

# Lecture 07 — ARRAYS IN C++ (FIRST-THOUGHT • INDUSTRY READY)

“Array samajh aaya = Programming ka backbone samajh aaya”

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## FIRST THOUGHT — ARRAY KYU AYA?

### Soch ke dekho (Real Life Trigger)

Tum teacher ho 

Tumhe **5 students** ke marks store karne hain → easy

Kal **100 students**

Parson **10,000 students**

 Old thinking (human-style):

```
int m1, m2, m3, m4, m5; // bas yahin tak thik
```

 Jaise hi scale badha:

- 100 variables 
- Loop ka use 
- Average, max, min 
- Code maintain 
- Interview me reject 

### ROOT PROBLEM (First Thought Principle)

“Ek hi type ka data bahut zyada quantity me handle karna hai”

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## ARRAY — ENGINEERING SOLUTION

### Definition (Simple + Deep)

Array ek aisa container hai jo

- same type ke multiple elements
- ek saath
- continuous memory me  
store karta hai

```
int marks[5] = {70, 80, 40, 50, 90};
```

- ➡ Ek naam → multiple values
  - ➡ Machine-friendly structure
  - ➡ CPU cache-friendly
  - ➡ Foundation of vectors, strings, matrices
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## REAL-LIFE MAPPING (BEST INTUITION)

Real Life	Array
Train seats	arr[0], arr[1], arr[2]
Classroom benches	rollNo based access
Parking slots	Slot 0, 1, 2
Hostel rooms	Room indexing

- 👉 Har jagah **numbering**
  - 👉 Har jagah **direct access**
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## ARRAY STRUCTURE (HOW C++ SEES IT)

```
datatype arrayName[size];
```

Example:

```
int arr[5];
```

- **datatype** → kitna memory lagega
- **arrayName** → base address ka naam
- **size** → fixed (compile-time decision)

➡ **Array ka size badalna possible nahi hota**  
(Ye limitation hi aage vector ko janam deti hai)

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## ARRAY INDEX 0 SE KYU START?

 Ye section agar samajh gaya → 80% confusion khatam

### CPU-Level Truth (Industry Reality)

Array me:

- Sirf **first element ka address** store hota hai
- Baaki sab ka address **calculate** hota hai

### Golden Formula

```
arr[i] ka address =  
base_address + (i × size_of_datatype)
```

### Example (int = 4 bytes)

Base address = 1000

arr[0] → 1000 + 0×4 = 1000

arr[1] → 1000 + 1×4 = 1004

arr[2] → 1000 + 2×4 = 1008

 Ab socho:

- Agar index 1 se start hota
- Har access me **-1** karna padta
- Extra CPU instruction
- Slow execution

### Conclusion (Interview Gold)

"Array indexing starts from 0 because it directly maps to memory offset calculation."

## ARRAY + LOOP = REAL PROGRAMMING

```
for(int i = 0; i < 5; i++){  
    cout << marks[i];  
}
```

- 👉 Loop ke bina array **almost useless**
  - 👉 Loop ke saath array = **power**
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## 🧠 ARRAY MEMORY — WHAT EXACTLY IS STORED?

### ❓ Beginner doubt:

“Array ka naam kya store karta hai?”  
“Har element ka address alag hota hai ya ek hi?”

### ✅ Truth (Industry Level):

- 👉 Array ka naam sirf FIRST element ka address store karta hai
  - 👉 Baaki elements ka address formula se calculate hota hai
- 

## 📌 MEMORY ADDRESS PRINT — LIVE PROOF

```
int arr[5] = {10, 20, 30, 40, 50};

for(int i = 0; i < 5; i++){
    cout << "Address of arr[" << i << "] = " << &arr[i] << endl;
}
```

### 🔍 Sample Output (example):

```
Address of arr[0] = 0x61ff08
Address of arr[1] = 0x61ff0c
Address of arr[2] = 0x61ff10
Address of arr[3] = 0x61ff14
Address of arr[4] = 0x61ff18
```

### 🧠 Observation:

- Har next address **+4 bytes** (because **int = 4 bytes**)
- Ye prove karta hai:

**Array elements contiguous memory me store hote hain**

### ⚠ Interview Line

“Arrays store elements in contiguous memory locations, enabling constant-time access.”

