in another module. This address resolution must be performed also when one module refers to a data location defined in another module. Line 26 in Example 5-1 has a `jsr sub_1` instruction where `sub_1` is in the second, separately assembled module. The assembler has three directives that organize the information needed to accomplish this linking task.

XREF, XDEF, and XREFB

Every XREF must have a companion XDEF in some other module.

When a symbol is used in one module but is defined in another, it is said to be an *external* symbol. In this case, the *XREF* – *External Reference* directive is used. In the module where the symbol is defined, the *XDEF* – *External Definition* directive is used. In Example 5-1 the symbol `sub_1` is the address for the jump-to-subroutine instruction on line 26. Because that symbol is not defined in the main module, it is declared to be external by `XREF sub_1` on line 10. Similarly, the `sub_1` symbol must be made visible in the subroutine module. This is done by the `XDEF sub_1` on line 13. For every XREF in our program modules there must be a partner XDEF in some other module.³

Example 5-1 shows `XDEF Entry` at line 8. This makes the code entry point (where the `Entry` label is at line 21) visible to the linker.

There is another directive, *XREFB* – *External Reference for Symbols on the Direct Page* included in the assembler directives. This is used when an external symbol is on the direct addressing page and thus fewer addressing bits are required.

5.6 The Linker Parameter File

The linker is controlled by a *linker parameter file* (*.prm). You put the memory map information shown in Figure 5-1 in the linker parameter file so the linker can locate each section properly. Example 5-3 shows a linker parameter file used to link (and locate) the two modules in Example 5-1 and Example 5-2.

The linker parameter file allows us to specify easily where code and data are located.

Example 5-3 Linker Parameter File

```
NAMES
END

SECTIONS
    RAM = READ_WRITE 0x0800 TO 0x0FFF;
/* unbanked FLASH ROM */
    ROM_C000 = READ_ONLY 0xC000 TO 0xFEFF;
END

PLACEMENT
/* Place the code and constant sections in ROM */
    .init, MainCode, SubCode, MainConst, SubConst,
    DEFAULT_ROM INTO ROM_C000;
/* Place the Variable Data */
    MyData, MySubData, DEFAULT_RAM INTO RAM;
END
/* Define the stack size needed */
STACKSIZE 0x100
/* Specify the initial reset vector entry point */
VECTOR 0 Entry /* reset vector */
/* INIT is needed for assembly applications */
INIT Entry
```

³ If you forget to declare a symbol as external with XREF, the assembler will give you an error message. If you forget the XDEF, the linker will generate an error message.

Parameter File Commands

Although the CodeWarrior linker is very powerful with many features, we will describe only those few necessary to link assembly and C program modules. Keywords and commands in the parameter file must be capitalized. Each section in the parameter file starts with the command name in capitals and ends with the **END**. All commands and qualifiers are case sensitive.

Names:

The **NAMES** command lists the files building the application. For an application linked in the CodeWarrior tool set, the files that have been "added" to the application are automatically linked. Other modules may be linked by including their names in the **NAMES** section.

Sections:

```
SECTIONS
    RAM = READ_WRITE 0x0800 TO 0x0FFF;
/* unbanked FLASH ROM */
    ROM_C000 = READ_ONLY 0xC000 TO 0xFEFF;
END
```

The **SECTIONS** part defines the memory map for your application. It allows you to assign meaningful names to address ranges. These can be used in the subsequent placement section to increase the readability of your parameter file. The syntax of a **SECTIONS** entry is the following:

`<your_name> = MEMORY_TYPE address TO address;`

where

`<your_name>` is the name you wish to use for the particular address range, and **MEMORY_TYPE** defines the type of memory as shown by the qualifiers in Table 5-2.

Table 5-2 Memory Type Qualifiers

Qualifier	Initialized Variables	Non-Initialized Variables	Constants	Code
READ_ONLY	Not applicable. ⁴	Not applicable. ⁴	Content written to target address.	Content written to target address.
READ_WRITE	Content written into copy down area with information to allow it to be copied at startup. ⁵	Area contained in zero out information. ⁶	Content written into copy down area. ^{5,6}	To allocate code in a RAM area for system development, you should declare this area as READ_ONLY . In an embedded system, this is not applicable. ⁴
NO_INIT	Not applicable. ⁴	Handled as allocated.	Not applicable. ⁴	Not applicable. ⁴
PAGED	Not applicable. ⁴	Handled as allocated.	Not applicable. ⁴	Not applicable. ⁴

⁴ These cases are not intended but the linker does allow them. When used, the qualifier controls what is written into the application.

⁵ Initialized objects and constants in **READ_WRITE** must be initialized at the start of the program at run-time. In a C program this is done by the start-up code but in an assembly language program you must ensure this is done. The copy down area contains information to initialize the constants in the start-up code.

