

Assignment - 3

Compiler Design

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Section - E1

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Ques 1. Define a Syntax directed definition (SDD). Differentiate between S-attributed and L-attributed definitions with suitable examples.

Ans An SDD is a Context free grammar in which each grammar symbol is associated with attributes, and each production has rules that define how attribute values are computed. These rules specify how attributes are calculated from other attribute values.

Difference between S-attributed and L-attributed definitions:

- S-attributed: Use only synthesized attributes. Values are computed from the attributes of the children in the parse tree.
ex:- Evaluating Arithmetic Expressions
- L-attributed: Use both synthesized and inherited attributes. Values are computed using the attributes of the children/sibling.
ex:- Type checking where inherited attributes are passed down.

Ques 2. Discuss how bottom-up evaluation for S-attributed definitions is performed. Why is it not suitable for L-attributed definitions?

Ans Bottom-up parsing computes synthesized attributes from the leaves to the root. Since S-attributed definitions have only synthesized attributes, they suit bottom-up evaluation. Not suitable for L-attributed definitions requires inherited attributes, which need to be passed down or across sibling, making bottom-up evaluation impractical.

Ques 3. Described L-attributed definitions. Provide an example where synthesized and inherited attributes with are both used.

Ans 3. An L-attributed is a special kind of Syntax-directed definitions where each inherited attribute can be evaluated from left to right using the inherited attributes of parent nodes or synthesized attributes of left siblings.

ex:- In an expression grammar, synthesized attributes calculate the expression value, while inherited attributes calculate the expression value, while inherited attributes manage variable types and scope.

Ques 4. What are the challenge involved in top-down translation of Syntax-directed definitions? Explain with an example.

1. Circular Dependencies: Attributes may depend on each other, causing evaluation conflicts.
2. Inherited Attribute management: Passing attributes down the parse tree is complex.
3. Complexity in non-LL Grammars: Difficult for left-recursive or ambiguous grammars.

ex:- In expression evaluation, managing inherited attributes for operator precedence can be problematic.

Ques 5. How does the direction of attribute evaluation affect the implementation of translation Schemes in Compilers?

- Ans 5.
1. Top-down (L-attributed): Efficient for inherited attributes but struggles with synthesized ones.
 2. Bottom-up (S-attributed): Efficient for synthesized attributes but not for inherited one.
 3. Hybrid approach: Combining both for complex cases.
 4. Efficiency Consideration: Direct influence on parsing strategy (LL for top-down, LR for bottom-up)
 5. Implementation Complexity: Direction affects how easily attributes are propagated.

Ques 6.7. Consider grammar with the following translation rules and ϵ as start symbol.

$$E \rightarrow E \# T \quad \{ E.val = E.val * T.val \}$$

$$E \rightarrow T \quad \{ E.val = T.val \}$$

$$T \rightarrow T \& F \quad \{ T.value = T.val + F.val \}$$

$$T \rightarrow F \quad \{ T.val = F.val \}$$

$$F \rightarrow \text{num} \quad \{ F.val = \text{num.val} \}$$

Compute $E.value$ for the root of the parse tree for the expression: $2 \# 3 \& 5 \# 6 \& 4$.

