1.2.3.4

The IP building blocks

Understanding the IP Protocol

IP Address

- Layer 3 property
- Can be set automatically or statically
- Network and Host portion
- 4 bytes in IPv4 - 32 bits

Network vs Host

- a.b.c.d/x (a.b.c.d are integers) x is the network bits and remains are host
- Example 192.168.254.0/24
- The first 24 bits (3 bytes) are network the rest 8 are for host
- This means we can have 2^24 (16777216) networks and each network has 2^8 (255) hosts
- Also called a subnet

Subnet Mask

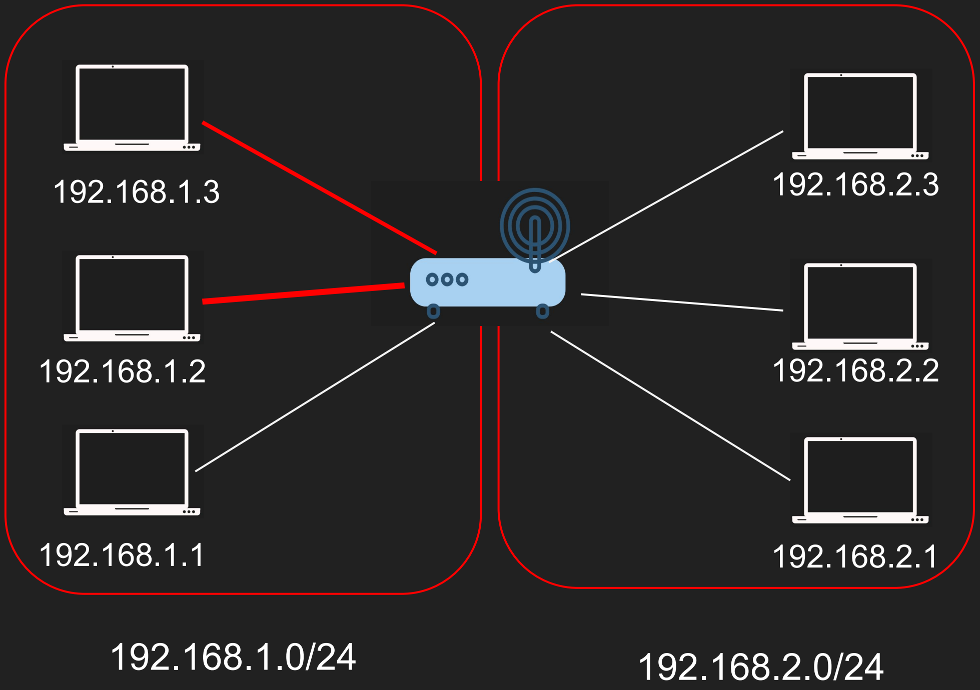
- 192.168.254.0/24 is also called a subnet
- The subnet has a mask 255.255.255.0
- Subnet mask is used to determine whether an IP is in the same subnet

Default Gateway

- Most networks consists of hosts and a Default Gateway
- Host A can talk to B directly if both are in the same subnet
- Otherwise A sends it to someone who might know, the gateway
- The Gateway has an IP Address and each host should know its gateway

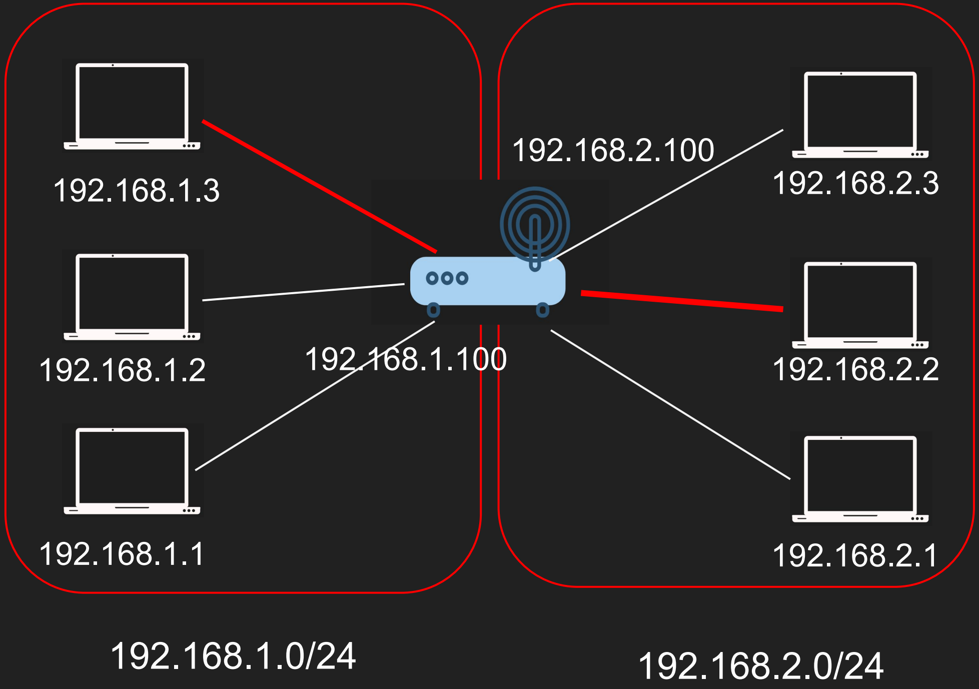
E.g. Host 192.168.1.3 wants to talk to 192.168.1.2

- 192.168.1.3 applies subnet mask to itself and the destination IP 192.168.1.2
- 255.255.255.0 & 192.168.1.3 = 192.168.1.0
- 255.255.255.0 & 192.168.1.2 = 192.168.1.0
- Same subnet ! no need to route



E.g. Host 192.168.1.3 wants to talk to 192.168.2.2

- 192.168.1.3 applies subnet mask to itself and the destination IP 192.168.2.2
- 255.255.255.0 & 192.168.1.3 = 192.168.1.0
- 255.255.255.0 & 192.168.2.2 = 192.168.2.0
- Not the subnet ! The packet is sent to the Default Gateway 192.168.1.100



Summary

- IP Address
- Network vs Host
- Subnet and subnet mask
- Default Gateway

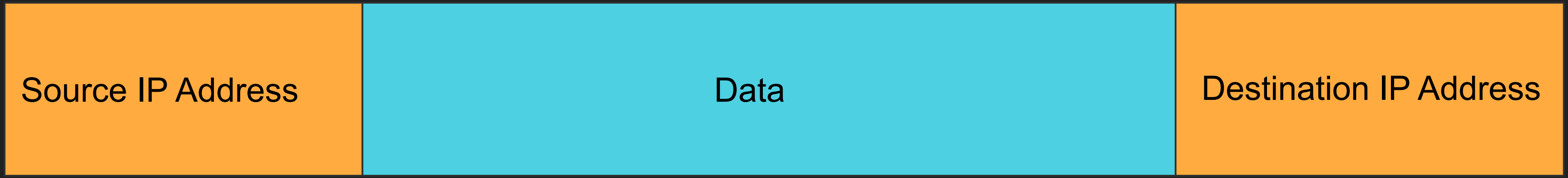
The IP Packet

Anatomy of the IP Packet

IP Packet

- The IP Packet has headers and data sections
- IP Packet header is 20 bytes (can go up to 60 bytes if options are enabled)
- Data section can go up to 65536

IP Packet to the Backend Engineer



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Actual IP Packet

Offsets	Octet	0								1								2								3							
Octet	**Bit**	**0**	**1**	**2**	**3**	**4**	**5**	**6**	**7**	**8**	**9**	**10**	**11**	**12**	**13**	**14**	**15**	**16**	**17**	**18**	**19**	**20**	**21**	**22**	**23**	**24**	**25**	**26**	**27**	**28**	**29**	**30**	**31**
0	**0**	Version				IHL				DSCP						ECN		Total Length															
4	**32**	Identification																Flags			Fragment Offset												
8	**64**	Time To Live								Protocol								Header Checksum															
12	**96**	Source IP Address																															
16	**128**	Destination IP Address																															
20	**160**	Options (if IHL > 5)																															
⋮	⋮																																
56	**448**																																
		Data																															

https://datatracker.ietf.org/doc/html/rfc791
https://en.wikipedia.org/wiki/IPv4

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Version - The Protocol version

Offsets	Octet	0								1								2								3							
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
0	0	Version				IHL				DSCP						ECN		Total Length															
4	32	Identification																Flags			Fragment Offset												
8	64	Time To Live								Protocol								Header Checksum															
12	96	Source IP Address																															
16	128	Destination IP Address																															
20	160	Options (if IHL > 5)																															
⋮	⋮																																
56	448																																
		Data																															

Internet Header Length - Defines the Options length

Offsets	Octet	0								1								2								3							
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
0	0	Version				IHL				DSCP						ECN		Total Length															
4	32	Identification																Flags			Fragment Offset												
8	64	Time To Live								Protocol								Header Checksum															
12	96	Source IP Address																															
16	128	Destination IP Address																															
20	160	Options (if IHL > 5)																															
⋮	⋮																																
56	448																																

Data

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Total Length - 16 bit Data + header

Offsets	Octet	0								1								2								3							
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
0	0	Version				IHL				DSCP						ECN		Total Length															
4	32	Identification																Flags			Fragment Offset												
8	64	Time To Live								Protocol								Header Checksum															
12	96	Source IP Address																															
16	128	Destination IP Address																															
20	160	Options (if IHL > 5)																															
⋮	⋮																																
56	448																																

Data

Fragmentation - Jumbo packets

Offsets	Octet	0								1								2								3							
Octet	**Bit**	**0**	**1**	**2**	**3**	**4**	**5**	**6**	**7**	**8**	**9**	**10**	**11**	**12**	**13**	**14**	**15**	**16**	**17**	**18**	**19**	**20**	**21**	**22**	**23**	**24**	**25**	**26**	**27**	**28**	**29**	**30**	**31**
0	**0**	Version				IHL				DSCP						ECN		Total Length															
4	**32**	Identification																Flags			Fragment Offset												
8	**64**	Time To Live								Protocol								Header Checksum															
12	**96**	Source IP Address																															
16	**128**	Destination IP Address																															
20	**160**	Options (if IHL > 5)																															
⋮	⋮																																
56	**448**																																
		Data																															

Time To Live - How many hops can this packet survive?

Offsets	Octet	0								1								2								3							
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
0	0	Version				IHL				DSCP						ECN		Total Length															
4	32	Identification																Flags			Fragment Offset												
8	64	Time To Live								Protocol								Header Checksum															
12	96	Source IP Address																															
16	128	Destination IP Address																															
20	160	Options (if IHL > 5)																															
⋮	⋮																																
56	448																																

Data

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Protocol - What protocol is inside the data section?

