



Lecture 07 — ARRAYS IN C++ (FIRST-THOUGHT • INDUSTRY READY)

“Array samajh aaya = Programming ka backbone samajh aaya”



FIRST THOUGHT — ARRAY KYU AYA?

? Soch ke dekho (Real Life Trigger)

Tum teacher ho 🧑

Tumhe **5 students** ke marks store karne hain → easy

Kal **100 students**

Parson **10,000 students**

✗ Old thinking (human-style):

```
int m1, m2, m3, m4, m5;    // bas yahin tak thik
```

🔧 Jaise hi scale badha:

- 100 variables ✗
- Loop ka use ✗
- Average, max, min ✗
- Code maintain ✗
- Interview me reject ✗

📌 ROOT PROBLEM (First Thought Principle)

“Ek hi type ka data bahut zyada quantity me handle karna hai”



ARRAY — ENGINEERING SOLUTION

🧠 Definition (Simple + Deep)

Array ek aisa container hai jo

- same type ke multiple elements
- ek saath
- continuous memory me store karta hai

```
int marks[5] = {70, 80, 40, 50, 90};
```

- 📌 Ek naam → multiple values
 - 📌 Machine-friendly structure
 - 📌 CPU cache-friendly
 - 📌 Foundation of vectors, strings, matrices
-

🌐 REAL-LIFE MAPPING (BEST INTUITION)

Real Life	Array
Train seats	arr[0], arr[1], arr[2]
Classroom benches	rollNo based access
Parking slots	Slot 0, 1, 2
Hostel rooms	Room indexing

- 👉 Har jagah **numbering**
 - 👉 Har jagah **direct access**
-

📦 ARRAY STRUCTURE (HOW C++ SEES IT)

```
datatype arrayName[size];
```

Example:

```
int arr[5];
```

- **datatype** → kitna memory lagega
- **arrayName** → base address ka naam
- **size** → fixed (compile-time decision)

📌 **Array ka size badalna possible nahi hota**
(Ye limitation hi aage vector ko janam deti hai)

ARRAY INDEX 0 SE KYU START?

 Ye section agar samajh gaya → 80% confusion khatam

CPU-Level Truth (Industry Reality)

Array me:

- Sirf **first element ka address** store hota hai
- Baaki sab ka address **calculate** hota hai

Golden Formula

`arr[i] ka address =`

`base_address + (i × size_of_datatype)`


Example (int = 4 bytes)

`Base address = 1000`

`arr[0] → 1000 + 0×4 = 1000`

`arr[1] → 1000 + 1×4 = 1004`

`arr[2] → 1000 + 2×4 = 1008`

 Ab socho:

- Agar index 1 se start hota
- Har access me **-1** karna padta
- Extra CPU instruction
- Slow execution

Conclusion (Interview Gold)

“Array indexing starts from 0 because it directly maps to memory offset calculation.”

ARRAY + LOOP = REAL PROGRAMMING

```
for(int i = 0; i < 5; i++){  
    cout << marks[i];  
}
```

- 👉 Loop ke bina array **almost useless**
 - 👉 Loop ke saath array = **power**
-

🧠 ARRAY MEMORY — WHAT EXACTLY IS STORED?

? Beginner doubt:

“Array ka naam kya store karta hai?”
“Har element ka address alag hota hai ya ek hi?”

✅ Truth (Industry Level):

- 👉 Array ka naam sirf **FIRST** element ka address store karta hai
 - 👉 Baaki elements ka address **formula se calculate** hota hai
-

📌 MEMORY ADDRESS PRINT — LIVE PROOF

```
int arr[5] = {10, 20, 30, 40, 50};

for(int i = 0; i < 5; i++){
    cout << "Address of arr[" << i << "] = " << &arr[i] << endl;
}
```

🔍 Sample Output (example):

```
Address of arr[0] = 0x61ff08
Address of arr[1] = 0x61ff0c
Address of arr[2] = 0x61ff10
Address of arr[3] = 0x61ff14
Address of arr[4] = 0x61ff18
```

🧠 Observation:

- Har next address **+4 bytes** (because `int = 4 bytes`)
- Ye prove karta hai:

Array elements contiguous memory me store hote hain

Interview Line

“Arrays store elements in contiguous memory locations, enabling constant-time access.”

