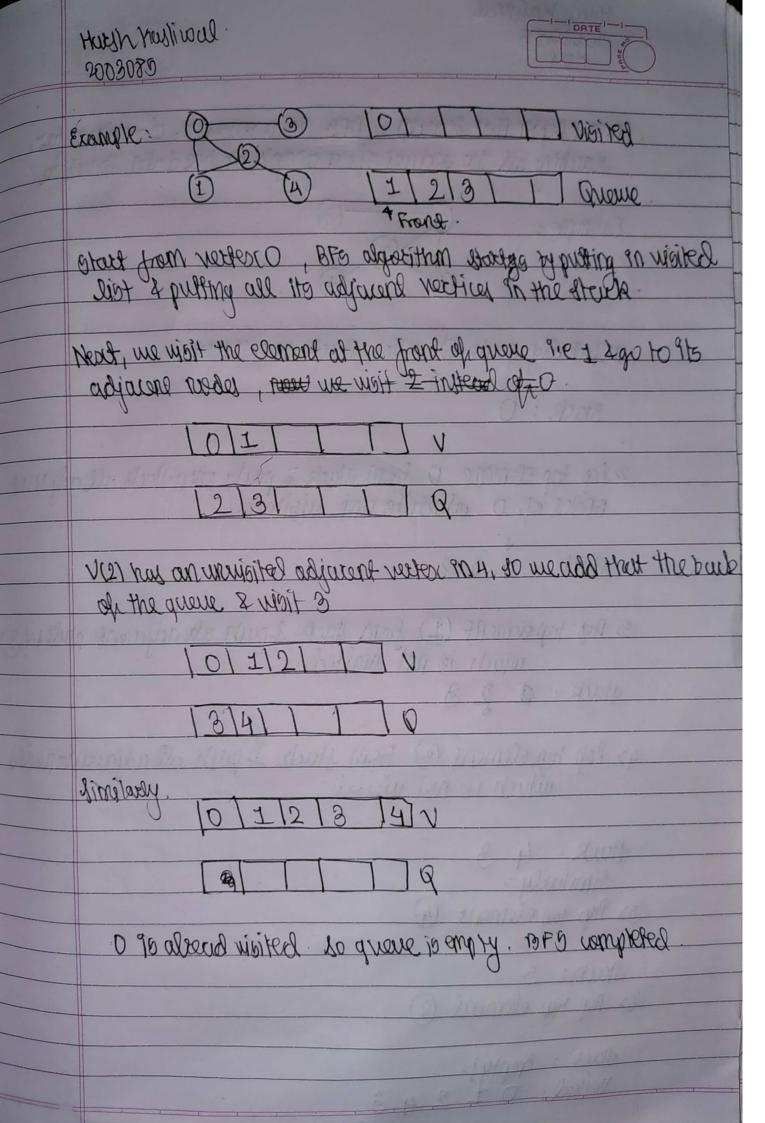
Hatch Kashwal 2003085
Experiment 9
Aim: Implement Graph truleval mathads
Theory:
A Graph to a non linear data structure consisting of nodes and edges. The nodes are sumetimes also reflected to as vertical and the edges are lines or ares that connect any two nodes in the graph.  (2)  (3)  (4)  (3)  (4)  (5)
Vertica: 50,1,2,3,43.  Edges: 501, 42,28,34,04,14,133.  Two needes are adjacent of they are agranuled by an edge.  Croaphs can be discred or analisabled.
In undirected graph, edges do not have any direction as what with them.  inc pair (u,v), & (v,u) represents the same edge.  eg.  2 9 8 12,23,34,413 ~ enty 4

	Harsh Naylineal DATE DATE
•	In diserted graph, edges forms an ordered pair 9.0 (4,4) and (4,4) and (4,4)
	3 C
	Terninologies:  1> Path: can be defined as sequence of westices (4,42-1) shows a way that the are (4,142). (42,49) edges in G(E).
	De pongen connected graph. If a path exist between each paix of
	2) Indepose: of a rode (u) 90 no of edges that teapinates af y out dapose: of a rode (u) 90 no of edges that egginates at u 5) begree: of a rode (u) 90 sum of endegree & andagree of that of langth: No of edges in the path.  7) yeld: A path that starts and ends of same nodes.
1300 J. 1.	BFG (Breaden Front Search): A stundar BFG implementation puls  each vertex of the graph into one of two categories.  - Vioited  - Not visited.
	The purpose of the algorithm is to mark each nextex as visited white



	Hayh Kayliwal 2008085 C21
	DFG ( Depth Fix) + Security : DFG 96 a securitive algorithm for searching all the various of a graph or tree data structure.
100000	Example 3  (1) (2) (3)
et of	15 Totally push O onto struk as follows
	2> Pep top element o from stack & push onto stack all adjusted.  noclas of o which is not wisiked.
dust which	Stuck . 1 2 3
	3> Pap toperant (I) from start & pun all adjacent pedesoft which is not vigited.
	4> Pop top element @ from strub 2 push all adjacent pedes which is not waited
	Struk: 4 3. Similarly:  5> Pep top element (4)
	5) hop top element (3)
	Hack: Empty! Visited: 0 1 2 4 3

Horsh Hashival 2003786 conclusion: BFG is more suitable for sentiling wedges which are closer to the given soutre while DFG is more suitable when there are solution, away from soutre.