Offsets	Octet	0								1								2								3							
Octet	**Bit**	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
0	**0**	Version				IHL				DSCP						ECN		Total Length															
4	**32**	Identification																Flags			Fragment Offset												
8	**64**	Time To Live								Protocol								Header Checksum															
12	**96**	Source IP Address																															
16	**128**	Destination IP Address																															
20	**160**	Options (if IHL > 5)																															
⋮	⋮																																
56	**448**																																
		Data																															

Source and Destination IP

Offsets	Octet	0								1								2								3							
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
0	0	Version				IHL				DSCP						ECN		Total Length															
4	32	Identification																Flags			Fragment Offset												
8	64	Time To Live								Protocol								Header Checksum															
12	96	Source IP Address																															
16	128	Destination IP Address																															
20	160	Options (if IHL > 5)																															
⋮	⋮																																
56	448																																
		Data																															

Explicit Congestion Notification

Offsets	Octet	0								1								2								3							
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
0	0	Version				IHL				DSCP						ECN		Total Length															
4	32	Identification																Flags			Fragment Offset												
8	64	Time To Live								Protocol								Header Checksum															
12	96	Source IP Address																															
16	128	Destination IP Address																															
20	160	Options (if IHL > 5)																															
⋮	⋮																																
56	448																																

Data

Summary

- The IP Packet has headers and data sections
- IP Packet header is 20 bytes (can go up to 60 bytes if options are enabled)
- Data section can go up to 65536
- Packets need to get fragmented if it doesn’t fit in a frame

ICMP

Internet Control Message Protocol

ICMP

- Stands for Internet Control Message Protocol
- Designed for informational messages
  - Host unreachable, port unreachable, fragmentation needed
  - Packet expired (infinite loop in routers)
- Uses IP directly
- PING and traceroute use it
- Doesn’t require listeners or ports to be opened

ICMP header

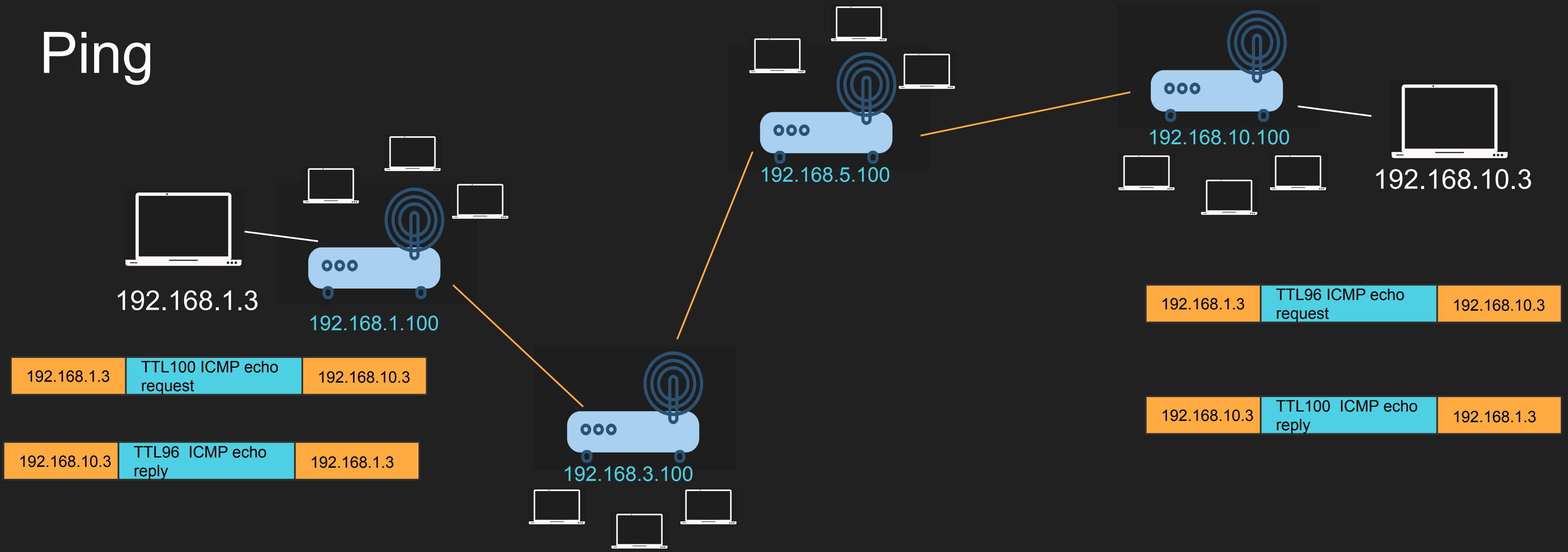
Offsets	Octet	0								1								2								3							
Octet	**Bit**	**0**	**1**	**2**	**3**	**4**	**5**	**6**	**7**	**8**	**9**	**10**	**11**	**12**	**13**	**14**	**15**	**16**	**17**	**18**	**19**	**20**	**21**	**22**	**23**	**24**	**25**	**26**	**27**	**28**	**29**	**30**	**31**
0	**0**	Type								Code								Checksum															
4	**32**	Rest of header																															

https://en.wikipedia.org/wiki/Internet_Control_Message_Protocol
https://datatracker.ietf.org/doc/html/rfc792

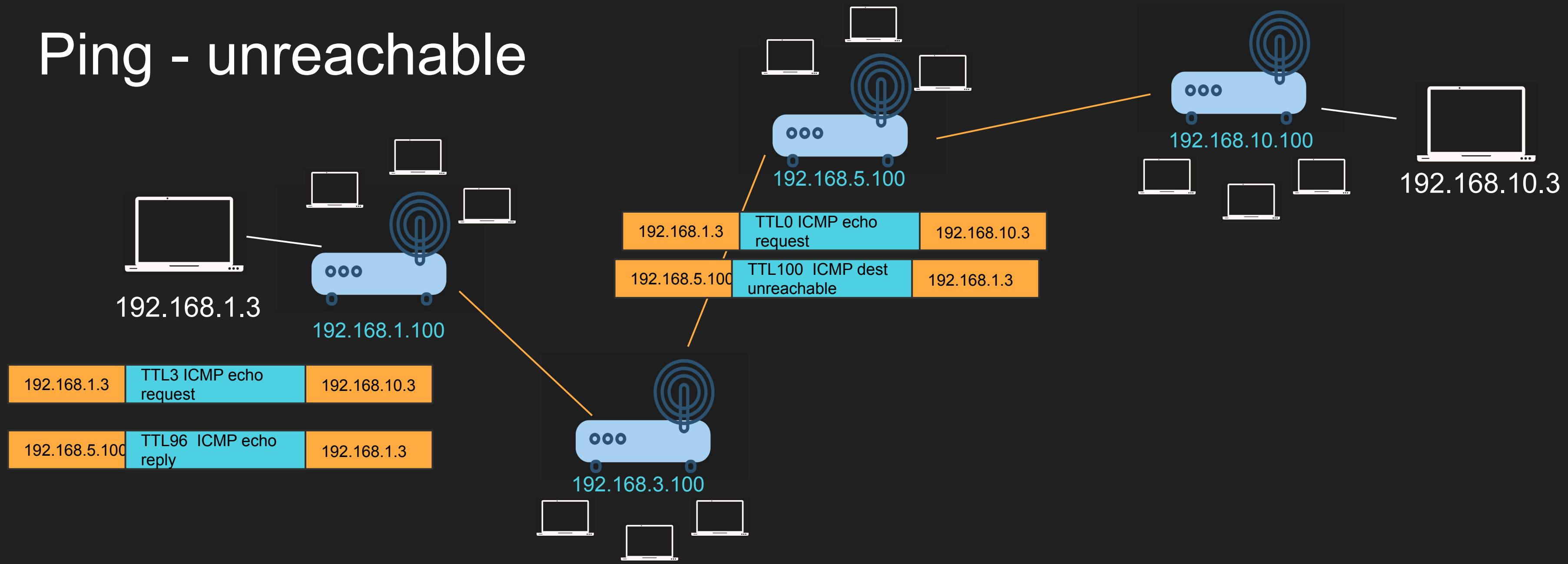
ICMP

- Some firewalls block ICMP for security reasons
- That is why PING might not work in those cases
- Disabling ICMP also can cause real damage with connection establishment
  - Fragmentation needed
- PING demo

Ping


192.168.1.3
192.168.1.100
192.168.3.100
192.168.5.100
192.168.10.100
192.168.10.3
192.168.1.3
TTL100 ICMP echo request
192.168.10.3
192.168.10.3
TTL96 ICMP echo reply
192.168.1.3
192.168.1.3
TTL96 ICMP echo request
192.168.10.3
192.168.10.3
TTL100 ICMP echo reply
192.168.1.3

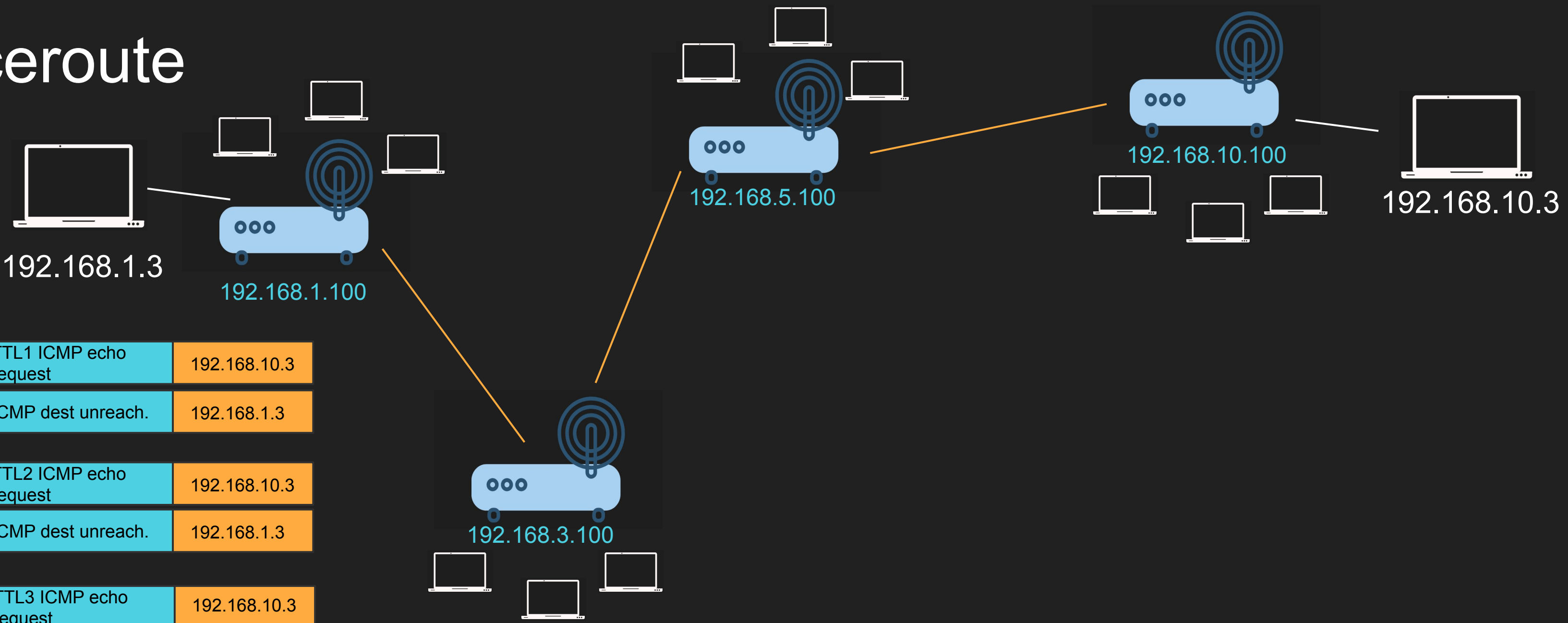
Ping - unreachable


192.168.1.3
192.168.1.100
192.168.3.100
192.168.5.100
192.168.10.100
192.168.10.3
192.168.1.3 | TTL3 ICMP echo request | 192.168.10.3
192.168.5.100 | TTL96 ICMP echo reply | 192.168.1.3
192.168.1.3 | TTL0 ICMP echo request | 192.168.10.3
192.168.5.100 | TTL100 ICMP dest unreachable | 192.168.1.3

