BIRLA INSTITUTE OF TECHNOLOGY DEOGHAR CAMPUS

| | | | QUIZ: F | IRST | | |
|--|---|---------------------------------------|---------------------------------|-------------------------------|-----------------|---|
| Name : Roll. No: | | | | | | |
| | | | | | | e: 50 minutes |
| Sem.: | | Branch: CSE | | Session: SP/24 Full Ma | | Marks: 10 |
| Instru | ctions: | | | | | |
| 1. | Attempts all | the Questions. | | | | |
| 2. Overwriting means Cancellation of question. | | | | | | Set: C |
| Q 1. W | hich of the fo | ollowing is NO | Γ a characteristic | | | [1mark] |
| | | as its own schedu n move between d | lling policy ifferent queues | b) Queues ard d) System ad | | assigned to processes lefine queue priorities |
| Q. 2. T | | - | | | | |
| • | | | | | | [1mark] |
| Q.3. TI | _ | e, and other inform | nation to control in | | | · · · · · · · · · · · · · · · · · · · |
| Q. 4. | Thread shares | with other thread | s belonging to the s | ame process it | S | [1 Mark] |
| | (a) code section and data section (1 | | | | program Co | . , |
| | (c) register set and stack | | | (d) thread id | | |
| | () | J | | () | | |
| Q. 5. | Suppose that the listed amo | | esses arrive for exe | cution at the ti | mes indicated | . Each process will run |
| | Proce | | rival Time | CPU Burs | st Time | Priority |
| | P1 | | 0 | 11 | 90 11110 | 2 |
| | P2 | | 5 | 28 | | 0 |
| | P3 | | 12 | 2 | | 3 |
| | P4 | | 2 | 10 | | 1 |
| | P5 | | 9 | 16 | | 4 |
| | | ollowing question | | | | |
| | | | late the average Tu | ırn-around tim | e for the follo | wing |
| | Scheduling algorithms. | | | | | [2] |
| | (a) Priority (Preemptive) Scheduling algorithm(b) Priority (Non Preemptive) Scheduling algorithm | | | | | [2marks] |
| | (b) | FIIOTHY (NOII PI | eempuve) Schedu | ing aigorium | | [2marks] |
| Q. 6. | Fill in the hl | anks in the foll | owing statement | | | |
| ٧. u. | | | _ | | lulino mav | not be able to avoid |
| | | | | | | |
| | deadline over | erruns if the to | otal fraction of C | PU time us | ea by the p | rocesses according to |

Eq.(7.4).... $\sum_{i=1\cdots n} \frac{x_i}{T_i} \le 1$

utilization if it is to meet deadlines of processes."

Eq. (7.4) exceeds, where *m* is the number of processes. This expression has a lower bound of, which implies that if an application has a large number of processes, RM scheduling may not be able to achieve more thanpercent CPU

[2marks]