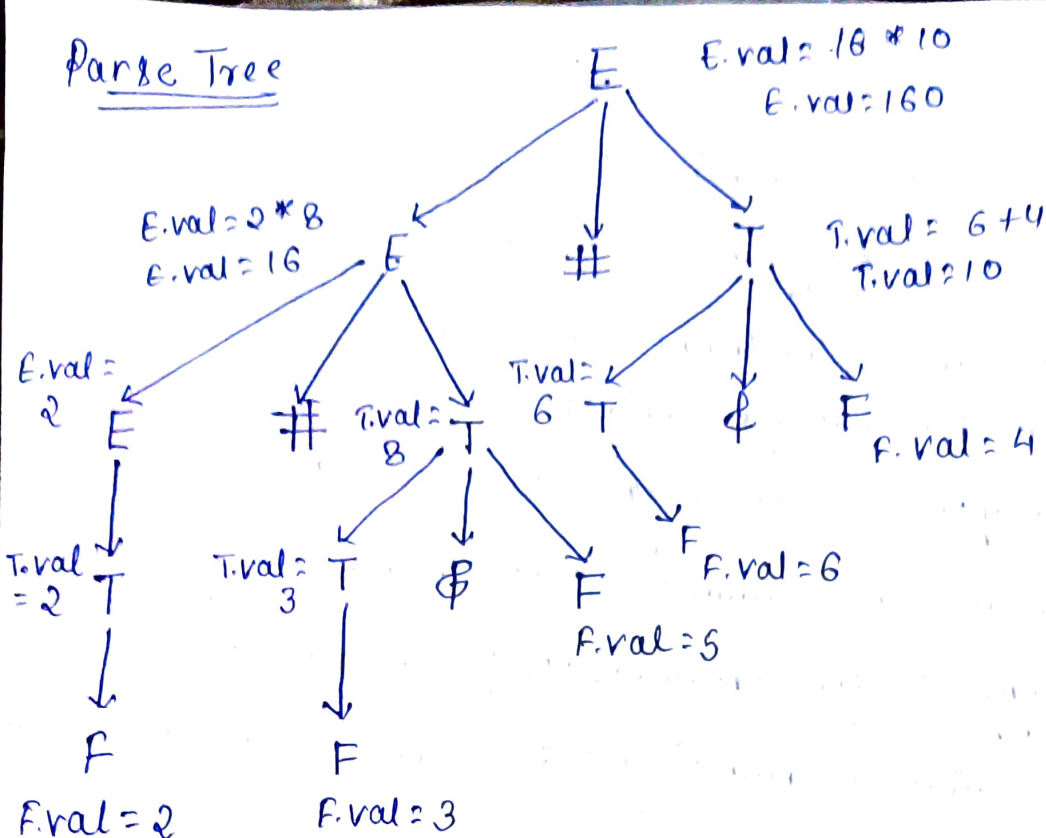
a) 200

b) 180

c) 160

d) 40

Parse Tree



So, the value of the expression is 160

Ques 6. Consider the translation scheme shown below:

S → TR

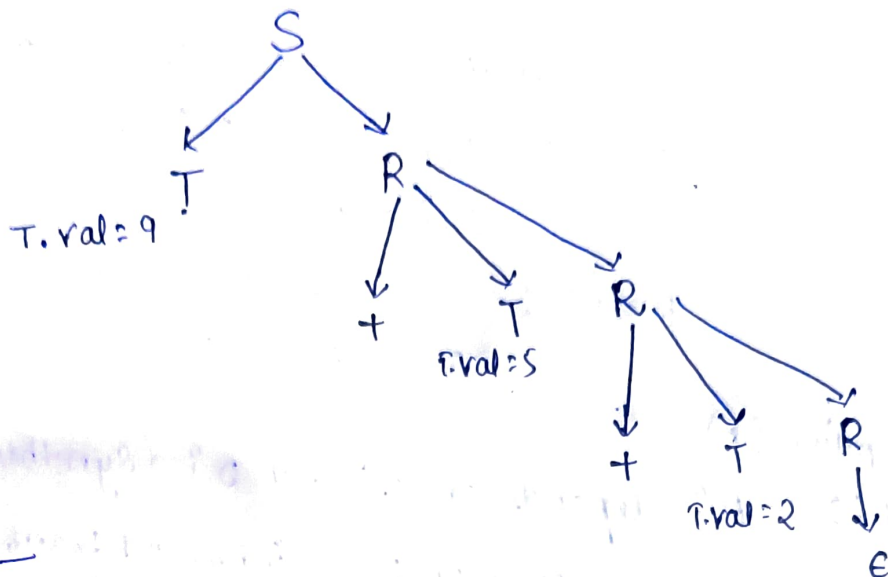
$R \rightarrow + T \mid \text{Print}(' + '); \mid R \mid \epsilon$

$T \rightarrow \text{num} \mid \text{print}(\text{num.val});$

Here num is a token and num.val is a integer value for input string 9+5+2, the translation will print.