TraceRoute

- Can you identify the entire path your IP Packet takes?
- Clever use of TTL
- Increment TTL slowly and you will get the router IP address for each hop
- Doesn’t always work as path changes and ICMP might be blocked

Traceroute



192.168.1.3	TTL1 ICMP echo request	192.168.10.3
192.168.1.100	ICMP dest unreach.	192.168.1.3
192.168.1.3	TTL2 ICMP echo request	192.168.10.3
192.168.3.100	ICMP dest unreach.	192.168.1.3
192.168.1.3	TTL3 ICMP echo request	192.168.10.3
192.168.5.100	ICMP dest unreach.	192.168.1.3
192.168.1.3	TTL4 ICMP echo request	192.168.10.3
192.168.10.100	ICMP dest unreach	192.168.1.3
192.168.1.3	TTL5 ICMP echo request	192.168.10.3
192.168.10.3	ICMP Echo reply	192.168.1.3

Summary

- ICMP is an IP level protocol used for information messages
- Critical to know if the host is available or port is opened
- Used for PING and TraceRoute
- Can be blocked which can cause problems

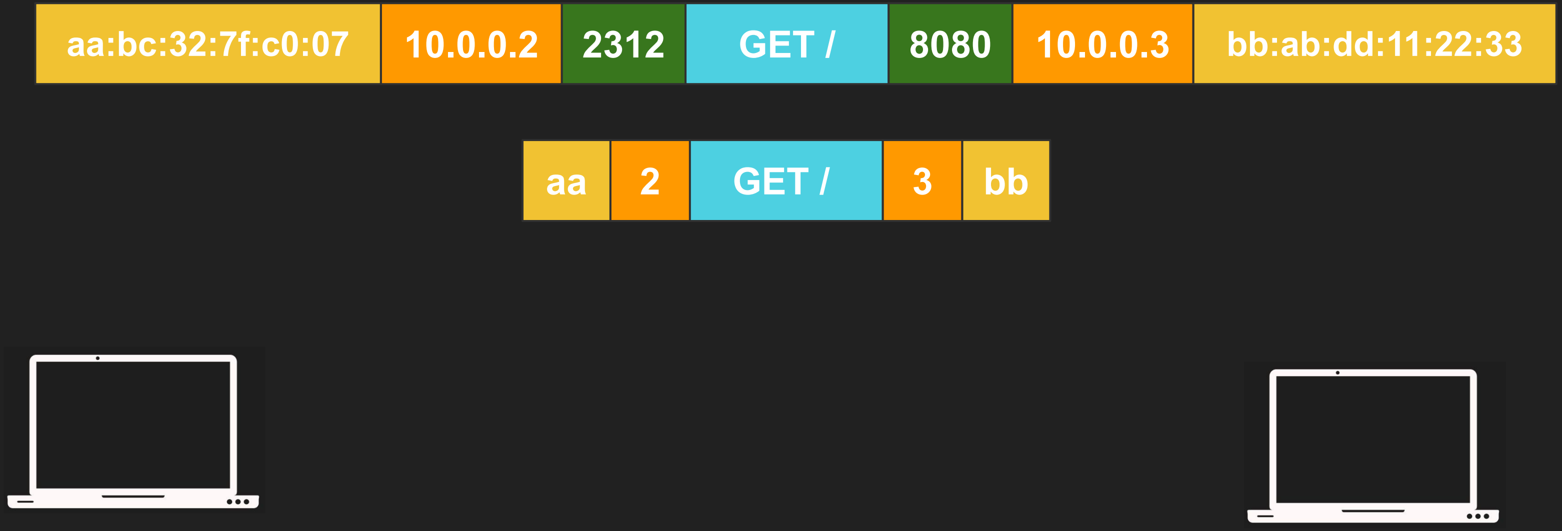
ARP

Address Resolution Protocol

Why ARP?

- We need the MAC address to send frames (layer 2)
- Most of the time we know the IP address but not the MAC
- ARP Table is cached IP->Mac mapping

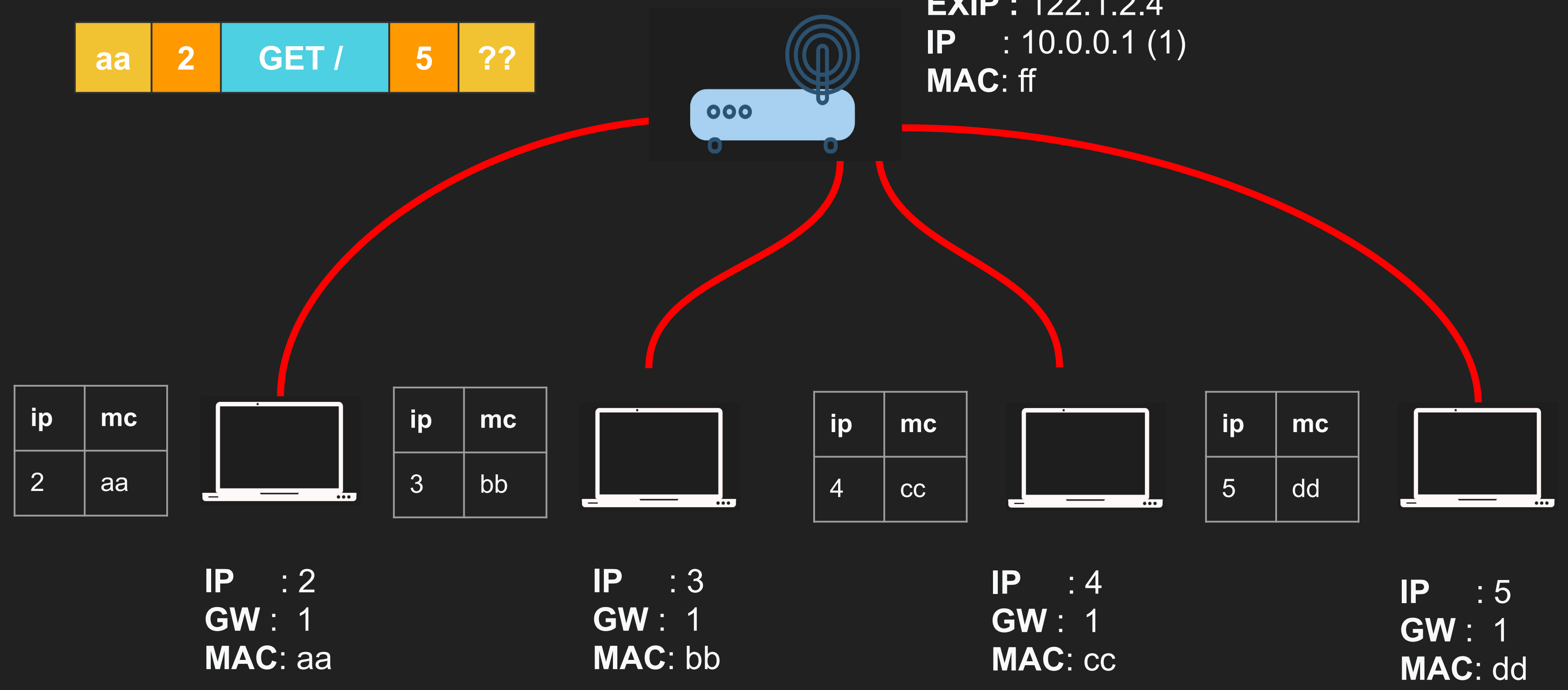
Network Frame



IP : 10.0.0.2
MAC: aa:bc:32:7f:c0:07

IP : 10.0.0.3
MAC: bb:ab:dd:11:22:33
Port: 8080

- IP 10.0.0.2 (**2**) wants to connect to IP 10.0.0.5 (**5**)
- Host 2 checks if host 5 is within its subnet, it is.
- Host **2** needs the MAC address of host **5**
- Host 2 checks its ARP tables and its not there



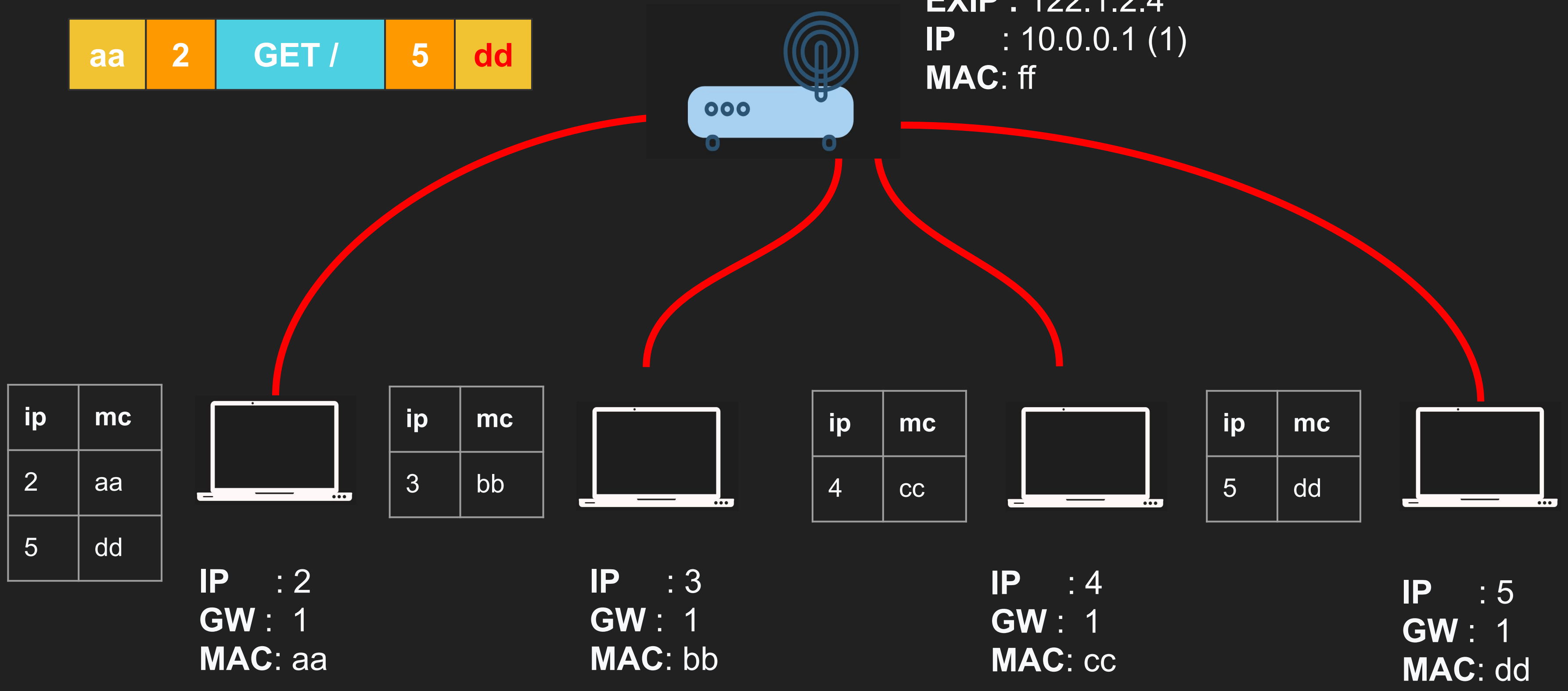
ip	mc
2	aa

ip	mc
3	bb

ip	mc
4	cc

ip	mc
5	dd

- Host 2 sends an ARP request broadcast to all machines in its network
- Who has IP address 10.0.0.5?
- Host 5 replies with dd
- Host 2 updates its ARP Table