⁶ The zero out information specifies areas that must be initialized with 0 at startup.

Placement:

```

PLACEMENT
/* Place the code and constant sections in ROM */
    .init, MainCode, SubCode, MainConst, SubConst,
    DEFAULT_ROM INTO ROM_C000;
/* Place the Variable Data */
    MyData, MySubData, DEFAULT_RAM INTO RAM;
END

```

The *PLACEMENT* section places your various sections into the proper type of memory as defined in the *SECTIONS* section. You must include *DEFAULT_RAM* and *DEFAULT_ROM* placements. The linker, by default, will place your sections into the proper placement. For example, the linker knows to place any memory allocation (DS) parts of the program into *DEFAULT_RAM* and any code and constants (DC) into *DEFAULT_ROM*. You may, for the sake of documentation readability or to control where sections are placed, "place" the named sections from your program as shown in Example 5-3. This is optional but is a good documentation practice.

Stacksize:

```

/* Define the stack size needed */
STACKSIZE 0x100

```

An important part of your code development is the allocation of storage and placement of the stack. As we will learn in Chapter 6, the stack is an area of RAM used for temporary variable storage and for return addresses for subroutines. Because both the stack and the variable data storage sections are in *READ_WRITE* memory, it is important that neither one interfere with or overwrite the other. You must estimate the amount of stack storage that is needed and specify at least this amount. Because stack overflow (using more stack memory than is allocated) can cause many problems in your running program, most embedded system engineers are very conservative and allocate plenty of stack space. In this example, *STACKSIZE* allocates \$0100 bytes, places the stack into RAM and defines the variable *__SEG_END_SSTACK* to be used in your program to initialize the stack pointer register (as shown in line 22 in Example 5-1.)

Vector:

```

/* Specify the initial reset vector entry point */
VECTOR 0 Entry /* reset vector */

```

The vector section allows you to initialize vectors. This example shows how to initialize the reset vector to allow the microcomputer to be *vectored* to the start of the program when power is turned on. We will see other vectors and their uses for interrupts in Chapter 12.

Init:

```

/* INIT is needed for assembly applications */
INIT Entry

```

When linking an assembly application, an *entry point* is required. The *INIT* command with its operand, in this case *Entry*, provides this information.

Comments:

Comments are entered into the parameter file using standard C or C++ syntax for comments.

5.7 The Linker Output Files

Two output files created by the linker are important to us after we have located and linked our object files. These are the linker map file, which shows us where the linker has placed our code, and the absolute code file to be loaded into the microcontroller's memory.

The Linker Map File

The *Linker Map File* shows us where our code, constant and variable data are located, along with other information. Following the linking controlled by the linker parameter file in Example 5-3, the linker map file shows the following sections.

The *linker map file* shows us where all elements of our program are located.

Target Section:

TARGET SECTION

```
Processor      : Motorola HC12
Memory Model   : SMALL
File Format    : ELF\Dwarf 2.0
Linker        : SmartLinker V-5.0.22 Build 4047, Feb 17 2004
```

The target section shows the processor, the memory model (more important when programming in C), the file format, and the linker version.

File Section:

FILE SECTION

rel_ex_1a_main.asm.o	Model: SMALL,	Lang: Assembler
rel_ex_1a_sub_1.asm.o	Model: SMALL,	Lang: Assembler

This lists all files that have been linked.

Startup Section:

STARTUP SECTION

Entry point: 0xC000 (Entry)

The startup section shows the entry point, i.e. where the program will start.

Section-Allocation Section:

Section Name	Size	Type	From	To	Segment
.init	21	R	0xC000	0xC014	ROM_C000
SubCode	7	R	0xC015	0xC01B	ROM_C000
MainConst	1	R	0xC01C	0xC01C	ROM_C000
SubConst	1	R	0xC01D	0xC01D	ROM_C000
MainData	1	R/W	0x800	0x800	RAM
SubData	1	R/W	0x801	0x801	RAM
.stack	256	R/W	0x802	0x901	RAM
.vectSeg0_vect	2	R	0xFFFFE	0xFFFFF	.vectSeg0

Summary of section sizes per section type:

READ_ONLY (R):	20 (dec:	32)
READ_WRITE (R/W):	102 (dec:	258)

Each of the software components that have been allocated into the various sections is described. Note that the code in the module that is first executed, that is the main module, is placed into a section named *.init*. In a C program, the startup code will be placed in *.init*. All other named sections, such as SubCode (in the subroutine), MainConst, MainData and Subdata are listed with their size, type of memory, the range of address, and the segment assignment.

Vector-Allocation Section:

VECTOR-ALLOCATION SECTION		
Address	InitValue	InitFunction
0xFFFFE	0xC000	Entry

This section shows how vectors have been initialized. We will see more of this when we discuss interrupts in Chapter 12.