GARBAGE VALUE — BIG BEGINNER TRAP

Case 1: No initialization

```
int arr[5];

for(int i = 0; i < 5; i++){
    cout << arr[i] << " ";
}
```


Output (example):

```
4196720 -858993460 32767 0 144
```

WHY?

- Compiler ne memory allocate kar di
- **Par initialize nahi ki**
- Jo purani memory values thi → wahi print ho gayi

Isse kehte hain **Garbage Value**

 Garbage value = memory me pehle se pada random data

Case 2: Partial initialization (VERY IMPORTANT)

```
int arr[5] = {10};

for(int i = 0; i < 5; i++){
    cout << arr[i] << " ";
}
```

✅ Output:

10 0 0 0 0

🧠 WHY?

- First value diya
- Baaki sab automatically **0 se fill**
- Compiler safety rule apply karta hai

📌 Golden Rule

“Partial initialization → remaining elements become zero”

✅ Case 3: Full initialization (BEST PRACTICE)

```
int arr[5] = {10, 20, 30, 40, 50};
```

- ✓ No garbage
 - ✓ Predictable
 - ✓ Industry-safe
-

🧠 COMPARISON — TO CLEAR CONFUSION

Situation	Output
No value given	Garbage values ❌
Partial value	Remaining = 0 ✅
Full values	Clean output ✅

BEGINNER MISTAKES (MUST READ)

- ✗ Assume karna ki array automatically zero hota hai
- ✗ Uninitialized array use karna
- ✗ Interview me bol dena “array safe hota hai”

- ✓ Always initialize
- ✓ Ya loop se input lo
- ✓ Ya `{0}` use karo

```
int arr[10] = {0};    // safest trick
```

REAL LIFE ANALOGY

 Socho:

- Tumne cupboard khola
- Kuch rakha hi nahi
- Jo pehle ka saman pada tha → wahi dikhega

 Garbage value exactly wahi hai

INITIALIZATION CONCEPT

Full initialization

```
int arr[5] = {1,2,3,4,5};
```

Partial initialization

```
int arr[4] = {12};
```

```
// Result: 12 0 0 0
```


 Reason:

- Compiler baaki memory zero se fill karta hai
-

MEMORY CONCEPT (CRUCIAL)


- Normal variables → scattered memory
- Array elements → **contiguous memory**

```
cout << &arr[i];
```

 Benefits:

- Fast traversal
 - Cache friendly
 - Predictable performance
 - Industry preference
-


LINEAR SEARCH (BASIC ALGO)


 Idea:

“Ek-ek element check karo”

```
bool found = false;
```

```
for(int i=0;i<n;i++){  
    if(arr[i] == key){  
        found = true;  
        break;  
    }  
}
```

 Worst case: $O(n)$

 Base of binary search

MAX / MIN LOGIC (GREEDY THINKING)

 Mindset:

“Pehle element ko answer maan lo”


```
int mx = arr[0];

for(int i=1;i<n;i++){

    if(arr[i] > mx) mx = arr[i];

}
```

👉 Ye thinking:

- Competitive programming
- System design
- ML preprocessing
tak jaati hai

ARRAY REVERSE — TWO APPROACHES

① Reverse WITHOUT Two Pointer (Extra Space)

🧠 Idea:

“Naya array banao, ulta copy karo”

```
int rev[n];

for(int i=0;i<n;i++){

    rev[i] = arr[n-1-i];

}
```

- 📌 Time: $O(n)$
- 📌 Space: $O(n)$
- 📌 Simple but **memory costly**

② Reverse WITH Two Pointer (Industry Preferred)

🧠 Idea:

- Start = 0
- End = $n-1$
- Swap till meet

```
int s = 0, e = n-1;

while(s < e){


    swap(arr[s], arr[e]);

    s++;

    e--;

}
```

📌 Time: $O(n)$

📌 Space: $O(1)$ 

📌 Used in interviews & real systems

ARRAY KYA SOLVE KARTA HAI?

- Repetition → Loop
- Scaling → Single structure
- Performance → Contiguous memory
- Logic clarity → Index-based access
- Industry use → Foundation Data Structure

🚩 FINAL SUMMARY (NEVER FORGET)

- Array = same type + continuous memory
- Index 0 = memory offset logic
- Base address + formula = fast access
- Loop ke bina array incomplete
- Reverse ke 2 tareeke: space vs optimization
- Array name stores base address
- Elements stored in contiguous memory
- Uninitialized array → garbage values
- Partial initialization → remaining zero
- Always initialize arrays

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