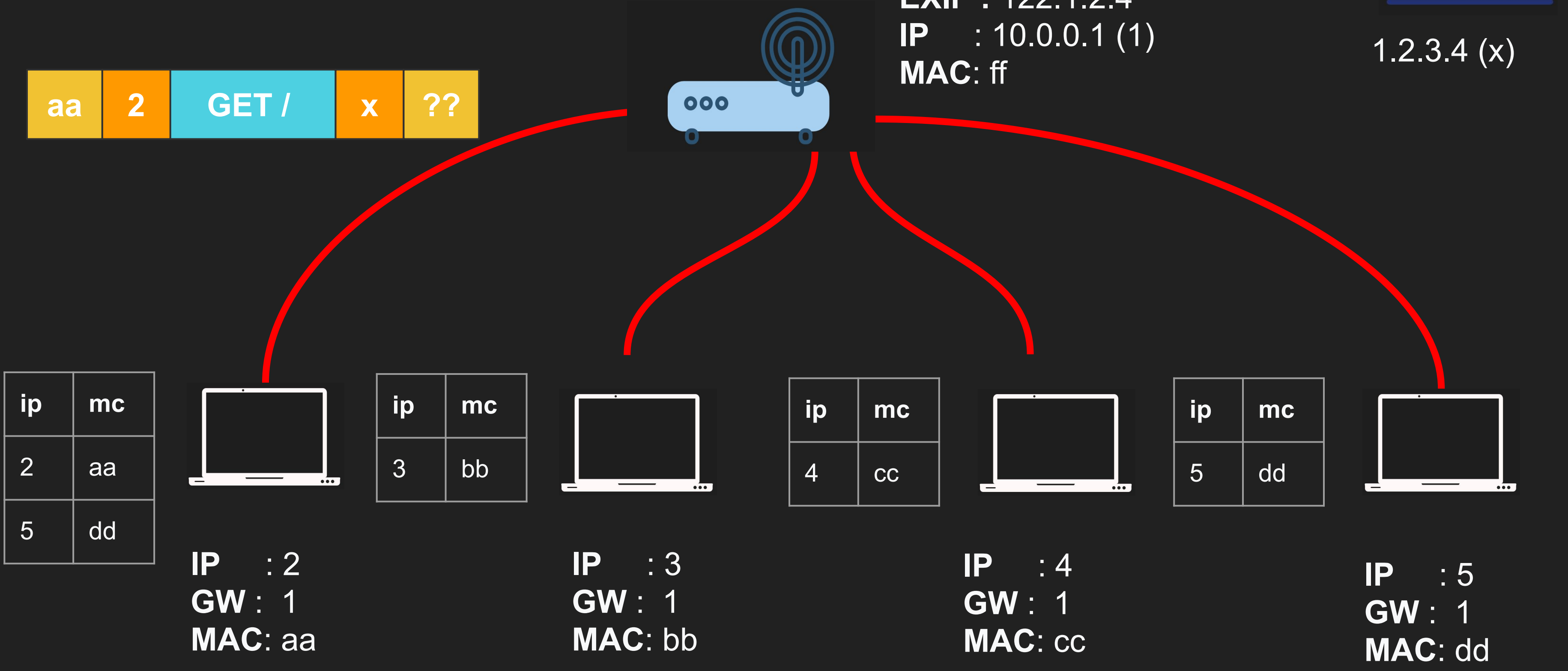
ip	mc
2	aa
5	dd

ip	mc
3	bb

ip	mc
4	cc

ip	mc
5	dd

- IP 10.0.0.2 (**2**) wants to connect to IP 1.2.3.4 (x)
- Host 2 checks if 1.2.3.4 is within its subnet, it is NOT!
- Host 2 needs to talk to its gatway
- Host **2** needs the MAC address of the gateway



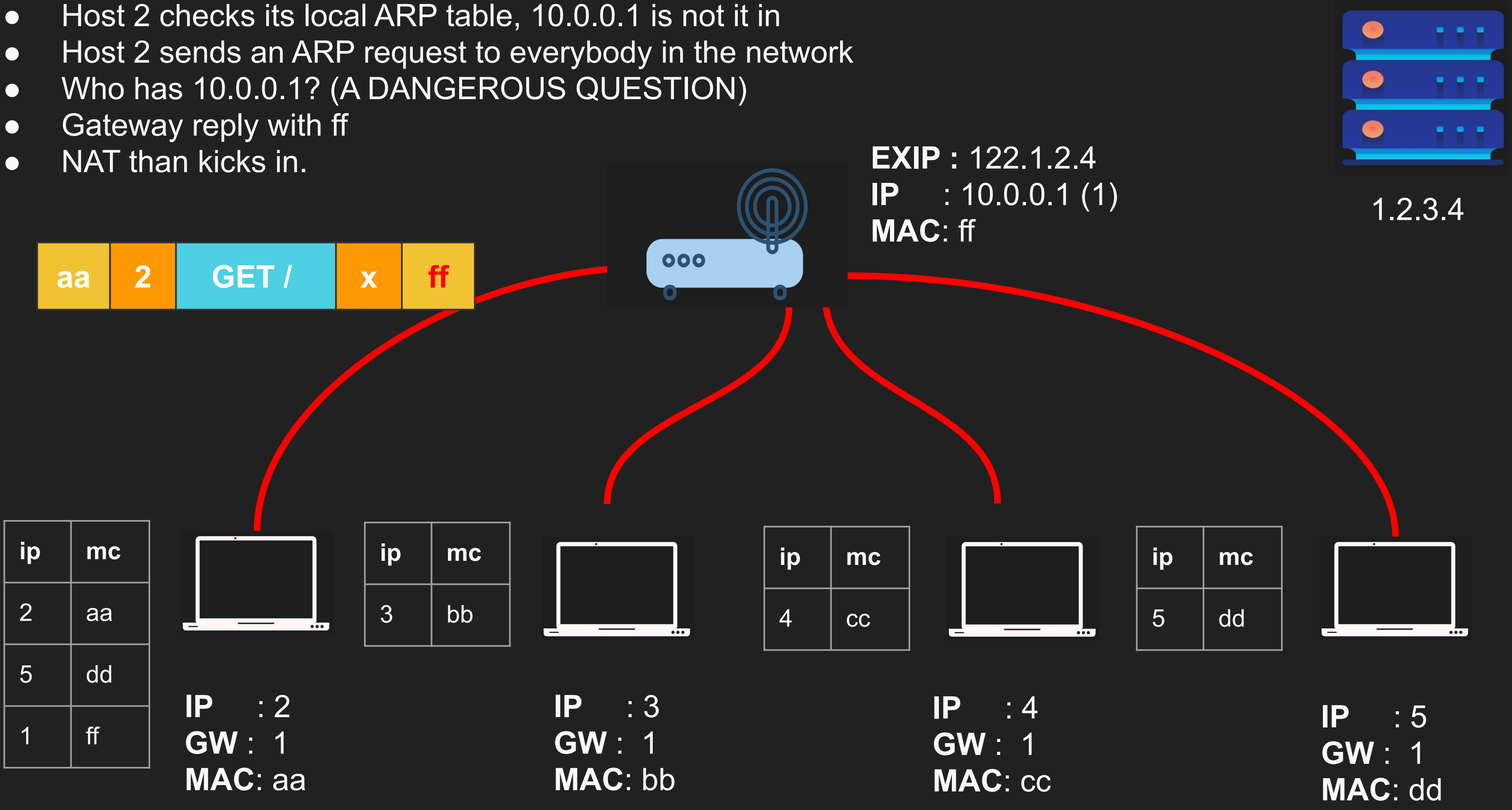
ip	mc
2	aa
5	dd

ip	mc
3	bb

ip	mc
4	cc

ip	mc
5	dd

- Host 2 checks its local ARP table, 10.0.0.1 is not it in
- Host 2 sends an ARP request to everybody in the network
- Who has 10.0.0.1? (A DANGEROUS QUESTION)
- Gateway reply with ff
- NAT than kicks in.