- a) $9+5+2$
b) $95+2+$
c) $95\ 2++$
d) $++9\ 5\ 2$

Parse Tree



output :-

9 + 5 + 2

Ques 8. Consider the following translation scheme

$S \rightarrow ER$

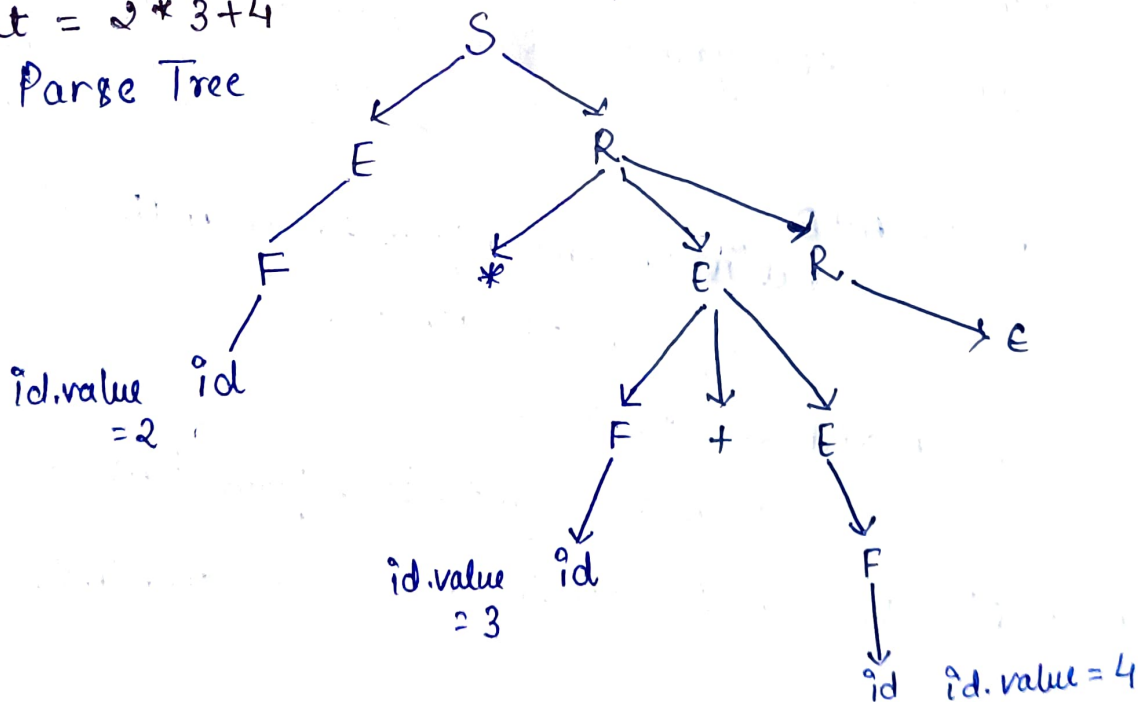
$R \rightarrow *E \mid \text{print}(*); \mid R \mid \epsilon$

$E \rightarrow F + E \mid \text{print}(+); \mid F$

$F \rightarrow (s) \mid \text{id} \mid \text{print}(\text{id.value});$

Input = 2 * 3 + 4

Ans Parse Tree



Output :- $2 * + 3 4$

Ques 9. Let the 2 production be $P \rightarrow CD$ and $P \rightarrow AB$

Consider rules: Rule 1: $P.i = C.i + D.i$ and

$$C.i = D.i - 1$$

Rule 2: $A.i = P.i * 2$ and

$$B.i = A.i + P.i$$

Ans 9. option (b)

Rule 1: $P.i$ depends on $C.i$ and $D.i$ {Synthesized from child}
 $C.i$ depends on $D.i$ {inherited from sibling}

Rule 1: is not L- attributed as $C.i$ is defined using $D.i$ which is not a parent child relation.

Rule 2: $A.i$ depends on $P.i$ {Synthesized from parent}
 $B.i$ depends on $A.i$ and $P.i$ {both synthesized from parent}

Rule 2: is L- attributed

Ques 10. Consider the grammar: $S \rightarrow A+B$ { $S.val = A.val + B.val$ }
 $A \rightarrow id$ { $A.val = id.val$ }
 $B \rightarrow id$ { $B.val = id.val$ }

Ans 10. In $S \rightarrow A+B$

$S.val$ is Synthesized from $A.val$ and $B.val$

Synthesized attribute

In $A \rightarrow id$ and $B \rightarrow id$

Both are synthesized from $id.val$

ans is (c) SDT is L- attributed and S- attributed.