Object-Allocation Section:

OBJECT-ALLOCATION SECTION							
Name	Module	Addr	hSize	dSize	Ref	Section	RLIB
MODULE: -- rel_ex_la_main.asm.o --							
- PROCEDURES:							
Entry		C000	9	9	0	.init	
loop		C009	C	12	1	.init	
- VARIABLES:							
main_data		800	1	1	3	MainData	
main_const		C01C	1	1	1	MainConst	
- LABELS:							
__SEG_END_SSTACK		902	0	0	1		
MODULE: -- rel_ex_la_sub_1.asm.o --							
- PROCEDURES:							
sub_1		C015	7	7	1	SubCode	
- VARIABLES:							
sub_const		C01D	1	1	1	SubConst	
sub_data		801	1	1	1	SubData	

One of the most important sections is the object-allocation section. It shows, for each of the modules (in this case the main and subroutine) where the procedures, labels, and data elements are located. We will need this information when it comes time to debug our program. The hSize and dSize give the number of bytes in hexadecimal and decimal.

Module Statistic:

MODULE STATISTIC			
Name	Data	Code	Const
rel_ex_la_main.asm.o	2	42	2
rel_ex_la_sub_1.asm.o	2	14	2
other	256	2	0

This shows for each module the number of code, data, and constants bytes that have been allocated. In this listing, *other* refers to the space allocated to the stack.

Section Use in Object-Allocation Section:

SECTION USE IN OBJECT-ALLOCATION SECTION	
SECTION: ".init"	
Entry	loop
SECTION: "SubCode"	
sub_1	
SECTION: "MainConst"	
main_const	
SECTION: "SubConst"	
sub_const	
SECTION: "MainData"	
main_data	
SECTION: "SubData"	
sub_data	

This section shows which elements (or objects) of our program use which memory section.

Object List Sorted by Address:

OBJECT LIST SORTED BY ADDRESS						
Name	Addr	hSize	dSize	Ref	Section	RLIB
main_data	800	1	1	3	MainData	
sub_data	801	1	1	1	SubData	
Entry	C000	9	9	0	.init	
loop	C009	C	12	1	.init	
sub_1	C015	7	7	1	SubCode	
main_const	C01C	1	1	1	MainConst	
sub_const	C01D	1	1	1	SubConst	

The address of each object (denoted by a label) is given here.

Unused Object Section:

This lists any objects found in the object files that were not linked, such as a label not referenced.

Copydown Section:

This section lists all blocks that are copied from ROM to RAM at program startup. This is more relevant in C programming.

Object-Dependencies Section:

OBJECT-DEPENDENCIES SECTION	
Entry	USES __SEG_END_SSTACK main_const main_data
loop	USES main_data sub_1 loop
sub_1	USES sub_data sub_const

This section lists the names of global objects which every function and variable.

Dependency Tree:

An object dependency tree shows in a tree format all detected dependencies between functions.

Statistic Section:

STATISTIC SECTION	
ExeFile:	

Number of blocks to be downloaded: 5	
Total size of all blocks to be downloaded: 32	

This section gives the number of bytes of code in the application.

Absolute Files

The linker can produce two kinds of output files to be loaded in the microcontroller memory. These are absolute, *.abs*, and Motorola S files.

Absolute files, which have the extension *.abs*, contain the program code plus some debugging information. We will see how to use this information in Chapter 8. Motorola S files, which have extensions *.s1*, *.s2*, *.s3* or *.s19*, contain all code from the READ_ONLY sections of the application. This, then, can be downloaded to the microcontroller or programmed into the ROM.

5.8 Remaining Questions

In this chapter we have described enough of the features of the linker to get started in assembly language programming. We still have a number of issues to cover and we will

do that in succeeding chapters. Some of these details yet to be explained include the following topics:

- What is startup code? *Startup code is the code that is executed when the processor is first powered-up or the program is first run. In C programs a standard startup code sequence is used. We will discuss this issue further in Chapter 8.*

5.9 Conclusion and Chapter Summary Points

In this chapter we have discussed the operation of the linker. We have included enough details to be able to assemble and link assembly language programs using separately assembled relocatable modules.

- There are two kinds of memory in a microcontroller system – RAM and ROM.
- There can be two types of ROM – EEPROM and FLASH EEPROM.
- You must know what the memory map is to be able to locate the various parts of the program in the proper memory types.
- Variable data storage and the stack use RAM.
- The program code and constants use ROM.
- The linker is controlled by a linker parameter file.
- It is very easy to change the parameter file to relocate the code for different hardware configurations.

5.10 Bibliography and Further Reading

Cady, F. M., *Microcontrollers and Microprocessors*, Oxford University Press, New York, 1997.

Smart Linker, Metrowerks Corporation, Austin, TX, 2003.

5.11 Problems

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