ip	mc
2	aa
5	dd
1	ff

ip	mc
3	bb

ip	mc
4	cc

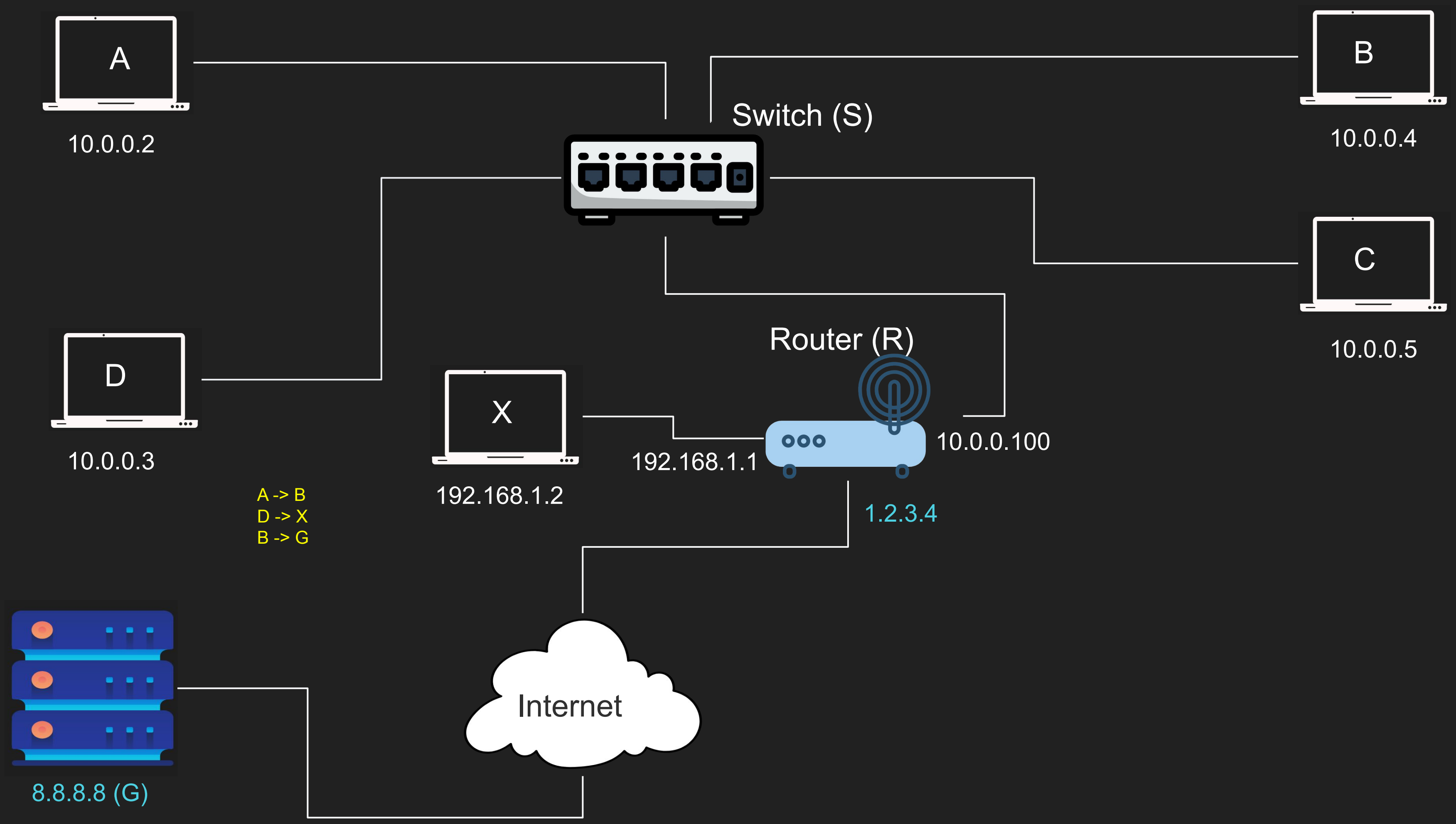
ip	mc
5	dd

Summary

- ARP stands for Address resolution protocol
- We need MAC address to send frames between machines
- Almost always we have the IP address but not the MAC
- Need a lookup protocol that give us the MAC from IP address
- Attacks can be performed on ARP (ARP poisoning)

Routing Example

How IP Packets are routed in Switches and Routers


A
10.0.0.2
B
10.0.0.4
Switch (S)
C
10.0.0.5
D
10.0.0.3
X
192.168.1.2
Router (R)
192.168.1.1
10.0.0.100
1.2.3.4
Internet
8.8.8.8 (G)

A -> B
D -> X
B -> G

UDP

User Datagram Protocol

UDP

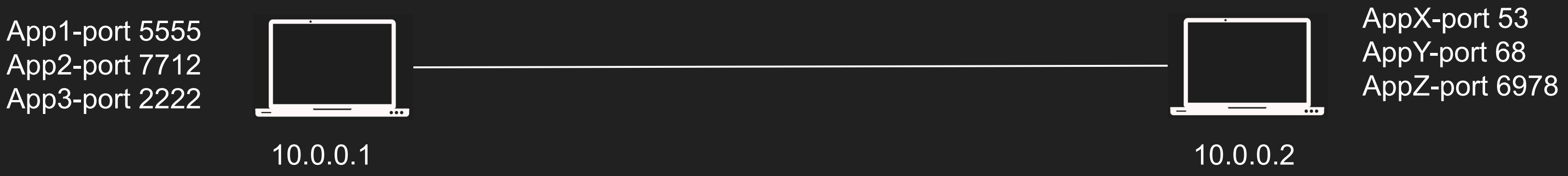
- Stands for User Datagram Protocol
- Layer 4 protocol
- Ability to address processes in a host using ports
- Simple protocol to send and receive data
- Prior communication not required (double edge sword)
- Stateless no knowledge is stored on the host
- 8 byte header Datagram

UDP Use cases

- Video streaming
- VPN
- DNS
- WebRTC

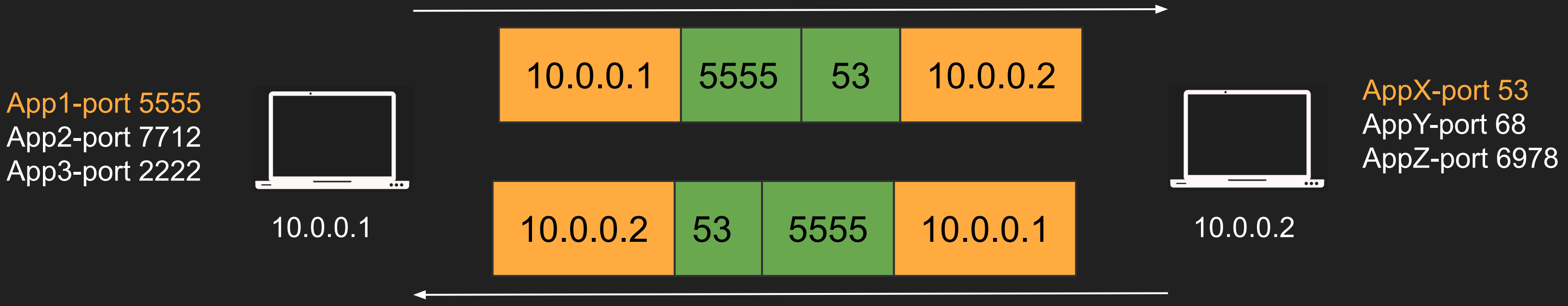
Multiplexing and demultiplexing

- IP target hosts only
- Hosts run many apps each with different requirements
- Ports now identify the “app” or “process”
- Sender multiplexes all its apps into UDP
- Receiver demultiplex UDP datagrams to each app



Source and Destination Port

- App1 on 10.0.0.1 sends data to AppX on 10.0.0.2
- Destination Port = 53
- AppX responds back to App1
- We need Source Port so we know how to send back data
- Source Port = 5555



Summary

- UDP is a simple layer 4 protocol
- Uses ports to address processes
- Stateless

UDP Datagram

The anatomy of the UDP datagram

UDP Datagram

- UDP Header is 8 bytes only (IPv4)
- Datagram slides into an IP packet as “data”
- Port are 16 bit (0 to 65535)

UDP Datagram header

Offsets	Octet	0								1								2								3							
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
0	0	Source port																Destination port															
4	32	Length																Checksum															
		Data																															

https://www.ietf.org/rfc/rfc768.txt
https://en.wikipedia.org/wiki/User_Datagram_Protocol

Source Port and Destination Port

Offsets	Octet	0								1								2								3							
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
0	0	Source port																Destination port															
4	32	Length																Checksum															
		Data																															

Length & Checksum

Offsets	Octet	0								1								2								3							
Octet	Bit	0	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28	29	30	31
0	0	Source port																Destination port															
4	32	Length																Checksum															

Data

UDP Pros and Cons

The power and drawbacks of UDP

UDP Pros

- Simple protocol
- Header size is small so datagrams are small
- Uses less bandwidth
- Stateless
- Consumes less memory (no state stored in the server/client)
- Low latency - no handshake , order, retransmission or guaranteed delivery

UDP Cons

- No acknowledgement
- No guarantee delivery
- Connection-less - anyone can send data without prior knowledge
- No flow control
- No congestion control
- No ordered packets
- Security - can be easily spoofed

TCP

Transmission Control Protocol

TCP

- Stands for Transmission Control Protocol
- Layer 4 protocol
- Ability to address processes in a host using ports
- “Controls” the transmission unlike UDP which is a firehose
- Connection
- Requires handshake
- 20 bytes headers Segment (can go to 60)
- Stateful

TCP Use cases

- Reliable communication
- Remote shell
- Database connections
- Web communications
- Any bidirectional communication

TCP Connection

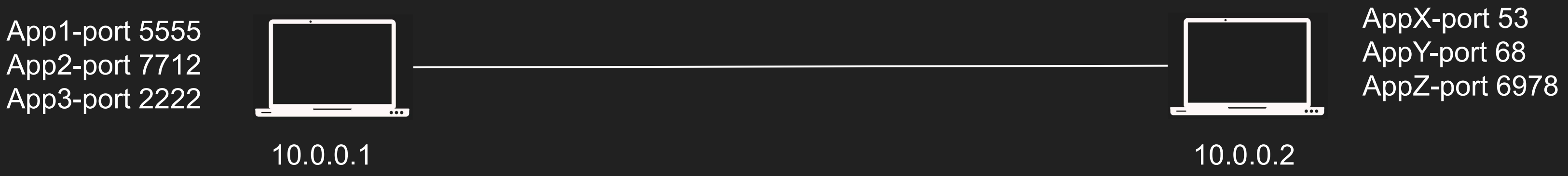
- Connection is a Layer 5 (session)
- Connection is an agreement between client and server
- Must create a connection to send data
- Connection is identified by 4 properties
  - SourceIP-SourcePort
  - DestinationIP-DestinationPort

TCP Connection

- Can’t send data outside of a connection
- Sometimes called socket or file descriptor
- Requires a 3-way TCP handshake
- Segments are sequenced and ordered
- Segments are acknowledged
- Lost segments are retransmitted

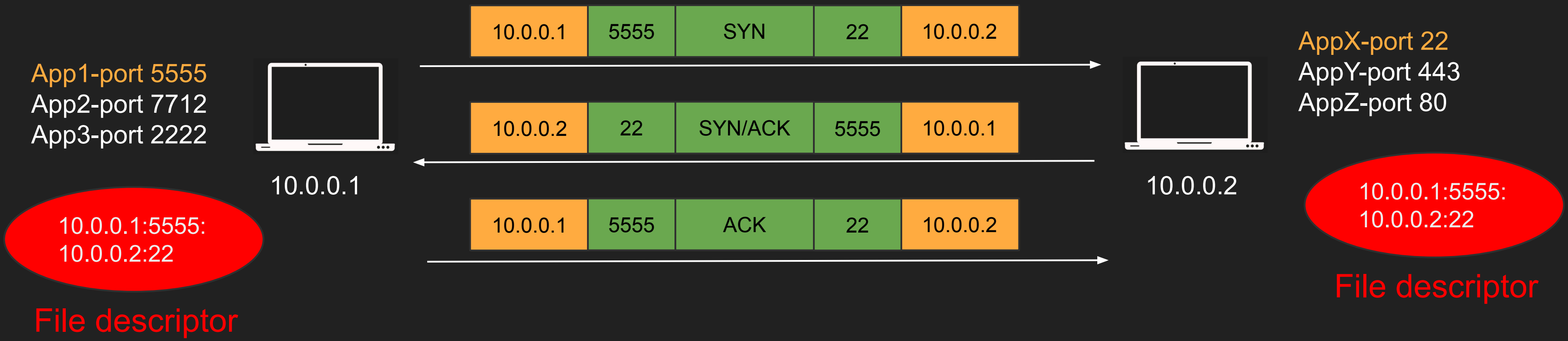
Multiplexing and demultiplexing

- IP target hosts only
- Hosts run many apps each with different requirements
- Ports now identify the “app” or “process”
- Sender multiplexes all its apps into TCP connections
- Receiver demultiplex TCP segments to each app based on connection pairs



