SOLUTION:

Code Sequence:

C=A+B

D=A-E

F=C+D

I. Accumulator Architecture

Load A
Add B
Store C
Load A
Sub E
Store D
Load C
Add D
Store F

Code	Destroyed	Overhead	Bits(Code Size)	Size of Memory
	Variable			Data
Load A			40	32
Add B	Α		40	32
Store C			40	32
Load A	С	1	40	32
Sub E	Α		40	32
Store D			40	32
Load C			40	32
Add D		1	40	32
Store F			40	32
Total		2x4=8 bytes	360 bits	288 bits

II. Memory-Register Architecture

Load R1,A
Add R2,R1,B
Store R2,C
Sub R3,R1,E
Store R3,D
Add R4,R2,D
Store R4,F

Code	Destroyed	Overhead	Bits (Code	Size of Memory
	Variable		Size)	Data
Load R1,A			48	32
Add R2,R1,B			56	32
Store R2,C			48	32
Sub R3,R1,E			56	32
Store R3,D			48	32
Add R4,R2,D			56	32
Store R4,F			48	32
Total			360 bits	224 bits

III. Register Architecture

Load R1,A

Load R2,B

Add R3,R1,R2

Store R3,C

Load R4,E

Sub R5,R1,R4

Store R5,D

Add R6,R3,R5

Store R6,F

Code	Destroyed	Overhead	Bits (Code	Size of Memory
	Variable		Size)	Data
Load R1,A			48	32
Load R2,B			48	32
Add R3,R1,R2			32	
Store R3,C			48	32
Load R4,E			48	32
Sub R5,R1,R4			32	
Store R5,D			48	32
Add R6,R3,R5			32	
Store R6,F			45	
Total			384 bits	192 bits

SOLUTION:

1.)

Size of foo struct = 1+1+4+8+2+4+8+8+4= 48 bytes

2.)

I
Data Size on 64 bit
machine(In bytes)
1
1
2
4
8
2
2
4
8
8
8
4
0
52

Required memory size is 52 bytes.

3.)

Re arranging size of struct as Highest to lowest we have:

Data Type	Data size on 64-bit machine (bytes)
double	8
double	8
pointer	8
pointer	8
int	4
int	4
float	4
short	2
char	1
bool	1
trailing pad	0
Total	48

So, We would need minimum 48 bytes.

3.

SOLUTION:

Remaining Instructions make up another 5%

So,

$$CPI = (0.15 + 0.25) \times 4 + 0.15 \times 3 + (0.35 + 0.05) \times 1 + 0.05 \times 12 = 1.6 + 0.45 + 0.4 + 0.6$$

$$= 3.05$$

Since Cycles can not be in fraction. So, 3 cycles per instruction.

SOLUTION:

```
Since A B C D are double precision floating points we must use Risc V D
extension.
Firstly load A and B in memory
Fld f1, A;
                       // load A in f1
Fld f2, B;
                       // load B in f2
Fsub f3,f1,f2;
                        // Compute value C and store in f3
//Now for D = 2-A + B
Fadd f7,f1,f2;
                       // storing a+b in reg f7
Fsub f4, 2,f7
                        // storing 2-A+B in D as f4
//Now we first need to load integer I and J from memory:
                       // load I in X5
Lw x5, 150[I]
Lw x6, 200[J]
                       // load J in X6
// now after having I and j stored in x5 and x6 respectively we can start if else computation.
                       // If(i==j)
Bne x5,x6,Else;
Fadd f1,f1,f2;
                       // A=A+B
Beq x0,x0,Exit;
                       // unconditional exit
Else: Fsub f2,f2,f1;
                       // B=B-A
Exit:
                       // Assembler calculates address.
```

SOLUTION:

```
1.)
```

```
// load d in x10
Ld x10,d;
Addi x11,x0,0; // i=0 in register x11
Ld x9,0(x0);
                       // X[0] is in & O(x0) is in reg x9
Lstart:
                      // load X[i] as x12
Ld x12,0(x9)
Add x12,x12,x10;
                       // X[i]+=d;
Addi x9,x9,8;
                       // & x[i++]
Addi x11,x11,1; // i++
Addi x13, x0,100;
                       // x13= 100
Blt x11,x13,Lstart;
                       //loop between current I and I=100
```

2.)

We assumed opcode is 8 bit, memory is 64 bit, and register addresses are 6 bits So,

- 1) 8+6+6+64= 84 bits
- 2) 8+6+6+6=26 bits
- 3) 84 bits
- 4) 84 bits
- 5) 8+6+6+6= 26 bits
- 6) 26 bits
- 7) 26 bits
- 8) 26 bits
- 9) 26 bits

Total code size in bits = 408 bits

3.)

Total number of load instructions : = 2+100=102

Branches = 100 instructions

ALU: 400 +1 = 401instructions

Now the total number of instructions are 603 instructions

CPI =
$$((102/603) \times 4) + ((100/603) \times 3) + ((401/603) \times 1)$$

= $0.67 + 0.49 + 0.66$
= 1.82

Approximately 2 cycles per instruction