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## ⚠ GARBAGE VALUE — BIG BEGINNER TRAP

### ✗ Case 1: No initialization

```
int arr[5];  
  
for(int i = 0; i < 5; i++){  
    cout << arr[i] << " ";  
}
```

### 💡 Output (example):

4196720 -858993460 32767 0 144

### 🧠 WHY?

- Compiler ne memory allocate kar di
- **Par initialize nahi ki**
- Jo purani memory values thi → wahi print ho gayi

### ⚠ Isse kehte hain Garbage Value

⚠ Garbage value = memory me pehle se pada random data

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## ✓ Case 2: Partial initialization (VERY IMPORTANT)

```
int arr[5] = {10};  
  
for(int i = 0; i < 5; i++){  
    cout << arr[i] << " ";  
}
```

## Output:

10 0 0 0 0

### WHY?

- First value diya
- Baaki sab automatically **0 se fill**
- Compiler safety rule apply karta hai

### Golden Rule

“Partial initialization → remaining elements become zero”

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## Case 3: Full initialization (BEST PRACTICE)

```
int arr[5] = {10, 20, 30, 40, 50};
```

- ✓ No garbage
  - ✓ Predictable
  - ✓ Industry-safe
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### COMPARISON — TO CLEAR CONFUSION

Situation	Output
No value given	Garbage values
Partial value	Remaining = 0
Full values	Clean output



## BEGINNER MISTAKES (MUST READ)

- ✗ Assume karna ki array automatically zero hota hai
- ✗ Uninitialized array use karna
- ✗ Interview me bol dena "array safe hota hai"
  
- ✓ Always initialize
- ✓ Ya loop se input lo
- ✓ Ya `{0}` use karo

```
int arr[10] = {0}; // safest trick
```

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## REAL LIFE ANALOGY

 Socho:

- Tumne cupboard khola
- Kuch rakha hi nahi
- Jo pehle ka saman pada tha → wahi dikhega

 Garbage value exactly wahi hai

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## INITIALIZATION CONCEPT

### Full initialization

```
int arr[5] = {1, 2, 3, 4, 5};
```

### Partial initialization

```
int arr[4] = {12};
```

```
// Result: 12 0 0 0
```

 Reason:

- Compiler baaki memory zero se fill karta hai
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## MEMORY CONCEPT (CRUCIAL)

- Normal variables → scattered memory
- Array elements → **contiguous memory**

```
cout << &arr[i];
```

### Benefits:

- Fast traversal
  - Cache friendly
  - Predictable performance
  - Industry preference
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## LINEAR SEARCH (BASIC ALGO)

### Idea:

“Ek-ek element check karo”

```
bool found = false;
```

```
for(int i=0;i<n;i++){  
    if(arr[i] == key){  
        found = true;  
        break;  
    }  
}
```

- 📌 Worst case:  $O(n)$
- 📌 Base of binary search

## MAX / MIN LOGIC (GREEDY THINKING)

### Mindset:

“Pehle element ko answer maan lo”

```
int mx = arr[0];  
  
for(int i=1;i<n;i++){  
  
    if(arr[i] > mx) mx = arr[i];  
  
}  
  
}
```

👉 Ye thinking:

- Competitive programming
- System design
- ML preprocessing  
tak jaati hai

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## ⌚ ARRAY REVERSE — TWO APPROACHES

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### 🟡 ① Reverse WITHOUT Two Pointer (Extra Space)

🧠 Idea:

“Naya array bao, ulta copy karo”

```
int rev[n];  
  
for(int i=0;i<n;i++){  
  
    rev[i] = arr[n-1-i];  
  
}
```

- 👉 Time: O(n)
- 👉 Space: O(n)
- 👉 Simple but **memory costly**

### 🟢 ② Reverse WITH Two Pointer (Industry Preferred)

🧠 Idea:

- Start = 0
- End = n-1
- Swap till meet

```
int s = 0, e = n-1;  
  
while(s < e){  
  
    swap(arr[s], arr[e]);  
  
    s++;  
  
    e--;  
  
}
```

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- ❖ Time: O(n)
  - ❖ Space: O(1) ✓
  - ❖ Used in interviews & real systems

## ARRAY KYA SOLVE KARTA HAI?

- Repetition → Loop
  - Scaling → Single structure
  - Performance → Contiguous memory
  - Logic clarity → Index-based access
  - Industry use → Foundation Data Structure
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## 🏁 FINAL SUMMARY (NEVER FORGET)

- Array = same type + continuous memory
- Index 0 = memory offset logic
- Base address + formula = fast access
- Loop ke bina array incomplete
- Reverse ke 2 tareeke: space vs optimization
- Array name stores base address
- Elements stored in contiguous memory
- Uninitialized array → garbage values
- Partial initialization → remaining zero
- Always initialize arrays