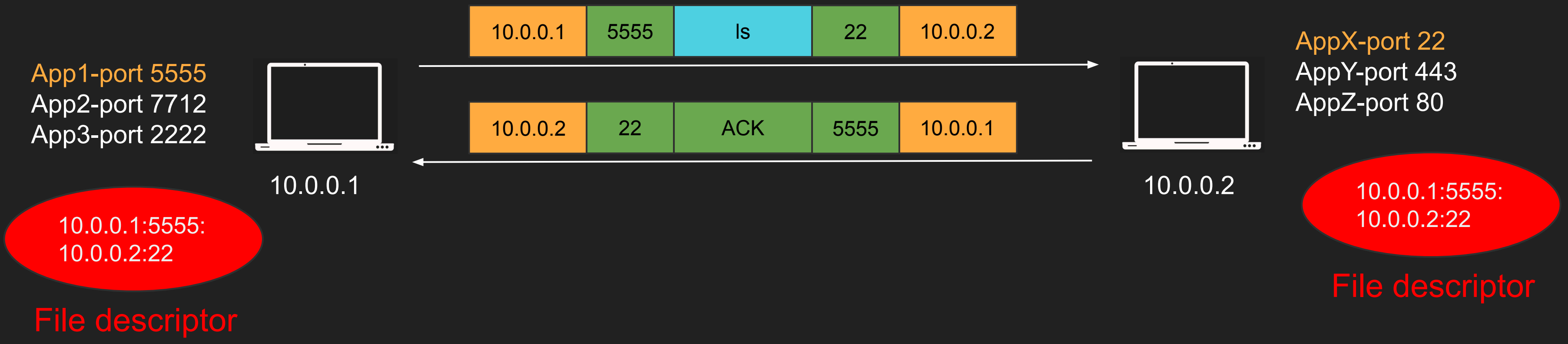
Connection Establishment

- App1 on 10.0.0.1 want to send data to AppX on 10.0.0.2
- App1 sends SYN to AppX to synchronous sequence numbers
- AppX sends SYN/ACK to synchronous its sequence number
- App1 ACKs AppX SYN.
- Three way handshake



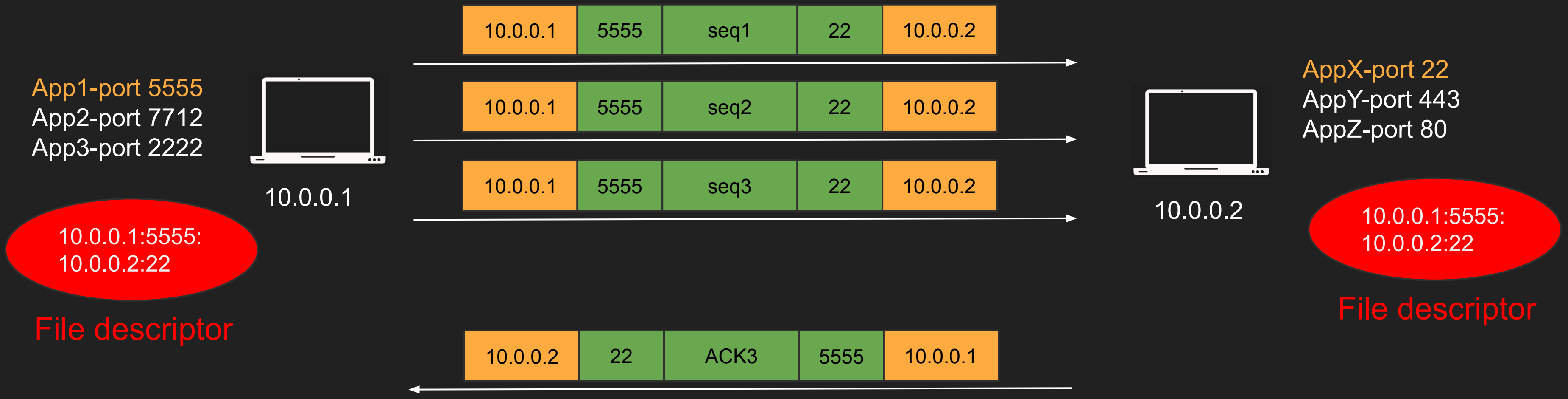
Sending data

- App1 sends data to AppX
- App1 encapsulate the data in a segment and send it
- AppX acknowledges the segment
- Hint: Can App1 send new segment before ack of old segment arrives?



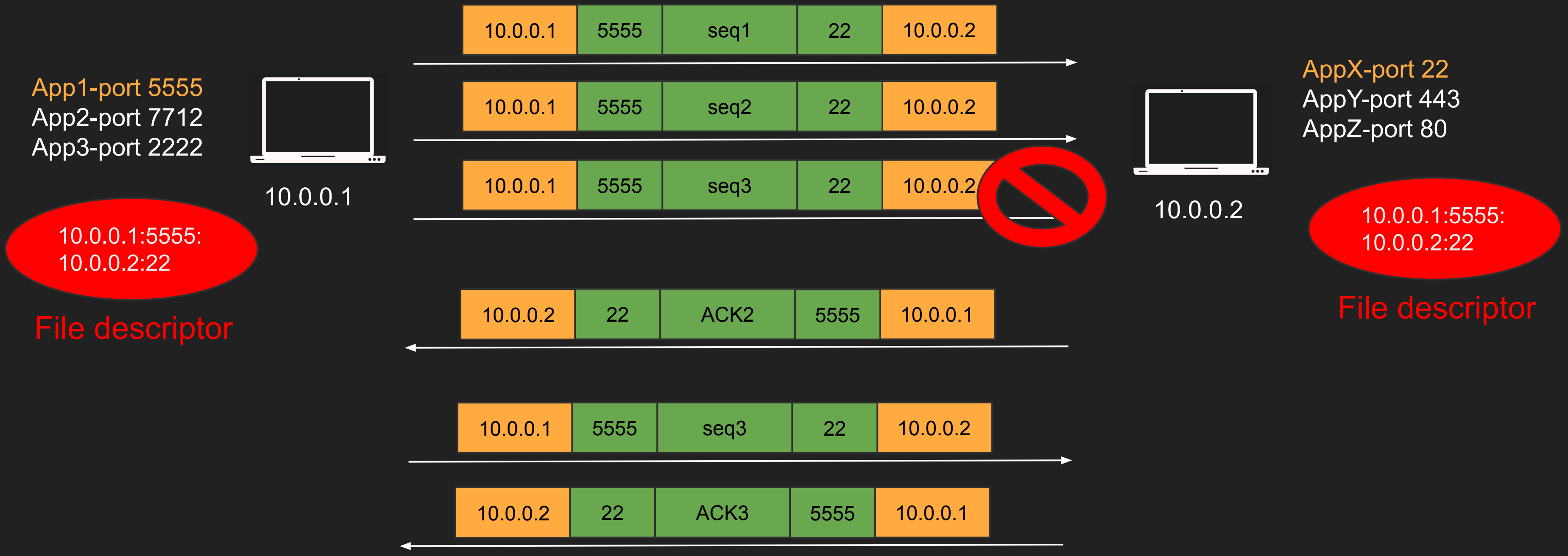
Acknowledgment

- App1 sends segment 1,2 and 3 to AppX
- AppX acknowledge all of them with a single ACK 3



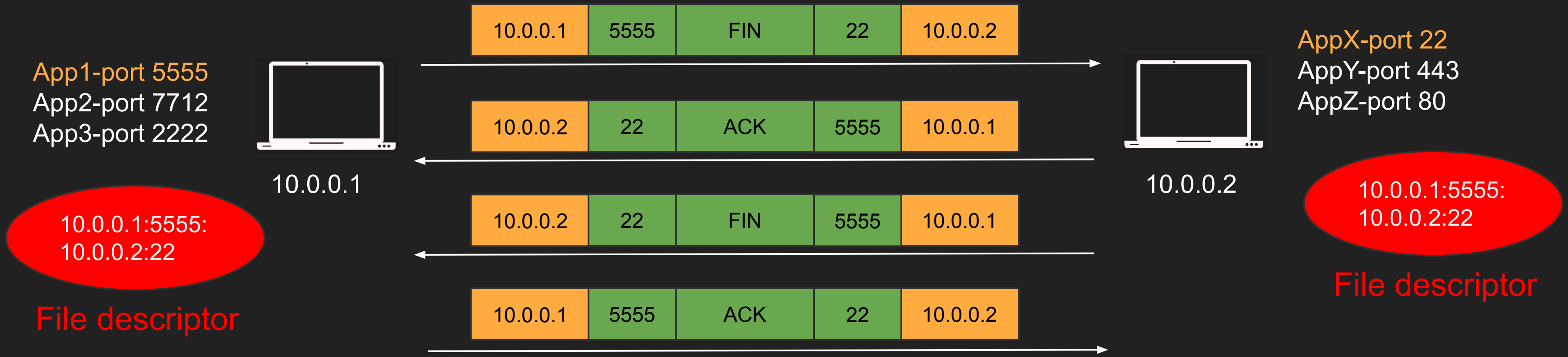
Lost data

- App1 sends segment 1,2 and 3 to AppX
- Seg 3 is lost, AppX acknowledge 3
- App1 resend Seq 3


App1-port 5555
App2-port 7712
App3-port 2222
10.0.0.1
10.0.0.1:5555:
10.0.0.2:22
File descriptor
10.0.0.1 | 5555 | seq1 | 22 | 10.0.0.2
10.0.0.1 | 5555 | seq2 | 22 | 10.0.0.2
10.0.0.1 | 5555 | seq3 | 22 | 10.0.0.2
10.0.0.2 | 22 | ACK2 | 5555 | 10.0.0.1
10.0.0.1 | 5555 | seq3 | 22 | 10.0.0.2
10.0.0.2 | 22 | ACK3 | 5555 | 10.0.0.1
AppX-port 22
AppY-port 443
AppZ-port 80
10.0.0.2
10.0.0.1:5555:
10.0.0.2:22
File descriptor

Closing Connection

- App1 wants to close the connection
- App1 sends FIN, AppX ACK
- AppX sends FIN, App1 ACK
- Four way handshake



Summary

- Stands for Transmission Control Protocol
- Layer 4 protocol
- “Controls” the transmission unlike UDP which is a firehose
- Introduces Connection concept
- Retransmission, acknowledgement, guaranteed delivery
- Stateful, connection has a state