## **PROGRAM:**

Write a program to implement infix to postfix conversion using stack.

## Code:

```
#include <iostream>
#include <conio.h>
using namespace std;
int q[20],top=-1,front=-1,rear=-1,a[20][20],vis[20],stack[20];
int deleteQ();
void add(int item);
void bfs(int s,int n);
void dfs(int s,int n);
void push(int item);
int pop();
int main()
int n,i,s,ch,j;
char c,;
cout<<"ENTER THE NUMBER VERTICES ";</pre>
cin>>n;
for(i=1;i<=n;i++)
for(j=1;j<=n;j++)
cout<<"ENTER 1 IF "<<i<<" HAS A NODE WITH "<<j<<" ELSE 0 ";</pre>
cin>>a[i][j];
cout<<"THE ADJACENCY MATRIX IS\n";</pre>
for(i=1;i<=n;i++)</pre>
for(j=1;j<=n;j++)
cout<<"\t"<<a[i][j];
cout<<"\n";</pre>
do
for(i=1;i<=n;i++)</pre>
vis[i]=0;
cout<<"\nENTER YOUR CHOICE";</pre>
cout<<"\n1.BFS Traversal";</pre>
cout<<"\n2.DFS Traversal\n";</pre>
cin>>ch;
```

```
cout<<"ENTER THE SOURCE VERTEX :";</pre>
cin>>s;
switch(ch)
case 1:bfs(s,n);
break;
case 2:
dfs(s,n);
break;
cout<<"\nDO U WANT TO CONTINUE(Y/N) ? ";</pre>
cin>>c;
}while((c=='y')||(c=='Y'));
void bfs(int s,int n)
int p,i;
add(s);
vis[s]=1;
p=deleteQ();
if(p!=0)
cout<<p<<"\t";</pre>
while(p!=0)
for(i=1;i<=n;i++)
if((a[p][i]!=0)&&(vis[i]==0))
add(i);
vis[i]=1;
p=deleteQ();
if(p!=0)
cout<<p<<"\t";</pre>
for(i=1;i<=n;i++)
if(vis[i]==0)
bfs(i,n);
void add(int item)
if(rear==19)
cout<<"QUEUE FULL";</pre>
else
```

```
if(rear==-1)
q[++rear]=item;
front++;
else
q[++rear]=item;
int deleteQ()
int k;
if((front>rear)||(front==-1))
return(0);
else
k=q[front++];
return(k);
void dfs(int s,int n)
int i,k;
push(s);
vis[s]=1;
k=pop();
if(k!=0)
cout<<k<<"\t";
while(k!=0)
for(i=1;i<=n;i++)
if((a[k][i]!=0)&&(vis[i]==0))
push(i);
vis[i]=1;
k=pop();
if(k!=0)
cout<<k<<"\t";
for(i=1;i<=n;i++)
if(vis[i]==0)
dfs(i,n);
void push(int item)
```

```
{
if(top==19)
cout<<"Stack overflow ";
else
stack[++top]=item;
}
int pop()
{
int k;
if(top==-1)
return(0);
else
{
k=stack[top--];
return(k);
}
return 0;
}</pre>
```

## **OUTPUT:**

```
PS D:\Harsh\SEM 3\DS\CODES> cd "d:\Harsh\SEM 3\DS\CODES\" ; if ($?)
ENTER THE NUMBER VERTICES 3
ENTER 1 IF 1 HAS A NODE WITH 1 ELSE 0 1
ENTER 1 IF 1 HAS A NODE WITH 2 ELSE 0 1
ENTER 1 IF 1 HAS A NODE WITH 3 ELSE 0 0
ENTER 1 IF 2 HAS A NODE WITH 1 ELSE 0 1
ENTER 1 IF 2 HAS A NODE WITH 2 ELSE 0 0
ENTER 1 IF 2 HAS A NODE WITH 3 ELSE 0 1
ENTER 1 IF 3 HAS A NODE WITH 1 ELSE 0 0
ENTER 1 IF 3 HAS A NODE WITH 2 ELSE 0 1
ENTER 1 IF 3 HAS A NODE WITH 3 ELSE 0 1
THE ADJACENCY MATRIX IS
       1
               1
                       0
       1
              0
                       1
       0
             1
                      1
ENTER YOUR CHOICE
1.BFS Traversal
2.DFS Traversal
1
ENTER THE SOURCE VERTEX: 2
      1
               3
DO U WANT TO CONTINUE(Y/N) ? y
ENTER YOUR CHOICE
1.BFS Traversal
2.DFS Traversal
ENTER THE SOURCE VERTEX: 2
      3
              1
DO U WANT TO CONTINUE(Y/N) ? n
PS D:\Harsh\SEM 3\DS\CODES>
```