TCP Segment

The anatomy of the TCP Segment

TCP Segment

- TCP segment Header is 20 bytes and can go up to 60 bytes
- TCP segments slides into an IP packet as “data”
- Port are 16 bit (0 to 65535)
- Sequences, Acknowledgment, flow control and more

husseinnasser

TCP Segment

Offsets	Octet	0								1								2								3							
Octet	Bit	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0
0	0	Source port																Destination port															
4	32	Sequence number																															
8	64	Acknowledgment number (if ACK set)																															
12	96	Data offset				Reserved 0 0 0			NS	CWR	ECE	URG	ACK	PSH	RST	SYN	FIN	Window Size															
16	128	Checksum																Urgent pointer (if URG set)															
20	160	Options (if *data offset* > 5. Padded at the end with "0" bits if necessary.)																															
⋮	⋮																																
60	480																																

https://en.wikipedia.org/wiki/Transmission_Control_Protocol
https://datatracker.ietf.org/doc/html/rfc793

Ports

Offsets	Octet	0								1								2								3							
Octet	**Bit**	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0
0	**0**	Source port																Destination port															
4	**32**	Sequence number																															
8	**64**	Acknowledgment number (if ACK set)																															
12	**96**	Data offset				Reserved **0 0 0**			NS	CWR	ECE	URG	ACK	PSH	RST	SYN	FIN	Window Size															
16	**128**	Checksum																Urgent pointer (if URG set)															
20	**160**	Options (if *data offset* > 5. Padded at the end with "0" bits if necessary.)																															
⋮	**⋮**																																
60	**480**																																

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Sequences and ACKs

Offsets	Octet	0								1								2								3							
Octet	Bit	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0
0	0	Source port																Destination port															
4	32	Sequence number																															
8	64	Acknowledgment number (if ACK set)																															
12	96	Data offset				Reserved 0 0 0			NS	CWR	ECE	URG	ACK	PSH	RST	SYN	FIN	Window Size															
16	128	Checksum																Urgent pointer (if URG set)															
20	160	Options (if *data offset* > 5. Padded at the end with "0" bits if necessary.)																															
⋮	⋮																																
60	480																																

Flow Control Window Size

Offsets	Octet	0								1								2								3							
Octet	Bit	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0
0	0	Source port																Destination port															
4	32	Sequence number																															
8	64	Acknowledgment number (if ACK set)																															
12	96	Data offset				Reserved 0 0 0			NS	CWR	ECE	URG	ACK	PSH	RST	SYN	FIN	Window Size															
16	128	Checksum																Urgent pointer (if URG set)															
20	160	Options (if *data offset* > 5. Padded at the end with "0" bits if necessary.)																															
⋮	⋮																																
60	480																																

9 bit flags

Offsets	Octet	0								1								2								3							
Octet	**Bit**	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0	7	6	5	4	3	2	1	0
0	**0**	Source port																Destination port															
4	**32**	Sequence number																															
8	**64**	Acknowledgment number (if ACK set)																															
12	**96**	Data offset				Reserved **0 0 0**			NS	CWR	ECE	URG	ACK	PSH	RST	SYN	FIN	Window Size															
16	**128**	Checksum																Urgent pointer (if URG set)															
20	**160**	Options (if *data offset* > 5. Padded at the end with "0" bits if necessary.)																															
⋮	**⋮**																																
60	**480**																																

Maximum Segment Size

- Segment Size depends the MTU of the network
- Usually 512 bytes can go up to 1460
- Default MTU in the Internet is 1500 (results in MSS 1460)
- Jumbo frames MTU goes to 9000 or more
- MSS can be larger in jumbo frames cases

TCP Pros and Cons

The power and drawbacks of TCP

TCP Pros

- Guarantee delivery
- No one can send data without prior knowledge
- Flow Control and Congestion Control
- Ordered Packets no corruption or app level work
- Secure and can’t be easily spoofed

TCP Cons

- Large header overhead compared to UDP
- More bandwidth
- Stateful - consumes memory on server and client
- Considered high latency for certain workloads (Slow start/ congestion/ acks)
- Does too much at a low level (hence QUIC)
  - Single connection to send multiple streams of data (HTTP requests)
  - Stream 1 has nothing to do with Stream 2
  - Both Stream 1 and Stream 2 packets must arrive
- TCP Meltdown
  - Not a good candidate for VPN

Overview of Popular Networking Protocols

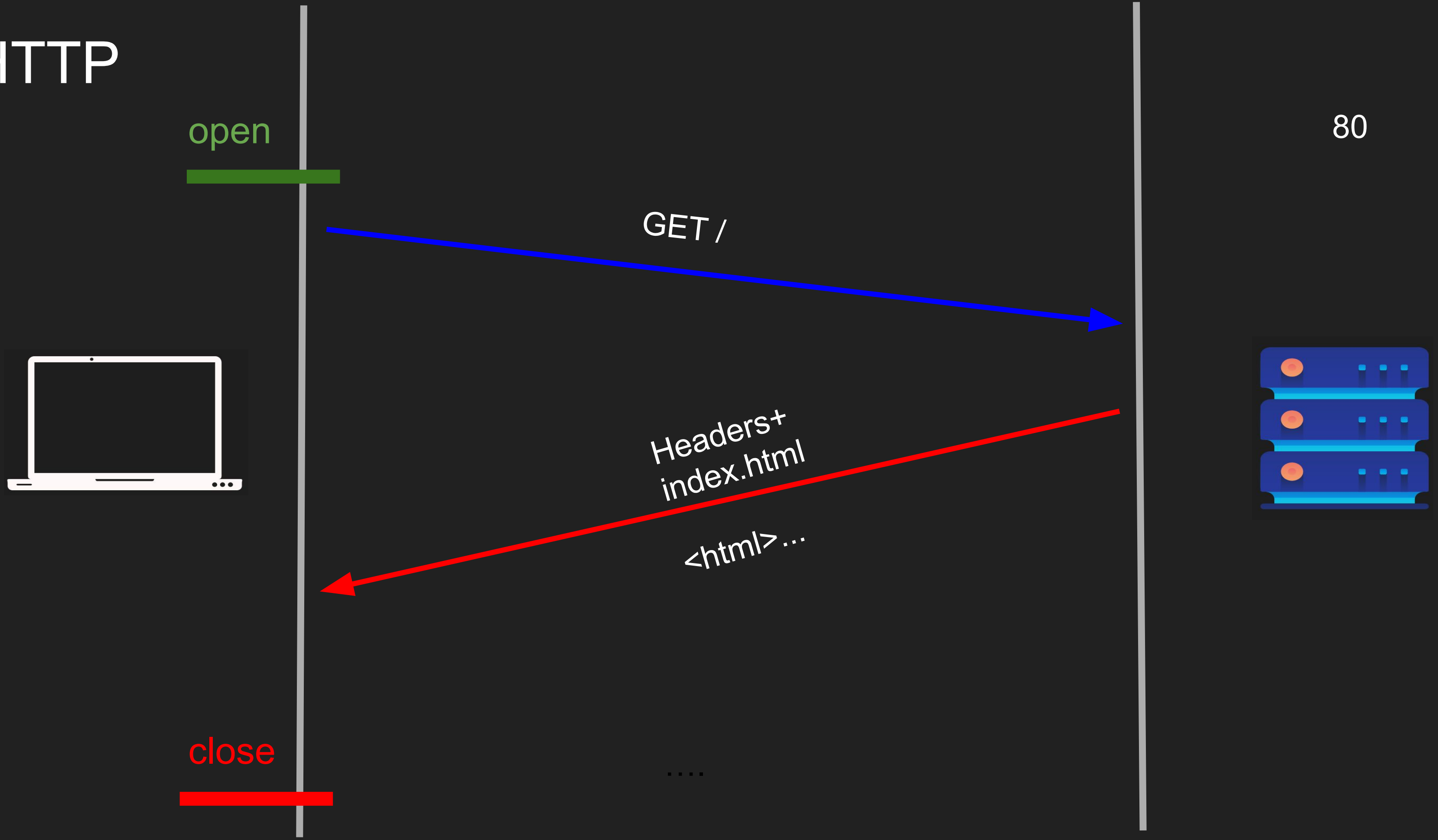
TLS

Transport Layer Security

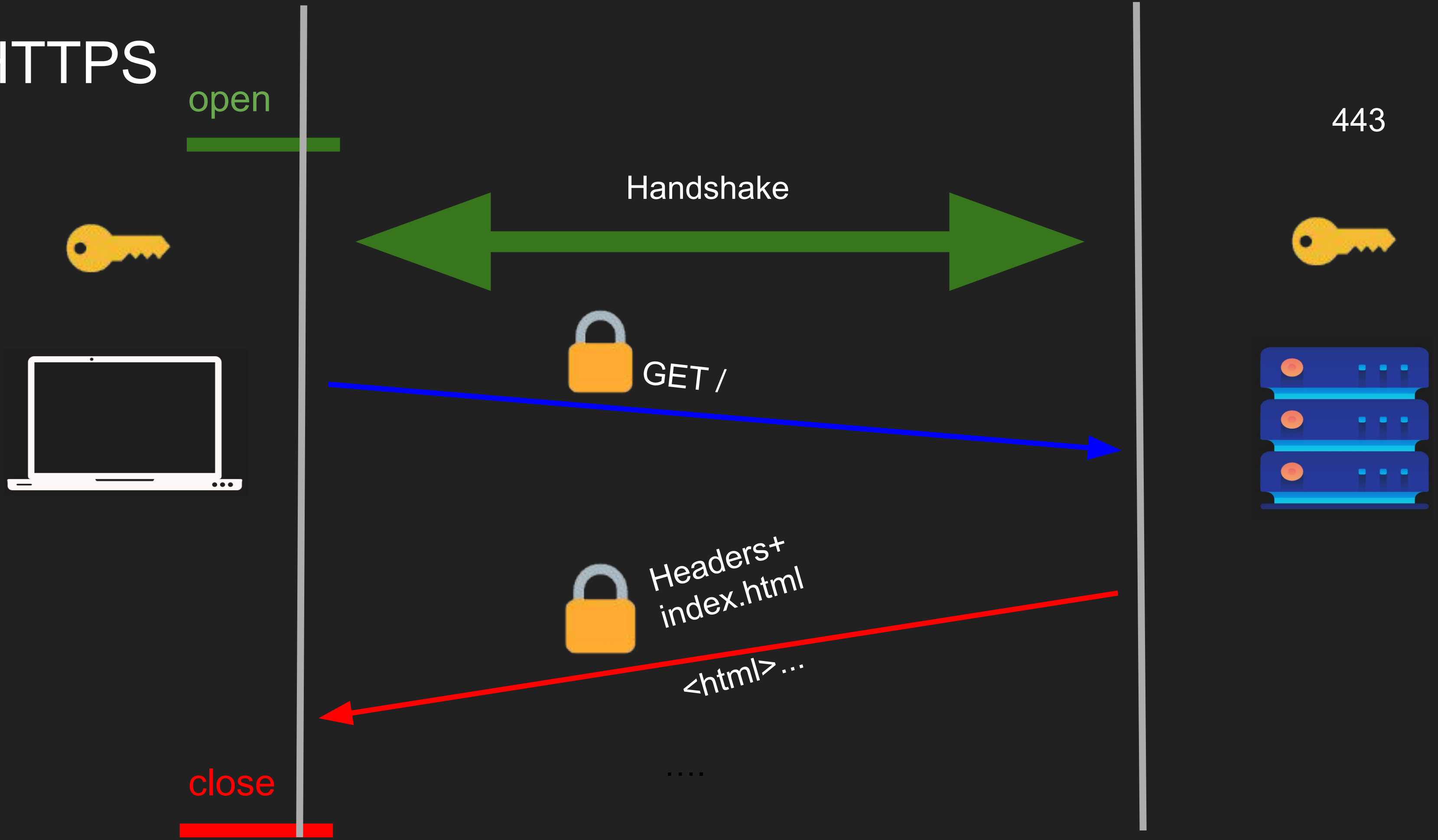
TLS

- Vanilla HTTP
- HTTPS
- TLS 1.2 Handshake
- Diffie Hellman
- TLS 1.3 Improvements

HTTP



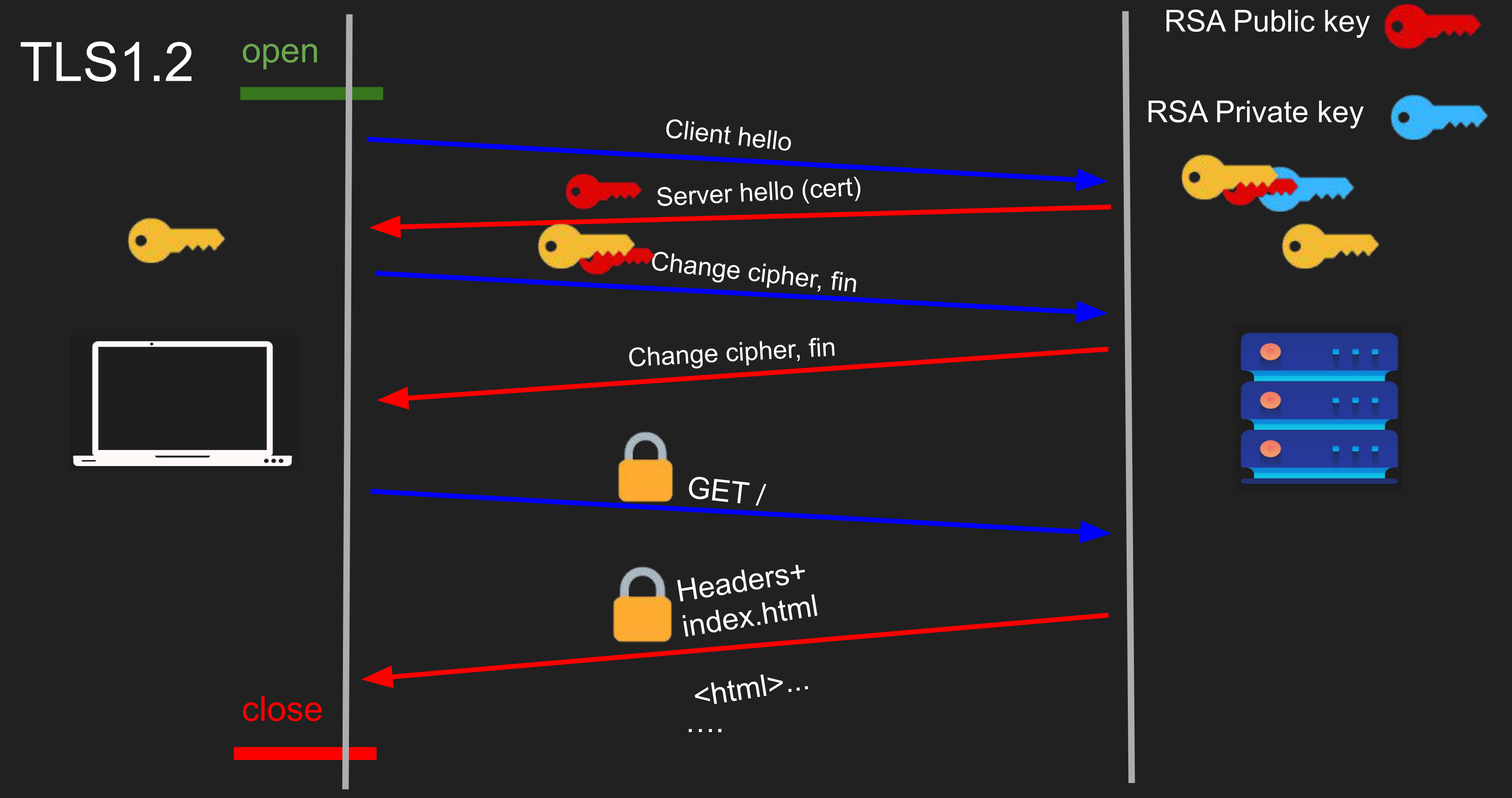
HTTPS



Why TLS

- We encrypt with symmetric key algorithms
- We need to exchange the symmetric key
- Key exchange uses asymmetric key (PKI)
- Authenticate the server
- Extensions (SNI, preshared, 0RTT)

TLS1.2


open
RSA Public key
RSA Private key
Client hello
Server hello (cert)
Change cipher, fin
Change cipher, fin
GET /
Headers+
index.html
<html>...
....
close

Diffie Hellman

Private x

+

Public g,n

+

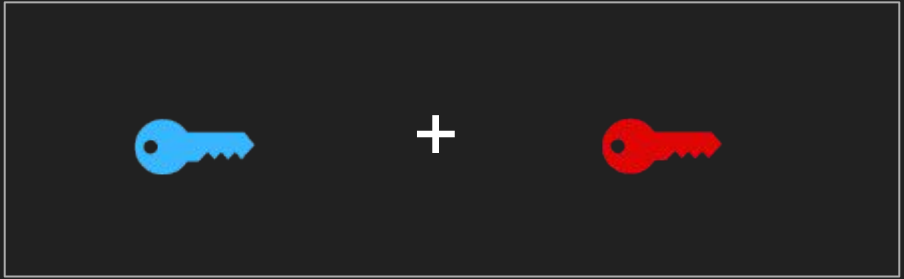
Private y

=

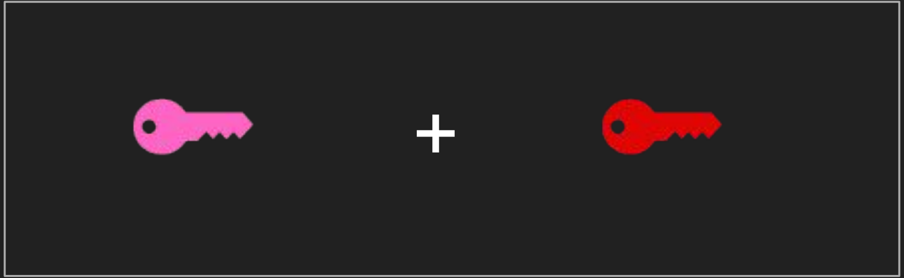
Symmetric key

Diffie Hellman

Public/
Unbreakable
/can be shared
g^x % n

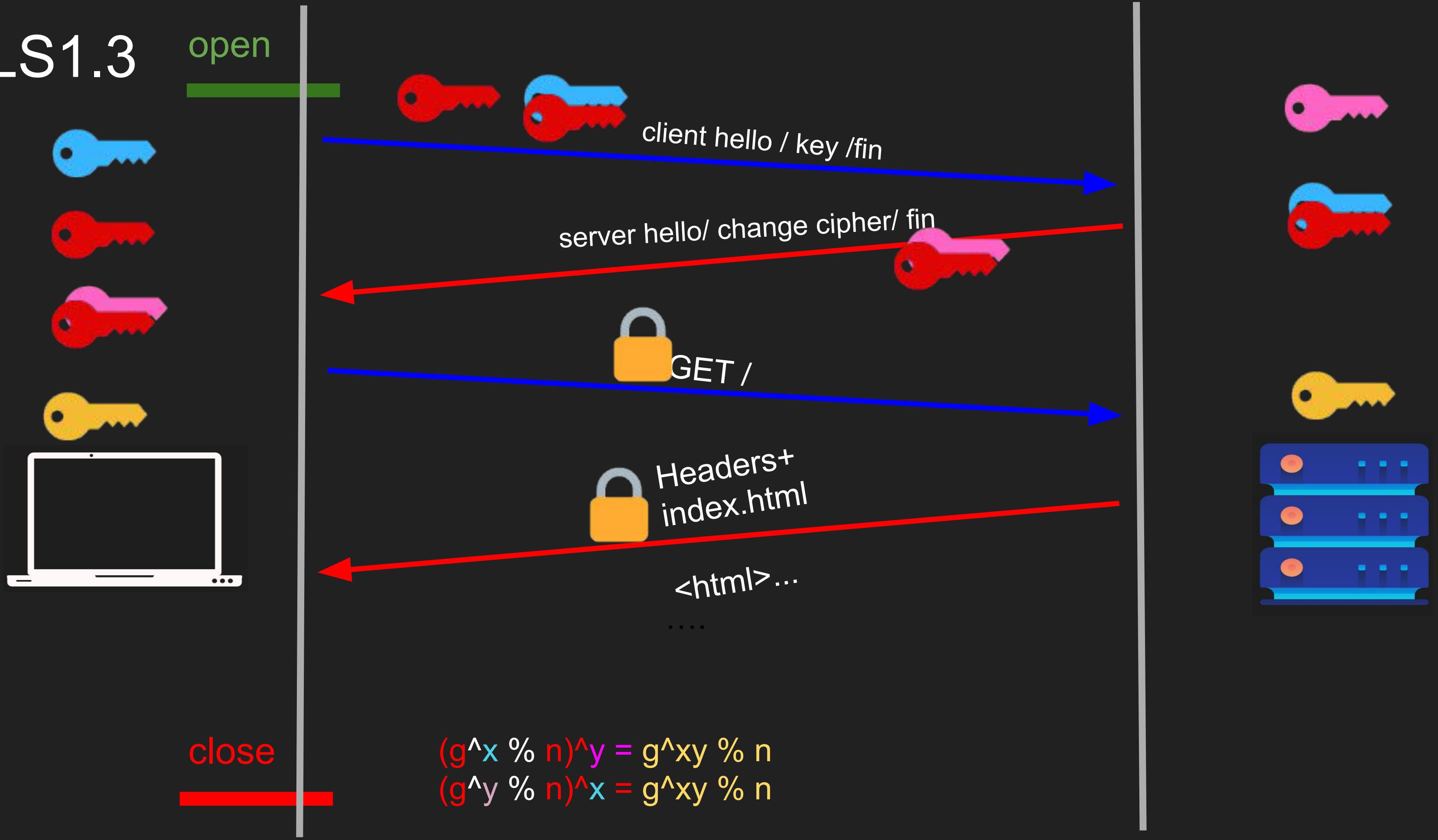


Public/
Unbreakable
/can be shared
g^y % n



(g^x % n)^y = g^xy % n
(g^y % n)^x = g^xy % n

TLS1.3



$$(g^x \% n)^y = g^{xy} \% n$$
$$(g^y \% n)^x = g^{xy} \% n$$

TLS Summary

- Vanilla HTTP
- HTTPS
- TLS 1.2 Handshake (two round trips)
- Diffie Hellman
- TLS 1.3 Improvements (one round trip can be zero)